CS 61BL Lab 20

Ryan Purpura

Announcements

- BearMaps due tomorrow at 11:59 PM! Remember the extra-credit portion (Auto-complete with Tries) is worth 5 extra credit points.
- The final is next Thursday! Remember there is no dead week over summer. Start studying now!

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 - $\log(n!) = \log 1 + \log 2 + \dots + \log n$
 - $n \log n = \log n + \log n + \dots + \log n$
 - So $\log(n!) \le n \log n \implies \log(n!) \in O(n \log n)$

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 - $\log(n!) = \log 1 + \log 2 + ... + \log n$

$$\log(n!) \ge \log\left(\frac{n}{2}\right) + \log\left(\frac{n}{2} + 1\right) + \dots + \log(n)$$

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 - $\log(n!) \ge \log\left(\frac{n}{2}\right) + \log\left(\frac{n}{2}\right) + \dots + \log\left(\frac{n}{2}\right)$
 - $\log(n!) \ge \frac{n}{2} \log\left(\frac{n}{2}\right)$
 - $\frac{n}{2}\log\left(\frac{n}{2}\right)\in\Theta(n\log n)$ so we know that $\log(n!)\in\Omega(n\log n)$

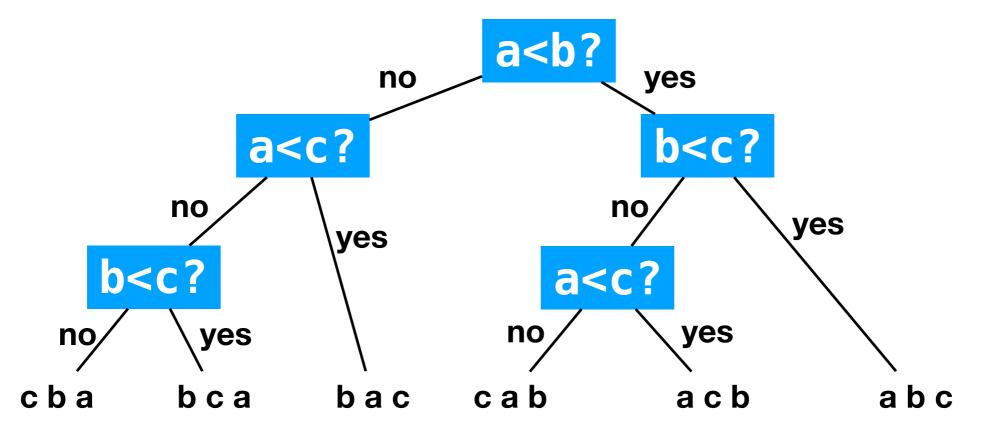
- Is $\log(n!) \in \Theta(n \log n)$?
- We know $\log(n!) \in \Omega(n \log n)$ and $\log(n!) \in O(n \log n)$
- So yes, $\log(n!) \in \Theta(n \log n)$

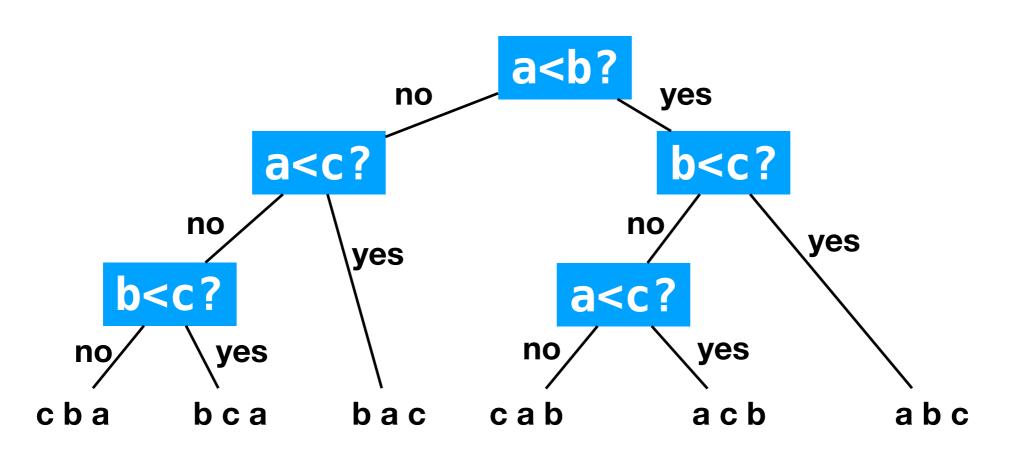
ullet Given an array of N distinct items, how many ways can it be ordered?

• Answer: *N*!

