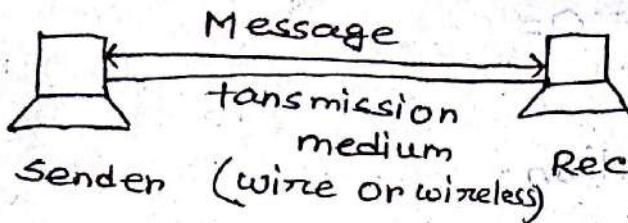


30-04-17  
1st (B) day

## Data communication:-



## Protocol (rules)

IP → Internet protocol

http → hyper text transfer protocol

- 1. Sender
  - 2. Receiver
  - 3. Transmission medium
  - 4. Message (या पाठ्यक्रम, धृत करा)
  - 5. Protocol (किंवद्दनकी नियम)
- 5 घोषणाएँ असंगत  
Data communication

## Transmission Modes:-

1. Simplex
2. Half Duplex
3. Full Duplex

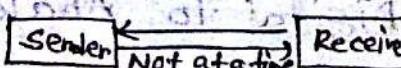
1. Simplex : One way communication



Example: Keyboard and Monitor Relation

Scanner → file

2. Half Duplex : Two way communication (Not at a time)



उदाहरण (Example)

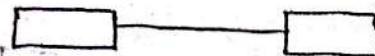
3. Full duplex: Two way communication (at a time simultaneously)

Example: Mobile phone

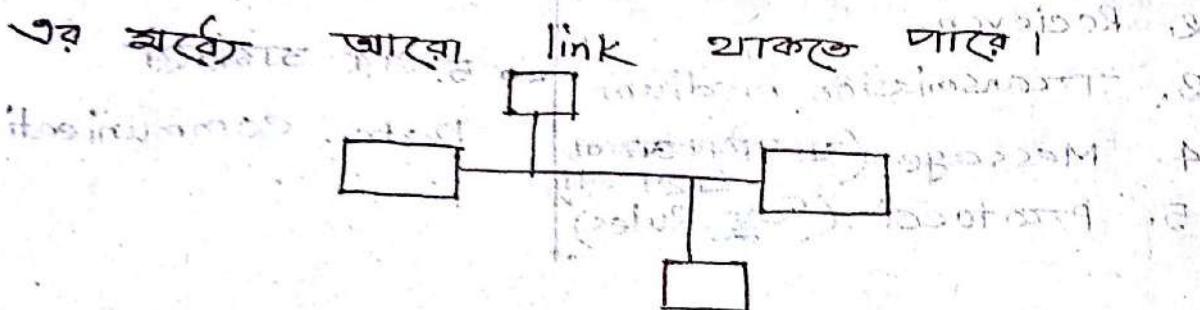
1. Point to point connection:
2. Multipoint connection:

1. Sender & Receiver

आलाप Connection याकर्त्ता | One-to-one link.



2. One to many link. Sender & Receiver



10101  
Digital data

Digital to  
Digital

Digital signal

Analog signal

Analog to  
Digital

10101

Digital data

Analog

Analog

PSK, FSK, ASK → Digital to Analog techniques

## Computer model तकारा

### Systems

#### 1. OSI model (Open Systems Interconnection)

#### 2. TCP/IP

03-05-17  
1st (D) day

Data communication :- Data communications are on the exchange of data between two devices via some form of transmission medium such as a wire cable.

The effectiveness of data communication :-

(i) Delivery

(ii) Accuracy

(iii) Timeliness

(iv) Jitter

→ 4 factors

- (i) The system must deliver data to the correct destination.
- (ii) The system must deliver data accurately.
- (iii) The system must deliver data in a timely manner.
- (iv) Jitter refers to the variation in packet arrival time.

Physical structure :-

Type of connection :-

(i) Point to point

(ii) Wide link → two or more devices share a multipoint connection

**Physical TOPOLOGY:** Refers to the way in which a network is laid out physically. Different way to make network design.

1) Mesh TOPOLOGY

2) star

3) Bus

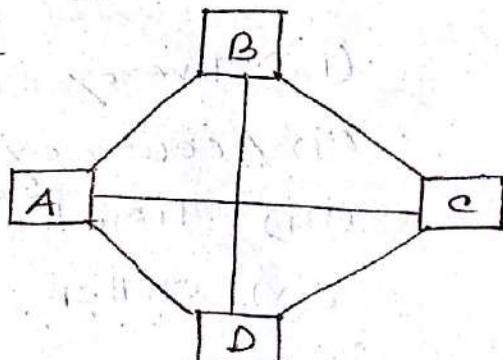
4) Ring

5) Hybrid

1) Mesh TOPOLOGY:-

every device has point to point connection.

$$\text{Here No. of link} = \frac{4(4-1)}{2} \\ = \frac{4 \cdot 3}{2} \\ = 6$$



$$\text{Mesh} = \frac{n(n-1)}{2}$$

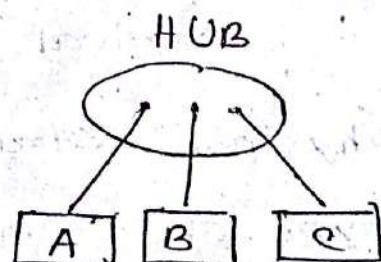
- Storage connection
- costly (because point to point)
- एकले link वर्षे इसे टेलिडिजिटल system दे कर्या ना-
- Data transmission fast होते
- ये link जो नष्ट होते तो detect करा असह

2) star TOPOLOGY:-

✓ Switch वर्षे गांविधर connected

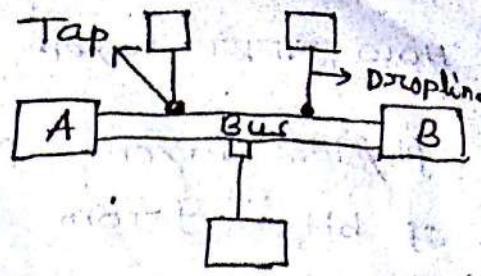
→ time रुपी नागर्वे

→ link कडा नागर्त्तु (costly वर्षे)



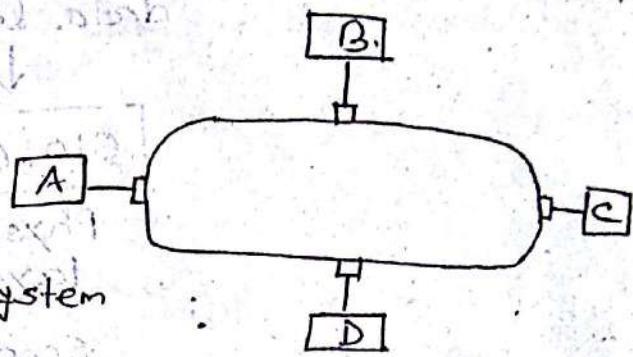
### 3) Bus TOPOLOGY:-

Multipoint connection



### 4) Ring TOPOLOGY:-

We connect hence by link



⇒ प्रकृते link नम्हे इले अवधु System  
link नम्हे इस्तु ना।

⇒ यद्यु link दे नम्हे इस्तु, तो detect करा अट्टा

~~Host~~

06-05-2017  
8 1st (E) day

Network model:-

अवधु basic model हन्न OSI model.

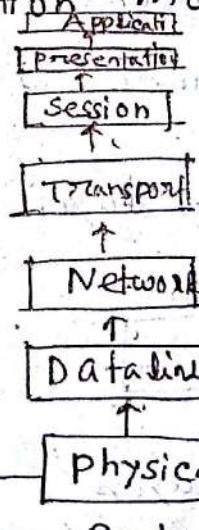
OSI → Open system Interconnection model

7 से layer तेसु हैं,

1. Application
2. Presentation
3. Session
4. Transport
5. Network
6. Datalink
7. Physical

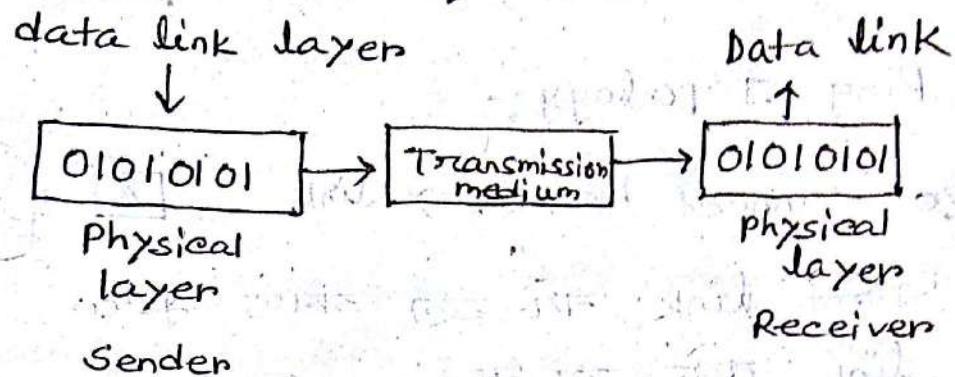
Sender

Transmission medium



Now आवर्णी layer शुल्काय करने हेतु

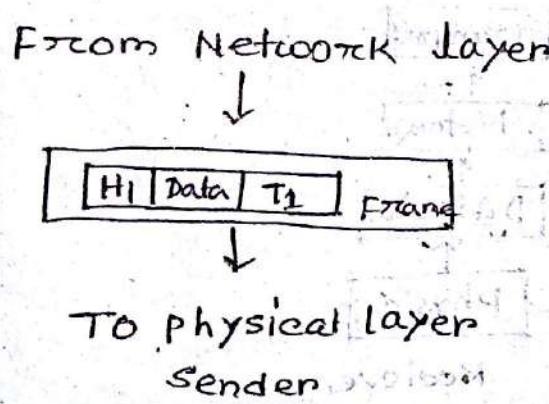
**Physical Layer**: - Physical layer is responsible for movements of bits from one hop (node) to the next.



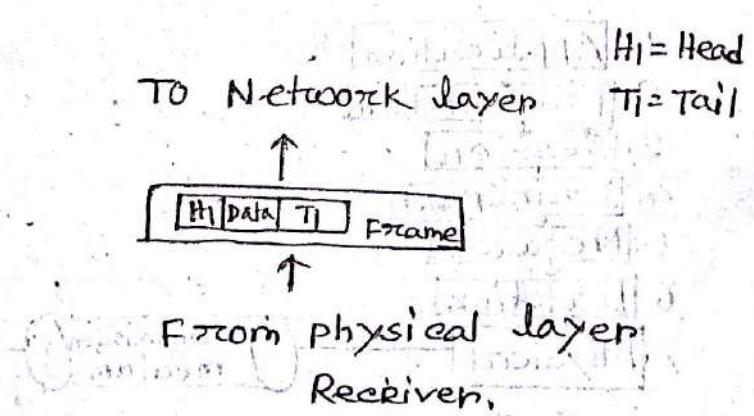
ଆହୋ କାଣ୍ଠ ।



**Data Link Layer:-** Data link layer is responsible for moving frames from one hop (node) to the next.



### Other Responsibilities:-



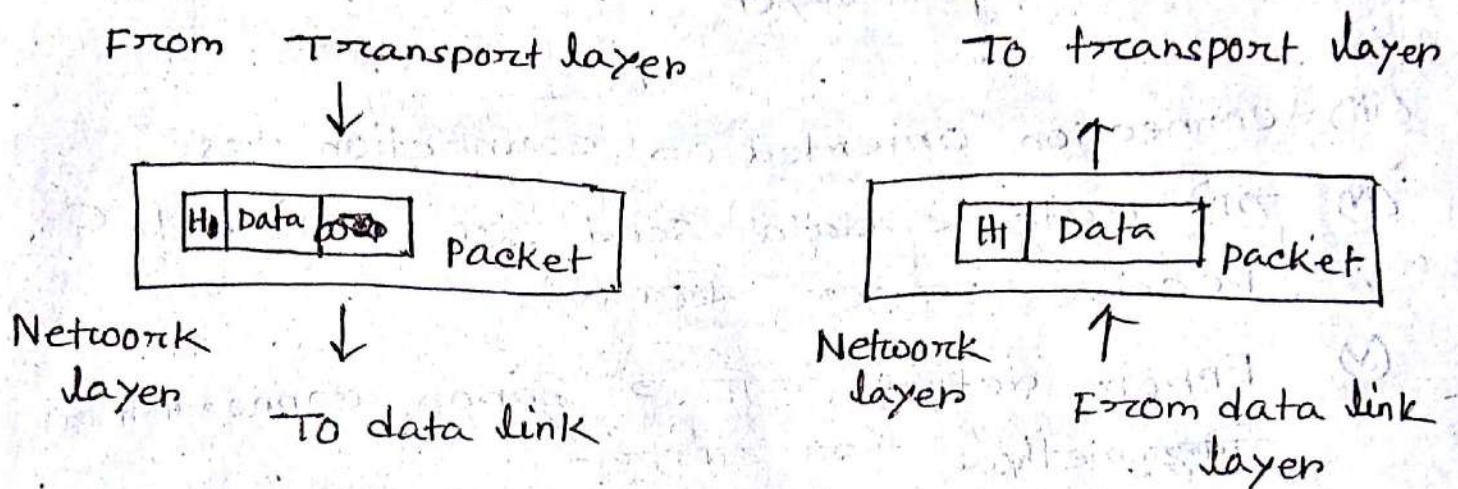
→ Examining

- Access control (इस वाले अपनी PC पर सभी पास्टर Access भिन्न भिन्न same डिस्ट्रिब्यूट करते हैं। जो डिस्ट्रिब्यूट करते हैं।)
- physical Addressing ( $H_1, T_1$  add करता है) उल्लेखित PC पर सभी पास्टर Access भिन्न भिन्न same डिस्ट्रिब्यूट करते हैं। जो डिस्ट्रिब्यूट करते हैं।)
- Flow control (sender की rate का send करते हैं, Receiver की rate का receive करते हैं।)

## Error control

Error control: Error detect तथा, Error correction.

Network layer:- Network layer is responsible for the delivery of individual packets from the source host to the destination host.



Other responsibilities:-

- Logical addressing (प्रायः IP Address)
- Routing (कठमुलो network interconnection रूप से एकत्रितीयास्तुति हस्त)

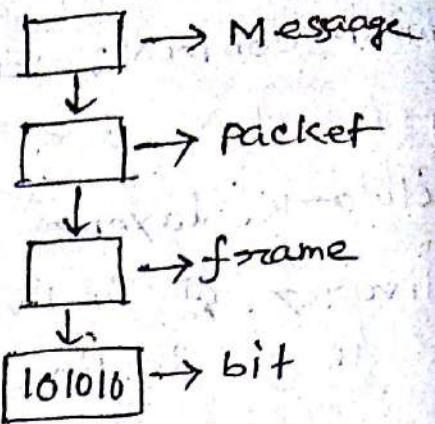
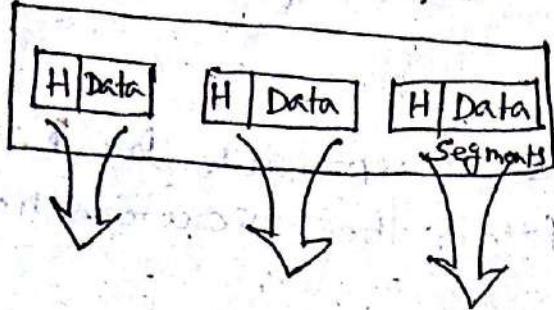
08-05-17  
2nd (B) day

Transport layer:- Transport layer is responsible for the delivery of message from one process to another.

Other responsibilities:-

- (i) service point addressing
- (ii) Segmentation and reassembly (एक message को अलग अलग segment में बदलना तथा उन्हें फिर से जोड़ना)
- (iii) connection control
- (iv) flow control
- (v) Error control.

From session layer



TO Transport layer

- (iii) Connection oriented or connection less
- (iv) Rate of data send तथा रेट रेट  
Receive करने वाले
- (v) Error detect करना & Error correction  
दो पद्धति ग्रे कार्यक्रम

(i) Logical → Network layer जैसा करना  
address } Physical → Data link जैसा करना  
Port → Transport, supply करना  
message जैसा करना authentic करने के लिए  
Port address ग्रे कार्यक्रम।

Session layer:- Session layer is responsible for dialog control and synchronization.

• Dialog control:- The session layer allows two system to enter a dialog. It allows the communication between two processes to take place in either half duplex mode or full duplex mode.

Synchronization: Message टोक्स Segmentation ए पार्श्व-प्रस्तुति  
100 दे पर्याप्त ए पार्श्व करना चाहे। प्रयोग 10 or 20 or  
30 जे यात्राव पर यात्रे वाके शूल वा वाकिलिए  
पाठानाव कल यात्रे हवे। 30 डोर पर यदि न यात्रा  
हवे तुम्हारे पर problem आए। अधन त्रैथान्यात्  
आवाहन send करव।

### Presentation layer:

presentation layer is responsible  
for (1) Translation (2) Compression & (3) Encryption

(1) एक device कर्तव्य operating system वा कर्तव्य  
अधन एवं एक device एक आखरी device  
वा पाठान, compiler त्रितीय operating system  
ज्ञान्यासी translate करवे।

(2) Message टोक्स कर्तव्य zip करवे पाठाने याए।

(3) आवश्यान message टोक्स encrypt करवाए  
जा hack न हो।

### Application layer:

Application layer is responsible  
for providing service to the user.

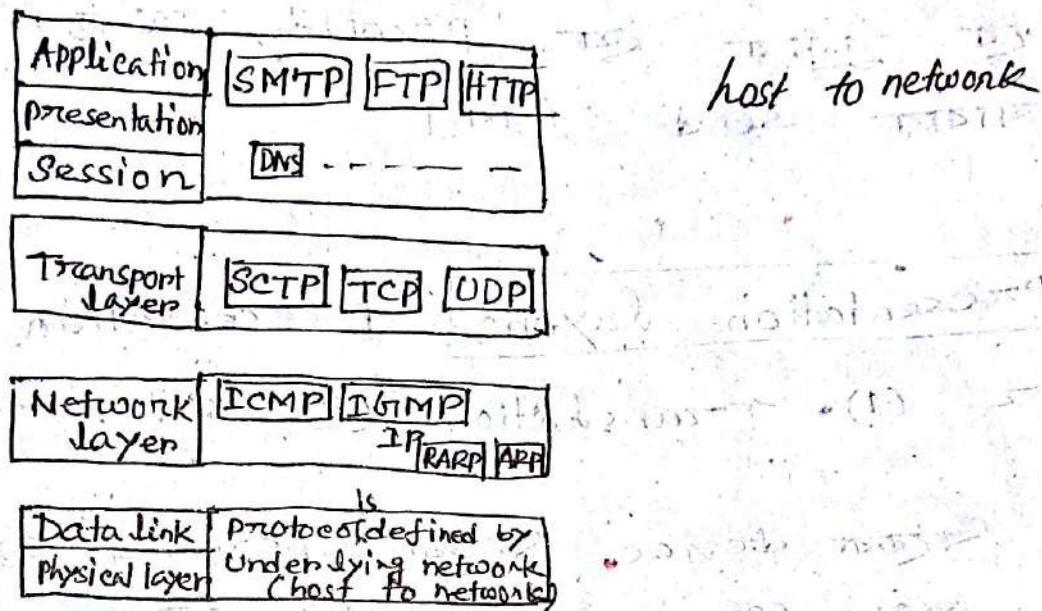
(i) File transfer, Access, Management

(ii) Access Mail service

(iii) Directory service

13-05-17  
2nd (d) day

TCP/IP Model :- ~~Top Top~~ Layer 6 fixed protocol  
21/22



TCP → Transmission control protocol

IP → Internet protocol

Protocol: protocol defined by underlying network (host to Network)

At physical and data link layer, TCP/IP does not define any specific protocol.

Network layer:

Internet protocol (IP): - (Logical Address protocol)

Address Resolution protocol (ARP):  
- is used to associate a logical address with a physical address.

Reverse Address Resolution protocol (RARP)

→ Allow a host to discover its internet address when it knows only its physical address.

ICMP → Internet control Message protocol

IGMP → Internet group

UDP → User datagram protocol.

Process to process transmission control

করতে।

TCP → Transmission control protocol

গোড়া করা transmission হবে, যাবে  
যাবে। কোন problem না হয়, provides  
full transport layer service to the  
Application layer.

SCTP → Stream control transmission protocol

চারটি active হবে, যাবে নতুন function  
or feature add হবে অন্য বিষয়ে

FTP → File transfer protocol

DNS → Domain Name server

14-05-17

2nd (E) day

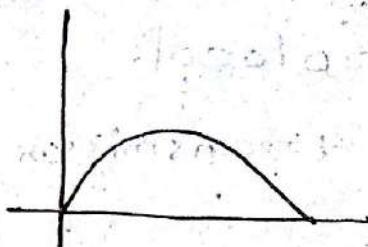
## Ch 3: Data & Signal

Analog Data :- Continuous time, continuous value नियन्त्रित Analog data

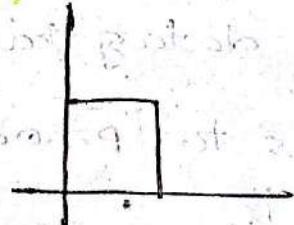
Digital :- Discrete, discrete value नियन्त्रित Digital data

Analog Signal :- Analog Data शुल्का नियन्त्रित से Signal

Digital :- Digital



Analog data  
with signal



Digital data  
with signal

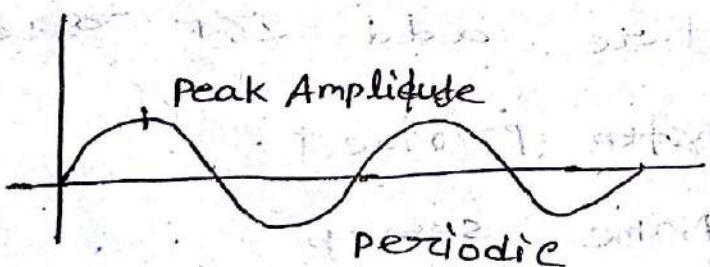
Periodic signal : नियन्त्रित time पर पर एक period के लिए आवृत्ति

Non periodic signal : नियन्त्रित time पर period नहीं आवृत्ति

एकमें complete cycle देखी करते हैं तो उसे time period

एक अवधि के अन्तराल में complete signal इसे ज्ञान-

जाता है frequency यही।



Signal characteristics :-

1) Amplitude

Amplitude

2) Frequency

A Frequency

3) Phase (एकोन अवधि signal पर अवधि)

Signal वर्तु highest value तक वजा असे Peak Amplitude आहे.

$$100\text{ms} = 100 \times 10^{-3} \text{s}$$

$$100\text{ms} = 100 \times 10^{-3} \times 10^6 \text{μs}$$

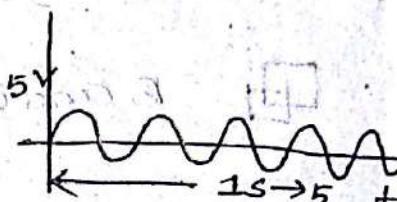
$$f = \frac{1}{T} = \frac{1}{10^5}$$

$$\text{वाव्हाणी लंबाई } (\lambda) \rightarrow \lambda = \frac{C \text{ ms}}{f \text{ s}} = m$$

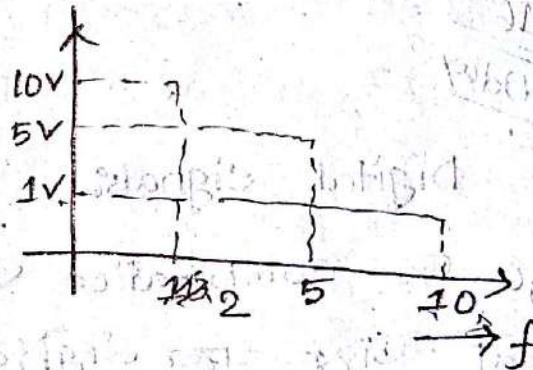
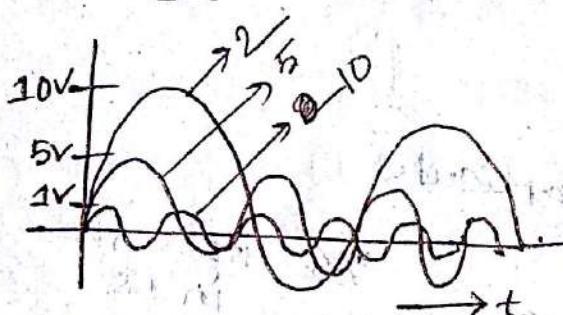
~~N, V, I~~  
D-Time Domain & Frequency Domain तो? Same from  
तो तरी,

i) Time वर्तु respect व तो Signal  
पद्धति complex रूपायास।

ii) Frequency  $\rightarrow$   $\rightarrow$   $\rightarrow$   $\rightarrow$   $\rightarrow$   
पद्धति अचूक रूपाचि देखाय ते Time domain तो



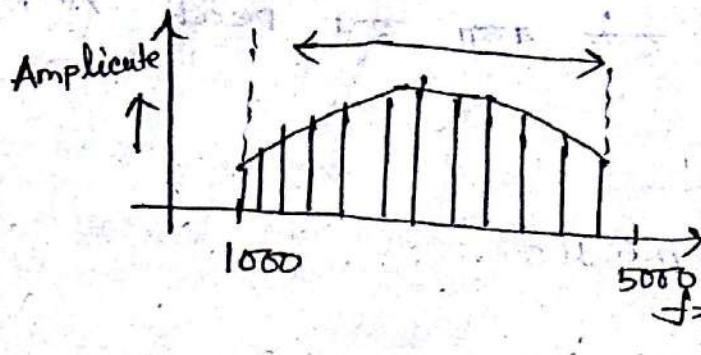
उभनास।



उक्तसे same frequency तरी तुटे देश-

अनेक शूल Signal Add करू ये Signal पाऊया यास आणि Composite Signal वाले।

Bandwidth:- The bandwidth of a composite signal is referred to as the difference between the highest frequencies and lowest frequencies contained in that signal.



$$\text{Bandwidth} = (5000 - 1000) \text{ Hz} = 4000 \text{ Hz}$$

Frequency 100, 300, 500, 700, 900 Hz

$$\text{Bandwidth} = (900 - 100) \text{ Hz} = 800 \text{ Hz}$$

Bandwidth 20 Hz, highest = 60 Hz, lowest = ?

