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1 Basic
1.1 .vimrc
svn on
se ai nu rnu ru cul mouse=a
se cin et ts=2 sw=2 sts=2
so $VIMRUNTIME/mswin.vim
colo desert
filet plugin indent on
no <F5> :!./a.out<CR>
no <F9> :!g++ -02 -std=gnu++14 -lm % -g -fsanitize=
     undefined -Wall -Wextra -Wshadow -Wno-unused-result<
1.2 hash.sh
cpp -dD -P -fpreprocessed $1 | tr -d '[:space:]'| md5sum
      Icut -c-6
1.3 Custom Hash
struct custom_hash {
  static uint64_t splitmix64(uint64_t x) {
    x += 0x9e3779b97f4a7c15;
    x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
return x ^ (x >> 31);
  size_t operator()(uint64_t x) const {
   static const uint64_t FIXED_RANDOM = chrono::
         steady_clock::now().time_since_epoch().count();
     return splitmix64(x + FIXED_RANDOM);
|};
1.4 python-related
from fractions import Fraction
from decimal import Decimal, getcontext
getcontext().prec = 250 # set precision
itwo = Decimal(0.5)
two = Decimal(2)
N = 200
def angle(cosT):
   """given cos(theta) in decimal return theta"""
   for i in range(N):
  cosT = ((cosT + 1) / two) ** itwo

<math>sinT = (1 - cosT * cosT) ** itwo
  return sinT * (2 ** N)
pi = angle(Decimal(-1))
     flow
2
2.1 ISAP
#define SZ(c) ((int)(c).size())
struct Maxflow {
  static const int MAXV = 50010;
  static const int INF = 1000000;
  struct Edge {
    int v, c, r;
Edge(int _v, int _c, int _r):
   v(_v), c(_c), r(_r) {}
  int s, t;
  vector<Edge> G[MAXV];
  int iter[MAXV], d[MAXV], gap[MAXV], tot;
  void init(int n, int _s, int _t) {
    tot = n, s = _s, t = _t;
for(int i = 0; i <= tot; i++) {
       G[i].clear();
       iter[i] = d[i] = gap[i] = 0;
```

```
void add_edge(int u, int v, int c) {
    G[u].push_back(Edge(v, c, SZ(G[v])));
G[v].push_back(Edge(u, 0, SZ(G[u]) - 1));
  int DFS(int p, int flow) {
    if(p == t) return flow;
    for(int &i = iter[p]; i < SZ(G[p]); i++) {</pre>
      Edge &e = G[p][i];
       if(e.c > 0 \& d[p] == d[e.v]+1)
         int f = DFS(e.v, min(flow, e.c));
         if(f) {
           e.c -= f;
           G[e.v][e.r].c += f;
           return f;
      }
    if( (--gap[d[p]]) == 0) d[s] = tot;
    else {
      d[p]++; iter[p] = 0; ++gap[d[p]];
    return 0;
  int flow() {
    int res = 0;
    for(res=0,gap[0]=tot;d[s]<tot;res+=DFS(s,INF));</pre>
    return res;
} flow;
```

2.2 MinCostFlow

```
struct zkwflow {
  static const int maxN = 10000;
  struct Edge { int v, f, re; ll w; };
  int n, s, t, ptr[maxN];
  bool vis[maxN]; ll dis[maxN];
  vector<Edge> E[maxN];
  void init(int _n, int _s, int _t) {
  n = _n, s = _s, t = _t;
  for (int i = 0; i < n; i++) E[i].clear();</pre>
  void add_edge(int u, int v, int f, ll w) {
    E[u].push_back({v, f, (int)E[v].size(), w});
E[v].push_back({u, 0, (int)E[u].size()-1, -w});
  bool SPFA() {
    fill_n(dis, n, LLONG_MAX);
    fill_n(vis, n, false);
    dis[s] = 0;
    queue<int> q; q.push(s);
    while (!q.empty()) {
      int u = q.front(); q.pop();
vis[u] = false;
for (auto &it : E[u]) {
         if (it.f > 0 && dis[it.v] > dis[u] + it.w) {
           dis[it.v] = dis[u] + it.w;
            if (!vis[it.v]) {
              vis[it.v] = true;
              q.push(it.v);
           }
         }
      }
    }
    return dis[t] != LLONG_MAX;
  int DFS(int u, int nf) {
    if (u == t) return nf;
    int res = 0; vis[u] = true;
    for (int &i = ptr[u]; i < (int)E[u].size(); i++) {</pre>
       auto &it = E[u][i];
       if (it.f > 0 && dis[it.v] == dis[u] + it.w && !vis
            [it.v]) {
         int tf = DFS(it.v, min(nf, it.f));
         res += tf, nf -= tf, it.f -= tf;
         E[it.v][it.re].f += tf;
         if (nf == 0) {
           vis[u] = false;
           break;
```

```
return res;
}
pair<int, ll> flow() {
  int flow = 0; ll cost = 0;
  while (SPFA()) {
    fill_n(ptr, n, 0);
    int f = DFS(s, INT_MAX);
    flow += f;
    cost += dis[t] * f;
}
return { flow, cost };
}
flow;
```

2.3 Dinic

```
struct Dinic{
  static const int MXN = 10000;
  struct Edge{ int v,f,re; };
  int n,s,t,level[MXN];
  vector<Edge> E[MXN];
  void init(int _n, int _s, int _t){
    n = _n;    s = _s;    t = _t;
    for (int i=0; i<n; i++) E[i].clear();</pre>
  void add_edge(int u, int v, int f){
    E[u].push_back({v,f,(int)E[v].size()});
     E[v].push_back({u,0,(int)}E[u].size()-1});
  bool BFS(){
     for (int i=0; i<n; i++) level[i] = -1;</pre>
     queue<int> que;
     que.push(s);
     level[s] = 0;
     while (!que.empty()){
       int u = que.front(); que.pop();
       for (auto &it : E[u]){
         if (it.f > 0 && level[it.v] == -1){
  level[it.v] = level[u]+1;
            que.push(it.v);
       }
     }
     return level[t] != -1;
  int DFS(int u, int nf){
     if (u == t) return nf;
     int res = 0;
     for (auto &it : E[u]){
       if (it.f > 0 && level[it.v] == level[u]+1){
         int tf = DFS(it.v, min(nf,it.f));
          res += tf; nf_-= tf; it.f_-= tf;
          E[it.v][it.re].f += tf;
         if (nf == 0) return res;
     if (!res) level[u] = -1;
     return res;
  int flow(int res=0){
     while (BFS()) res += DFS(s,2147483647);
     return res:
}flow;
```

2.4 Kuhn Munkres

```
struct KM{ // max weight, for min negate the weights
  static const int MXN = 2001; // 1-based
  static const ll INF = 0x3f3f3f3f;
  int n, mx[MXN], my[MXN], pa[MXN];
  ll g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
  bool vx[MXN], vy[MXN];
  void int(int _n) {
    n = _n;
    for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0);
  }
  void addEdge(int x, int y, ll w) {g[x][y] = w;}
  void augment(int y) {
    for(int x, z; y; y = z)
        x=pa[y], z=mx[x], my[y]=x, mx[x]=y;</pre>
```

```
void bfs(int st) {
     for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;</pre>
     queue<int> q; q.push(st);
     for(;;) {
       while(q.size()) {
          int x=q.front(); q.pop(); vx[x]=1;
for(int y=1; y<=n; ++y) if(!vy[y]){</pre>
            ll t = lx[x]+ly[y]-g[x][y];
            if(t==0){
              pa[y]=x
               if(!my[y]){augment(y);return;}
               vy[y]=1, q.push(my[y]);
            }else if(sy[y]>t) pa[y]=x,sy[y]=t;
         }
       for(int y=1; y<=n; ++y)</pre>
          if(!vy[y]&&cut>sy[y]) cut=sy[y];
       for(int j=1; j<=n; ++j){
  if(vx[j]) lx[j] -= cut;
  if(vy[j]) ly[j] += cut;</pre>
         else sy[j] -= cut;
       for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){
   if(!my[y]){augment(y);return;}</pre>
          vy[y]=1, q.push(my[y]);
    }
  ll solve(){
    fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)
       lx[x] = max(lx[x], g[x][y]);
     for(int x=1; x<=n; ++x) bfs(x);</pre>
    11 \text{ ans} = 0;
     for(int y=1; y<=n; ++y) ans += g[my[y]][y];
     return ans;
}graph;
2.5 SW min-cut
const int INF=0x3f3f3f3f3f:
template<typename T>
struct stoer_wagner{// 0-base
  static const int MAXN=501;
  T g[MAXN][MAXN], dis[MAXN];
  int nd[MAXN],n,s,t;
  void init(int _n){
     n=_n;
     for(int i=0;i<n;++i)</pre>
       for(int j=0;j<n;++j)g[i][j]=0;</pre>
  void add_edge(int u,int v,T w){
    g[u][v]=g[v][u]+=w;
  T min_cut(){
     T ans=INF;
     for(int i=0;i<n;++i)nd[i]=i;</pre>
     for(int ind,tn=n;tn>1;--tn){
       for(int i=1;i<tn;++i)dis[nd[i]]=0;</pre>
       for(int i=1;i<tn;++i){</pre>
          ind=i;
          if(dis[nd[ind]]<dis[nd[j]])ind=j;</pre>
         swap(nd[ind],nd[i]);
       if(ans>dis[nd[ind]])
          ans=dis[t=nd[ind]],s=nd[ind-1];
       for(int i=0;i<tn;++i)
          g[nd[ind-1]][nd[i]] = g[nd[i]][nd[ind-1]]
                                  +=g[nd[i]][nd[ind]];
```

```
int n , m;
struct Edge { int v , w , c , r; };
vector<Edge> g[ MAXN ];
int dis[ MAXN ] , prv[ MAXN ] , prve[ MAXN ];
   bool vis[ MAXN ];
   int ans;
  void init( int _n , int _m ) : n(_n), m(_m) {}
void adde( int u , int v , int w , int c ) {
  g[u].push_back({v,w,c,(int)g[v].size()});
     g[v].push_back({u,-w,0,(int)}g[u].size()-1);
   bool poscyc() {
     fill( dis , dis+n+1 , 0 );
     fill( prv , prv+n+1 , 0 );
     fill( vis , vis+n+1 , 0 );
int tmp = -1;
     for(int t=0;t<=n;t++) {</pre>
        for(int i=1;i<=n;i++) {</pre>
          for(int j=0;j<(int)g[i].size();j++) {</pre>
             Edge& e = g[ i ][ j ];
if( e.c && dis[ e.v ] < dis[ i ]+e.w ) {
  dis[ e.v ] = dis[ i ]+e.w;
                prv[ e.v ] = i; prve[ e.v ] = j;
                if( t == n ) {
                  tmp = i;
                  break;
     int cur = tmp;
     while( !vis[ cur ] ) {
        vis[ cur ] = 1;
cur = prv[ cur ];
     int now = cur, cost = 0, df = 100000;
        Edge &e = g[prv[now]][prve[now]];
        df = min(df, e.c);
        cost += e.w;
        now = prv[ now ];
     }while( now != cur );
     ans += df*cost; now = cur;
       Edge &e = g[ prv[ now ] ][ prve[ now ] ];
Edge &re = g[ now ][ e.r ];
e.c -= df; re.c_+= df;
        now = prv[ now ];
     }while( now != cur );
     return 1;
|} circ;
2.7 Gomory-Hu Tree
//n,Dinic::flow must be filled
//result:e[u][v]=u-v mincut;p[u]:u's parent on cut tree
int n,e[MXN][MXN],p[MXN];
void gomory_hu(){
  fill(p, p+n, 0);
fill(e[0], e[n], INF);
for(int s = 1; s < n; s++){
     int t = p[s];
Dinic F; F.init(n,s,t);
     copy(flow.E,flow.E+MXN,F.E);
     int tmp = F.flow();
     for( int i = 0 ; i < s ; i++ )
     e[s][i] = e[i][s] = min(tmp, e[t][i]);
for( int i = s+1; i < n; i++)</pre>
        if ( p[i] == t && F.level[i]!=-1 ) p[i] = s;
}
2.8 Max flow with lower/upper bound
// Max flow with lower/upper bound on edges
// use with ISAP
 int in[N],out[N],1[M],r[M],a[M],b[M]
int solve(int n, int m, int s, int t){
   flow.init( n );
   for( int i = 0 ; i < m ; i ++ ){
  in[ r[ i ] ] += a[ i ];</pre>
```

struct MaxCostCirc {

static const int MAXN = 33;

2.6 Max Cost Circulation

return ans;

```
out[ l[ i ] ] += a[ i ];
flow.addEdge( l[ i ] , r[ i ] , b[ i ] - a[ i ] );
// flow from l[i] to r[i] must in [a[ i ], b[ i ]]
int nd = 0;
for( int i = 0 ; i <= n ; i ++ ){
  if( in[ i ] < out[ i ] ){
    flow.addEdge( i , flow.t , out[ i ] - in[ i ] );
    nd += out[ i ] - in[ i ];
}</pre>
    if( out[ i ] < in[ i ] )</pre>
       flow.addEdge( flow.s , i , in[ i ] - out[ i ] );
 // original sink to source
flow.addEdge( t , s , INF );
if( flow.solve() != nd ) return -1; // no solution
int ans = flow.G[ s ].back().c; // source to sink
flow.G[ s ].back().c = flow.G[ t ].back().c = 0;
// take out super source and super sink
for(size_ti=0;i<flow.G[flow.s].size();i++){
  flow.G[ flow.s ][ i ].c = 0;
  Maxflow::Edge &e = flow.G[ flow.s ][ i ];</pre>
    flow.G[e.v][e.r].c = 0;
for(size_ti=0;i<flow.G[flow.t].size();i++){</pre>
   flow.G[ flow.t ][ i ].c = 0;
Maxflow::Edge &e = flow.G[ flow.t ][ i ];
    flow.G[ e.v ][ e.r ].c = \bar{0};
flow.addEdge( flow.s , s , INF );
flow.addEdge( t , flow.t , INF );
flow.reset(); // set iter,d,gap to 0
return ans + flow.solve();
```

2.9 HLPPA

```
template <int MAXN, class T = int>
struct HLPP {
  const T INF = numeric_limits<T>::max();
  struct Edge { int to, rev; T f; };
  int n, s, t;
  vector<Edge> adj[MAXN];
  deque<int> lst[MAXN];
  vector<int> gap[MAXN];
  T ef[MAXN];
  int ptr[MAXN],h[MAXN],cnt[MAXN],work,hst=0; // highest
  void init(int _n, int _s, int _t) {
    n=_n+1; s = _s; t = _t;
    for(int i=0;i<n;i++) adj[i].clear();</pre>
  void add_edge(int u,int v,T f,bool isDir = true){
  adj[u].push_back({v,adj[v].size(),f});
    adj[v].push_back({u,adj[u].size()-1,isDir?0:f});
  void updHeight(int v, int nh) {
    work++;
     if(h[v] != n) cnt[h[v]]--;
    h[v] = nh;
    if(nh == n) return;
     cnt[nh]++, hst = nh; gap[nh].push_back(v);
     if(ef[v]>0) lst[nh].push_back(v), ptr[nh]++;
  void globalRelabel() {
    work = 0; fill(h, h+n, n); fill(cnt, cnt+n, 0);
     for(int i=0; i<=hst; i++)</pre>
    lst[i].clear(), gap[i].clear(), ptr[i] = 0;
queue<int> q({t}); h[t] = 0;
    while(!q.empty()) {
       int v = q.front(); q.pop();
for(auto &e : adj[v])
         if(h[e.to] == n \&\& adj[e.to][e.rev].f > 0)
            q.push(e.to), updHeight(e.to, h[v] + 1);
       hst = h[v];
    }
  }
  void push(int v, Edge &e) {
    if(ef[e.to] == 0)
       lst[h[e.to]].push_back(e.to), ptr[h[e.to]]++;
    T df = min(ef[v], e.f);
e.f -= df, adj[e.to][e.rev].f += df;
ef[v] -= df, ef[e.to] += df;
```

```
void discharge(int v) {
     int nh = n;
     for(auto &e : adj[v]) {
       if(e.f > 0) {
         if(h[v] == h[e.to] + 1) {
            push(v, e);
            if(ef[v] <= 0) return;</pre>
         else nh = min(nh, h[e.to] + 1);
     if(cnt[h[v]] > 1) updHeight(v, nh);
     else {
       for(int i = h[v]; i < n; i++) {
         for(auto j : gap[i]) updHeight(j, n);
gap[i].clear(), ptr[i] = 0;
    }
  T flow() {
     fill(ef, ef+n, 0); ef[s] = INF, ef[t] = -INF;
     globalRelabel();
     for(auto &e : adj[s]) push(s, e);
     for(; hst >= 0; hst--) {
  while(!lst[hst].empty())
         int v=lst[hst].back(); lst[hst].pop_back();
         discharge(v);
if(work > 4 * n) globalRelabel();
     return ef[t] + INF;
  }
};
```

2.10 Flow Method

To reconstruct the minimum vertex cover, dfs from each unmatched vertex on the left side and with unused edges only. Equivalently, dfs from source with unused edges only and without visiting sink. Then, a vertex is chosen iff. it is on the left side and without visited or on the right side and visited through dfs.

Minimum Weighted Bipartite Edge Cover:

Construct new bipartite graph with n+m vertices on each side:

for each vertex u, duplicate a vertex u' on the other side

for each edge (u,v,w), add edges (u,v,w) and (v',u',w)

for each vertex u, add edge (u,u',2w) where w is min edge connects to u

then the answer is the minimum perfect matching of the new graph (KM)

Maximum density subgraph (\sum{W_e}+ \sum{W_v}) / IVI

Binary search on answer:

```
Requiring subgraph: all vertex can be reached from
    source with
edge whose cap > 0.
Maximum closed subgraph
1. connect source with positive weighted vertex(capacity
    =weight)
connect sink with negitive weighted vertex(capacity=-
   weiaht)
make capacity of the original edges = inf
4. ans = sum(positive weighted vertex weight) - (max
    flow)
Minimum Path Cover of DAG
1. For each vertex v, split it to v_in and v_out.
2. For each edge (u->v), add an edge between u_out and
    v_in
3. |Minimum Path Cover| = |V| - |Maximum Matching| of
    the new bipartite graph
```

3 Math

3.1 FFT

