# Tracking Errors through Types in Quantum Programs

Kesha Hietala, *Robert Rand*, Michael Hicks Programming Languages for Quantum Computing

### Overview

- Errors, particularly those introduced by gate application, will be prevalent in near-term quantum machines, so need to be taken into account when writing programs.
- It might be useful to have a notion of errors at the programming language level
- We have designed a simple type system to track errors in quantum programs, implemented in QWIRE

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- A linear type system for enforcing no-cloning
- A denotational semantics in terms of density matrices

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**W2** 

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Inductive Circuit (W : WType) : Type :=
 output : Pat W -> Circuit W
 gate
           Gate W1 W2 -> Pat W1 ->
            (Pat W2 -> Circuit W) -> Circuit W
 lift : Pat Bit -> (bool -> Circuit W) ->
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                      <del>------</del> р : W
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$$|0\rangle_E = (|000\rangle + |111\rangle) \otimes (|000\rangle + |111\rangle) \otimes (|000\rangle + |111\rangle)$$
  
$$|1\rangle_E = (|000\rangle - |111\rangle) \otimes (|000\rangle - |111\rangle) \otimes (|000\rangle - |111\rangle)$$

# Adding Errors

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Inductive Gate : nat -> WType -> WType -> Set := \mid U \mid : (U : Unitary k W) -> Gate k W (map W (k + sum_err W)) \mid init0 : Gate \epsilon_i One (Qubit \epsilon_i) \mid new0 : Gate 0 One Bit \mid meas : Gate \epsilon_m (Qubit n) Bit \mid discard : Gate 0 Bit One \mid EC : Gate \epsilon_e (Qubit n) (Qubit 0).
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#### Fault Tolerance

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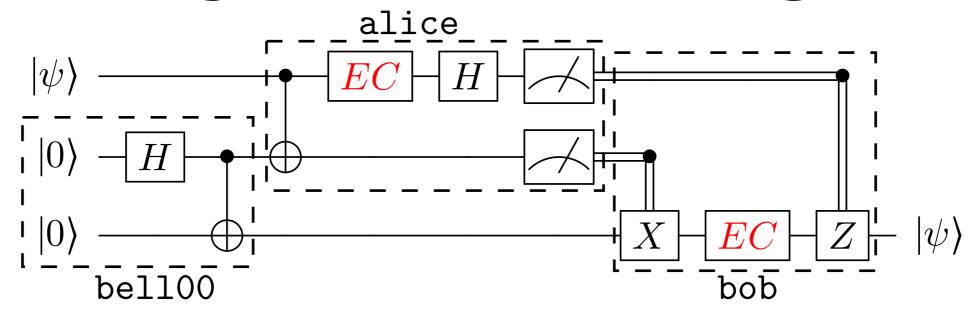
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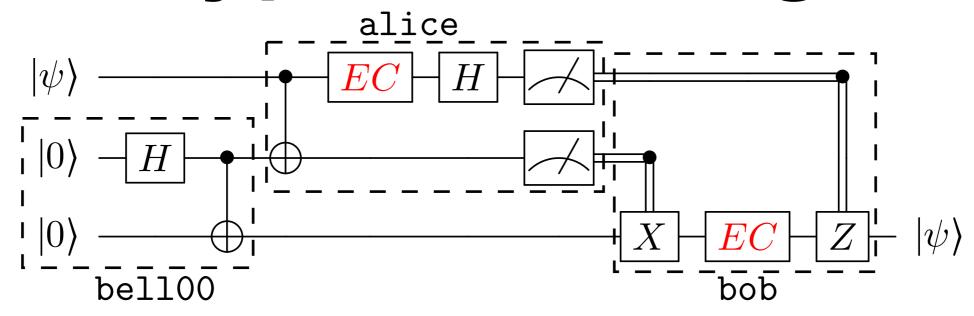
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 implies  $decode(EC; | \psi \rangle) = decode(| \psi \rangle)$ 