$$20 = 60 - l\omega$$

$$\Rightarrow l\omega = 40 \text{ Hz}$$

20-05-16  
3rd (d) day

### Digital Signals:-

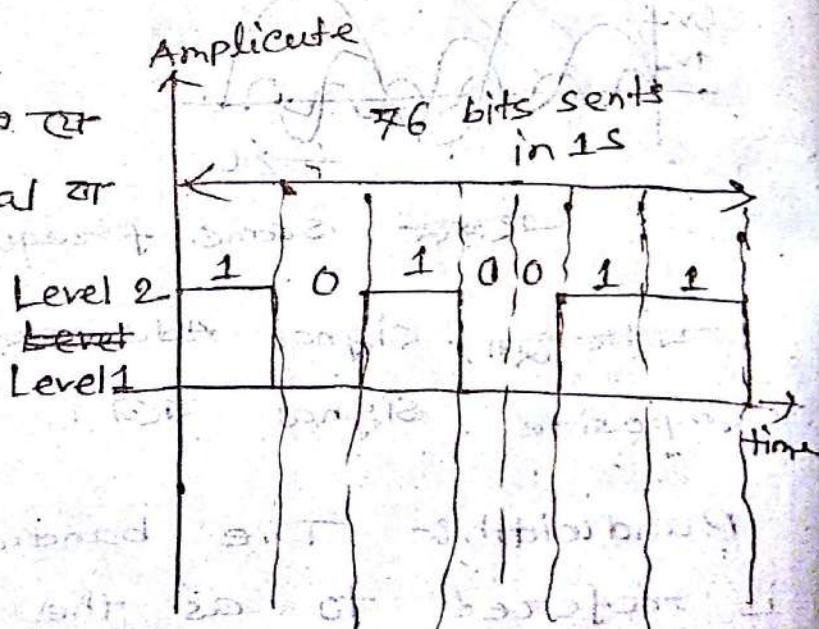
10 Gb combination वर्ते

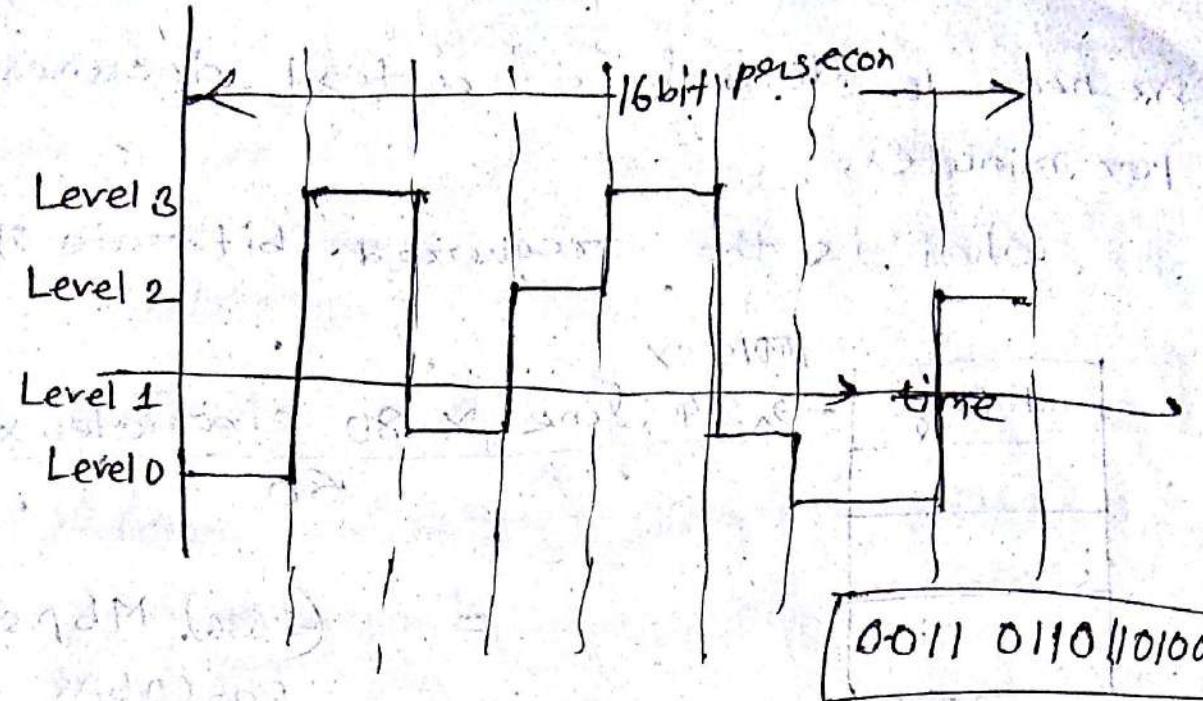
signal जाके यहाँ digital वर

Discrete signals.

info बताने  $\rightarrow 1$

" ना "  $\rightarrow 0$





$$2^r = 4 \rightarrow 00 \rightarrow 0$$

$$01 \rightarrow 1$$

$$10 \rightarrow 2$$

$$11 \rightarrow 3$$

at a time 3 bit send करते जाएं  $\log_2 L$   
यद्यपि,  $\log_2 L = 8$   $= \log_2^3$   
 $= 3$

8 के level हल्के at a time 3 bit send करते  
एकमें level के integer होते, ये 6 bit पास जाते  
अलग integer होते।

$$\log_2 9. हल्के = 3.2, \text{ यद्यपि} - \text{ Consider करें } 4.6$$

$$= 0 =$$

$$000 \rightarrow 0$$

$$001 \rightarrow 1$$

$$010 \rightarrow 2$$

यद्यपि 4 bit का होता है

$$0000 \rightarrow 0$$

$$00$$

We need to download a text document 100 pages per minutes,

What is the required bit rate of the channel

1 Page	100 pages
	$\rightarrow 24 \text{ line} \times 80 \text{ character} \times 8 \text{ bit}$
	$\frac{60}{25600 \text{ bps}}$

$$= (1.636) \text{ Mbps}$$
$$25600 \text{ bps}$$

Bit rate  $\rightarrow$  No of bits per second

Bit length  $\rightarrow$  propagation speed  $\times$  bit duration

$$\rightarrow \text{ms}^{-1} \times s$$

$$\rightarrow [ \text{m} - \text{Sec} ]$$

Transmission impairment :- (वायन त्रिकृति द्वारा यार्ड बार्कटर signal को perfectly Received न हो सकता)

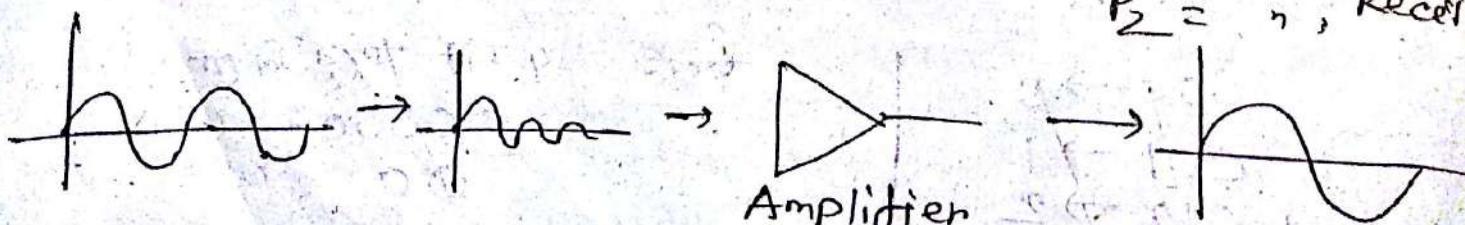
- i) Attenuation
- ii) Distortion
- iii) Noise

(loss of energy)

energy के change इसलिए  
इस signal देख

$P_1$  = Power of sender

$P_2$  = , Receiver



$$\text{decibel, } dB = 10 \log_{10} \frac{P_1}{P_2}$$

Example

$$\begin{aligned} dB &= 10 \log_{10} \frac{P_1}{0.5 P_1} \\ &= 10 \log_{10} \left( \frac{1}{0.5} \right) \\ &= 3.01 \text{ dB} \end{aligned}$$

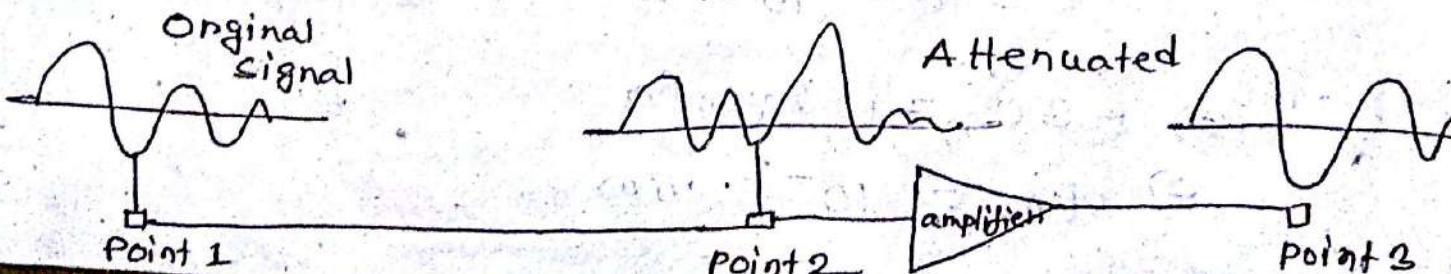
$$\begin{aligned} P_1 &= P_1 \\ P_2 &= \sum P_i \\ &= 0.5 P_1 \end{aligned}$$

Positive वार्ताएँ, चाहे वार्ता loss positive, loss वार्ता  
loss . negative वार्ताएँ, loss वार्ता

21-05-2017  
3rd (E) day

Transmission Impairment:- Signals travel through transmission media, which are not perfect. The perfection causes signal impairment. This means that the signal at the begining of the medium is not the same as the signal at the end of the medium. What is sent is not what is received. Three causes of impairment are attenuation, distortion and noise.

Attenuation: Attenuation means a loss of energy when a signal, simple or composite, travels through a medium, it losses some of its energy in overcoming the resistance of the medium.



$$dB = 10 \log_{10} \frac{P_2}{P_1}$$

Variables  $P_1$  and  $P_2$  are the powers of signal at Point 1 & Point 2 respectively.

### Example

$$\begin{aligned} dB &= 10 \log_{10} \frac{P_2}{P_1} \\ &= 10 \log_{10} \frac{0.5 P_1}{P_1} \\ &= -3 \text{ dB} \end{aligned}$$

$$\begin{aligned} P_1 &= P_1 \\ P_2 &= \frac{1}{2} P_1 \end{aligned}$$

वर्ग लोस  $\rightarrow$

### Another example

$$\begin{aligned} \text{Loss of energy, } dB &= 10 \log_{10} \frac{10}{1} \\ &= 10 \times 1 \\ &= 10 \text{ dB.} \end{aligned}$$

वर्ग गेन  $\rightarrow$  लोस  $\rightarrow$  (+) अरे

$$P_2 = 10 P_1$$

$$P_1 = P_1$$

### Another example

$$\text{Power in milliwatt, } P_m = \frac{P_2}{P_1}$$

$$\begin{aligned} dB_m &= 10 \log_{10} \frac{P_2}{P_1} \\ &= 10 \log_{10} P_m \end{aligned}$$

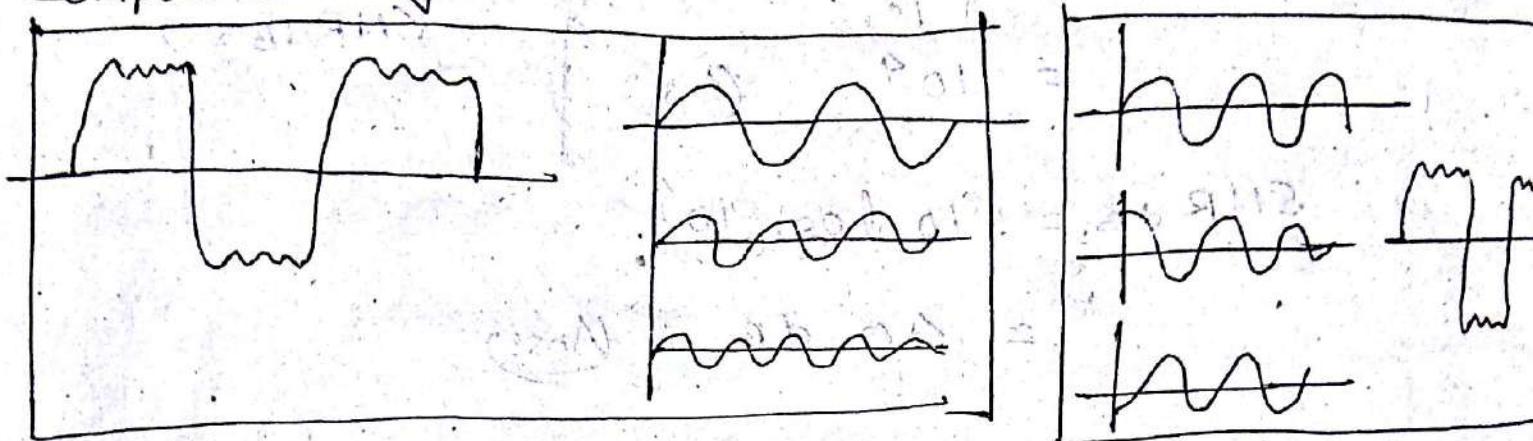
$$\Rightarrow -30 = 10 \log_{10} P_m$$

$$\Rightarrow P_m = 10^{-3} \text{ mW}$$

$$dB_m = -30$$

$$P_m = \frac{P_2}{P_1} = ?$$

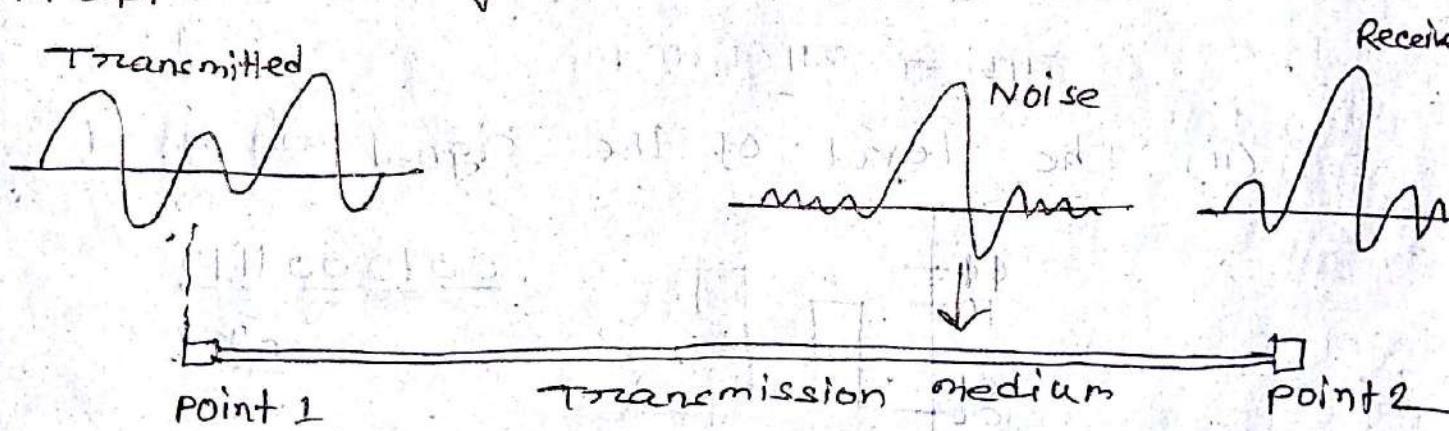
Distortion: Distortion means that the signal changes its form or shape. Distortion can occur in a composite signal made of different frequency.



At the sender

At the receiver

Noise: Noise is another cause of impairment. Several types of noise such as thermal noise, induced noise, cross talk noise and impulse noise, may corrupt the signal.



The signal to noise relation is defined as,

$$\text{SNR} = \frac{\text{Avg. Signal power}}{\text{Average Noise power}} \\ = 10 \log_{10} \text{SNR}$$

SNR  $\Rightarrow$  high SNR  $\Rightarrow$  less noise  $\Rightarrow$  better,  $\text{SNR} > 1$   
 " "  $\Rightarrow$  low SNR  $\Rightarrow$  more noise  $\Rightarrow$  worse,  $\text{SNR} < 1$

### Example

$$SNR = \frac{P_1}{P_2}$$

$$= \frac{10^4}{10} = 10^4 \text{ (Ans:)}$$

Signal power,  $P_1 = 10 \text{ mW}$   
 $= 10^4 \mu\text{W}$   
 Noise,  $P_2 = 1 \mu\text{W}$

$$SNR = ?$$

$$SNR_{dB} = ?$$

$$SNR_{dB} = 10 \log_{10} 10^4$$

$$= 40 \text{ dB} \text{ (Ans:)}$$

03-07-17  
4th (B) day

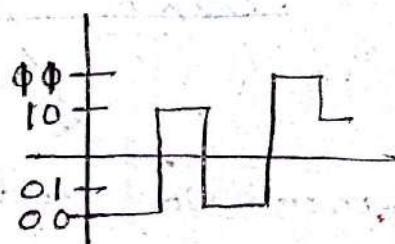
Data rate Limit:- 3 factors needs:-

(i) Bandwidth of the channel:

Bandwidth का आकलन data rate

वाल्याना यावे ता।

(ii) The level of the signal we used



00 10 00 11 10

$2^5$

channel

signal : (the level of noise)

noise एवं इस quality धारणा

तो यह तरीका जानना चाहिए

Bit rate:-

No. of Bit per second.

Noiseless channel :- Nyquist Bit Rate

Maximum Bitrate,

$$\text{Bitrate} = 2 \times \text{Bandwidth} \times \log_2 L$$

[Level to power of 2 तक तक]

Noisy channel :- Shannon capacity.

Theoretical highest data Rate formula,

$$\text{Capacity} = \text{Bandwidth} \times \log_2 (1 + \text{SNR})$$

Signal noise Ratio

$$\text{SNR} = \frac{\text{Avg signal power}}{\text{noise}}$$



प्र० अंका आर०,

$$\underline{\text{BW}} = 3000 \text{ Hz}$$

$$L = 8$$

$$\text{Bitrate} = ?$$

$$\text{Bitrate} = 2 \times \text{BW} \times \log_2 L$$

$$= 2 \times 3000 \times \log_2 8$$

$$= 18000 \text{ bps}$$

Ans.



प्र० अंका आर०,

$$\text{Bit rate} = 265 \text{ Kbps} = 265000 \text{ bps}$$

$$\text{BW} = 20 \text{ kHz} = 20000 \text{ Hz}$$

$$L = ?$$

$$265000 = 2 \times 20000 \times \log_2 L$$

$$\Rightarrow L = 2^{6.625}$$

$$\Rightarrow \log_2 L = 6.625$$

$$\Rightarrow L = 2^{6.625} = 98.7 \text{ level}$$

एवं level के 128 गे निम्न याव,

$$\text{Now, Bit rate} = 2 \times 20,000 \times \log_2^{128}$$
$$= 2,80,000 \text{ bps}$$
$$= 280 \text{ Kbps}$$

यद्यपि level के 64 गे निम्न याव,

$$\text{Bit rate} = 2 \times 20,000 \times \log_2^6$$
$$= 2,4900 \text{ bps}$$
$$= 240 \text{ Kbps}$$

level के एवं floating value  $\frac{2}{2^9} \text{ ता } \frac{1}{2^9}$   
 $= 0 =$

यदि  $SNR = 0$  तथा

$$\text{capacity} = BW \times \log_2 (1+SNR)$$
$$= BW \times \log_2 (1+0)$$
$$= BW \times \log_2 1$$
$$= 0$$

अतः कोर नोइस याकृत ना। data transmit  
हो ला। Noise किंवा याकृत हो।

জোম্পা আরে,

We know,

$$\text{Capacity} = B \times \text{BW} \times \log_2^{(1+\text{SNR})}$$
$$= 2 \times 2 \times 10^6 \times \log_2^{(1+10^{3.6})}$$
$$= 2 \times 10^6 \times \log_2^{(1+3981)}$$
$$= 24 \text{ Mbps}$$

(Ans.)

$$\text{SNR}_{dB} = 10 \log_{10}^{\text{SNR}}$$
$$= 36$$

$$36 = 10 \log_{10}^{\text{SNR}}$$

$$\Rightarrow \log_{10}^{\text{SNR}} = 3.6$$

$$\Rightarrow \text{SNR} = 10^{3.6}$$

$$= 3981$$

$$\text{BW} = 2 \text{ MHz}$$

$$\text{Capacity} = ?$$

= 0 =

Throughput:- Throughput is a measure of how fast we actually send data through a network.

যদি কোনো নেটওর্ক এর Bit rate = 10 Mbps  
কান সেন্স 12000 frame per minutes each frame  
কান কার্য 10,000 bit. Throughput = ?

$$\text{Throughput} = \frac{12000 \times 10000}{60 \times 10^6 \text{ sec}}$$
$$= 2 \text{ Mbps}$$

Propagation time: 1st bit source to dest  
⇒ সৌধার দ্বা অবস্থ হয়।

$$\text{Propagation time} = \frac{\text{Distance}}{\text{propagation speed}}$$

$$\text{Propagation time} = \frac{12000 \times 10^3 \text{ m}}{2.4 \times 10^8 \text{ m/s}} = 50 \text{ ms}$$

distance = ~~12 km~~ 12000 km

propagation speed = ~~2.4 × 10^8 m/s~~

$$2.4 \times 10^8 \text{ m/s}$$

propagation time = ?

Transmission time: ~~whole message~~  
Sender ~~center~~

time: ~~whole message~~ ~~center~~  
destn ~~center~~ ~~message~~ ~~center~~

$$\text{Transmission time} = \frac{\text{Message size}}{\text{Bandwidth}}$$

exam → ~~total~~ ~~total~~ propagation time (3)

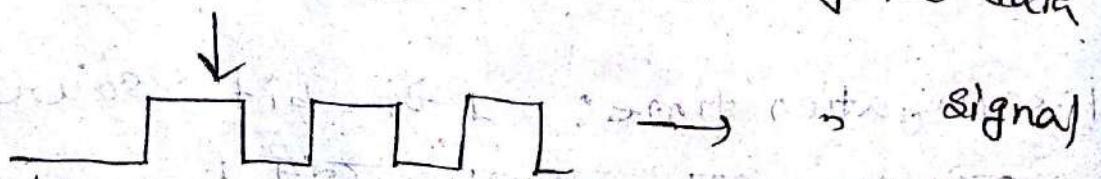
Transmission time ~~total~~ ~~total~~ ~~total~~ ~~total~~

$$= 0 =$$

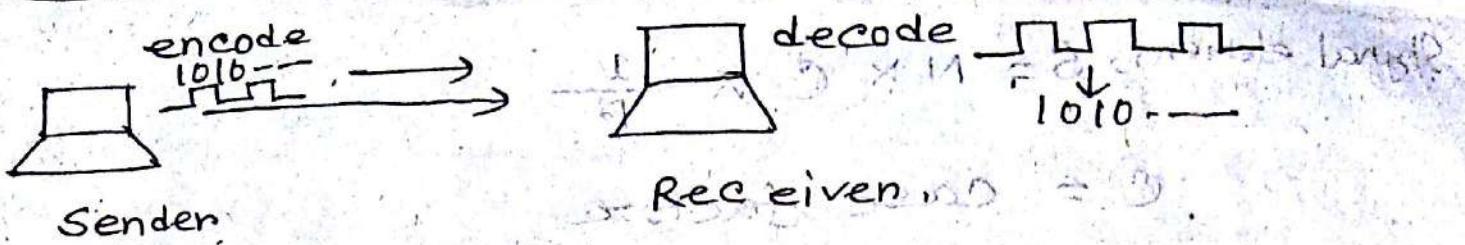
### Digital Transmission

Digital to Digital conversion:-

101011001010 ----- → Digital data

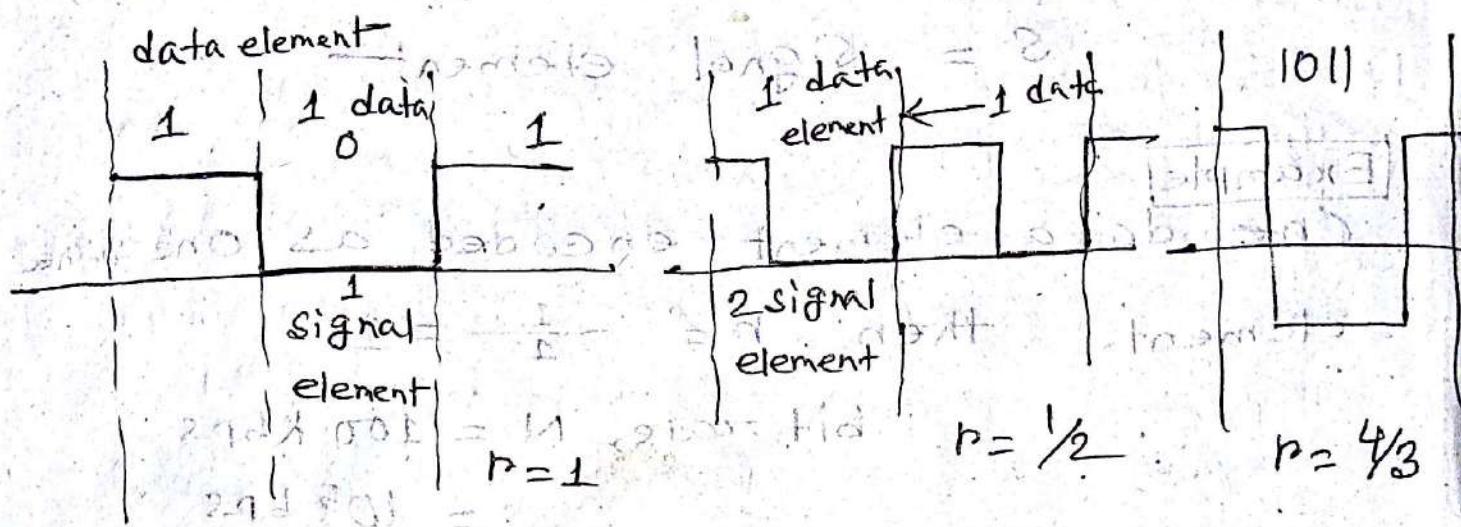


Digital data ~~center~~ Digital signal ~~center~~  
D. to D. conversion.



Line coding technique: one kind of technique where we can do D. to D. conversion.

Data element vs signal element :-



1 Signal element: 1 से data पाठें रख।

2, 5 2, 0 2, 5 2, 3

Data element = bps

(What we need to send)

Signal element = bound rate / pulse rate/modulation rate

(1 sec. second ~ कठशुल्य) data pulse sent इ)

(canvien, यारे घार्व) वा send करा य)

$$\text{Signal elements } S = N \times C \times \frac{1}{R}$$

$C$  = Case factor

$N$  = Number of bits per second

$$R = 4 \text{ bit} \quad \text{then } R = \frac{4}{3} = \frac{\text{bit}}{\text{number of signal}}$$

→ 1 second → 1 bit

$S$  = Signal element

### Example

One data element encoded as one signed element, then  $R = \frac{1}{1} = 1$

$$\text{bit rate, } N = 100 \text{ Kbps} \\ = 10^5 \text{ bps}$$

case factor between 0 & 1

$$\text{Now, } C = 0.5$$

Avg. Signal rate  $S = ?$

$$S = N \times C \times \frac{1}{R}$$

$$C = 0.5$$

$$S = 10^5 \times 0.5 \times 1$$

worst case

$$= 50000 \text{ baud}$$

$$C = 1$$

Best case

Avg case

$$C = \frac{10000}{2}$$

$$= 0.5$$

Band width  
Minimum Bit rate:

$$B_{\min} = C \times N \times \frac{1}{n}$$

$$= 0 =$$

D. D to D. conversion problem:-

- i) Baseline wandering
  - ii) DC component
  - iii) Self synchronization
- } Factor of Digital to Digital conversion

(i) (i) & (ii) द्वारा अनेक घिन आहे।

00000001000000

~~1111101111~~

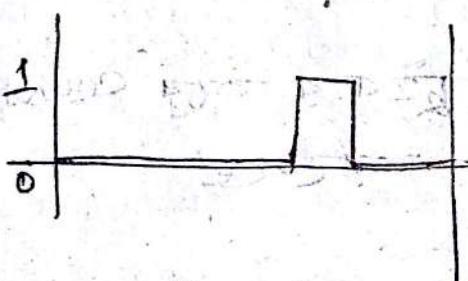
String एवं bit sequence से यदि अलगशुल्ला

एवं पर 1 वा अलग 1 एवं पर 0 अको 0। तरी

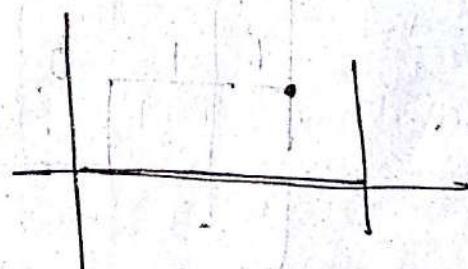
avg value count, करते, ज्ञान 1 तरी 0

संकेत त्रैमा Receiver, ज्ञान signal आके

अव 0 इम 1



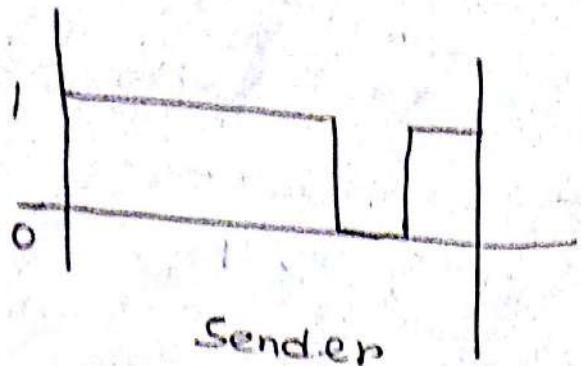
Sender



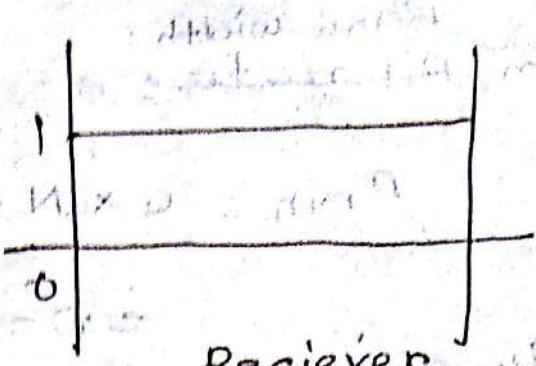
Receiver

एटा Baseline wandering problem.

