1.3

- Values: 1 ✓ 2 ✗ 3 ✗ 4 ✗
  5 ✗ 6 ✓ 7 ✗ 8 ✗

- "is-val" relation : " e val "

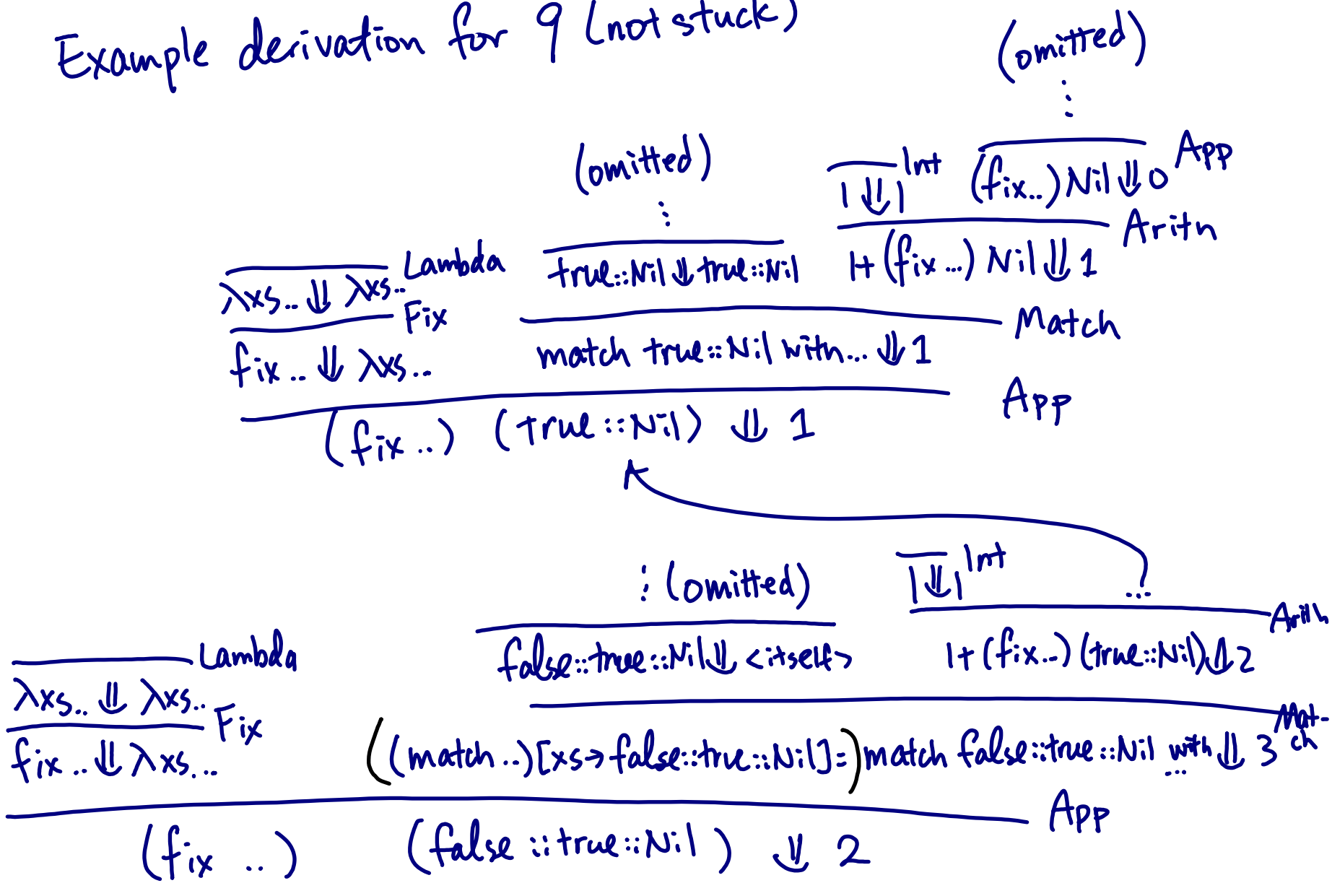
Rules :

$$\frac{}{i \text{ val}}\ \text{V-Int} \qquad \frac{}{(\lambda x.\, e)\ \text{val}}\ \text{V-Lambda}$$

