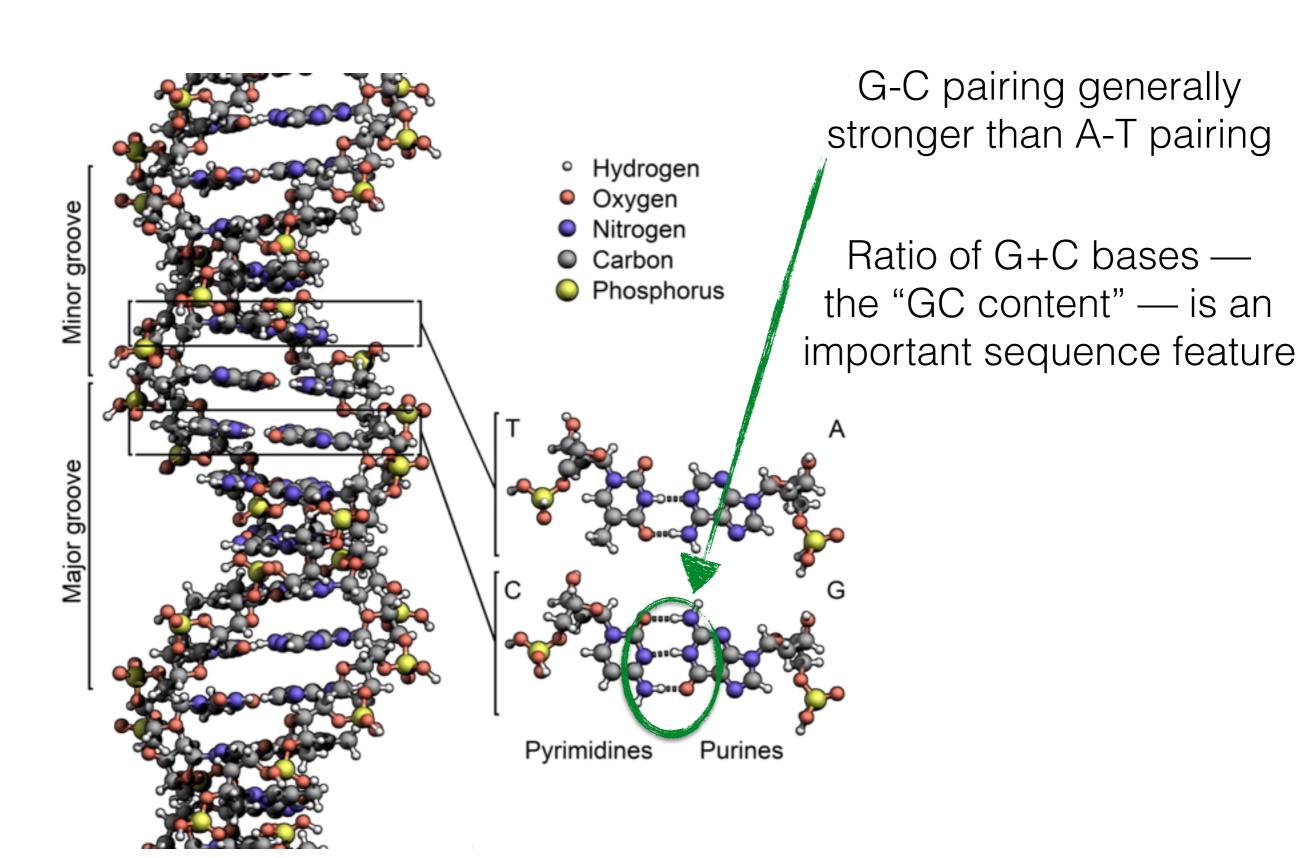
DNA (the genome)



DNA (the genome)

gene — will go on to become a protein

"non-coding DNA" — may or may not produce transcripts (e.g. functional non-coding RNA)

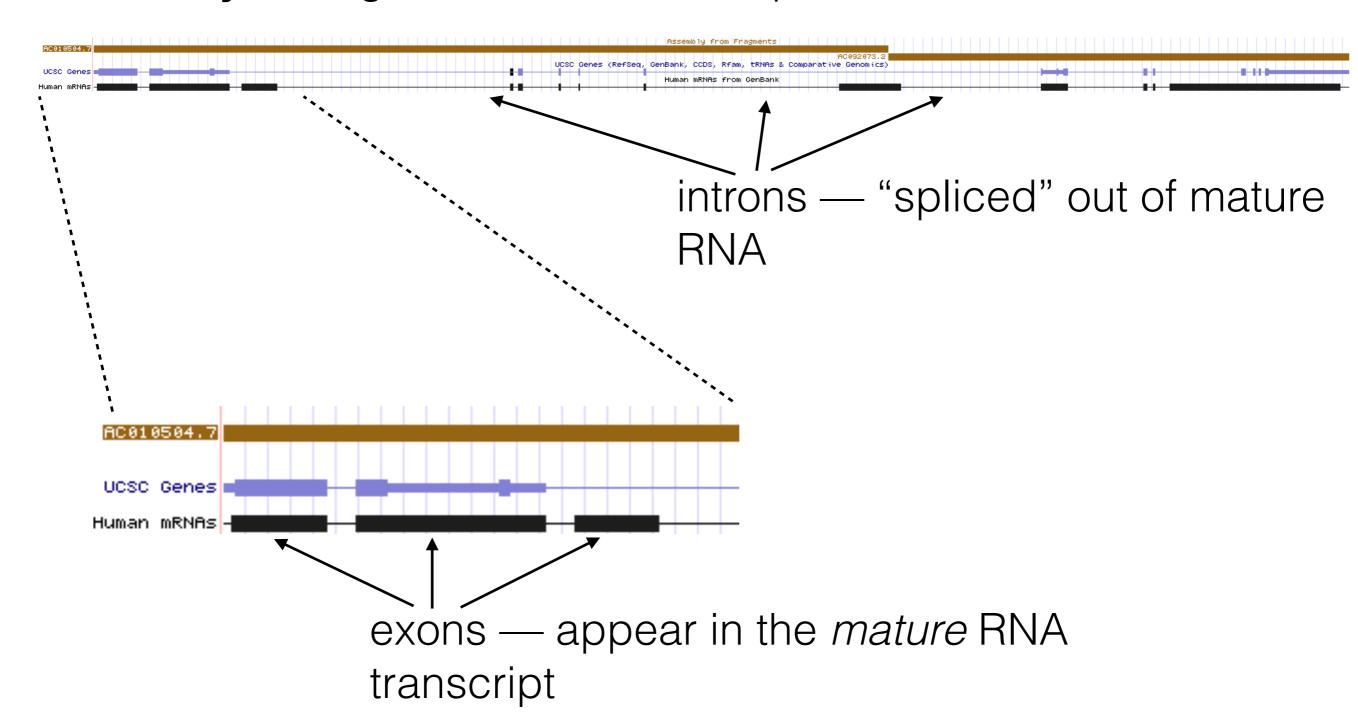
In humans, most DNA is "non-coding" ~98%

In typical bacterial genome, only small fraction — ~2% — of DNA is "non-coding"

