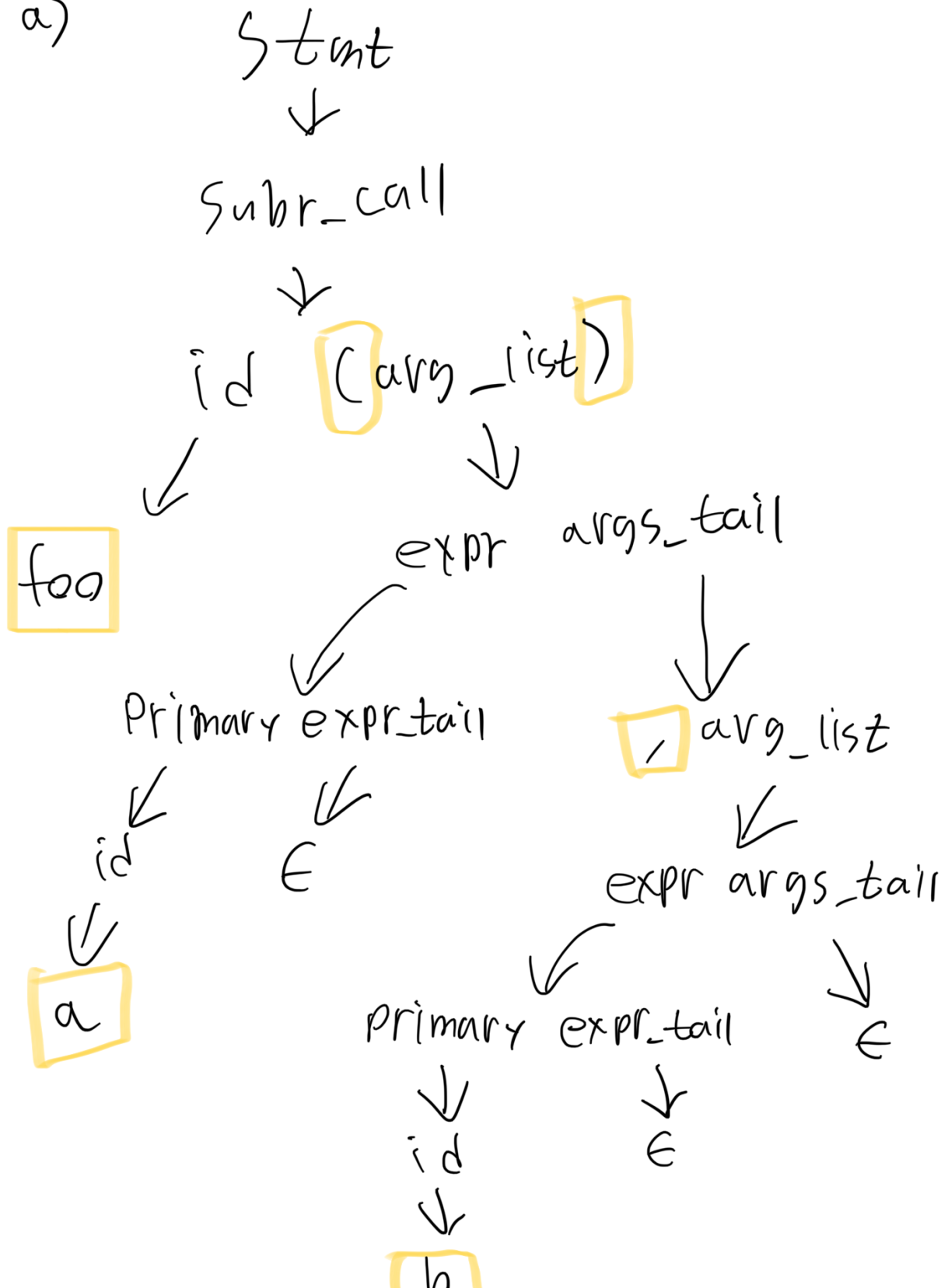


# HW2

Page 108-112: 2.13, 2.17, 2.18 and 2.40 (use python)

2.13

a)





(b)

foo(a, b)

stmt	→	subr_call
	→	id ( arg_list )
	→	id ( expr args_tail )
	→	id ( expr , arg_list )
	→	id ( expr , expr args_tail )
	→	id ( expr , expr $\epsilon$ )
	→	id ( expr , primary expr_tail
$\epsilon$ )		
	→	id ( expr , primary $\epsilon \epsilon$ )
	→	id ( expr , id $\epsilon \epsilon$ )
	→	id ( primary expr_tail , id $\epsilon \epsilon$ )
	→	id ( id $\epsilon$ , id $\epsilon \epsilon$ )

(c)

It can't be LL(1). Because there is an ambiguity in the "expr" non-terminal to "id". We cannot determine whether to expand "expr" into "primary" followed by an "id" or a "subr call".

