

Signed Multiplier (Booth)

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Signed Number

Here, Sign number will mean 2's complement number.

For example,

$$\begin{aligned} &1011 \text{ (-5)} \\ &\quad \text{is} \\ &0100 \text{ (1's complement) + 1} \\ &\quad = \\ &0101 \text{ (+5) 's Negative version.} \end{aligned}$$

Signed Multiplication Example

For Signed number,

$$\begin{array}{r} 1011 \quad (-5) \\ 1101 \quad (-3) \\ \hline 00001011 \\ 0000000X \\ 001011XX \\ 01011XXX \\ \hline 10001111 \quad (-113) \end{array}$$

The “Binary”
Multiplication
Table

*	0	1
0	0	0
1	0	1

We can see that multiplying 2 4-bit signed binary numbers do not work like unsigned numbers.

Signed Multiplication Example

For Signed number,

1011 (-5)

1101 (-3)

This numbers are not signed numbers. Since multiplying 2 4-bit numbers result in 8 bit product, so multiplicand and multiplier must be sign extended.

11111011 (-5)

11111101 (-3)

Since numbers are negative numbers, they are sign extended with 1s.

Signed Multiplication Example

For Signed number,

$$\begin{array}{r}
 11111011 \quad (-5) \\
 11111101 \quad (-3) \\
 \hline
 11111011 \\
 0000000X \\
 111011XX \\
 11011XXX \\
 1011XXXX \\
 101XXXXX \\
 10XXXXXX \\
 1XXXXXXXX \\
 \hline
 101 \mid 00001111 \quad (+15)
 \end{array}$$

The “Binary”
Multiplication
Table

*	0	1
0	0	0
1	0	1

We can **calculate it like unsigned numbers** just by **sign extending it** and also by **performing twice many steps** (8 steps vs 4 steps) than unsigned numbers.

Booth Multiplication Algorithm

One of most efficient algorithm to perform multiplication with signed numbers is **Booth Multiplication Algorithm**.

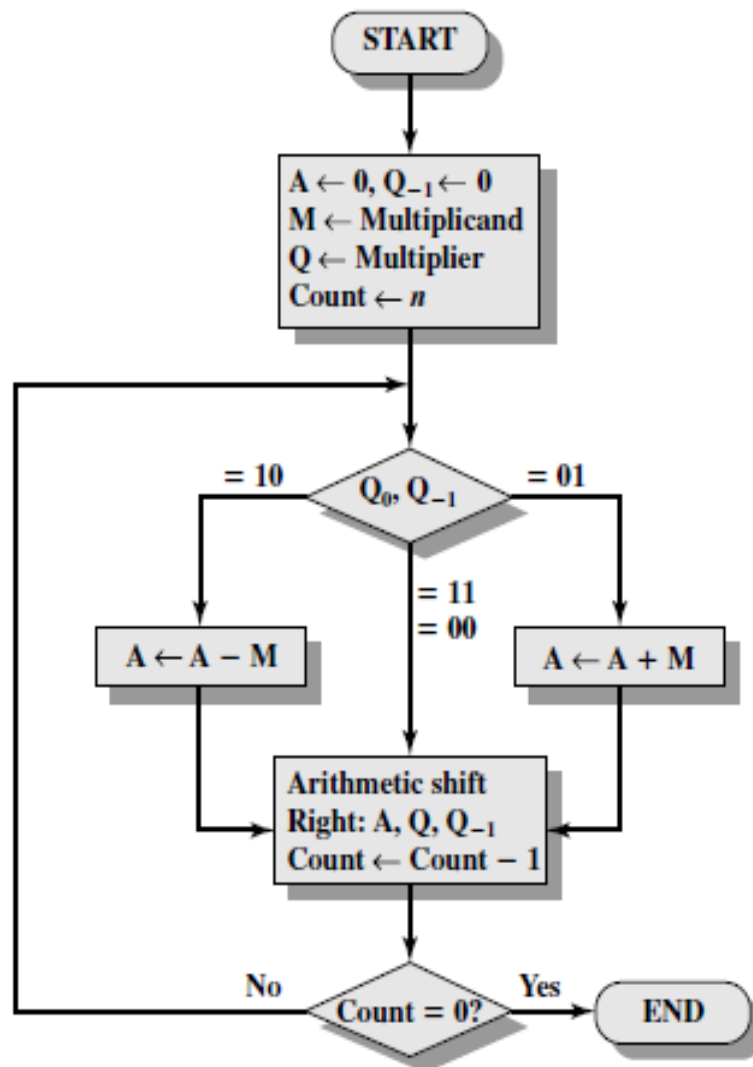


Figure: Booth's Algorithm
for Two's Complement Multiplication

Here,
 $A - M = \text{Signed (2's Complement) Subtraction}$

Booth Multiplication Example

Table: Example of Booth Multiplication (Type 1)

Multiplicand, **M** = 1011 (4-bit)

And Multiplier, **Q** = 1101 (4-bit)

So, Product, **P = A,Q** = 00001111 (8-bit)

A	Q	Q ₁	M	Operation	Cycle
0000	110 <u>1</u>	<u>0</u>	1011	Initial Value	
0101 0010	1101 111 <u>0</u>	0 <u>1</u>	1011 1011	A = A-M Shift A,Q to Right	1
1101 1110	1110 111 <u>1</u>	1 <u>0</u>	1011 1011	A = A+M Shift A,Q to Right	2
0011 0001	1111 111 <u>1</u>	0 <u>1</u>	1011 1011	A = A-M Shift A,Q to Right	3
0001 0000	1111 1111	1 1	1011 1011	A = A Shift A,Q to Right	4