```
const int MAXN = 262144;
// (must be 2^k)
// before any usage, run pre_fft() first
typedef long double ld;
typedef complex<ld> cplx;
const ld PI = acosl(-1);
const cplx I(0, 1);
cplx omega[MAXN+1];
void pre_fft(){
  for(int i=0; i<=MAXN; i++)</pre>
    omega[i] = exp(i * 2 * PI / MAXN * I);
// n must be 2^k
void fft(int n, vector<cplx> &a, bool inv=false){
  int basic = MAXN / n;
  int theta = basic;
  for (int m = n; m >= 2; m >>= 1) {
    int mh = m >> 1;
for (int i = 0; i < mh; i++) {</pre>
       cplx w = omega[inv ? MAXN-(i*theta%MAXN)]
                            : i*theta%MAXN];
       for (int j = i; j < n; j += m) {
         int k = j + mh;
         cplx x = a[j] - a[k];
         a[j] += a[k];
         a[k] = w * x;
    theta = (theta * 2) % MAXN;
  int i = 0;
  for (int j = 1; j < n - 1; j++) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
     if (j < i) swap(a[i], a[j]);</pre>
  if(inv) for (i = 0; i < n; i++) a[i] /= n;
}
```

3.2 NTT

```
/* p=a*2^k+1
                                         root
  998244353
                          119
                                 23
                                         3
   2013265921
                          15
                                 27
                                         31
   2061584302081
                          15
                                 37
   2748779069441
                                 39
   1945555039024054273
                          27
                                 56
                                         5 */
template<ll P,ll root,int MAXK,int MAXN>
struct NTT{
  static ll powi(ll a,ll b){
    ll ret=1;
    for(;b;b>>=1,a=mul(a, a, P)){
      if(b&1) ret=mul(ret, a, P);
    return ret;
  static ll inv(ll a,ll b){
    if(a==1) return 1;
```

```
return (((a-inv(b%a,a))*b+1)/a)%b; // overflow
  11 omega[MAXK+1],inv_omega[MAXK+1];
  NTT(){
    omega[MAXK]=powi(root,(P-1)>>MAXK);
    for(int i=MAXK-1;i>=0;i--)
      omega[i]=mul(omega[i+1], omega[i+1], P);
     for(int i=0;i<=MAXK;i++)
      inv_omega[i]=inv(omega[i],P);
  void tran(int n,ll a[],bool inv_ntt=false){//n=2^i
    for(int i=1,j=0;i<n;i++){</pre>
      for(int k=n>>1;!((j^=k)&k);k>>=1);
      if(i<j) swap(a[i],a[j]);</pre>
    11 *G=(inv_ntt?inv_omega:omega);
    for(int k=2,t=1;k<=n;k<<=1){</pre>
      int k2=k>>1;ll dw=G[t++];
      for(int j=0;j<n;j+=k){</pre>
        11 w=1;
        for(int i=j;i<j+k2;i++){</pre>
           ll x=a[i], y=mul(a[i+k2], w, P);
          a[i]=x+y; if(a[i]>=P) a[i]-=P
          a[i+k2]=x-y; if(a[i+k2]<0) a[i+k2]+=P;
          w=mul(w, dw, P);
        }
      }
    }
    if(inv_ntt){
      ll inv_n=inv(n,P);
      for(int i=0;i<n;i++) a[i]=mul(a[i], inv_n, P);</pre>
  }
};
const ll P=2013265921,root=31;
const int MAXN=4194304, MAXK=22; //MAXN=2^k
NTT<P,root,MAXK,MAXN> ntt;
```

3.3 Fast Walsh Transform

```
/* xor convolution:
 * x = (x0,x1) , y = (y0,y1)
* z = (x0y0 + x1y1 , x0y1 + x1y0 )
 * x' = (x0+x1, x0-x1), y' = (y0+y1, y0-y1)
* z' = ((x0+x1)(y0+y1), (x0-x1)(y0-y1))
* z = (1/2) * z''
 * or convolution:
 * x = (x0, x0+x1), inv = (x0, x1-x0) w/o final div
 * and convolution:
 * x = (x0+x1, x1), inv = (x0-x1, x1) w/o final div
 * ternery xor convolution:
   x = (x0+x1+x2,x0+x1w+x2w^2,x0+x1w^2+x2w)
 * inv = (1/3) * (x0+x1+x2,x0+x1w^2+x2w,x0+x1w+x2w^2)
 * where w^3=1 and w^2=-w-1 */
typedef long long ll;
const int MAXN = (1 < 20) + 10;
const ll MOD = 1e9+7;
inline ll pw( ll x , ll k ) {
  ll res = 1;
  for( ll bs = x ; k ; k >>= 1, bs = (bs * bs)%MOD )
  if( k&1 ) res = ( res * bs ) % MOD;
  return res:
inline ll invf( ll x )
  return pw(x, MOD-2);
inline void fwt( ll x[ MAXN ] , int N , bool inv=0 ) {
  for( int d = 1 ; d < N ; d <<= 1 ) {
    int d2 = d<<1;</pre>
     for( int s = 0 ; s < N ; s += d2 )
        for( int i = s , j = s+d ; i < s+d ; i++, j++ ){
    ll ta = x[ i ] , tb = x[ j ];</pre>
          x[i] = ta+tb;
x[j] = ta-tb;
if(x[i] >= MOD ) x[i] -= MOD;
if(x[j] < 0 ) x[j] += MOD;</pre>
  ll invN = invf( N );
  if( inv )
     for( int i = 0 ; i < N ; i++ ) {</pre>
```

static ll lnb[MAXN], c[MAXN], tmp[MAXN];

```
x[ i ] *= invN;
x[ i ] %= MOD;
                                                                        assert(a[0] == 0); // dont know exp(a[0]) mod P if (n == 1) \{b[0] = 1; return;\}
                                                                        Exp((n+1)/2, a, b);
                                                                        fill(b+(n+1)/2, b+n, 0);
}
                                                                        Ln(n, b, lnb);
                                                                        fill(c, c+n, \emptyset); c[\emptyset] = 1;
3.4 Poly operator
                                                                        FOR(i, n) {
 c[i] += a[i] - lnb[i];
struct PolyOp {
                                                                           if (c[i] < 0) c[i] += P
#define FOR(i, c) for (int i = 0; i < (c); ++i)
                                                                           if (c[i] \rightarrow P) c[i] \rightarrow P;
  NTT<P, root, MAXK, MAXN> ntt;
  static int nxt2k(int x) {
    int i = 1; for (; i < x; i <<= 1); return i;</pre>
                                                                        Mul(n, b, n, c, tmp);
                                                                        copy(tmp, tmp+n, b);
  void Mul(int n, ll a[], int m, ll b[], ll c[]) {
   static ll aa[MAXN], bb[MAXN];
                                                                      bool Sqrt(int n, ll a[], ll b[]){
                                                                        // Square root of a : b * b = a ( mod x^n )
     int N = nxt2k(n+m);
    copy(a, a+n, aa); fill(aa+n, aa+N, 0); copy(b, b+m, bb); fill(bb+m, bb+N, 0);
                                                                        // bb = a mod x^(n/2)
                                                                        // ( bb - a )^2 = 0 mod x^n
                                                                        // ( bb + a )^2 = 4 bba
    ntt.tran(N, aa); ntt.tran(N, bb);
                                                                        // ( ( bb + a ) / 2b )^2 = a // sqrt(a) = <math>b / 2 + a / 2b
    FOR(i, N) c[i] = aa[i] * bb[i] % P;
ntt.tran(N, c, 1);
                                                                        static ll c[MAXN];
                                                                        int ind=0,x,y,p=1
  void Inv(int n, ll a[], ll b[]) {
    // ab = aa^{-1} = 1 \mod x^{(n/2)}
                                                                        while(a[ind]==0) ind++;
    // (b - a^-1)^2 = 0 mod x^n
                                                                        for(int i=0;i<n;i++) a[i]=a[i+ind];</pre>
                                                                        if((ind&1)||!solve(a[0],mod,x,y)) // discrete sqrt
    // bb + a^-2 - 2 ba^-1 = 0
    // bba + a^{-1} - 2b = 0
                                                                        b[0]=min(x,y);
    // a^{-1} = 2b - bba
    static ll tmp[MAXN];
                                                                        while(p<n) p<<=1;</pre>
                                                                        for(int t=2;t<=p;t<<=1){
   Inv(t,b,c); Mul(t,a,t,c,c);</pre>
    if (n == 1) {b[0] = ntt.inv(a[0], P); return;}
Inv((n+1)/2, a, b);
    int N = nxt2k(n*2);
                                                                           for(int i=0;i<t;i++)</pre>
    copy(a, a+n, tmp);
fill(tmp+n, tmp+N, 0);
                                                                             b[i]=(b[i]+c[i])*inv(2)%mod;
    fill(b+n, b+N, 0);
                                                                        if(ind){
                                                                           for(int i=p-1;i>=ind/2;i--) b[i]=b[i-ind/2];
    ntt.tran(N, tmp); ntt.tran(N, b);
                                                                           for(int i=0;i<ind/2;i++) b[i]=0;</pre>
       ll t1 = (2 - b[i] * tmp[i]) % P;
                                                                           for(int i=p-1;i>=ind;i--) a[i]=a[i-ind];
       if (t1 < 0) t1 += P
                                                                           for(int i=0;i<ind;i++) a[i]=0;</pre>
       b[i] = b[i] * t1 % P;
                                                                        }
                                                                      }
                                                                   } polyop;
    ntt.tran(N, b, 1);
    fill(b+n, b+N, 0);
                                                                    3.5 Linear Recurrence
  void Div(int n, ll a□, int m, ll b□, ll d□, ll r□)
                                                                    // Usage: linearRec({0, 1}, {1, 1}, k) //k'th fib
    // Ra = Rb * Rd mod x^{n-m+1}
                                                                    typedef vector<ll> Poly
    // Rd = Ra * Rb^{-1} mod
                                                                    ll linearRec(Poly&& S, Poly&& tr, ll k) {
                                                                      int n=tr.size();
    static ll aa[MAXN], bb[MAXN], ta[MAXN], tb[MAXN];
                                                                      auto combine=[&](Poly& a, Poly& b) {
    if (n < m) {copy(a, a+n, r); fill(r+n, r+m, 0);
                                                                        Poly res(n*2+1);
                                                                        for(int i=0;i<=n;i++) for(int j=0;j<=n;j++)</pre>
    // d: n-1 - (m-1) = n-m (n-m+1 terms)
    copy(a, a+n, aa); copy(b, b+m, bb);
reverse(aa, aa+n); reverse(bb, bb+m);
                                                                         res[i+j]=(res[i+j]+a[i]*b[j])%mod;
for(int i=2*n;i>n;--i) for(int j=0;j<n;j++)
    Inv(n-m+1, bb, tb);
Mul(n-m+1, ta, n-m+1, tb, d);
                                                                           res[i-1-j]=(res[i-1-j]+res[i]*tr[j])%mod;
                                                                        res.resize(n+1);
    fill(d+n-m+1, d+n, 0); reverse(d, d+n-m+1);
                                                                        return res;
                                                                      }; // a * b mod (x^n-tr)
     // r: m-1 - 1 = m-2 (m-1 terms)
    Poly pol(n+1), e(pol);
                                                                      pol[0]=e[1]=1;
         += P; }
                                                                      for (++k;k;k/=2) {
                                                                        if(k%2)pol=combine(pol,e);
  void dx(int n, ll a[], ll b[]) { for(int i=1;i<=n-1;i</pre>
                                                                        e=combine(e,e);
       ++) b[i-1] = i * a[i] % P;
                                                                      ll res=0;
  void Sx(int n, ll a[], ll b[]) {
    b[0] = 0;
                                                                      for(int i=0;i<n;i++) res=(res+pol[i+1]*S[i])%mod;</pre>
     FOR(i, n) b[i+1] = a[i] * ntt.inv(i+1,P) % P;
                                                                      return res;
  void Ln(int n, ll a[], ll b[]) {
   // Integral a' a^-1 dx
                                                                    3.6 BerlekampMassey
     static ll a1[MAXN], a2[MAXN], b1[MAXN];
     int N = nxt2k(n*2);
                                                                    // find shortest linear recurrence relation 0(n^2)
     dx(n, a, a1); Inv(n, a, a2);
                                                                    // example: BM({1,1,2,3,5,8,13,21})
    Mul(n-1, a1, n, a2, b1);
Sx(n+n-1-1, b1, b);
                                                                    // 2*len terms for uniqueness
                                                                    inline vector<ll> BM(const vector<ll> &x) {
                                                                      vector<ll> ls, cur;
     fill(b+n, b+N, 0);
                                                                      int lf; ll ld;
for(int i=0;i<x.size();++i) {</pre>
  void Exp(int n, ll a[], ll b[]) {
    // Newton method to solve g(a(x))=\ln(b(x))-a(x)=0
                                                                        11 t=0:
    // b' = b - g(b(x)) / g'(b(x))
// b' = b (1 - lnb + a)
                                                                        for(int j=0;j<cur.size();++j)</pre>
                                                                           t=(t+x[i-j-1]*cur[j])%mod;
```

if((t-x[i])%mod==0) continue;

```
if(!cur.size()) {
    cur.resize(i+1);lf=i;ld=(t-x[i])%mod;continue;
}
ll k=-(x[i]-t)*inv(ld, mod)%mod;
vector<ll> c(i-lf-1); c.push_back(k);
for(auto j:ls) c.push_back(-j*k%mod);
if(c.size()<cur.size()) c.resize(cur.size());
for(int j=0;j<cur.size();++j)
    c[j]=(c[j]+cur[j])%mod;
if(i-lf+(int)ls.size()>=(int)cur.size())
    ls=cur,lf=i,ld=(t-x[i])%mod;
cur=move(c);
}
for(auto& xx:cur) xx=(xx%mod+mod)%mod;
return cur;
}
```

3.7 Miller Rabin

```
// n < 4,759,123,141
// n < 1,122,004,669,633
                                   2, 7, 61
2, 13, 23, 1662803
                              3:
                                     6 : pirmes <= 13
// n < 3,474,749,660,383
// n < 2^64
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
bool witness(ll a,ll n,ll u,int t){
  if(!a) return 0;
  11 x=mypow(a,u,n);
  for(int i=0;i<t;i++) {</pre>
    ll nx=mul(x,x,n)
    if(nx==1&&x!=1&&x!=n-1) return 1;
    x=nx;
  return x!=1;
bool miller_rabin(ll n,int s=100) {
  // iterate s times of witness on n
  // return 1 if prime, 0 otherwise
  if(n<2) return 0;</pre>
  if(!(n&1)) return n == 2;
  ll u=n-1; int t=0;
  while(!(u&1)) u>>=1, t++;
  while(s--){
    ll a=randll()%(n-1)+1;
    if(witness(a,n,u,t)) return 0;
  return 1;
```

3.8 Simplex

```
/*taraet:
  \max \sum_{j=1}^n A_{0,j}*x_j
condition:
  \sum_{j=1}^n A_{i,j}*x_j \ll A_{i,0} i=1~m
  x_j >= 0 | j=1\sim n
VDB = vector<double>*/
template<class VDB>
VDB simplex(int m,int n,vector<VDB> a){
  vector<int> left(m+1), up(n+1);
  iota(left.begin(), left.end(), n);
  iota(up.begin(), up.end(), 0);
auto pivot = [&](int x, int y){
  swap(left[x], up[y]);
     auto k = a[x][y]; a[x][y] = 1;
     vector<int> pos;
for(int j = 0; j <= n; ++j){
    a[x][j] /= k;</pre>
        if(a[x][j] != 0) pos.push_back(j);
     for(int i = 0; i <= m; ++i){
        if(a[i][y]==0 | | i == x) continue;
       k = a[i][y], a[i][y] = 0;
for(int j : pos) a[i][j] -= k*a[x][j];
  for(int x,y;;){
     for(int i=x=1; i <= m; ++i)
  if(a[i][0]<a[x][0]) x = i;</pre>
     if(a[x][0]>=0) break;
     for(int j=y=1; j <= n; ++j)</pre>
     if(a[x][j] < a[x][y])'y = j;

if(a[x][y] >= 0) return VDB();//infeasible
```

3.9 Faulhaber

```
/* faulhaber's formula -
 * cal power sum formula of all p=1\sim k in O(k^2) */
#define MAXK 2500
const int mod = 1000000007;
int b[MAXK]; // bernoulli number
int inv[MAXK+1]; // inverse
int cm[MAXK+1] [MAXK+1]; // combinatorics
int co[MAXK] [MAXK+2]; // coeeficient of x^j when p=i
inline int getinv(int x) {
  int a=x,b=mod,a0=1,a1=0,b0=0,b1=1;
  while(b) {
    int q,t;
    q=a/b; t=b; b=a-b*q; a=t;
t=b0; b0=a0-b0*q; a0=t;
     t=b1; b1=a1-b1*q; a1=t;
  return a0<0?a0+mod:a0;</pre>
inline void pre() {
  for(int i=0;i<=MAXK;i++) {</pre>
     cm[i][0]=cm[i][i]=1;
     for(int j=1;j<i;j++)
   cm[i][j]=add(cm[i-1][j-1],cm[i-1][j]);</pre>
  for(int i=1;i<=MAXK;i++) inv[i]=getinv(i);</pre>
  b[0]=1; b[1]=getinv(2); // with b[1] = 1/2
for(int i=2;i<MAXK;i++) {
     if(i&1) { b[i]=0; continue; }
     b[i]=1;
     for(int j=0;j<i;j++)</pre>
       b[i]=sub(b[i], mul(cm[i][j],mul(b[j], inv[i-j+1]))
  }
/* faulhaber */
  // sigma_x=1~n \{x^p\} =
        1/(p+1) * sigma_j=0~p {C(p+1,j)*Bj*n^(p-j+1)}
  for(int i=1;i<MAXK;i++) {</pre>
     co[i][0]=0;
     for(int j=0;j<=i;j++)</pre>
       co[i][i-j+1]=mul(inv[i+1], mul(cm[i+1][j], b[j]));
/* sample usage: return f(n,p) = sigma_x=1\sim (x^p) */
inline int solve(int n,int p) {
  int sol=0,m=n;
  for(int i=1;i<=p+1;i++)</pre>
     sol=add(sol,mul(co[p][i],m));
    m = mul(m, n);
  return sol;
```

3.10 Chinese Remainder

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
    ll g = __gcd(m1, m2);
    if((x2 - x1) % g) return -1;// no sol
    m1 /= g; m2 /= g;
    pair<ll,ll> p = gcd(m1, m2);
    ll lcm = m1 * m2 * g;
```

```
ll res=p.first*(x2-x1)%lcm*m1%lcm+x1; // overflow m^3
  return (res % lcm + lcm) % lcm;
}
```