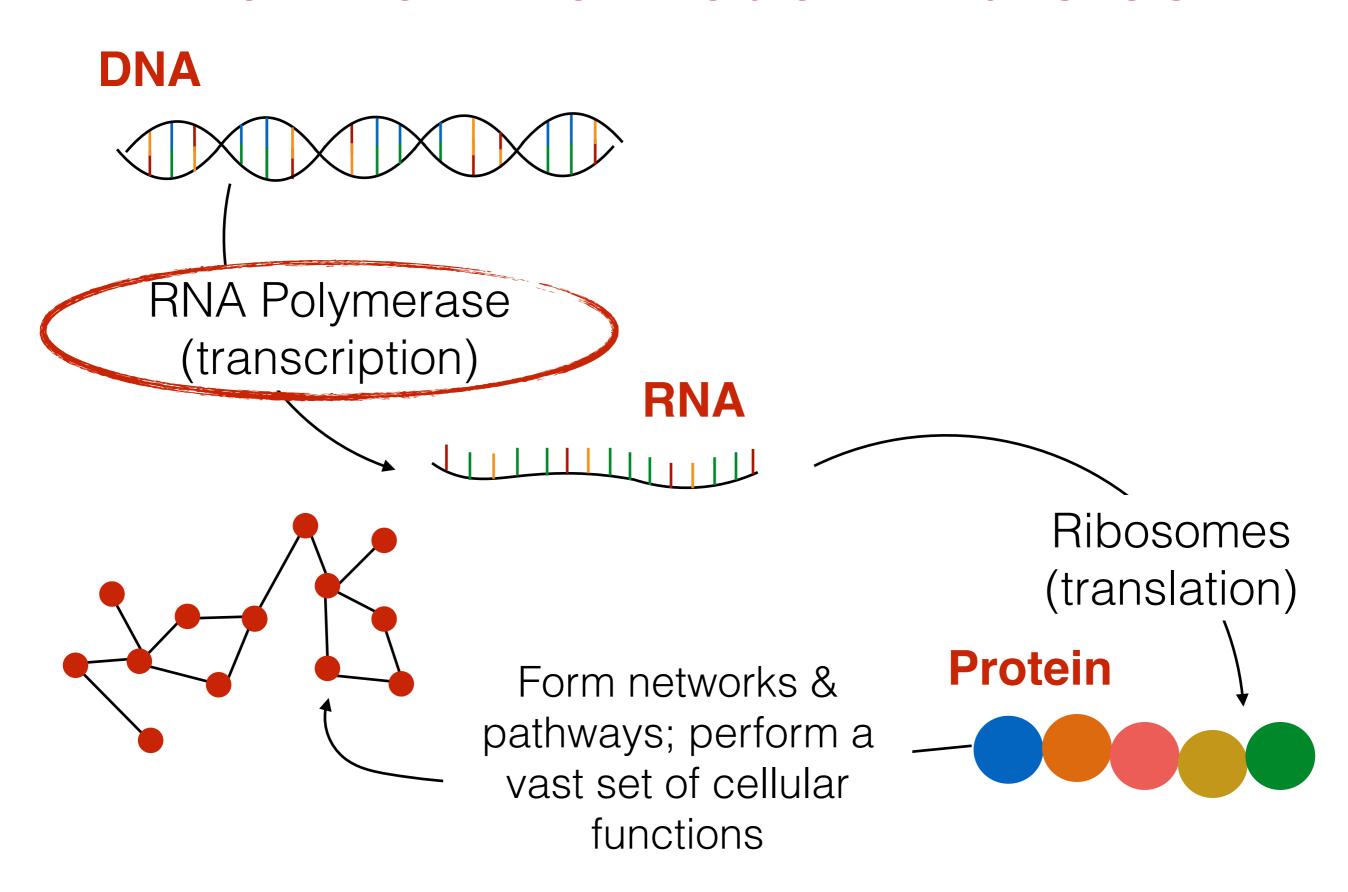
Sometimes referred to as "junk" DNA — much is not, in any way, "junk"

DNA (the genome)

In **prokaryotes**, genes are typically contiguous DNA segment In **eukaryotes**, genes can have complex structure



"Flow" of information in the cell

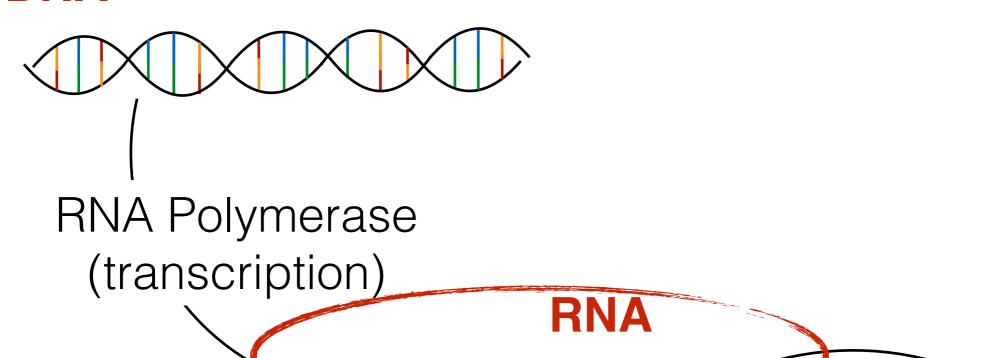


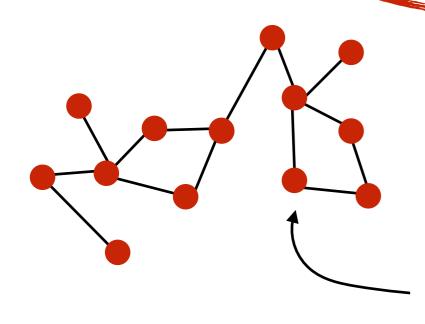
DNA → RNA : Transcription



"Flow" of information in the cell

DNA





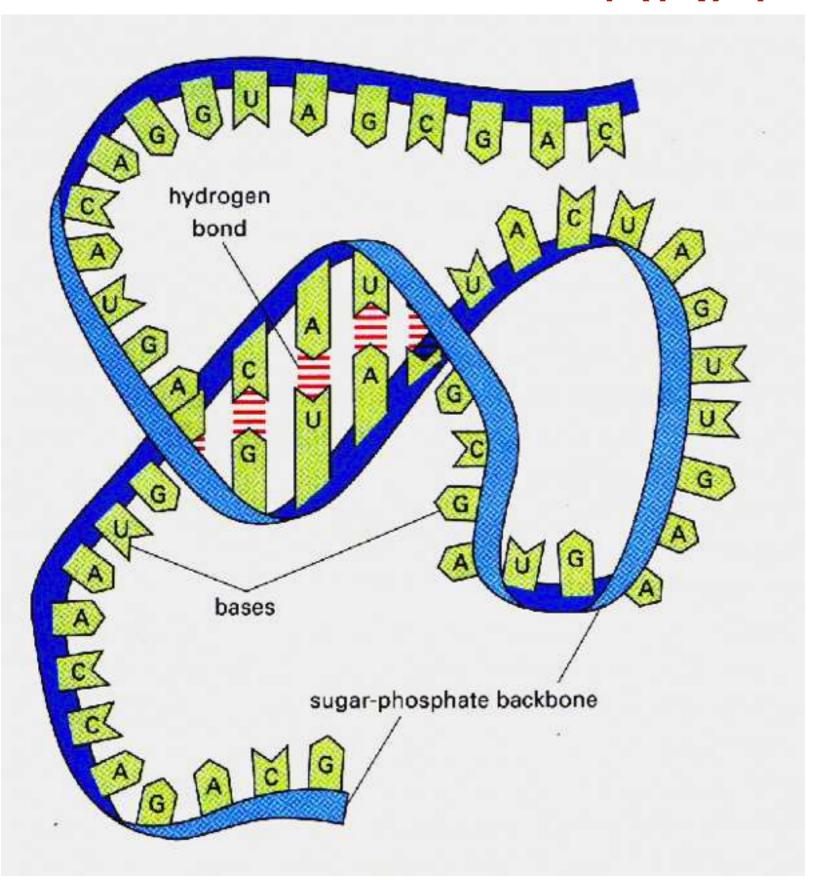
Form networks & pathways; perform a vast set of cellular functions

