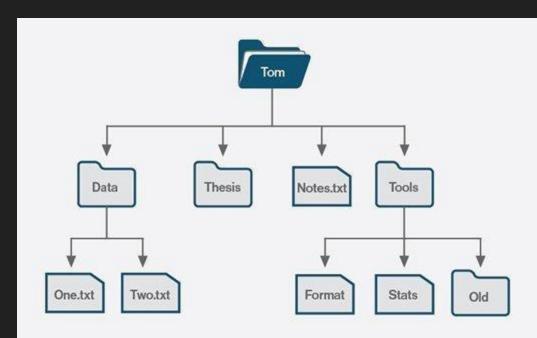
CSCI 452 Presentation

Raymond Healy, Koen Komeya, Ian Campbell

What we set out to accomplish

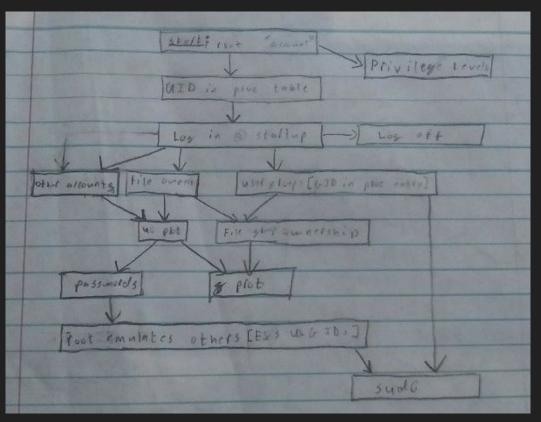
- Permission system Raymond
- File System Ian & Raymond
- Disk Driver Koen



Permission System - Overview

- Generally designed to ape unix divided its permissions
 - Permissions divided into 3 rings, the owner, the owning group, and others
 - Permissions at each level are also divided
 - Read/Write (files, FS traversal)
 - Meta (change file ownership, location, and permissions)
 - Syscall execution (e.g. only root, sudoers, and owners kill a proc)
- When trying to create a plan I tried to think through logical extensions of existing elements.
 - E.g. "Protections based on UID" to "Changing UID should be restricted to root" to "the login shell must be spawned as root"

Permission System



Permission System - Early progress

```
typedef struct pcb s {
  // Start with these eight bytes, for easy access in assembly
   context t *context;  // pointer to context save area on stack
   stack t *stack; // pointer to process stack
   inode id t wDir: // ID of the working directory's inode
   int32 t exit status; // termination status, for parent's use
   event t event;  // what this process is waiting for
   pid t pid;  // unique PID for this process
   pid t ppid: // PID of the parent
   gid t gid;  // group ID of this process
   uid t uid: // user ID of this process
   state t state; // current state (see common.h)
   prio t priority;  // process priority (MLQ queue level)
   uint8 t quantum; // quantum for this process
   uint8 t ticks:  // ticks remaining in current slice
   fd t files[MAX OPEN FILES]; // File descriptors
} pcb t;
```

```
(U33) is this land made for you and me?
(000)
(031) Nobody living can ever stop me,
(037) As I go walking that freedom highway;
(040) Nobody living can ever make me turn back
(035) This land was made for you and me.
$ test 1 1
M User 1.0 reports uid 100
M User 1.0 reports gid 0
$ list 0
        Baseline tests (bank 0):
                A, B, C: Print out name a finite nr of times w/ delays
                   D, E: Like A-C, but also checks write's return
                   F, G: Report, sleep (10 and 20 secs) and exit
                      H: Test orphan reparenting
                      I: Test killing children
                      J: Test overflow proc table
                   K, L: Iterate spawning user X and sleeping
                   M, N: Test various syscalls through children
                      P: Iterates reporting time and sleeping
                      Q: Tests a bad syscall
                      R: Echo SIO forever (NO EXIT!!!)
                      S: Sleep forever (NO EXIT!!!)
                   T, U: Spawn user W copies and then wait or kill them
                      V: Play with process priority
$ test 0 F
FFF
```

Permission System - Later progress

- Following this I largely shifted focus to FS until it was nearing completion
- At that point, I returned to my earlier work
 - Improved what was already there (group selection, logon with names and passwords)
 - Extended the reach of permissions forward (e.g. implementing file permissions)
- In testing these modifications (and with a fuller drive), I also ran into and helped to fix many of the bugs remaining in FS, as a bonus.

```
Enter your username: Ray
Enter your password:
Test shell started
Try help for a list of commands
$ ls
    dwrwrwr
    dwrwrwr
                          . shadow
    ---wrwr
                         testLongName
    -wrwrwr
    --rwrwr
               00
                   00
                          .groups
$ test 2 3
Test FS 3.0: Opened "/test.txt" as channel 2
Editing the new file by writing to it
Test FS 3.0: Closed "/test.txt"
(025) HELLO WORLD IN A NEW FILE
$ ls
    dwrwrwr
    dwrwrwr
                          . shadow
    ---wrwr
                          testLongName
    -wrwrwr
    --rwrwr
                          .groups
               64
                   00
                          test.txt
    -wrwrwr
```