3.11 Pollard Rho

```
// does not work when n is prime
ll f(ll x, ll mod){ return add(mul(x,x,mod),1,mod); }
ll pollard_rho(ll n) {
   if(!(n&1)) return 2;
   while(true){
      ll y=2, x=rand()%(n-1)+1, res=1;
      for(int sz=2; res==1; sz*=2) {
        for(int i=0; i<sz && res<=1; i++) {
            x = f(x, n);
            res = __gcd(abs(x-y), n);
      }
      y = x;
    }
   if (res!=0 && res!=n) return res;
}</pre>
```

3.12 ax+by=gcd

```
pair<ll, ll> gcd(ll a, ll b){
   if(b == 0) return {1, 0};
   pair<ll, ll> q = gcd(b, a % b);
   return {q.second, q.first - q.second * (a / b)};
}
```

3.13 Discrete sqrt

```
void calcH(int &t, int &h, const int p) {
  int tmp=p-1; for(t=0;(tmp&1)==0;tmp/=2) t++; h=tmp;
// solve equation x^2 \mod p = a where p is a prime
bool solve(int a, int p, int &x, int &y) {
  if(p == 2) { x = y = 1; return true; }
int p2 = p / 2, tmp = mypow(a, p2, p);
if (tmp == p - 1) return false;
  if ((p + 1) \% 4 == 0) {
     x=mypow(a,(p+1)/4,p); y=p-x; return true;
  } else {
     int t, h, b, pb; calcH(t, h, p);
if (t >= 2) {
        do \{b = rand() \% (p - 2) + 2;
        } while (mypow(b, p / 2, p) != p - 1);
       pb = mypow(b, h, p);
int s = mypow(a, h / 2, p);
     for (int step = 2; step <= t; step++) {
  int ss = (((ll)(s * s) % p) * a) % p;</pre>
        for(int i=0;i<t-step;i++) ss=mul(ss,ss,p);</pre>
     if (ss + 1 == p) s = (s * pb) % p;
pb = ((ll)pb * pb) % p;
} x = ((ll)s * a) % p; y = p - x;
  } return true;
```

3.14 Romberg

3.15 Simpson

```
template<class F>
ld quad(ld a, ld b, F f, const int n = 1000) {
    ld h = (b-a)/2/n, v=f(a)+f(b);
```

```
for(int i = 1; i < n*2; ++i)
    v += f(a+i*h) * (i&1 ? 4 : 2);
return v * h / 3;</pre>
```

3.16 Prefix Inverse

```
void solve( int m ){
  inv[ 1 ] = 1;
  for( int i = 2 ; i < m ; i ++ )
     inv[ i ] = ((ll)(m - m / i) * inv[m % i]) % m;
}</pre>
```

3.17 Roots of Polynomial

```
const double eps = 1e-12;
const double inf = 1e+12;
double a[ 10 ], x[ 10 ]; // a[0..n](coef) must be filled
int n; // degree of polynomial must be filled
int sign( double x ){return (x < -eps)?(-1):(x>eps);}
double f(double a[], int n, double x){
  double tmp=1, sum=0;
  for(int i=0;i<=n;i++)</pre>
  { sum=sum+a[i]*tmp; tmp=tmp*x; }
  return sum;
double binary(double l,double r,double a[],int n){
  int sl=sign(f(a,n,l)),sr=sign(f(a,n,r));
  if(sl==0) return l; if(sr==0) return r;
  if(sl*sr>0) return inf;
  while(r-l>eps){
     double mid=(l+r)/2;
     int ss=sign(f(a,n,mid));
     if(ss==0) return mid;
     if(ss*sl>0) l=mid; else r=mid;
  return 1;
void solve(int n,double a[],double x[],int &nx){
  if(n==1){ x[1]=-a[0]/a[1]; nx=1; return; }
  double da[10], dx[10]; int ndx;
  for(int i=n;i>=1;i--) da[i-1]=a[i]*i;
  solve(n-1,da,dx,ndx);
  nx=0:
  if(ndx==0){
     double tmp=binary(-inf,inf,a,n);
     if (tmp<inf) x[++nx]=tmp;</pre>
  double tmp;
  tmp=binary(-inf,dx[1],a,n);
  if(tmp<inf) x[++nx]=tmp;
for(int i=1;i<=ndx-1;i++){</pre>
     tmp=binary(dx[i],dx[i+1],a,n);
     if(tmp<inf) x[++nx]=tmp;</pre>
  tmp=binary(dx[ndx],inf,a,n);
  if(tmp<inf) x[++nx]=tmp;</pre>
} // roots are stored in x[1..nx]
```

3.18 Primes and μ function

```
/* 12721, 13331, 14341, 75577, 123457, 222557, 556679
* 999983, 1097774749, 1076767633, 100102021, 999997771
* 1001010013, 1000512343, 987654361, 999991231
* 999888733, 98789101, 987777733, 999991921, 1010101333
* 1010102101, 1000000000039, 100000000000037
* 2305843009213693951, 4611686018427387847
* 9223372036854775783, 18446744073709551557 */
int mu[ N ] , p_tbl[ N ]; // mobius, min prime factor vector<int> primes;
void sieve() { // calculate multiplicative function f
    mu[ 1 ] = p_tbl[ 1 ] = 1;
    for( int i = 2 ; i < N ; i ++ ){
        if( !p_tbl[ i ] ){
            p_tbl[ i ] = i;
            primes.push_back( i );
            mu[ i ] = -1; // f(i)=... where i is prime
    }
    for( int p : primes ){
        int x = i * p;
        if( x >= N ) break;
```

```
p_tbl[ x ] = p;
mu[ x ] = -mu[ i ];
if( i % p == 0 ){ // f(x)=f(i)/f(p^(k-1))*f(p^k)
          mu[x] = 0;
          break;
        } // else f(x)=f(i)*f(p) where gcd(i,p)=1
  }
vector<int> factor( int x ){
  vector<int> fac{ 1 };
while( x > 1 ){
  int fn = fac.size(), p = p_tbl[ x ], pos = 0;
  while( x % p == 0 ){
        x \neq p;
        for( int i = 0 ; i < fn ; i ++ )
           fac.push_back( fac[ pos ++ ] * p );
  return fac;
```

3.19 Subset Convolution

```
// h(s) = \sum_{s' \leq s} f(s')g(s \cdot s')
vector<int> SubsetConv(int n, const vector<int> &f,
      const vector<int> &g) {
   const int m = 1 \ll n;
   vector<vector<int>> a(n + 1, vector<int>(m)), b=a;
   for (int i = 0; i < m; ++i) {
    a[__builtin_popcount(i)][i] = f[i];</pre>
      b[__builtin_popcount(i)][i] = g[i];
   for (int i = 0; i <= n; ++i) {
  for (int j = 0; j < n; ++j) {
    for (int s = 0; s < m; ++s) {
      if (s >> j & 1) {
        a[i][s] += a[i][s ^ (1 << j)];
      b[i][s] += b[i][s ^ (1 << j)];
}</pre>
   vector<vector<int>> c(n + 1, vector<int>(m));
   for (int s = 0; s < m; ++s) {
  for (int i = 0; i <= n; ++i) {</pre>
         for(intj=0; j<=i;++j)c[i][s]+=a[j][s]*b[i-j][s];</pre>
   for (int i = 0; i <= n; ++i) {
  for (int j = 0; j < n; ++j) {
    for (int s = 0; s < m; ++s) {</pre>
            if (s >> j & 1) c[i][s] -= c[i][s ^ (1 << j)];</pre>
   vector<int> res(m);
   for (int i = 0; i < m; ++i)
      res[i] = c[__builtin_popcount(i)][i];
   return res;
}
```

3.20 Result

- Lucas' Theorem : For $n,m\in\mathbb{Z}^*$ and prime P, $C(m,n)\mod P=\Pi(C(m_i,n_i))$ where m_i is the i-th digit of m in base P.
- 1st Stirling Numbers(permutation |P|=n with k cycles): $\begin{array}{l} S(n,k) = \text{coefficient of } x^k \text{ in } \Pi_{i=0}^{n-1}(x+i) \\ S(n+1,k) = nS(n,k) + S(n,k-1) \end{array}$
- 2nd Stirling Numbers(Partition n elements into k non-empty set): $S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n$ S(n+1,k) = kS(n,k) + S(n,k-1)
- Calculate f(x+n) where $f(x) = \sum\limits_{i=0}^{n-1} a_i x^i$: $f(x+n) = \sum\limits_{i=0}^{n-1} a_i (x+n)^i = \sum\limits_{i=0}^{n-1} x^i \cdot \frac{1}{i!} \sum\limits_{j=i}^{n-1} \frac{a_j}{j!} \cdot \frac{n^{j-i}}{(j-i)!}$
- Calculate $c[i-j]+=a[i] \times b[j]$ for a[n],b[m] 1. a=reverse(a); c=mul(a,b); c=reverse(c[:n]); b=reverse(b); c=mul(a,b); c=rshift(c,m-1);
- Eulerian number(permutation $1 \sim n$ with m a[i] > a[i-1]): $A(n,m) = \sum_{i=0}^{m} (-1)^{i} {\binom{n+1}{i}} (m+1-i)^{n}$ A(n,m) = (n-m)A(n-1,m-1) + (m+1)A(n-1,m)
- $D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n$

• Pick's Theorem : A=i+b/2-1• Euler Characteristic: planar graph: V-E+F-C=1 convex polyhedron: V-E+F=2V,E,F,C: number of vertices, edges, faces(regions), and components - number of spanning tree of undirected graph: - number of spanning tree of undirected graph. degree matrix $D_{ii} = deg(i)$, $D_{ij} = 0$ adjacency matrix $G_{ij} = \# \ of \ (i,j) \in E$, $G_{ii} = 0$, let A = D - G, delete any one row, one column, and cal det(A') - number of spanning tree of directed graph: in-degree matrix $D_{ii}^{in} = indeg(i)$, $D_{ij}^{in} = 0$ out-degree matrix $D_{ii}^{out} = outdeg(i)$, $D_{ij}^{out} = 0$ let $L^{in} = D^{in} - G$, $L^{out} = D^{out} - G$, delete the i-th row and column $det(L_i^{in})$ and $det(L_i^{out})$ is the number of spanning tree from/to root i• Burnside Lemma: $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$ • Polya theorem: $|Y^x/G| = \frac{1}{|G|} \sum_{g \in G} m^{c(g)}$ $m=\left|Y\right|$: num of colors, c(g) : num of cycle • Anti SG (the person who has no strategy wins) : first player wins iff either 1. SG value of ALL subgame ≤ 1 and SG value of the game =0 2. SG value of some subgame >1 and SG value of the game $\neq 0$ • Möbius inversion formula : $g(n) = \sum\limits_{d \mid n} f(d)$ for every integer $n \geq 1$, then $f(n) = \sum_{d \mid n}^{d \mid n} \mu(d) g(\frac{n}{d}) = \sum_{d \mid n} \mu(\frac{n}{d}) g(d) \text{ for every integer } n \geq 1$ Dirichlet convolution : $f * g = g * f = \sum_{d \mid n} f(d) g(\frac{n}{d}) = \sum_{d \mid n} f(\frac{n}{d}) g(d)$ $\begin{array}{l} g=f*1 \Leftrightarrow f=g*\mu \text{, } \epsilon=\mu*1 \text{, } Id=\phi*1 \text{, } d=1*1 \text{, } \sigma=Id*1=\phi*d \text{,} \\ \sigma_k=Id_k*1 \text{ where } \epsilon(n)=[n=1] \text{, } 1(n)=1 \text{, } Id(n)=n \text{, } Id_k(n)=n^k \text{,} \\ d(n)=\#(divisor) \text{, } \sigma(n)=\sum divisor, \ \sigma_k(n)=\sum divisor^k \end{array}$ • Find a Primitive Root of n: n has primitive roots iff $n=2,4,p^k,2p^k$ where p is an odd prime. 1. Find $\phi(n)$ and all prime factors of $\phi(n)$, says $P=\{p_1,...,p_m\}$ 2. $\forall g \in [2,n)$, if $g^{\frac{\phi(n)}{p_i}} \neq 1, \forall p_i \in P$, then g is a primitive root. 3. Since the smallest one isn't too big, the algorithm runs fast. 4. n has exactly $\phi(\phi(n))$ primitive roots. Sum of Two Squares Thm (Legendre): For a given positive integer N, let $D1=(\#\mbox{ of }d\in N\mbox{ dividing }N\mbox{ that }d=1\mbox{ (mod 4)})$ $D3=(\#\ of\ d\in N\ dividing\ N\ that\ d=3\ (\mod\ 4))$ then N can be written as a sum of two squares in exactly R(N)=4(D1-D3) ways. • Difference of D1-D3 Thm: let $N=2^t\times[p_1^{e_1}\times\ldots\times p_r^{e_r}]\times[q_1^{f_1}\times\ldots\times q_s^{f_s}]$ where $p_i\in mod\ 4=1\ prime$, $q_i\in mod\ 4=3\ prime$ then $D1-D3= \begin{cases} (e1+1)(e2+1)...(er+1) & if \ f_i \ all \ even \\ 0 & if \ any \ f_i \ is \ odd \end{cases}$ • Sherman-Morrison formula: suppose $A \in \mathbb{R}^{n \times n}$ is invertible and $u,v \in \mathbb{R}^n$ $A + uv^T$ is invertible if and only if $1 + v^TA^{-1}u \neq 0$ $(A + uv^T)^{-1} = A^{-1} - \frac{A^{-1}uv^TA^{-1}}{1 + v^TA^{-1}u}$ • Pohlig-Hellman algorithm (discrete log): Given an order n group, generator g, element h, find x s.t. $g^x = h$. * let $x_0=0$, $\gamma=g^{p^{e-1}}$ where γ has order p . * for $k=0\sim e-1$: 1. let $h_k = (g^{-x_k}h)^{p^e-1-k}$ whose order divide $p \implies h_k \in \langle \gamma \rangle$. 2. find d_k s.t. $\gamma^{d_k} = h_k$ with baby-step giant-step in $O(\sqrt{p})$. 3. set $x_{k+1}=x_k+p^kd_k$ * return x_e in total time complexity $O(e\sqrt{p})$

4 Geometry

* return $x = CRT(\{x_i \mod p_i^{e_i}\})$

If $n = \prod_{i=1}^r p_i^{e_i}$:

Intersection of 2 lines

```
Pt LLIntersect(Line a, Line b) {
  Pt p1 = a.s, p2 = a.e, q1 = b.s, q2 = b.e;
ld f1 = (p2-p1)^{(q1-p1)}, f2 = (p2-p1)^{(p1-q2)}, f;
   if(dcmp(f=f1+f2) == 0)
     return dcmp(f1)?Pt(NAN,NAN):Pt(INFINITY,INFINITY);
  return q1*(f2/f) + q2*(f1/f);
```

* for each i=1 $\underset{e_i}{\sim} r$:

1. let $g_i=g^{n/p_i}$ having order $p_i^{e_i}$, $h_i=h^{n/p_i^{e_i}}$ where $h_i\in\langle g_i\rangle$.

2. find x_i s.t. $g_i^{x_i}=h_i$ using above algorithm.

4.2 halfPlaneIntersection

```
// for point or line solution, change > to >=
bool onleft(Line L, Pt p) {
  return dcmp(L.v^(p-L.s)) > 0;
// assume that Lines intersect
vector<Pt> HPI(vector<Line>& L) {
  sort(L.begin(), L.end()); // sort by angle
int n = L.size(), fir, las;
Pt *p = new Pt[n];
  Line *q = new Line[n];
  q[fir=las=0] = L[0];
for(int i = 1; i < n; i++) {</pre>
    while(fir < las && !onleft(L[i], p[las-1])) las--;</pre>
    while(fir < las && !onleft(L[i], p[fir])) fir++;</pre>
     q[++las] = L[i];
     if(dcmp(q[las].v^q[las-1].v) == 0) {
       las-
       if(onleft(q[las], L[i].s)) q[las] = L[i];
     if(fir < las) p[las-1] = LLIntersect(q[las-1], q[las</pre>
          ]);
  while(fir < las && !onleft(q[fir], p[las-1])) las--;</pre>
  if(las-fir <= 1) return {};</pre>
  p[las] = LLIntersect(q[las], q[fir]);
  \frac{1}{1} int m = 0;
  vector<Pt> ans(las-fir+1);
  for(int i = fir ; i <= las ; i++) ans[m++] = p[i];</pre>
  return ans;
```

4.3 Intersection of 2 segments

```
bool onseg(Pt p, Line L) {
  Pt x = L.s-p, y = L.e-p;
  return dcmp(x^y) == 0 && dcmp(x*y) <= 0;
} // inseg: dcmp(x^y)==0&&dcmp(x*y)<0

// assume a.s != a.e != b.s != b.e
Pt SSIntersect(Line a, Line b) {
  Pt p = LLIntersect(a, b);
  if(isinf(p.x)&&(onseg(a.s,b)||onseg(a.e,b)||onseg(b.s, a)||onseg(b.e,a))) return p; // overlap
  if(isfinite(p.x)&&onseg(p,a)&&onseg(p,b)) return p;
  return {NAN,NAN}; // non-intersect
}</pre>
```

4.4 Banana

4.5 Intersection of circle and line

```
vector<Pt> CLInter(const Line &a,const Circle &c){
   Pt p=a.s+(c.o-a.s)*a.v/norm2(a.v)*a.v;
   ld d=c.r*c.r-norm2(c.o-p);
   if(d<-eps) return {};
   if(d<eps) return {p};
   Pt v=a.v/norm(a.v)*sqrt(d);
   return {p+v,p-v};
}</pre>
```

