CS2010 PS7 - A Trip to Supermarket v4 (with R-option)

Released: Sunday, 02 November 2014, 8pm

Due: Thursday, 13 November 2014, 08.59am (Happy Birthday Grace)

Collaboration Policy. You are encouraged to work with other students or teaching staffs (inside or outside this module) on solving this problem set. However, you must write Java code by yourself. In addition, when you write your Java code, you must list the names of every collaborator, that is, every other person that you talked to about the problem (even if you only discussed it briefly). This list may include certain posts in CS2010 Facebook group. If you have access to your seniors' CS2010 files (that is, problem sets version 1/2/3), please refrain from looking at their code verbatim. Any deviation from this policy will be considered as cheating. If the offender is caught beyond reasonable doubt, he/she will be punished severely, including referral to the NUS Board of Discipline. It is not worth it to cheat just to get 19% when you will lose out in the other 81%.

R-option. This PS has R-option at the back. This additional task usually requires understanding of data structure or algorithm *beyond* CS2010. A self research on those relevant additional topics will be needed but some pointers will be given. CS2010R students *have to* attempt this R-option. CS2010 students can choose to attempt this R-option too for higher points, or simply leave it.

2011 Story. Before Jane's birth, Grace could not travel far and could not carry heavy stuffs (see PS4). After Jane's birth, Grace needs a lot of time to rest and to breastfeed Jane. Therefore, Steven has to help Grace with a certain household chore that he thought he does not need to do anymore after he got married: Grocery Shopping. (To the guys reading this, get ready:D).

The Actual Problem. Today, Steven has to visit a supermarket at West Coast area to buy groceries (name omitted to avoid indirect advertising). Steven has visited this place numerous times and therefore he knows the location of various N items in that supermarket. Steven has estimated the *direct* walking time from one point to every other points in that supermarket (in seconds) and store that information in a 2D table T of size $(N+1) \times (N+1)$. Given a list of K items to be bought today, he wants to know what is the minimum amount of time to complete the shopping duty of that day. As mentioned several time in lectures, we have to make some simplifying assumptions in order for this problem to be solvable...:

- 1. Steven always starts at vertex 0, the entrance (+ cashier section) of that supermarket.
- 2. There are V = N + 1 vertices due to the presence of this special vertex 0. The other N vertices corresponds to the N items. The vertices are numbered from [0..N].
- 3. The direct walking time graph that is stored in a 2D table T is a **complete graph**. T is a symmetric square adjacency matrix with $T[i][i] = 0 \ \forall i \in [0..N]$.
- 4. Steven has to grab all K items (numbered from [1..K]) that he has to buy that day, or he will have hard time explaining (to Grace) why he does not buy certain items... $1 \le K \le N$.
- 5. Steven is very efficient, once he arrives at the point that stores item i, he can grab item i into his shopping bag in 0 seconds (e.g. for item 'banana', he does not have to compare the price

of 'banana A' versus 'banana B' and he does not have to select which banana looks nicer, etc...). So, Steven's shopping time is only determined by the total walking time inside that supermarket to grab all the K items.

- 6. In this problem, Steven ends his shopping duty when he arrives at cashier (also at vertex 0) after grabbing all the required K items.
- 7. Steven can grab the K items in **any order**, but the total walking time (i.e. time to walk from vertex $0 \to \text{to}$ various points in the that supermarket in order to grab all the K items \to back to vertex 0) **must be minimized**. Steven can choose to bypass a certain point that is *not in his shopping list* or even *revisit* a point that contains item that he already grab (he does not need to re-grab it again) if this leads to faster overall shopping time.