Type - 1

Booth Multiplication Example

Table: Example of Booth Multiplication (Type 1)

Multiplicand, **Y** = 1011 (4-bit)

And Multiplier, **X** = 1101 (4-bit)

So, Product, **P** = 00001111 (8-bit)

P	Y	X, X ₋₁	Operation	Cycle
0000 0000	1111 1011		Initial Value	1
0000 0101 0000 0101	1111 1011 1111 0110	110 <u>10</u>	P = P-Y Shift Y to Left	
1111 1011 1111 1011	1111 0110 1110 1100	11 <u>01</u> 0	P = P+Y Shift Y to Left	
0000 1111 0000 1111	1110 1100 1101 1000	1 <u>10</u> 10	P = P-Y Shift Y to Left	
0000 1111 0000 1111	1101 1000 1011 0000	<u>11</u> 010	P = P Shift Y to Left	

Type – 2

We are going to create a Booth Multiplier based on Type-2

Booth Multiplier Building Block (Cell B)

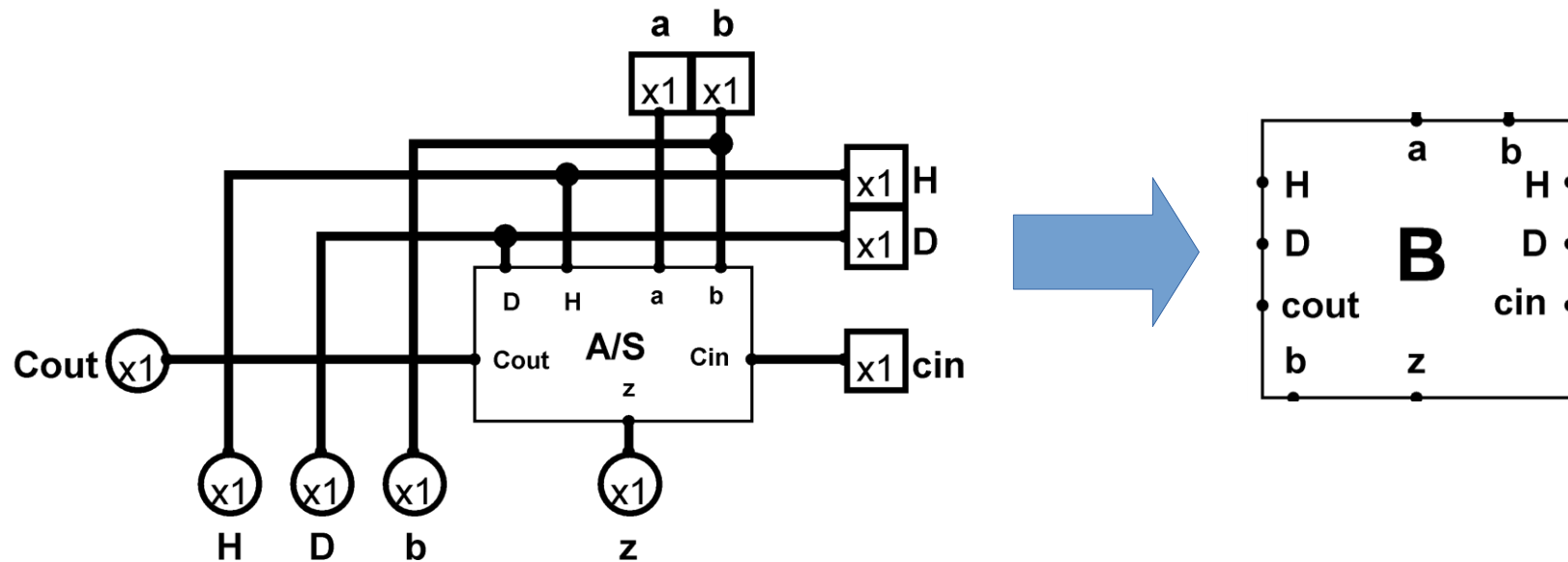


Figure: Cell B (Building Block)

Here,

1. A/S is combination of Full Adder and Full Subtractor and will switch in between them based on H and D value.
2. Value of H, D and b will propagate to next Block which means they are also outputs.

We are going to use a building block (Cell B) to create Booth multiplier₁₀

Booth Multiplier Building Block (Cell B)