4.6 Intersection of polygon and circle

```
|ld PCIntersect(vector<Pt> v, Circle cir) {
```

```
for(int i=0;i<(int)v.size();++i) v[i]=v[i]-cir.o;</pre>
ld ans=0,r=cir.r;
int n=v.size();
for(int i=0;i<n;++i) {</pre>
  Pt pa=v[i],pb=v[(i+1)%n];
  if(norm(pa)<norm(pb)) swap(pa,pb);</pre>
  if(dcmp(norm(pb))==0) continue;
  ld s,h,theta;
  ld a=norm(pb),b=norm(pa),c=norm(pb-pa);
  ld cosB=(pb*(pb-pa))/a/c, B=acos(cosB);
  if(cosB>1) B=0
  else if(cosB<-1) B=PI;</pre>
  ld cosC=(pa*pb)/a/b,C=acos(cosC);
  if(cosC>1) C=0;
  else if(cosC<-1) C=PI;</pre>
  if(a>r){
    s=(C/2)*r*r
    h=a*b*sin(C)/c;
    if(h<r&&B<PI/2) s-=(acos(h/r)*r*r-h*sqrt(r*r-h*h))
  else if(b>r) {
    theta=PI-B-asin(sin(B)/r*a);
    s=0.5*a*r*sin(theta)+(C-theta)/2*r*r;
  else s=0.5*sin(C)*a*b;
  ans+=abs(s)*dcmp(v[i]^v[(i+1)\%n]);
return abs(ans);
```

4.7 Intersection of 2 circles

```
vector<Pt> CCinter(Circle& a, Circle& b){
  Pt o1=a.o,o2=b.o; ld r1=a.r,r2=b.r;
  if(norm(o1-o2)>r1+r2) return {};
  if(norm(o1-o2)<max(r1,r2)-min(r1,r2)) return {};
  ld d2=(o1-o2)*(o1-o2),d=sqrt(d2);
  if(d>r1+r2) return {};
  Pt u=(o1+o2)*0.5+(o1-o2)*((r2*r2-r1*r1)/(2*d2));
  ld A=sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2+d));
  Pt v=Pt(o1.y-o2.y,-o1.x+o2.x)*A/(2*d2);
  return {u+v,u-v};
}
```

4.8 Circle cover

```
#define N 1021
struct CircleCover{
  int C; Circle c[N];
  bool g[N][N], overlap[N][N];
              : area covered by at least i circles
  // Area[i]
  ld Area[N];
  void init(int _C){ C = _C; }
  struct Teve {
    Pt p; ld ang; int add;
Teve() {}
    Teve(Pt \_a, 1d \_b, int \_c):p(\_a), ang(\_b), add(\_c){}
    bool operator<(const Teve &a)const
    {return ang < a.ang;}
  }eve[N * 2];
  // strict: x = 0, otherwise x = -1
bool disjuct(Circle& a, Circle &b, int x)
  {return sign(norm(a.o - b.o) - a.r - b.r) > x;}
  bool contain(Circle& a, Circle &b, int x)
{return sign(a.r - b.r - norm(a.o - b.o)) > x;}
  bool contain(int i, int j){
    contain(c[i], c[j], -1);
  void solve(){
    for(int i = 0; i <= C + 1; i++) Area[i] = 0;
    for(int i = 0; i < C; i++)
for(int j = 0; j < C; j++)</pre>
         overlap[i][j] = contain(i, j);
    for(int i = 0; i < C; i++)
for(int j = 0; j < C; j++)
         g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                       disjuct(c[i], c[j], -1));
    for(int i = 0; i < C; i++){
```

```
int E = 0, cnt = 1;
for(int j = 0; j < C; j++)</pre>
         if(j != i && overlap[j][i])
           cnt++;
       for(int j = 0; j < C; j++)
  if(i != j && g[i][j]){</pre>
           vector<Pt> v=CCinter(c[i], c[j]);
           ld A=atan2(v[0].y - c[i].o.y, v[0].x - c[i].o.
           ld B=atan2(v[1].y - c[i].o.y, v[1].x - c[i].o.
                x);
           eve[E++] = Teve(v[1], B, 1);
           eve[E++] = Teve(v[0], A, -1);
           if(B > A) cnt++;
       if(E == 0) Area[cnt] += pi * c[i].r * c[i].r;
       else{
         sort(eve , eve + E);
         eve[E] = eve[0];
         for(int j = 0; j < E; j++){
            cnt += eve[j].add;
           Area[cnt] += (eve[j].p \wedge eve[j + 1].p) * .5;
           ld theta = eve[j + 1].ang - eve[j].ang;
if (theta < 0) theta += 2. * pi;</pre>
           Area[cnt] +=
              (theta - sin(theta)) * c[i].r*c[i].r * .5;
```

4.9 Li Chao Segment Tree

```
struct LiChao_min{
   struct line{
     ll m, c;
line(ll _m=0, ll _c=0) { m = _m; c = _c; }
   };
   struct node{
     node *l, *r; line f;
     node(line v) { f = v; l = r = NUll; }
   typedef node* pnode;
pnode root; int sz;
#define mid ((l+r)>>1)
  void insert(line &v, int l, int r, pnode &nd){
  if(!nd) { nd = new node(v); return; }
     ll trl = nd->f.eval(l), trr = nd->f.eval(r);
     ll vl = v.eval(l), vr = v.eval(r);
     if(trl <= vl && trr <= vr) return;</pre>
     if(trl > vl && trr > vr) { nd->f = v; return; }
     if(trl > vl) swap(nd->f, v);
     if(nd->f.eval(mid) < v.eval(mid)) insert(v, mid + 1,</pre>
           r, nd->r);
     else swap(nd->f, v), insert(v, l, mid, nd->l);
  il query(int x, int l, int r, pnode &nd){
  if(!nd) return LLONG_MAX;
     if(l == r) return nd->f.eval(x);
     if(mid >= x) return min(nd->f.eval(x), query(x, l,
          mid, nd->1));
     return min(nd->f.eval(x), query(x, mid + 1, r, nd->r
          ));
  /* -sz <= ll query_x <= sz */
  void init(int _sz){ sz = _sz + 1; root = NUll; }
void add_line(ll m, ll c){ line v(m, c); insert(v, -sz
         sz, root); }
  11 query(ll x) { return query(x, -sz, sz, root); }
};
```

4.10 Convex Hull trick

```
/* Given a convexhull, answer querys in O(\lg N)
CH should not contain identical points, the area should
be > 0, min pair(x, y) should be listed first */
double det( const Pt& p1 , const Pt& p2 )
{ return p1.x * p2.y - p1.y * p2.x; }
struct Conv{
  int n;
  vector<Pt> a, upper, lower;
  Conv(vector<Pt> _a) : a(_a){
    n = a.size(); int ptr = 0;
    for(int i=1; i<n; ++i) if (a[ptr] < a[i]) ptr = i;</pre>
```

```
for(int i=0; i<=ptr; ++i) lower.push_back(a[i]);
for(int i=ptr; i<n; ++i) upper.push_back(a[i]);</pre>
  upper.push_back(a[0]);
int sign( ll x ){ // fixed when changed to double
  return x < 0 ? -1 : x > 0; }
pair<ll,int> get_tang(vector<Pt> &conv, Pt vec){
  int l = 0, r = (int)conv.size() - 2;
  while(l + 1 < r){
     int mid = (l + r) / 2;
     if(sign(det(conv[mid+1]-conv[mid],vec))>0)r=mid;
     else l = mid;
  }
  void upd_tang(const Pt &p, int id, int &i0, int &i1){
  if(det(a[i0] - p, a[id] - p) > 0) i0 = id;
  if(det(a[i1] - p, a[id] - p) < 0) i1 = id;</pre>
void bi_search(int l, int r, Pt p, int &i0, int &i1){
  if(l == r) return;
  upd_tang(p, l % n, i0, i1);
int sl=sign(det(a[l % n] - p, a[(l + 1) % n] - p));
  while(l + 1 < r) {
     int mid = (l + r) / 2;
     int smid=sign(det(a[mid%n]-p, a[(mid+1)%n]-p));
     if (smid == sl) l = mid;
     else r = mid;
  upd_tang(p, r % n, i0, i1);
int bi_search(Pt u, Pt v, int l, int r) {
  int sl = sign(det(v - u, a[l % n] - u));
  while(l + 1 < r) 
     int mid = (l + r) / 2;
int smid = sign(det(v - u, a[mid % n] - u));
     if (smid == sl) l = mid;
     else r = mid;
  }
  return 1 % n;
^{-}// 1. whether a given point is inside the CH
bool contain(Pt p) {
  if (p.x < lower[0].x || p.x > lower.back().x) return
  int id = lower_bound(lower.begin(), lower.end(), Pt(
        p.x, -INF)) - lower.begin();
  if (lower[id].x == p.x) {
  if (lower[id].y > p.y) return 0;
}else if(det(lower[id-1]-p,lower[id]-p)<0)return 0;
}else if(det(lower[id-1]-p,lower[id]-p)<0)return 0;</pre>
  id = lower_bound(upper.begin(), upper.end(), Pt(p.x,
  INF), greater<Pt>()) - upper.begin();
if (upper[id].x == p.x) {
  if (upper[id].y < p.y) return 0;
}else if(det(upper[id-1]-p,upper[id]-p)<0)return 0;</pre>
  return 1;
// 2. Find 2 tang pts on CH of a given outside point
// return true with i0, i1 as index of tangent points
// return false if inside CH
bool get_tang(Pt p, int &i0, int &i1) {
  if (contain(p)) return false;
  i0 = i1 = 0;
  int id = lower_bound(lower.begin(), lower.end(), p)
        lower.begin();
  bi_search(0, id, p, i0, i1);
bi_search(id, (int)lower.size(), p, i0, i1);
  id = lower_bound(upper.begin(), upper.end(), p,
       greater<Pt>()) - upper.begin();
  bi_search((int)lower.size() - 1, (int)lower.size() -
         1 + id, p, i0, i1);
  bi_search((int)lower.size() - 1 + id, (int)lower.
       size() - 1 + (int)upper.size(), p, i0, i1);
  return true;
// 3. Find tangent points of a given vector
// ret the idx of vertex has max cross value with vec
int get_tang(Pt vec){
  pair<ll, int> ret = get_tang(upper, vec);
  ret.second = (ret.second+(int)lower.size()-1)%n;
  ret = max(ret, get_tang(lower, vec));
```

```
return ret.second;
}
// 4. Find intersection point of a given line
// return 1 and intersection is on edge (i, next(i))
// return 0 if no strictly intersection
bool get_intersection(Pt u, Pt v, int &i0, int &i1){
  int p0 = get_tang(u - v), p1 = get_tang(v - u);
  if(sign(det(v-u,a[p0]-u))*sign(det(v-u,a[p1]-u))<0){
    if (p0 > p1) swap(p0, p1);
    i0 = bi_search(u, v, p0, p1);
    i1 = bi_search(u, v, p1, p0 + n);
    return 1;
}
return 0;
}
```

4.11 Tangent line of two circles

```
vector<Line> go(const Circle& c1, const Circle& c2, int
  // sign1 = 1 for outer tang, -1 for inter tang
  vector<Line> ret:
  double d_{sq} = norm2(c1.o - c2.o);
  if(d_sq < eps) return ret;</pre>
  double d = sqrt(d_sq);
  Pt v = (c2.0 - c1.0) / d;
  double c = (c1.r - sign1 * c2.r) / d;
if(c * c > 1) return ret;
double h = sqrt(max(0.0, 1.0 - c * c));
for(int sign2 = 1; sign2 >= -1; sign2 -= 2){
    Pt n = { v.x * c - sign2 * h * v.y ,
                 v.y * c + sign2 * h * v.x };
     Pt p1 = c1.o + n * c1.r;
     Pt p2 = c2.0 + n * (c2.r * sign1);
     if(fabs(p1.x - p2.x) < eps and
         fabs(p1.y - p2.y) < eps)
        p2 = p1 + perp(c2.o - c1.o);
     ret.push_back({p1, p2});
  return ret;
```

4.12 Tangent line of point and circle

```
vector<Line> PCTangent(const Circle& C, const Pt& P) {
  vector<Line> ans;
  Pt u = C.o - P;
  double dist = norm(u);
  if(dist < C.r) return ans;
  else if(abs(dist) < eps) {
    ans.push_back({P, P+rotate(u, M_PI/2)});
    return ans;
  }
  else {
    double ang = asin(C.r/dist);
    ans.push_back({P, P+rotate(u, -ang)});
    ans.push_back({P, P+rotate(u, +ang)});
    return ans;
  }
}</pre>
```

4.13 Min distance of two convex

4.14 Poly Union

```
struct PY{
  int n; Pt pt[5]; double area;
  Pt& operator[](const int x){ return pt[x]; }
  void init(){ //n,pt[0~n-1] must be filled
    area=pt[n-1]^pt[0];
    for(int i=0;i<n-1;i++) area+=pt[i]^pt[i+1];</pre>
    if((area/=2)<0)reverse(pt,pt+n),area=-area;</pre>
PY py[500];
pair<double,int> c[5000];
inline double segP(Pt &p,Pt &p1,Pt &p2){
 if (dcmp(p1.x-p2.x)==0) return (p.y-p1.y)/(p2.y-p1.y);
 return (p.x-p1.x)/(p2.x-p1.x);
double polyUnion(int n){ //py[0~n-1] must be filled
  int i,j,ii,jj,ta,tb,r,d;
  double z,w,s,sum,tc,td;
  for(i=0;i<n;i++) py[i][py[i].n]=py[i][0];</pre>
  sum=0:
  for(i=0;i<n;i++){</pre>
    for(ii=0;ii<py[i].n;ii++){</pre>
      r=0;
      c[r++]=make_pair(0.0,0);
      c[r++]=make_pair(1.0,0);
      for(j=0;j<n;j++){</pre>
        tb=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj+1])
              );
          if(ta==0 \&\& tb==0){
            if((py[j][jj+1]-py[j][jj])*(py[i][ii+1]-py[i
              ][ii])>0 && j<i){
c[r++]=make_pair(segP(py[j][jj],py[i][ii],
                   py[i][ii+1]),1);
              c[r++] = make\_pair(segP(py[j][jj+1],py[i][ii])
                  ],py[i][ii+1]),-1);
          }else if(ta>=0 && tb<0){</pre>
            tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
            td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
            c[r++]=make\_pair(tc/(tc-td),1);
          }else if(ta<0 && tb>=0)
            tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
            td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
            c[r++]=make_pair(tc/(tc-td),-1);
        }
      sort(c,c+r);
      z=min(max(c[0].first,0.0),1.0);
      d=c[0].second; s=0;
      for(j=1; j<r; j++)</pre>
        w=min(max(c[j].first,0.0),1.0);
        if(!d)_s+=w-z
        d+=c[j].second; z=w;
      sum+=(py[i][ii]^py[i][ii+1])*s;
    }
  return sum/2;
```

4.15 Lower Concave Hull

```
const ll is_query = -(1LL<<62);
struct Line {
    ll m, b;
    mutable function<const Line*()> succ;
    bool operator<(const Line& rhs) const {
        if (rhs.b != is_query) return m < rhs.m;
        const Line* s = succ();
        return s ? b - s->b < (s->m - m) * rhs.m : 0;
    }
}; // maintain upper hull for maximum
struct HullDynamic : public multiset<Line> {
    bool bad(iterator y) {
        auto z = next(y);
        if (y == begin()) {
```

```
if (z == end()) return 0;
      return y->m == z->m && y->b <= z->b;
    auto x = prev(y);
    if(z==end())return y->m==x->m&y->b<=x->b;
    return (x->b-y->b)*(z->m-y->m)>=
            (y->b-z->b)*(y->m-x->m);
  void insert_line(ll m, ll b) {
    auto y = insert({m, b});
    y->succ = [=]{return next(y)==end()?0:&*next(y);};
    if(bad(y)) {erase(y); return; }
    while(next(y)!=end()&&bad(next(y)))erase(next(y));
    while(y!=begin()&&bad(prev(y)))erase(prev(y));
  ll eval(ll x) {
    auto l = *lower_bound((Line) {x, is_query});
    return l.m * x + l.b;
};
```

```
4.16
         Delaunay Triangulation
/* Delaunay Triangulation:
Given a sets of points on 2D plane, find a
triangulation such that no points will strictly
inside circumcircle of any triangle.
find : return a triangle contain given point
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
Region of triangle u: iterate each u.edge[i].tri, each points are u.p[(i+1)%3], u.p[(i+2)%3]
Voronoi diagram: for each triangle in triangulation,
the bisector of all its edges will split the region.
nearest point will belong to the triangle containing it
const int N = 100000 + 5;
const type inf = 2e3;
type eps = 1e-6; // 0 when integer
type sqr(type x) { return x*x; }
// return p4 is in circumcircle of tri(p1,p2,p3)
bool in_cc(const Pt& p1, const Pt& p2, const Pt& p3,
     const Pt& p4){
  type u11 = p1.x - p4.x; type u12 = p1.y - p4.y;
  type u21 = p2.x - p4.x; type u22 = p2.y - p4.y;
type u31 = p3.x - p4.x; type u32 = p3.y - p4.y;
type u33 = sqr(p1.x)-sqr(p4.x)+sqr(p1.y)-sqr(p4.y);
  type u23 = sqr(p2.x)-sqr(p4.x)+sqr(p2.y)-sqr(p4.y);
  type u33 = sqr(p3.x)-sqr(p4.x)+sqr(p3.y)-sqr(p4.y);
type det = -u13*u22*u31 + u12*u23*u31 + u13*u21*u32
                -u11*u23*u32 - u12*u21*u33 + u11*u22*u33;
  return det > eps;
type side(const Pt& a, const Pt& b, const Pt& p)
{ return (b - a) ^ (p - a); }
typedef int SdRef;
struct Tri;
typedef Tri* TriRef;
struct Edge {
  TriRef tri; SdRef side;
  Edge():tri(0), side(0){}
  Edge(TriRef _tri, SdRef _side):tri(_tri), side(_side)
struct Tri {
  Pt p[3];

  Edge edge[3];
  TriRef chd[3];
  Tri() {}
  Tri(const Pt& p0, const Pt& p1, const Pt& p2) {
  p[0] = p0; p[1] = p1; p[2] = p2;
  chd[0] = chd[1] = chd[2] = 0;
  bool has_chd() const { return chd[0] != 0; }
  int num_chd() const {
    return chd[0] == 0? 0
           : chd[1] == 0 ? 1
           : chd[2] == 0 ? 2 : 3;
```

```
bool contains(Pt const& q) const { for( int i = 0 ; i < 3 ; i ++ ) if( side(p[i], p[(i + 1) \% 3] , q) < -eps )
         return false;
    return true;
} pool[ N * 10 ], *tris;
void edge( Edge a, Edge b ){
  if(a.tri) a.tri->edge[a.side] = b;
  if(b.tri) b.tri->edge[b.side] = a;
struct Trig { // Triangulation
  Trig(){
     the_root = // Tri should at least contain all points
       new(tris++)Tri(Pt(-inf,-inf),Pt(+inf+inf,-inf),Pt
             (-inf,+inf+inf));
  TriRef find(Pt p)const{ return find(the_root,p); }
  void add_point(const Pt& p){ add_point(find(the_root,p)
       ),p); }
  TriRef the_root;
  static TriRef find(TriRef root, const Pt& p) {
    while( true ){
       if( !root->has_chd() )
         return root:
       for( int i = 0; i < 3 && root->chd[i] ; ++i )
  if (root->chd[i]->contains(p)) {
            root = root->chd[i];
            break;
    assert( false ); // "point not found"
  void add_point(TriRef root, Pt const& p) {
    TriRef tab, tbc, tca;
     /* split it into three triangles */
     tab=new(tris++) Tri(root->p[0],root->p[1],p);
    tbc=new(tris++) Tri(root->p[1],root->p[2],p);
tca=new(tris++) Tri(root->p[2],root->p[0],p);
edge(Edge(tab,0), Edge(tbc,1));
     edge(Edge(tbc,0), Edge(tca,1));
     edge(Edge(tca,0), Edge(tab,1))
    edge(Edge(tab,2), root->edge[2]);
edge(Edge(tbc,2), root->edge[0]);
     edge(Edge(tca,2), root->edge[1]);
     root->chd[0] = tab;
     root->chd[1] = tbc;
    root->chd[2] = tca;
     flip(tab,2);
     flip(tbc,2);
     flip(tca,2);
  void flip(TriRef tri, SdRef pi) {
    TriRef trj = tri->edge[pi].tri;
     int pj = tri->edge[pi].side;
    if (!trj) return;
     if (!in_cc(tri->p[0],tri->p[1],tri->p[2],trj->p[pj])
          ) return;
      * flip edge between tri,trj */
     TriRef trk = new(tris++) Tri(tri->p[(pi+1)%3], trj->
          p[pj], tri->p[pi]);
     TriRef trl = new(tris++) Tri(trj->p[(pj+1)%3], tri->
    p[pi], trj->p[pj]);
edge(Edge(trk,0), Edge(trl,0));
edge(Edge(trk,1), tri->edge[(pi+2)%3]);
edge(Edge(trk,2), trj->edge[(pj+1)%3]);
    edge(Edge(trl,1), trj->edge[(pj+2)%3]);
edge(Edge(trl,2), tri->edge[(pi+1)%3]);
tri->chd[0]=trk; tri->chd[1]=trl; tri->chd[2]=0;
    trj->chd[0]=trk; trj->chd[1]=trl; trj->chd[2]=0;
    flip(trk,1); flip(trk,2);
flip(trl,1); flip(trl,2);
vector<TriRef> triang; // vector of all triangle
set<TriRef> vst;
void go( TriRef now ){ // store all tri into triang
  if( vst.find( now ) != vst.end() )
    return
  vst.insert( now );
  if( !now->has_chd() ){
    triang.push_back( now );
```

```
return;
}
for( int i = 0 ; i < now->num_chd() ; i ++ )
    go( now->chd[ i ] );
}
void build( int n , Pt* ps ){ // build triangulation
    tris = pool; triang.clear(); vst.clear();
    random_shuffle(ps, ps + n);
    Trig tri; // the triangulation structure
    for(int i = 0; i < n; ++ i)
        tri.add_point(ps[i]);
    go( tri.the_root );
}</pre>
```

