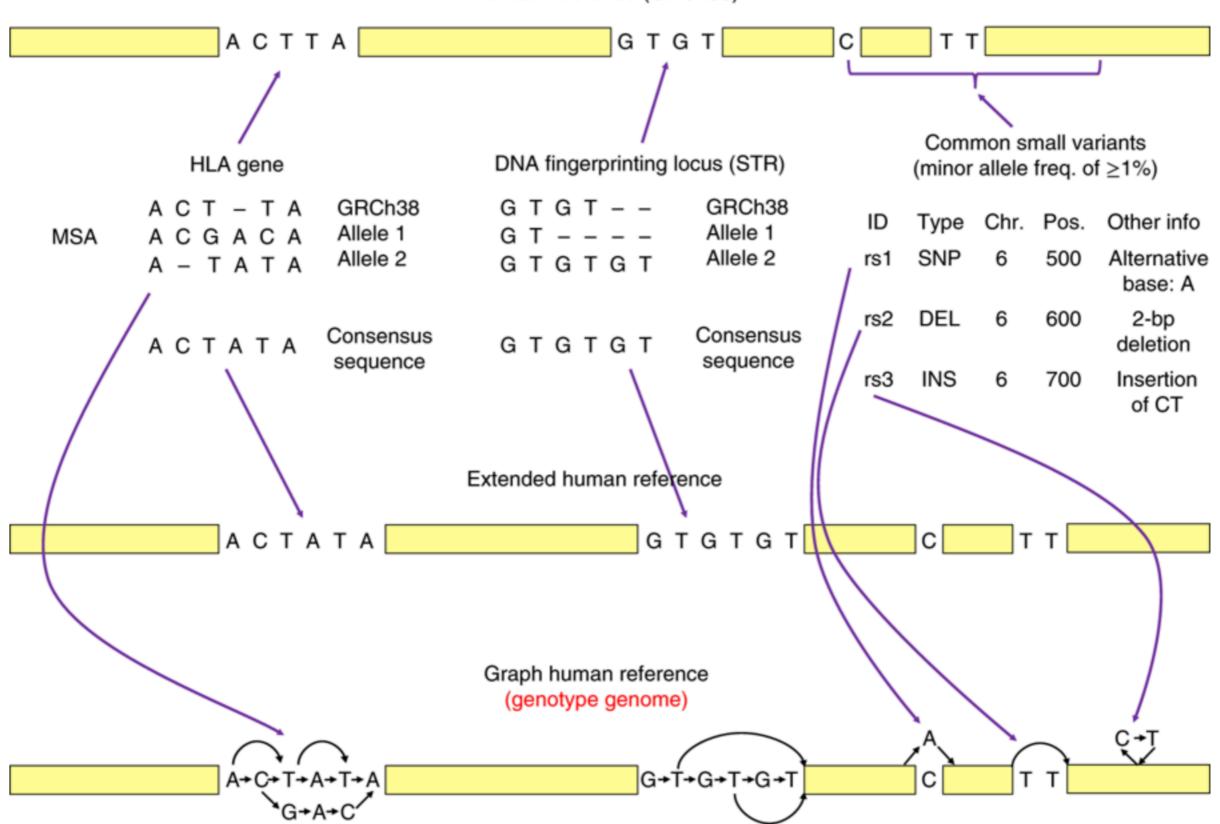


Idea / motivation : No sample is the reference

We have spent a lot of effort characterizing major human variants, yet most aligners simply map against a single human reference genome that doesn't even have the most likely variant at each locus.

HISAT2 is one of a new breed of "graph" aligners, that views the genome as a graph rather than a simple string. This framework allows encoding variants as alternative "paths" through the genome.

Human reference (GRCh38)



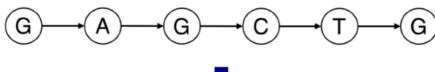
#### Graph FM-Index

Siren, J., Valimaki, N. & Makinen, V. Indexing graphs for path queries with applications in genome research. *IEEE-ACM Trans. Comput. Biol. Bioinform.* **11**, 375–388 (2014).