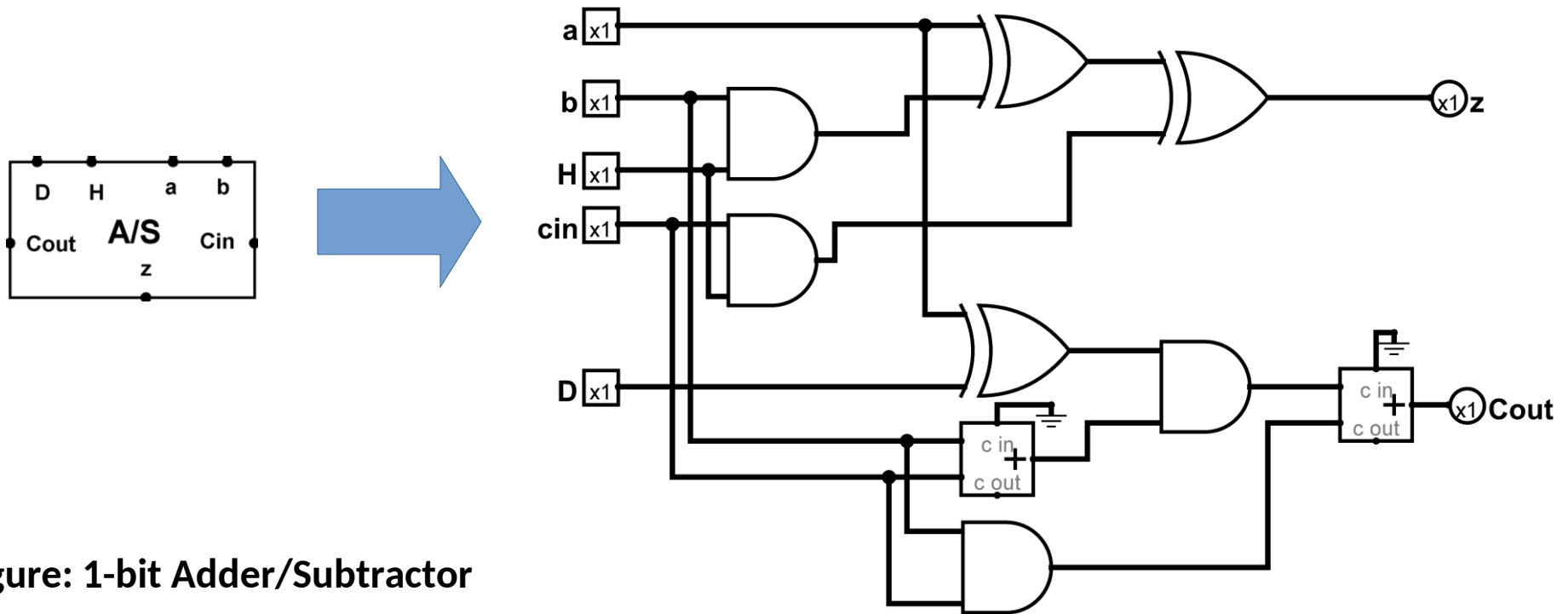


Figure: 1-bit Adder/Subtractor

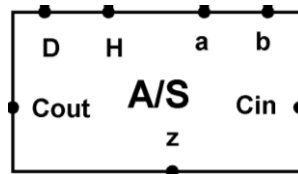
$$z = a \text{ XOR } (b \text{ AND } H) \text{ XOR } (c_{in} \text{ AND } H)$$

$$C_{out} = (a \text{ XOR } D) \text{ AND } (b \text{ OR } c_{in}) \text{ OR } (b \text{ AND } c_{in})$$

Here,

1. H and D are control inputs which will turn A/S into Full Adder or Full Subtractor.
2. Circuit will turn into Full Adder based on H and D value where a , b and c_{in} (Carry in) are inputs, c_{out} (Carry out) and z (Sum) are outputs.
3. Circuit will also turn into Full Subtractor based on H and D value where a , b and c_{in} (Borrow in) are inputs, c_{out} (Borrow out) and z (Difference) are outputs.

Booth Multiplier Building Block (Cell B)



$$z = a \text{ XOR } (b \text{ AND } H) \text{ XOR } (cin \text{ AND } H)$$

$$C_{out} = (a \text{ XOR } D) \text{ AND } (b \text{ OR } cin) \text{ OR } (b \text{ AND } cin)$$

Figure: 1-bit Adder/Subtractor

Table: Function Table for 1-bit Adder/Subtractor

H	D	z and C _{out}	Function
0	0	$z = a \text{ XOR } 0 \text{ XOR } 0 = a$ $C_{out} = a \text{ AND } (b \text{ OR } cin) \text{ OR } (b \text{ AND } cin)$ $= (a \text{ AND } b) \text{ OR } (a \text{ AND } cin) \text{ OR } (b \text{ AND } cin)$	$z = a$ (no operation)
0	1	$z = a \text{ XOR } 0 \text{ XOR } 0 = a$ $C_{out} = \sim a \text{ AND } (b \text{ OR } cin) \text{ OR } (b \text{ AND } cin)$ $= (\sim a \text{ AND } b) \text{ OR } (\sim a \text{ AND } cin) \text{ OR } (b \text{ AND } cin)$	$z = a$ (no operation)
1	0	$z = a \text{ XOR } b \text{ XOR } cin$ $C_{out} = a \text{ AND } (b \text{ OR } cin) \text{ OR } (b \text{ AND } cin)$ $= (a \text{ AND } b) \text{ OR } (a \text{ AND } cin) \text{ OR } (b \text{ AND } cin)$	$C_{out}, z = a + b + c$ (add) [Full Adder]
1	1	$z = a \text{ XOR } b \text{ XOR } cin$ $C_{out} = \sim a \text{ AND } (b \text{ OR } cin) \text{ OR } (b \text{ AND } cin)$ $= (\sim a \text{ AND } b) \text{ OR } (\sim a \text{ AND } cin) \text{ OR } (b \text{ AND } cin)$	$C_{out}, z = a - b - c$ (subtract) [Full (Unsigned) Subtractor] (Not 2's Complement Subtraction)

Booth Multiplier Building Block (Cell C)

