

2nd cycle (c Day)

JK

Ref.

Propositions

Boolean

$P \rightarrow$ Today is Friday

$\neg P \rightarrow$ Today is not Friday

* Disjunction ($P \vee q$)

* Conjunction ($P \wedge q$)

* Exclusive or ($P \oplus q$) $\rightarrow \overline{P}q + \overline{q}P$

$$\frac{F \cdot T}{F} + \frac{F \cdot T}{F} = F$$

$$F + T \rightarrow T$$

P	q	$P \oplus q$
T	T	F
F	T	T
T	F	T
T	T	F

Premise

↑ hypothesis

$P \rightarrow q \rightarrow$ consequence

↓ Conclusion

If in T \rightarrow F out then only F

- otherwise always T.

If it is sunny today, then we will go to beach

P

T \rightarrow F then it's F otherwise its true.

You can access the internets from
campus only if you are a
computer science major or you are not
a freshman.

$$a \rightarrow (cvt_{\text{fresh}})$$

(p v q) - without p
(p ^ q) - without p

$$q \bar{p} + p \bar{q} \leftarrow (p \oplus q) \text{ is unboxed}$$

$$\begin{array}{c} T: F \\ F: T \\ \hline F: T + F \end{array}$$

		p \oplus q	1	0	1
		0	1	0	1
		1	0	1	0
0	0	0	0	0	0
0	1	1	1	0	1
1	0	1	0	1	0
1	1	0	1	1	0

grouped $\rightarrow p + q$

standard \checkmark

T also want two T $\rightarrow T \oplus T$

T again standard

so

load of up the multibit you need it in the

segment T standard T \oplus T

so far 241

Lab JK

2nd cycle (1) day

4.4.10

Propositional Equivalences.

- ① * Tautology \rightarrow Every Statement is always alltime True.
- ② A Contradiction
- ③ * Contingency

①	P	$\neg P$	$P \vee \neg P$	$P \wedge \neg P$
	T	F	T	F
F	T	T	F	

Bi-conditional

P	q	$P \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

Logical Equivalences

$$(P \rightarrow q) = = (\neg P \vee q)$$

P	q	$P \rightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

P	q	$\neg P$	$\neg P \vee q$
T	T	F	T
T	F	F	F
F	T	T	T
F	F	T	T

— — — — —

$$P \vee (q \wedge r) \Leftrightarrow (P \vee q) \wedge (P \vee r)$$

$$P \vee (q \wedge r) \Leftrightarrow (P \vee q) \wedge (P \vee r)$$

P	q	r	$P \vee (q \wedge r)$	$(P \vee q) \wedge (P \vee r)$
T	T	F	F	F
T	F	F	F	F
F	T	F	F	F
F	F	T	T	T
F	F	F	F	F

$$\neg(P \vee (\neg P \wedge q)) \Leftrightarrow (\neg P \wedge \neg q)$$

$$\neg P \wedge \neg(\neg P \wedge q)$$

$$\neg P \wedge (P \vee \neg q)$$

$$\Rightarrow (\neg P \wedge P) \vee (\neg P \wedge \neg q)$$

$$\Rightarrow F \vee (\neg P \wedge \neg q)$$

$$= \neg P \wedge \neg q$$

Predicates & Quantifiers

$P(x)$: x is greater than three

$$P(x) : x > 3$$

Universe of dis courses

Domain

$$P(n, y), \quad n = y + 3$$

$$P(1, 2)$$

$$P(5, 2)$$

* Universal Quantification:

* Existential

$P(x) \rightarrow$ true (for all $\rightarrow \forall x$) int the U of D .

$$x < x + 1 \quad x < 0$$

Every student in dis. this class has studied

Calculus.

$P(x) \rightarrow x$ has studied calculus.