(d)

<i>stmt</i>	→	<i>id stmt_tail</i>
<i>stmt_tail</i>	→	<i>id := expr</i>
	→	<i>( arg list )</i>
<i>expr</i>	→	<i>primary expr_tail</i>
<i>expr_tail</i>	→	<i>op expr</i>
	→	$\varepsilon$
<i>primary</i>	→	<i>id primary_tail</i>
<i>primary</i>	→	<i>( arg_list )</i>
	→	<i>( expr )</i>
<i>primary_tail</i>	→	$\varepsilon$
<i>op</i>	→	<i>+   -   *   /</i>
<i>arg_list</i>	→	<i>expr args_tail</i>
<i>args_tail</i>	→	<i>, arg list</i>
	→	$\varepsilon$

2.17

<i>program</i>	→	<i>stmt list \$\$</i>
<i>stmt_list</i>	→	<i>stmt list_stmt   <math>\varepsilon</math></i>
<i>stmt</i>	→	<i>id := expr</i>
		<i>  read id</i>
		<i>  write expr</i>
		<i>  if condition then stmt_list fi</i>

| while condition do stmt\_list od  
 condition  $\rightarrow$  expr comp\_op expr  
 comp\_op  $\rightarrow$  < | <= | > | >= | == | !=  
 expr  $\rightarrow$  term expr\_tail  
 expr\_tail  $\rightarrow$  add\_op term expr\_tail |  $\epsilon$   
 term  $\rightarrow$  factor term\_tail  
 term\_tail  $\rightarrow$  mult\_op factor term\_tail |  $\epsilon$   
 factor  $\rightarrow$  ( expr )  
           | id  
           | number  
 add\_op  $\rightarrow$  + | -  
 mult\_op  $\rightarrow$  \* | /

2.18

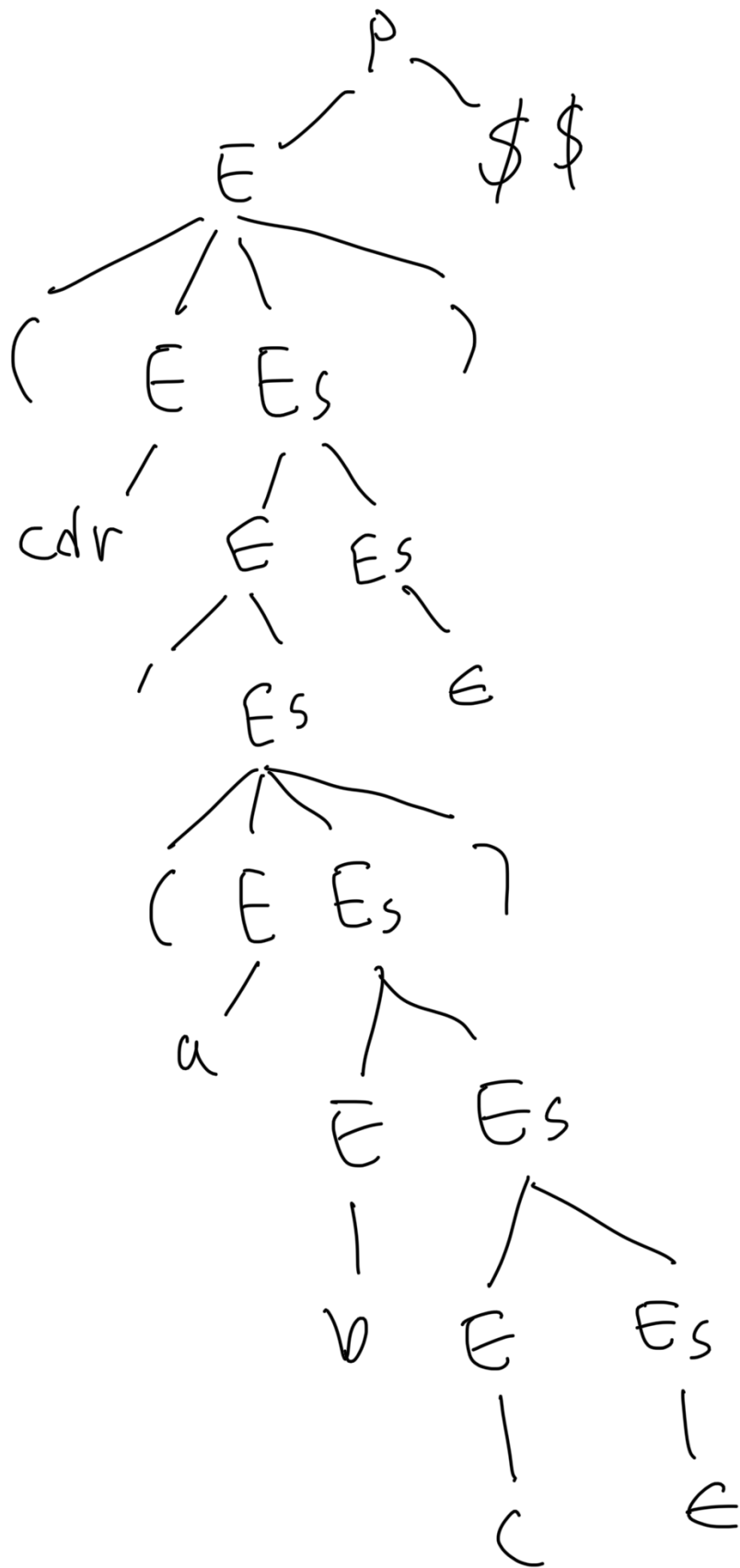
(a)