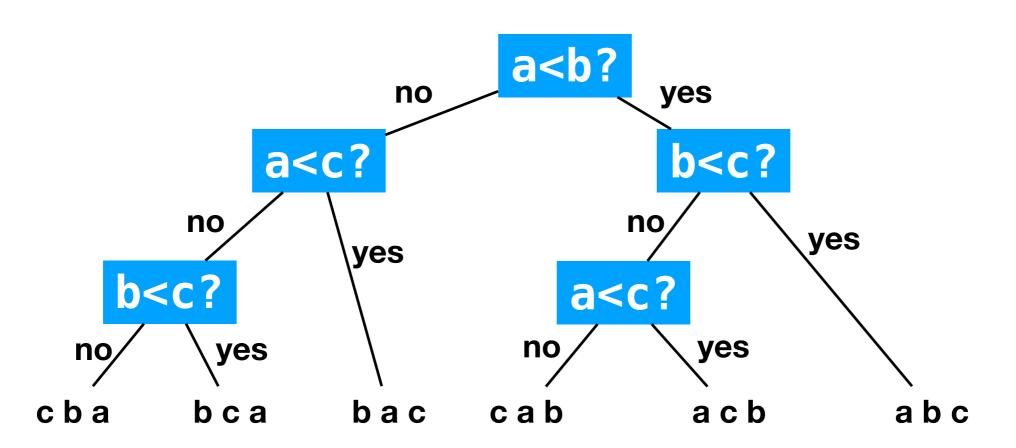
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- When we sort a list, we perform comparisons (yes-or-no questions involving two items) in order to choose which permutation is the "sorted" one consistent with the comparisons.

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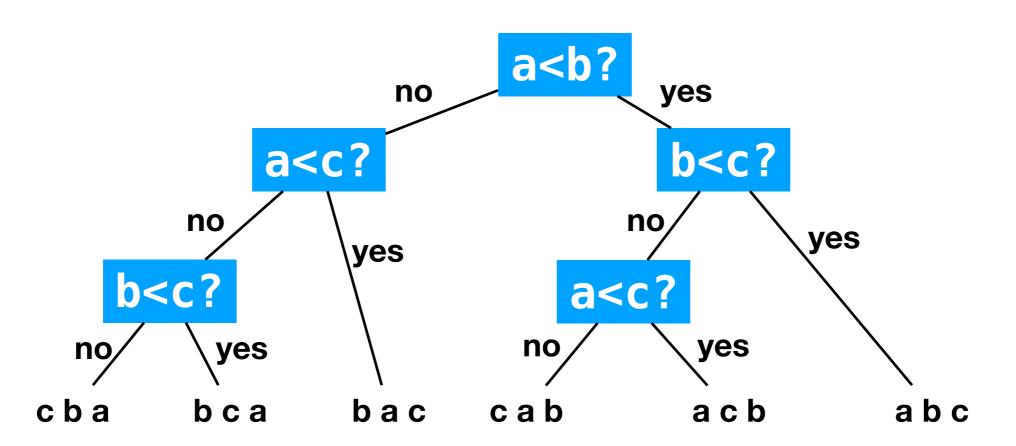




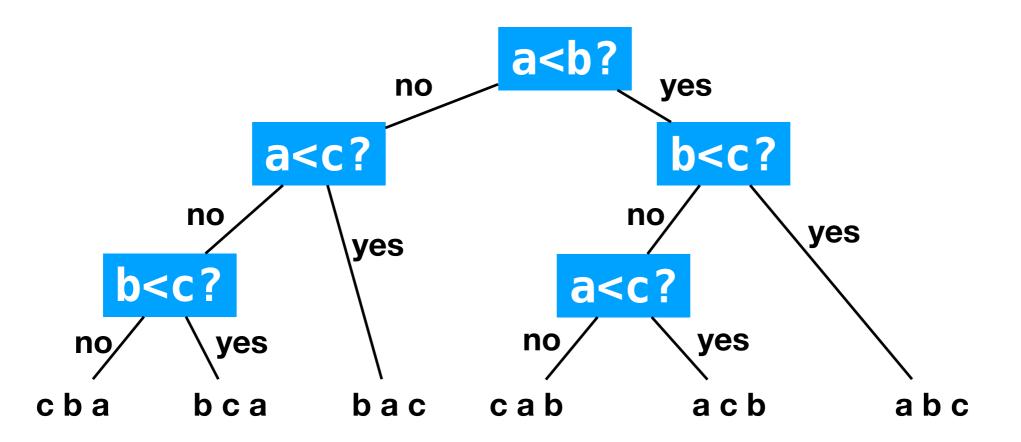
• If we have n! leaves of our tree, how tall must our tree be?