U of $D \rightarrow$ All the students in this class

$\forall x P(x)$

$\cup \text{of } D = \{1, 2, 3, 4\}$

$\forall x \in P(x)$

$\exists x P(x)$

$\forall P(1) \wedge P(2) \wedge P(3) \wedge P(4)$

$\exists P(1) \vee P(2) \vee P(3) \vee P(4)$

$\forall x (C(x) \vee \exists y (CC(y) \wedge F(x, y)))$

$C(x) \rightarrow x \text{ has a computer}$

$F(x, y) \rightarrow x \text{ & } y \text{ are friend}$

$\cup \text{of } D \rightarrow (x, y) \in \{\text{all student in the class}\}$

1. All lions are fierce

2. Some lions do not drink coffee.

3. Some fierce creatures do not drink coffee.

$P(x) \rightarrow n$ is a lion

$Q(n) \rightarrow n$ is fierce

$R(x) \rightarrow n$ drink coffee

$$\forall_n (P(n) \rightarrow Q(n))$$

$$\exists_n (P(n) \wedge \neg R(n))$$

3rd cycle A day

6. 4. 10

Sets

$$A = \{1, 2, 3, 4, 5\}$$

elements
number

$$R = \{x \mid x \text{ is a real number}\}$$

$$A \subseteq B$$

$$\forall x \{x \in A \rightarrow x \in B\}$$

$S \rightarrow$ finite set

$n \rightarrow$ distinct element

cardinality of S

$$A = \underbrace{\{1, 3, 5, 7, 9\}}_{\begin{matrix} < 10 \\ n \text{ distinct element} \end{matrix}} = 5$$

$$|A| = 5$$

Identity law, deviation law.

$$\overline{A \cap B} = \overline{A} \cup \overline{B}$$

$$\overline{A \cap B} = \{x \mid x \notin A \cap B\}$$

$$= \{x \mid \neg(x \in A \wedge x \in B)\}$$

$$= \{x \mid x \notin A \vee x \notin B\}$$

$$= \{x \mid \overline{A} \cup \overline{B}\}$$

- o -

$$\overline{A \cup (B \cap C)} = (\overline{C} \cup \overline{B}) \cap \overline{A}$$

$$\overline{A \cap B}$$

$$= \overline{A} \cup \overline{B}$$

$$\overline{A \cup B} = \overline{B} \cup \overline{A}$$

$$\overline{A} \cap (\overline{B \cap C})$$

$$= \overline{A} \cap (\overline{B} \cup \overline{C})$$

$$= (\overline{C} \cup \overline{B}) \cap \overline{A}$$

JL

Functions

3rd cycle C day

One to One

Onto

One to one correspondence

1 to 1 plus onto, bijection.

* Inverse function

* Invertible \rightarrow 1-to-1 comes

* Not " \rightarrow not 1-to-1 "

Composition of function:

$$(f \circ g)a = f(g(a))$$

$$f(x) = 2x + 3$$

$$g(x) = 3x + 7$$

$$(f \circ g)x = f(g(x))$$

$$= f(3x + 7)$$

$$= 2(3x + 7) + 3$$

$$= 6x + 17$$

$$\begin{aligned}
 (g \circ f)(x) &= g(f(x)) \\
 &= g(2x+3) \\
 &= 3(2x+3)+7 \\
 &= 6x+16
 \end{aligned}$$

$\Sigma x 17.$

$$\begin{aligned}
 * \text{floor } [3.4] &= 3 \\
 * \text{Ceiling } \lceil 3.7 \rceil &= 4 \\
 \lceil 3.4 \rceil &= 4
 \end{aligned}$$

A. 7. Sequences & Sum Self.

let,

$$y = \sqrt{3-x}$$

$$x = \{$$

1
2
3

$$y = ax^n + bx + c$$

Read, a, b, c.