(ii) conse. 11111011111



Sender



Receiver

अपश्लासा के 1 receive count तक

एटोडे D के component problem.

(iii) sender के अवस्था वर्णन send करते

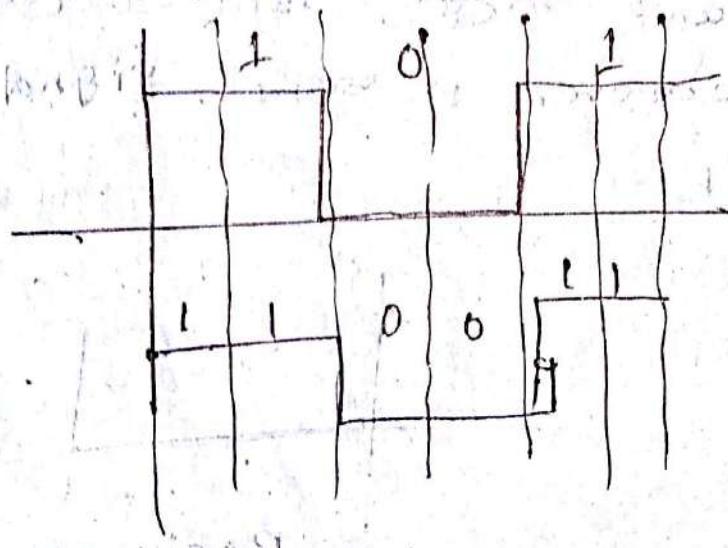
Receiver के अवस्था वर्णन- Receive

काले ना पारवे यही synchronization है

मार्क उत्तर भारक (iii) तले। एक्स

delay यह, पारवे इसे 1 में

पारवे पारवे इसे 0 और Receive करते।



इसका काले count  
वर्णन

08-07-17  
5th (A) day  
4th (E) day

Line coding schemes:- (disadvantages, advantages, uses)

1. Unipolar  $\rightarrow$  NRZ (Non return to zero)
2. Polar  $\rightarrow$  NRZ, RZ, and biphase (Manchester)
3. Bipolar  $\rightarrow$  AMI and pseudoternary Differential Manchester
4. Multilevel  $\rightarrow$  2B/1T, 8B/6T and 4D-PAM
5. Multitransition  $\rightarrow$  MLT-3

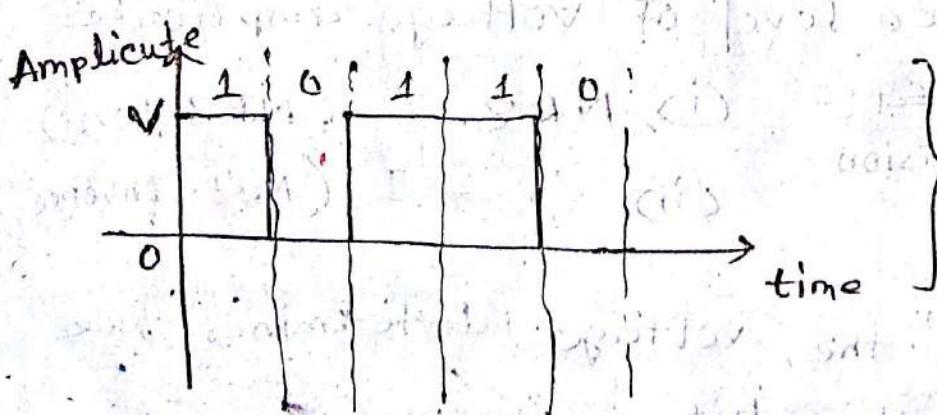
### ① Unipolar scheme:-

All the signal levels are on one side of the time axis either above or below.

Non return to zero (NRZ)

— Positive voltage define bit 1 & zero voltage define bit 0

— It is called NRZ because the signal does not return to zero at the middle of the bit.



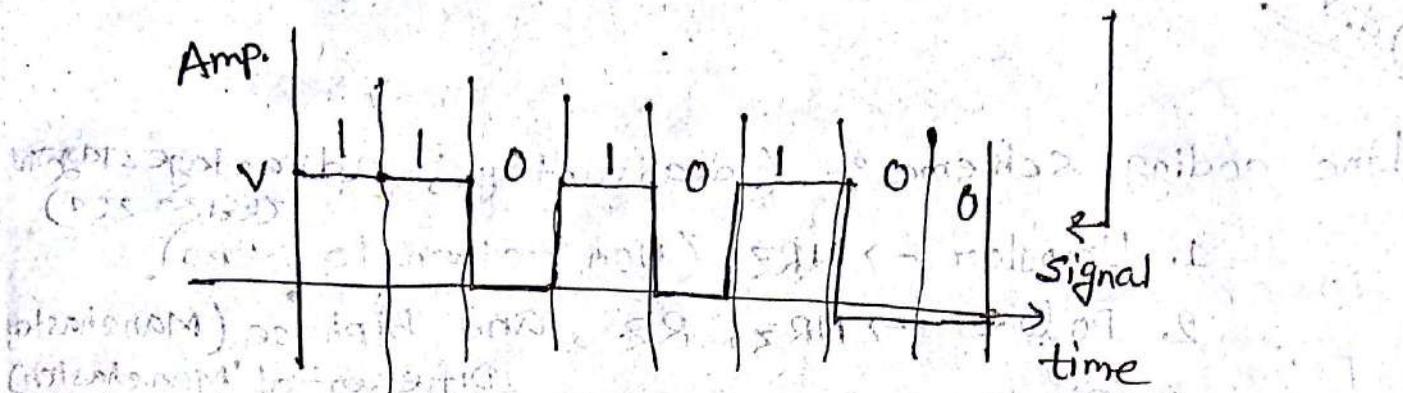
Positive फॉर्म वाले  
वेंसुन्तर consider करा  
से 1 से कानून रखिए  
आकर्षण - Signal दे  
either (+) or (-)

प्रयोग से signal एवं bit एवं ब्यास  
transition असे ना। एकल वला  
इस NRZ

फॉर्म आकर्षण। छाई  
unipolar scheme.

2

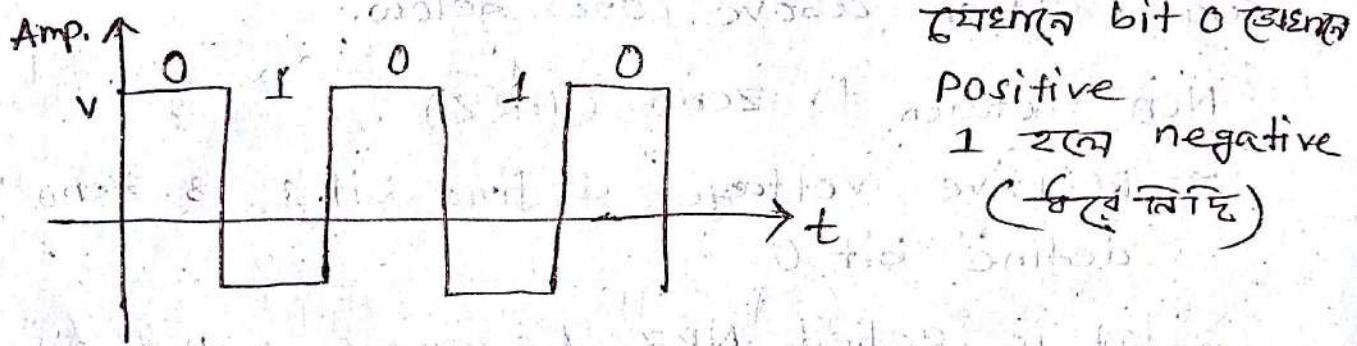
11 01 0100  $\rightarrow$  unipolar scheme



2

Polar Scheme:-

- the voltages are on the both sides of the time axis. For example, the voltage level for 0 can be positive & the voltage level for 1 can be negative.



Polar NRZ:-

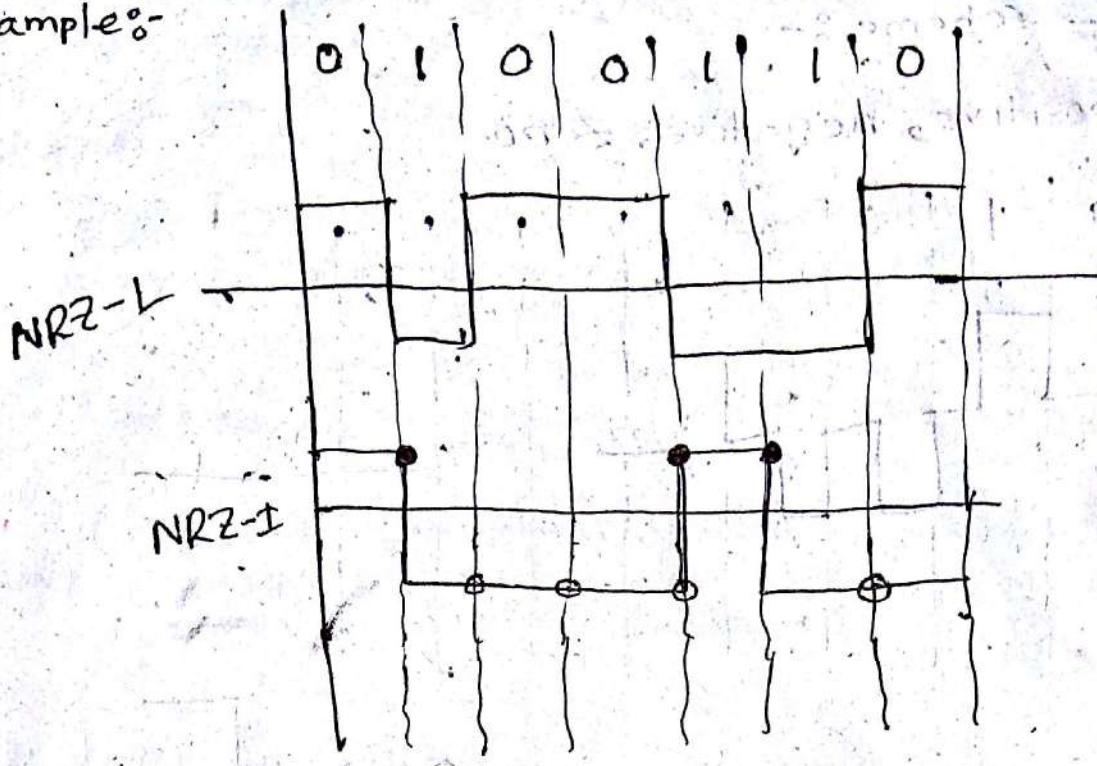
We use two level of voltage amplitude.

दो रेल परिवर्तन :- (i) NRZ-L (NRZ-Level)  
(ii) " - I (NRZ-Inverse)

(i) the level of the voltage determines the value of the bit.

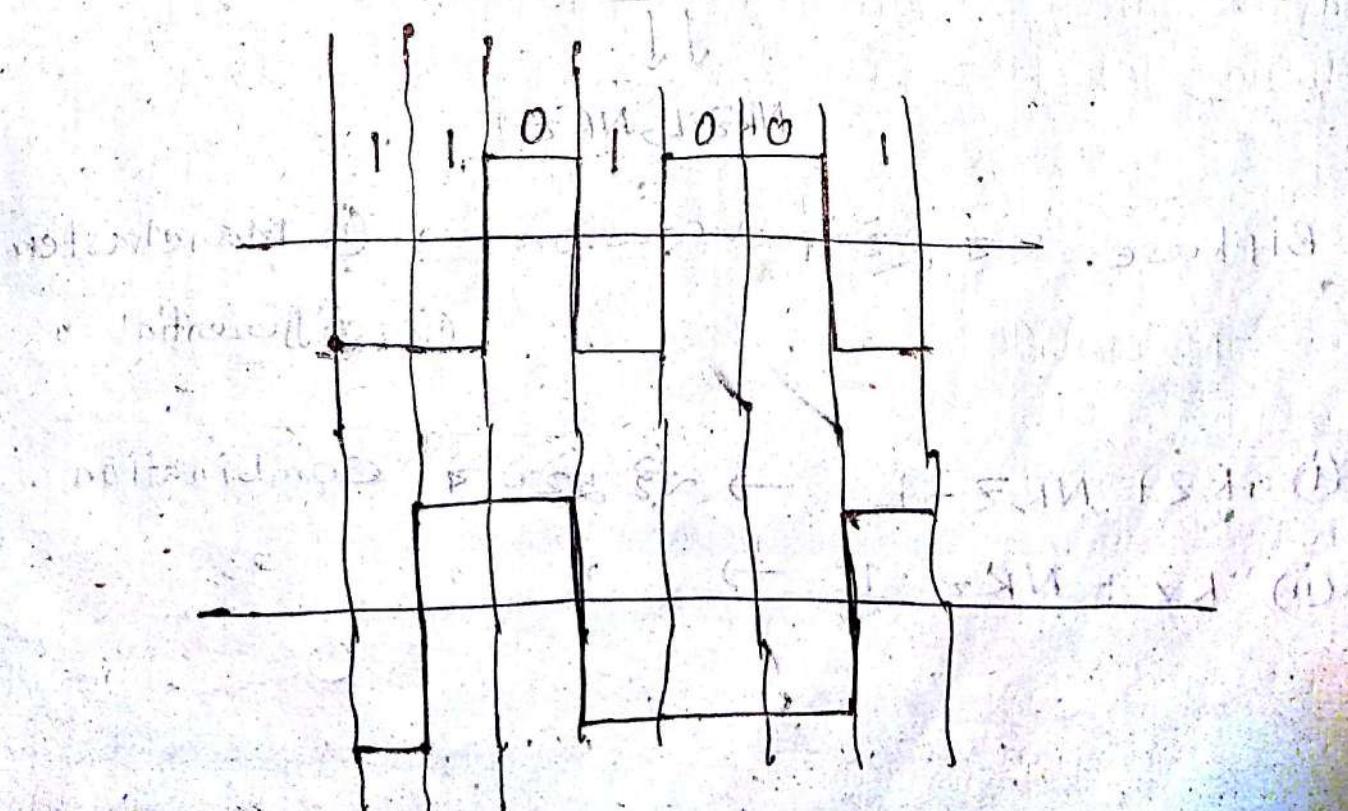
(ii) If there is no change the bit is 0  
" 3 " a 3 3 3 1

Example:-



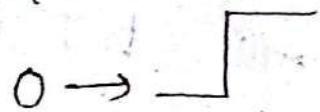
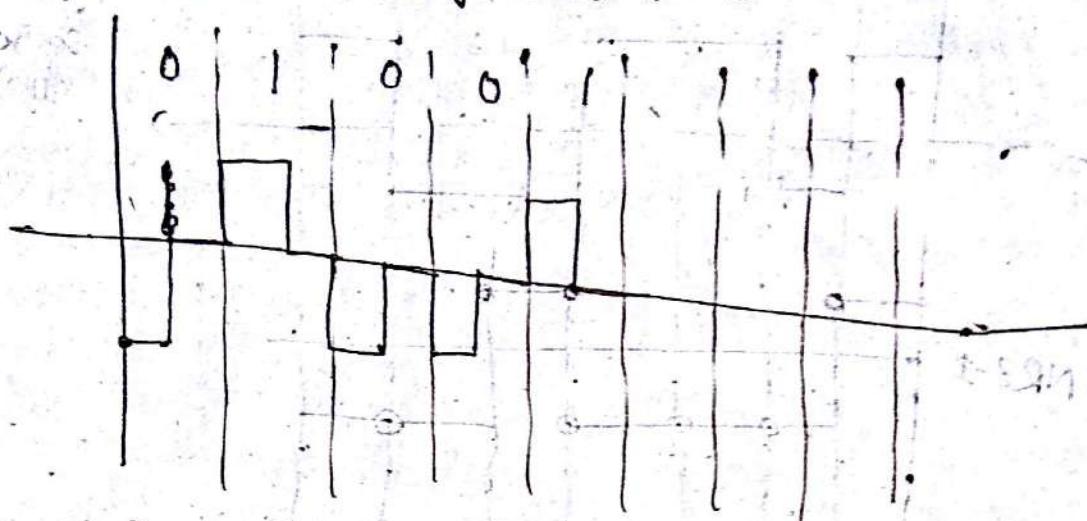
NRZ-L → No inversion  
NRZ-I → next bit inversion  
NRZ-I → Next positive pulse start

Another example: 1 1 0 1 0 0 1



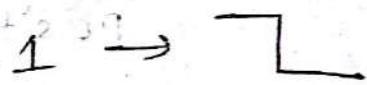
Polar RZ - Scheme :-

→ Positive, negative, zero.



0 पाइले negative तरफ 0 level.

1 " positive  $\rightarrow$  0 "



Polar biphasic ( $\frac{NRZ+RZ}{J\downarrow}$ )

NRZ-L, NRZ-I

Biphase एवं इसकी version → (i) Manchester  
(ii) differential

(i) RZ + NRZ-L  $\rightarrow$  इसकी Combination

(ii) RZ + NRZ-I  $\rightarrow$  " "



RZ एवं level शृङ्खला रख  $\rightarrow$  positive

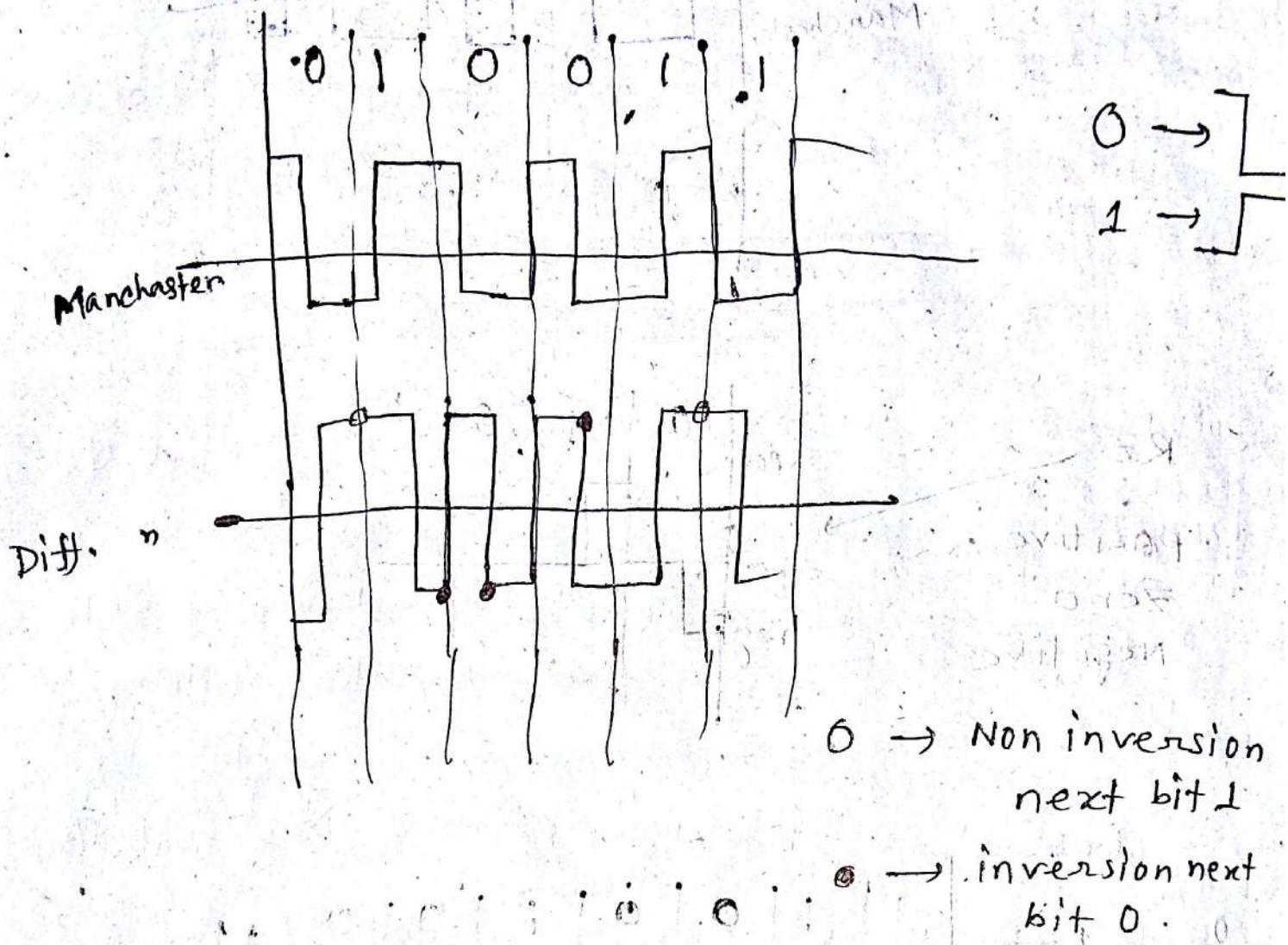
zero

Negative

pos.

0

-  
neg.

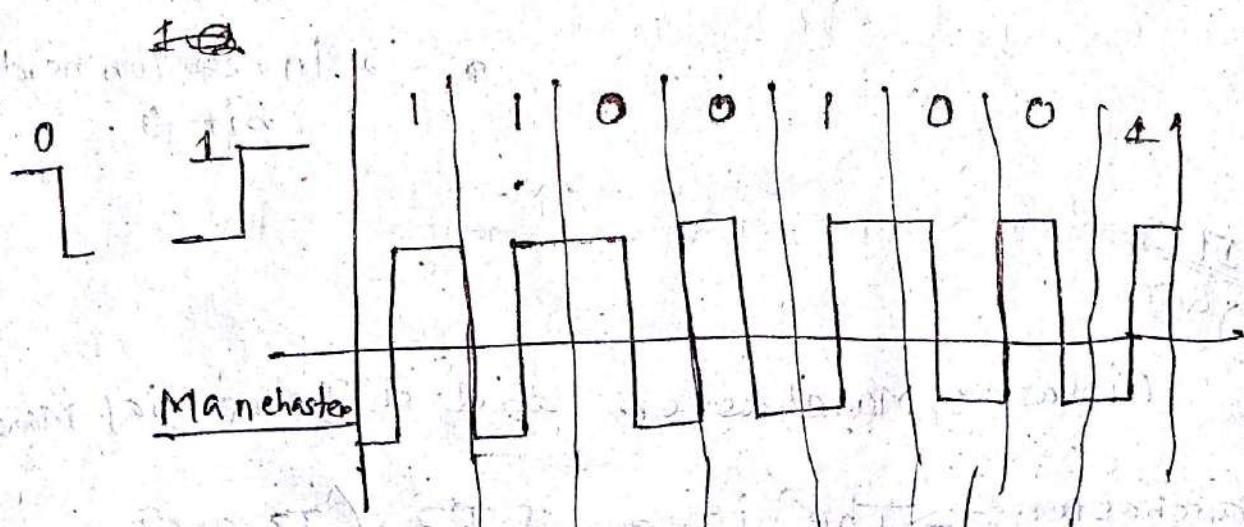
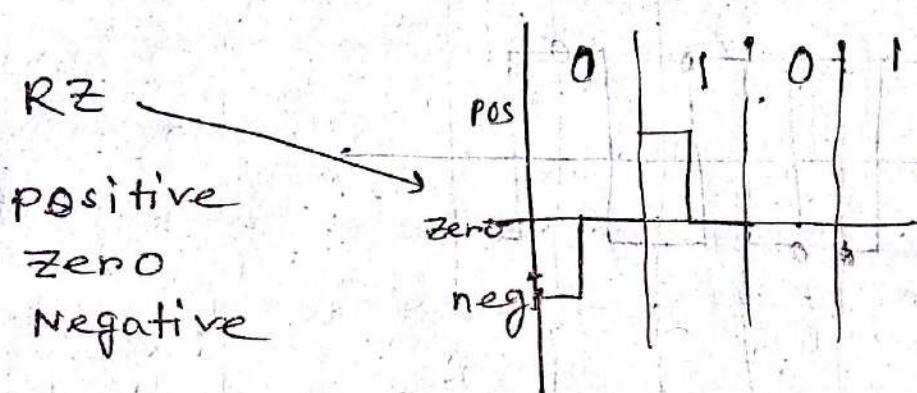
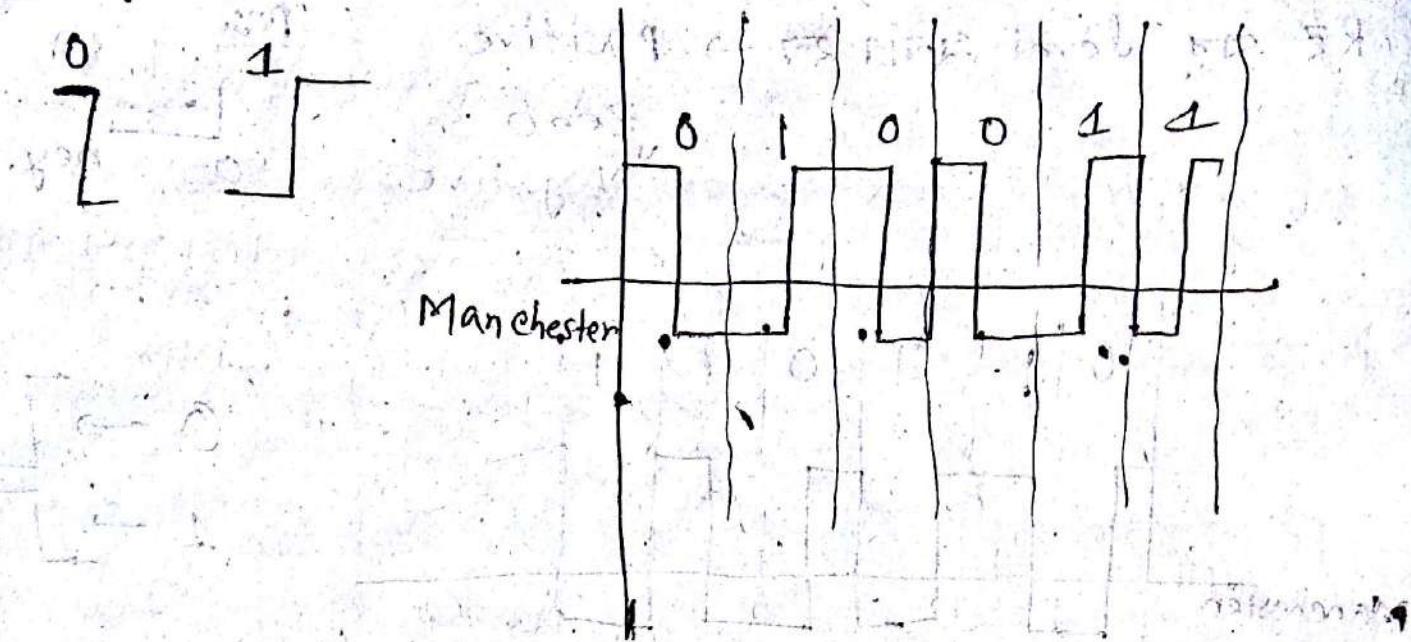


10-07-17  
5th (B) day

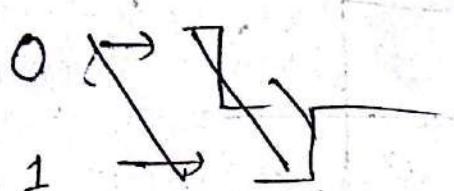
Biphase: Manchester and differential Manchester

Manchester :-  $\Rightarrow$  The idea of RZ (Transition at the middle of the bit) and the idea of NRZ-L are combined into the Manchester scheme.