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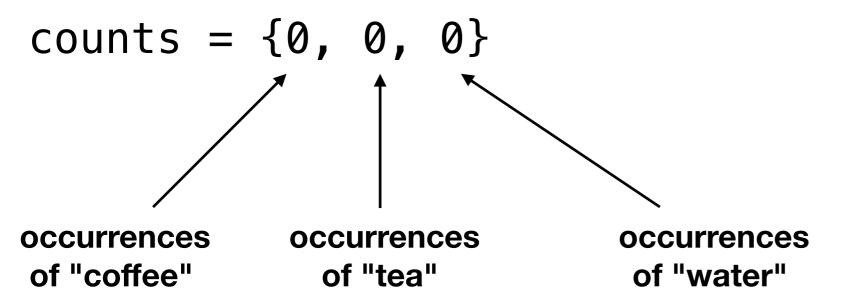


Sorting with Limited Variety of Items

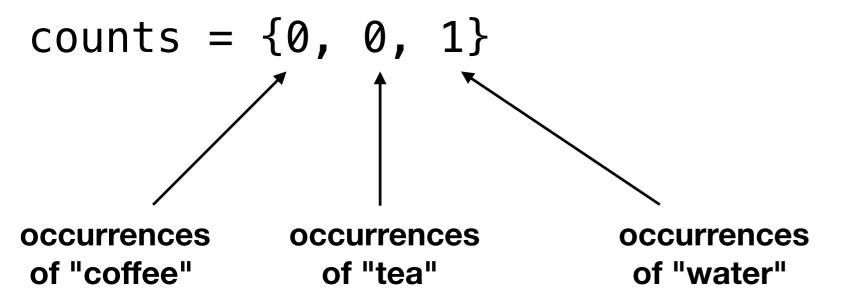
- What if the the list we are sorting only has a limited number of varieties?
- For example, we have a list of Strings where the only values are "water", "coffee", and "tea" and want to sort in alphabetical order.
- Can we sort this in better than $n \log n$ time?

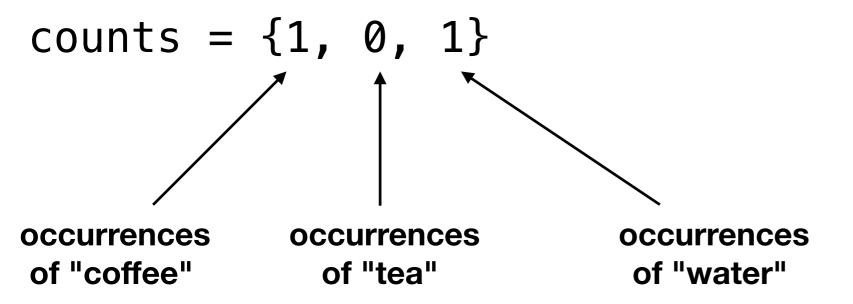
Counting Sort

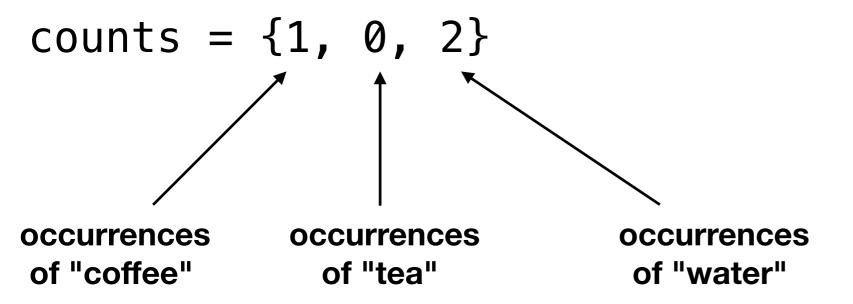
- Idea: count occurrences of each. Then we know where each group starts!