put value of x

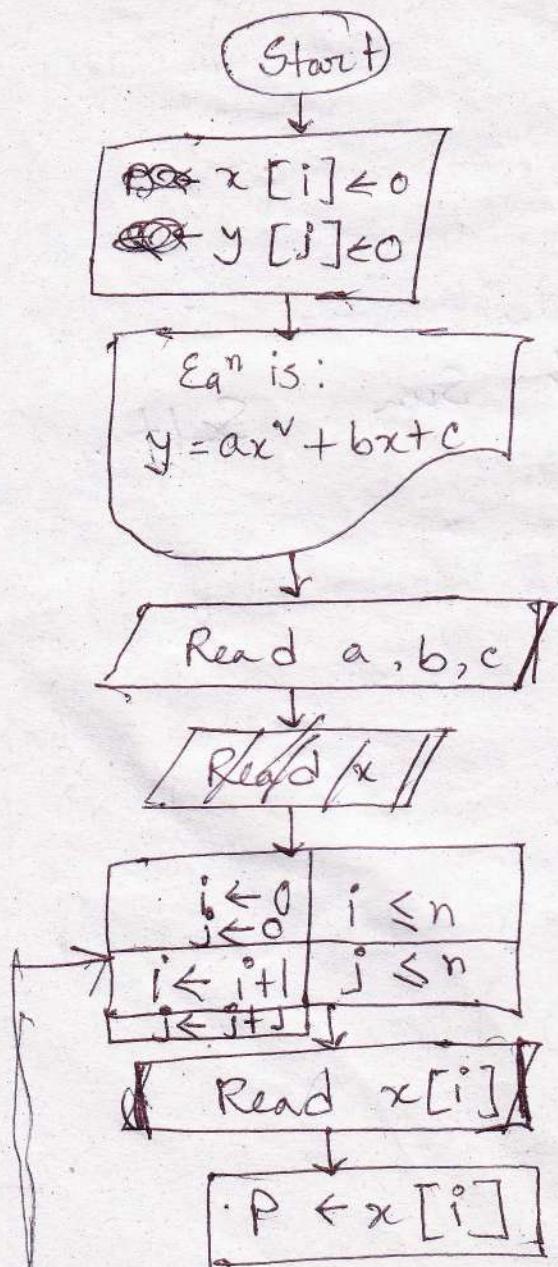
$$y = \boxed{x^{\frac{n}{2}} + 1}$$

$$y =$$

$$x[2] \rightarrow 2$$

$$y[]$$

$$y = a x^n + b x + c$$



$$y[j] = a \cdot x[i] \cdot x[i] + b \cdot x[i] + c$$

$$y[i] \leftarrow y[i+1]$$

$$x[i] \leftarrow x[i+1]$$

$$x[i] : y[i]$$

one to one

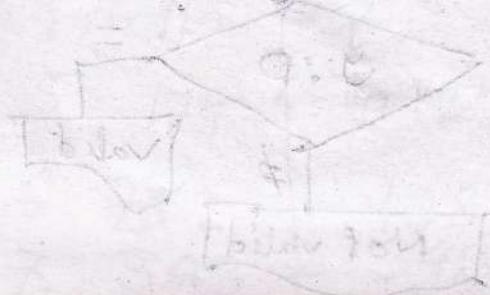
not one to one

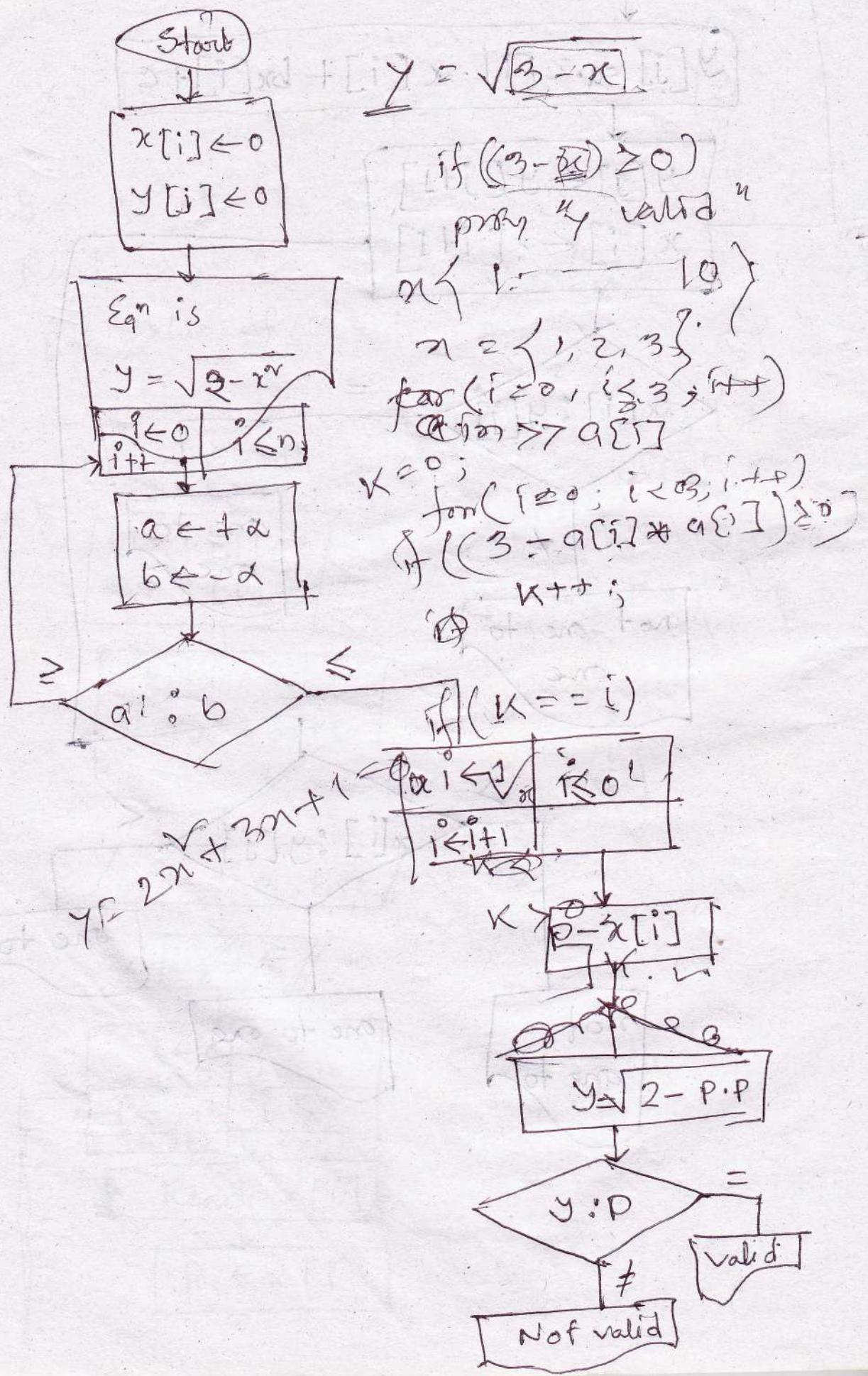
$$x[i] : y[i]$$

one to

not one to

one to one





9 cycle C Day

MJK
Cot

1.8 growth of f^n .

$$f(n) = x^n + \dots \approx ax^n + bx + c.$$

Big O Notation

$|f(x)| \leq c|g(x)|$
 $x > k$ $f(x)$ is big-O of $g(x)$

pair (c, k)

$$f(n) = x^n + 2x + 1$$

$O(n^k)$

$$0 \leq x^n + 2x + 1$$

$x > 1$

$$x^n + 2x^n + x^n = 4x^n$$

$$c = 4$$

$$k = 1$$

$(4, 1)$

$(3, 2)$

$$7x^n \rightarrow O(n^3)$$

	$7x^n$	x^3
$n=1$	7	1
2	28	8
3	63	27
4	112	64
5	175	125
6	252	198
7	343	343

$$x = 8 - k =$$

$$x^3 > 7x^n$$

$$x > 7$$

Why $O(7x^n) \rightarrow x^3$

$$\begin{aligned}
 f(x) &= |a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0| \\
 &\leq |a_n x^n| + |a_{n-1} x^{n-1}| + \dots + |a_1 x| + |a_0| \\
 &= x^n \left| a_n + a_{n-1} \frac{x}{x^n} + \dots + \frac{a_1}{x^{n-1}} \right| + |a_0| \\
 &\leq x^n \left| a_n + a_{n-1} \frac{x}{x^n} \right| \geq \dots
 \end{aligned}$$

$$6! = 6 \cdot 5 \cdot 4 \cdot 3 \cdot 2 \cdot 1$$

$$n! = 1 \cdot 2 \cdot 3 \cdots n$$

2 3 2 . . . 2

$$= n^n$$

$$e^1 = e^6.$$

— 8 —

Fig. 1. A photograph of the head of a female *Leucostethus* with the mouthparts removed.

卷之三

18. 19. 20.