Ribosomes (translation)





RNA



Less regular structure than DNA

Generally a single-stranded molecule

Secondary & tertiary structure can affect function

Act as transcripts for protein, but also perform important functions themselves

Same "alphabet" as DNA, except thymine replaced by uracil

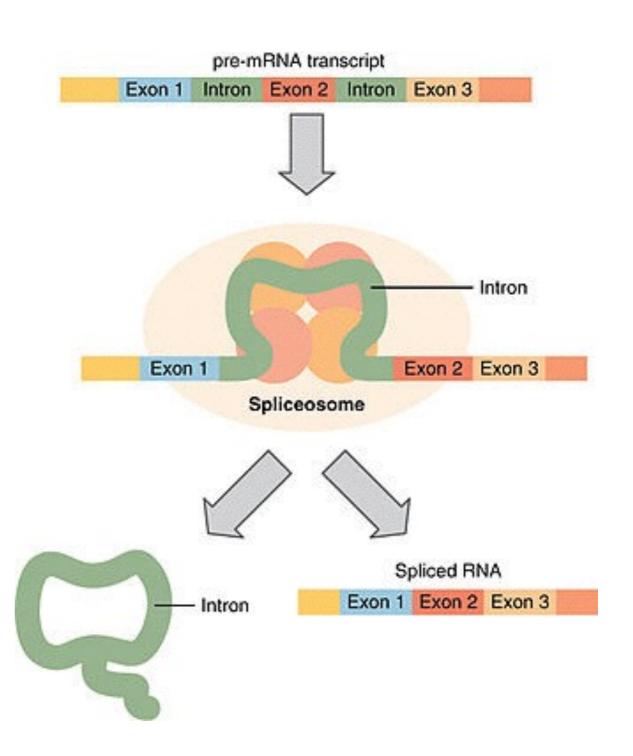
RNA Splicing

DNA transcribed into pre-mRNA

Some "processing occurs" capping & polyadenylation

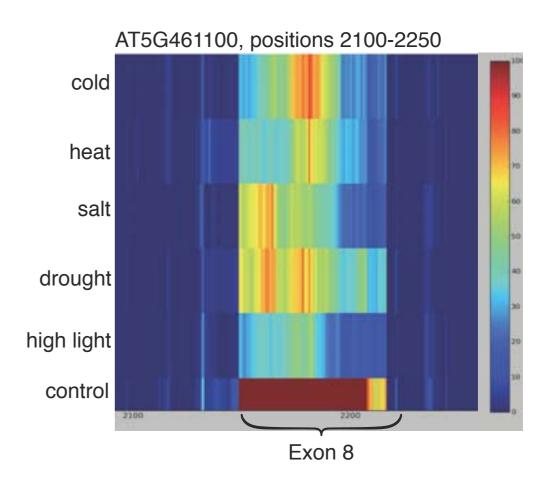
Introns removed from pre-mRNA

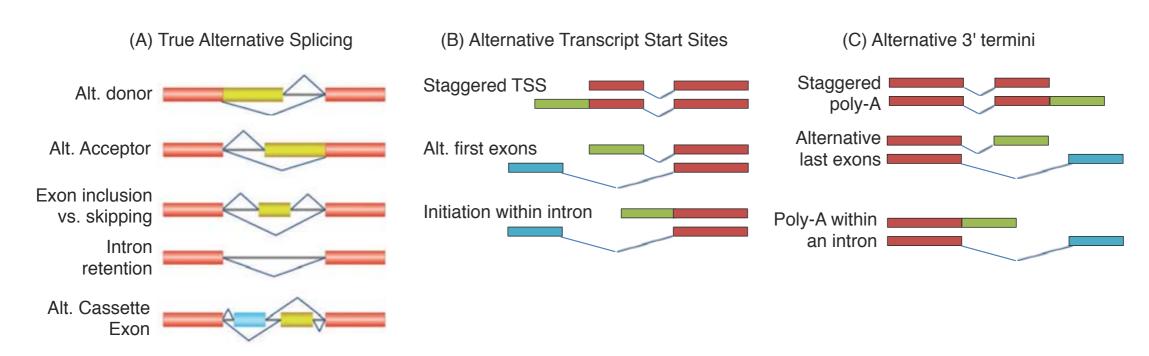
Introns removed resulting in mature mRNA



Alternative Splicing & Isoform Expression

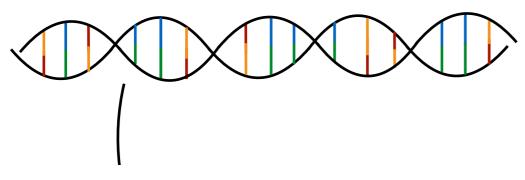
- Expression of genes can be measured via RNA-seq (sequencing transcripts)
- Sequencing gives you short (35-300bp length reads)





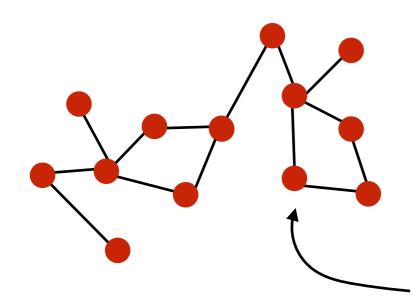
"Flow" of information in the cell

DNA

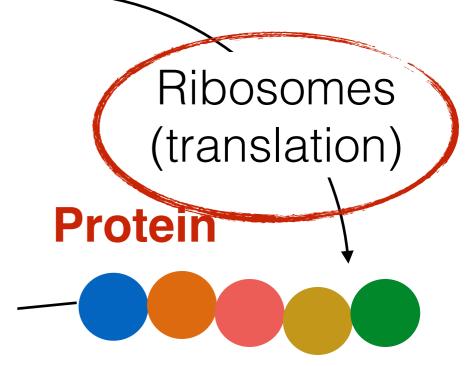


RNA Polymerase (transcription)