File System - Overview

- Very basic system to read and write to a device
- Reads and writes 512 byte blocks
- Currently functions on a RAMdisk and a disk driver emulated via QEMU.
- Reading, writing, opening, closing, creation, moving, and removal of files possible!
 - Done with Syscalls
 - Repurpose given read/write
- Each PCB contains array of File Descriptors
- Can do basic CAT and LS style operations.

```
* This structure will contain a variety of data
                                                                               #define NUM DIRECT POINTERS 14
 * relating to a file/directory in the FS.
                                                                                typedef uint32 t block t;
typedef struct inode s {
   // Meta Data (16 bytes)
                                                                                typedef struct {
   inode id t id;
                                                                                   uint32 t devID : 8;
   uint32 t nBlocks;
                                                                                    uint32 t idx : 24;
   uint32 t nBytes; // Either number of file bytes or number of subdirectories
                                                                                } inode id t;
   uint32 t nRefs; // Number of inodes referencing this one
                                                                                typedef struct {
   // Permission information (8 bytes)
                                                                                        char name[12];
   uint32 t permissions: 24;
                                                                                        inode id t inode;
   uint32 t nodeType: 8;
                                                                                } dirEnt t;
   uid t uid;
   gid t gid;
                                                                                // Each is 16 bytes
                                                                                typedef union {
   // Lock + Padding (4 bytes)
                                                                                    block t blocks[4];
   uint8 t lock; // 1 byte
                                                                                    dirEnt t dir;
   uint8 t pad[3]; // 3 bytes
                                                                                } data u;
   // Indirect Pointers 4 bytes
   block t extBlock; // Points to a block
   data u direct pointers[NUM DIRECT POINTERS]; // 32 + 14 * 16 = 256 bytes per inode
} inode t:
```

File System - Philosophies

- Keep it simple
 - Designed inode early
 - Identified block sizes and inode sizes 2 inodes per block
 - Don't get fancy currently no extension blocks can only have 14 direct blocks
- Test early & plug and play
 - RAMdisk
 - Allows us to test everything other than reboot persistence
 - Blessing in disguise
 - Make a device interface standard for all disk drivers
 - Register
 - Fill a structure with two function pointers (read & write)

File System - Issues we had

- Unintended recursion stack smashing on invaild read
- Write is far more complicated than read
- Allocating and freeing blocks has many points of failure
- To get a basic FS working requires a lot of helper functions to even be possible
 - o Alloc/Free Blocks, Alloc/Free inodes, getSubDir, rmDirEnt, addDirEnt, etc...

File Systems - If we had more time

- Extension Blocks
- FS Caching scheme

Disk Driver - The Final Product

- PATA Driver
- Based on PATA revision 6,0
- Compatible with Multiple Disks
- Works on 512 byte sectors
- In theory, supports both hard disk drives and solid state drives,

```
for (int ch = 0; ch < 2; ch++) {
    for (int dr = 0; dr < 2; dr++) {
        // Select drive
        ide write(ch, ATA REG HDDEVSEL, 0xA0 | (dr << 4));
        sleep1ms();
        // Identify command
        ide write(ch, ATA REG COMMAND, ATA CMD IDENTIFY);
        sleep1ms();
        // Ensure this is ATA
        if ( ide read(ch, ATA REG STATUS) == 0)
            continue:
        int fail = 0:
        unsigned char status;
        do {
            status = ide read(ch, ATA REG STATUS);
            if ((status & ATA SR ERR)) {fail = 1; break;}
        } while ((status & ATA SR BSY) || !(status & ATA SR DRQ))
```

Disk Driver - Overview

 First, talk to the PCI (Peripheral Component Interconnect) bus to get the PATA devices.

 PATA stands for Parallel AT Attachment, where AT referred to an IBM PC model, short for "Advanced Technology"

 Based on Western Digital's IDE, which stands for Integrated Drive Electronics. Because of their development alongside each other, they are used interchangeably to refer to the same thing.

 PATA is basically one of the protocols for communicating with the device that interfaces with the actual hard disk.



Disk Driver - Issues I faced

 Was hoping to implement AHCI but realized I probably wouldn't be able to finish it, so I switched course to PATA.

 Didn't work while testing due to misconfigured QEmu. Solution - include the hardware components in the command line, but only what you need.

Mysterious interrupts firing on the Hard Disk ISR even when interrupts were

disabled.



Image Sources

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