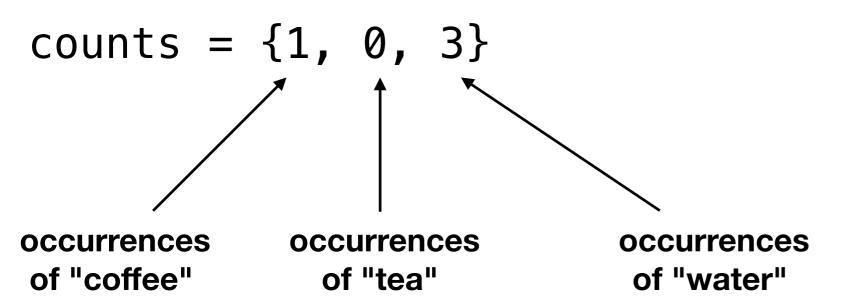


myDrinks = {"water", "coffee", "water", "vater", "coffee", "tea"}

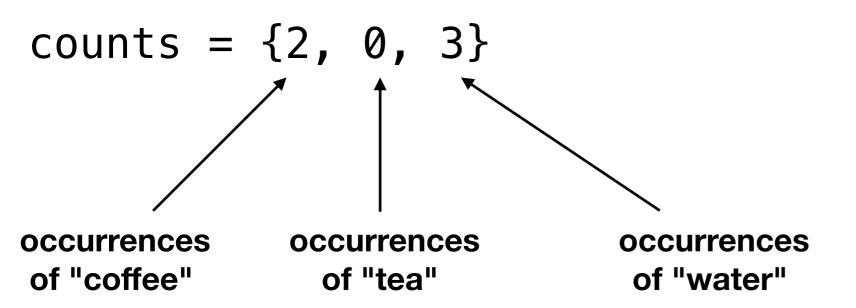




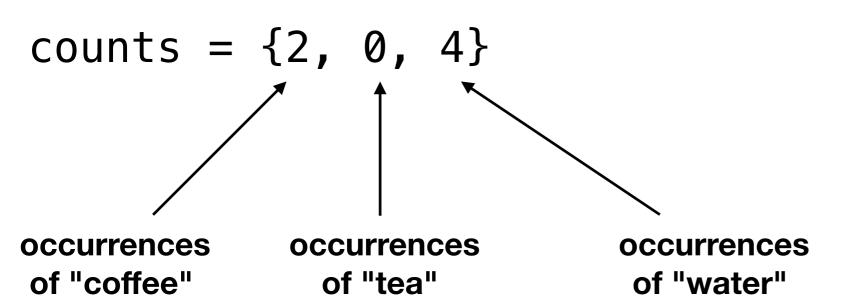




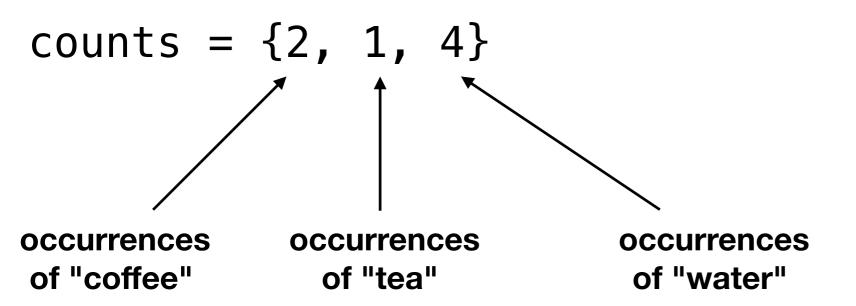
myDrinks = {"water", "coffee", "water", "coffee", "tea"}
"coffee", "water", "tea", "coffee", "tea"}



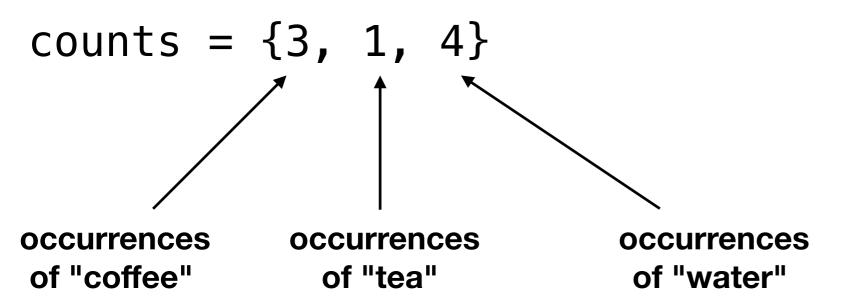
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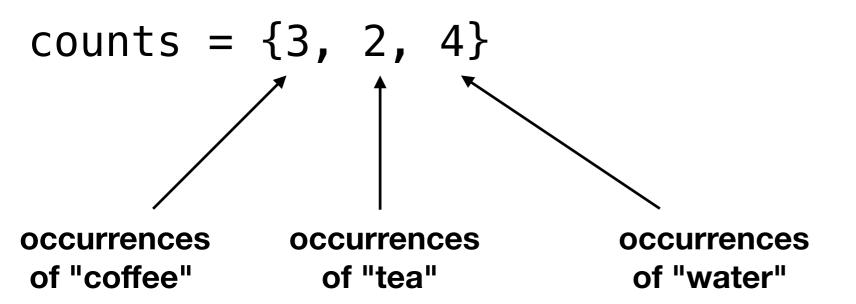
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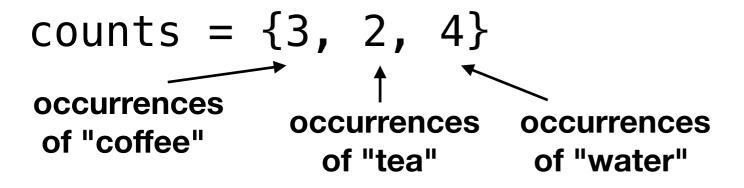