Construction of graph FM index relies on creation of prefix-range-sorted automata

- Key property needed for backward search:
  - For list (u,v) of outgoing edges, sorted by pairs (p(u), p(v)), I(u)p(v) must be sorted by sequences
  - For any c, all outgoing edges from nodes labeled with c are lexicographically adjacent and are sorted by the prefix p(v) of the destination node
  - All occurrences of c in BWT encode an incoming edge from a node with label c, and thus are sorted by prefix p(v) of the destination
  - Hence, incoming edge labeled by nth occurrence of c is the same as the outgoing edge for rank C[c]+j

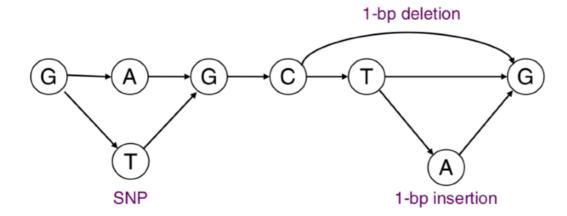
1. Reference sequence (6 bp long)



1

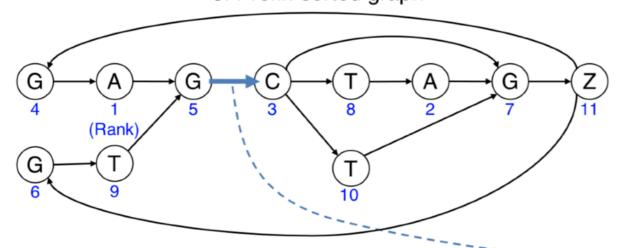
single-nucleotide variant (A/T), a 1-bp deletion (T) and a 1-bp insertion (A)

2. Graphical representation (original graph)



Prefix doubling and pruning

3. Prefix-sorted graph

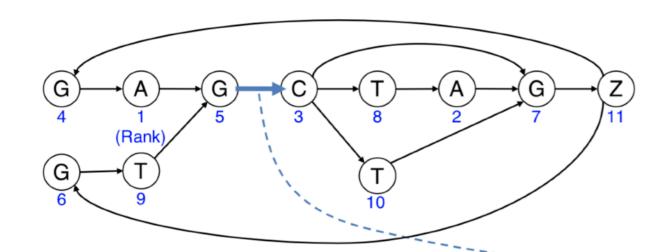


4. Tabular representation of the prefix-sorted graph

Outg edg	joing e(s)			ming e(s)
Node rank	First		Last	Node rank
1	Α		G	1
2	Α		Т	2
	С		G	3
3	С		Z	4
	С		Α	5
4	G		Т	
5	G /		Z	6
6	G		Α	
7	G		С	7
8	Т		Т	
9	Т	-	С	8
10	Т	1	G	9
4.4	Z	1	С	10
11	Z	$\dot{I}$	G	11
,/				

Siren, J., Valimaki, N. & Makinen, V. Indexing graphs for path queries with applications in genome research. *IEEE-ACM Trans. Comput. Biol. Bioinform.* **11**, 375–388 (2014).

(the transformation from 2 -> 3 is crucial to allow indexing)



Outg edg			Inco edg	ming e(s)
Node rank	First		Last	Node rank
1	Α		G	1
2	Α		Т	2
	С		G	3
3	С		Z	4
	С		Α	5
4	G		Т	
5	G /		Z	6
6	G		Α	
7	G	-	С	7
8	Т	i i	Т	
9	Т	-	С	8
10	Т	1	G	9
11	Z	1	С	10
11	Z	,	G	11

Note: We have an LF mapping here, just like a normal BWT

This 1-to-1 correspondence isn't possible without the graph transformation.

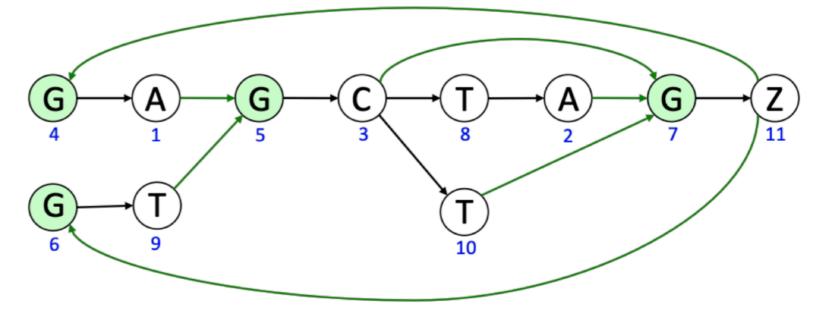
Hint: try and search for the pattern "GTG"