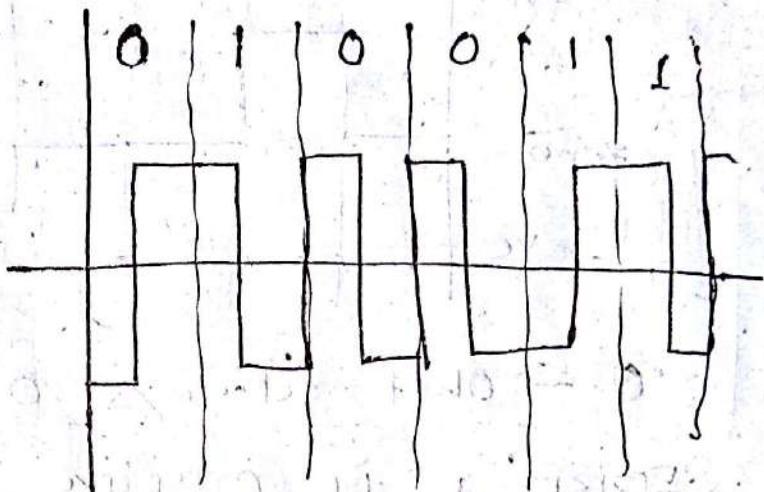
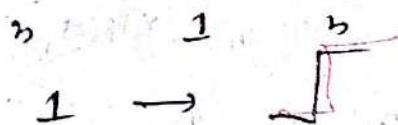
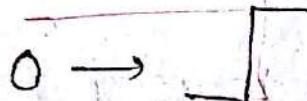
$\Rightarrow$  The voltage remains at one level during the first half and moves to other level in the second half.



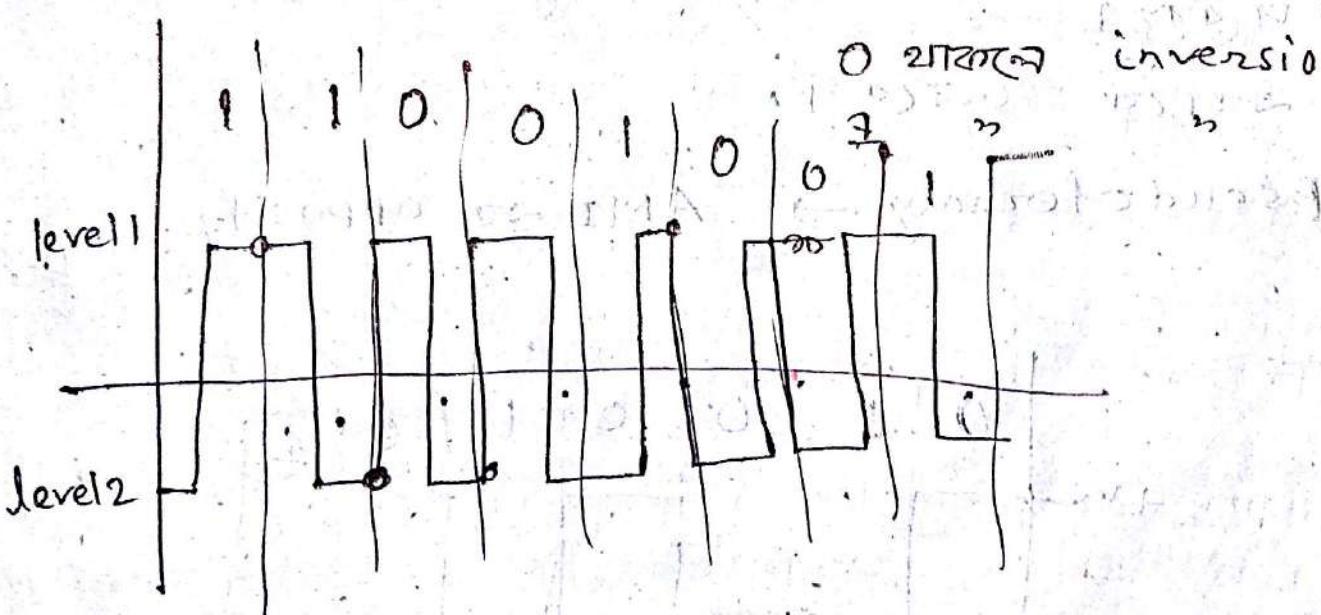
Differential Manchester:- Diff. Man. combines the ideas of RZ and NRZ-I. There is always a transition at the middle of the bit, but the bit values are determined at the beginning of the bit. If the next bit is zero, there is a transition.



ପ୍ରଥମ । ୦ ଶାକଳେ



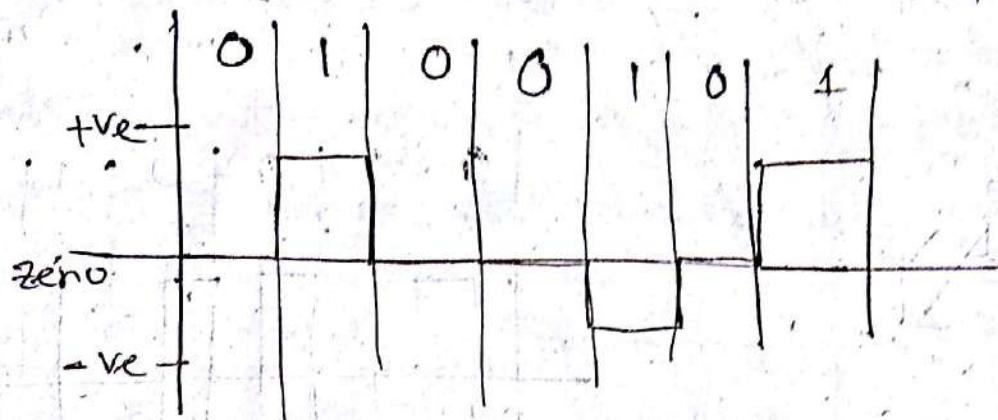
- 0 → no inversion, next bit → 1
  - → change, next bit → 0



## Bipolar Schemes:-

- There are three voltage level - positive, negative, zero

AMI (Alternate Mark Inversion)



0 → 0 bit always, 0

परिवर्तन 1 के positive level के बारे में प्रत्येक एक

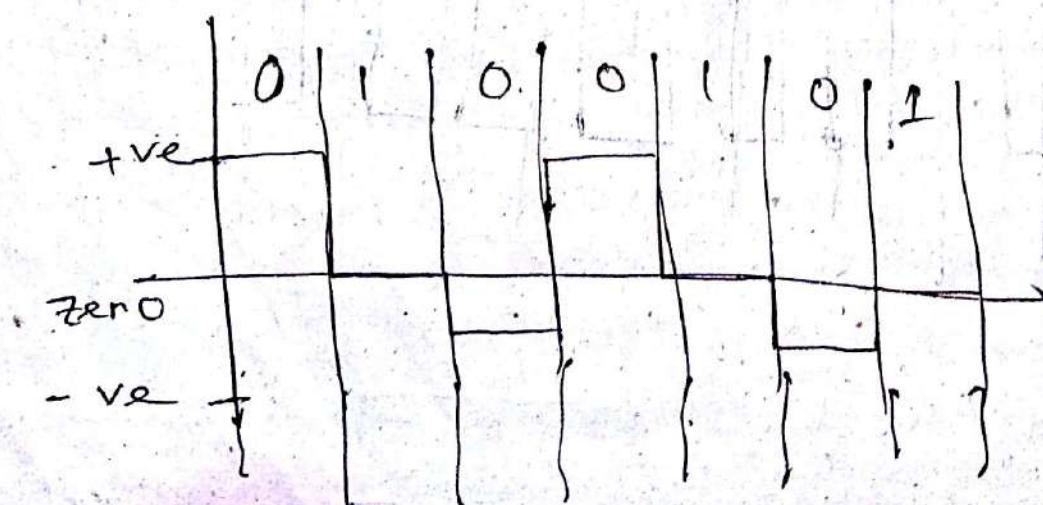
परिवर्तन 1 के negative के बारे में प्रत्येक

परिवर्तन 1 के अधिकारी positive के बारे में

परिवर्तन 1 के अधिकारी negative के बारे में

प्रत्येक चलने।

Pseudoternary → AMI के opposite



1  $\rightarrow$  always 0 level

0 पर change होते अब AMI नहीं 1 का पाठ्यक्रम

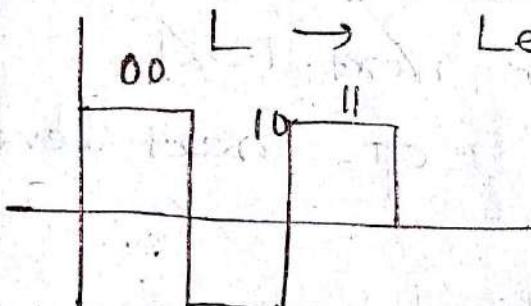
12-07-17  
5th (d) day

Multilevel Scheme:- 0, 1 लिए एकाविक डाटों पाठ्यक्रम  
at a time 2 बिट करते send करते 4 से combination  
अस्तव ।

mBnL scheme:- A pattern of m data elements  
encoded as a pattern of n signal elements in  
which  $2^m < L^n$

$m \rightarrow$  data element

$n \rightarrow$  signal



$\rightarrow$  एकल -  $m = 2$   
 $n = 1$

2 बिट करते पाठ्यक्रम

signal

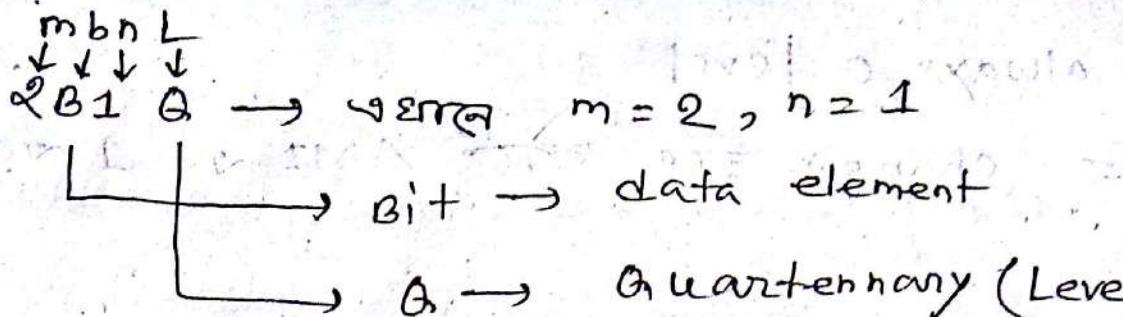
2 बिट करते पाठ्यक्रम  $\rightarrow$

00  
01  
10  
11

Possible level 4 से  
 $L = 4$

$$2^m = 2^1 = 2$$

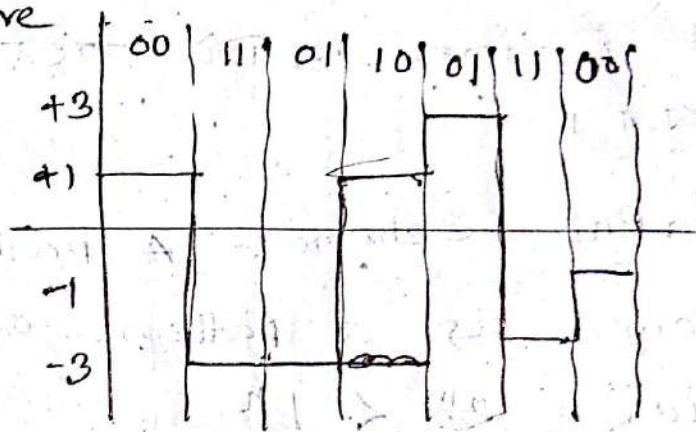
$$L^n = 4^1 = 4$$



ग्रेड 2 का 2B1Q Scheme.

पूर्वी लेवल कोहार्ड ने perform करते।  
 Previous level positive previous level negative

Next bits	Next level	Next Level
00	+1	-1
01	+3	-3
10	-1	+1
11	-3	+3



Transition table.

प्रारंभिक 00 जब 11, Next level 11 तक 2/3  
 next bit 01 then next level

पूर्वी कोहार्ड रेटिंग 2/3, next level 0.

Previous level.

प्रारंभिक positive रेटिंग start करता।

Multiline transmission → (M QLT.-3) लेव

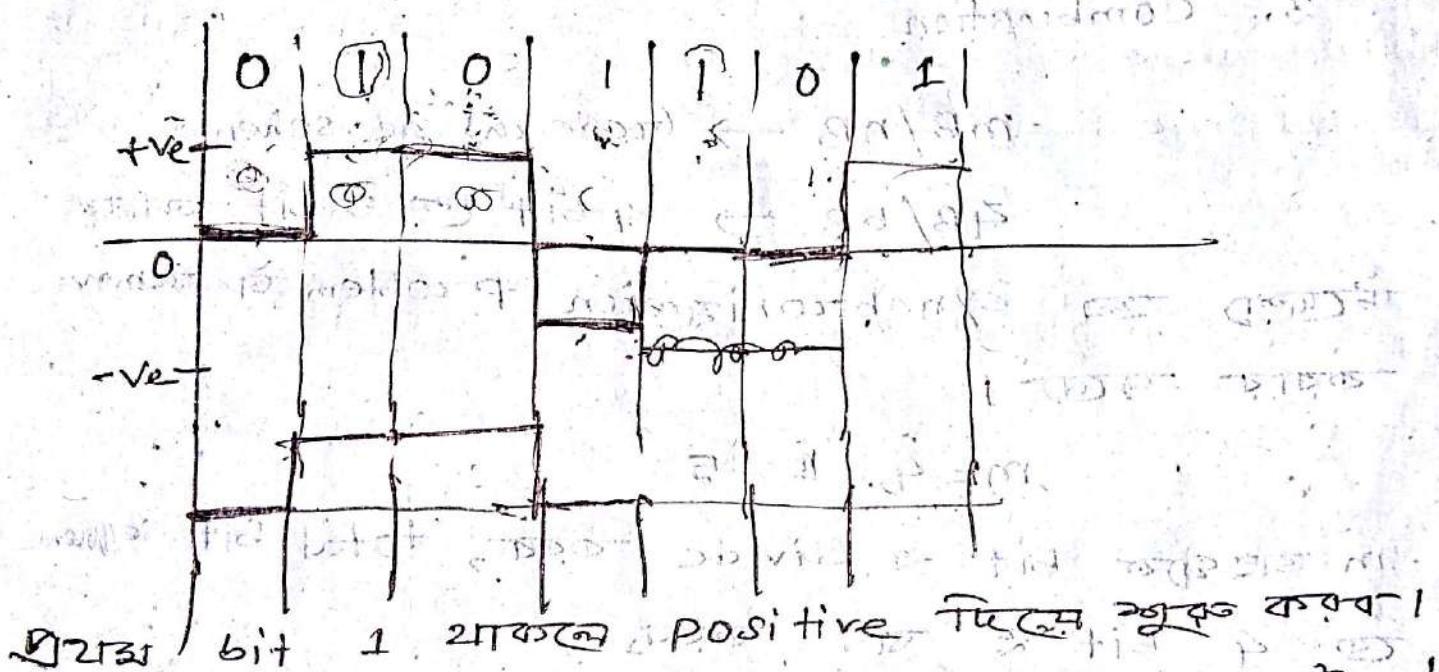
3 levels, (i) + V<sub>OH</sub>

(ii) - V<sub>OL</sub>

(iii) 0 V<sub>OL</sub>

3 conditions:-

1. If the next bit is 0, there is no transition.
2. If the next bit is 1 and the current level is not 0, the next level is 0.
3. If the next bit is 1, and the current level is 0, the next level is the opposite of the last non zero level.



17-07-17  
6th (B) day

Block coding: Digital to Digital conversion that means Digital data to Digital signal

Bit sequence  
4B/5B  $\rightarrow$  1010|1000|...

4 bit कर्ते जा सकते।

Step शैली हैः-

1. Division

2. Substitution

3. Combination

mB/nB  $\rightarrow$  Generalize scheme

4B/5B  $\rightarrow$  4 bit के 5 bit रूपान्वय

उद्देश्य सन्धारणा problem को remove कराए गए।

$$m=4, n=5$$

m अवधारणा bit  $\rightarrow$  divide करते, total bit sequence के 4 bit के जा सकते,

एवं प्रत्येक substitute करते, substitute m bit to n bit, 4 bit के 5 bit के रूपान्वय।

परन्तु 5 bit के sequence के combined करते network न होते।

4 bit कर्ते पाठ्यालय possible 16 के combination आते। 4 bit के 5 bit के रूपान्वय निम्नलिखित हैं।

(i) प्रथम 3 में 1 पाठ्यालयीय याकले परन्तु 0 रखा यास्त। यथाः 1111, लेधा यास 11101

(ii) परंपरा 3. से 0 थाक्ट्स 1 रेखा यास। फैशनः-  
0001 रेखा यास 00011

विट्टे table से देखते हैं।

Combined sequence जैसे NRZ-I APPLY होनी कहे digital signal कहे करवते। उपर्युक्त signal जैसे bit ने रूपान्तरित होय। 5 bit वाले sequence जैसे 4 bit ने रूपान्तरित करते हैं bit वाले यास।

Scrambling → करे जैसे रेखा, रेखा  
→ एको एको line coding scheme.

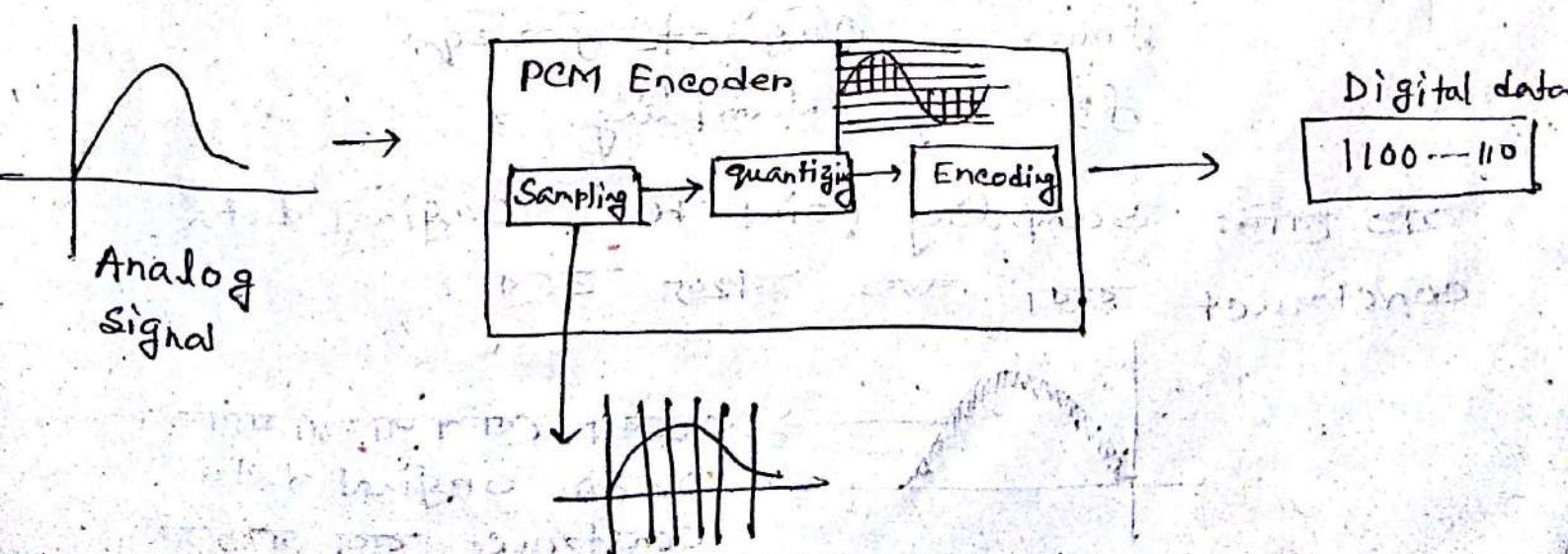
Line coding एवं scheme एवं advantage & disadvantages

देखते हैं।

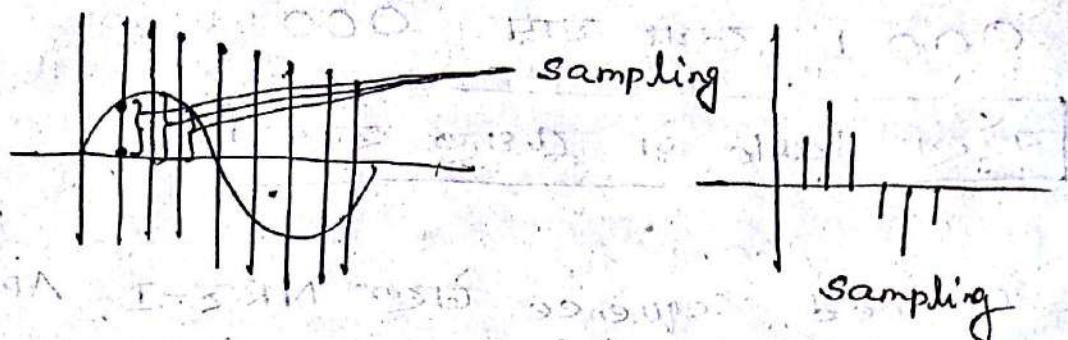
गोपनीय scheme एवं disadvantage आए, उठाले scheme कैसे remove करा यारे?

Analog to Digital conversion:-

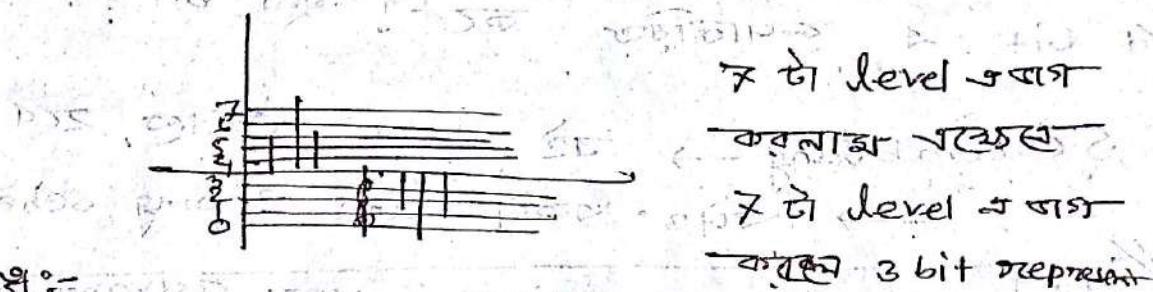
Pulse code Modulation (PCM) :-



Sampling: No. of collection sample.



Quantizing: Sampling के horizontally leveling करते। एकलेवल level के बीच इनमें से एक रिज्यूलेशन होता है।



Encoding:-

एकलेवल sample इनमें आहेत,

5 7 . 6 2 0 1

सातलेवल bit representation करते। आहे एजेंट Encoding

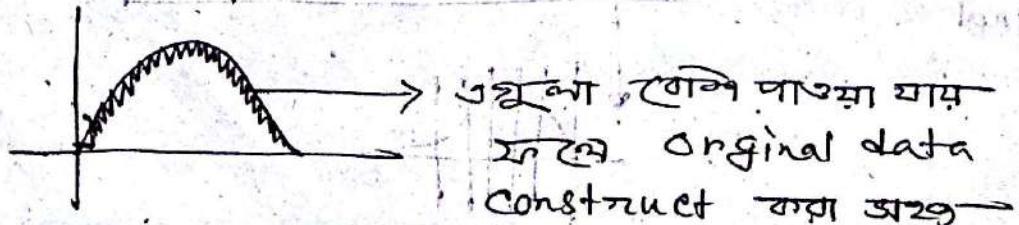
Sampling rate के लक्षण होते ताकि Nyquist theorem:- एडे theorem

Sampling Rate  $f_s \geq 2f_{\max}$

$f_{\max}$  = Original freq.

$f_s \geq$  Sampling

यात्रा किंवा sampling point मिळाले Original data construct करता आहे।

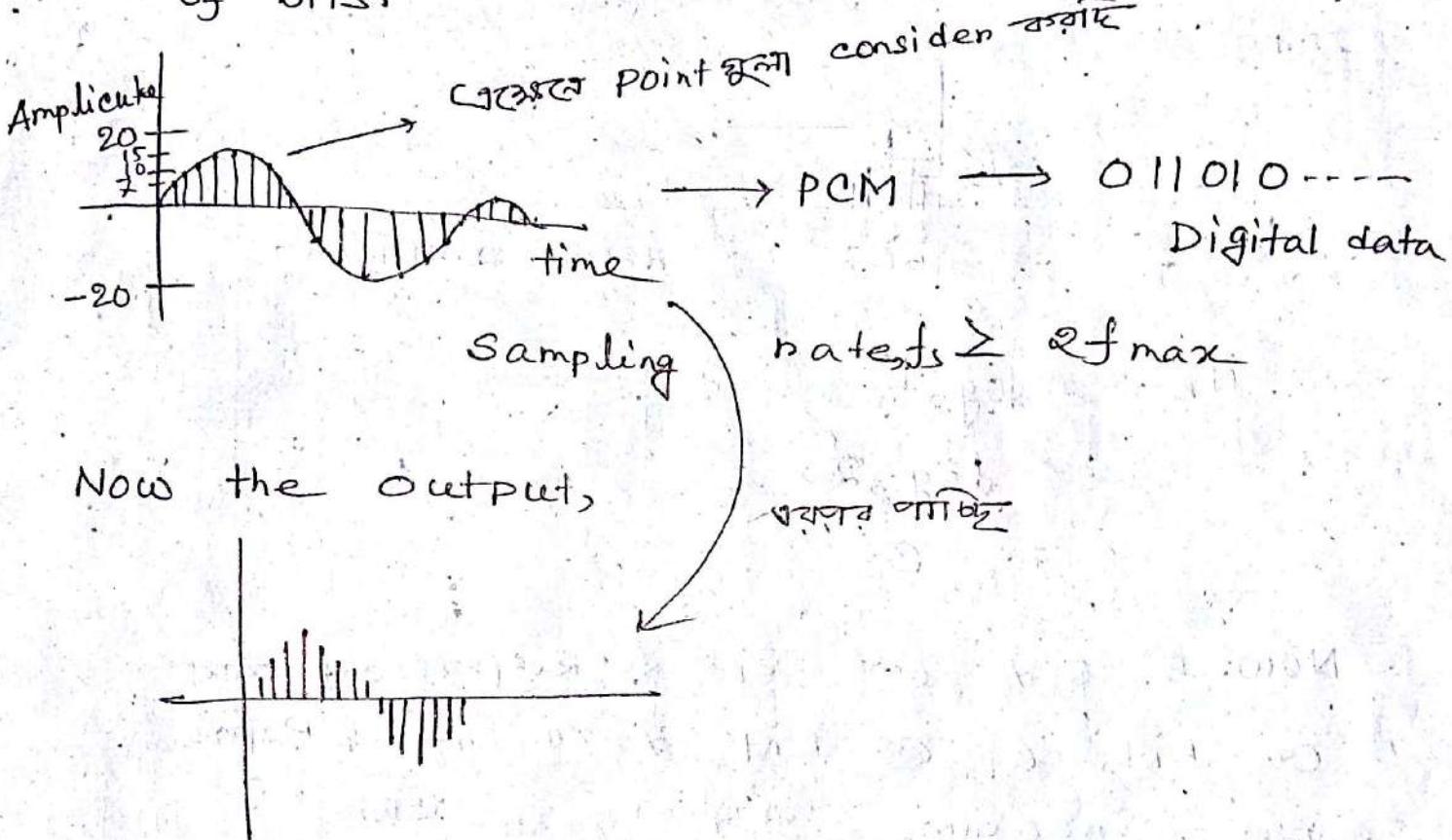


19-07-17  
6th (d) day

## Sampling

### Steps of PCM :-

1. The Analog signal is sampled.
2. The sample signal is quantized.
3. The quantized values are encoded as streams of bits.



### Steps of Quantization:-

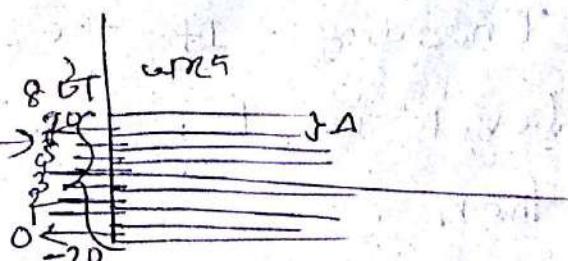
$L \rightarrow$  division

$$L = 8$$

$$A = \frac{V_{\max} - V_{\min}}{L}$$

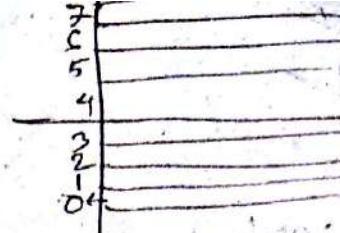
$$= \frac{20 - (-20)}{8}$$

$$= 5 \text{ V}$$



0-L-1 पर्याप्त masking करवा.