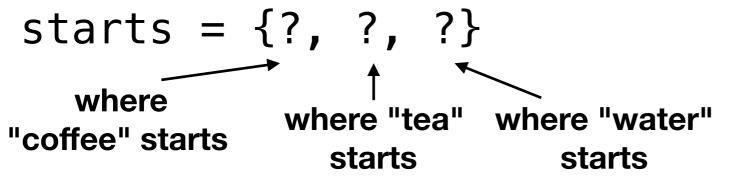
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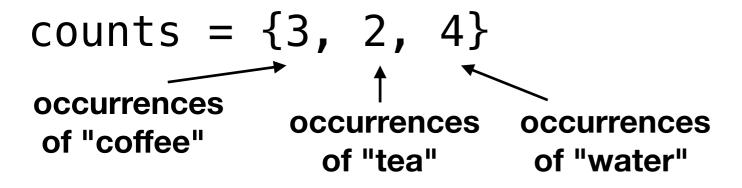


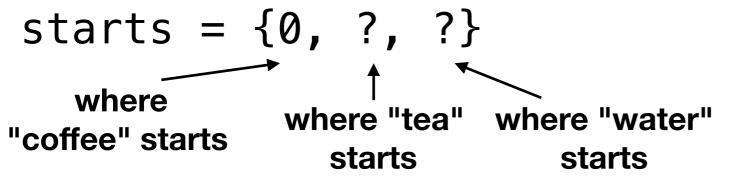
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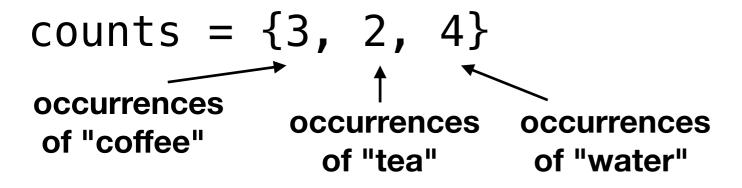