4.17 Min Enclosing Circle

```
struct Mec{
  // return pair of center and r
  static const int N = 101010;
  int n;
  Pt p[N], cen;
  double r2;
  void init( int _n , Pt _p[] ){
    memcpy( p , _p , sizeof(Pt) * n );
  double sqr(double a){ return a*a; }
  Pt center(Pt p0, Pt p1, Pt p2) {
    Pt a = p1-p0;
    Pt b = p2-p0;
    double c1=norm2( a ) * 0.5;
     double c2=norm2( b ) * 0.5;
     double d = a \wedge b;
    double x = p0.x + (c1 * b.y - c2 * a.y) / d;
     double y = p0.y + (a.x * c2 - b.x * c1) / d;
    return Pt(x,y);
  pair<Pt,double> solve(){
    random_shuffle(p,p+n);
     for (int i=0; i<n; i++){</pre>
       if (norm2(cen-p[i]) <= r2) continue;</pre>
       cen = p[i];
       r2 = 0;
       for (int j=0; j<i; j++){
  if (norm2(cen-p[j]) <= r2) continue;
  cen=Pt((p[i].x+p[j].x)/2,(p[i].y+p[j].y)/2);</pre>
         r2 = norm2(cen-p[j]);
for (int k=0; k<j; k++){
    if (norm2(cen-p[k]) <= r2) continue;
            cen = center(p[i],p[j],p[k]);
            r2 = norm2(cen-p[k]);
    return {cen,sqrt(r2)};
} mec;
```

4.18 Min Enclosing Ball

```
#define N 202020
int n, nouter; Pt pt[ N ], outer[4], res;
double radius,tmp;
double det(double m[3][3]){
    return m[0][0]*m[1][1]*m[2][2]
          + m[0][1]*m[1][2]*m[2][0]
+ m[0][2]*m[2][1]*m[1][0]
- m[0][2]*m[1][1]*m[2][0]
          - m[0][1]*m[1][0]*m[2][2]
           - m[0][0]*m[1][2]*m[2][1];
void ball() {
  Pt q[3]; double m[3][3], sol[3], L[3], d;
  int i,j;
            res.x = res.y = res.z = radius = 0;
  switch ( nouter ) {
    case 1: res=outer[0]; break;
case 2: res=(outer[0]+outer[1])/2; radius=norm2(res,
           outer[0]); break;
    case 3:
       for(i=0; i<2; ++i) q[i]=outer[i+1]-outer[0];</pre>
```

```
for(i=0; i<2; ++i) for(j=0; j<2; ++j) m[i][j]=(q[i
      ] * q[j])*2;
for(i=0; i<2; ++i) sol[i]=(q[i] * q[i]);
      if(fabs(d=m[0][0]*m[1][1]-m[0][1]*m[1][0])<eps)</pre>
      L[0]=(sol[0]*m[1][1]-sol[1]*m[0][1])/d;
L[1]=(sol[1]*m[0][0]-sol[0]*m[1][0])/d;
      res=outer[0]+q[0]*L[0]+q[1]*L[1];
      radius=norm2(res, outer[0]);
      break;
    case 4:
      for(i=0; i<3; ++i) q[i]=outer[i+1]-outer[0], sol[i
    ]=(q[i] * q[i]);</pre>
      for(i=0;i<3;++i) for(j=0;j<3;++j) m[i][j]=(q[i] *
           q[j])*2;
      d=det(m);
      if(fabs(d)<eps) return;</pre>
      for(j=0; j<3; ++j)
         for(i=0; i<3; ++i) m[i][j]=(q[i] * q[j])*2;</pre>
      } res=outer[0];
      for(i=0; i<3; ++i ) res = res + q[i] * L[i];
radius=norm2(res, outer[0]);</pre>
void minball(int n){ ball();
  if(nouter < 4) for(int i = 0; i < n; i ++)
    if(norm2(res, pt[i]) - radius > eps){
      outer[nouter ++] = pt[i]; minball(i); --nouter;
      if(i>0){ Pt Tt = pt[i];
  memmove(&pt[1], &pt[0], sizeof(Pt)*i); pt[0]=Tt;
double solve(){
 // n points in pt
  random_shuffle(pt, pt+n); radius=-1;
  for(int i=0;i<n;i++) if(norm2(res,pt[i])-radius>eps)
    nouter=1, outer[0]=pt[i], minball(i);
  return sqrt(radius);
```

4.19 Minkowski sum

```
vector<Pt> minkowski(vector<Pt> p, vector<Pt> q){
 int n = p.size() , m = q.size();
 Pt c = Pt(0, 0);
  for( int i = 0; i < m; i ++) c = c + q[i];
 c = c / m;
  for( int i = 0; i < m; i ++) q[i] = q[i] - c;
 int cur = -1;
for( int i = 0; i < m; i ++)</pre>
    if( (q[i] ^ (p[0] - p[n-1])) > -eps)
     if( cur == -1 || (q[i] \( (p[0] - p[n-1])) > (q[cur] \( (p[0] - p[n-1])) \)
        cur = i;
 vector<Pt> h;
 p.push_back(p[0]);
 for( int i = 0; i < n; i ++)
   while( true ){
     h.push_back(p[i] + q[cur]);
     int nxt = (cur + 1 == m ? 0 : cur + 1);
     else break;
 for (auto &&i : h) i = i + c;
 return convex_hull(h);
```

4.20 Min dist on Cuboid

4.21 Heart of Triangle

```
Pt inCenter( Pt &A, Pt &B, Pt &C) { // 内心 double a = norm(B-C), b = norm(C-A), c = norm(A-B); return (A * a + B * b + C * c) / (a + b + c); }
Pt circumCenter( Pt &a, Pt &b, Pt &c) { // 外心 Pt bb = b - a, cc = c - a; double db=norm2(bb), dc=norm2(cc), d=2*(bb ^ cc); return a-Pt(bb.Y*dc-cc.Y*db, cc.X*db-bb.X*dc) / d; }
Pt othroCenter( Pt &a, Pt &b, Pt &c) { // 垂心 Pt ba = b - a, ca = c - a, bc = b - c; double Y = ba.Y * ca.Y * bc.Y, A = ca.X * ba.Y - ba.X * ca.Y, x0= (Y+ca.X*ba.Y*b.X-ba.X*ca.Y*c.X) / A, y0= -ba.X * (x0 - c.X) / ba.Y + ca.Y; return Pt(x0, y0); }
```

5 Graph

5.1 DominatorTree

```
const int MAXN = 100010;
struct DominatorTree{
#define REP(i,s,e) for(int i=(s);i<=(e);i++)</pre>
#define REPD(i,s,e) for(int i=(s);i>=(e);i--)
  int n , m , s;
  vector< int > g[ MAXN ] , pred[ MAXN ];
vector< int > cov[ MAXN ];
int dfn[ MAXN ] , nfd[ MAXN ] , ts;
  int par[ MAXN ];
int sdom[ MAXN ] , idom[ MAXN ];
int mom[ MAXN ] , mn[ MAXN ];
inline bool cmp( int u , int v )
  { return dfn[ u ] < dfn[ v ]; } int eval( int u ){
     if( mom[ u ] == u ) return u;
     int res = eval( mom[ u ] );
if(cmp( sdom[ mn[ mom[ u ] ] ] , sdom[ mn[ u ] ] ))
        mn[ u ] = mn[ mom[ u ] ];
     return mom[ u ] = res;
  void init( int _n , int _m , int _s ){
     ts = 0; n = _n; m = _m; s = _s;
REP( i, 1, n ) g[ i ].clear(), pred[ i ].clear();
  void addEdge( int u , int v ){
  g[ u ].push_back( v );
  pred[ v ].push_back( u );
  void dfs( int u ){
     dfn['u ] = ts;
nfd[ ts ] = u;
     for( int v : g[ u ] ) if( dfn[ v ] == 0 ){
        par[ v ] = u;
        dfs( v );
  void build(){
     REP( i , 1 , n ){
   dfn[ i ] = nfd[ i ] = 0;
        cov[ i ].clear();
        mom[i] = mn[i] = sdom[i] = i;
```

```
dfs( s );
     REPD( i , n , 2 ){
  int u = nfd[ i ];
       if( u == 0 ) continue ;
for( int v : pred[ u ] ) if( dfn[ v ] ){
          eval( v );
          if( cmp( sdom[ mn[ v ] ] , sdom[ u ] ) )
  sdom[ u ] = sdom[ mn[ v ] ];
       cov[ sdom[ u ] ].push_back( u );
mom[ u ] = par[ u ];
       for( int w : cov[ par[ u ] ] ){
          eval( w );
          if( cmp( sdom[ mn[ w ] ] , par[ u ] ) )
            idom[w] = mn[w];
          else idom[ w ] = par[ u ];
       cov[ par[ u ] ].clear();
     REP( i , 2 , n ){
       int u = nfd[ i ];
       if( u == 0 ) continue ;
       if( idom[ u ] != sdom[ u ] )
          idom[u] = idom[idom[u]];
     }
} domT;
```

5.2 Directed MST(ElogE)

```
struct RollbackUF {
  vi e; vector<pii> st;
  RollbackUF(int n) : e(n, -1) {}
int size(int x) {return -e[find(x)]; }
  int find(int x) {return e[x] < 0? x : find(e[x]);}
  int time() {return st.size();}
  void rollback(int t) {
    for(int i = time(); i --> t;)
      e[st[i].first] = st[i].second;
    st.resize(t);
  bool join(int a, int b)
    a = find(a), b = find(b);
    if(a == b) return false;
    if(e[a] > e[b]) swap(a, b);
    st.push_back({a, e[a]}); st.push_back({b, e[b]});
    e[a] += e[b]; e[b] = a;
    return true;
};
struct Edge {int a, b; ll w;};
struct Node { // lazy skew heap node 
Edge key; Node *1, *r; ll d;
  void prop() {
    key.w+=d; if(l) l->d+=d; if(r) r->d+=d; d=0;
  Edge top() {prop(); return key;}
Node *merge(Node *a, Node *b) {
  if(!a | | !b) return a ?: b;
  a->prop(), b->prop();
  if(a\rightarrow key.w > b\rightarrow key.w) swap(a, b):
  swap(a->l, (a->r = merge(b, a->r)));
  return a;
void pop(Node*& a) {a->prop(); a=merge(a->l, a->r);}
pair<ll, vi> dmst(int n, int r, vector<Edge>& g) {
  RollbackUF uf(n); vector<Node*> pq(n);
  for(Edge e:g) pq[e.b]=merge(pq[e.b], new Node{e});
  ll res = 0; vi seen(n, -1), path(n), par(n);
  seen[r] = r;
  vector<Edge> Q(n), in(n, \{-1,-1\});
  deque<tuple<int, int, vector<Edge>>> cycs;
  rep(s,0,n) {
    int u = s, qi = 0, w;
while(seen[u] < 0) {</pre>
       Edge e = pq[u]->top();
pq[u]->d -= e.w, pop(pq[u]);
      Q[qi] = e, path[qi++] = u, seen[u] = s;
      res += e.w, u = uf.find(e.a);
      if(seen[u] == s) { // found cycle, contract
```

```
Node* cyc = 0; int end = qi, t = uf.time();
    do cyc = merge(cyc, pq[w = path[--qi]]);
    while(uf.join(u, w));
    u = uf.find(u), pq[u] = cyc, seen[u] = -1;
        cycs.push_front({u, t, {&Q[qi], &Q[end]}});
    }
    rep(i,0,qi) in[uf.find(Q[i].b)] = Q[i];
}
for(auto& [u,t,comp] : cycs) {// restore sol
    uf.rollback(t); Edge inEdge = in[u];
    for(auto& e : comp) in[uf.find(e.b)] = e;
    in[uf.find(inEdge.b)] = inEdge;
}
rep(i,0,n) par[i] = in[i].a;
return {res, par};
}
```

5.3 MaximalClique

```
#define N 80
struct MaxClique{ // 0-base
  typedef bitset<N> Int;
  Int lnk[N] , v[N];
  int n;
  void init(int _n){
    n = _n;
    for(int i = 0; i < n; i ++){
      lnk[i].reset(); v[i].reset();
  void addEdge(int a , int b)
{ v[a][b] = v[b][a] = 1; }
  int ans , stk[N], id[N] , di[N] , deg[N];
  Int cans:
  void dfs(int elem_num, Int candi, Int ex){
    if(candi.none()&ex.none()){
      cans.reset();
      for(int i = 0)
                        i < elem_num ; i ++)
        cans[id[stk[i]]] = 1;
      ans = elem_num; // cans is a maximal clique
      return;
    int pivot = (candilex)._Find_first();
    Int smaller_candi = candi & (~lnk[pivot]);
    while(smaller_candi.count()){
      int nxt = smaller_candi._Find_first();
       candi[nxt] = smaller_candi[nxt] = 0;
      ex[nxt] = 1;
      stk[elem_num] = nxt;
      dfs(elem_num+1,candi&lnk[nxt],ex&lnk[nxt]);
    }
  int solve(){
    for(int i = 0; i < n; i + +){
      id[i] = i; deg[i] = v[i].count();
    sort(id , id + n_, [\&](int id1, int id2){
           return deg[id1] > deg[id2]; });
    for(int i = 0; i < n; i + +) di[id[i]] = i;
    for(int i = 0 ; i < n ; i ++)</pre>
      for(int j = 0; j < n; j ++)
  if(v[i][j]) lnk[di[i]][di[j]] = 1;</pre>
    ans = 1; cans.reset(); cans[0] = 1;
dfs(0, Int(string(n,'1')), 0);
    return ans;
} solver;
```

5.4 MaxCliqueDyn

```
#define N 150
struct MaxClique{ // Maximum Clique
  bitset<N> a[N],cs[N];
  int ans,sol[N],q,cur[N],d[N],n;
  void init(int _n){
    n=_n; for(int i=0;i<n;i++) a[i].reset();
  }
  void addEdge(int u,int v){a[u][v]=a[v][u]=1;}
  void csort(vector<int> &r,vector<int> &c){
    int mx=1,km=max(ans-q+1,1),t=0,m=r.size();
    cs[1].reset(); cs[2].reset();
```

```
for(int i=0;i<m;i++){
  int p=r[i],k=1;</pre>
       while((cs[k]&a[p]).count()) k++;
       if(k>mx){ mx++; cs[mx+1].reset();}
       cs[k][p]=1;
       if(k<km) r[t++]=p;</pre>
     c.resize(m);
     if(t) c[t-1]=0;
     for(int k=km;k<=mx;k++){</pre>
       for(int p=cs[k]._Find_first();p<N;p=cs[k].</pre>
            _Find_next(p)){
         r[t]=p; c[t]=k; t++;
    }
  }
  void dfs(vector<int> &r,vector<int> &c,int 1,bitset<N>
    while(!r.empty()){
       int p=r.back(); r.pop_back(); mask[p]=0;
       if(q+c.back()<=ans) return;</pre>
       cur[q++]=p;
       vector<int> nr,nc; bitset<N> nmask=mask&a[p];
       for(int i:r) if(a[p][i]) nr.push_back(i);
       if(!nr.empty()){
         if(1<4){
            for(int i:nr) d[i]=(a[i]&nmask).count();
           sort(nr.begin(),nr.end(),[&](int x,int y){
                return d[x]>d[y];});
         csort(nr,nc); dfs(nr,nc,l+1,nmask);
       else if(q>ans){
         ans=q; copy(cur,cur+q,sol);
       c.pop_back(); q--;
    }
  int solve(bitset<N> mask=bitset<N>(string(N,'1'))){ //
        vertex mask
    vector<int> r,c; ans=q=0;
    for(int i=0;i<n;i++) if(mask[i]) r.push_back(i);
for(int i=0;i<n;i++) d[i]=(a[i]&mask).count();</pre>
     sort(r.begin(),r.end(),[&](int i,int j){return d[i]>
         d[j];});
    csort(r,c); dfs(r,c,1,mask);
return ans; // sol[0 ~ ans-1]
}graph;
```