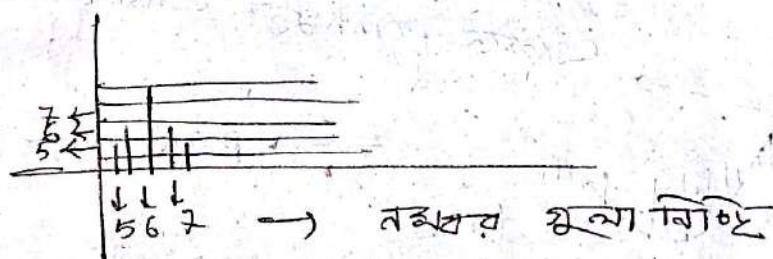
0-7 प्राप्ति याएँ।



$$\text{Normalization} = \frac{\text{amplitude}}{A} = \frac{20}{5} = 4$$

$\Rightarrow$  Quantization part code शुल्क का विषय होने वाले छेद-

जैसे



$$n_b = \log_2 L$$

$$= \log_2 2^3$$

$$= 3$$

Now 5, 6, 7 के bit व Represent करवा  
5 के bit, 6 के bit, 7 के bit व Represent  
करते Encoding प्रक्रिया याएँ।

Encoding:- If the number of quantization level is  $L$ , the number of bit is  $n_b = \log_2 L$ . then

Bit rate = Sampling rate  $\times$  number of bits per sample

$$= f_s \times n_b$$

We want to digitize human voice? What is the bit rate of human voice. Analog, एंडॉलॉगिक digital कारबत,

Ans:

Assume that 8 bits per sample.

$$f_s = 2 \times f_{\max} \\ = 2 \times 4000 = 8000$$

now, bit rate =  $f_s \times n_b$

$$= 8000 \times 8 \\ \approx 64000 \text{ bps.} \\ = 64 \text{ kbps}$$

= 0 =

$$\text{Minimum Bit, } B_{\min} = C \times N \times \frac{1}{n}$$

$$= C \times n_b \times f_s \times \frac{1}{n}$$

$$\approx C \times n_b \times 2B_{\text{analog}} \times \frac{1}{n}$$

$$= @ \frac{1}{2} \times n_b \times 2B_{\text{analog}} \times 1$$

$$\approx n_b \times B_{\text{analog}}$$

$$f_{\max} = 4000 \text{ Hz}$$

$$f_s = ?$$

$$C = \text{Case}$$

$$= \text{average case}$$

$$= 0.5$$

$$\frac{1}{n} = \frac{1}{p} = 1$$

1 bit 1 signal

फिल्टर

$$f_s = 2 \times f_{\max}$$

PCM द्वारा Minimum Band width,  $B_{\min} =$

$$n_b \times B_{\text{analog}}$$

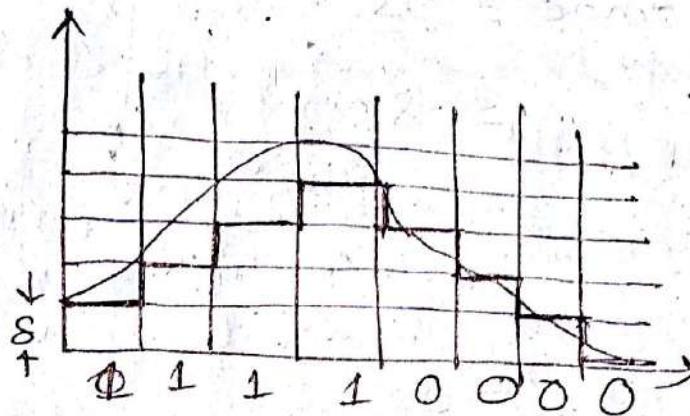
$$= 2 \times B_{\text{analog}}$$

আগের অংকটোর যদি  $B_{\text{analog}} = 4000 \text{ Hz}$   
 $n_b = 8$

$$\text{Minimum } B_w = n_b \times B_{\text{analog}} \\ = 32000$$

Minimum  $32000 \text{ Hz}$   $B_w$  পার্স দরবার  
 তা না এম  $\text{PCM}$  গে pass হবে না।

### ~~Da~~ Delta Modulation (ΔDM)



increase to  
 1 ক্ষণীয় denote  
 decrease to  
 0 ক্ষণীয় denote

Analog টাৰ digital version = 11110000

PCM = analog  $\rightarrow$  digital data  
 DM  $\rightarrow$  analog  $\rightarrow$  digital data

PCM টো ধুৰ complex way টো analog  
 এবতে digital গে পাৰিব কৰে।

29-07-17  
6th (E2day)

### Transmission mode

(a) Parallel

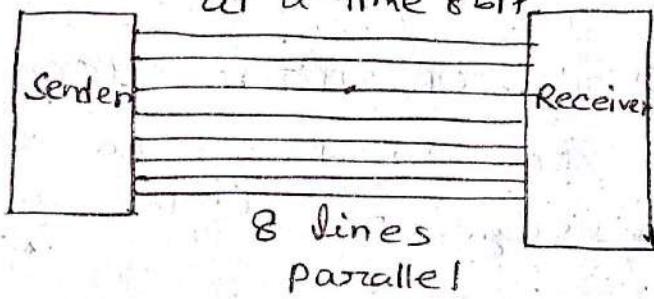
(b) Serial

i) Asynchronous

ii) Synchronous

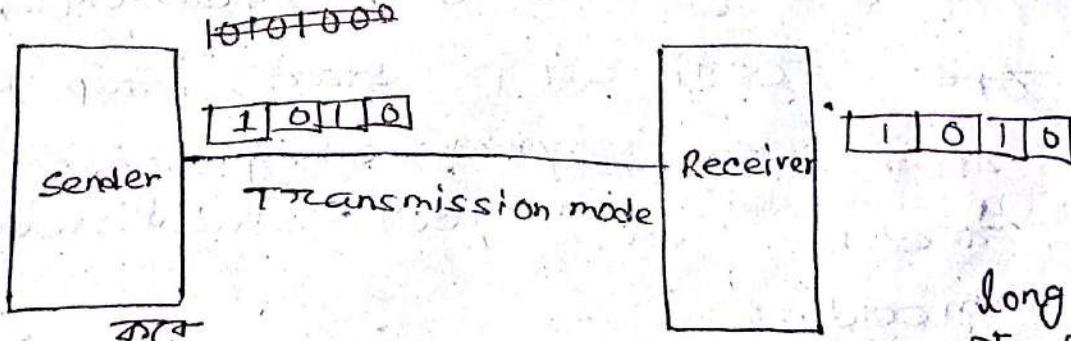
iii) Isochronous

(a) Sender कर्ते Reciever 6 at a time, 8 bit कर्ते  
पाठ्यक्रम होते, उन्हें parallel 8 की line करका,  
फिर cast द्वारा किसी data  
at a time 8bit



transmission fast  
short  
long distance  
use 24  
parallel

(b)



एकटे एकटे 1 bit यांसे serially

long distance  
use 24

① पहले start bit 1 then 8 bit send करते, then  
एकटे stop bit पाठ्यक्रम, then एकटे gap करते  
then आगे आगे करते 8 bit, जैसे पाठ्यक्रम है

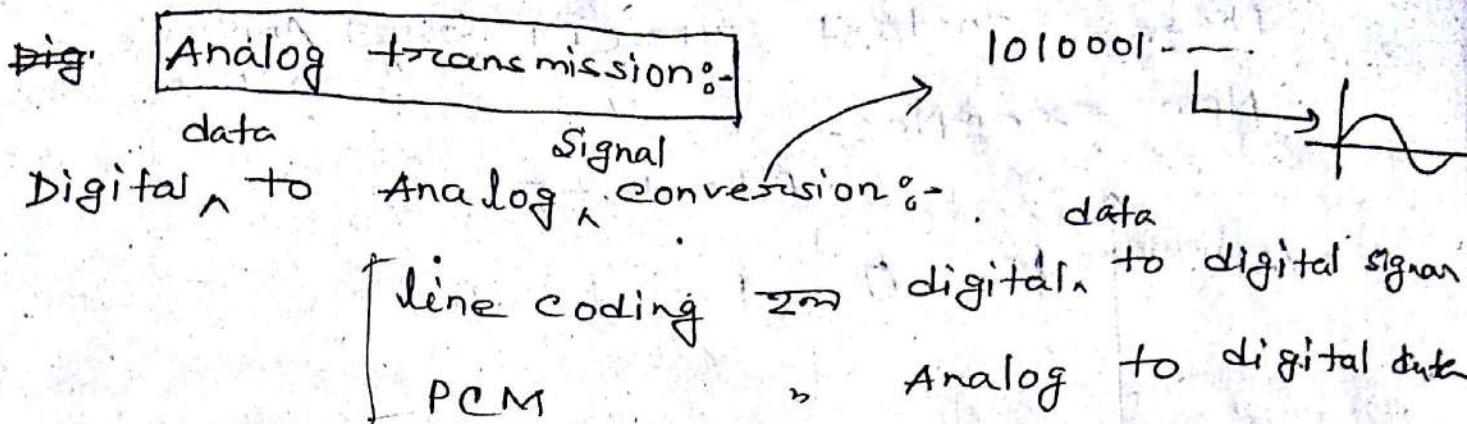
- ⇒ data sender द्वारा Receiver missing जुड़ाव अप्पावता कर।
- ⇒ यदि missing है, उपर आप message देखावास अनुबिधः पाण्डा है तो, यदि missing हो जाएँगे याहे।
- ⇒ ④ इसे slot व एक extra करे bit add कर सynchronization भाकर ना बदलें। यह नाम Asynchronous.

(ii) Start, stop bit, gap आकर्ष ना रखे एकटे करे bit पाण्डा है। synchronization अवश्य दूर करे है एवं शारीर्य।

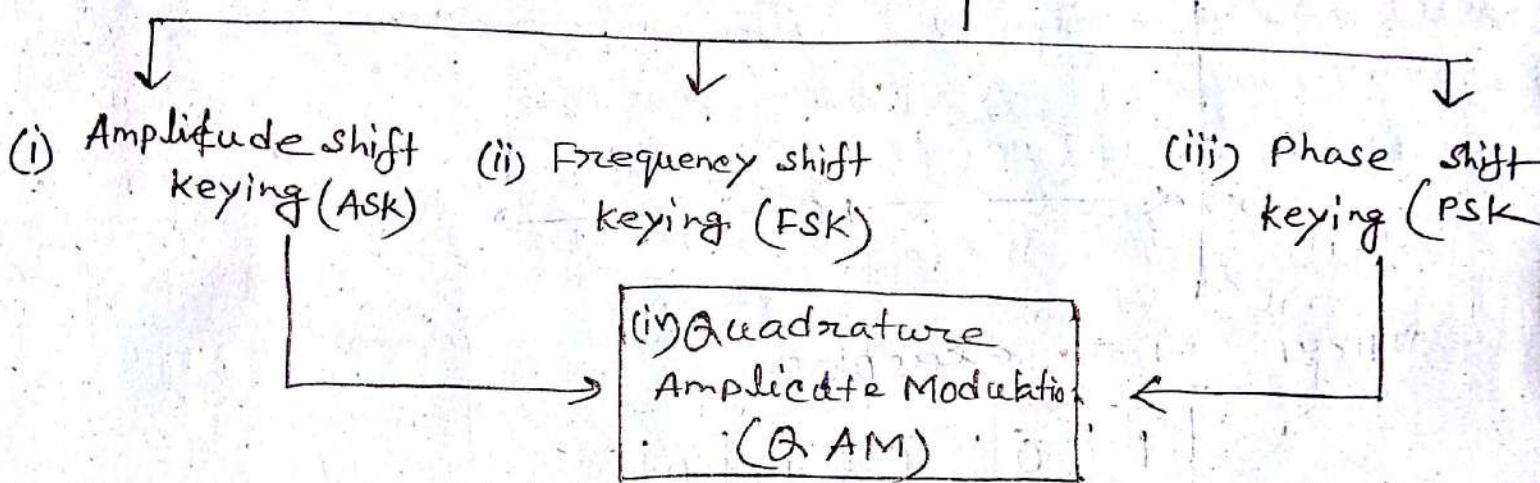
जो bit इन्हे ना आकर्ष करने वाले message दे जावाब send जाएँ हैं, यदि message न करने missing आएँ।

(iii) Synchronous व Asynchronous एवं अनुबिधः सही है। Start, stop bit व gap आकर्ष नये synchronization एवं maintain करे, एवं well, established transmission mood.

## ~~ii~~ Analog



3 methods for digital data to analog signals:

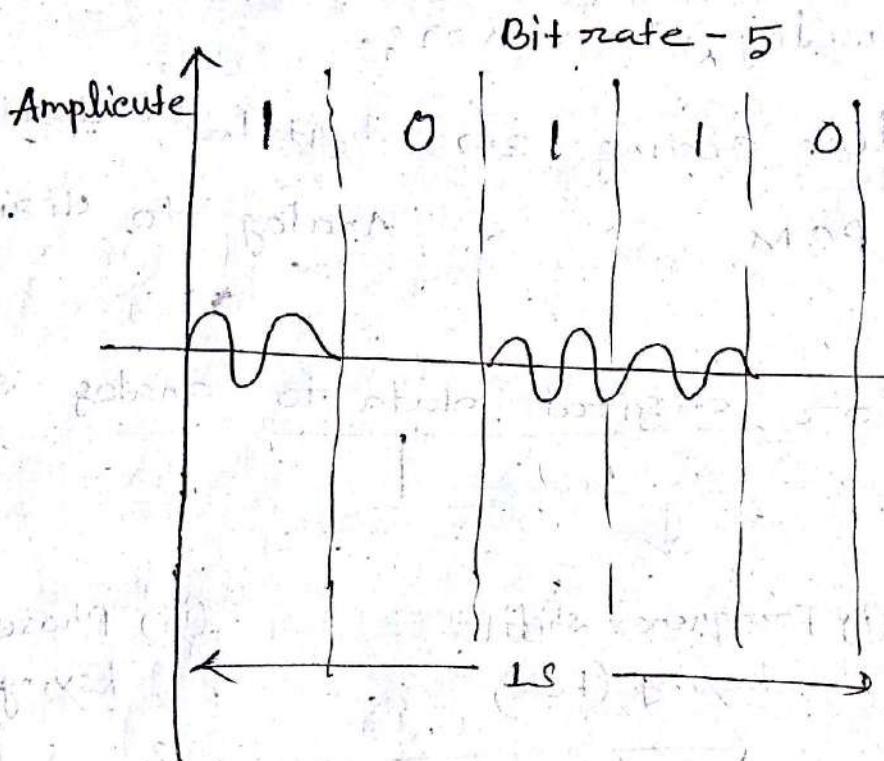


Carrier Signals: In analog transmission, the sending device produces a high frequency signal that act as a base for the information signal. This base signal is called carrier signal or carrier frequency.

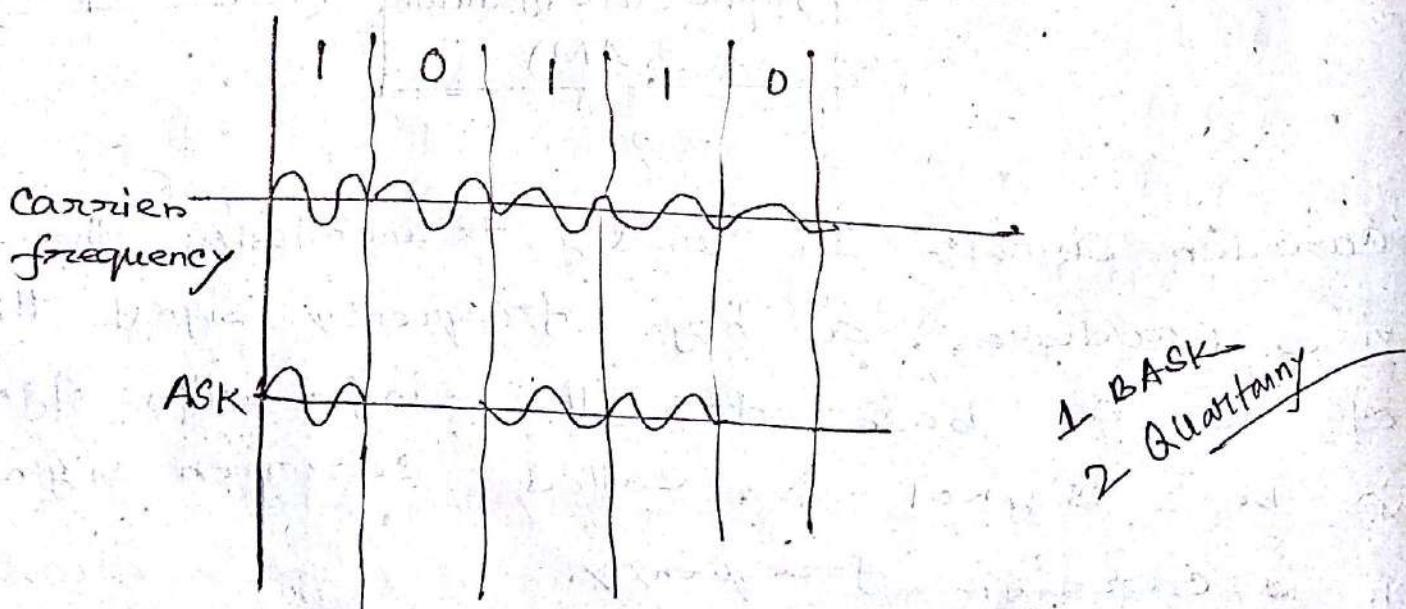
- (i) Amplitude
- (ii) Frequency
- (iii) Phase

→ यहाँ से जो भिन्नी लाई जाए तो उसका नाम बनता है।

(i) यदि Amplitude change हो, frequency नहीं.  
 Phase constant रहता है। यदि Amplitude हो,  
 Now example:-



आर्थिक उदाहरण example:-



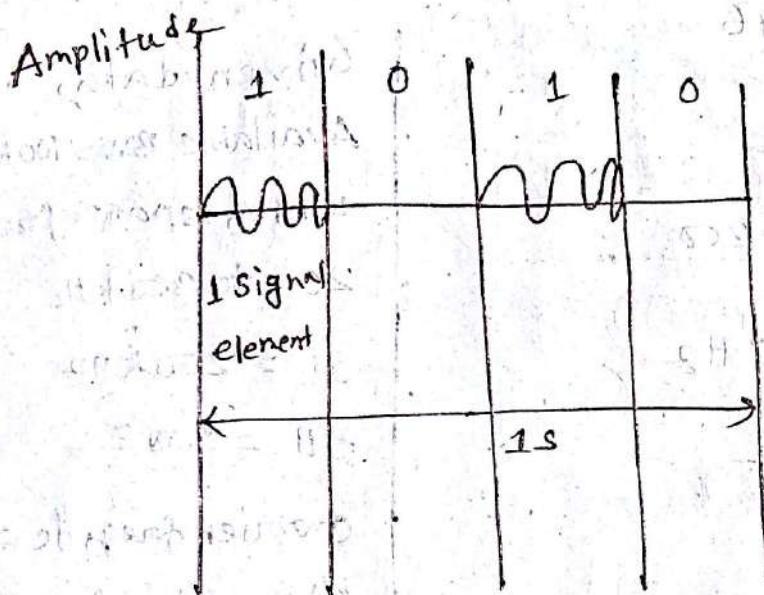
यदि एक बारे में Binary ASK. At a time 1 bit  
 करने send करे जाए। At a time 2 bit करने  
 send करने करने Quarternary.

31-07-2017  
7th (B) day

## ASK

Binary ASK: Binary ASK is normally implemented using only two levels. This is referred to as binary Amplitude shift keying (BASK) or on-off keying (OKK).

- The peak amplitude of one signal level is 0; the other is the same as the amplitude of the carrier frequency.



2<sup>nd</sup> level of binary ASK

Signal rate vs और bit rate का relationship

$$S = N \times 1/r$$

$r = 1$  bit प्रति संकेत  
वाली signal passen

(ratio)

$$S = N \times 1/r$$

$$B = (1+d)S$$

$$d = 0 \text{ to } 1$$

Worst signal

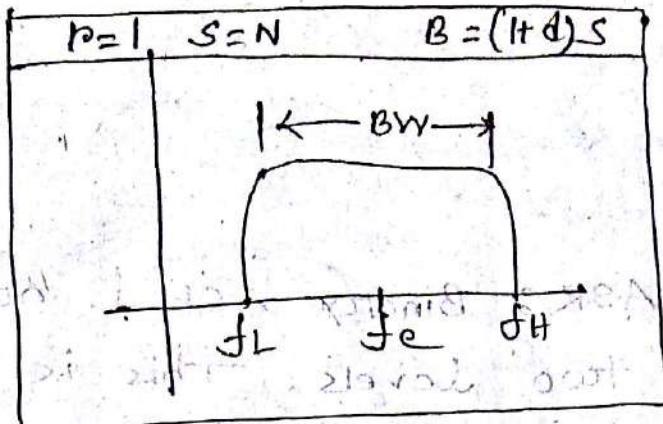
$$d = 0$$

Best signal (G)

$$d = 1$$

Avg. case G d = 0.5

normally d = 1



Example:-

Page-146

$$\begin{aligned} f_c &= \frac{f_L + f_H}{2} \\ &= \frac{200 + 300}{2} \\ &= 250 \text{ Hz} \end{aligned}$$

$$S = N \times \frac{1}{r}$$

Here,

$$B = (1+d)S$$

$$= (1+1)S$$

$$\text{bit rate} = \frac{2S}{2 \times N \times \frac{1}{r}}$$

$$1000 = 2 \times N \times \frac{1}{r}$$

$$\therefore N = 500 \text{ kbps}$$

Given data,

Available BW = 100 kHz  
which spans from  
200 to 300 kHz

$$f_L = 200 \text{ kHz}$$

$$f_H = 300 \text{ kHz}$$

CARRIER freq,  $f_c = ?$

Bit rate if we modulate  
our data by ASK  
(with  $d=1$  and  
 $r=1$ )

bit rate,  $N = ?$

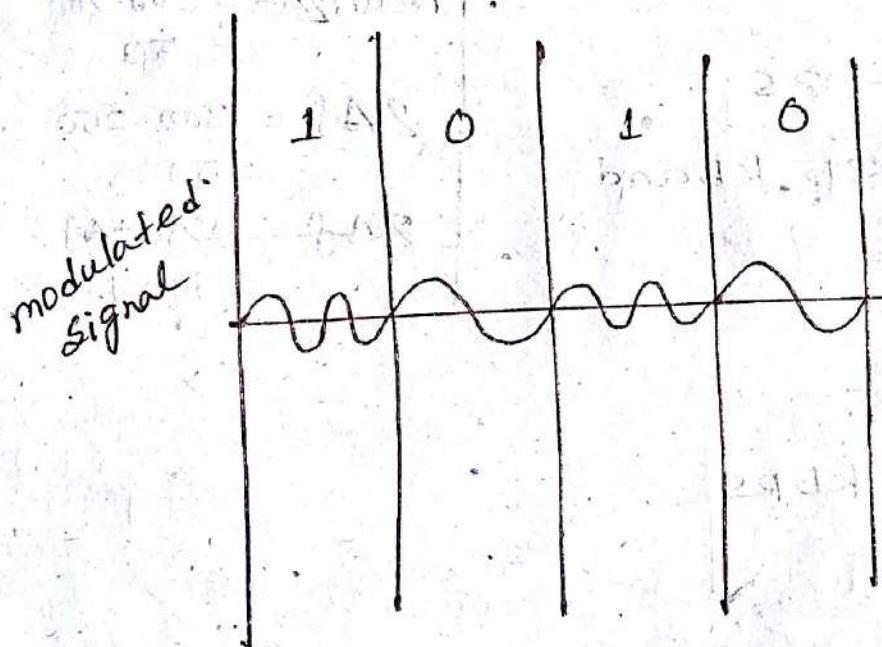
## Frequency Shift Keying (FSK) :-

In FSK, the frequency of the carried signal is varied to represent data.

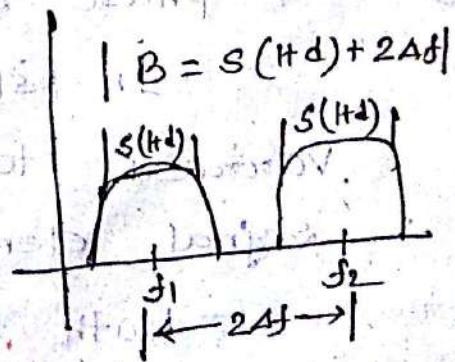
### Binary FSK (BFSK) :-

- Consider two carrier frequencies  $f_1$  and  $f_2$
- We use the 1st carrier if the data element is zero (0)

We use the 2nd carrier if the data element is one (1)



$$B = S(Hd) + 2Af$$



$$f_2 - f_1 = Af$$

Example

$$B = (1+d)S + 2Af$$

$$S = N \times \frac{1}{r}$$

$$B = (1+1)S + 2Af$$

$$= 2S + 2Af$$

$$100 = 2S + 2Af$$

$$\Rightarrow 100 = 2S + 50$$

$$\Rightarrow 100 - 50 = 2S$$

$$\therefore S = 25 \text{ Kband}$$

$$S = N \times \frac{1}{r}$$

$$\Rightarrow S = N \times \frac{1}{r}$$

$$\therefore N = 25 \text{ kbps}$$

$$BW = 100 \text{ kHz}$$

$$\text{Span} = 200-300 \text{ kHz}$$

What should be the carrier frequency,  $f_c = ?$

Bit rate = ?

If we modulated our data by FSK with  $d=1, r=1$

$$f_c = \frac{200+300}{2} \\ = 250$$

$$\text{Now, } 2Af = 300 - 200 \\ = 100$$

$$2Af = 300 - 250 \\ = 50$$

$$2Af = 2Af + Af \\ = 100 \text{ Hz}$$

02-08-17  
7th (D) day

Phase shifting keying (PSK):-

- In PSK, the phase of the carrier signal varied to represent two or more different signal elements.

- Both peak amplitude and frequency remain constant as two phase changes.

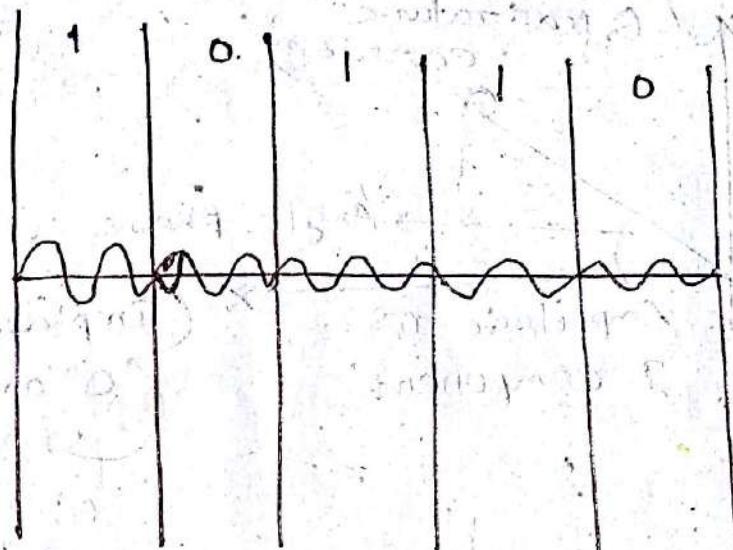
## Binary PSK (BPSK) :-

- we have only two signal elements, one with a phase  $0^\circ$  and the other with a phase  $180^\circ$

\*\*\* PSK और FSK उ अन्यान्य इनका एकत्र उल्लंघन का साधन आहे.

Amplitude व Frequency change करा असणे किंवा Phase change करा करिता, फले PSK दो best. Phase के define करा आता, Noise द्वारा का प्रभावित रप्पी, ताते less susceptible (अवैज्ञानिक).