Form networks & pathways; perform a vast set of cellular functions

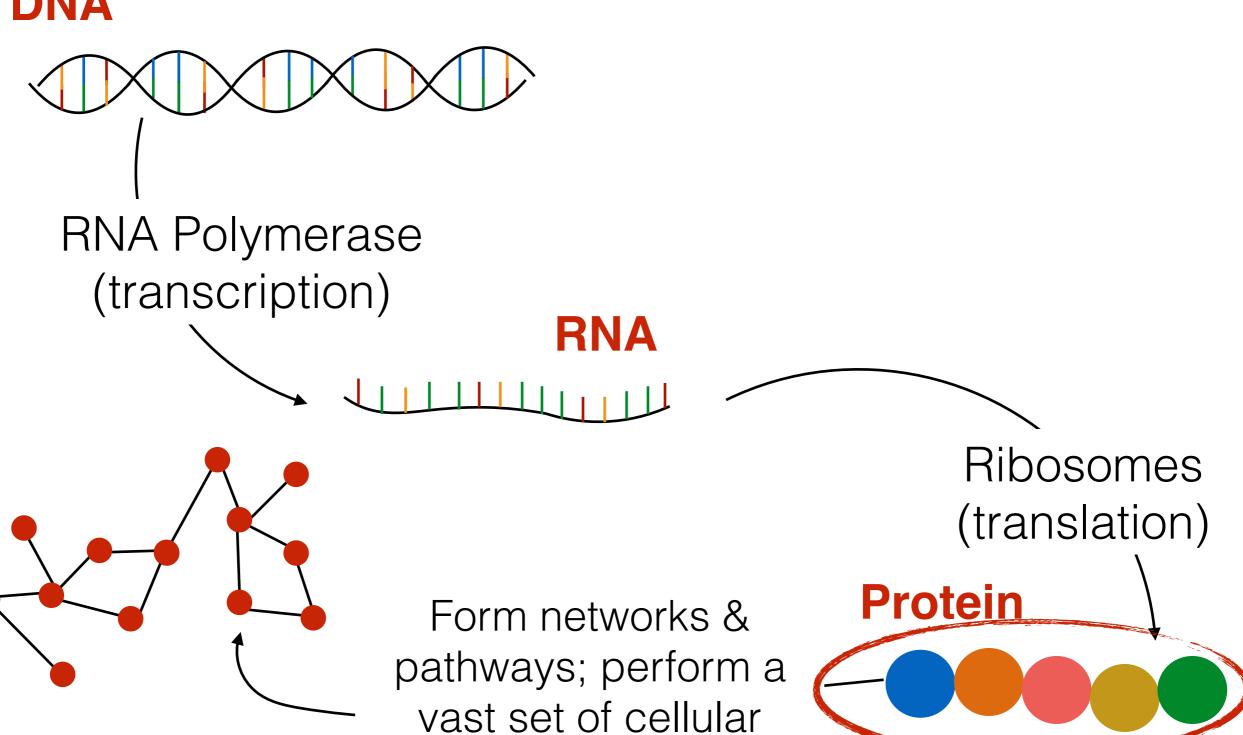


mRNA→ Protein: Translation

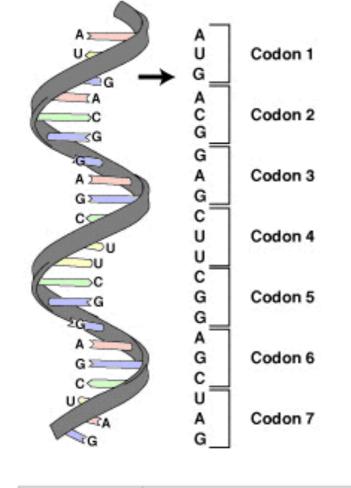


"Flow" of information in the cell

DNA



functions



Protein

Triplets of mRNA bases (codons) correspond to specific amino acids

This mapping is known as the "genetic code" — an *almost* law of molecular Biology

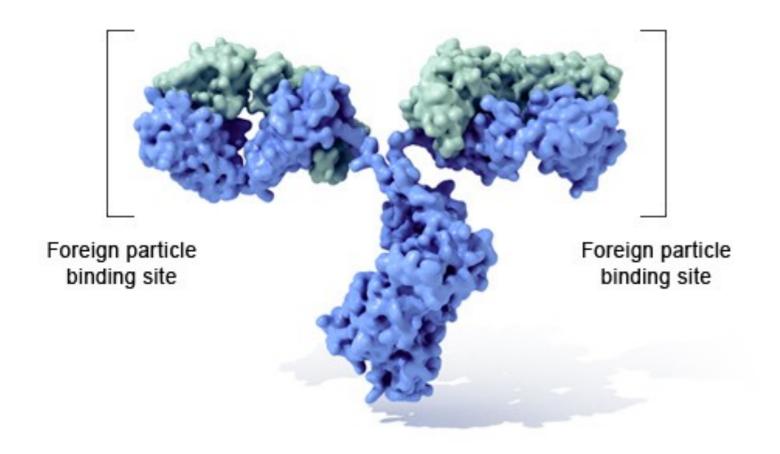
Inverse table (compressed using IUPAC notation)

Amino acid	Codons	Compressed	Amino acid	Codons	Compressed	
Ala/A	GCU, GCC, GCA, GCG	GCN	Leu/L	UUA, UUG, CUU, CUC, CUA, CUG	YUR, CUN	
Arg/R	CGU, CGC, CGA, CGG, AGA, AGG	CGN, MGR	Lys/K	AAA, AAG	AAR	
Asn/N	AAU, AAC	AAY	Met/M	AUG		
Asp/D	GAU, GAC	GAY	Phe/F	UUU, UUC	UUY	
Cys/C	UGU, UGC	UGY	Pro/P	CCU, CCC, CCA, CCG	CCN	
Gln/Q	CAA, CAG	CAR	Ser/S	UCU, UCC, UCA, UCG, AGU, AGC	UCN, AGY	
Glu/E	GAA, GAG	GAR	Thr/T	ACU, ACC, ACA, ACG	ACN	
Gly/G	GGU, GGC, GGA, GGG	GGN	Trp/W	UGG		
His/H	CAU, CAC	CAY	Tyr/Y	UAU, UAC	UAY	
lle/l	AUU, AUC, AUA	AUH	Val/V	GUU, GUC, GUA, GUG	GUN	
START	AUG		STOP	UAA, UGA, UAG	UAR, URA	

en.wikipedia.org: CC BY-SA 3.0

Protein

Immunoglobulin G (IgG)