Table: Relationship between X_i , X_{i-1} and H, D

X_i	X_{i-1}	H	D	Operation
0	0	0	X	No Operation
0	1	1	0	Add
1	0	1	1	Subtract
1	1	0	X	No Operation

$$H = X_i \text{ XOR } X_{i-1}$$
$$D = X_i \text{ AND } (\sim X_{i-1})$$

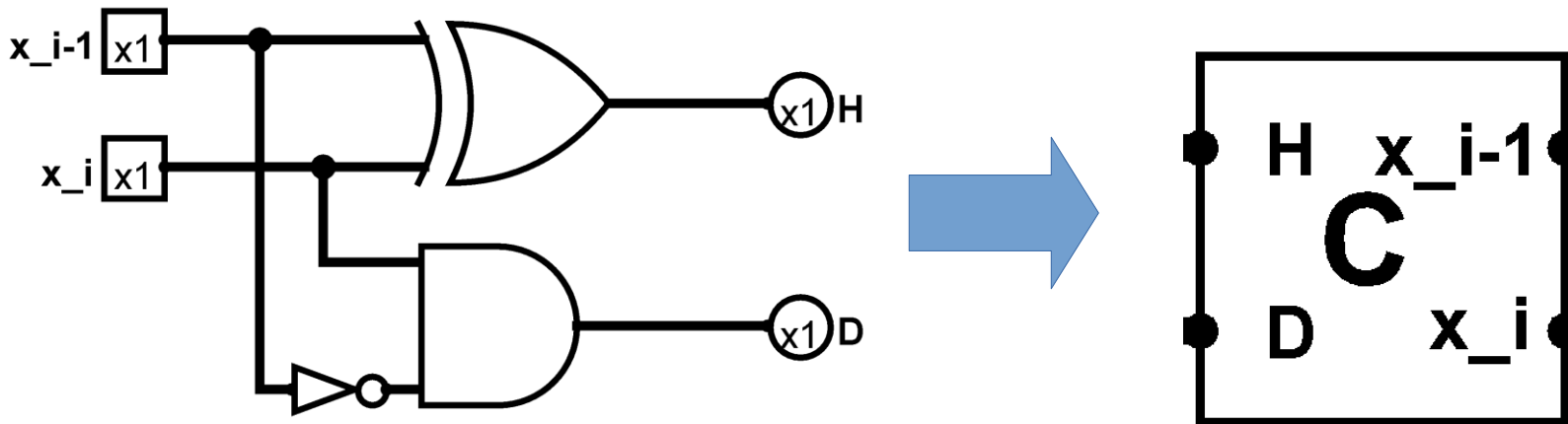
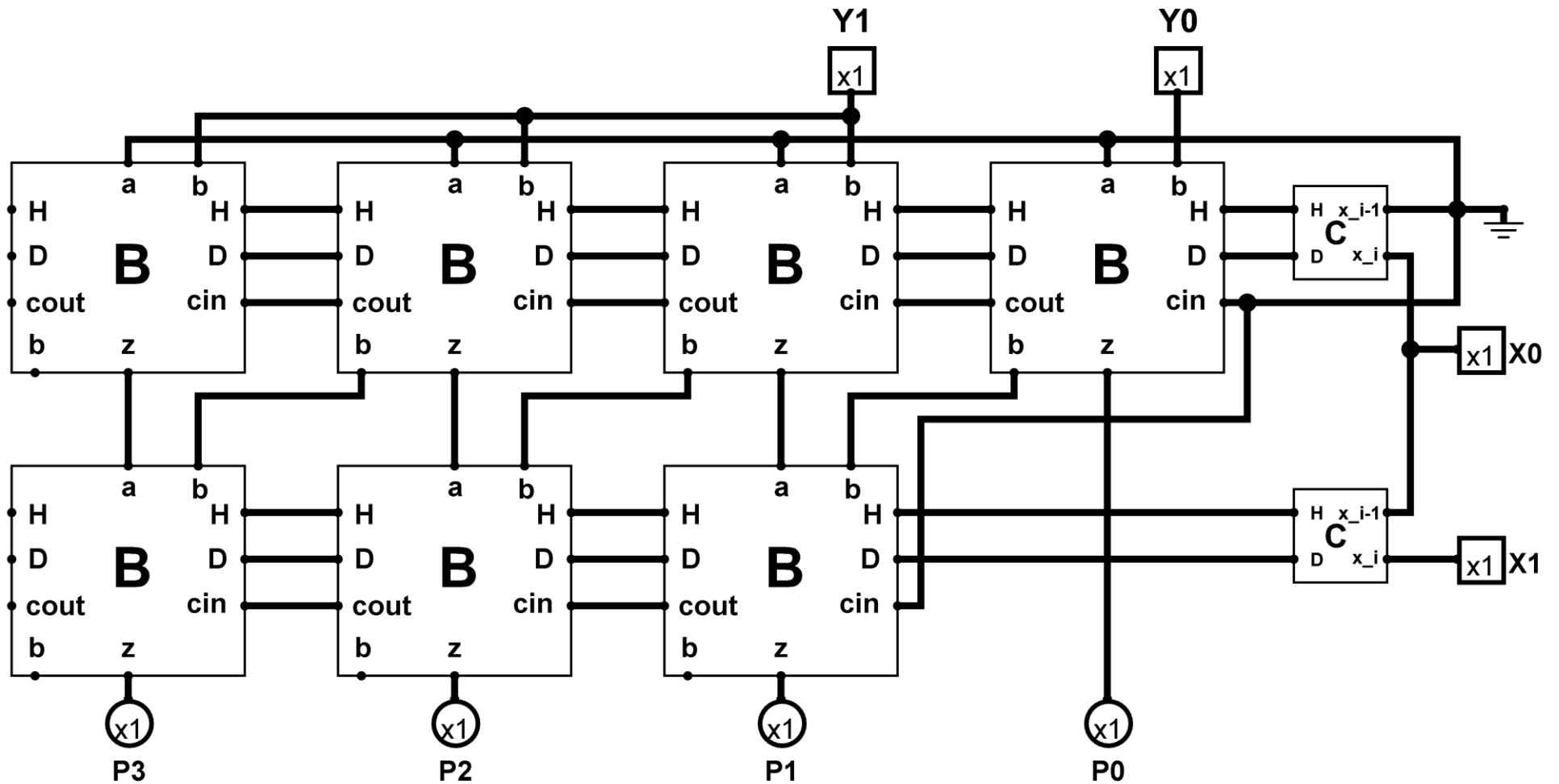