1 *	joing e(s)	Inco edg	ming e(s)
Node rank	First	Last	Node rank
1	A 0	G <sub>0</sub>	1
2	A <sub>1</sub>	Т 0	2
	C <sub>0</sub>	G <sub>1</sub>	3
3	C <sub>1</sub>	Z <sub>0</sub>	4
	C <sub>2</sub>	A 0	5
4	G <sub>0</sub>	Т 1	
5	G <sub>1</sub>	Z <sub>1</sub>	6
6	G <sub>2</sub>	A <sub>1</sub>	
7	G <sub>3</sub>	A <sub>1</sub> C <sub>0</sub>	7
8	T <sub>0</sub>	Т <sub>2</sub>	
9	T <sub>1</sub>	C <sub>1</sub>	8
10	Т 2	G <sub>2</sub>	9
4.4	z <sub>o</sub>	C 2	10
11	Z <sub>0</sub> Z <sub>1</sub>	G <sub>3</sub>	11

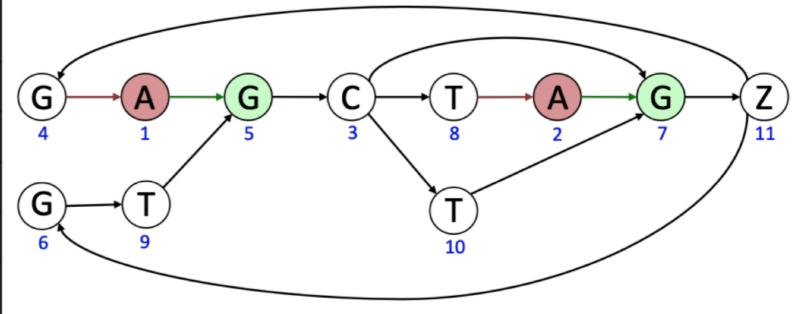
_	joing e(s)		ı	ming e(s)
Node rank	First		Last	Node rank
1	A 0		G <sub>0</sub>	1
2	A <sub>1</sub>		Т 0	2
	C <sub>0</sub>		G <sub>1</sub>	3
3	C 1		Z <sub>0</sub>	4
	C 2		A 0	5
4	G <sub>0</sub>		Т 1	
5	G <sub>1</sub>		Z <sub>1</sub>	6
6	G <sub>2</sub>		A <sub>1</sub> C <sub>0</sub>	
7	G <sub>3</sub>		c <sub>o</sub>	7
8	T <sub>0</sub>		Т <sub>2</sub>	
9	T <sub>1</sub>		C <sub>1</sub>	8
10	T <sub>2</sub>		G <sub>2</sub>	9
11	Z <sub>0</sub> Z <sub>1</sub>		C 2	10
11	Z <sub>1</sub>	٠	G <sub>3</sub>	11

Outgoing edge(s)	
Node rank	First
1	<b>A</b> 0
2	A <sub>1</sub>
	C <sub>0</sub>
3	C <sub>1</sub>
	C <sub>2</sub>
4	Go
5	G <sub>1</sub>
6	G <sub>2</sub>
7	G <sub>3</sub>
8	T <sub>0</sub>
9	T <sub>1</sub>
10	Т <sub>2</sub>
11	Z <sub>0</sub> Z <sub>1</sub>

Incoming edge(s)		
Last	Node rank	
G <sub>0</sub>	1	
T <sub>0</sub>	2	
G <sub>1</sub>	3	
Z <sub>0</sub>	4	
A 0	5	
T <sub>1</sub>		
Z <sub>1</sub>	6	
A <sub>1</sub> C <sub>0</sub>	7	
Т <sub>2</sub>		
C <sub>1</sub>	8	
G <sub>2</sub>	9	
C 2	10	
G <sub>3</sub>	11	

