Calculus: Early Transcendentals, 9th Ed. Diagnostic A: Algebra

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Evaluate each expression without using a calculator.

1.
$$(-3)^4 = 81$$

$$2. -3^4 = -81$$

$$3. \ 3^{-4} = 1/3^4 = 1/81$$

4.
$$5^{23}/5^{21} = 5^{23-21} = 5^2 = 25$$

5.
$$(2/3)^{-2} = (2/3)^{-1*2} = (3/2)^2 = 9/4$$

6.
$$16^{-3/4} = (16^{\frac{3}{4}})^{-1} = (\frac{1}{16^{\frac{3}{4}}}) = \frac{1}{(\sqrt[4]{16})^3} = \frac{1}{2^3} = \frac{1}{8}$$

Problem 2

Simplify each expression. Write your answer without negative exponents.

$$1. \ \sqrt{200} - \sqrt{32} = \sqrt{100*2} - \sqrt{16*2} = 100*\sqrt{100*2} - 8*\sqrt{16*2} = 10*\sqrt{2} - 4*sqrt2 = 6*\sqrt{2}$$

- (a) note that 32 is a power of 2; we can get a clean factor
- (b) since 100 * 2 is also a perfect square factor of 200
- (c) many factors of 200: 1, 2, 4, 5, 10, 20, 40, 50, 100, 200

2.
$$(3a^3b^3)(4ab^2)^2 = 3*4*4*a^3*a*a*b^3*b^2*b^2 = 48a^5b^7$$

3.
$$\left(\frac{3x^{3/2}y^3}{x^2y^{\frac{-1}{2}}}\right)^{-2} = \frac{(x^2y^{\frac{-1}{2}})^2}{(3x^{3/2}y^3)^2} = \frac{x^4y^{-1}}{9x^3y^6} = \frac{x}{9y^7}$$

Problem 3

Expand and simplify

1.
$$3(x+6) + 4(2x-5) = 3x + 3 * 6 + 8x - 20 = 11x - 2$$

2.
$$(\sqrt{a} + \sqrt{b})(\sqrt{a} - \sqrt{b}) = a - b$$

3.
$$(x+3)(4x-5) = 4x^2 + 7x - 15$$

4.
$$(2x+3)^2 = 4x^2 + 12x + 9$$

5.
$$(x+2)^3 == (x+2)(x+2)(x+2) = x^3 + 6x^2 + 12x + 8$$

Problem 4

Factor each expression.

1.
$$4x^2 - 25 = (2x - 5)(2x + 5)$$

2.
$$2x^2 + 5x - 12 = (2x - 3)(1x + 4)$$
 - solved by factor and brute force

(a) System of equations... hmm - how can we solve below?

(b)
$$a, b, c, d : b * d = -12, a * c = 2, a * d + b * c = 5$$

3.
$$x^3 - 3x^2 - 4x + 12$$

Give an appropriate positive constant c such that $f(n) \leq c \cdot g(n)$ for all n > 1.

1.
$$f(n) = n^2 + n + 1$$
, $g(n) = 2n^3$

2.
$$f(n) = n\sqrt{n} + n^2$$
, $g(n) = n^2$

3.
$$f(n) = n^2 - n + 1$$
, $g(n) = n^2/2$

Solution

We solve each solution algebraically to determine a possible constant c.

Part One

$$n^{2} + n + 1 =$$

$$\leq n^{2} + n^{2} + n^{2}$$

$$= 3n^{2}$$

$$\leq c \cdot 2n^{3}$$

Thus a valid c could be when c = 2.

Part Two

$$n^{2} + n\sqrt{n} =$$

$$= n^{2} + n^{3/2}$$

$$\leq n^{2} + n^{4/2}$$

$$= n^{2} + n^{2}$$

$$= 2n^{2}$$

$$\leq c \cdot n^{2}$$

Thus a valid c is c = 2.

Part Three

$$n^{2} - n + 1 =$$

$$\leq n^{2}$$

$$\leq c \cdot n^{2}/2$$

Thus a valid c is c = 2.

Let $\Sigma = \{0,1\}$. Construct a DFA A that recognizes the language that consists of all binary numbers that can be divided by 5.

Let the state q_k indicate the remainder of k divided by 5. For example, the remainder of 2 would correlate to state q_2 because 7 mod 5 = 2.

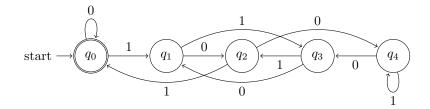


Figure 1: DFA, A, this is really beautiful, ya know?

Justification

Take a given binary number, x. Since there are only two inputs to our state machine, x can either become x0 or x1. When a 0 comes into the state machine, it is the same as taking the binary number and multiplying it by two. When a 1 comes into the machine, it is the same as multiplying by two and adding one.

Using this knowledge, we can construct a transition table that tell us where to go:

	$x \mod 5 = 0$	$x \mod 5 = 1$	$x \mod 5 = 2$	$x \mod 5 = 3$	$x \mod 5 = 4$
x0	0	2	4	1	3
x1	1	3	0	2	4

Therefore on state q_0 or $(x \mod 5 = 0)$, a transition line should go to state q_0 for the input 0 and a line should go to state q_1 for input 1. Continuing this gives us the Figure 1.

