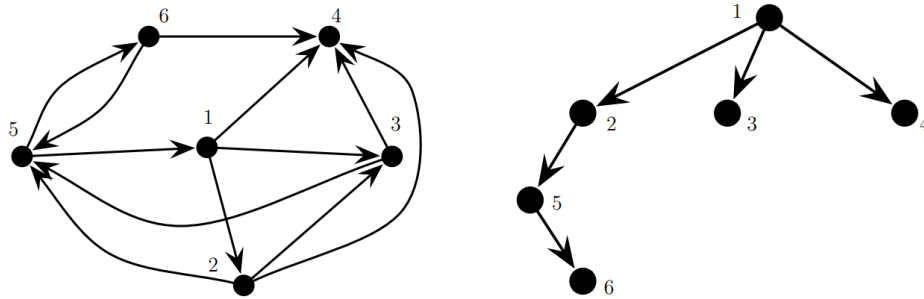


Flooding con pipe

Studiare il funzionamento delle pipe in linux. E' possibile realizzare per il grafo in figura l'algoritmo Flooding in C utilizzando un processo (o thread) per ogni agente ed una **Pipe** (bloccante o non bloccante?) per ogni canale di comunicazione? Si considerino due pipe per ogni agente (in-pipe, out-pipe) e si implementino con precisione le funzioni **msg** ed **stf**.



Synchronous Network: $\mathcal{S} = (\{1, \dots, n\}, E_{\text{cmm}})$

Distributed Algorithm: FLOODING

Alphabet: $\mathbb{A} = \{\alpha, \dots, \omega\} \cup \text{null}$

Processor State: $w = (\text{parent}, \text{data}, \text{snd-flag})$, where

parent	$\in \{0, \dots, n\},$	initially: $\text{parent}^{[1]} = 1,$ $\text{parent}^{[j]} = 0$ for all $j \neq 1$
data	$\in \mathbb{A},$	initially: $\text{data}^{[1]} = \mu,$ $\text{data}^{[j]} = \text{null}$ for all $j \neq 1$
snd-flag	$\in \{\text{false}, \text{true}\},$	initially: $\text{snd-flag}^{[1]} = \text{true},$ $\text{snd-flag}^{[j]} = \text{false}$ for $j \neq 1$

function **msg**(w, i)

```

1: if (parent  $\neq i$ ) AND (snd-flag = true) then
2:   return data
3: else
4:   return null

```

function **stf**(w, y)

```

1: case
2:   (data = null) AND (y contains only null messages):
      % The node has not yet received the token
3:     new-parent := null
4:     new-data := null
5:     new-snd-flag := false
6:   (data = null) AND (y contains a non-null message):
      % The node has just received the token
7:     new-parent := smallest UID among transmitting in-neighbors
8:     new-data := a non-null message
9:     new-snd-flag := true
10:  (data ≠ null):
      % If the node already has the token, then do not re-broadcast it
11:    new-parent := parent
12:    new-data := data
13:    new-snd-flag := false
14: return (new-parent, new-data, new-snd-flag)

```