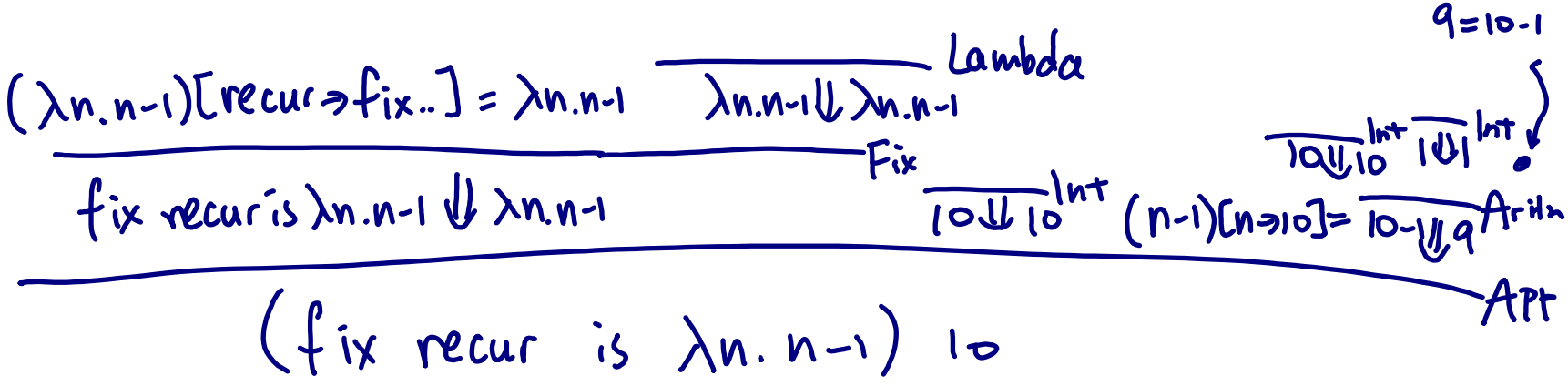
$$\frac{}{\text{Nil val}}\ \text{V-Nil} \qquad \frac{e_1 \text{ val} \quad e_2 \text{ val}}{(e_1 :: e_2)\ \text{val}}\ \text{V-Cons}$$

$e \Downarrow v$: Check v using the CSIL reference interpreter

Example derivation for 9 (not stuck)



11 :


(λn.n-1)[recur→fix..] = λn.n-1
Lambda
λn.n-1 ⇓ λn.n-1
Fix
fix recur is λn.n-1 ⇓ λn.n-1
10 ⇓ 10 Int
(n-1)[n→10] = 10-1 ⇓ 9 Arith
9=10-1
10 ⇓ 10 Int
1 ⇓ 1 Int
App
(fix recur is λn. n-1) 10

What-if:

App-Alt 1: "$(\lambda x.\lambda y.\ x+y)\ 3\ 2$" gets stuck

App-Alt 2: "$(\lambda x.\ 0)\ (1+\text{true})$" will not get stuck

"$(\lambda x.\ x+x)\ (1+2)$" will have a larger tree.

Let-Alt:
- "let x = 1+true in 0" gets stuck with $\Downarrow$ but not $\Downarrow_1$.
- If e gets stuck with $\Downarrow_1$, e will also get stuck with $\Downarrow$.
- let x = 2*3 in x+x"

MatchCons-Alt:

$$\frac{e_1 \Downarrow h::t \quad (h \overset{\text{optional}}{\Downarrow} v_1)(t \overset{\text{optional}}{\Downarrow} v_2) \quad e_3[x \mapsto v_1][y \mapsto v_2] \Downarrow v}{\text{match } e_1 \text{ with } Nil \to e_2 \mid x::y \to e_3 \text{ end} \Downarrow v}$$

(or $[x \mapsto h][y \mapsto t]$)

$$\frac{}{e_1 :: e_2 \quad \text{val}} \; \text{V-ConsAlt}$$