Problem 7

Write part of Quick-Sort(list, start, end)

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1: function Quick-Sort(list, start, end)
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- 2: **if** $start \ge end$ **then**
- 3: return
- 4: end if
- 5: $mid \leftarrow Partition(list, start, end)$
- 6: Quick-Sort(list, start, mid 1)
- 7: QUICK-SORT(list, mid + 1, end)
- 8: end function

Algorithm 1: Start of QuickSort

Suppose we would like to fit a straight line through the origin, i.e., $Y_i = \beta_1 x_i + e_i$ with i = 1, ..., n, $E[e_i] = 0$, and $Var[e_i] = \sigma_e^2$ and $Cov[e_i, e_j] = 0$, $\forall i \neq j$.

Part A

Find the least squares esimator for $\hat{\beta}_1$ for the slope β_1 .

Solution

To find the least squares estimator, we should minimize our Residual Sum of Squares, RSS:

$$RSS = \sum_{i=1}^{n} (Y_i - \hat{Y}_i)^2$$
$$= \sum_{i=1}^{n} (Y_i - \hat{\beta}_1 x_i)^2$$

By taking the partial derivative in respect to $\hat{\beta}_1$, we get:

$$\frac{\partial}{\partial \hat{\beta}_1}(RSS) = -2\sum_{i=1}^n x_i(Y_i - \hat{\beta}_1 x_i) = 0$$

This gives us:

$$\sum_{i=1}^{n} x_i (Y_i - \hat{\beta}_1 x_i) = \sum_{i=1}^{n} x_i Y_i - \sum_{i=1}^{n} \hat{\beta}_1 x_i^2$$
$$= \sum_{i=1}^{n} x_i Y_i - \hat{\beta}_1 \sum_{i=1}^{n} x_i^2$$

Solving for $\hat{\beta}_1$ gives the final estimator for β_1 :

$$\hat{\beta_1} = \frac{\sum x_i Y_i}{\sum x_i^2}$$

Part B

Calculate the bias and the variance for the estimated slope $\hat{\beta}_1$.

Solution

For the bias, we need to calculate the expected value $E[\hat{\beta}_1]$:

$$\begin{aligned} \mathbf{E}[\hat{\beta}_1] &= \mathbf{E}\left[\frac{\sum x_i Y_i}{\sum x_i^2}\right] \\ &= \frac{\sum x_i \mathbf{E}[Y_i]}{\sum x_i^2} \\ &= \frac{\sum x_i (\beta_1 x_i)}{\sum x_i^2} \\ &= \frac{\sum x_i^2 \beta_1}{\sum x_i^2} \\ &= \beta_1 \frac{\sum x_i^2 \beta_1}{\sum x_i^2} \\ &= \beta_1 \end{aligned}$$

Thus since our estimator's expected value is β_1 , we can conclude that the bias of our estimator is 0.

For the variance:

$$\operatorname{Var}[\hat{\beta_1}] = \operatorname{Var}\left[\frac{\sum x_i Y_i}{\sum x_i^2}\right]$$

$$= \frac{\sum x_i^2}{\sum x_i^2 \sum x_i^2} \operatorname{Var}[Y_i]$$

$$= \frac{\sum x_i^2}{\sum x_i^2 \sum x_i^2} \operatorname{Var}[Y_i]$$

$$= \frac{1}{\sum x_i^2} \operatorname{Var}[Y_i]$$

$$= \frac{1}{\sum x_i^2} \sigma^2$$

$$= \frac{\sigma^2}{\sum x_i^2}$$

Problem 9

Prove a polynomial of degree k, $a_k n^k + a_{k-1} n^{k-1} + \ldots + a_1 n^1 + a_0 n^0$ is a member of $\Theta(n^k)$ where $a_k \ldots a_0$ are nonnegative constants.

Proof. To prove that $a_k n^k + a_{k-1} n^{k-1} + \ldots + a_1 n^1 + a_0 n^0$, we must show the following:

$$\exists c_1 \exists c_2 \forall n \ge n_0, \ c_1 \cdot g(n) \le f(n) \le c_2 \cdot g(n)$$

For the first inequality, it is easy to see that it holds because no matter what the constants are, $n^k \le a_k n^k + a_{k-1} n^{k-1} + \ldots + a_1 n^1 + a_0 n^0$ even if $c_1 = 1$ and $n_0 = 1$. This is because $n^k \le c_1 \cdot a_k n^k$ for any nonnegative constant, c_1 and a_k .

Taking the second inequality, we prove it in the following way. By summation, $\sum_{i=0}^{k} a_i$ will give us a new constant, A. By taking this value of A, we can then do the following:

$$a_k n^k + a_{k-1} n^{k-1} + \ldots + a_1 n^1 + a_0 n^0 =$$

$$\leq (a_k + a_{k-1} \ldots a_1 + a_0) \cdot n^k$$

$$= A \cdot n^k$$

$$\leq c_2 \cdot n^k$$

where $n_0 = 1$ and $c_2 = A$. c_2 is just a constant. Thus the proof is complete.

Evaluate $\sum_{k=1}^{5} k^2$ and $\sum_{k=1}^{5} (k-1)^2$.

Problem 19

Find the derivative of $f(x) = x^4 + 3x^2 - 2$

Problem 6

Evaluate the integrals $\int_0^1 (1-x^2) \mathrm{d}x$ and $\int_1^\infty \frac{1}{x^2} \mathrm{d}x$.