Perform vast majority of intra & extra cellular functions

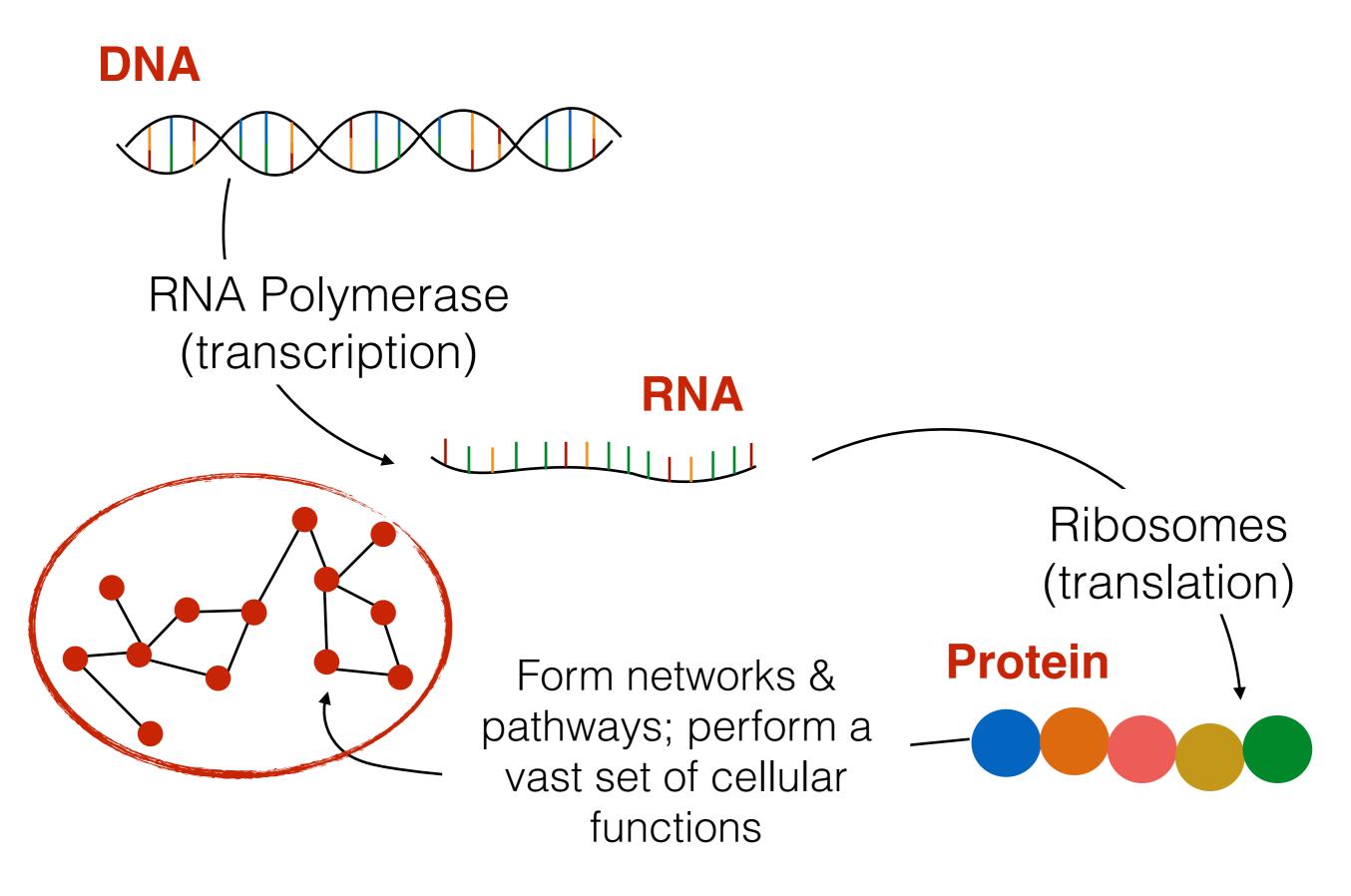
Can range from a few amino acids to *very* large and complex molecules

Can bind with other proteins to form protein complexes

U.S. National Library of Medicine

The shape or *conformation* of a protein is intimately tied to its function. Protein shape, therefore, is strongly conserved through evolution — even moreso than sequence. A protein can undergo sequence mutations, but fold into the same or a similar shape and still perform the same function.

"Flow" of information in the cell



Glycolysis Pathway

Converts glucose → pyruvate phosphoglucose isomerase Generates ATP ("energy currency" of the cell) this is an **example**, no need to memorize this Bio. Pyruvate Phosphoglucose Glucose 6-phosphate Fructose 6-phosphate Phosphofructokinase Phosphoenolpyruvate 2-phosphoglycerate Phosphoglycerate Fructose 1,6-bisphosphate Legend Fructose bisphosphate aldolase Phosphoglycerate 3-phosphoglycerate Hydrogen Glyceraldehyde 3-phosphate Glyceraldehyde phosphate Oxygen dehydrogenase Phosphate group diphosphate Triosephosphate isomerase H. PO, Inorganic phosphate rreversible reaction (highly exergonic) Magnesium ion (cofactor) Reversible reaction Nicotinamide adenine 1,3-bisphosphoglycerate Dihydroxyacetone phosphate

en.wikipedia.org: CC BY-SA 3.0

Hexokinase Enzyme

Some Interesting Facts

Organism	Genome size	# of genes
ф Х174 (<i>E. coli</i> virus)	~5kb	11
E. coli K-12	~4.6Mb	~4,300
Fruit Fly	~122Mb	~17,000
Human	~3.3Gb	~21,000
Mouse	~2.8Gb	~23,000
P. abies (a spruce tree)	~19.6Gb	~28,000

No strong link between genome size & phenotypic complexity Plants can have **huge** genomes (adapt to environment while stationary!)

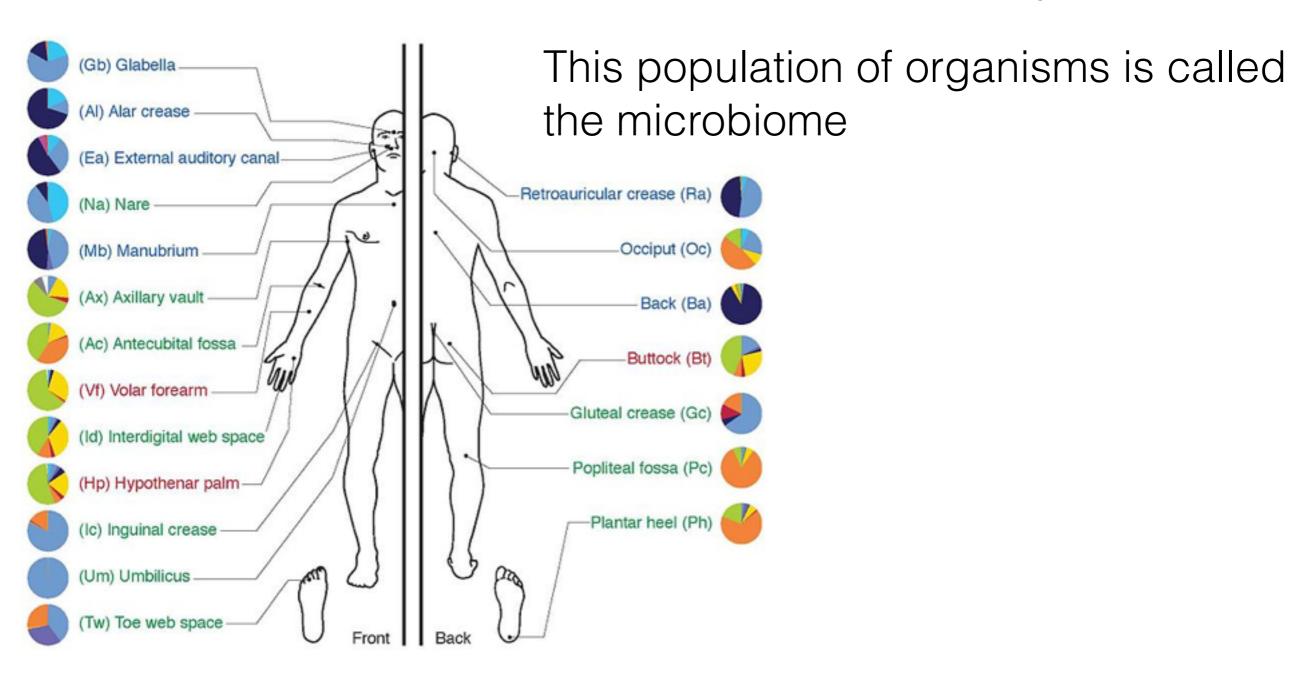
http://users.rcn.com/jkimball.ma.ultranet/BiologyPages/G/GenomeSizes.html