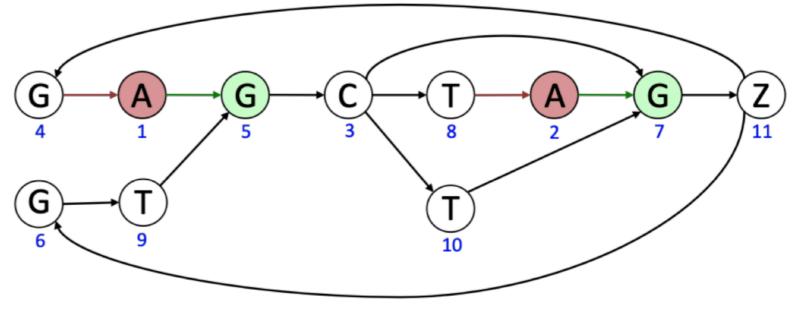


Outg edg	joing e(s)		l .	ming e(s)
Node rank	First		Last	Node rank
1	A <sub>0</sub>		G	1
2	A <sub>1</sub>		T <sub>0</sub>	2
	C 0		G <sub>1</sub>	3
3	C <sub>1</sub>	,,,,,	Z <sub>0</sub>	4
	C <sub>2</sub>	,,,,,	A <sub>0</sub>	5
4	G <sub>0</sub>		T <sub>1</sub>	
5	G <sub>1</sub>	1	Z <sub>1</sub>	6
6	G <sub>2</sub>		A <sub>1</sub>	
7	G <sub>3</sub>		C 0	7
8	T <sub>0</sub>		T 2	
9	T <sub>1</sub>		C <sub>1</sub>	8
10	Т 2		G <sub>2</sub>	9
11	Z <sub>0</sub> Z <sub>1</sub>		C 2	10
''	Z <sub>1</sub>		G <sub>3</sub>	11

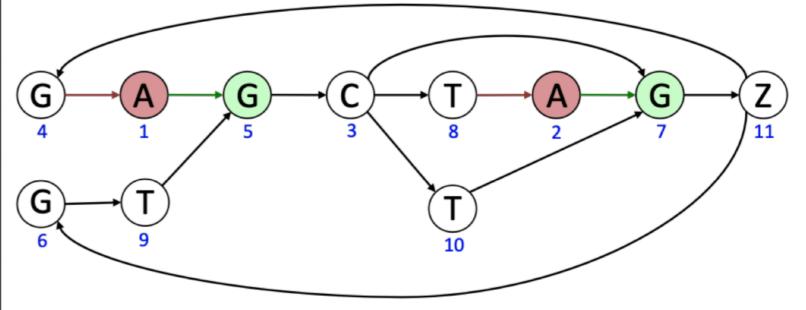


Outg edg	
Node rank	First
1	A <sub>0</sub>
2	A <sub>1</sub>
	C <sub>0</sub>
3	C <sub>1</sub>
	C <sub>2</sub>
4	G <sub>0</sub>
5	G <sub>1</sub>
6	G <sub>2</sub>
7	G <sub>3</sub>
8	T <sub>0</sub>
9	T <sub>1</sub>
10	Т 2
11	Z <sub>0</sub> Z <sub>1</sub>

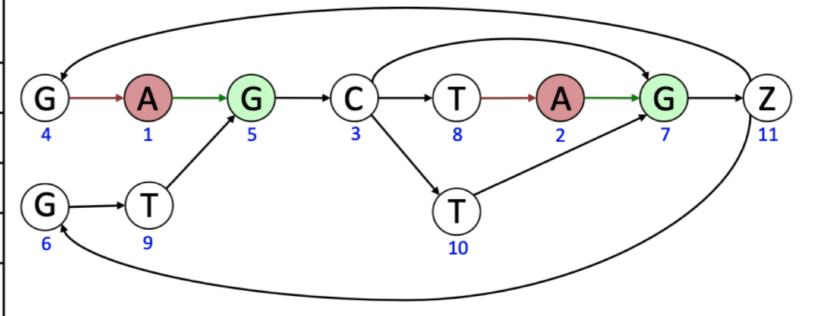
Incoming edge(s)	
Last	Node rank
G <sub>0</sub>	1
Τ <sub>0</sub>	2
G <sub>1</sub>	3
Z <sub>0</sub>	4
A 0	5
T <sub>1</sub>	
Z <sub>1</sub>	6
A 1	
C 0	7
T <sub>2</sub>	
C <sub>1</sub>	8
G <sub>2</sub>	9
C 2	10
G <sub>3</sub>	11



Outg edg			1	ming e(s)
Node rank	First		Last	Node rank
1	A 0		G <sub>o</sub>	1
2	A <sub>1</sub>		T <sub>0</sub>	2
	C <sub>0</sub>	2	G <sub>1</sub>	3
3	C <sub>1</sub>		Z <sub>0</sub>	4
	C <sub>2</sub>	,	A 0	5
4	G <sub>0</sub>		T <sub>1</sub>	
5	G <sub>1</sub>		Z <sub>1</sub>	6
6	G <sub>2</sub>		A 1	
7	G <sub>3</sub>		C 0	7
8	T <sub>0</sub>		Т 2	
9	T <sub>1</sub>		C <sub>1</sub>	8
10	Т2		G <sub>2</sub>	9
11	Z <sub>0</sub> Z <sub>1</sub>		C 2	10
11	Z <sub>1</sub>		G <sub>3</sub>	11