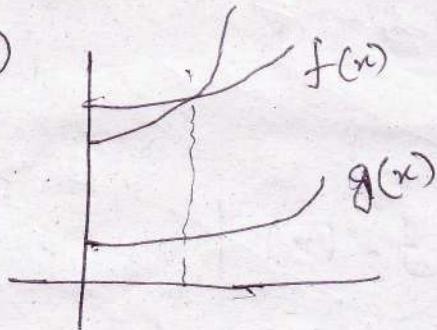
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S cycle C Day

$$f(x) \leq c(g(x)) \quad O(g(x))$$

$x > k$

(c, k)



$$(c, k) \rightarrow O(g(x))$$

$$f_1(x) + f_2(x) \rightarrow O(g(x))$$

$$f_1(x) f_2(x) \rightarrow O(g(x))$$

$$x^n + 3 \rightarrow O(2x^n)$$

$$|g| = 1 \quad 4 \quad 2$$

$$|c=2| \quad 7 \quad 8$$

$$|c=3| \quad 12 \quad 18$$

$$\begin{matrix} 1+2 \\ = 3 \end{matrix} < 4$$

$$\begin{matrix} 1+2+3 \\ = 6 \end{matrix} < 9$$

$$1+2+3+4 <$$

$$10 < 16$$

$$\textcircled{O} \quad 1+2+\dots+n \leq n^2$$

$$1 \leq 1^2$$

$$x^3 + 3 \textcircled{O} \quad O(x^3)$$

$$\begin{aligned}
 f_1(x) + f_2(x) &\leq c_1 |g_1(x)| + c_2 |g_2(x)| \\
 &\leq c_1 |g(x)| + c_2 |g(x)| \\
 &= g(x) (c_1 + c_2) \\
 &= C |g(x)|
 \end{aligned}$$

$$g(x) = \max(g_1(x), g_2(x))$$

$$\begin{aligned}
 |f_1(x) + f_2(x)| &= |f_1(x)| |f_2(x)| \\
 &\leq |c_1 g_1(x)| |c_2 g_2(x)| \\
 &\leq c_1 c_2 |g_1(x)| |g_2(x)| \\
 &\leq C |g(x)|
 \end{aligned}$$

$$f(x) = 3x \log x! + (x^n + 3) \log x$$

$$3x \log x!$$

$$3x \log x^*$$

$$3x^n \log x$$

$$(x^n + 3) \log x$$

$$x^n \log x$$

$$g(x) = x^n \log x$$

$$f(x) = (x+1) \log (x^n + 1) + 3x^n$$

$$x^n + 1 \leq 2x^n$$

$$\log (2x^n)$$

$$= \log 2 + \log x^n$$

$$= \log 2 + 2 \log x \leq 3 \log$$

$O(x) \rightarrow f(x) \leq c|g(x)|$ upper limit

$\Omega(x) \rightarrow f(x) \geq c|g(x)|$ lower limit

$\Theta(x) \rightarrow = \leq f(x) \leq \leq$

$$(3x+7) < x$$

$$-13x < -7$$

$$3 < 1$$

$$6 < 2$$

$$9 < 3$$

$$12 < 4$$

$$1 \textcircled{b} f(x) = 3x + 7$$

$$3x + 7 \leq x$$

$$x=1$$

$$x=2$$

$$x=3$$

$$3x + 7 \cancel{\leq x} = 10 \cancel{\leq} \\ \cancel{10} \cancel{\leq} 13 \cancel{x} \\ 16 \cancel{\leq} 1$$

So $f(x)$ is not $\mathcal{O}(x)$.