Actinobacteria Corynebacterineae Propionibacterineae Micrococcineae Other Actinobacteria Bacteroidetes Cyanobacteria Firmicutes Other Firmicutes Staphylococcaceae Proteobacteria Divisions contributing < 1% Unclassified

Some Interesting Facts

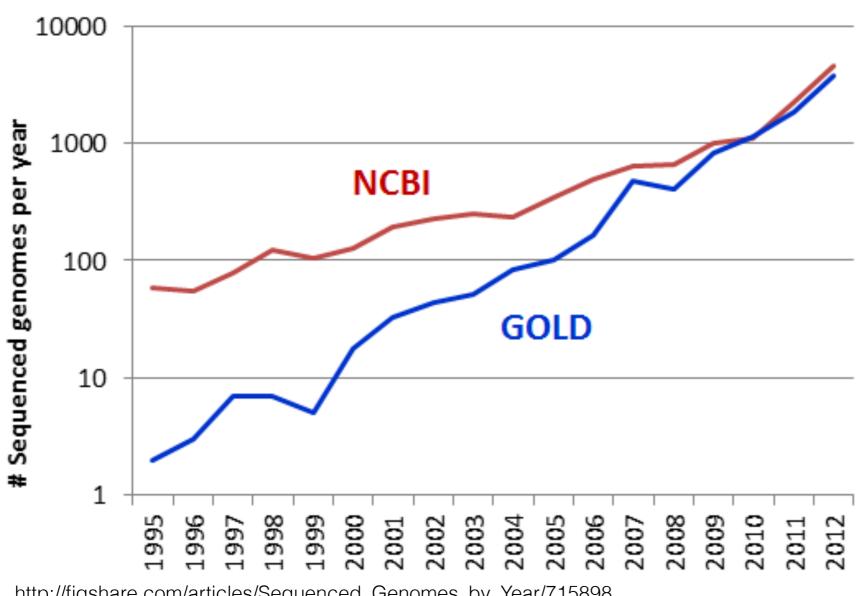
You are mostly bacteria, fungi & arches

Non-human cells outnumber human cells ~10:1 in the human body



en.wikipedia.org: public domain

Some Interesting Facts



http://figshare.com/articles/Sequenced Genomes by Year/715898

... Out of 8.7 \pm 1.3 Mil*

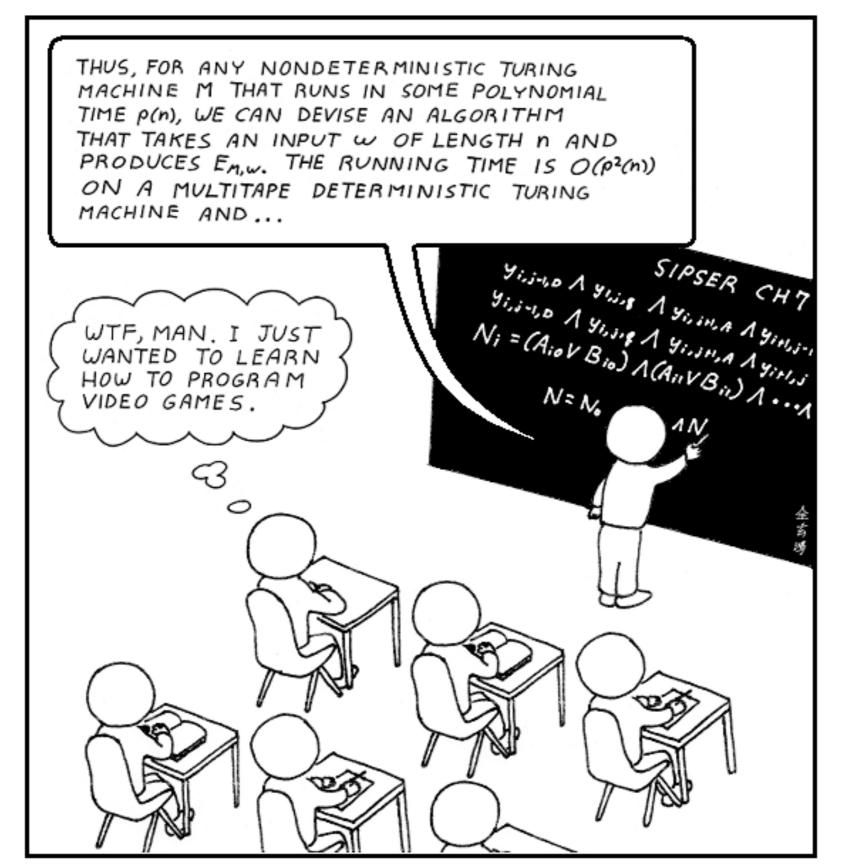
Vast majority of species unsequenced & can not be cultivated in a lab (motivation for metagenomics)

^{*}Mora, Camilo, et al. "How many species are there on Earth and in the ocean?." PLoS biology 9.8 (2011): e1001127.

CSE 549: Computational Biology

Computer Science for Biologists Biology





Not actually simple to define constructively

Still debate whether certain areas constitute CS

Computer science is the scientific and practical approach to computation and its applications. It is the systematic study of the feasibility, structure, expression, and mechanization of the methodical procedures (or algorithms) that underlie the acquisition, representation, processing, storage, communication of, and access to information* ...

What isn't Computer Science?

Don't install operating systems (may develop them)

Don't set up the office network (may study / design network protocols)

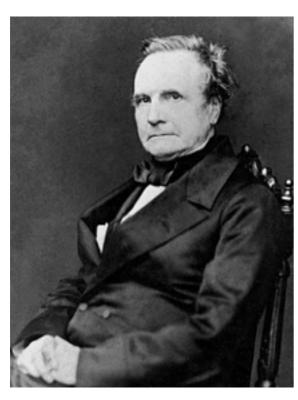
Not about Hacking together a program or learning a web-framework — programming ≠ CS (may study formal languages and develop new programming languages)

^{*}http://www.cs.bu.edu/AboutCS/WhatIsCS.pdf

Started as a branch of Mathematics — early computing machines

Charles Babbage (1791-1871)

Ada Lovelace (1791-1871)



Difference engine → Analytical engine*



Commonly considered the first "programmer"; developed an algorithm for the analytical engine to compute the Bernoulli numbers

*Analytical engine (would have been the first Turing-complete, general purpose computer) was never completed