5.5 Strongly Connected Component

```
void dfs(int i){
    V[i]=low[i]=++ts,stk[top++]=i,instk[i]=1;
    for(auto x:E[i]){
        if(!V[x])dfs(x),low[i]=min(low[i],low[x]);
        else if(instk[x])low[i]=min(low[i],V[x]);
    }
    if(V[i]==low[i]){
        int j;
        do{j = stk[--top], instk[j] = 0, scc[j] = i;
        }while(j != i);
    }
}
```

5.6 Dynamic MST

```
/* Dynamic MST O( Q lg^2 Q )
n nodes, m edges, Q query
(u[i], v[i], w[i])->edge
(qid[i], qw[i])->chg weight of edge No.qid[i] to qw[i]
delete an edge: (i, \infty)
add an edge: change from \infty to specific value */
const int SZ=M+3*MXQ;
int a[N],*tz;
int find(int x){
   return x==a[x]?x:a[x]=find(a[x]);
}
bool cmp(int aa,int bb){ return tz[aa]<tz[bb]; }
int kx[N],ky[N],kt, vd[N],id[M], app[M], cur;
long long answer[MXQ]; // answer after ith query
bool extra[M];</pre>
```

```
if(Q==1){
    for(int i=1;i<=n;i++) a[i]=0;</pre>
    z[ qx[0] ]=qy[0]; tz = z;
for(int i=0;i<m1;i++) id[i]=i;
    sort(id,id+m1,cmp); int ri,rj;
    for(int i=0;i<m1;i++){</pre>
      ri=find(x[id[i]]); rj=find(y[id[i]]);
      if(ri!=rj){ ans+=z[id[i]]; a[ri]=rj; }
        answer[cur++]=ans;
    return;
  int ri,rj;
  //contract
  kt=0;
  for(int i=1;i<=n;i++) a[i]=0;</pre>
  for(int i=0;i<Q;i++){</pre>
    ri=find(x[qx[i]]); rj=find(y[qx[i]]); if(ri!=rj) a[
        ri]=rj;
  }
  int tm=0;
  for(int i=0;i<m1;i++) extra[i]=true;</pre>
  for(int i=0;i<Q;i++) extra[ qx[i] ]=false;</pre>
  for(int i=0;i<m1;i++) if(extra[i]) id[tm++]=i;</pre>
  tz=z; sort(id,id+tm,cmp);
  for(int i=0;i<tm;i++){</pre>
    ri=find(x[id[i]]); rj=find(y[id[i]]);
    if(ri!=rj){
      a[ri]=rj; ans += z[id[i]];
      k\bar{x}[k\bar{t}]=\bar{x}[id[i]]; k\bar{y}[k\bar{t}]=\bar{y}[id[i]]; kt++;
    }
  for(int i=1;i<=n;i++) a[i]=0;</pre>
  for(int i=0;i<kt;i++) a[ find(kx[i]) ]=find(ky[i]);</pre>
  int n2=0;
  for(int i=1;i<=n;i++) if(a[i]==0)</pre>
  vd[i]=++n2;
  for(int i=1;i<=n;i++) if(a[i])</pre>
  vd[i]=vd[find(i)];
  int m2=0, *Nx=x+m1, *Ny=y+m1, *Nz=z+m1;
  for(int i=0;i<m1;i++) app[i]=-1;</pre>
  for(int i=0;i<Q;i++) if(app[qx[i]]==-1){</pre>
    app[qx[i]]=m2; m\bar{2}++;
  for(int i=0;i<Q;i++){ z[ qx[i] ]=qy[i]; qx[i]=app[qx[i</pre>
      ]]; }
  for(int i=1;i<=n2;i++) a[i]=0;</pre>
  for(int i=0;i<tm;i++){</pre>
    ri=find(vd[ x[id[i]] ]); rj=find(vd[ y[id[i]] ]);
    if(ri!=rj){
      a[ri]=rj; Nx[m2]=vd[ x[id[i]] ];
      Ny[m2]=vd[ y[id[i]] ]; Nz[m2]=\bar{z}[id[i]]; m2++;
    }
  int mid=Q/2;
  solve(qx,qy,mid,n2,Nx,Ny,Nz,m2,ans);
  solve(qx+mid,qy+mid,Q-mid,n2,Nx,Ny,Nz,m2,ans);
int u[SZ],v[SZ],w[SZ],qid[MXQ],qw[MXQ],n,m,Q;
void work(){if(Q) cur=0,solve(qid,qw,Q,n,u,v,w,m,0);}
```

5.7 Maximum General graph Matching

```
// should shuffle vertices and edges
const int N = 100005, E = (2e5) * 2 + 40;
struct Graph{
  int to[E],bro[E],head[N],e;
  int lnk[N],vis[N],stp,n;
  void init( int _n ){
    stp = 0; e = 1; n = _n;
    for( int i = 1; i <= n; i ++ )
        head[i] = lnk[i] = vis[i] = 0;
}
void add_edge(int u,int v){
    to[e]=v,bro[e]=head[u],head[u]=e++;
    to[e]=u,bro[e]=head[v],head[v]=e++;
}
bool dfs(int x){</pre>
```

```
vis[x]=stp;
    for(int i=head[x];i;i=bro[i]){
       int v=to[i];
      if(!lnk[v]){
         lnk[x]=v, lnk[v]=x;
         return true;
    } for(int i=head[x];i;i=bro[i]){
      int v=to[i];
      if(vis[lnk[v]]<stp){</pre>
         int w=lnk[v]
         lnk[x]=v, lnk[v]=x, lnk[w]=0;
         if(dfs(w)) return true;
         lnk[w]=v, lnk[v]=w, lnk[x]=0;
      }
    }
    return false;
  int solve(){
    int ans = 0;
    for(int i=1;i<=n;i++) if(!lnk[i])</pre>
         stp++, ans += dfs(i);
    return ans;
  }
} graph;
```

5.8 Minimum General Weighted Matching

```
struct Graph {
  // Minimum General Weighted Matching (Perfect Match)
  static const int MXN = 105;
  int n, edge[MXN][MXN];
  int match[MXN],dis[MXN],onstk[MXN];
  vector<int> stk;
  void init(int _n) {
    n = _n;
for( int i = 0 ; i < n ; i ++ )</pre>
      for( int j = 0 ; j < n ; j ++ )
  edge[ i ][ j ] = 0;</pre>
  void add_edge(int u, int v, int w)
  \{ edge[u][v] = edge[v][u] = w; \}
  bool SPFA(int u){
    if (onstk[u]) return true;
    stk.push_back(u);
    onstk[u] = 1;
for (int v=0; v<n; v++){
      if (u != v && match[u] != v && !onstk[v]){
        int m = match[v];
if (dis[m] > dis[u] - edge[v][m] + edge[u][v]){
           dis[m] = dis[u] - edge[v][m] + edge[u][v];
           onstk[v] = 1;
           stk.push_back(v);
           if (SPFA(m)) return true;
          stk.pop_back();
          onstk[v] = 0;
        }
      }
    onstk[u] = 0;
    stk.pop_back();
    return false;
  int solve() {
    // find a match
    for (int i=0; i<n; i+=2){
      match[i] = i+1;
      match[i+1] = i;
    while (true){
      int found = 0;
      for( int i = 0 ; i < n ; i ++ )</pre>
        onstk[ i ] = dis[ i ] = 0;
      for (int i=0; i< n; i++){
        stk.clear();
        if (!onstk[i] && SPFA(i)){
           found = 1;
          while ((int)stk.size()>=2){
             int u = stk.back(); stk.pop_back();
             int v = stk.back(); stk.pop_back();
             match[u] = v;
```

match[v] = u;

```
}
}
if (!found) break;
}
int ret = 0;
for (int i=0; i<n; i++)
    ret += edge[i][match[i]];
    ret /= 2;
    return ret;
}
}graph;</pre>
```

5.9 Maximum General Weighted Matching

```
struct WeightGraph {
  static const int INF = INT_MAX;
  static const int N = 514;
  struct edge{
    int u,v,w; edge(){}
    edge(int ui,int vi,int wi)
      :u(ui),v(vi),w(wi){}
  int n,n_x;
  edge g[N*2][N*2];
  int lab[N*2];
  int match[N*2],slack[N*2],st[N*2],pa[N*2];
  int flo_from[N*2][N+1],S[N*2],vis[N*2];
  vector<int> flo[N*2];
  queue<int> q;
int e_delta(const edge &e){
    return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
  void update_slack(int u,int x){
    if(!slack[x]||e_delta(g[u][x])<e_delta(g[slack[x]][x</pre>
        ]))slack[x]=u;
  void set_slack(int x){
    slack[x]=0;
    for(int u=1;u<=n;++u)</pre>
      if(g[u][x].w>0&&st[u]!=x&&S[st[u]]==0)
        update_slack(u,x);
  void q_push(int x){
    if(x<=n)q.push(x);</pre>
    else for(size_t i=0;i<flo[x].size();i++)
  q_push(flo[x][i]);</pre>
  void set_st(int x,int b){
    st[x]=b;
    if(x>n)for(size_t i=0;i<flo[x].size();++i)</pre>
      set_st(flo[x][i],b);
  int get_pr(int b,int xr){
    int pr=find(flo[b].begin(),flo[b].end(),xr)-flo[b].
        begin();
    if(pr%2==1){
      reverse(flo[b].begin()+1,flo[b].end());
      return (int)flo[b].size()-pr;
    }else return pr;
  void set_match(int u,int v){
    match[u]=g[u][v].v;
    if(u<=n) return;</pre>
    edge e=g[u][v];
    int xr=flo_from[u][e.u],pr=get_pr(u,xr);
    for(int i=0;i<pr;++i)set_match(flo[u][i],flo[u][i</pre>
         ^1]);
    set_match(xr,v);
    rotate(flo[u].begin(),flo[u].begin()+pr,flo[u].end()
  void augment(int u,int v){
    for(;;){
      int xnv=st[match[u]];
      set_match(u,v);
      if(!xnv)return;
      set_match(xnv,st[pa[xnv]]);
      u=st[pa[xnv]],v=xnv;
  int get_lca(int u,int v){
```

```
static int t=0;
  for(++t;ullv;swap(u,v)){
    if(u==0)continue;
    if(vis[u]==t)return u;
    vis[u]=t;
    u=st[match[u]];
    if(u)u=st[pa[u]];
  }
  return 0;
void add_blossom(int u,int lca,int v){
  int b=n+1;
  while(b \le n_x \&st[b])++b;
  if(b>n_x)++n_x;
  lab[b]=0,S[b]=0;
  match[b]=match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for(int x=u,y;x!=lca;x=st[pa[y]])
    flo[b].push_back(x),flo[b].push_back(y=st[match[x
         ]]),q_push(y)
  reverse(flo[b].begin()+1,flo[b].end());
for(int x=v,y;x!=lca;x=st[pa[y]])
    flo[b].push_back(x),flo[b].push_back(y=st[match[x
         ]]),q_push(y);
  set_st(b,b);
  for(int x=1;x<=n_x;++x)g[b][x].w=g[x][b].w=0;
  for(int x=1;x<=n;++x)flo_from[b][x]=0;</pre>
  for(size_t i=0;i<flo[b].size();++i){</pre>
    int xs=flo[b][i];
    for(int x=1;x<=n_x;++x)</pre>
      if(g[b][x].w==0|le_delta(g[xs][x])<e_delta(g[b][
           (([x
        g[b][x]=g[xs][x],g[x][b]=g[x][xs];
    for(int x=1;x<=n;++x)</pre>
      if(flo_from[xs][x])flo_from[b][x]=xs;
  set_slack(b);
}
void expand_blossom(int b){
  for(size_t i=0;i<flo[b].size();++i)</pre>
    set_st(flo[b][i],flo[b][i])
  int xr=flo_from[b][g[b][pa[b]].u],pr=get_pr(b,xr);
  for(int i=0;i<pr;i+=2){</pre>
    int xs=flo[b][i],xns=flo[b][i+1];
    pa[xs]=g[xns][xs].u;
    S[xs]=1,S[xns]=0;
    slack[xs]=0,set_slack(xns);
    q_push(xns);
  S[xr]=1,pa[xr]=pa[b];
  for(size_t i=pr+1;i<flo[b].size();++i){</pre>
    int xs=flo[b][i];
    S[xs]=-1, set\_slack(xs);
  st[b]=0;
bool on_found_edge(const_edge &e){
  int u=st[e.u],v=st[e.v];
  if(S[v]=-1){
    pa[v]=e.u,S[v]=1;
    int nu=st[match[v]]
    slack[v]=slack[nu]=0;
    S[nu]=0,q_push(nu);
  }else if(S[v]==0){
    int lca=get_lca(u,v);
    if(!lca)return augment(u,v),augment(v,u),true;
    else add_blossom(u,lca,v);
  }
  return false;
bool matching(){
  memset(S+1,-1,sizeof(int)*n_x);
  memset(slack+1,0,sizeof(int)*n_x);
  q=queue<int>();
  for(int x=1;x<=n_x;++x)</pre>
    if(st[x]==x\&\&!match[x])pa[x]=0,S[x]=0,q_push(x);
  if(q.empty())return false;
  for(;;){
    while(q.size()){
      int u=q.front();q.pop();
      if(S[st[u]]==1)continue;
```

```
for(int v=1;v<=n;++v)
  if(g[u][v].w>0&&st[u]!=st[v]){
                                                                                      dst[ i ][ j ] = INF;
dst[ i ][ i ] = 0;
                if(e_delta(g[u][v])==0){
                                                                                    }
                   if(on_found_edge(g[u][v]))return true;
                                                                                 }
                                                                                 void add_edge( int ui , int vi , int wi ){
  dst[ ui ][ vi ] = min( dst[ ui ][ vi ] , wi );
  dst[ vi ][ ui ] = min( dst[ vi ][ ui ] , wi );
                }else update_slack(u,st[v]);
        int d=INF;
        for(int b=n+1;b<=n_x;++b)</pre>
                                                                                 void shortest_path(){
                                                                                    for( int i = 0 ; i < n ; i ++ )
for( int j = 0 ; j < n ; j ++ )
if( i != j && dst[ i ][ j ] != INF )
           if(st[b]==b&&S[b]==1)d=min(d,lab[b]/2);
        for(int x=1;x<=n_x;++x)</pre>
           if(st[x]==x\&slack[x]){
                                                                                    dst[ i ][ j ] += w[ i ];
for( int k = 0 ; k < n ; k ++ )</pre>
             if(S[x]==-1)d=min(d,e_delta(g[slack[x]][x]));
             else if(S[x]==0)d=min(d,e_delta(g[slack[x]][x
                                                                                      for (int i = 0; i < n; i ++)
                   ])/2);
                                                                                         for( int j = 0; j < n; j ++ )

dst[ i ][ j ] = min( dst[ i ][ j ],

dst[ i ][ k ] + dst[ k ][ j ]);
        for(int u=1;u<=n;++u){
  if(S[st[u]]==0){</pre>
                                                                                    for( int i = 0; i < n; i ++ )
for( int j = 0; j < n; j ++ )
  if( dst[ i ][ j ] != INF )
    dst[ i ][ j ] += w[ j ];</pre>
             if(lab[u]<=d)return 0;</pre>
             lab[u]-=d;
          }else if(S[st[u]]==1)lab[u]+=d;
        for(int b=n+1;b<=n_x;++b)
  if(st[b]==b){</pre>
                                                                                 int solve( const vector<int>& ter ){
                                                                                    int t = (int)ter.size();
             if(S[st[b]]==0)lab[b]+=d*2;
                                                                                    for( int i = 0 ; i < (1 << t ) ;
for( int j = 0 ; j < n ; j ++ )
dp[ i ][ j ] = INF;
             else if(S[st[b]]==1)lab[b]-=d*2;
                                                                                                                                 i ++ )
        q=queue<int>();
                                                                                    for( int i = 0; i < n; i ++ )
dp[0][i] = 0;
        for(int x=1;x<=n_x;++x)</pre>
           if(st[x]==x&&slack[x]&&st[slack[x]]!=x&&e_delta(
                                                                                    for( int msk = 1 ; msk < ( 1 << t ) ; msk ++ ){
                g[slack[x]][x])==0)
                                                                                      if( msk == (msk & (-msk))){
             if(on_found_edge(g[slack[x]][x]))return true;
                                                                                         int who = __[g( msk );
for( int i = 0 ; i < n ; i ++ )
  dp[ msk ][ i ] = dst[ ter[ who ] ][ i ];</pre>
        for(int b=n+1;b<=n_x;++b)</pre>
           if(st[b]==b\&\&S[b]==1\&\&lab[b]==0)expand_blossom(b)
                                                                                         continue;
     return false;
                                                                                       for( int i = 0 ; i < n ; i ++ )</pre>
  pair<long long,int> solve(){
                                                                                         for( int submsk = ( msk - 1 ) & msk ; submsk ;
    submsk = ( submsk - 1 ) & msk )
     memset(match+1,0,sizeof(int)*n);
                                                                                               dp[ msk ][ i ] = min( dp[ msk ][ i ],
     n x=n:
                                                                                                                    dp[ submsk ][ i ] +
dp[ msk ^ submsk ][ i ] - w[
     int n_matches=0;
     long long tot_weight=0;
                                                                                                                           i]);
     for(int u=0;u<=n;++u)st[u]=u,flo[u].clear();</pre>
                                                                                      for( int i = 0 ; i < n ; i ++ ){
  tdst[ i ] = INF;
  for( int j = 0 ; j < n ; j ++ )
    tdst[ i ] = min( tdst[ i ],</pre>
     int w_max=0;
     for(int u=1;u<=n;++u)</pre>
        for(int v=1;v<=n;++v){</pre>
           flo_from[u][v]=(u==v?u:0);
                                                                                                          w_max=max(w_max,g[u][v].w);
     for(int u=1;u<=n;++u)lab[u]=w_max;</pre>
                                                                                      for( int i = 0 ; i < n ; i ++
  dp[ msk ][ i ] = tdst[ i ];</pre>
     while(matching())++n_matches;
     for(int u=1;u<=n;++u)</pre>
        if(match[u]&&match[u]<u)</pre>
           tot_weight+=g[u][match[u]].w;
                                                                                    int ans = INF;
                                                                                    for( int i = 0; i < n; i ++)
     return make_pair(tot_weight,n_matches);
                                                                                      ans = min(ans, dp[(1 << t) - 1][i]);
  void add_edge( int ui , int vi , int wi ){
                                                                                    return ans;
     g[ui][vi].w = g[vi][ui].w = wi;
                                                                              } solver;
  void init( int _n ){
     n = _n;
for(int u=1;u<=n;++u)</pre>
                                                                               5.11 BCC based on vertex
        for(int v=1;v<=n;++v)</pre>
                                                                              struct BccVertex {
                                                                                 int n,nBcc,step,dfn[MXN],low[MXN];
vector<int> E[MXN],bccv[MXN];
          g[u][v]=edge(u,v,0);
} graph;
                                                                                 // vector<pair<int,int>> bcce[MXN];
                                                                                 int top,stk[MXN];
                                                                                 // pair<int,int> estk[MXM]; // max edge number
// int id[MXN],pos[MXN];
          Minimum Steiner Tree
5.10
                                                                                 void init(int _n) {
  n = _n;  nBcc = step = 0;
  for (int i=0; i<n; i++) E[i].clear();</pre>
// Minimum Steiner Tree O(V 3^T + V^2 2^T)
// shortest_path() should be called before solve()
// w:vertex weight, default 0
struct SteinerTree{
#define V 66
                                                                                 void addEdge(int u, int v)
#define T 10
                                                                                 { E[u].push_back(v); E[v].push_back(u); }
#define INF 1023456789
                                                                                 void DFS(int u, int f) {
  int n , dst[V][V] , dp[1 << T][V] , tdst[V] , w[V]; void init( int _{\rm n} ){
                                                                                    dfn[u] = low[u] = step++;
                                                                                    stk[top++] = u;
     n = _n; fill( w , w + n , 0 );
for( int i = 0 ; i < n ; i ++ ){
  for( int j = 0 ; j < n ; j ++ )</pre>
                                                                                    for (auto v:E[u]) {
```