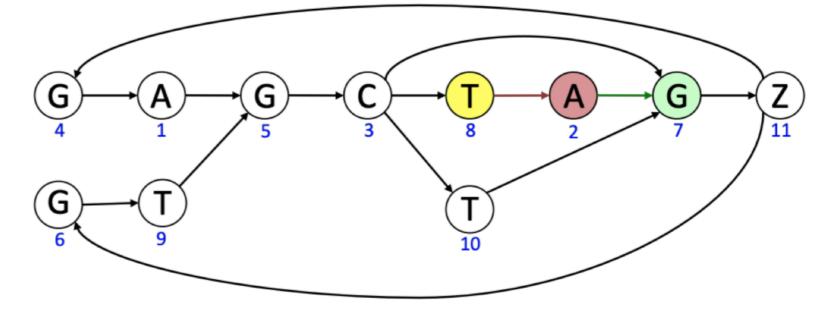


Outg edg			Inco edg	ming e(s)
Node rank	First		Last	Node rank
1	A 0	2	G <sub>0</sub>	1
2	A <sub>1</sub>	3	Τ 0	2
	C <sub>0</sub>		G <sub>1</sub>	3
3	C <sub>1</sub>		Z <sub>0</sub>	4
	C <sub>2</sub>		A 0	5
4	G <sub>0</sub>		T <sub>1</sub>	
5	G <sub>1</sub>		Z <sub>1</sub>	6
6	G <sub>2</sub>		A 1	
7	G <sub>3</sub>		C 0	7
8	T <sub>0</sub>		Т 2	
9	T <sub>1</sub>		C <sub>1</sub>	8
10	Т2		G <sub>2</sub>	9
11	Z <sub>0</sub> Z <sub>1</sub>		C 2	10
11	Z <sub>1</sub>	_	G <sub>3</sub>	11



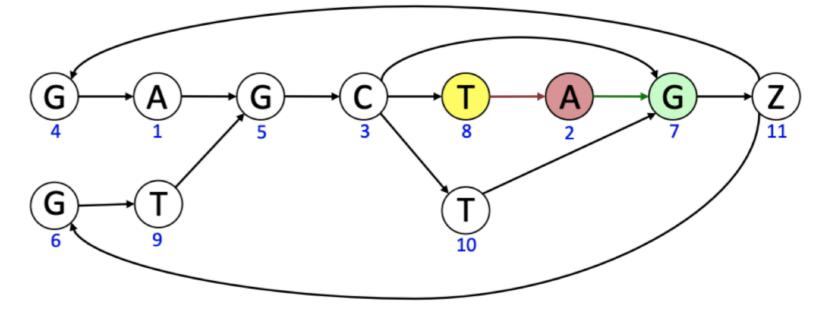
Outg edg	
Node rank	First
1	A 0
2	A <sub>1</sub>
	C <sub>0</sub>
3	C <sub>1</sub>
	C <sub>2</sub>
4	G <sub>0</sub>
5	G <sub>1</sub>
6	G <sub>2</sub>
7	G <sub>3</sub>
8	T <sub>0</sub>
9	T <sub>1</sub>
10	Т 2
11	Z <sub>0</sub> Z <sub>1</sub>

	Incoming edge(s)				
	Last	Node rank			
	G <sub>0</sub>	1			
	T <sub>0</sub>	2			
	G <sub>1</sub>	3			
	Z <sub>0</sub>	4			
	A 0	5			
· '	T <sub>1</sub>				
	Z <sub>1</sub>	6			
	A 1				
	C 0	7			
	T 2				
	C <sub>1</sub>	8			
	G <sub>2</sub>	9			
	C 2	10			
	G <sub>3</sub>	11			