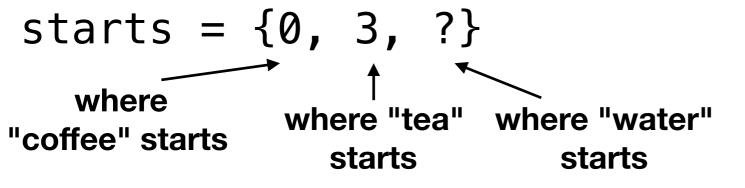


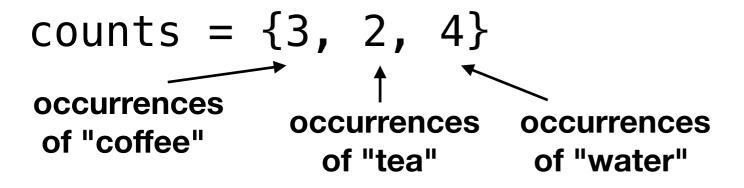


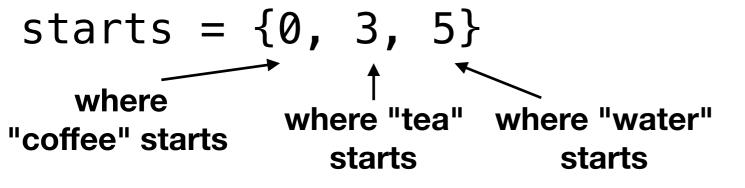




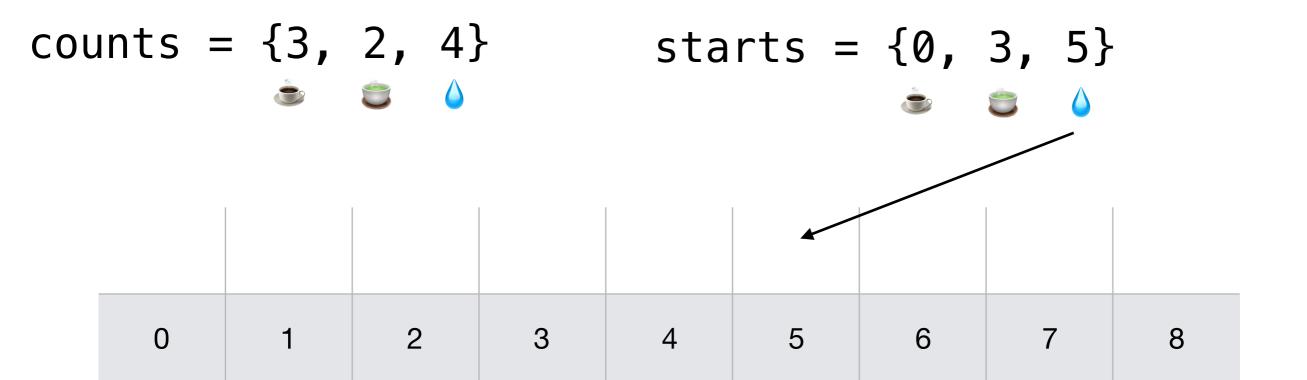








myDrinks = {"water", "coffee", "water", "vater", "coffee", "water", "tea", "coffee", "tea"}



counts =
$$\{3, 2, 4\}$$

starts =
$$\{0, 3, 5\}$$

					water			
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$

starts =
$$\{0, 3, 6\}$$



					water			
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{0, 3, 6\}$

coffee					water			
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{1, 3, 6\}$

coffee					water			
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{1, 3, 6\}$

coffee					water	water		
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{1, 3, 7\}$

coffee					water	water		
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{1, 3, 7\}$

coffee					water	water	water	
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{1, 3, 8\}$

coffee					water	water	water	
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{1, 3, 8\}$

coffee	coffee				water	water	water	
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{2, 3, 8\}$

coffee	coffee				water	water	water	
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{2, 3, 8\}$

coffee	coffee				water	water	water	water	
0	1	2	3	4	5	6	7	8	

counts =
$$\{3, 2, 4\}$$
 starts = $\{2, 3, 9\}$

coffee	coffee				water	water	water	water
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{2, 3, 9\}$

coffee	coffee		tea		water	water	water	water
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{2, 4, 9\}$

coffee	coffee		tea		water	water	water	water
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{2, 4, 9\}$

coffee	coffee	coffee	tea		water	water	water	water
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{3, 4, 9\}$

coffee	coffee	coffee	tea		water	water	water	water
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{3, 4, 9\}$

coffee	coffee	coffee	tea	tea	water	water	water	water
0	1	2	3	4	5	6	7	8

counts =
$$\{3, 2, 4\}$$
 starts = $\{3, 5, 9\}$

coffee	coffee	coffee	tea	tea	water	water	water	water	
0	1	2	3	4	5	6	7	8	

counts =
$$\{3, 2, 4\}$$
 starts = $\{3, 5, 9\}$

coffee	coffee	coffee	tea	tea	water	water	water	water
0	1	2	3	4	5	6	7	8

stable?

counts =
$$\{3, 2, 4\}$$
 starts = $\{3, 5, 9\}$

coffee	coffee	coffee	tea	tea	water	water	water	water
0	1	2	3	4	5	6	7	8

stable? Yes!

counts =
$$\{3, 2, 4\}$$
 starts = $\{3, 5, 9\}$

coffee	coffee	coffee	tea	tea	water	water	water	water
0	1	2	3	4	5	6	7	8

Runtime? Let # distinct elements = K.

counts =
$$\{3, 2, 4\}$$
 starts = $\{3, 5, 9\}$

coffee	coffee	coffee	tea	tea	water	water	water	water
0	1	2	3	4	5	6	7	8

Runtime? Let # distinct elements = K. $\Theta(N+K)$

Extending Counting Sort

- Although counting sort only works if there are a small set of values, there are many cases that items to be sorted are composed of elements that can only take on a small set of values.
- For example, numbers in base 10 like 1283642 are made up of digits, which can only be one of ten values (0 through 9).
- Can we use counting sort on each digit to quickly sort a list of numbers?

- Let's sort from least significant digit to most significant digit.
- We'll use counting sort, but only using one digit at a time as our search key.

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321	32 1
534	36 1
163	16 3
654	81 3
361	53 4
813	65 4
499	49 9

- Let's sort from least significant digit to most significant digit.
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		4
321	32 1	8 1 3
534	36 1	3 2 1
163	16 3	5 3 4
654	81 3	6 5 4
361	53 4	3 6 1
813	65 4	1 6 3
499	49 9	4 9 9

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- We'll use counting sort, but only using one digit at a time as our search key.