Figure: Cell C

Figure: 2 *2 Booth Multiplier



Booth Multiplier Simulation

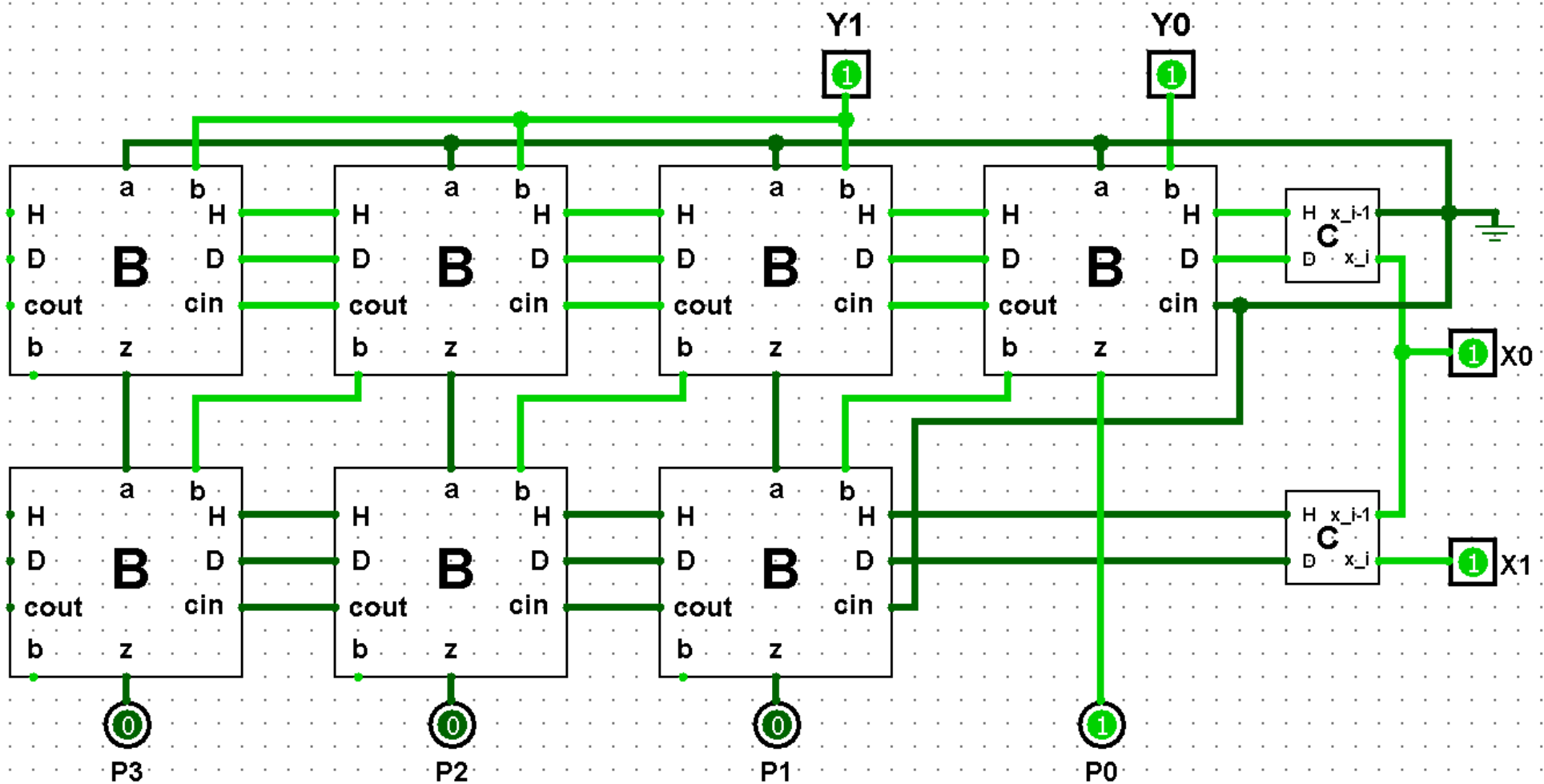


Figure: 2 * 2 Booth Multiplier Simulation for input X = 11 (-1) and Y = 11 (-1). Output is P = 0001 (+1)