We check these using QWIRE's type system

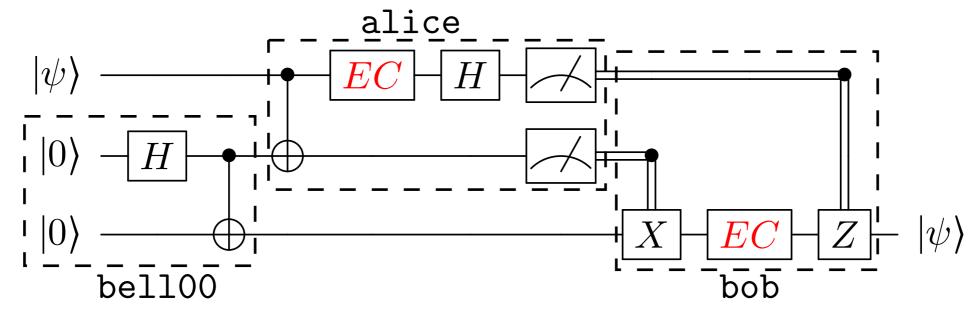
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- We benefit from linearity we cannot introduce an error onto an existing qubit, but only produce a new qubit with some number of errors
- However, checking fault tolerance is orthogonal to checking linearity.



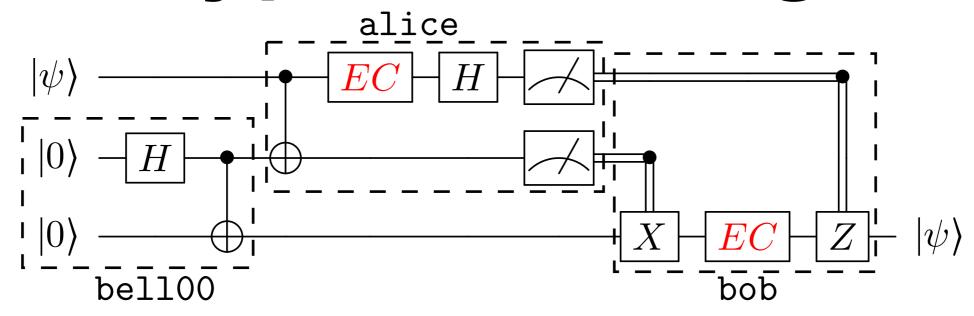


**Definition** bell : Box One (Qubit 2 ⊗ Qubit 2) := ...



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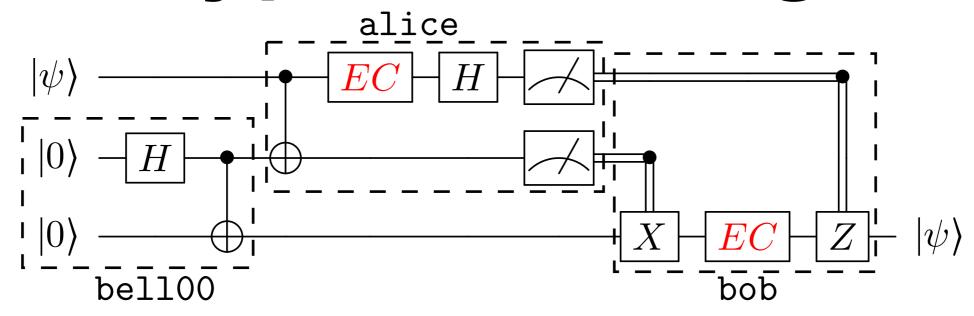
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Definition bob : Box (Bit ⊗ Bit ⊗ Qubit 2) (Qubit 1) := ...