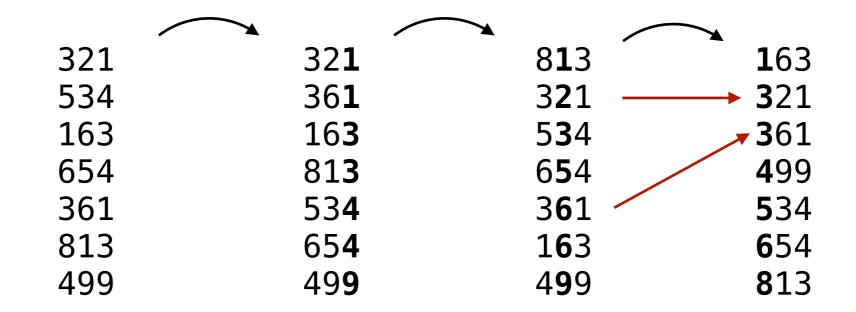
321	32 1	8 1 3	1 63
534	36 1	3 2 1	3 21
163	16 3	5 3 4	3 61
654	81 3	6 5 4	4 99
361	53 4	3 6 1	5 34
813	65 4	1 6 3	6 54
499	49 9	4 9 9	8 13

- Let's sort from least significant digit to most significant digit.
- We'll use counting sort, but only using one digit at a time as our search key.

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534	36 1	3 2 1	→ 3 21
163	16 3	5 3 4	→3 61
654	81 3	6 5 4	4 99
361	53 4	3 6 1	5 34
813	65 4	1 6 3	6 54
499	49 9	4 9 9	8 13

Notice how it is essential that counting sort is stable!

- Let's sort from least significant digit to most significant digit.
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Runtime? Let D be max # of digits and K be # of distinct possible digits.

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Runtime? Let D be max # of digits and K be # of distinct possible digits. E.g. for base 10, K = 10.

$$\Theta(D(N+K))$$

- What about from most significant digit to least significant digit?
- We can't do the same thing as before, since the most-recently applied sort takes priority.

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21

34

63

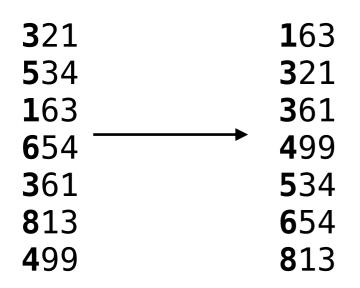
54

61

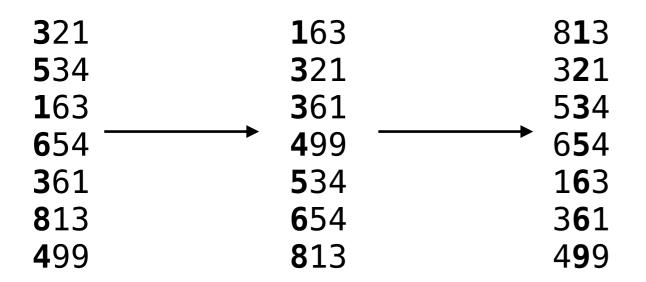
13

99

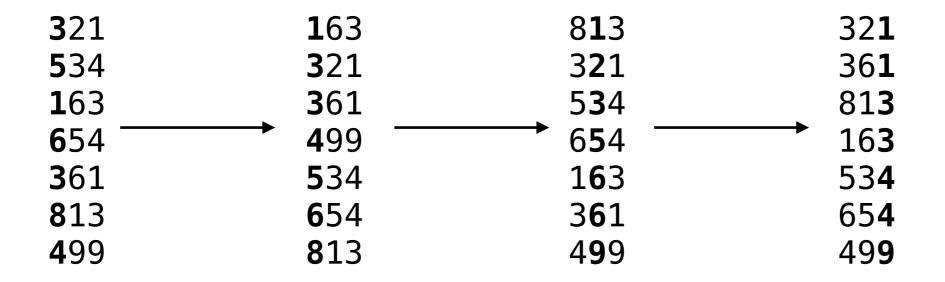
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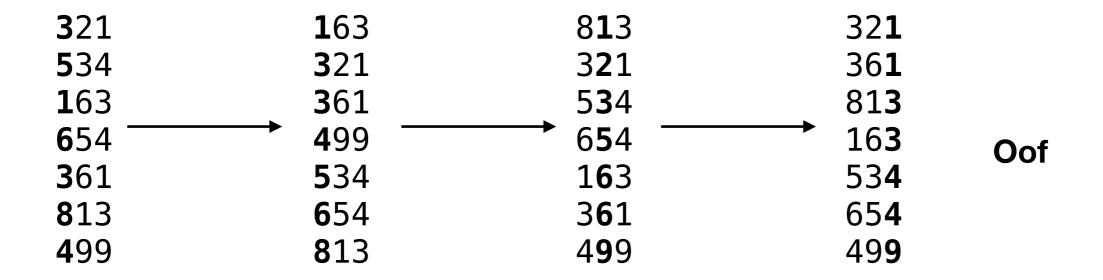
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 We have to group items with the same digit into buckets and then sort the buckets recursively.