Booth Multiplier Simulation

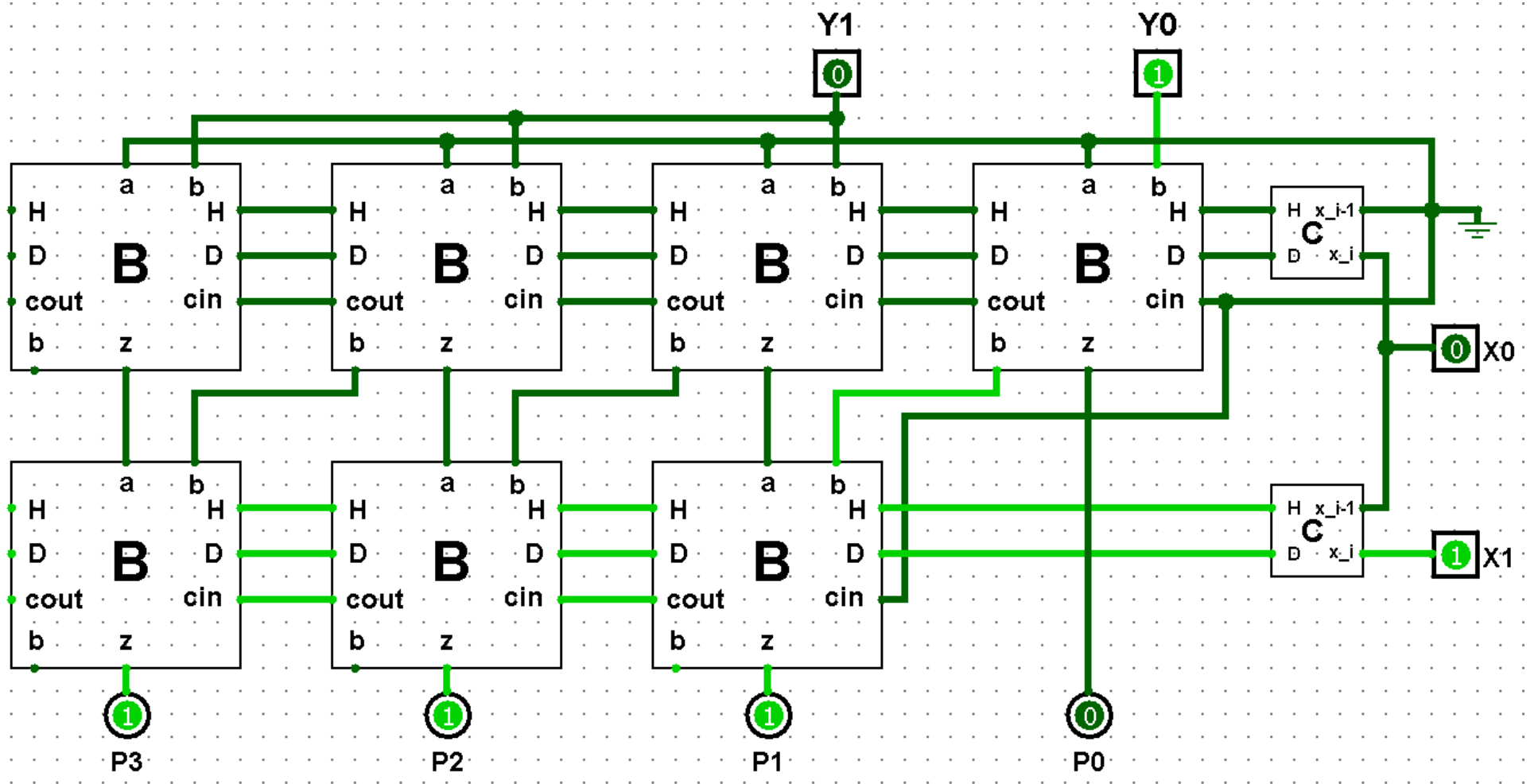