$$1. \textcircled{c} x^2 + x + 1 = f(x)$$

$$x^2 + x + 1 \leq x$$

Not $\mathcal{O}(x)$

$$\textcircled{a} f(x) = 10$$

$$10 \leq x$$

$$x=11, k=10, c=1$$

$$\textcircled{b} f(x) = 3x + 7$$

$$0 \leq 3x + 7 \leq 3x + x = 4x \rightarrow c=4$$

$$3x + 7 \leq x$$

$$k=7$$

$$① f(x) = \log x$$

$$0 \leq \log x$$

$$x=1, \log x = 0 \leq \log x$$

$$K=1.$$

$$② f(x) = \lfloor \log x \rfloor$$

$$c=1, K=0.$$

$$1) f(x) = \left\lceil \frac{x}{2} \right\rceil \\ = 1,$$

$$2) f(x) = x^n + 1000$$

$$0 \leq x^n + 1000$$

$$0 \leq x^n + x^n = 2x^n$$

$$c=2$$

$$x^n + 1000 \leq x^n$$

~~Not O(n)~~

$$K=0.$$

$$③ f(x) = x \log x \\ = \log x^x$$

$$c=1$$

$$\log x^x < n$$

$$x=1 \quad 0 \leq 1$$

$$x=2 \quad 0.60 \leq 2$$

$$x=3 \quad 1.43 \leq 3$$

$$c=1, k=3$$

$$2d) f(x) = 2^x$$

$$c=1$$

$$2^x \leq x^v$$

$$x=1 \quad 2 \leq 1$$

$$x=2 \quad 4 \leq 4$$

$$x=3 \quad 8 \leq 9$$

$$c=1, k=3,$$

$$3. f(x) = x^4 + 9x^3 + 4x + 7$$

$$x^4 + 9x^3 + 4x + 7 \leq 9x^4$$

$$c=9$$

$$k=9$$

$$4. f(x) = 2^x + 17$$

$$2^x + 17 \leq 3^x$$

$$K = 3, C = 1$$

$$5. f(x) = \frac{x^n + 1}{x + 1} \quad C = 1$$

$$\frac{x^n + 1}{x + 1} \leq x$$

$$x = 1, \quad x^n + 1 \leq x^n + x$$

$$K = 1$$

$$6. f(x) = \frac{x^3 + 2x}{2x + 1}$$

$$\frac{x^3 + 2x}{2x + 1} \leq x^n$$

$$C = 1, \quad K = 3$$

6 cycle C Day

Seq Search (L, r, x) Indexd
 L is an Array, n entries, $x \rightarrow$ key to be
 searched

O/P: if x is in L , then O/P is its
 index else O/P 0.

Pseudo code: $index = 1$

while ($index \leq n \ \&\ & L[index] \neq x$)

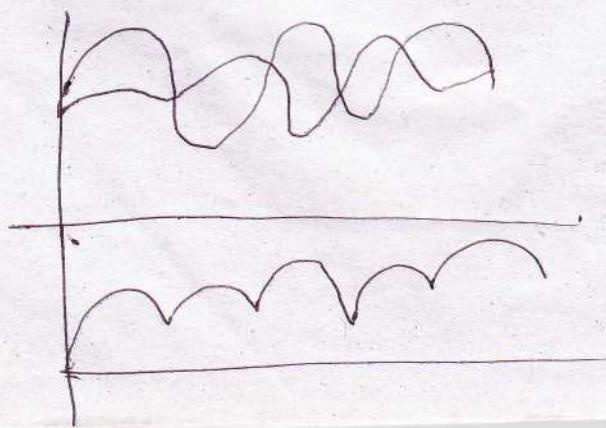
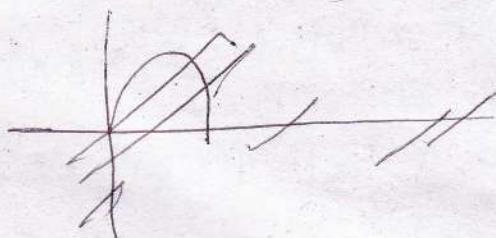
$index = index + 1$;

 if ($index > n$), then $index = 0$

 return $index$;

* Time
 * Space

big - Oh \rightarrow worst
 case



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$$f(x) \geq c g(x)$$

$$x > k$$

$$f(x) = 3x^{\alpha} - 3x - 5 \quad g(x^{\alpha})$$

$$3x^{\alpha} - 3x - 5 > x^{\alpha}$$

$f(x)$	$g(x^{\alpha})$
0 → -5	0
1 → -5	1
2 → 1	4
3 → 13	9
4 → 31	16

$$0 \leq c_1 g(x) \leq f(x) \leq c_2 g(x)$$

* Algorithm

* Computational Complexity

→ Time (Matrix)

→ Space (Database)

Logic 1.1

Propositional Equiv 1.2

Predicates 1.3

Set 1.4

Set op 1.5

Functions 1.6

Seq & Seq 1.7

Growth of f^n 1.8

7 cycle C Day

30.05.10

+ 2.2 complexity of alg.

