

Topos Theory I: Presheaf Categories and Related Constructions

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Contents

Some text.

- Presheaf Categories

Contents

Some text.

- Presheaf Categories
- Pullbacks

Contents

Some text.

- Presheaf Categories
- Pullbacks
- Subobject Classifiers

Contents

Some text.

- Presheaf Categories
- Pullbacks
- Subobject Classifiers
- Colimits

Contents

- Functor Categories

Contents

- Functor Categories
- Pullbacks

Contents

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- Pullbacks
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- Pullbacks
- Subobject Classifiers
- Colimits

Functor Categories at a Distance

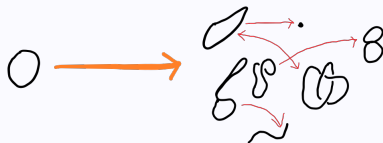
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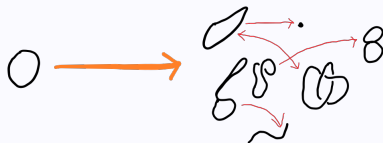
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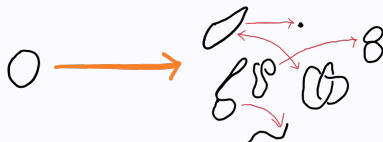
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- *Analogously...* a functor $F : \mathbf{J} \rightarrow \mathbf{C}$ is a “drawing” of \mathbf{J} in \mathbf{C} .
- A functor category $\mathbf{C}^{\mathbf{J}}$ considers all such “drawings” along with compatible maps between them.

Category of Presheaves

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- ▶ $\mathbf{Set}^{\mathbf{N}} = \hat{\mathbf{N}^{\text{op}}}$, where \mathbf{N} is the total order category of \mathbb{N} .

$$\begin{array}{ccc} p_n & \longrightarrow & p_m \\ \downarrow & & \downarrow \\ q_n & \longrightarrow & q_m \end{array} \quad \begin{array}{cccc} p_0 & \longrightarrow & p_1 & \longrightarrow \dots \\ \downarrow & & \downarrow & \\ q_0 & \longrightarrow & q_1 & \longrightarrow \dots \end{array}$$

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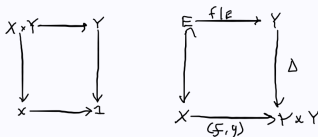
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$$(X \times_B Y)(C) \cong X(C) \times_{B(C)} Y(C)$$
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- All finite limits can be constructed from equalizers and products:



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- If we define $\text{true} : 1 \hookrightarrow 2$, then S can be recovered from ϕ_S by pulling back along true .

$$\begin{array}{ccc} X \times 1 & \longrightarrow & 1 \\ \downarrow & & \downarrow \\ X & \xrightarrow{\phi_S} & 2 \end{array}$$

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- *Idea:* generalize the set 2 to an object $\Omega \in \mathbf{C}$ of “truth values” for a general category \mathbf{C} .

Subobject Classifiers

Definition

In a category \mathbf{C} with finite limits, a **subobject classifier** is a monomorphism $true : 1 \rightarrow \Omega$ such that for every monomorphism $S \rightarrowtail X$, we get a unique morphism ϕ forming a pullback square

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 - ▶ $\text{Sub}_{\mathbf{C}}(X)$ is the set of subobjects $S \rightarrowtail X$.
 - ▶ For $f : S' \rightarrow S$, $\text{Sub}_{\mathbf{C}}(f)(S \rightarrowtail X)$ is the (monic) pullback $S' \rightarrowtail X'$ of $S \rightarrowtail X$ along f .

Subobject Classifiers

Proposition

A locally small category \mathbf{C} with finite limits has a subobject classifier iff $\text{Sub}_{\mathbf{C}}$ is representable: there exists an object Ω and a natural isomorphism:

$$\theta_X : \text{Sub}_{\mathbf{C}}(X) \cong \text{Hom}_{\mathbf{C}}(X, \Omega).$$

If this holds, then \mathbf{C} is well-powered: $\text{Sub}_{\mathbf{C}}(X)$ is isomorphic to a small set for all X .