## QPSK (Quadrature phase shift keying) :-



इते विवरण phase

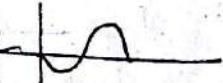
व आदे आटे

पठाणे BPSK

$0^\circ$  &  $180^\circ$

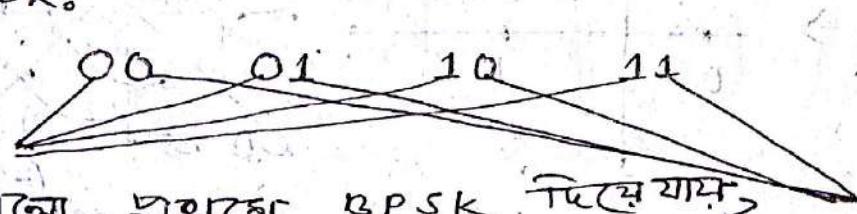
$1 \rightarrow 0^\circ$  phase

$0 \rightarrow 180^\circ$  phase

$180^\circ \rightarrow$  

$0^\circ \rightarrow$  

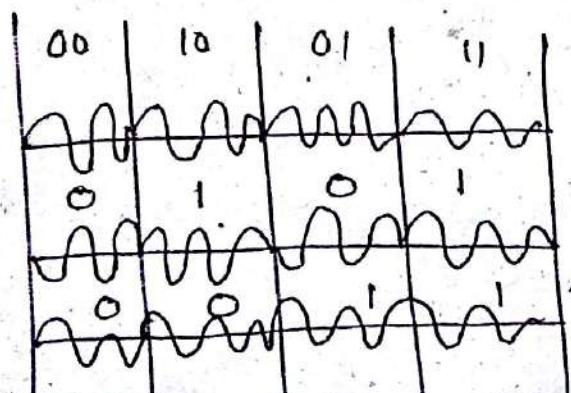
## QPSK:



एवढा प्रधारण BPSK ठिक्यायाप्प,

पठाणे पठाणे combined करा इत्ता । 1001 + 1100

पठाणे एवढा BPSK



$1 \rightarrow 0$  phase  
 $0 \rightarrow 180^\circ$  phase

पठाणे एवढे combined इत्ता ;

द्वारा 4 घे Phase पाठ्य →  $45^\circ, 135^\circ, -45^\circ, -135^\circ$

00 →  $-135^\circ$

10 →  $-45^\circ$

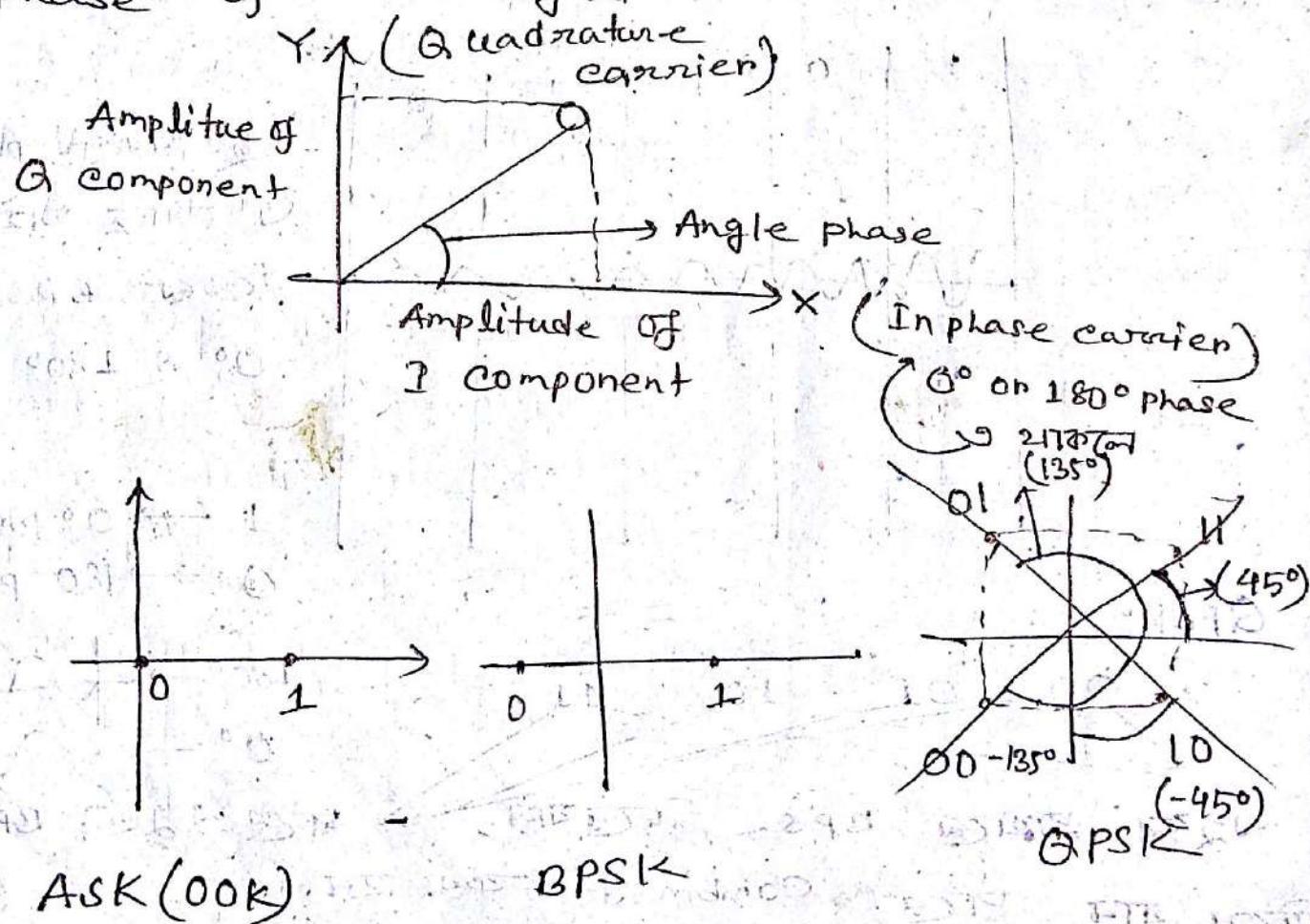
01 →  $135^\circ$

11 →  $45^\circ$

वर्तन द्वारा एक आदि एको diagram.

### \* Constellation Diagrams

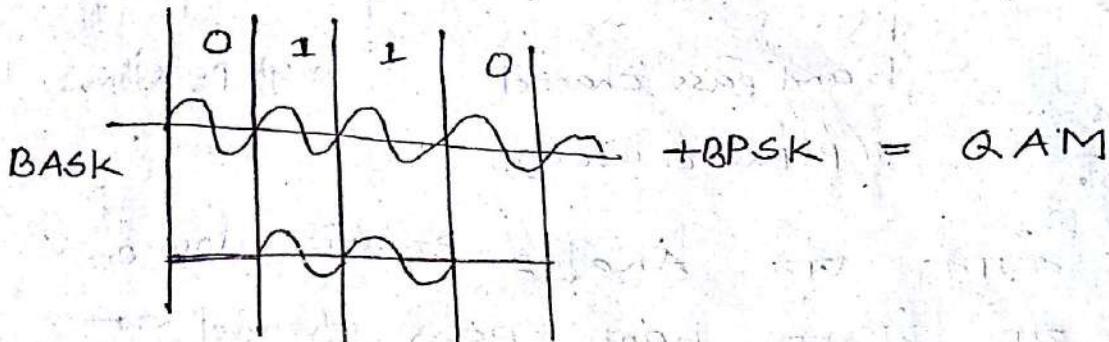
— helps us to define the Amplitude and phase of a signal element.



05-08-2017  
7th (E) day

QAM → Quadrature Amplitude Modulation  
→ Combination of ASK & PSK

Amplitude वरे Frequency द्वारा बदलने के change होते।



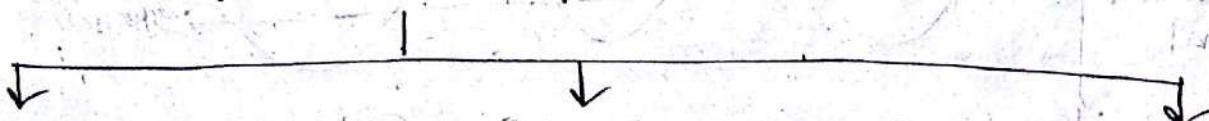
जुटाविंग:-

- i) efficient. योग्यता
- ii) Noise छाड़ा अणवित होता करता
- iii) Strong signal के पास उपर्युक्त योग्यता
- iv) information के loss असंभव possibility करता

अड्डाविंग:-

- i) उपर्युक्त complex, नक्कल pattern  $\rightarrow$  device योग्यता current, voltage योग्यता लाए।

Analog to Analog conversion:-



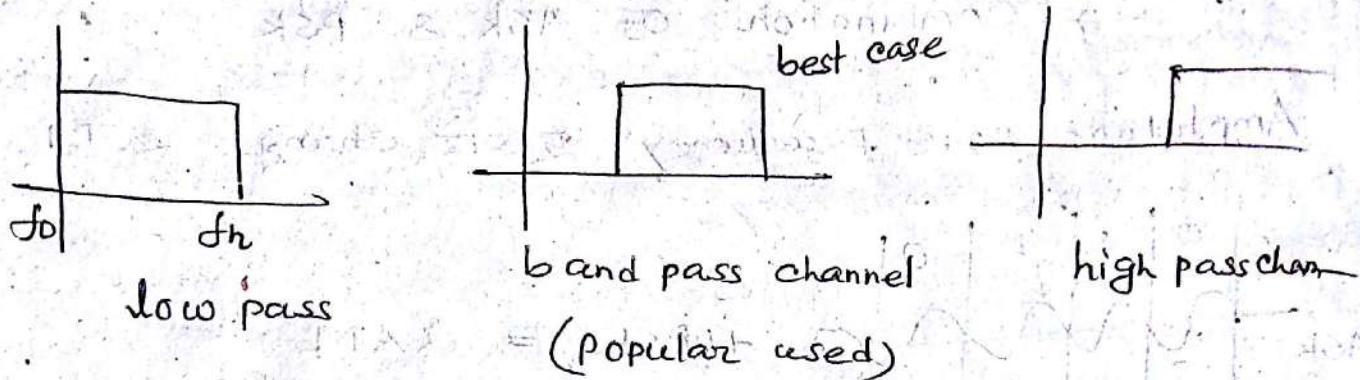
- i) Amplitude Modulation (AM)
- ii) Frequency Modulation (FM)
- iii) Phase Modulation (PM)

Amplitude vary वाले 2 से ज्यादा ग्राहक आवश्यक

Phase vary वाले याकेपुरी लिए आवश्यक

## Q) Modulation कैने करका?

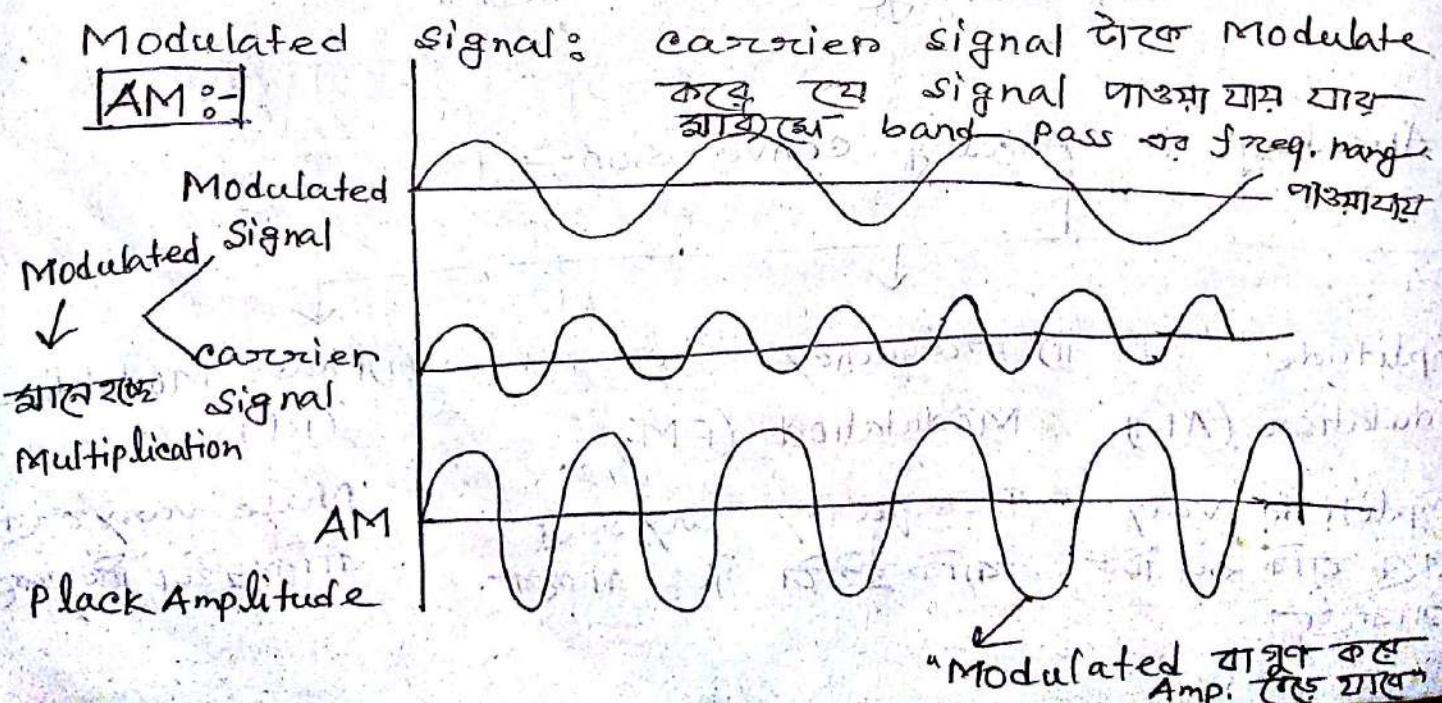
channel एवं bandwidth एवं आरें related, इसले कहा, band pass channel



Modulation करका पर एवं Analog से यहाँ low or high pass एवं यहाँ band pass channel एवं बार्कुप pass करके low वृ high एवं बार्कुप range एवं अंतर्भूत याद।

\*\* Analog signal के लिए modulation करका यहाँ band pass channel एवं range एवं range की तिक्टोड़ आज्ञा याद, ताकि हमें best case पाओगा याद।

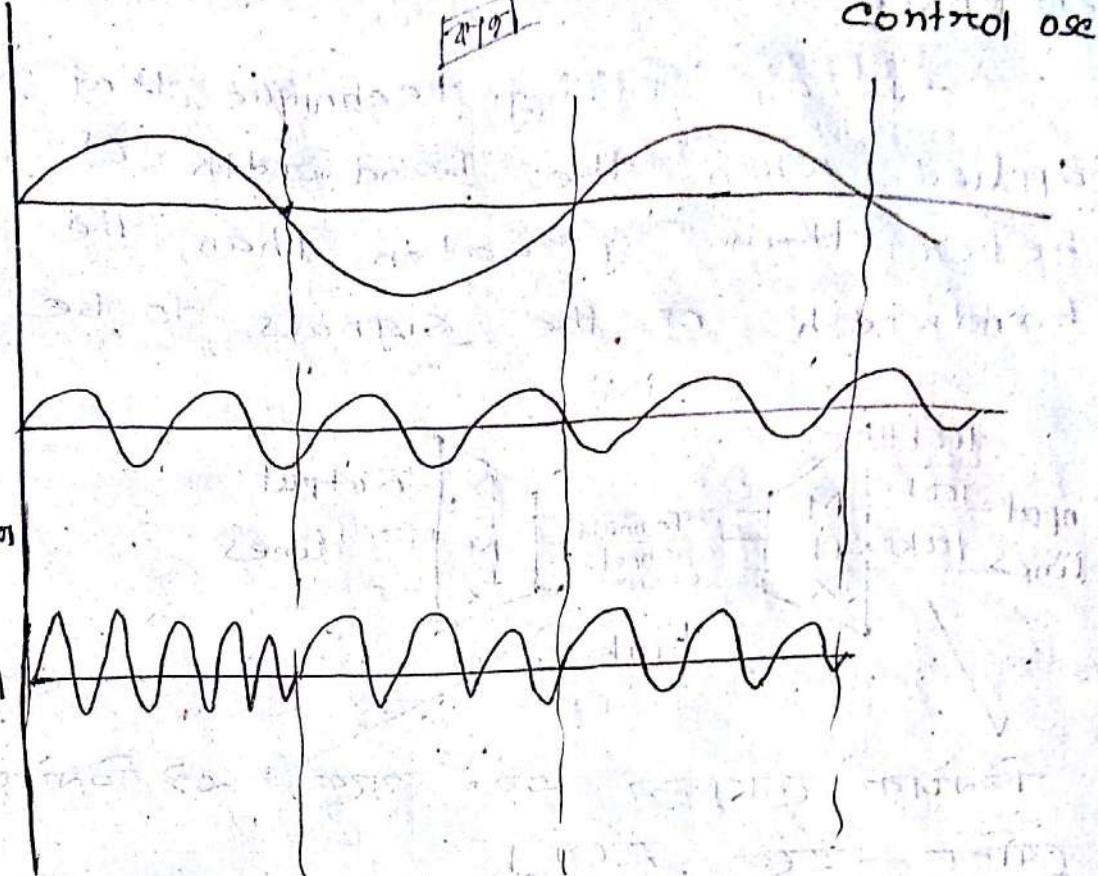
$\Rightarrow$



**FM**

VCO  $\rightarrow$  voltage control oscillator

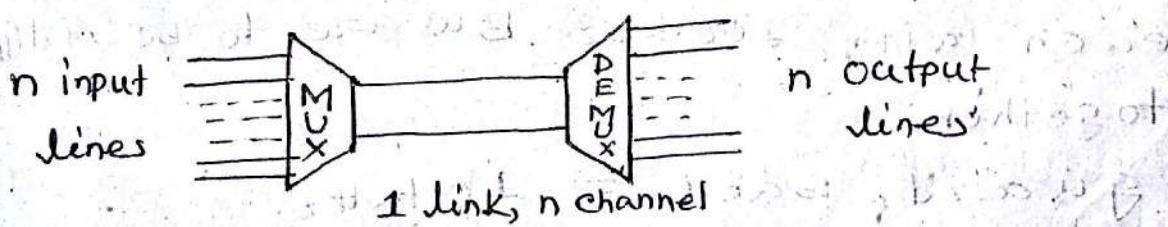
Modulated signal  
Modulated carrier signal  
VCO Pass band



07-08-2017  
8th (B) day

## MULTIPLEXING :-

— Multiplexing is a set of techniques that allows the simultaneous transmission of multiple signals across a single data link.

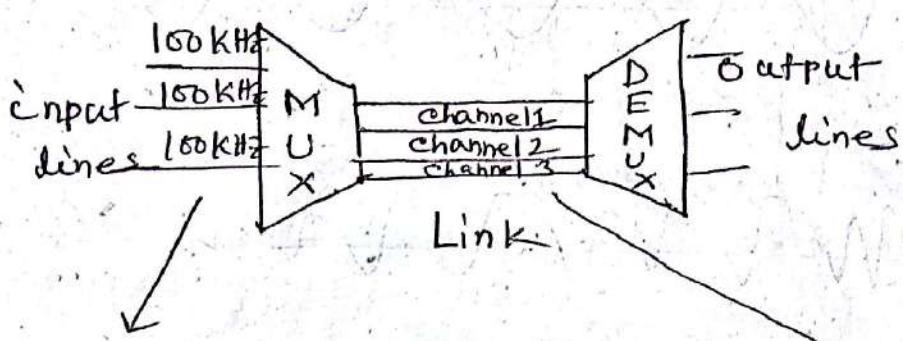


For further multiplexing techniques  $\Rightarrow$

- (i) Frequency division (ii) wavelength division (iii) Time division
- (Analog) Multiplexing (FDMA) Parallel Multiplexing (WDM) Multiplexing (TDM)

(i) FDM:-

- is an analog technique that can be applied when the bandwidth of a link is greater than or equal to the combined bandwidth of the signals to be transmitted.



তিনটির যোগফল এক রেখে  $\Rightarrow$  এই তিনটির যোগফল  
ব্রেডিং হত হবে।

গুড়ি Guard band  $\rightarrow$  কাঁচে interface overlapping  
না হওয়ার কাবলে channel এর মাঝে স্থাপিত  
Guard band use হয়। Guard band  $\rightarrow$   
এক চিহ্নের band width.  
যদি channel  $n$  হ, তবে Guard band  $n-1$

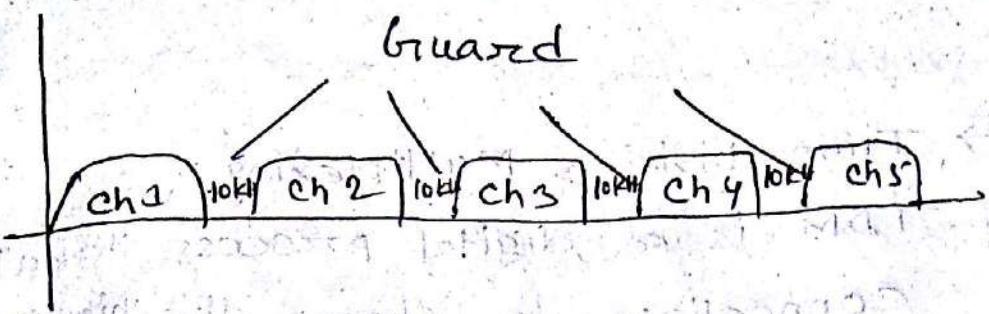
Example:- page  $\rightarrow$  164 159

5 channel  $\rightarrow$

each with 100 kHz BW are to be multiplexed  
together

guard band = 10 kHz

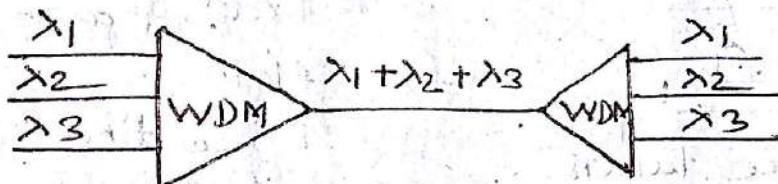
Minimum band width = ?



$$\begin{aligned}\text{Minimum bandwidth} &= 5 \times 100 + 4 \times 10 \\ &= 540 \text{ kHz}\end{aligned}$$

### WDM :-

- WDM is designed to use high data rate capabilities of fiber optic cable. The optic fiber cable data rate is higher than the data rate of metallic transmission cable.



WDM technique follow ~~prism~~ prism

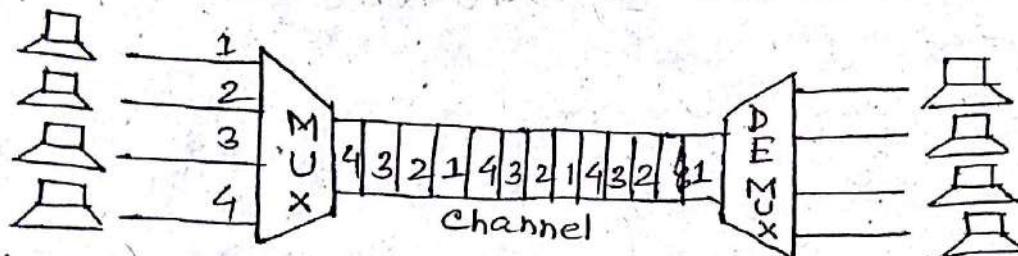
\* Higher efficiency than FDM

(n-1)

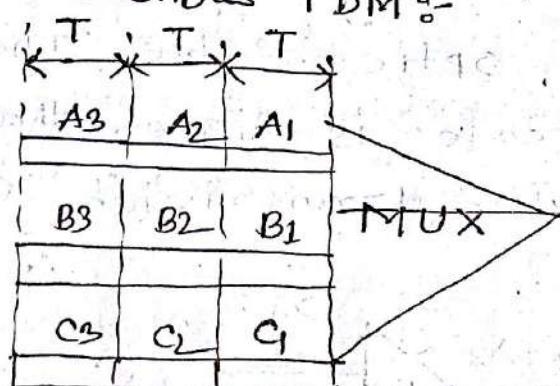
12-08-2014  
8th (E) date

TDM  $\rightarrow$  Time Division Multiplexing

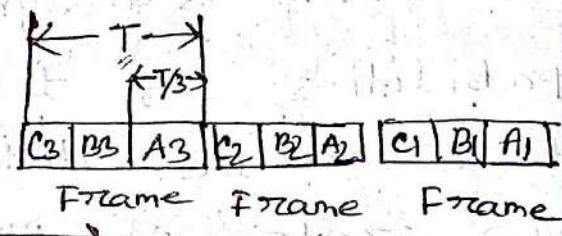
- TDM is a digital process that allows several connections to share the high bandwidth of a link.



(i) Synchronous TDM :-



Data are taken from each line every  $T_s$



Each frame is 3 time slots

Each time slot duration is  $T/3$

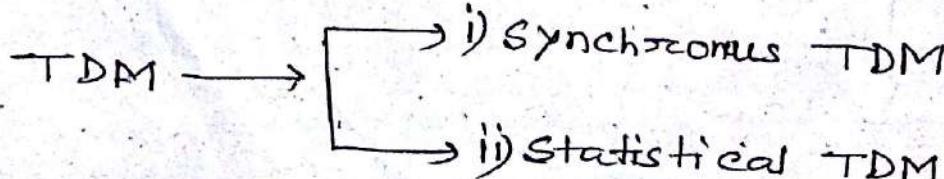
TDM  $\rightarrow$  ज्ञात share the high bandwidth, digital pro

FDM  $\rightarrow$  " " " " " नाइ, Analog

Data Rate = Bit rate

WDM ओ " "

BW वाले Bit rate येकु प्राप्त, एवज- एकात्म time slot  $\hookrightarrow$  उपर करा हम । ऐसा एको TDM बना हम ।



(i) n अंधक device याकले n अंधक time slot पर जग करते। २ यह छिपे तो device ताके ३ टा time slot ए जग रहे।

प्रथम  $\boxed{C|B_1|A_1}$  frame टो send होना चाहिए जो २ से A, then B<sub>1</sub> then C<sub>1</sub> send होता। ऐसे frame टो send होना T time में जगले A<sub>1</sub> or B<sub>1</sub> or C<sub>1</sub> send करना प्रत्यक्ष time जागता  $T/3$ .

एकको slot वा time division =  $T/n$

n = Number of devices

Transmission और उपयोगी use करता है TDM, जिसे बताते हैं at a time data send होते वे then at a time data receive होते

उपयोगी user वे दो जाके ३ टो slot पर जगते हैं

BW टोल्ड याके उपयोगी प्राविधि वाले यास Multiplexing Technique use करता है

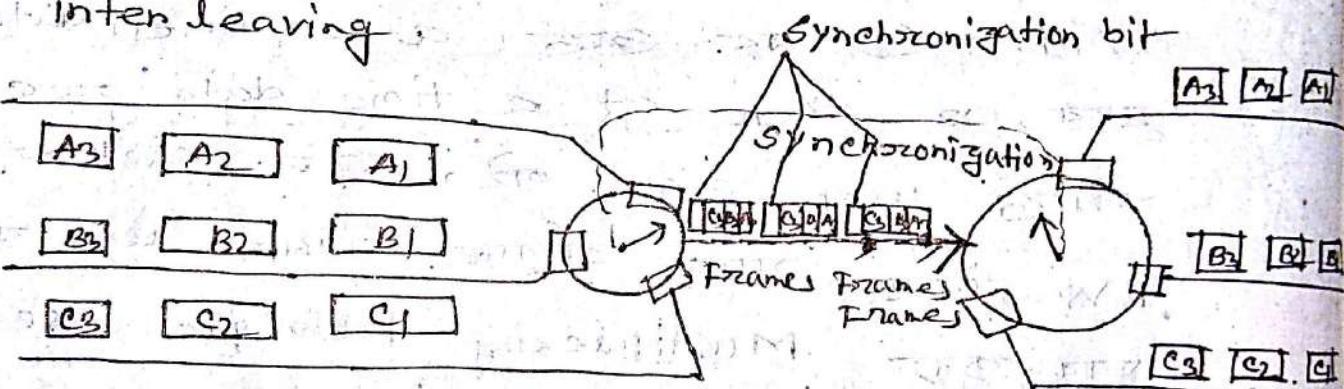
**वही घटक (i) का Math गुणान देखते हैं।**



20-08-17  
9th (D) day

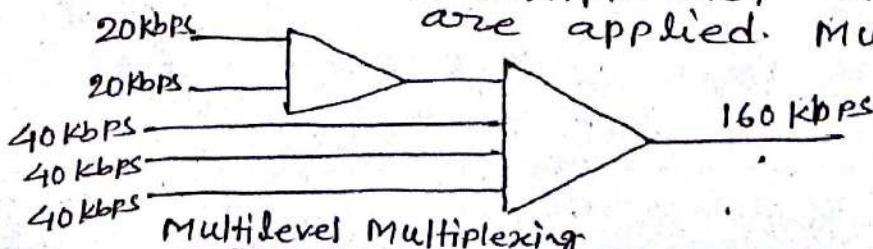
## Interleaving:-

- TDM can be visualized as two rotating switches, one on the multiplexing side and the other on the demultiplexing side. The switches are synchronized and rotate at the same speed, but in opposite direction.
- On the multiplexing side, as the switch opens in front of a connection, that connection has the opportunity to send a unit onto the path. This process is called interleaving.