2.1 Algorithms

LSA (Linear Search Alg.)

Procedure : linear search (x : int, a_1, a_2, a_n : distinct integer)

while $i = 1$ ($i \leq n$ and $x \neq a_i$)

$i = i + 1$
if ($i \leq n$) then location = i

else location = 0.

$i = 2i + 1$

I/P

Output or Time $2n+2$

O/P

correctness

Binary Search ~~algo~~ $\log_2 n$
+ Complexity

finiteness

Average Case

2.3

Prime

$P > 1$

$P \rightarrow P, 1$

a, b, c $b = ac$ \Rightarrow remainder
factor a/b multiple of a
 a/b \rightarrow integers

1. a/b & a/c then $a/(bc)$
2. a/b the a/bc & c
3. a/b & b/c then a/c .

* Fundamental theorem of Arith

* integer \rightarrow product of primes

8cycle B Day (15.01.01) GCD

Primes :

$$* P > 1$$

Positive factor of Primes $\{ \}$ & P

* Composite [Product of 2 Prime numbers]

$\downarrow n \rightarrow$ Prime division $\leq \sqrt{n}$

* Mersenne primes

$$a = d q + r$$

↓ ↓ ↓

dividend division quotient

remainder (0.01, 001)

M.D.

$2^P - 1 = () \rightarrow$ Prime হায়তে
Otherwise not

3900

GCD

LCM

(24, 34)

$(abm)_{H.C.F} \equiv P.S$

$$24 = 2^3 \cdot 3 = 12 \cdot 2$$

$$36 = 2^2 \cdot 3^2 = 12 \cdot 3$$

$$\boxed{S = 2 - f}$$

$$GCD(17, 22) = 1$$

✓
relatively
Prime

$$GCD(17, 34)$$

$\neq R.P.$

$$\text{GCD}(10, 17, 21) \stackrel{\text{def}}{=} R.P$$

$$120 = 2^3 \cdot 3 \cdot 5$$

L < 9 *

$$500 = 2^2 \cdot 5^3 \rightarrow P_1 = 2$$

$$P_1 = \min(3, 2) = 2, P_2 = \min(1, 3) = 1$$

$$2^2 \cdot 5^1$$

$$= 20$$

$$\text{GCD}(500, 120) = 20$$

L C M

Modular Arithmetic:

$$\begin{aligned} r &= a \pmod{m} & m &\mid a - b \\ \text{Congruent} \quad a &\equiv b \pmod{m} & \rightarrow & \\ \downarrow & a \pmod{m} = b \pmod{m} & & \\ \text{Ex. } 17 &\equiv 5 \pmod{6} & & \text{MOD} \\ 24 &\equiv 14 \pmod{6} & & \end{aligned}$$

$$\boxed{17 - 5 = 12} \quad \begin{array}{l} S \cdot S_1 = 8 \cdot 8 = 64 \\ E \cdot S_1 = 8 \cdot 5 = 40 \end{array}$$

$$(64, 40) \text{ GCD}$$

9. 9 *

$$L = (64, 40) \text{ GCD}$$

Method
ans

Goyal B Day

$$c = (P+K) \bmod 26$$

$$\therefore = \bmod (26 + \underline{\underline{?}})$$

$0 \rightarrow 9$

$$r = (\underline{\underline{?}}) \bmod 10$$

$$= (\underline{\underline{?}}) \bmod 11$$

$$a \equiv b \pmod{m}$$

$$m \mid a - b$$

$m = (+) ve$
If & only if $f \rightarrow k$
such that $a = b + km$

$$a \equiv b \pmod{m}$$

$$b \equiv d \pmod{m}$$

$$\text{then, } a+c \equiv b+d \pmod{m}$$

$$\& ac \equiv bd \pmod{m}$$

$$a = b + km$$

$$c = d + tm$$

$$\frac{(a+c)}{n} \doteq \frac{(b+d)+tm}{y} \frac{(k+t)}{z}$$

$$x = y + zm$$

The Euclidean Alg.