FIRST(Es) : atom, ' (, ),  $\epsilon$

FOLLOW (E): atom, ' (, ), \$\$

PREDICT (Es  $\rightarrow \epsilon$ ): )

(b)



(c)

P → E \$\$  
 → ( E Es ) \$\$  
 → ( cdr E s ) \$\$  
 → ( cdr E Es ) \$\$  
 → ( cdr ' E E s ) \$\$  
 → ( cdr ' ( E Es ) Es ) \$\$  
 → ( cdr ' ( a Es ) Es ) \$\$  
 → ( cdr ' ( a E Es ) Es ) \$\$  
 → ( cdr ' ( a b Es ) Es ) \$\$  
 → ( cdr ' ( a b E Es ) Es ) \$\$  
 → ( cdr ' ( a b c Es ) Es ) \$\$  
 → ( cdr ' ( a b c ) Es ) \$\$  
 → ( cdr ' ( a b c ) ) \$\$

(d)

Parse stack	Input stream	Comment
P	(cdr ' (a b c)) \$\$	
E\$\$	(cdr ' (a b c)) \$\$	Predict P → E \$\$
( E Es ) \$\$	(cdr ' (a b c)) \$\$	Predict E → ( E Es )

E Es ) \$\$	cdr ' (a b c)) \$\$	Match (
atom Es ) \$\$	cdr ' (a b c)) \$\$	Predict E → atom
Es ) \$\$	' (a b c)) \$\$	Match atom
E Es ) \$\$	' (a b c)) \$\$	Predict Es →E Es
' E Es ) \$\$	' (a b c)) \$\$	Predict E → ' E
E Es ) \$\$	(a b c)) \$\$	Match '
( E Es ) Es ) \$\$	(a b c)) \$\$	Predict E → ( E Es )
E Es ) Es ) \$ \$	a b c)) \$\$	Match (
atom Es ) Es ) \$\$	a b c)) \$\$	Predict E → atom
Es ) Es ) \$\$	b c)) \$\$	Match atom
E Es ) Es ) \$ \$	b c)) \$\$	Predict Es → E Es
atom Es ) Es ) \$\$	b c)) \$\$	Predict E → atom
Es ) Es ) \$\$	c)) \$\$	Match atom



E Es ) Es ) \$ \$	c)) \$\$	Predict Es → E Es
atom Es ) Es ) \$\$	c)) \$\$	Predict E → atom
Es ) Es ) \$\$	)) \$\$	Match atom
) Es ) \$\$	)) \$\$	Predict Es → $\epsilon$
Es ) \$\$	) \$\$	Match )
) \$\$	) \$\$	Predict Es → $\epsilon$
\$\$	\$\$	Match )

(e)

"P", "E", "Es", "match"

2.40

Easier. I think that Python's amazing readability is the main reason why Python is so popular. The whitespaces makes the code neat and the lack of braces and parentacies

meke the code cleaner.