if (v == f) continue;
if (dfn[v] == -1) {

```
// estk[etop++]={u,v};
        DFS(v<sub>,</sub>u);
        low[u] = min(low[u], low[v]);
        if (low[v] >= dfn[u]) {
          bccv[nBcc].clear(); // bcce[nBcc].clear();
          // pair<int,int> ez;
          // do {
          //
               ez=estk[--etop];
                bcce[nBcc].push_back(ez);
           // }while(ez.first!=u);
          do {
            z = stk[--top];
            // id[z]=nBcc; pos[z]=bccv[nBcc].size();
bccv[nBcc].push_back(z);
           } while (z != v);
          bccv[nBcc++].push_back(u);
      }else
        low[u] = min(low[u],dfn[v]);
        //if(dfn[v]<dfn[u]) estk[etop++]={u,v};</pre>
    }
  }
  vector<vector<int>> solve() {
    vector<vector<int>> res;
    for (int i=0; i<n; i++)
      dfn[i] = low[i] = -1;
    for (int i=0; i<n; i++)
      if (dfn[i] == -1)
        top = 0; etop = 0;
        DFS(i,i);
    for(int i=0;i<nBcc;i++) res.push_back(bccv[i]);</pre>
    return res;
  /* bccv[.first][{.first,.second}]={u,v}
  pair<int,pair<int,int>> getpos(int u,int v){
    if(dfn[u]>dfn[v]) swap(u,v);
    int cid=id[v]
    if(id[u]==cid) return {cid,{pos[v],pos[u]}};
    else return {cid,pos[v],{bccv[cid].size()-1}};
}graph;
```

5.12 Min Mean Cycle

```
* minimum mean cycle O(VE) */
struct MMC{
#define E 101010
#define V 1021
#define inf 1e9
#define eps 1e-6
  struct Edge { int v,u; double c; };
   int n, m, prv[V][V], prve[V][V], vst[V];
  Edge e[E];
  vector<int> edgeID, cycle, rho;
  double d[V][V];
  void init( int _n )
  \{ n = _n; m = 0; \}
  // WARNING: TYPÉ matters
  void addEdge( int vi , int ui , double ci )
  { e[ m ++ ] = { vi , ui , ci }; }
void bellman_ford() {
     for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {
  fill(d[i+1], d[i+1]+n, inf);
  for(int j=0; j<m; j++) {
    int y = 0[i] y = 0.</pre>
          int v = e[j].v, u = e[j].u;
if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
             d[i+1][u] = d[i][v]+e[j].c;
             prv[i+1][u] = v;
             prve[i+1][u] = j;
          }
       }
    }
  double solve(){
     // returns inf if no cycle, mmc otherwise
     double mmc=inf;
     int st = -1;
     bellman_ford();
     for(int i=0; i<n; i++) {</pre>
```

```
double avg=-inf;
      for(int k=0; k<n; k++) {</pre>
         if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i</pre>
             ])/(n-k));
        else avg=max(avg,inf);
      if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
    }
    if(st==-1) return inf;
    FZ(vst);edgeID.clear();cycle.clear();rho.clear();
    for (int i=n; !vst[st]; st=prv[i--][st]) {
      vstΓst]++;
      edgeID.push_back(prve[i][st]);
      rho.push_back(st);
    while (vst[st] != 2) {
      int v = rho.back(); rho.pop_back();
      cycle.push_back(v)
      vst[v]++;
    }
    reverse(ALL(edgeID));
    edgeID.resize((int)cycle.size());
    return mmc;
} mmc;
```

5.13 Directed Graph Min Cost Cycle

```
// works in O(N M)
#define INF 1000000000000000LL
#define N 5010
#define M 200010
struct edge{
  int to; LL w;
  edge(int a=0, LL b=0): to(a), w(b){}
struct node{
  LL d; int u, next;
  node(LL a=0, int b=0, int c=0): d(a), u(b), next(c){}
}b[M];
struct DirectedGraphMinCycle{
  vector<edge> g[N], grev[N];
  LL dp[N][N], p[N], d[N], mu;
  bool inq[N];
  int n, bn, bsz, hd[N];
  void b_insert(LL d, int u){
    int i = d/mu;
    if(i >= bn) return;
    b[++bsz] = node(d, u, hd[i]);
    hd[i] = bsz;
  void init( int _n ){
    n = _n;
for( int i = 1 ; i <= n ; i ++ )
g[ i ].clear();
  void addEdge( int ai , int bi , LL ci )
  { g[ai].push_back(edge(bi,ci)); }
  LL solve(){
    fill(dp[0], dp[0]+n+1, 0);
    for(int i=1; i<=n; i++){
      fill(dp[i]+1, dp[i]+n+1, INF);
for(int j=1; j<=n; j++) if(dp[i-1][j] < INF){
  for(int k=0; k<(int)g[j].size(); k++)</pre>
           dp[i][g[j][k].to] =min(dp[i][g[j][k].to]
                                       dp[i-1][j]+g[j][k].w);
      }
    mu=INF; LL bunbo=1;
    for(int i=1; i<=n; i++) if(dp[n][i] < INF){</pre>
       LL a=-INF, b=1;
       for(int j=0; j<=n-1; j++) if(dp[j][i] < INF){
  if(a*(n-j) < b*(dp[n][i]-dp[j][i])){</pre>
           a = dp[n][i]-dp[j][i];
           b = n-j;
         }
       if(mu*b > bunbo*a)
         mu = a, bunbo = b;
    if(mu < 0) return -1; // negative cycle</pre>
    if(mu == INF) return INF; // no cycle
```

```
if(mu == 0) return 0;
     for(int i=1; i<=n; i++)
       for(int j=0; j<(int)g[i].size(); j++)
g[i][j].w *= bunbo;</pre>
     memset(p, 0, sizeof(p));
     queue<int> q;
for(int i=1; i<=n; i++){</pre>
       q.push(i);
       inq[i] = true;
     while(!q.empty()){
       int i=q.front(); q.pop(); inq[i]=false;
       for(int j=0; j<(int)g[i].size(); j++){
  if(p[g[i][j].to] > p[i]+g[i][j].w-mu){
    p[g[i][j].to] = p[i]+g[i][j].w-mu;
}
             if(!inq[g[i][j].to]){
               q.push(g[i][j].to);
               inq[g[i][j].to] = true;
            }
          }
       }
     for(int i=1; i<=n; i++) grev[i].clear();
for(int i=1; i<=n; i++)</pre>
       for(int j=0; j<(int)g[i].size(); j++){</pre>
          g[i][j].w += p[i]-p[g[i][j].to]
          grev[g[i][j].to].push_back(edge(i, g[i][j].w));
     LL mldc = n*mu;
                                                                            }
     for(int i=1; i<=n; i++){</pre>
       bn=mldc/mu, bsz=0;
memset(hd, 0, sizeof(hd));
       fill(d+i+1, d+n+1, INF);
       b_insert(d[i]=0, i);
       for(int j=0; j<=bn-1; j++) for(int k=hd[j]; k; k=b</pre>
             [k].next){
          int u = b[k].u;
          LL du = b[k].d;
          if(du > d[u]) continue;
for(int l=0; l<(int)g[u].size(); l++) if(g[u][l</pre>
               1.to > i)
            if(d[g[u][l].to] > du + g[u][l].w){
               d[g[u][l].to] = du + g[u][l].w;
               b_insert(d[g[u][l].to], g[u][l].to);
            }
          }
       for(int j=0; j<(int)grev[i].size(); j++) if(grev[i])</pre>
                                                                                 ) + 1;
             |[i].to > i|
          mldc=min(mldc,d[grev[i][j].to] + grev[i][j].w);
     return mldc / bunbo;
|} graph;
5.14 K-th Shortest Path
// time: O(|E| \setminus g \mid E| + \mid V| \setminus g \mid V| + K)
// memory: 0(|E| \lg |E| + |V|)
struct KSP{ // 1-base
   struct nd{
     int u, v; ll d;
nd(int ui = 0, int vi = 0, ll di = INF)
     \{ u = ui; v = vi; d = di; \}
   struct heap{
    nd* edge; int dep; heap* chd[4];
   static int cmp(heap* a,heap* b)
   { return a->edge->d > b->edge->d; }
  struct node{
     int v; ll d; heap* H; nd* E;
```

node(){}

nd *nxt[N];

node(ll _d, int _v, nd* _E)
{ d =_d; v = _v; E = _E; }
node(heap* _H, ll _d)

friend bool operator<(node a, node b)</pre>

 $\{ H = _H; d = _d; \}$

{ return a.d > b.d; }

int n, k, s, t, dst[N];

```
vector<nd*> g[ N ], rg[ N ];
heap *nullNd, *head[ N ];
  void init( int _n , int _k , int _s , int _t ){
    n = _n; k = _k; s = _s; t = _t;
for( int i = 1 ; i <= n ; i ++ ){
    g[ i ].clear(); rg[ i ].clear();
    nxt[ i ] = NULL; head[ i ] = NULL;
    dst[ i ] = -1;</pre>
  void addEdge( int ui , int vi , ll di ){
  nd* e = new nd(ui, vi, di);
     g[ ui ].push_back( e );
     rg[ vi ].push_back( e );
  queue<int> dfsQ;
  void dijkstra(){
     while(dfsQ.size()) dfsQ.pop();
     priority_queue<node> Q;
     Q.push(node(0, t, NULL));
     while (!Q.empty()){
       node p = Q.top(); Q.pop();
       if(dst[p.v] != -1) continue;
dst[ p.v ] = p.d;
nxt[ p.v ] = p.E;
       dfsQ.push( p.v_);
       for(auto e: rg[ p.v ])
          Q.push(node(p.d + e->d, e->u, e));
  heap* merge(heap* curNd, heap* newNd){
     if(curNd == nullNd) return newNd;
    heap* root = new heap;
memcpy(root, curNd, sizeof(heap));
     if(newNd->edge->d < curNd->edge->d){
       root->edge = newNd->edge;
root->chd[2] = newNd->chd[2]
       root->chd[3] = newNd->chd[3];
       newNd->edge = curNd->edge;
newNd->chd[2] = curNd->chd[2];
       newNd - > chd[3] = curNd - > chd[3];
     if(root->chd[0]->dep < root->chd[1]->dep)
       root->chd[0] = merge(root->chd[0],newNd);
       root->chd[1] = merge(root->chd[1],newNd);
     root->dep = max(root->chd[0]->dep, root->chd[1]->dep
     return root;
  vector<heap*> V;
  void build(){
     nullNd = new heap;
     nullNd->dep = 0;
     nullNd->edge = new nd;
     fill(nullNd->chd, nullNd->chd+4, nullNd);
     while(not dfsQ.empty()){
       int u = dfsQ.front(); dfsQ.pop();
if(!nxt[ u ]) head[ u ] = nullNd;
       else head[ u ] = head[nxt[ u ]->v];
       V.clear();
       for( auto&& e : g[ u ] ){
          int v = e->v;
          if( dst[ v ] == -1 ) continue;
e->d += dst[ v ] - dst[ u ];
          if( nxt[ u ] != e ){
            heap* p = new heap;
            fill(p->chd, p->chd+4, nullNd);
            p->dep = 1;
            p->edge = e;
            V.push_back(p);
         }
       if(V.empty()) continue;
       make_heap(V.begin(), V.end(), cmp);
#define L(X) ((X<<1)+1)
#define R(X) ((X<<1)+2)
       for( size_t i = 0 ; i < V.size() ; i ++ ){</pre>
          if(L(i) < V.size()) V[i]->chd[2] = V[L(i)];
          else V[i]->chd[2]=nullNd;
          if(R(i) < V.size()) V[i]->chd[3] = V[R(i)];
          else V[i]->chd[3]=nullNd;
```

```
head[u] = merge(head[u], V.front());
     }
   }
   vector<ll> ans;
   void first_K(){
     ans.clear();
     priority_queue<node> Q;
     if( dst[ s ] == -1 ) return;
ans.push_back( dst[ s ] );
if( head[s] != nullNd )
        Q.push(node(head[s], dst[s]+head[s]->edge->d));
     for( int _ = 1 ; _ < k and not Q.empty() ; _ ++ ){
  node p = Q.top(), q; Q.pop();</pre>
        ans.push_back( p.d );
        if(head[ p.H->edge->v ] != nullNd){
          q.H = head[p.H->edge->v];
          q.d = p.d + q.H->edge->d;
          Q.push(q);
        for( int i = 0 ; i < 4 ; i ++ )
  if( p.H->chd[ i ] != nullNd ){
    q.H = p.H->chd[ i ];

             q.d = p.d - p.H->edge->d + p.H->chd[i]->edge
                  ->d;
             Q.push( q );
     }
   void solve(){ // ans[i] stores the i-th shortest path
     dijkstra();
     build();
     first_K(); // ans.size() might less than k
} solver;
```

5.15 Chordal Graph

```
struct Chordal {
  static const int MXN = 100010;
  vector<int> E[MXN], V[MXN];
  int n,f[MXN],rk[MXN],order[MXN],stk[MXN],nsz[MXN];
  bool vis[MXN], isMaximalClique[MXN];
  void init(int _n) {
    n = _n;
for(int i = 0; i <= n; ++i) {
    E[i].clear(), V[i].clear();</pre>
      f[i]=rk[i]=order[i]=vis[i]=0;
  void addEdge(int x, int y) {
    E[x].push_back(y), E[y].push_back(x);
  void mcs() {
    for(int i = 1; i <= n; ++i) V[0].push_back(i);</pre>
    for(int i = n, M = 0; i >= 1; --i) {
      for(;;) {
         while(V[M].size()&&vis[V[M].back()])
           V[M].pop_back();
         if(V[M].size()) break; else M--;
      auto x=V[M].back();order[i]=x;rk[x]=i;vis[x]=1;
for(auto y : E[x]) if(!vis[y])
         f[y]++, V[f[y]].push_back(y), M=max(M,f[y]);
    }
  bool isChordal() {
    for(int i = 0; i <= n; ++i) vis[i] = stk[i] = 0;
for(int i = n; i >= 1; --i) {
      int top = 0, cnt = 0, m = n+1;
      for(auto x : E[order[i]]) if(rk[x] > i)
         stk[top++]=x, vis[x]=1, m = min(m, rk[x]);
      if(m==n+1) continue:
      for(auto x : E[order[m]]) if(vis[x]) ++cnt;
      for(int j = 0; j < top; ++j) vis[stk[j]] = 0;</pre>
      if(cnt + 1 != top) return 0;
    return 1;
  void getMaximalClique() {
    for(int i = n; i >= 1; --i) {
      int M = n+1, w = order[i], v = 0;
```

```
nsz[w] = 0; isMaximalClique[w] = 1;
for(auto x : E[w]) if(rk[x] > i) {
       nsz[w]++;
       if(\bar{r}k[x] < M) M = rk[x], v = x;
     if(v)isMaximalClique[v]&=nsz[v]+1>nsz[w];
  }
}
int getMaximumClique() {
  int res = 0;
  for(int i = 1; i \le n; ++i) res=max(res,f[i]+1);
  return res;
int getMaximumIndependentSet() {
  for(int i = 0; i <= n; ++i) vis[i] = 0;</pre>
  int res = 0;
  for(int i = 1; i <= n; ++i) if(!vis[order[i]]) {
  res++, vis[order[i]] = 1;</pre>
     for(auto x : E[order[i]]) vis[x] = 1;
  return res;
```

5.16 Graph Method

Manhattan MST For each point, consider the points that surround it(8 octants). Then, connect it with the closest point. For example, consider 45~90. For each point p, the closest point is $min\{x+y \mid x-y >= p.x-p.y, x >= p.x$ Finally, the answer is this new graphs(E=4N) MST.