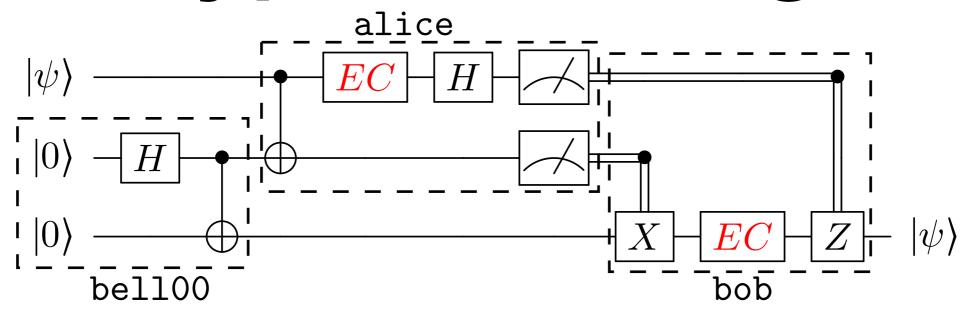


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Definition teleport : Box (Qubit 0) (Qubit 1) :=



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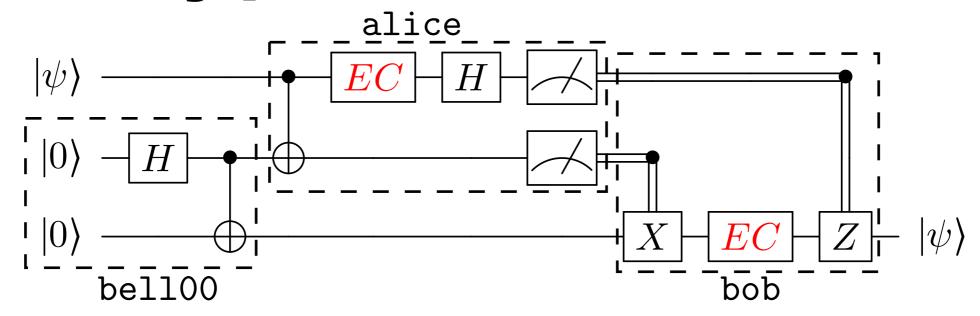
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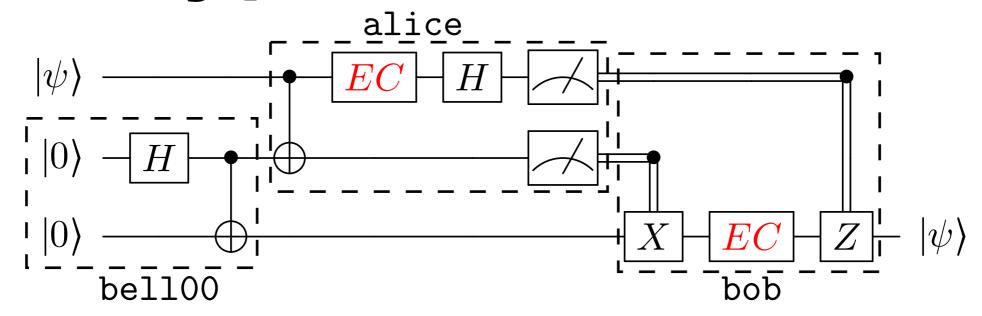
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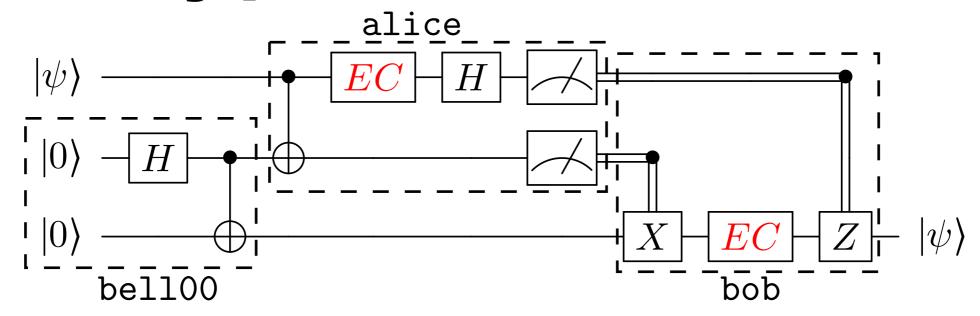
Lemma teleport_EC_WT : Typed_Box teleport 3.