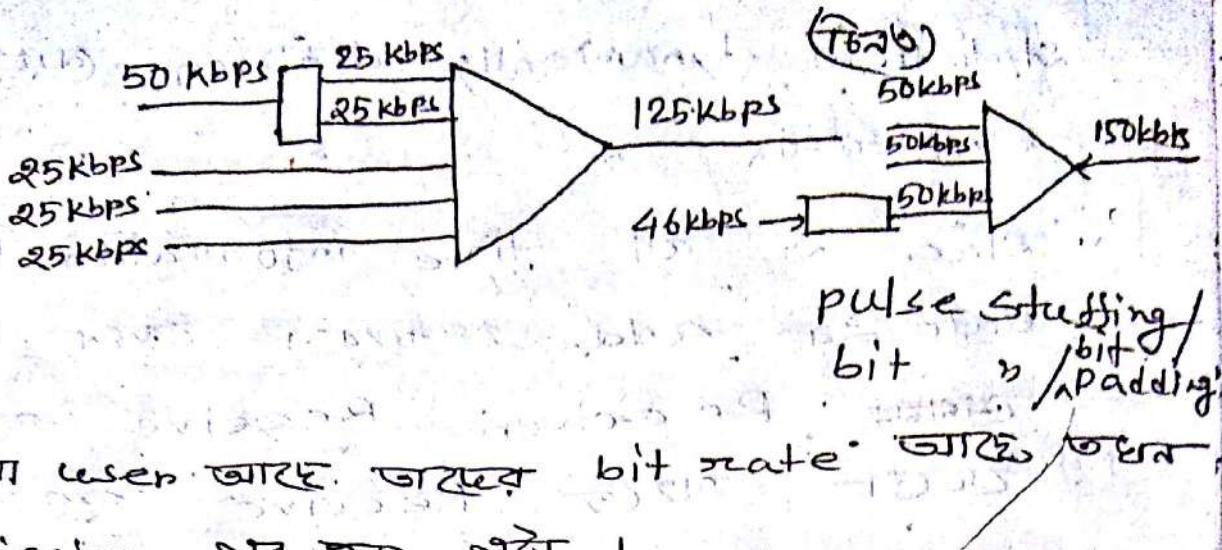


এতেকসে frame এবং আর্থে synchronization bit আর padding করে আগবং synchronization configuration করতে পারি।

Multilevel Multiplexing:- If data rates are not same, the three multiple slot allocation strategies are applied. Multi-level Multiplexing is one of them.



(क्रिय 2)



pulse stuffing  
bit → bit padding

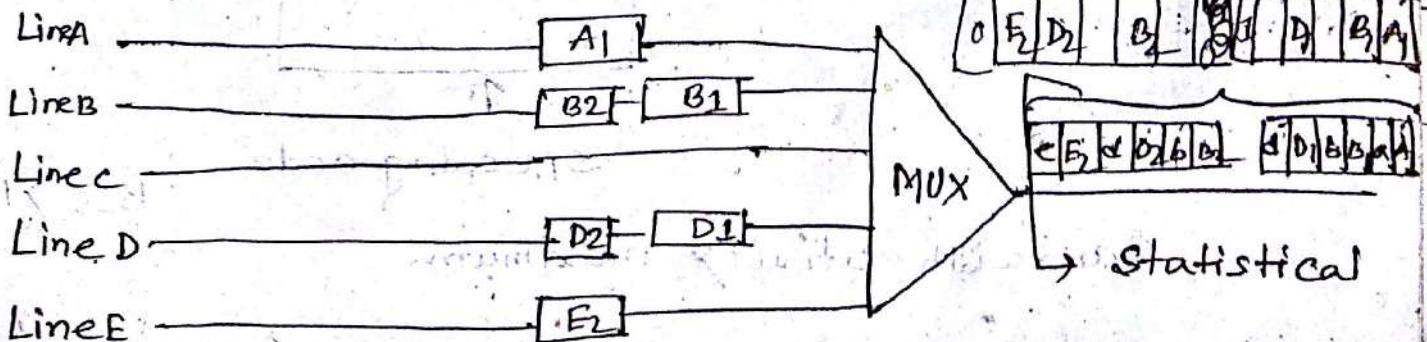
$$\text{पुराने में } 20 \text{ kbps} + 20 \text{ kbps} = 40 \text{ kbps} \text{ करने वाले थे।}$$

इन्हीं ने 50 kbps को demultiplexing करते हुए 25 kbps + 25 kbps करने वाले transmit करते हैं।

इन्हीं ने merge करते हुए अवास करते हैं और उन्हें अवास करने के लिए extra bit add करते हैं जिसका नाम 46 bit stuffing / bit padding है।

इन्हीं ने merge करते हुए अवास करने के लिए extra bit add करते हैं जिसका नाम 46 bit stuffing / bit padding है।

### Statistical TDM :-



Statistical dynamically decision विषय

Synchronous

Stas. & all time information विषय corresponding data एवं add. आकर्षणीय होता है आपके लिए.

Receivers Receive करते हैं अपने add. address और Receivers Receive करते हैं पास होते हैं।

Round Robin scheduling. इधार उसे अप्पे देखा जाएगा apply करो इस।

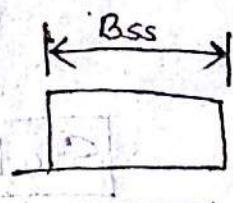
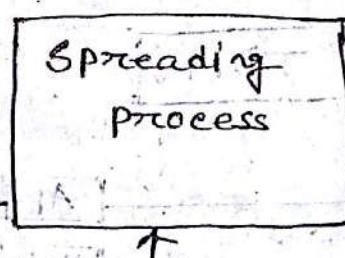
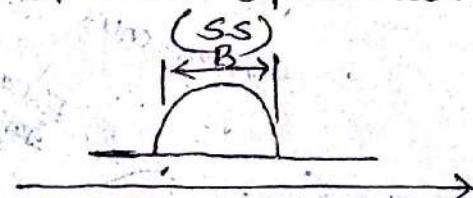
नियमित frame bit add करता चाहते synchronisation एवं शुद्धीकरण।

पास होते हैं। ~~syn~~ statistical & frame size fixed, वही syn. & frame size fixed.

Synchronous उपलब्ध समय में विपरीत।

21-08-2017  
9th (E) day

Spread Spectrum:-



spreading code

$B_{ss} > B$  (source BW)

Bandwidth utility maximum

SS achieves its goal through two principles:-

- 1) The bandwidth allocated to each station needs to be larger than what is needed. This allows redundancy.
- 2) The expanding of the original bandwidth  $B$  to the bandwidth  $B_{SS}$  must be done by a process that is independent of the original signal.

Two types:-

- 1) Frequency Hopping Spread Spectrum (FHSS)
- 2) Direct Sequence Spread Spectrum (DSSS)

Spread spectrum :-

उद्देश्य medium इन wireless, ज्यान तरं Receiver access पास, ज्यान Bandwidth को high करते हुए।  
याकृति कार्यालय interfrence को हटते, उद्देश्य नहीं है।  
spread spectrum के लिए goal को achieve कराते हुए principal के लिए यहां क्या कहा। कोई secure technique का उद्देश्य original code का आरथ करने को code division कार्यालय नहीं।  
यहां sender व receiver कार्यालय पास है। carrier Now two types बताये जाना:- frequency को mapping करते हुए उपर कहा है।  
मूल hacker कार्यालय नहीं, जबकि carrier

- 1) FHSS:- FDM के आरथ बनवाते हुए, जबकि carrier frequency use करते।

$K$  = spreading code

$M = 2^K$  = carrier frequency

K bit pattern

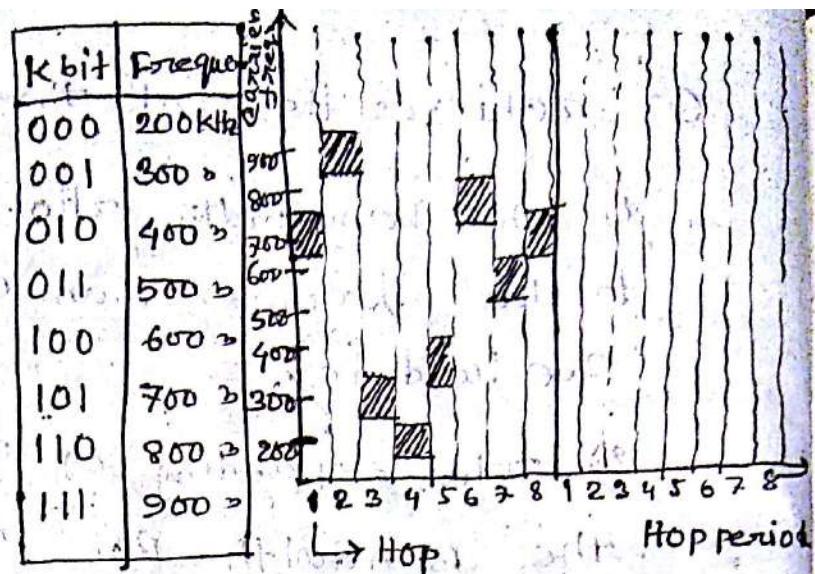
$$K=3 \rightarrow 7 \text{ bits}$$

101	111	001	000	010	110	011	100
-----	-----	-----	-----	-----	-----	-----	-----

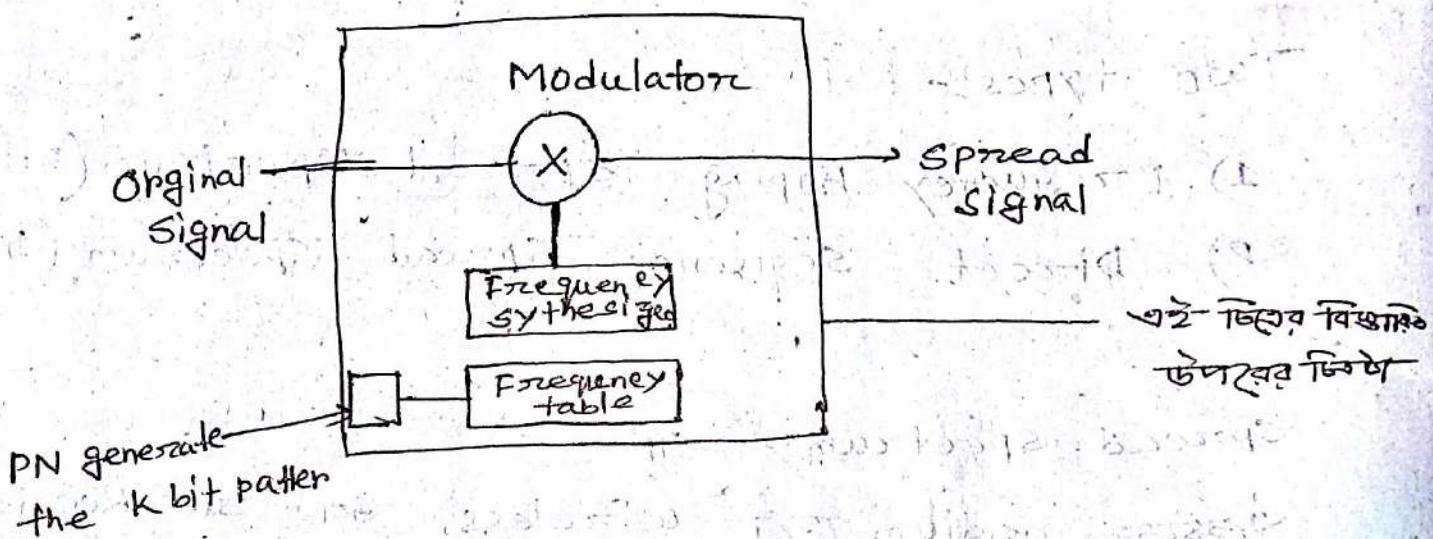
First Hop Selection

$$K=3$$

$$M=8$$



Frequency table

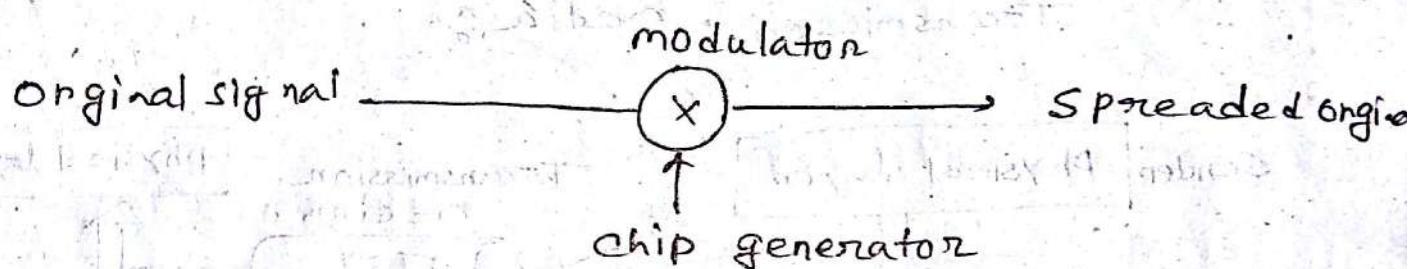


Pseudocode random code generator (PN), creates  
k bit pattern for every hopping period.

वाटन randomly bit घूमावने के लिए ट्रैक, फले जाएं  
corresponding carrier freq. पाठ्य यात्रा, जो carrier  
frequency Original freq. ने आरें mapping करेंगे  
करने करने, फले hacker करने का तरीका पाठ्य  
फले वाटन Safety generate कर सकते हैं।

## 2) DSSS :-

- DSSS techniques also expands the bandwidth of the original signal, but the process is different.
- In DSSS, we replace each data bit with  $n$  bits using a spreading code.
- In other words, each bit is assigned a code of  $n$  bits called chips, where the chip rate is  $n$  times that of the data bit.

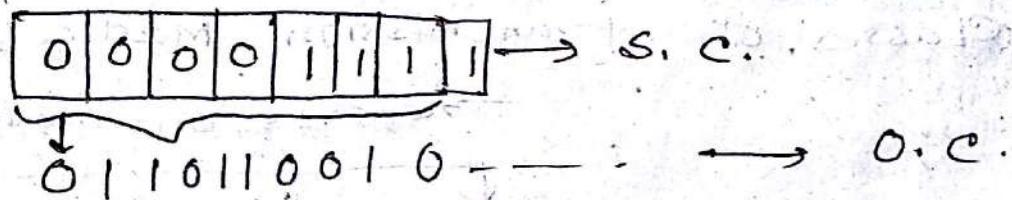


spreading code আবাবে, এগুলো code র মাঝে  
add রয়ে এবং co spreading code রয়ে

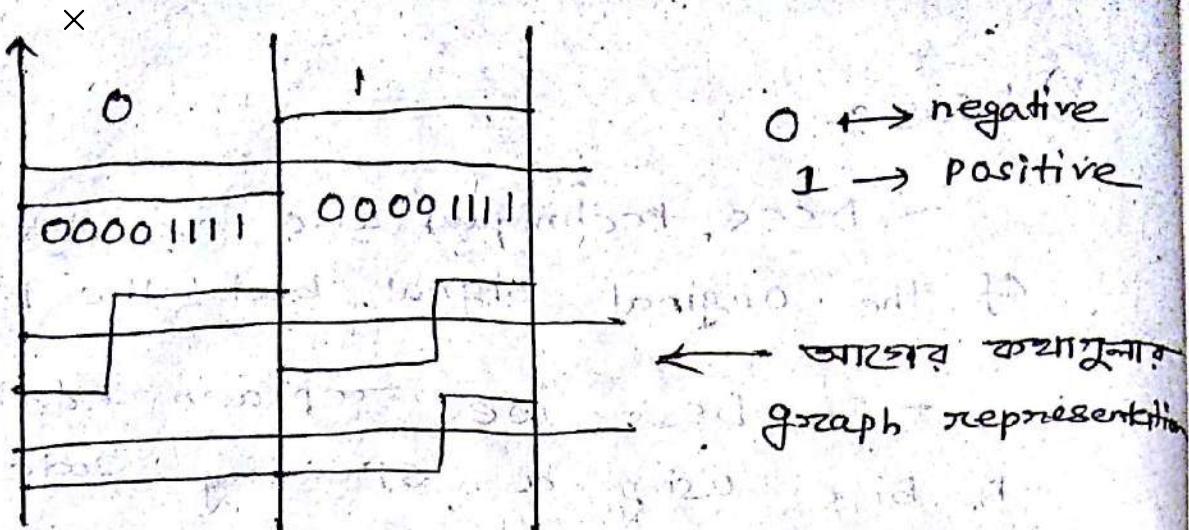
হবে কাবি, spreading code = 8 bit

original code = 32 bit

total =  $8 \times 32$  bit



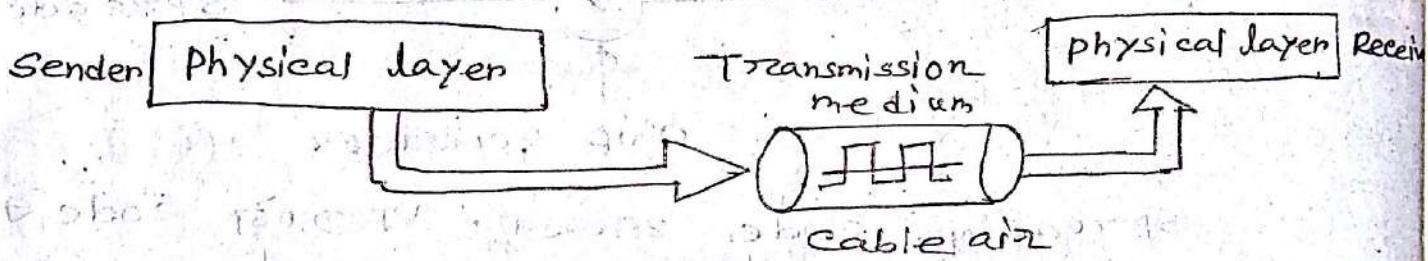
0000011110000111---- → spreaded  
original  
modulator



23-08-2017  
10th (B) day

## Chapter 7

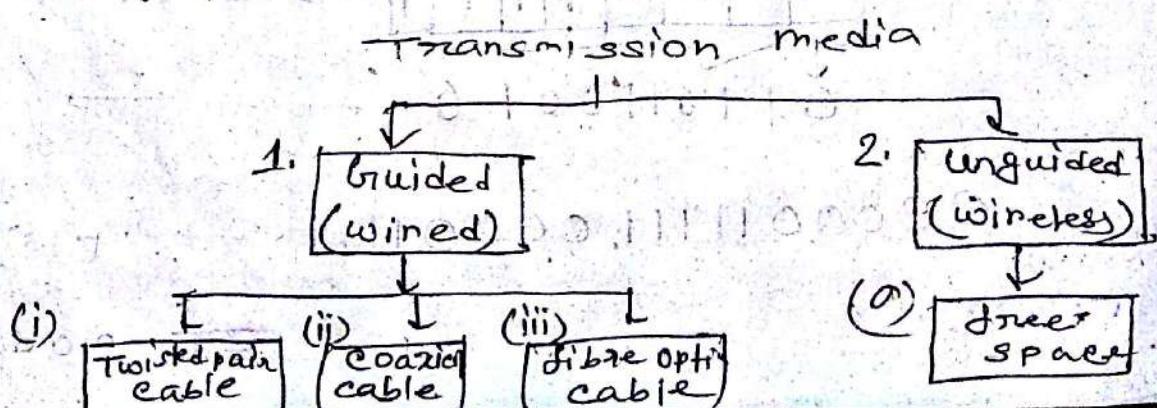
### Transmission media :-



- Transmission medium can be broadly defined as anything that carry information from source to destination.

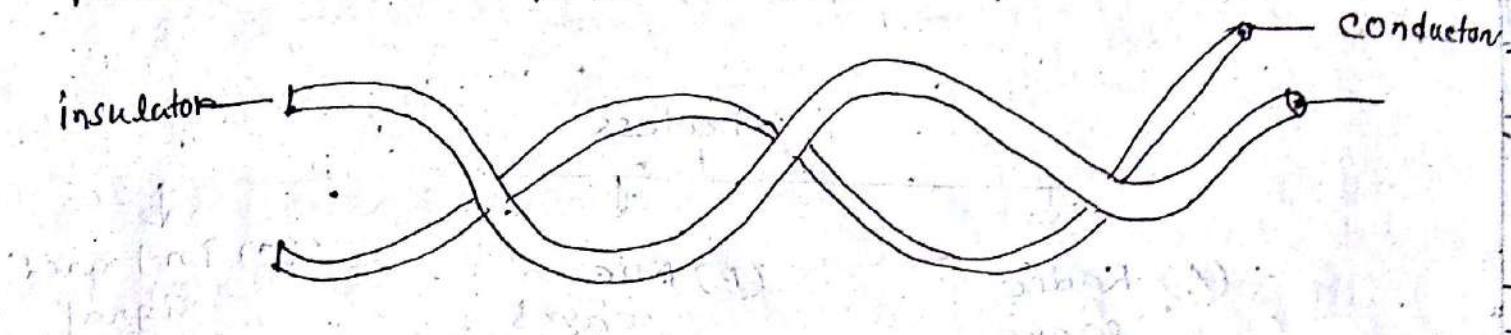
- Physical layer का काम क्या है ?

### Classes of transmission Media :-



### (i) Twisted pair cable:-

- A twisted pair cable consists of two conductors (normally copper), each with its own plastic insulation, twisted together.



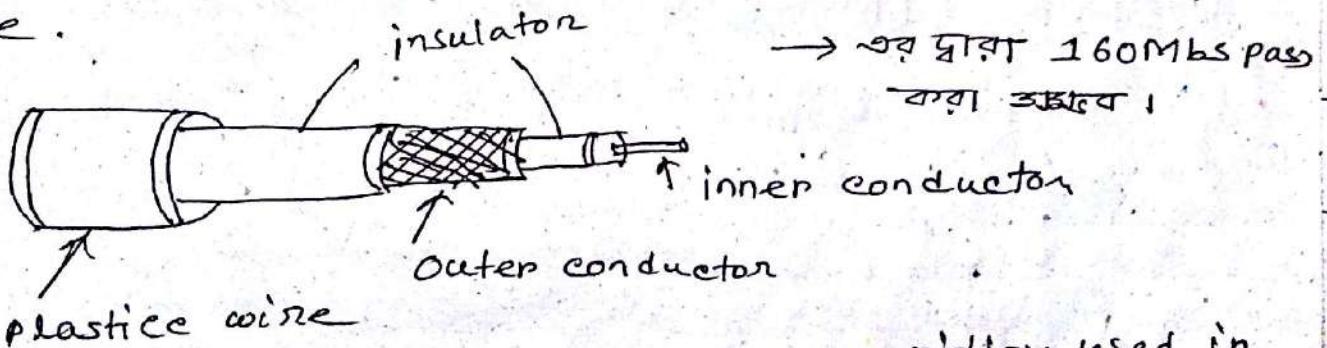
- One of the wires is used to carry signal to the receiver, and the other is used only as a ground reference.

### Application:-

Twisted pair cables are used in telephone lines to provide voice and data channels.

### (ii) Coaxial cable:- (Dish line)

- coaxial cable carries signals of higher frequency ranges than those in twisted pair cable.

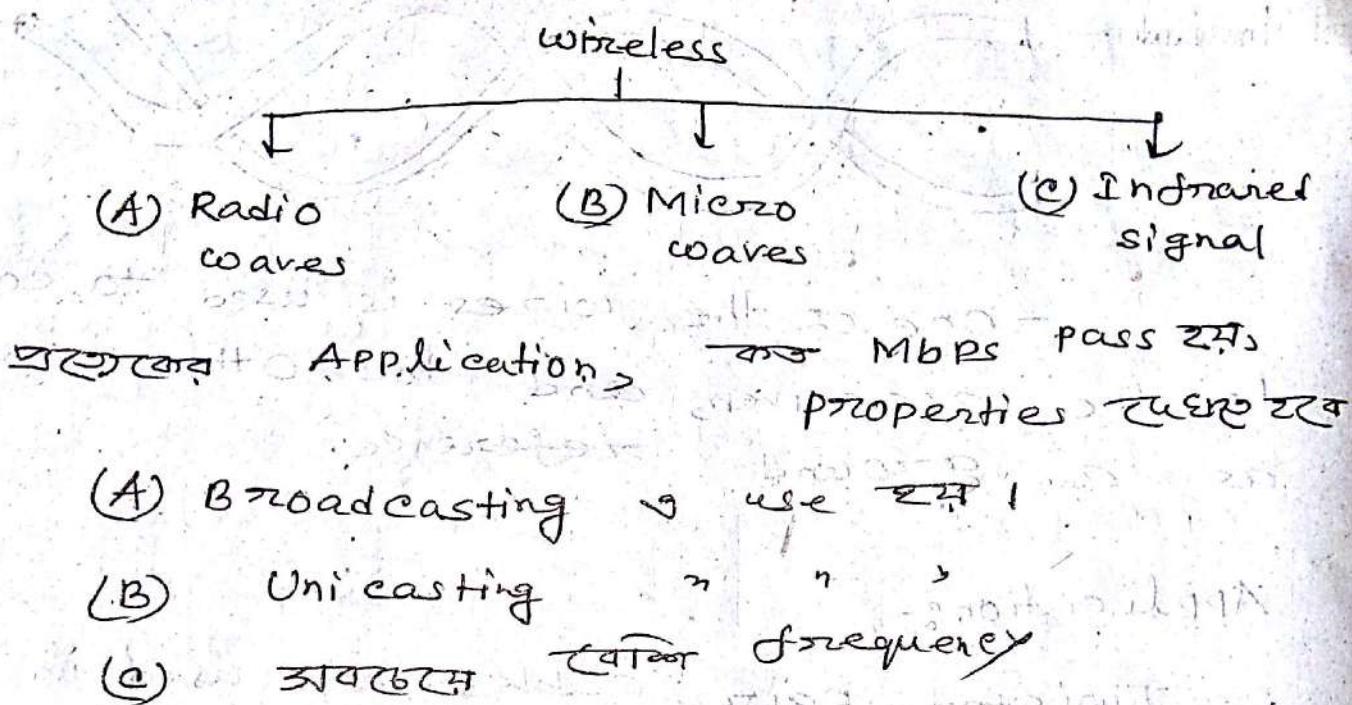


Application: coaxial cable was widely used in analog networks where a single coaxial network could carry 1000 voice signals.

- Cable TV network also used coaxial cables.