Kurt Gödel









Kurt Gödel



Alan Turing







Kurt Gödel



Alan Turing



Alonzo Church





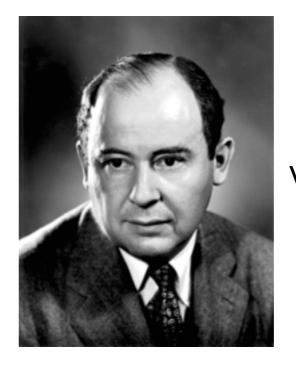
Kurt Gödel



Alan Turing



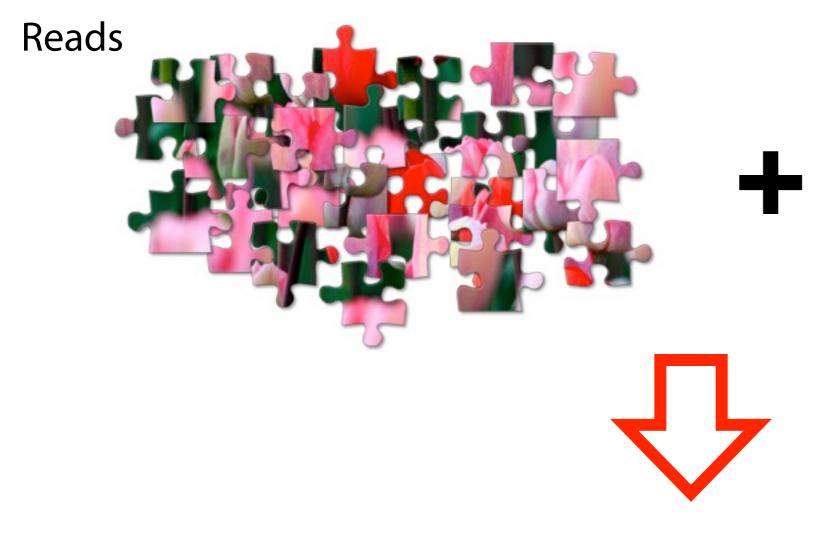
Alonzo Church



John von Neumann

Concerned with the development of provably **correct** and **efficient** computational procedures (algorithms & data structures) to answer **well-specified** problems.

To answer a computational question, we first need a well-formulated problem.



Input DNA



Reference genome.

How to assemble puzzle without the benefit of knowing what the finished product looks like?

Whole-genome "shotgun" sequencing starts by copying and fragmenting the DNA

("Shotgun" refers to the random fragmentation of the whole genome; like it was fired from a shotgun)

Input: GGCGTCTATATCTCGGCTCTAGGCCCCTCATTTTTT

Copy: GGCGTCTATATCTCGGCTCTAGGCCCCTCATTTTTT
GGCGTCTATATCTCGGCTCTAGGCCCCTCATTTTTT
GGCGTCTATATCTCGGCTCTAGGCCCCTCATTTTTT
GGCGTCTATATCTCGGCTCTAGGCCCCTCATTTTTT

Fragment: GGCGTCTA TATCTCGG CTCTAGGCCCTC ATTTTTT
GGC GTCTATAT CTCGGCTCTAGGCCCTCA TTTTTT
GGCGTC TATATCT CGGCTCTAGGCCCT CATTTTTT
GGCGTCTAT ATCTCGGCTCTAG GCCCTCA TTTTTT

Assume sequencing produces such a large # fragments that almost all genome positions are *covered* by many fragments...

Reconstruct this

CTAGGCCCTCAATTTTT
CTCTAGGCCCTCAATTTTT
GGCTCTAGGCCCCTCATTTTTT
CTCGGCTCTAGCCCCCTCATTTTT
TATCTCGACTCTAGGCCCCTCA
TATCTCGACTCTAGGCC
TCTATATCTCGGCTCTAGG
GGCGTCTATATCTCG
GGCGTCGATATCT
GGCGTCTATATCT

From these

→ GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT

...but we don't know what came from where

Reconstruct this

CTAGGCCCTCAATTTTT
GGCGTCTATATCT
CTCTAGGCCCTCAATTTTT
TCTATATCTCGGCTCTAGG
GGCTCTAGGCCCTCATTTTT
CTCGGCTCTAGCCCCTCATTTTT
TATCTCGACTCTAGGCCCTCA
GGCGTCGATATCT
TATCTCGACTCTAGGCC
GGCGTCTATATCTCG

From these

→ GGCGTCTATATCTCGGCTCTAGGCCCTCATTTTTT

Concerned with the development of provably **correct** and **efficient** computational procedures (algorithms & data structures) to answer **well-specified** problems.

To answer a computational question, we first need a well-formulated problem.

Given: a collection, R, of sequencing reads (strings)

Find: The genome (string), G, that generated them

Concerned with the development of provably **correct** and **efficient** computational procedures (algorithms & data structures) to answer **well-specified** problems.

To answer a computational question, we first need a well-formulated problem.

Given: a collection R, of sequencing reads (strings)

Find: The genome (string), G, that generated them

Not well-specified.

What makes one genome more likely than another? What constraints do we place on the space of solutions?

Concerned with the development of provably **correct** and **efficient** computational procedures (algorithms & data structures) to answer **well-specified** problems.

To answer a computational question, we first need a well-formulated problem.

Given: a collection, R, of sequencing reads (strings)

าร

Find: The shortest genome (string), G, that contains all of them

Shortest Common Superstring

Given: a collection, $S = \{s_1, s_2, \dots, s_k\}$, of sequencing reads (strings)

Find*: The shortest possible genome (string), G, such that s_1, s_2, \ldots, s_k are all substrings of G

How, might we go about solving this problem?

^{*}for reasons we'll explore later, this isn't actually a great formulation for genome assembly.

Shortest common superstring

Given a collection of strings *S*, find *SCS*(*S*): the shortest string that contains all strings in *S* as substrings

Without requirement of "shortest," it's easy: just concatenate them

Example: S: BAA AAB BBA ABA ABB BBB AAA BAB

SCS(S): AAABBBABAA

─10 **─**

AAA
ABB
ABB
BBA
BAB
ABA
BAA

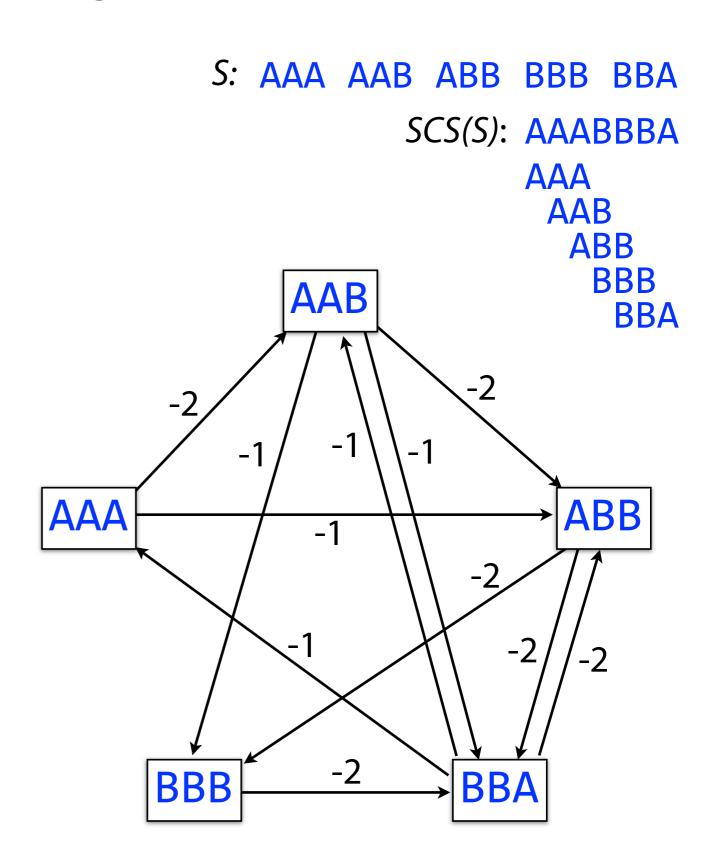
Shortest common superstring

Can we solve it?