Proof. type_check. Qed.
```



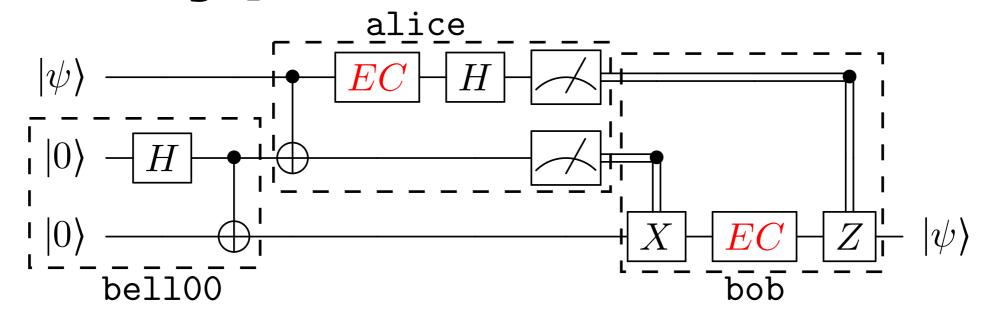


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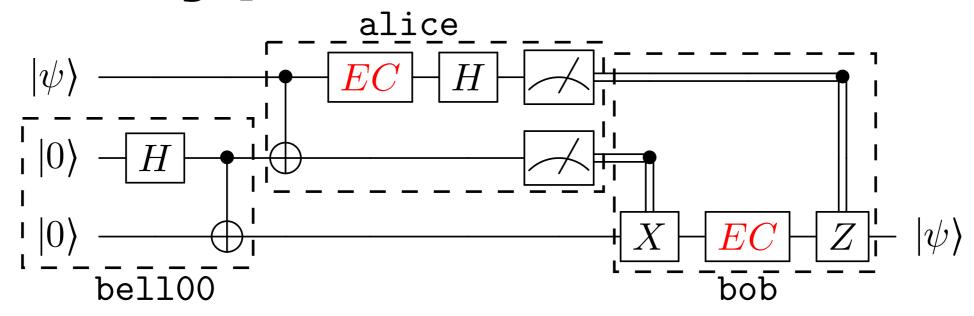
Check

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#### Check

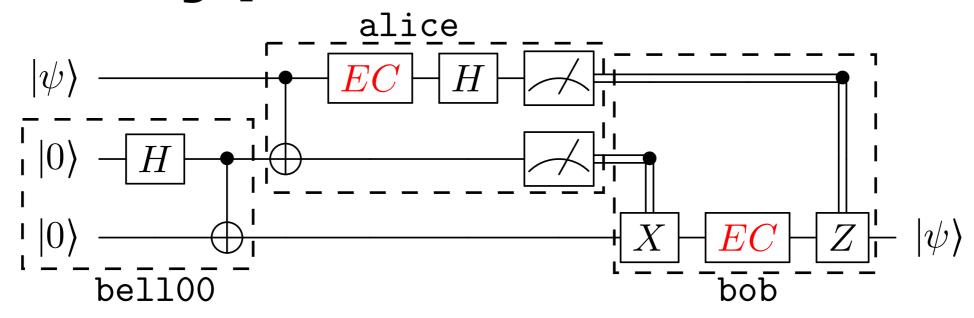
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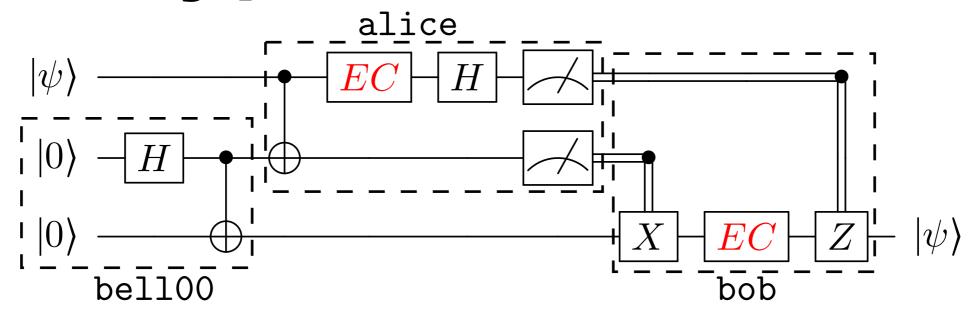
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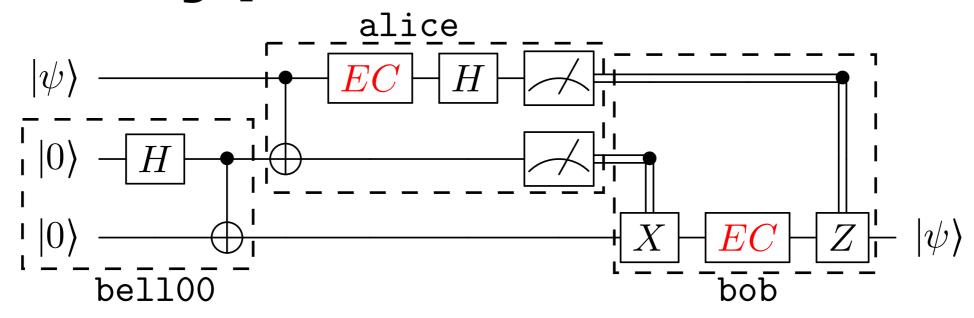
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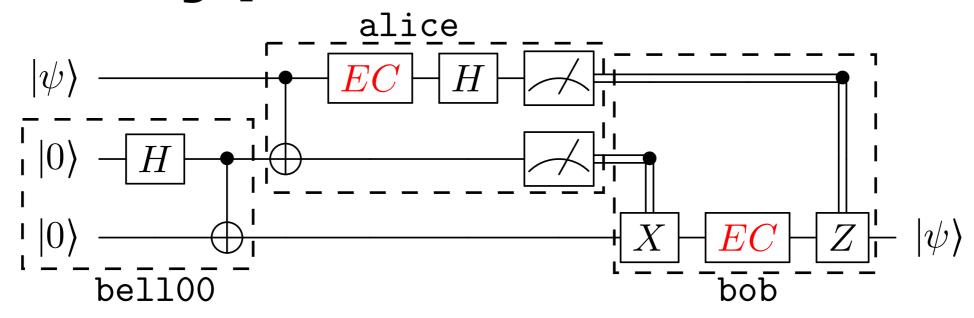