- Ethernet LAN

(iii) Fibre optic cable → ताँधे निष्ठा रखे



Radio wave	Infrared	Light wave
Micro wave	300 GHz	400 THz

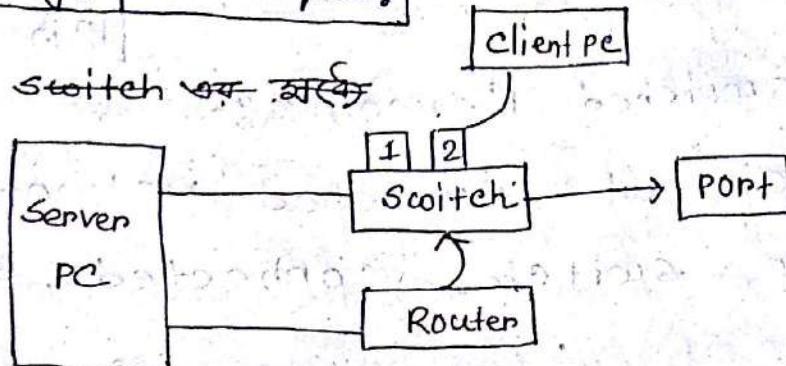
Wavelength  
radio wave

radio wave

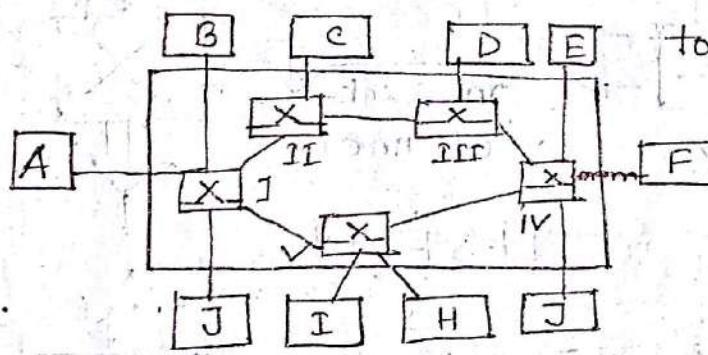
28-08-2017  
10th (E) day

## Switching :- techniques:-

एको switch एवं वर्तुल

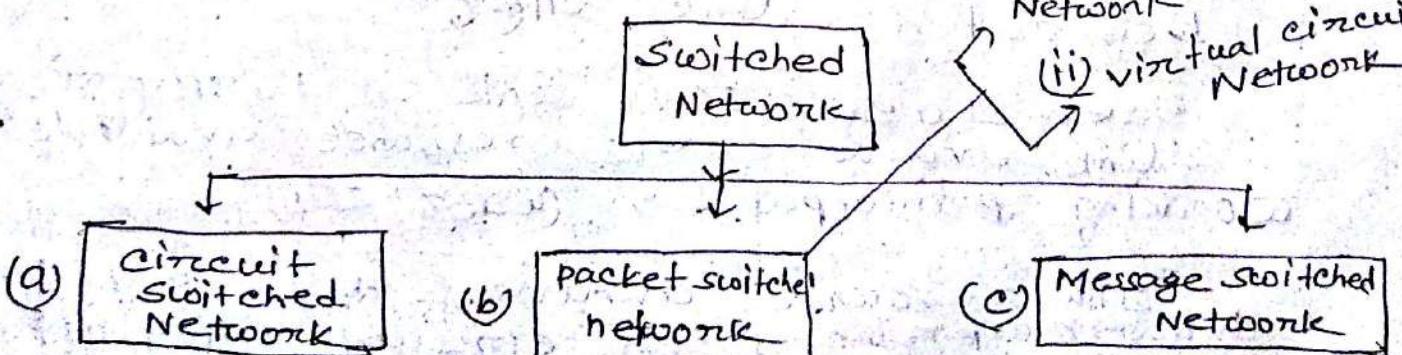


यस topology हरेक switching technique से better. topology को प्रत्येक आद्य dedicated line आए, एको device idle थाक्कले अर्थात् transmission उपरा properly, ओर अवधिगते दूर बहुत हम switch to point connection द्वारा Point dynamically decision लाने।



A यसको data पाठाले चाहिए, logical, & physical address एवं शारीरिक, switch जसे आवार्द्ध विशेष A यसको ले देयायः एक आरथ अलक्ष्यनाम switch एकावर्त्ती access कर्ता करते। switching काट करते। switching cost करते।

(i) Datagram Network  
(ii) virtual circuit Network



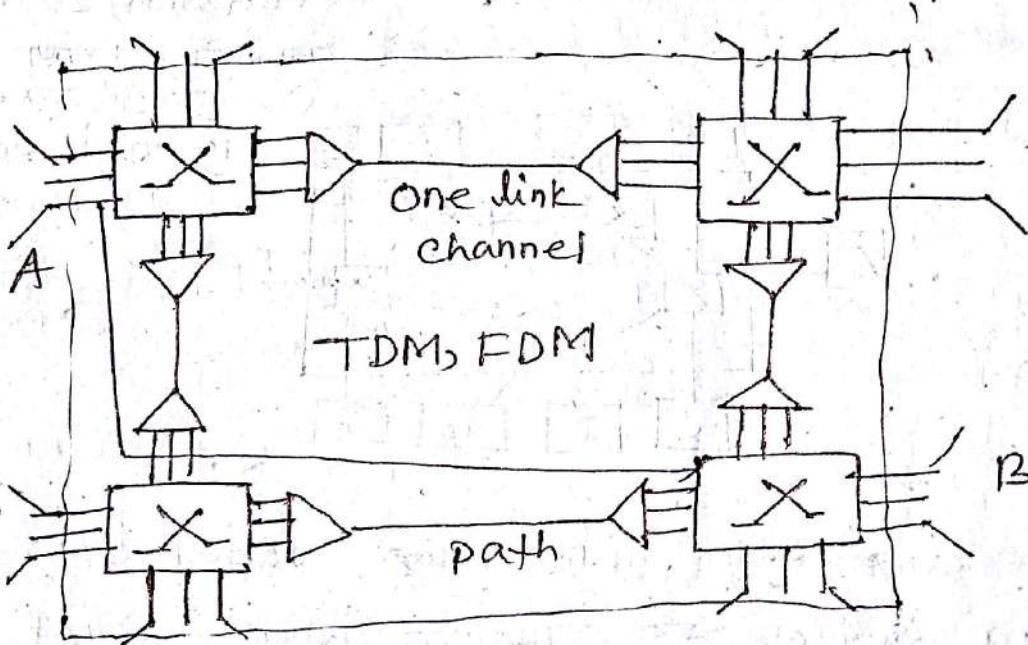
09-09-17  
got (11th(B)day)

Chapter 1, 3, 4, 5, 7 → Quiz (25 marks)

12th (E)day

## Circuit Switched Networks :-

- A circuit switched network consists of a set of switches connected by physical link.
- It is made of a set of switches connected by physical links in which each link is divided into  $n$  channels.



TDM, FDM वर्तमान समय में Established रखते हैं।  
cost बहुत कम है और Connection बहुत।

Dedicated link बनते हैं।

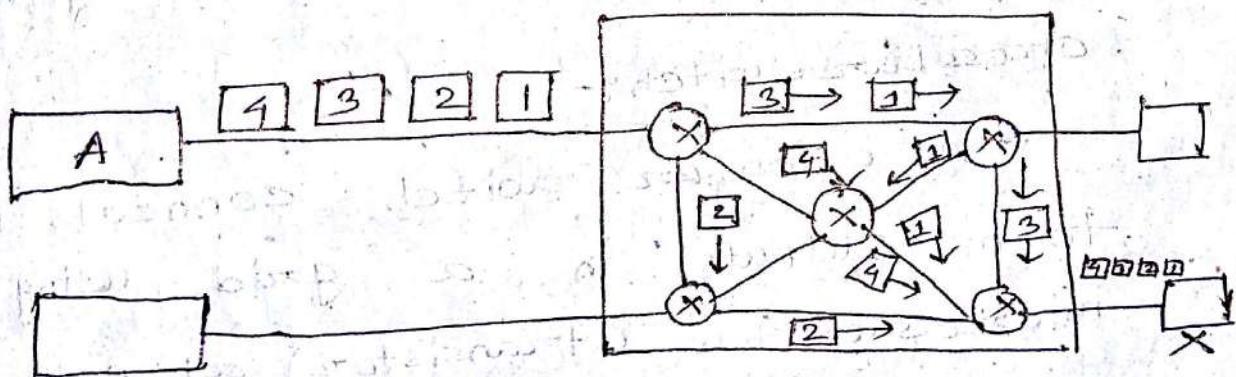
User 10 तक link की 10 तक संख्या लेते हैं।  
link available तक resource available होते हैं।  
working principle → यहाँ देखें।

dynamic resource available at यहाँ dynamic  
dynamic process का circuit switch network

fail करता है। Datagram, network के द्वारा disadvantages  
करता है।

## Datagram Network:-

- In packet switched networks, there is no procedure for resource reservation. resources are allocated on demand.
- In a datagram network, each packet is treated independently of all others. Even if a packet is part of a multipacket transmission, the network treats it as though it existed alone. packets in this approach are referred to as datagram.



- एप्पे पैकेट के बारे मानोर्ही इन्हें बातया जा सकता है, sequentially receive करके order में लगाया जा सकता है।
- यहाँ से link के access रेट, उपर लगे use का भी आवश्यक है।
- उदाहरण फिल्म के dedicated link में।

Virtual circuit Network: dedicated link  
dynamic एवं, किसी अतिरिक्त combination

11-09-2017  
11th (D) day

location निर्दिष्ट.

### Virtual circuit Network:-

points :- Circuit Switch Network & Data gram network  
को combined करके पाइजा याज़।

Characteristic गुण वाले घटक।

Structure of a switch:-

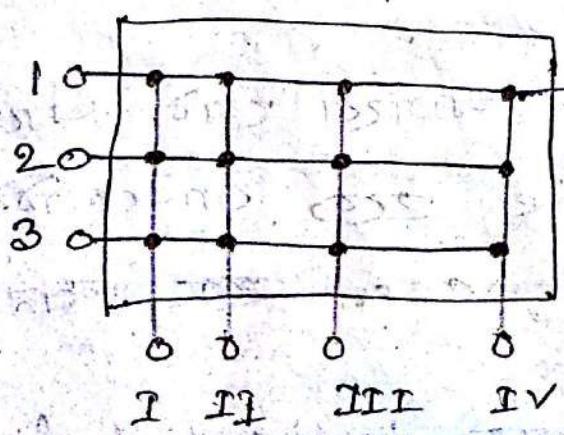
(i) Space division switch

(ii) Time

(iii) वर्तमान यथावत् cross bar switch

Crossbar switch:-

- A crossbar switch connects  $n$  inputs to  $m$  outputs in a grid, using electronic microswitches (transistor) at each crosspoint.



crosspoint

Crosspoint नंबर =  $n \times m$

1000 फैल उपर्युक्त आर्क्युलेटर connected

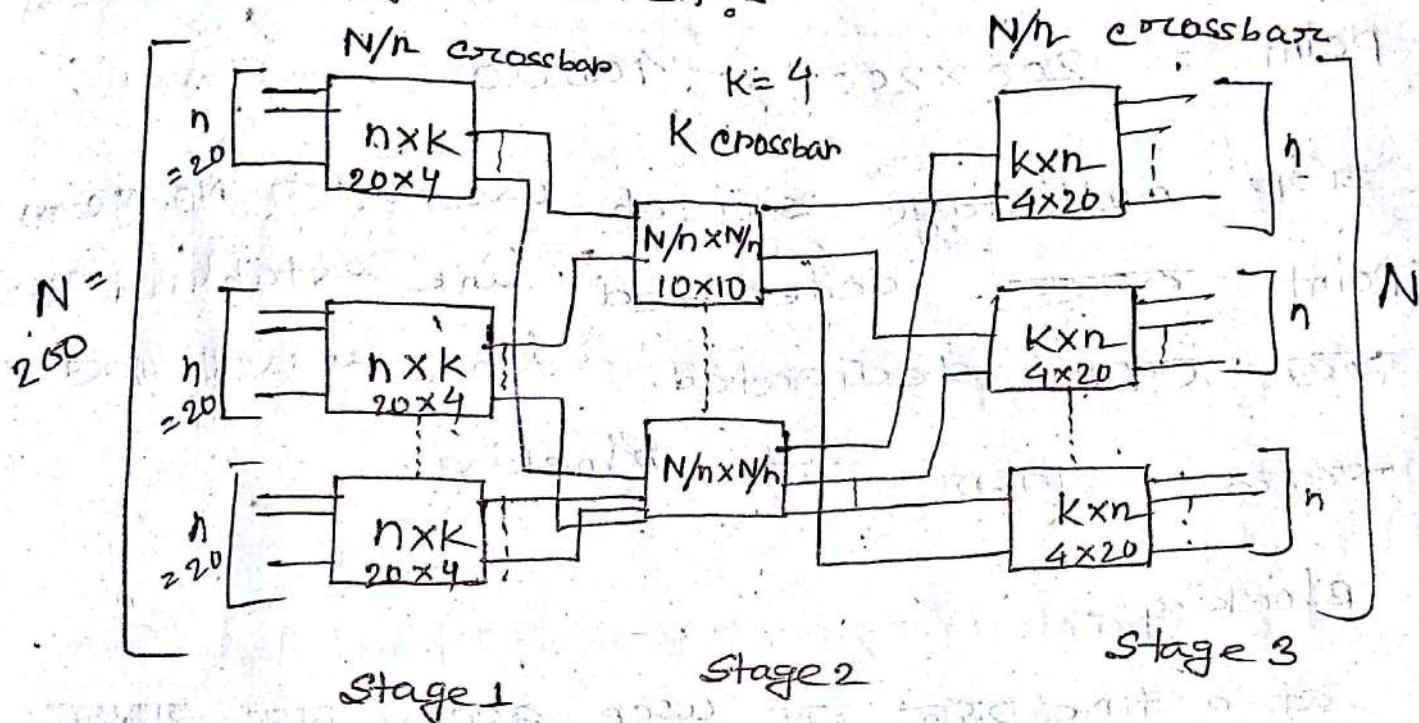
2000 फैल तो crosspoint =  $1000 \times 1000$

Disadvantage:-

No of cross point येते ज्ञान dedicated link established  
इयं प्रतिक्रिया अवधि 75% dedicated link बनावा  
आउत ना at a time.

जैसात् Overcome करावा तरी,

Multistage switch:-



$$200/20 = 10 \text{ } \overline{\text{cross bar}}$$

1. We divide the N input lines into groups, each of  $n$  lines for each group, we use one cross of size  $n \times k$ ,  $k$  is the number of cross in middle stage.

2. We use  $k$  cross bars, each of size  $(N/n) \times (N/n)$  in the middle.

3. We use  $N/n$  cross bars, each of size  $K \times n$  at the third stage.

$$\frac{N}{n} (n \times k) + k (\frac{N}{n} \times \frac{N}{n}) + \frac{N}{n} (k \times n) = 2KN \times k$$

$$= 10(20 \times 4) + 4 (10 \times 10) + 10 (4 \times 20)$$

$$= 2000$$

किन्तु cross bar switch use करने पर total cross

$$\text{Point} = 200 \times 200 = 40000$$

अब multistage switch use करने पर No. of cross

Point करते, dedicated link established रखके

फले लोन dedicated link unused रखते हैं।

इसका अवादा है Blocking.

Blocking:-

at a time यदि एक user access देता, उसके

अवादे access पाते तो अवादे at a time, उसके

wait करते हैं जैसे यह blocking.

इसका यह किन्तु cross point पाकाय रखना है

उसे blocking से overcome करना पास blocking

जैसे कि Clos algorithm जाने। (clos criteria

जाने।

- According clos criteria:

$$n = (\frac{N}{2})^{1/2}$$

$$K > 2n - 1$$

$$\text{Total Number of crosspoint} = \sum 4N (\frac{2N}{2})^{1/2} - 1$$

ଅଶ୍ଵା satisfied ଏହି blocking ଦ୍ୱାରା କଣେ ।

$$N = 200 \text{ ଲୋକୀ}$$

$$n = (200/2)^{1/2} = \sqrt{100} = 10$$

$$k = 2n - 1 = 19 \quad [\text{cross bar ପରିମାଣ}]$$

ଏହି ମିଡଲେ stage କୁ 19 ଟଙ୍କା cross ଯାଏ ଆବଶ୍ୟକ ।

then, ଏହିକାଟେ cross ଯାଏ 20 ଟଙ୍କା input, 20  
ଟଙ୍କା output, then total cross ଯାଏ middle

~~9500~~

$$19 (20 \times 20)$$

$$1^{\text{st}} \text{ stage } \rightarrow 20 \times (10 \times 19)$$

$$3^{\text{rd}} \text{ stage } \rightarrow 20 \times (19 \times 10)$$

$$\begin{aligned} \text{then, total cross bars, } & 19 (20 \times 20) + 20 \times (10 \times 19) \\ & + 20 \times (19 \times 10) \\ & = 9500 \text{ cross bar.} \end{aligned}$$

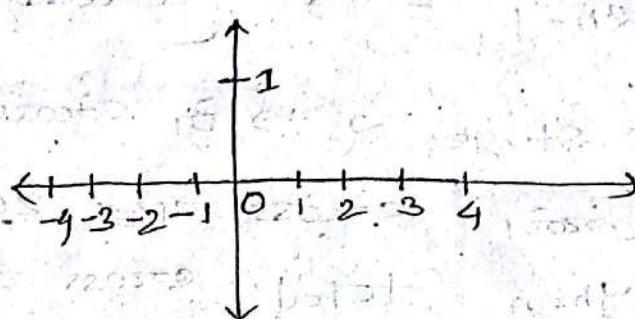
ଆମର ଏହି cross bar କୁ 2000, Now,  
 $= 9500$

ଫଳ୍ପ ବଲ୍କିଂ ଅନ୍ତର୍ଗତ ଦ୍ୱାରା ଏହି ଏକାକ୍ରମ ଏବଂ  
cross bar ଓ ଏକାକ୍ରମ 100% cross point  
ହେଲୁ ନା, ଆଜିଥା ଯାଏ 24%. cross point.

12-00  
11th(E) day

delta function: एक function जो अंतर्वाले क्षेत्र में output देता है

$$\delta[n] = \begin{cases} 1 & n=0 \\ 0 & \text{otherwise} \end{cases} \quad f(x) = \{1, 0, 1\}$$



↑ = origin

अब आवश्यक

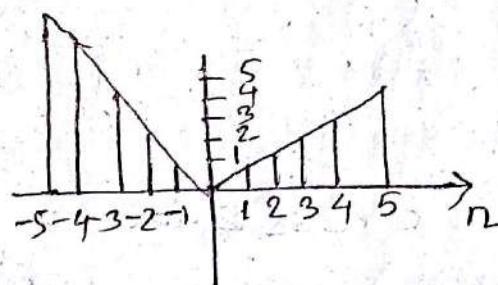
$x=0$  के लिए

$f(x)=1$

$$\delta[n] = \{ \dots, 0, 0, 0, 1, 0, 0, 0, 0, \dots \}$$

प्रृथम आवश्यक impulse response (जोड़ करते हुए)  
response i unit यादि इस)

Even signals (symmetric signals)



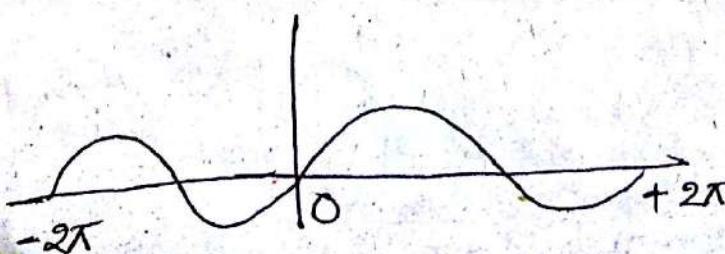
[Even digital signal]

$x(n) =$  n का correspond  
~ output

$$x(n) = x(-n)$$

n = positive लिए ए signal पाठ्य यादि, n = negative

लिए तो signal लिए तो even signal हो



[Even Analog signal]

$$x(n) = \{0, 1, 2, 3, 4, 5\}$$

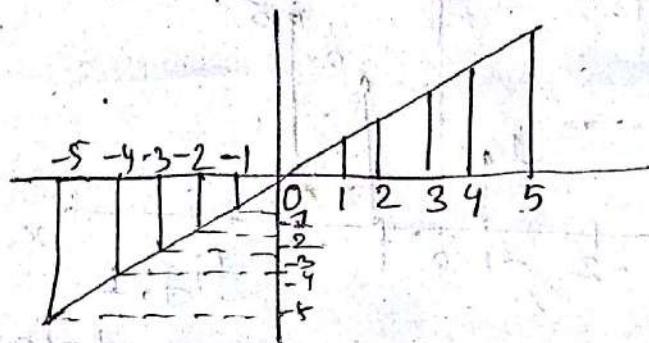
↑

$$x(-n) = \{0, 1, 2, 3, 4, 5\}$$

$$x(-n) = \{5, 4, 3, 2, 1, 0\}$$

even signal

Odd signal:- (Antisymmetric signal)



$$x(n) = -x(-n)$$

$$x(1) = +1$$

$$x(-1) = -1$$

मुख्य signal वर्त half portion graph एवं उक्ते हिंदे वाले half portion opposite हिंदे आवश्यक ताकि Odd

Co-correlation: एक छत्रसंघ आवश्यक समिक्षण करने वाला तथा विवेदन

Cross correlation: दो विवेदन वाले अलग अलग signal के क्रमानुसार co-related ताकि वह cross correlation,

$$\left. \begin{array}{l} x(n) \\ y(n) \end{array} \right\} \text{दो विवेदन signal}$$

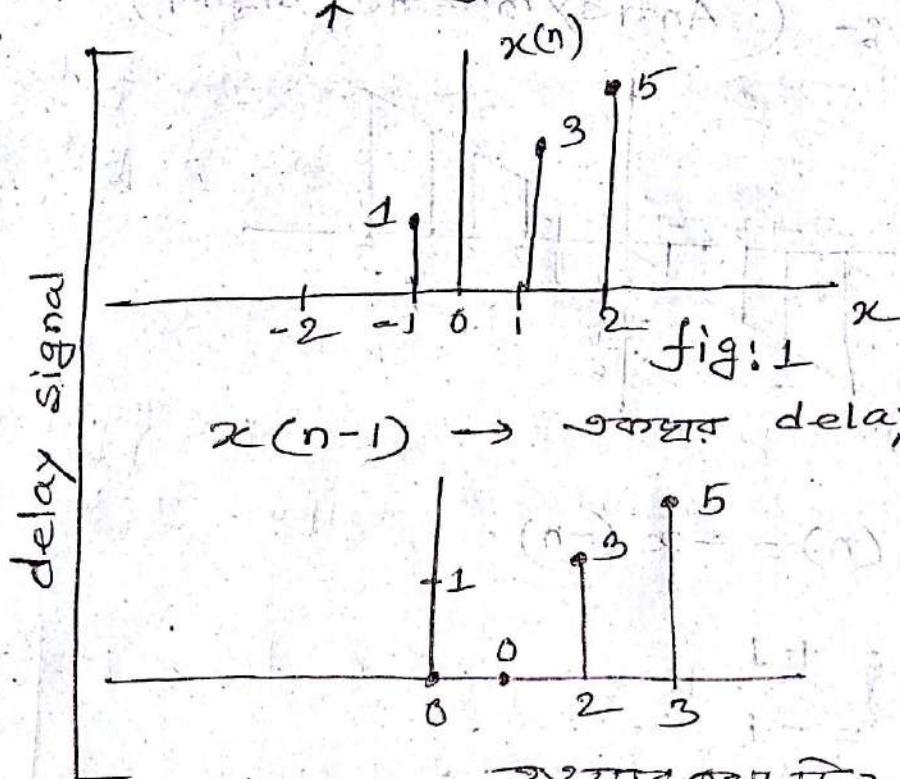
Auto correlation: ~~एक~~ एक signal के पाठ्यांक एवं या पाठ्यांक अपेक्षित जैसी समिक्षिति करने ताकि  $x(n) \rightarrow$  उक्त signal.

Shifting:-

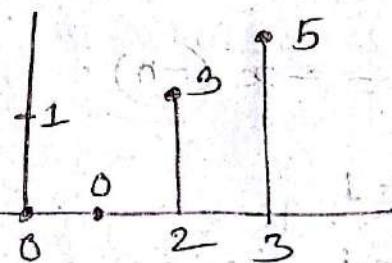
(i) delay signal

(ii) Advance

$$x(n) = \{1, 0, 3, 5\}$$



$x(n-1) \rightarrow$  একস্থান delay then,



মনে করো time  $\rightarrow$  স্থান  $\rightarrow$  সময় তাহা থেকে যদি বেশি  
অবস্থা স্থান অবস্থা তাহা থেকে delay signal বলে

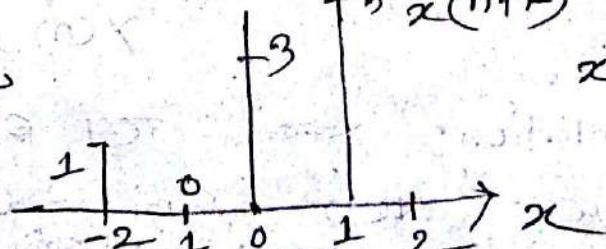
$$k = 1, 2, \dots \quad (n-k) \rightarrow k \text{ delay}$$

$$x(n-1) = \{0, 1, 0, 3, 5\}$$

(ii)  $x(n+1)$

এখন left shift হবে সত্ত্বেও  $x(n+k)$

fig:1 র মত হবে,



$$x(n+1) = \{0, 1, 0, 3, 5\}$$

23-09-17  
13th (B) day

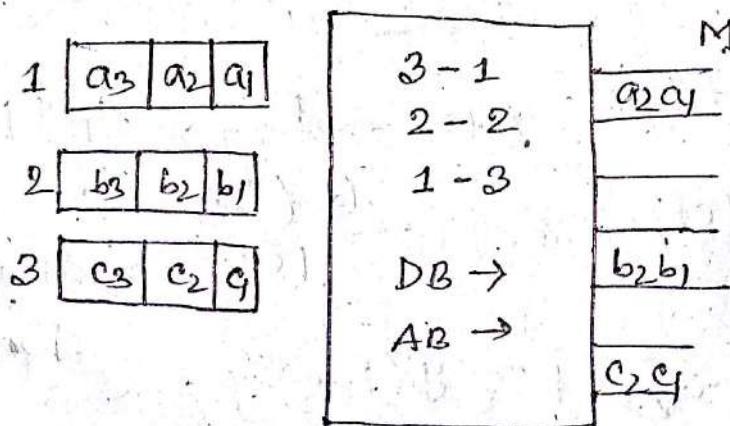
## ଓঁ আলোর lecture mobile এবং

Crossbar switch  $\rightarrow N \times M$

Multistage  $\rightarrow$  Crosspoint পারে মাঝের  
কারণ Hacking ক্ষেত্রে রে  
পারে নন blocking দ্বা  
রা কোনো clos criteria  
use কো ইয়ে।

TIME DIVISION Switching:-

\*\*\* Time slot interchange: (TSI)



Control unit

$\rightarrow$  যাই priority হচ্ছে  
ইটে অন্ত আগে send করত

Time slot interchange কি? কিভাবে করা করে?

Scaling property:-

$$f(x) = 2x \text{ or } f(t) = 2t$$

$$f(\frac{x}{2}) = \frac{x}{2} \text{ or } f(\frac{t}{2}) = \frac{t}{2}$$

$$t=0 \quad f(t)=0$$

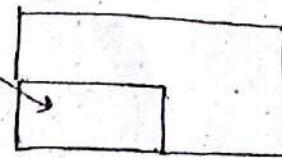
$$t=1 \quad f(t)=2$$

$$t=2 \quad f(t)=4$$

वर्तमान  $f(t) = 2t$  Scaling factor

Scaling factor = 2

Scaling factor = 0.5



Scaling factor एक नियंत्रित constant value, याकि बेस वाले signal को change करता है।

$$\begin{aligned} &= 0 = \\ \text{Scaling property} \quad & f(t) \cdot g(t) \xrightarrow{F} G(f) \\ F \cdot \{g(ct)\} &= \frac{G(f/c)}{|c|} \rightarrow \text{always positive} \end{aligned}$$

$$\text{Shifting property, } F\{g(t-a)\} = e^{-j2\pi fa} G(f)$$

Scaling factor  $> 0$  एक signal amplify.

Scaling  $< 0$   $\rightarrow$  signal reduce.

Signal और आवर्त यथा scaling factor युक्त एक

then उस scaling factor युक्त signal होता है।

$$\mathcal{F} \{ g(t) \} \longleftrightarrow \int_{-\infty}^{\infty} g(t) e^{-i2\pi ft} dt$$

$$\mathcal{F} \{ g(ct) \} \rightarrow \int_{-\infty}^{\infty} g(ct) e^{-j2\pi ft} dt$$

Substitute,

$$u = ct$$

$$\Rightarrow du = cdt$$

If  $c$  is positive

$$\mathcal{F} \{ g(ct) \} = \int_{-\infty}^{\infty} \frac{g(u)}{c} e^{-i2\pi f \frac{u}{c}} du$$

$$= \int_{-\infty}^{\infty} \frac{g(u)}{c} e^{-i2\pi f \frac{u}{c}} du$$

$$= \frac{G_i(f/c)}{|c|}$$

If  $c$  is negative,

$$\mathcal{F} \{ g(ct) \} = \int_{+\infty}^{-\infty} \frac{g(u)}{c} e^{-i2\pi f \frac{u}{c}} du$$

$$= - \int_{-\infty}^{\infty} \frac{g(u)}{c} e^{-i2\pi f \frac{u}{c}} du$$

$$= \frac{G_i(f/c)}{|c|}$$