See below for an example supermarket:

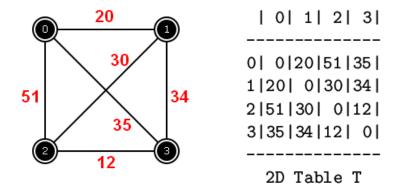


Figure 1: A Sample Supermarket Layout, N = 3, V = 3 + 1 = 4

If today, Steven has to buy all item 1, item 2, and item 3, then one of the best possible shopping route is like this: $0 \to 1 \to 2 \to 3 \to 0$ with a total walking time of: 20+30+12+35=97 seconds.

If tomorrow, Steven has to buy only item 1 and item 2, then one of the best possible shopping route is still: $0 \to 1 \to 2(\to 3) \to 0$ with a total walking time of: 20+30+(12+35)=97 seconds. Notice that although Steven does not have to buy item 3, taking sub path $2 \to 3 \to 0$ (where Steven will just bypass item 3) is faster than taking sub path $2 \to 0$.

If two days later, Steven has to buy only item 2, then one of the best possible shopping route is like this: $0(\to 3) \to 2(\to 3) \to 0$ with a total walking time of: (35+12)+(12+35) = 94 seconds. Steven bypass through item 3 twice!

If three days later, Steven has to buy only item 2 and item 3, then one of the best possible shopping route is like this: $0 \to 3 \to 2(\to 3) \to 0$ with a total walking time of: 35+12+(12+35)=94 seconds. See that Steven can revisit point 3 although he does not have to grab item 3 twice.

These four examples should be clear enough to describe this problem.

The skeleton program Supermarket.java is already written for you, you just need to implement one (or more) method(s)/function(s):

int Query()

You are given a 2D matrix T of size $(N+1) \times (N+1)$ and an array shoppingList of size K. Query these two data structures and answer the query as defined above.

• If needed, you can write additional helper methods/functions to simplify your code.

Subtask A (16 points). The supermarket is a very small convenience store and everything there have to be grabbed/bought. $1 \le K = N \le 9$.

Subtask B (Additional 35 points, i.e. one-level worth of EXP :). The supermarket is still a very small convenience store but slightly bigger than in Subtask A and everything there have to be grabbed/bought. $1 \le K = N \le 15$.

Subtask C (Additional 35 points, i.e. one-level worth of EXP :). The supermarket is a minimart that sells $\approx hundreds$ of grocery items. However, Grace only asks Steven to buy some items. $1 \le N \le 200, 1 \le K \le 15, K \le N$.

Subtask D (Additional 14 points, if you already have 2*7 = 14 points cushion from PS Bonus Subtask A+B, you can ignore this Subtask). The supermarket is a large supermarket that sells \approx thousands of grocery items. However, Grace only asks Steven to buy some items. $1 \le N \le 1000$, $1 \le K \le 15$, $K \le N$.

Subtask E in Mooshak (R-option). The supermarket is a large supermarket that sells \approx thousands of grocery items. However, Grace only asks Steven to buy some items. $1 \le N \le 1000$, $1 \le K \le 29$, $K \le N$ (Notice the new range of K).

WARNING: This Subtask E requires something that I do not teach in CS2010 and can be very frustrating (lots of TLE) if you do not use 'that technique(s)'. CS2010 students, please do NOT attempt this. CS2010R students, you need to figure out what is/are the required technique(s) to solve this problem by consulting CS2010R TA: Pang Yan Han as early as possible. The result of this Subtask E will likely determine the grade A versus A+ for CS2010R as most CS2010R students have successfully completed PS1R, PSBonus, and PS5R.

Note: The official test data has been uploaded to Mooshak online judge, but it is hidden from your view. The time limit setting in Mooshak online judge for Subtask A,B,C,D,E are 1,1,1,3,7 seconds, respectively. This is very strict – especially for the last Subtask E. You are encouraged to generate and post additional test data in Facebook group. Please use SupermarketVerifier.java to verify whether your custom-made test data conform with the required specifications. This program will test if:

- 1. The shopping items are within [1..N].
- 2. The given matrix contains non-negative integers, symmetric, and all entries in the main diagonal of the matrix are zeroes.