

### 8.32

# block offset bits =  $\log_2(\text{block size}) = \log_2(64\text{B}) = 6$  bits  
# blocks in the cache =  $32\text{K} / 64\text{B} = 32 \cdot 2^{10} / 64 = 512$   
# index bits =  $\log_2(\text{\#blocks}/\text{\#sets}) = \log_2(512/4) = 7$  bits  
# tag bits =  $32 - 6 - 7 = 19$  bits

### 8.33

a)

#sets = #blocks / associativity

#blocks = #sets \* associativity

#index bits =  $\log_2(\text{\#sets}) = 8$

#sets =  $2^8 = 256$

#blocks =  $256 * 2 = 512$

b)

#address bits = #tag bits + #index bits + #block offset bits  
=  $14 + 8 + 2 = 24$

24 bit addresses implies  $2^{24}$  bytes of memory

therefore

#memory blocks =  $2^{24} \text{ bytes} / 4 \text{ bytes per block} = 2^{22}$  memory blocks

c)

d)

### 8.34

8B per object  
128 blocks in the cache  
8 objects per block

*A*

Misses once every 8  $j$  cycles  
 $1024^2 j$  cycles / 8 misses per  $j$  cycle = 131072 misses

*B*

Misses every  $j$  cycle once every 8  $i$  cycles  
 $1024 j$  cycles \*  $1024 i$  cycles \*  $\frac{1}{8}$  misses per  $i$  cycle = 131072 misses

### 8.36

24 bit virtual addresses  
16 bit physical addresses  
1K page size

# virtual pages =  $2^{24} / 1K = 2^{24} / 2^{10} = 2^{14} = 16K$  pages  
# physical pages =  $2^{16} / 1K = 2^{16} / 2^{10} = 2^6 = 64$  pages  
# bits in virtual page number =  $\log_2(16K) = \log_2(2^{14}) = 14$  bits  
# bits in physical page number =  $\log_2(64) = \log_2(2^6) = 6$  bits  
# bits in page offset =  $\log_2(1K) = \log_2(2^{10}) = 10$  bits

### 8.37

*a)*  
access violation  
*b)*  
access violation  
*c)*  
access violation