$$\gcd(91, 287)$$

$$287 = \frac{a}{3} \cdot \frac{b}{91} + \underline{\underline{r}}$$

$$\rightarrow \gcd(91, 14)$$

$$91 = 14 \cdot 6 + \underline{\underline{7}}$$

$$\gcd(7, 14)$$

$$14 = 7 \cdot 2$$

$$\text{calculate } \gcd(91, 287) = \gcd(91, 14) = \gcd(7, 14) = 7$$

$$a = bq + r$$

$$a, b, q, r = \text{int}$$

$$\gcd(a, b) = \gcd(b, r)$$

$$a \geq b$$

$$a = bq + r_2 \quad 0 \leq r_2 < b$$

$$b = r_2 q_2 + r_3 \quad 0 \leq r_3 < r_2$$

$$r_2 = r_3 q_3 + r_4 \quad 0 \leq r_4 < r_3$$

$$0 \leq r_2 < r_{n-1}$$

$$r_{n-2} = r_{n-1} q_{n-1} + r_n$$

$$r_{n-1} = r_n q_n$$

$$10^2 \quad 10^1 \quad 10^0 \\ 2 \quad 8 \quad 7$$

$$n = a_k b^k + a_{k-1} b^{k-1} + \dots + a_0 b^0$$

a_i, b are integers & $b > 1$

b = base of no.

Chapter 2.5

$$a, b = (+) \text{ or } (-)$$

$$\gcd(a, b)$$

$$sa + tb$$

$$\gcd(252, 198) = 18$$

Linear Combination of 252 & 198

$$252 = 1 \cdot 198 + 54$$

$$198 = 3 \cdot 54 + 36$$

$$54 = 1 \cdot 36 + 18$$

$$36 = 2 \cdot 18$$

→ next to last

$$18 = 54 - 3 \cdot 18$$

$$36 = 198 - 3 \cdot 54$$

$$18 = 54 - 1(198 - 3 \cdot 54) = 4 \cdot 54 = 1 \cdot 198$$

$$54 = 252 - 1 \cdot 198$$

$$\text{Finally } 18 = 4(252 - 1 \cdot 198)$$

$$= 4 \cancel{252} - 4 \cdot 198$$

$a, b, c = (+) \text{ ve int}$

$\gcd(a, b) = 1 \quad \& \quad a/bc \text{ then } a/c$

Theory $m = (+) \text{ ve int}, (a, b, c) = \text{int}$

if $ac \equiv bc \pmod{m} \quad \& \quad \gcd(c, m) = 1$

then $a \equiv b \pmod{m}$

$$\rightarrow m/a(b-a)$$

R.P.C

$$m/(a-b)$$

(\rightarrow to simplified ratio)

$$P^2 + 8C1 \cdot 1 = 525$$

$$P^2 + P^2 \cdot 8 = 8P^2$$

$$31 + 3E \cdot 1 = P^2$$

$$31 \cdot 5 = 3E$$

sum of these

$$1 \cdot 5E + 1 \cdot 8 = 31$$

$$P^2 \cdot E - 8C1 = 3E$$

$$8C1 \cdot 1 = P^2 - P(P^2 \cdot E - 8C1) + P^2 = 31$$

$$8C1 \cdot 1 - 525 = P^2$$

9cycle C Day

Linear Congruences

$$ax \equiv b \pmod{m}$$

* $m = (+)$ ve int,

$a, b =$ int

$x =$ variable

$$\bar{a}x \equiv 1 \pmod{m}$$

\bar{a} is an inverse

relatively prime

of a module m

$$a \& m \Rightarrow RP \& (m > 1)$$

i.e. $\gcd(a, m) = 1$

inverse of $a \pmod{m}$

$$sa + tm = 1 \Rightarrow sa + tm \equiv 1 \pmod{m}$$

$$\therefore tm \equiv 0 \pmod{m}$$

$$sa \equiv 1 \pmod{m}$$

So, $\bar{a} = s$ that means inverse of a .

3 mod 7.

3 7

$$\gcd(3, 7) = 1$$

$a \& m \Rightarrow RP \& (m > 1)$ \Rightarrow inverse

not integer.

$$3(-2) + 7 \cdot 1 \Rightarrow 1 \pmod{7}$$

$$3(-2) + 7 \cdot 1 = 1$$

$$\therefore a + tm$$