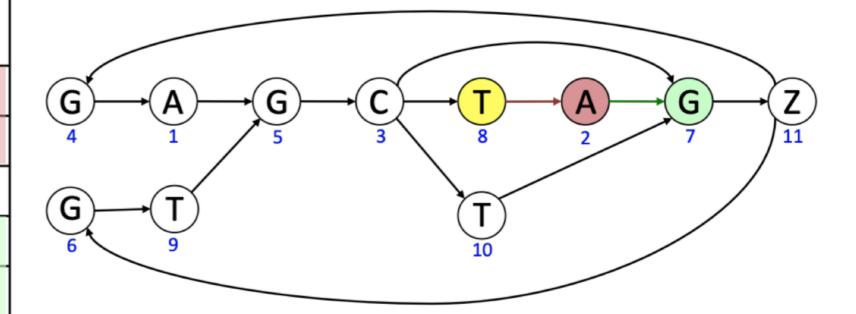
		_	
Outgoing edge(s)			
Node rank	First		
1	A 0		
2	A <sub>1</sub>		
	C <sub>0</sub>		
3	C <sub>1</sub>		
	C <sub>2</sub>		
4	G <sub>0</sub>		
5	G <sub>1</sub>		
6	G <sub>2</sub>	] ;	
7	G <sub>3</sub>		
8	T <sub>0</sub>	];  ;	
9	T <sub>1</sub>		
10	Т2		
11	Z <sub>0</sub> Z <sub>1</sub>		
''	Z <sub>1</sub>		

Incoming edge(s)				
Last	Node rank			
G <sub>0</sub>	1			
T <sub>0</sub>	2			
G <sub>1</sub>	3			
Z <sub>0</sub>	4			
A 0	5			
T <sub>1</sub>				
Z <sub>1</sub>	6			
A 1				
C 0	7			
T 2				
C <sub>1</sub>	8			
G <sub>2</sub>	9			
C 2	10			
G <sub>3</sub>	11			



Outg edg	joing e(s)		Incoming edge(s)		
Node rank	First		Last	Node rank	
1	A 0		G <sub>0</sub>	1	
2	A <sub>1</sub>	] ;	T <sub>0</sub>	2	
	C <sub>0</sub>		G <sub>1</sub>	3	
3	C <sub>1</sub>		Z <sub>0</sub>	4	
	C <sub>2</sub>		A 0	5	
4	G <sub>0</sub>	4	T <sub>1</sub>		
5	G <sub>1</sub>		Z <sub>1</sub>	6	
6	G <sub>2</sub>		A 1		
7	G <sub>3</sub>		C 0	7	
8	T <sub>0</sub>	[; ;	Т 2		
9	T <sub>1</sub>		C <sub>1</sub>	8	
10	T <sub>2</sub>		G <sub>2</sub>	9	
11	Z <sub>0</sub> Z <sub>1</sub>		C 2	10	
11	Z <sub>1</sub>		G <sub>3</sub>	11	

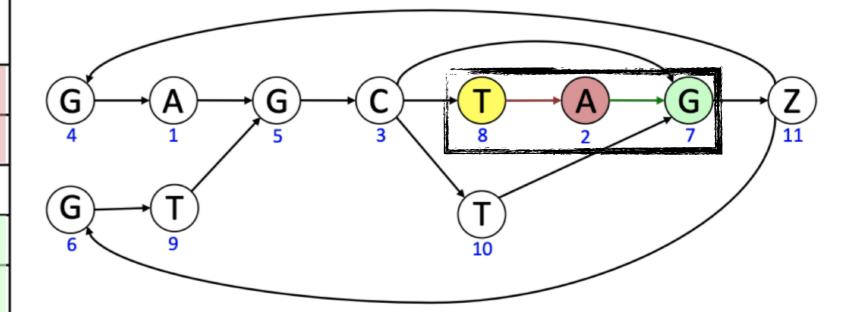
Query: TAG



The last step, to row 8, gives us the **ID** of the node corresponding to the prefix.

	joing e(s)		Incoming edge(s)	
Node rank	First		Last	Node rank
1	A 0		G <sub>0</sub>	1
2	A <sub>1</sub>	] ;	T <sub>0</sub>	2
	C <sub>0</sub>		G <sub>1</sub>	3
3	C 1		Z <sub>0</sub>	4
	C 2		A 0	5
4	G <sub>0</sub>	4	T <sub>1</sub>	
5	G <sub>1</sub>		Z <sub>1</sub>	6
6	G <sub>2</sub>		A 1	
7	G <sub>3</sub>		C 0	7
8	T <sub>0</sub>	<i>:</i>	Т 2	
9	T <sub>1</sub>		C <sub>1</sub>	8
10	T <sub>2</sub>		G <sub>2</sub>	9
11	Z <sub>0</sub> Z <sub>1</sub>		C 2	10
	Z <sub>1</sub>		G <sub>3</sub>	11

Query: TAG



We've found our pattern in the graph!

