

double circle make a token so 4b would be two tokens 4 and b

1. $E \longrightarrow T\ TT$
 - $\triangleright TT.st := T.val$ $\triangleright E.val := TT.val$
2. $TT_1 \longrightarrow +\ T\ TT_2$
 - $\triangleright TT_2.st := TT_1.st + T.val$ $\triangleright TT_1.val := TT_2.val$
3. $TT_1 \longrightarrow -\ T\ TT_2$
 - $\triangleright TT_2.st := TT_1.st - T.val$ $\triangleright TT_1.val := TT_2.val$
4. $TT \longrightarrow \epsilon$
 - $\triangleright TT.val := TT.st$
5. $T \longrightarrow F\ FT$
 - $\triangleright FT.st := F.val$ $\triangleright T.val := FT.val$
6. $FT_1 \longrightarrow *\ F\ FT_2$
 - $\triangleright FT_2.st := FT_1.st \times F.val$ $\triangleright FT_1.val := FT_2.val$
7. $FT_1 \longrightarrow /\ F\ FT_2$
 - $\triangleright FT_2.st := FT_1.st \div F.val$ $\triangleright FT_1.val := FT_2.val$
8. $FT \longrightarrow \epsilon$
 - $\triangleright FT.val := FT.st$
9. $F_1 \longrightarrow -\ F_2$
 - $\triangleright F_1.val := - F_2.val$
10. $F \longrightarrow (E)$
 - $\triangleright Eval := E.val$
11. $F \longrightarrow \text{const}$
 - $\triangleright Eval := \text{const}.val$

$\text{EPS}(a)$: if it products eps. then true

$\text{First}(a)$: the first terminal of the productions

$\text{Follow}(a)$: the first terminal comes after a
(derive non-terminals into terminals)

and if a is produced by non-terminal b then add $\text{Follow}(b)$ to $\text{Follow}(a)$

$\text{Predict}(a \rightarrow b)$: $\text{First}(b)$

if $\text{EPS}(b)$ true then add

$\text{Follow}(a)$

$\text{Program} \rightarrow \text{Stmts} \ \$\$$

$\text{Stmts} \rightarrow \text{Stmt} \ \text{Stmts_Tail}$

$\text{Stmt} \rightarrow \text{until} \ (\text{Stmts}, \text{id}) \mid \text{id}$

$\text{Stmts_Tail} \rightarrow ; \ \text{Stmt} \ \text{Stmts_Tail} \mid \epsilon$

short circuit evaluation check one thing

that matters

non-short

check everything

| | | |
|----------------------------|------------------|-------|
| EPS(Program) | | false |
| EPS(Stmts) | | false |
| EPS(Stmt) | | false |
| EPS(Stmts_Tail) | | true |
| First(Program) | until | id |
| First(Stmts) | until | id |
| First(Stmt) | until | id |
| First(Stmts_Tail) | ; | |
| Follow(Program) | empty set | |
| Follow(Stmts) | \$\$ | , |
| Follow(Stmt) | ; | |
| Follow(Stmts_Tail) | \$\$ | , |
| Predict(Stmts_Tail -> eps) | \$\$ | , |
| | EPS(eps) is true | |

Scoping

keep the track of globally and locally declared variables

static scoping

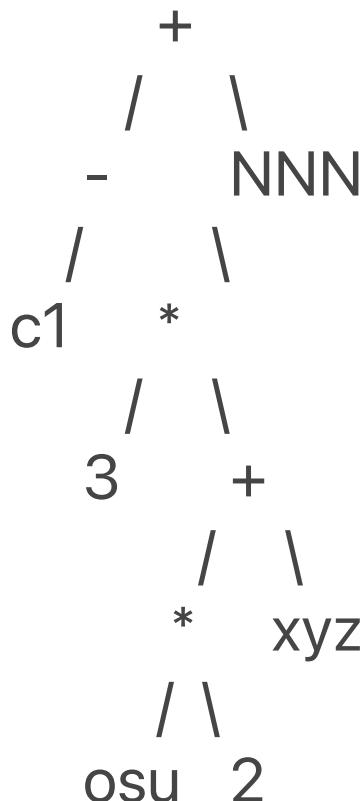
no locally declared var, looks for global var

dynamic scoping

no locally declared var, looks for parent function var

draw a tree for

$c1 - 3 * (osu * 2 + xyz) + NNN$



2. Suppose that we use the CFG and the attribute grammar in Fig. 8-18
its decoration for the expression

$$10 * (8 - 4)$$