99

91

34

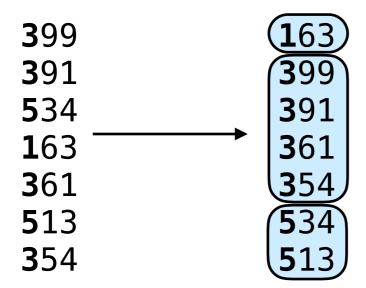
63

61

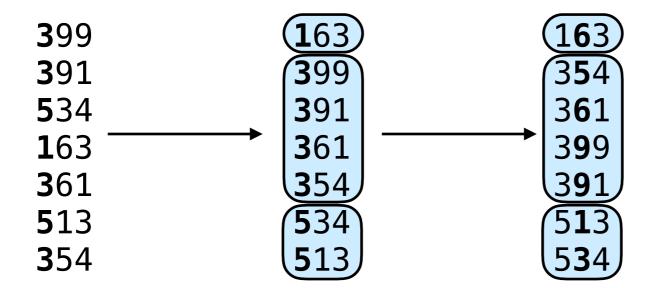
13

54

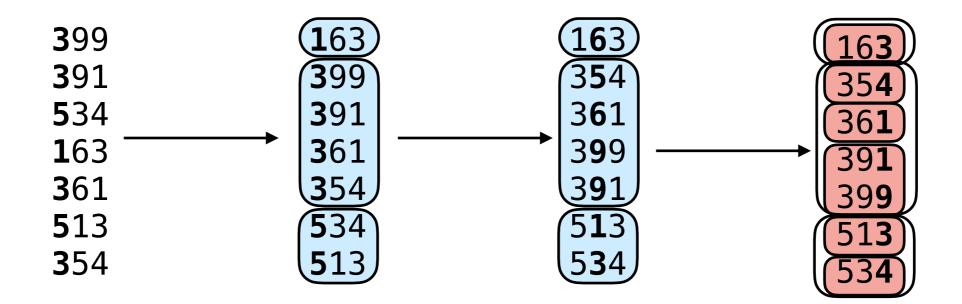
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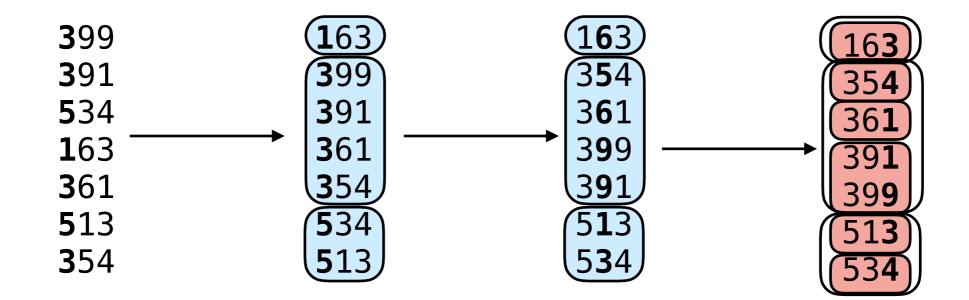
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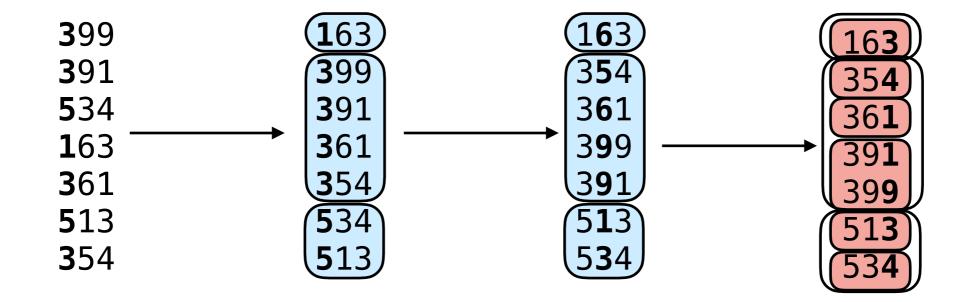


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It's possible that when we sort the left-most digit, all of the most significant digits are unique, so we're done in one pass, but it's also possible we have to recurse through all of the digits. So $\Omega(N+K)$, O(D(N+K))