Imagine a modified overlap graph where each edge has cost = - (length of overlap)

SCS corresponds to a path that visits every node once, minimizing total cost along path

That's the *Traveling Salesman Problem (TSP)*, which is NP-hard!

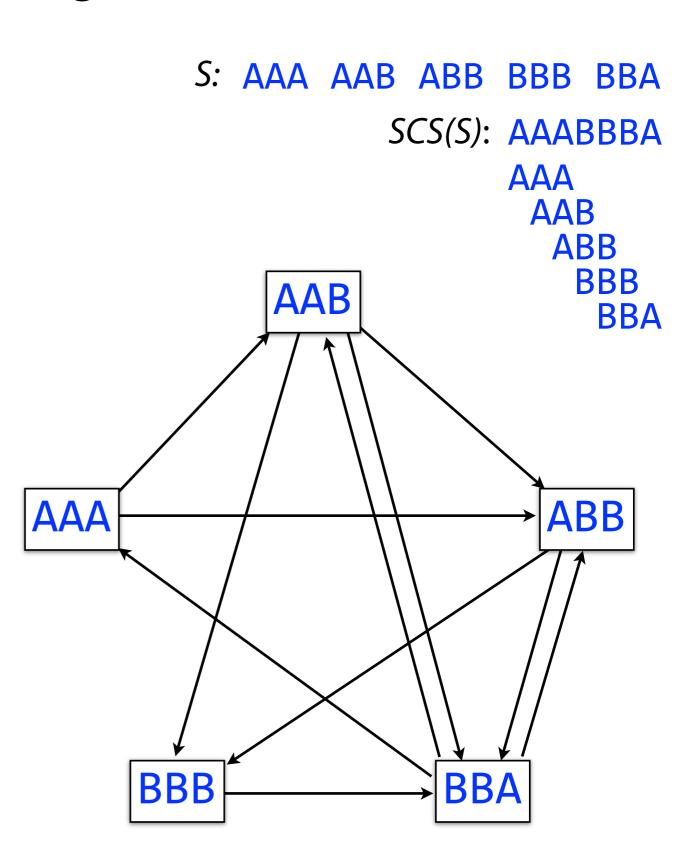


Shortest common superstring

Say we disregard edge weights and just look for a path that visits all the nodes exactly once

That's the *Hamiltonian Path* problem: NP-complete

Indeed, it's well established that SCS is NP-hard



Shortest common superstring & friends

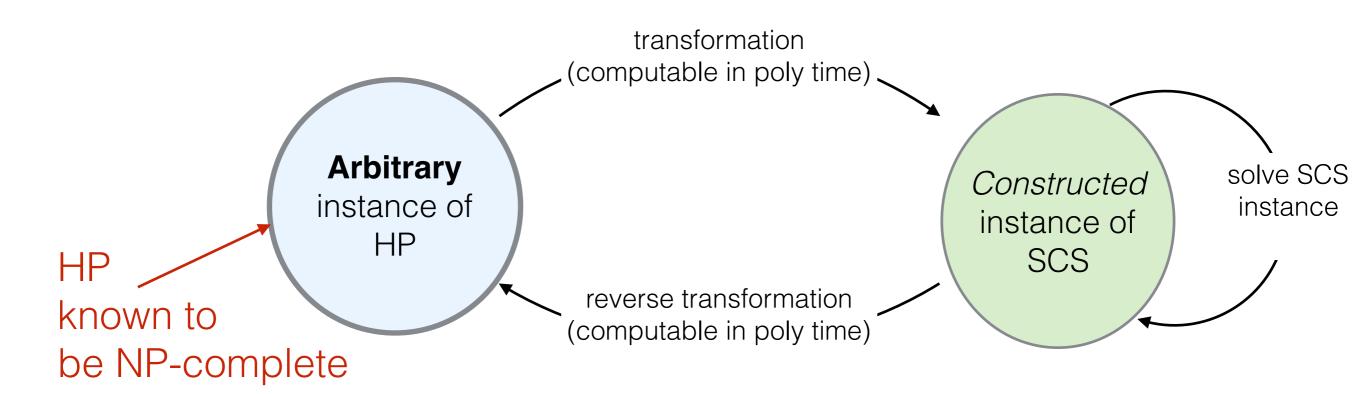
Traveling Salesman, Hamiltonian Path, and Shortest Common Superstring are all NP-hard

For refreshers on Traveling Salesman, Hamiltonian Path, NP-hardness and NP-completeness, see Chapters 34 and 35 of "Introduction to Algorithms" by Cormen, Leiserson, Rivest and Stein, or Chapters 8 and 9 of "Algorithms" by Dasgupta, Papadimitriou and Vazirani (free online: http://www.cs.berkeley.edu/~vazirani/algorithms)

Important note: The fact that we modeled SCS as NP-hard problems (TSP and HP) **does not** prove that (the decision version of) SCS is NP-complete. To do that, we must **reduce** a known NP-complete problem to **SCS**.

Given an instance I of a known hard problem, generate an instance I' of SCS such that if we can solve I' in polynomial time, then we can solve I in polynomial time. This *implies* that SCS is *at least* as hard as the hard problem.

This can be done e.g. with HAMILTONIAN PATH



Shortest Common Superstring

The fact that SCS is **NP-complete** means that it is unlikely that there exists *any* algorithm that can solve a general instance of this problem in time polynomial in n — the number of strings.

If we give up on finding the *shortest* possible superstring G, how does the situation change?

Shortest Common Superstring

There's a "greedy" *heuristic* that turns out to be an *approximation algorithm* (provides a solution within a constant factor of the the optimum)

At each step, chose the pair of strings with the maximum overlap, merge them, and return the merged string to the collection.

Greedy conjecture factor of 2-OPT *is* the worst case — proof for factor 3.5

Different approx. (not all greedy)

ratio	authors	year				
approximating SCS						
3	Blum, Jiang, Li, Tromp and Yannakakis [4]	1991				
$2\frac{8}{9}$	Teng, Yao [23]	1993				
$2\frac{5}{6}$	Czumaj, Gasieniec, Piotrow, Rytter [8]	1994				
$2\frac{50}{63}$	Kosaraju, Park, Stein [15]	1994				
$2\frac{3}{4}$	Armen, Stein [1]	1994				
$2\frac{50}{69}$	Armen, Stein [2]	1995				
$2\frac{2}{3}$	Armen, Stein [3]	1996				
$2\frac{25}{42}$	Breslauer, Jiang, Jiang [5]	1997				
$2\frac{1}{2}$	Sweedyk [21]	1999				
$2\frac{1}{2}$	Kaplan, Lewenstein, Shafrir, Sviridenko [12]	2005				
$2\frac{1}{2}$	Paluch, Elbassioni, van Zuylen [18]	2012				
$2\frac{11}{23}$	Mucha [16]	2013				