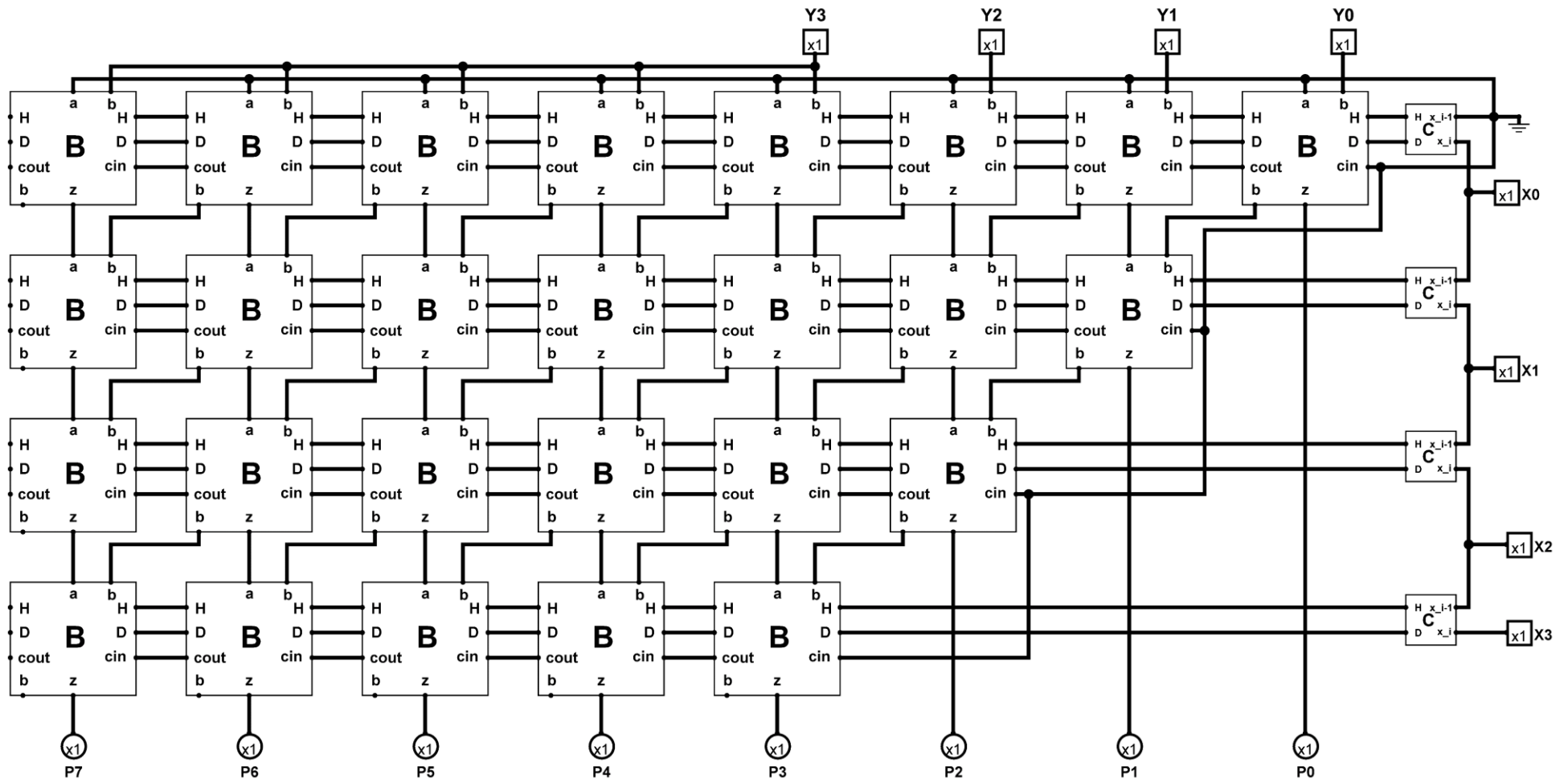
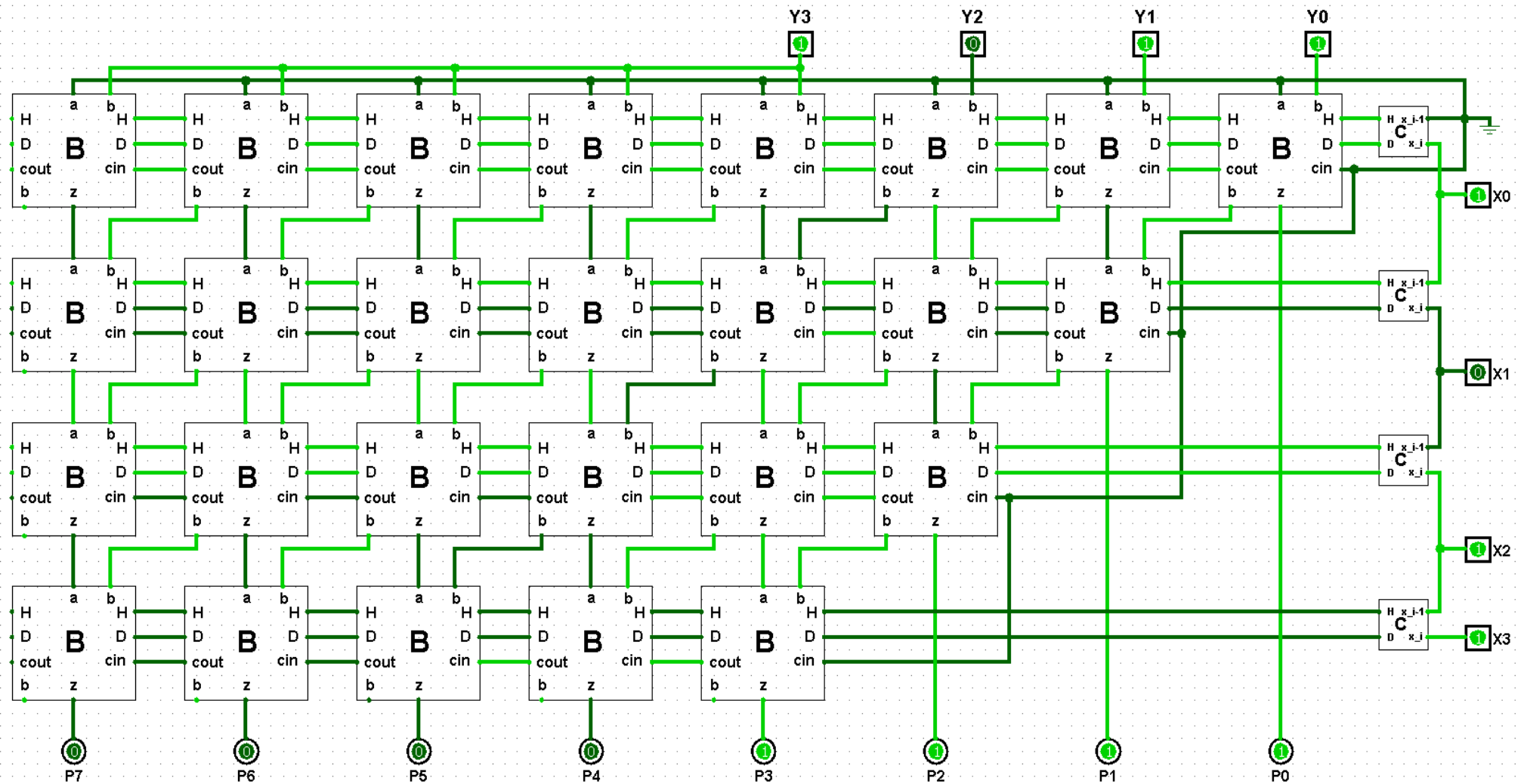


