1.3

o Values: 1 / 2 × 3 × 4 × 5 × 6 √ 7 × 8 ×

a is-val relation: "e val"

Rules:

i val

()x.e) val V-Lambda

Nil val V-Nil (e1::e2)val V-Cons

ellv: C	heck v using the CS	Il reference interpreter
Example deriv	vation for 9 Lnot stuck	(omitted)
	(omitted)	Wil H (fix) Nil U o App H (fix) Nil U 1
fix U >xs motch true:: N:  with U1		
	(fix) (true::Nil) U	App App
	: (omitte	ed) Illimt
Tambda  Axs III. Axs.	folse::true::NIU <	itself> It (fix) (true::Nil)(12
Tix U Xxs Fix	Wil]=)match false::true::Nil with [] 3 ch	
(fix)	(false ::true::Nil)	U2

11:

 $(\lambda n. n-1) \text{ [recurs fix...]} = \lambda n. n-1 \qquad \lambda n. n-1 \text{ [lambda]}$   $fix recurs \lambda n. n-1 \text{ [low lot]} \text{ [n-1) [n-10]} = 10-y \text{ [arthouse of the local of the$ 

```
What - if:
                  "()x. ly. x+y) 3 2 " gets stuck
   App- Al+1:
                  " ()x.0) (Ittrue)" will not get stuck
  App- Alt 2:
                  " ()x.x+x) (1+2)" will have a larger
  Let-Alt: « "let X= 1+true in D" gets stuck with U but not U1.
             · If e gets stuck with U1, e will also get stuck
                   with U.
             · let x= 2*3 in x+x"
                                               or[x+h]ty=t]
                     eill h::+ (hllvi)(tllvz) ez[x+vi][y=vz]llv
 Match Cons - Alt:
```

match e, with Nilsezl X:: yse, end ll v

e, :: ez val V-Cons Alt