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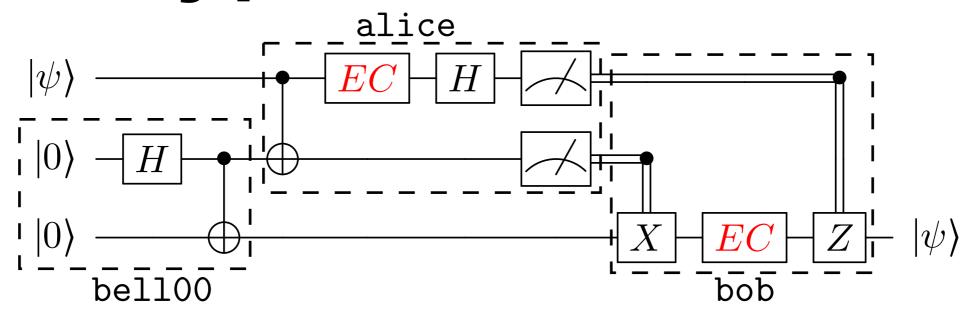
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: Box (Qubit 0) (Qubit 1)

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#### Questions

- How can we reflect our error handling back into QWIRE's denotational semantics?
- What about errors due to decoherence, that aren't captured by gate application?
- What are better ways of handling errors in a type system or broader reasoning system?