Figure: 2 * 2 Booth Multiplier Simulation for input X = 10 (-2) and Y = 01 (+1). Output is P = 1110 (-2)

Figure: 4*4 Booth Multiplier

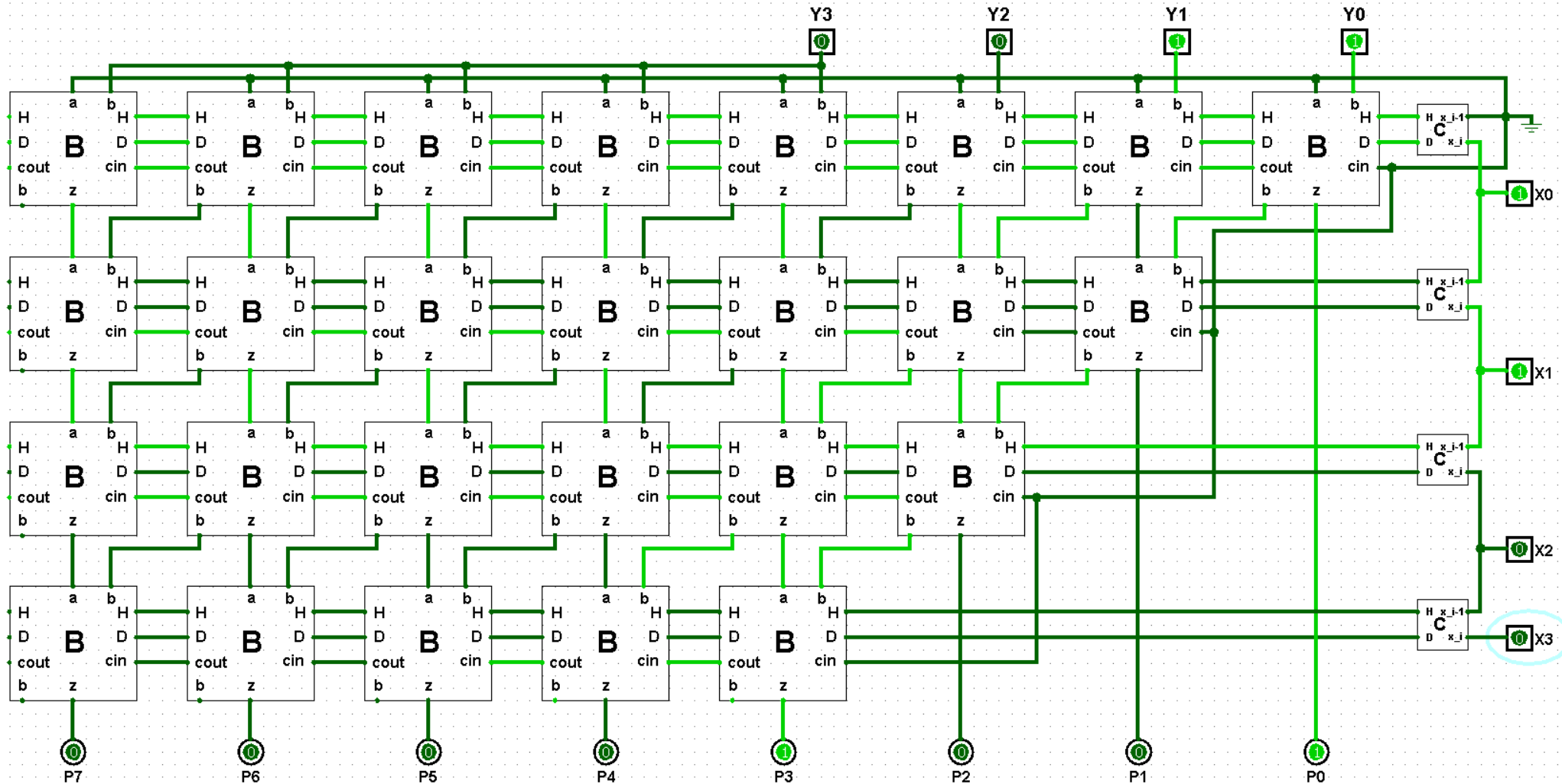


Booth Multiplier Simulation



**Figure: 2 * 2 Booth Multiplier Simulation for input X = 1101 (-3) and Y = 1011 (-5).
Output is P = 00001111 (+15)**

Booth Multiplier Simulation



**Figure: 2 * 2 Booth Multiplier Simulation for input X = 0011 (+3) and Y = 0011 (+3).
Output is P = 00001001 (+9)**

Exercises

1) Calculate multiplication of two numbers using booth algorithm when

- $X = 10$ and $Y = 11$ **OR** $X = 111$ and $Y = 100$ **OR** $X = 1001$ and $Y = 1111$ **OR**
- $X = 10$ and $Y = 1111$ **OR** $X = 1110$ and $Y = 1$ **OR**
- $X = -2$ and $Y = -3$ **OR** $X = +2$ and $Y = -3$

2) Multiply two signed (2's complement) binary numbers $1001 * 1010$ and design a circuit which can calculate this.

3) How does your computer do multiplication in program statement,

$Z = -X * Y$ or $Z = -2 * -3$ or $Z = 1001 * 1010$ (both are signed binary numbers).

Design a circuit and show how it calculates the result in each component.

4) Design a 2/3/4 bit signed/booth multiplier and show output of each circuit when

- $X = 10$ and $Y = 11$ **OR** $X = 111$ and $Y = 100$ **OR** $X = 1001$ and $Y = 1111$ **OR**
- $X = 10$ and $Y = 1111$ **OR** $X = 1110$ and $Y = 1$ **OR**
- $X = -2$ and $Y = -3$ **OR** $X = +2$ and $Y = -3$