String

6.1 PalTree

```
const int MXN = 1000010;
struct PalT{
  int nxt[MXN][26],fail[MXN],len[MXN];
  int tot,lst,n,state[MXN],cnt[MXN],num[MXN];
  int diff[MXN],sfail[MXN],fac[MXN],dp[MXN];
  char s[MXN]={-1};
int newNode(int l,int f){
    len[tot]=1,fail[tot]=f,cnt[tot]=num[tot]=0;
    memset(nxt[tot],0,sizeof(nxt[tot]));
diff[tot]=(l>0?l-len[f]:0);
    sfail[tot]=(l>0&&diff[tot]==diff[f]?sfail[f]:f);
    return tot++:
  int getfail(int x){
    while(s[n-len[x]-1]!=s[n]) x=fail[x];
    return x;
  int getmin(int v){
  dp[v]=fac[n-len[sfail[v]]-diff[v]];
    if(diff[v]==diff[fail[v]])
         dp[v]=min(dp[v],dp[fail[v]]);
    return dp[v]+1;
  int push(){
    int c=s[n]-'a',np=getfail(lst);
if(!(lst=nxt[np][c])){
       lst=newNode(len[np]+2,nxt[getfail(fail[np])][c]);
       nxt[np][c]=lst; num[lst]=num[fail[lst]]+1;
    fac[n]=n;
    for(int_v=lst;len[v]>0;v=sfail[v])
         fac[n]=min(fac[n],getmin(v));
    return ++cnt[lst],lst;
  void init(const char *_s){
    tot=lst=n=0;
    newNode(0,1), newNode(-1,1);
     for(;_s[n];) s[n+1]=_s[n],++n,state[n-1]=push();
    for(int i=tot-1;i>1;i--) cnt[fail[i]]+=cnt[i];
}palt;
```

6.2 SAIS

```
| const int N = 300010;
```

```
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )</pre>
#define REP1(i,a,b) for ( int i=(a); i <= int(b); i++)
    bool _t[N*2];
     int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
             hei[N], r[N];
    int operator [] (int i){ return _sa[i]; }
void build(int *s, int n, int m){
         memcpy(_s, s, sizeof(int) * n);
         sais(_s, _sa, _p, _q, _t, _c, n, m);
mkhei(n);
     void mkhei(int n){
         REP(i,n) r[\_sa[i]] = i;
         hei[0] = 0;
         REP(i,n) if(r[i]) {
              int ans = i>0 ? max(hei[r[i-1]] - 1, 0) : 0;
              hei[r[i]] = ans;
        }
     void sais(int *s, int *sa, int *p, int *q, bool *t,
               int *c, int_n, int z){
         bool uniq = t[n-1] = true, neq;
         int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
                   lst = -1;
#define MSO(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
memcpy(x, c, sizeof(int) * z); \
memcpy(x + 1, c, sizeof(int) * (z - 1)); 

REP(i,n) if(sa[i] && !t[sa[i]-1]) sa[x[s[sa[i]-1]]++] =
sa[i]-1; \
memcpy(x, c, sizeof(int) * z); \
for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i]-1])
          sa[--x[s[sa[i]-1]]] = sa[i]-1;
         MSO(c, z);
         REP(i,n) uniq &= ++c[s[i]] < 2;
        MAGIC(REP1(i,1,n-1) if(t[i] && !t[i-1]) sa[--x[s[i ]]]=p[q[i]=nn++]=i);
         REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {
              \label{eq:neq_lst_objective} \begin{subarray}{ll} neq=lst<0 | lmemcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa[l]) | lmemcmp(s+sa[i],s+lst,(p[sa[i]]+1]-sa[l]) | lmemcmp(s+sa[i],s+ls
                        il)*sizeof(int));
             ns[q[lst=sa[i]]]=nmxz+=neq;
         sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz
                   + 1);
         MAGIC(for(int i = nn - 1; i \ge 0; i--) sa[--x[s[p[
                   nsa[i]]]] = p[nsa[i]]);
}sa;
int H[N], SA[N], RA[N];
void suffix_array(int* ip, int len) {
     // should padding a zero in the back
     // ip is int array, len is array length
     // ip[0..n-1] != 0, and ip[len] = 0
     ip[len++] = 0;
    sa.build(ip, len, 128);
memcpy(H,sa.hei+1,len<<2);</pre>
    memcpy(SA,sa._sa+1,len<<2)</pre>
     for(int i=0; i<len; i++) RA[i] = sa.r[i]-1;</pre>
     // resulting height, sa array \in [0,len)
```

6.3 SuffixAutomata

```
// any path start from root forms a substring of S
// occurrence of P : iff SAM can run on input word P
// number of different substring : ds[1]-1
// total length of all different substring : dsl[1]
// max/min length of state i : mx[i]/mx[mom[i]]+1
// assume a run on input word P end at state i:
// number of occurrences of P : cnt[i]
// first occurrence position of P : fp[i]-IPI+1
// all position of P : fp of "dfs from i through rmom"
const int MXM = 1000010;
struct SAM{
   int tot, root, lst, mom[MXM], mx[MXM]; //ind[MXM]
```

```
int nxt[MXM][33]; //cnt[MXM],ds[MXM],dsl[MXM],fp[MXM]
  // bool v[MXM]
  int newNode(){
    int res = ++tot;
    fill(nxt[res], nxt[res]+33, 0);
    mom[res] = mx[res] = 0; //cnt=ds=dsl=fp=v=0
    return res;
  void init(){
    tot = 0; root = newNode(); lst = root;
  void push(int c){
    int p = lst;
    int np = newNode(); //cnt[np]=1
mx[np] = mx[p]+1; //fp[np]=mx[np]-1
    for(; p && nxt[p][c] == 0; p = mom[p])
       nxt[p][c] = np
     if(p == 0) mom[np] = root;
    else{
       int q = nxt[p][c];
       if(mx[p]+1 == mx[q]) mom[np] = q;
         int nq = newNode(); //fp[nq]=fp[q]
         mx[nq] = mx[p]+1;
         for(int i = 0; i < 33; i++)
           nxt[nq][i] = nxt[q][i];
         mom[nq] = mom[q]; mom[q] = nq; mom[np] = nq;
         for(; p && nxt[p][c] == q; p = mom[p])
           nxt[p][c] = nq;
    lst = np;
  void calc(){
    calc(root); iota(ind,ind+tot,1);
    sort(ind,ind+tot,[&](int i,int j){return mx[i]<mx[j</pre>
    for(int i=tot-1;i>=0;i--)
    cnt[mom[ind[i]]]+=cnt[ind[i]];
  void calc(int x){
    v[x]=ds[x]=1;dsl[x]=0; //rmom[mom[x]].push_back(x);
    for(int i=0;i<26;i++){</pre>
       if(nxt[x][i]){
         if(!v[nxt[x][i]]) calc(nxt[x][i]);
ds[x]+=ds[nxt[x][i]];
         dsl[x]+=ds[nxt[x][i]]+dsl[nxt[x][i]];
    }
  void push(char *str){
    for(int i = 0; str[i]; i++)
push(str[i]-'a');
} sam;
6.4 Z Value
void z_value(const char *s,int len,int *z){
  z[0]=len;
  for(int i=1,l=0,r=0;i<len;i++){</pre>
    z[i]=i<r?(i-l+z[i-l]<z[l]?z[i-l]:r-i):0;
    while(i+z[i]<len&&s[i+z[i]]==s[z[i]]) ++z[i];</pre>
     if(i+z[i]>r) l=i,r=i+z[i];
  }
}
6.5 BWT
struct BurrowsWheeler{
#define SIGMA 26
```

```
#define SIGMA 26
#define BASE 'a'
vector<int> v[ SIGMA ];
void BWT(char* ori, char* res){
   // make ori -> ori + ori
   // then build suffix array
}
void iBWT(char* ori, char* res){
   for( int i = 0 ; i < SIGMA ; i ++ )
      v[ i ].clear();
   int len = strlen( ori );
   for( int i = 0 ; i < len ; i ++ )</pre>
```

```
v[ ori[i] - BASE ].push_back( i );
vector<int> a;
for( int i = 0 , ptr = 0 ; i < SIGMA ; i ++ )
    for( auto j : v[ i ] ){
        a.push_back( j );
        ori[ ptr ++ ] = BASE + i;
    }
for( int i = 0 , ptr = 0 ; i < len ; i ++ ){
    res[ i ] = ori[ a[ ptr ] ];
    ptr = a[ ptr ];
    }
    res[ len ] = 0;
}
bwt;</pre>
```

6.6 ZValue Palindrome

```
void z_value_pal(char *s,int len,int *z){
    len=(len<<1)+1;
    for(int i=len-1;i>=0;i--)
        s[i]=i&1?s[i>>1]:'@';
    z[0]=1;
    for(int i=1,l=0,r=0;i<len;i++){
        z[i]=i<r?min(z[l+l-i],r-i):1;
        while(i-z[i]>=0&&i+z[i]<len&&s[i-z[i]]==s[i+z[i]])++
        z[i];
        if(i+z[i]>r) l=i,r=i+z[i];
    }
}
```

6.7 Smallest Rotation

```
//rotate(begin(s),begin(s)+minRotation(s),end(s))
int minRotation(string s) {
  int a = 0, N = s.size(); s += s;
  rep(b,0,N) rep(k,0,N) {
    if(a+k == b || s[a+k] < s[b+k])
      {b += max(0, k-1); break;}
    if(s[a+k] > s[b+k]) {a = b; break;}
  } return a;
}
```

6.8 Cyclic LCS

```
#define L 0
#define LU 1
#define U 2
const int mov[3][2]=\{0,-1,-1,-1,-1,0\};
int al,bl;
char a[MAXL*2],b[MAXL*2]; // 0-indexed
int dp[MAXL*2](MAXL];
char pred[MAXL*2][MAXL];
inline int lcs_length(int r) {
  int i=r+al,j=bl,l=0;
 while(i>r) {
    char dir=pred[i][j];
    if(dir==LU) l++;
    i+=mov[dir][0]; j+=mov[dir][1];
  return 1;
inline void reroot(int r) { // r = new base row
  int i=r, j=1;
  while(j<=bl&&pred[i][j]!=LU) j++;</pre>
 if(j>bl) return;
pred[i][j]=L;
while(i<2*al&&j<=bl) {</pre>
    if(pred[i+1][j]==U) {
  i++; pred[i][j]=L;
    } else if(j<bl&&pred[i+1][j+1]==LU) {</pre>
      i++; j++; pred[i][j]=L;
    } else j++;
 }
int cyclic_lcs() {
 \ensuremath{/\!/} a, b, al, bl should be properly filled
  // note: a WILL be altered in process
                 concatenated after itself
  char tmp[MAXL];
  if(al>bl) {
    swap(al,bl); strcpy(tmp,a);
    strcpy(a,b); strcpy(b,tmp);
```

```
strcpy(tmp,a); strcat(a,tmp);
// basic lcs
for(int i=0;i<=2*al;i++) {</pre>
   dp[i][0]=0; pred[i][0]=U;
for(int j=0;j<=bl;j++) {
  dp[0][j]=0; pred[0][j]=L;</pre>
for(int i=1;i<=2*al;i++) {
  for(int j=1;j<=bl;j++) {
  if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;</pre>
     else dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
if(dp[i][j-1]==dp[i][j]) pred[i][j]=L;
else if(a[i-1]==b[j-1]) pred[i][j]=LU;
     else pred[i][j]=U;
  }
}
// do cyclic lcs
int clcs=0;
for(int i=0;i<al;i++) {</pre>
  clcs=max(clcs,lcs_length(i)); reroot(i+1);
// recover a
a[al]='\0'
return clcs;
```

7 Data Structure

7.1 Link-Cut Tree

```
const int MEM = 100005;
struct Splay {
  static Splay nil, mem[MEM], *pmem;
  Splay *ch[2], *f;
int val, rev, size; // int sum, vir, tot;
  Splay (int _val=-1) : val(_val), rev(0), size(1)
  \{ f = ch[0] = ch[1] = &nil; \}
  bool isr()
  { return f->ch[0] != this && f->ch[1] != this; }
  int dir(){return f->ch[0] != this;}
  void setCh(Splay *c, int d){
    ch[d] = c; if (c != &nil) c->f = this; pull();
  void push(){
  if( !rev ) return;
    swap(ch[0], ch[1]);
if (ch[0] != &nil) ch[0]->rev ^= 1;
if (ch[1] != &nil) ch[1]->rev ^= 1;
    rev=0;
  void pull(){
    size = ch[0] -> size + ch[1] -> size + 1;
    // sum={ch[0]->sum,ch[1]->sum,val};
    // tot={sum,vir}
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
}Splay::nil,Splay::mem[MEM],*Splay::pmem=Splay::mem;
Splay *nil = &Splay::nil;
void rotate(Splay *x){
  Splay *p = x->f; int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f
  p->setCh(x->ch[!d], d); x->setCh(p, !d);
vector<Splay*> splayVec;
void splay(Splay *x){
  splayVec.clear();
  for (Splay *q=x;; q=q->f){
    splayVec.push_back(q);
    if (q->isr()) break;
  reverse(begin(splayVec), end(splayVec));
for (auto it : splayVec) it->push();
  while (!x->isr()) {
    if (x->f->isr()) rotate(x)
    else if (x->dir()==x->f->dir())
      rotate(x->f),rotate(x);
    else rotate(x), rotate(x);
```

```
int id(Splay *x) { return x - Splay::mem + 1; }
Splay* access(Splay *x){
  Splay *q = nil;
  for (;x!=nil;x=x->f){
    splay(x); // x->vir+={x->ch[0]->tot}-{q->tot};
    x \rightarrow setCh(q, 1); q = x;
  return q;
void chroot(Splay *x){
  access(x); splay(x); x->rev ^= 1;
void link(Splay *x, Splay *y){
 chroot(y); access(x); splay(x); y->f=x;
  // x->vir+={y->tot};
void cut_p(Splay *y) {
  access(y); splay(y); y->ch[0] = y->ch[0]->f = nil;
void cut(Splay *x, Splay *y){
 chroot(x); cut_p(y);
Splay* get_root(Splay *x) {
  x=access(x):
  for(; x \rightarrow ch[0] != nil; x = x \rightarrow ch[0]) x \rightarrow push();
  splay(x); return x;
bool conn(Splay *x, Splay *y) {
 return get_root(x) == get_root(y);
Splay* lca(Splay *x, Splay *y) {
 access(x); return access(y);
/* query(Splay *x,Splay *y){
  setroot(y), x=access(x); return x->size; // x->sum;
/* query(Splay *x,Splay *y){
 Splay *p=lca(x,y);
return 1+p->ch[1]->size+(x!=p?x->size:0);
  // {p->val,p->ch[1]->sum,x!=p?x->sum:0};
```

8 Others

8.1 Find max tangent(x,y is increasing)

```
const int MAXN = 100010;
Pt sum[MAXN], pnt[MAXN], ans, calc; inline bool cross(Pt a, Pt b, Pt c){
  return (c.y-a.y)*(c.x-b.x) > (c.x-a.x)*(c.y-b.y);
double find_max_tan(int n,int l,LL dy[]){
  int np, st, ed, now;
  sum[0].x = sum[0].y = np = st = ed = 0;
for (int i = 1, v; i <= n; i++)
    sum[i].x=i,sum[i].y=sum[i-1].y+dy[i-1];</pre>
  ans.x = now = 1, ans.y = -1;
  for (int i = 0; i <= n - 1; i++){
    while(np>1&&cross(pnt[np-2],pnt[np-1],sum[i]))
    if (np < now \&\& np != 0) now = np;
    pnt[np++] = sum[i];
    while(now<np&&!cross(pnt[now-1],pnt[now],sum[i+l]))</pre>
      now++;
    calc = sum[i + l] - pnt[now - 1]
    if (ans.y * calc.x < ans.x * calc.y)</pre>
      ans = calc,st = pnt[now - 1].x,ed = i + l;
  return (double)(sum[ed].y-sum[st].y)/(sum[ed].x-sum[st
       \exists .x);
```

8.2 Exact Cover Set

```
// given n*m 0-1 matrix
// find a set of rows s.t.
// for each column, there's exactly one 1
#define N 1024 //row
#define M 1024 //column
#define NM ((N+2)*(M+2))
char A[N][M]; //n*m 0-1 matrix
```

```
bool used[N]; //answer: the row used
int id[N][M]
int L[NM],R[NM],D[NM],U[NM],C[NM],S[NM],ROW[NM];
void remove(int c){
    L[R[c]]=L[c]; R[L[c]]=R[c];
  for( int i=D[c]; i!=c; i=D[i] )
    for( int j=R[i]; j!=i; j=R[j] )
      U[D[j]]=U[j]; D[U[j]]=D[j]; S[C[j]]--;
void resume(int c){
  for( int i=D[c]; i!=c; i=D[i] )
    for( int j=L[i]; j!=i; j=L[j] ){
      U[D[j]]=D[Ū[j]]=j; S[C[j]]++;
  L[R[c]]=R[L[c]]=c;
bool dfs(){
  if(R[0]==0) return 1;
  int md=100000000,c;
  for( int i=R[0]; i!=0; i=R[i] )
    if(S[i]<md){ md=S[i]; c=i; }</pre>
  if(md==0) return 0;
  remove(c);
  for( int i=D[c]; i!=c; i=D[i] ){
    used[ROW[i]]=1
    for( int j=R[i]; j!=i; j=R[j] ) remove(C[j]);
    if(dfs()) return 1;
    for( int j=L[i]; j!=i; j=L[j] ) resume(C[j]);
    used[ROW[i]]=0;
  resume(c);
  return 0;
bool exact_cover(int n,int m){
  for( int i=0; i<=m; i++ ){
   R[i]=i+1; L[i]=i-1; U[i]=D[i]=i;</pre>
    S[i]=0; C[i]=i;
  R[m]=0; L[0]=m;
  int t=m+1;
  for( int i=0; i<n; i++ ){</pre>
    int k=-1;
    for( int j=0; j<m; j++ ){
   if(!A[i][j]) continue;</pre>
      if(k==-1) L[t]=R[t]=t
      else{ L[t]=k; R[t]=R[k];
      k=t; D[t]=j+1; U[t]=U[j+1];
      L[R[t]]=R[L[t]]=U[D[t]]=D[U[t]]=t;
      C[t]=j+1; S[C[t]]++; ROW[t]=i; id[i][j]=t++;
    }
  for( int i=0; i<n; i++ ) used[i]=0;</pre>
  return dfs();
```

8.3 Binary Next Permutation

```
|ull next_perm(ull v){
   ull t=v|(v-1);
   return (t+1)|(((~t&-~t)-1)>>(__builtin_ctzll(v)+1));
}
```

8.4 Hilbert Curve

```
long long hilbert(int n, int x, int y) {
  long long res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x & s) > 0, ry = (y & s) > 0;
    res += s * 1ll * s * ((3 * rx) ^ ry);
    if (ry == 0) {
        if (rx == 1) x = s - 1 - x, y = s - 1 - y;
        swap(x, y);
    }
  }
  return res;
}
```