The last step, to row 8, gives us the **ID** of the node corresponding to the prefix.

### How to store the GFM efficiently

Outg edg		Incor edg	min e(s)
Node rank	First	Last	N ra
1	Α	G	
2	Α	Т	
	С	G	
3	С	Z	
	С	Α	
4	G	Т	
5	G	Z	
6	G	Α	
7	G	С	
8	Т	Т	
9	Т	С	
10	Т	G	
11	Z Z	С	
11	Z	 G	

	Incoming edge(s)				
Last	Node rank				
G	1				
Т	2				
G	3				
Z	4				
Α	5				
Т					
Z	6				
Α					
С	7				
Т					
С	8				
G	9				
С	10				
G	11				

Outgoing edge(s)	1	ming e(s)
Node rank	Last	Node rank
1	10	1
1	11	1
1	10	1
0	00	1
0	00	1
1	11	0
1	00	1
1	00	1
1	01	0
1	11	0
1	01	1
1	10	1
1	01	1
0	10	1

First				
Α	2			
С	3			
G	4			
Т	3			
Z	2			

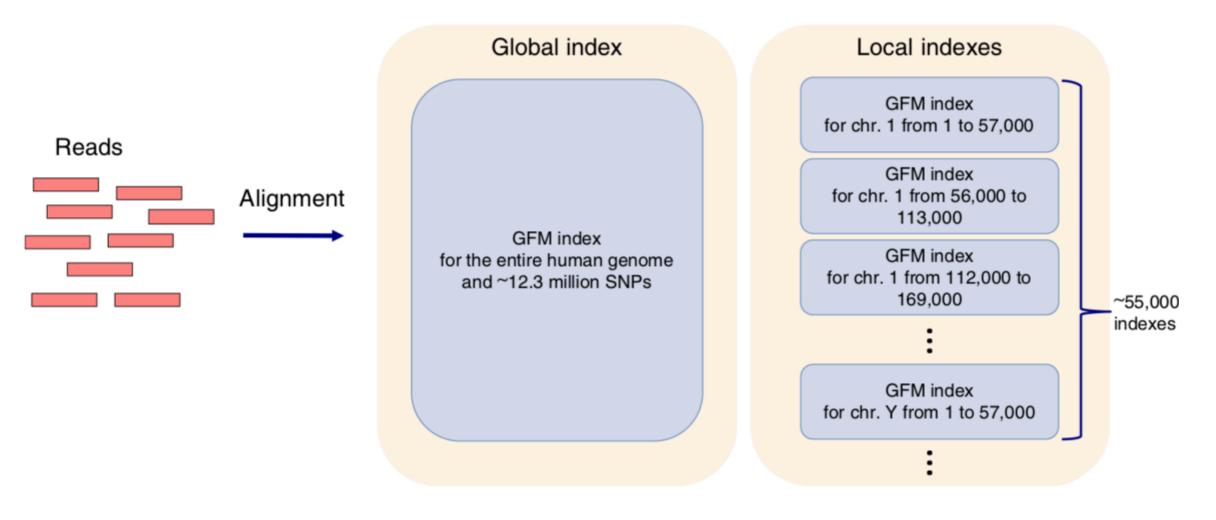
#### How to store the GFM efficiently

1 -	oing e(s)		Incoming edge(s)			Outgoing edge(s)			ming e(s)
Node rank	First		Last	Node rank		Node rank		Last	Node rank
1	Α		G	1		1		10	1
2	^			_ ^		11		11	1
	Think of the GFM index as an index over the edges of the prefix-sorted graph. The tables							1	
3				-	_	-		es	1
				the information associated with the grand ending vertices of the edges, and					1
4	1 16		dges are grouped by their destination					0	
5		respe	ectivel	ctively source) vertices in the two tables.				1	
6	_	The tables are tied together by the same LF						1	
7					a together by used in the li				0
8	l N	ndex		ut WC	uscu III tii <del>e</del> II	iicai tex			0
9	-							1	
10	Т		G	9		1		10	1
11	Z		С	10		1		01	1
11	Z		G	11		0		10	1

First				
A 2				
С	3			
G	4			
Т	3			
Z	2			

# Uses same idea as HISAT to make GFM Cache-efficient

#### 1. HFGM



# Uses same idea as HISAT to make GFM Cache-efficient

#### Special handling of repetitive sequences

