

Boosting Typical Query Patterns

PostgreSQL 18's Performance Enhancements

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25+ years experience with PostgreSQL and databases

- Previous: Principal Solutions Architect @ Crunchy Data
- Managed DBA and DevOps teams
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Personal

- Brazilian (Manaus) expatriate in Utah
- Interests: more computer stuff
- Photography, Snowboarding, Austrian Economics

Introduction
Asynchronous I/O
Index Improvements
Query Planning & Observability
dditional Features & Real Applications
Summary & Conclusion

About Crunchy Data and Snowflake Talk Overview





https://www.crunchydata.com/blog/crunchy-data-joins-snowflake

Talk Overview

- PostgreSQL 18 introduces significant performance improvements
- Foundational changes and real-world query patterns

What We'll Cover:

- Asynchronous I/O subsystem
- B-tree skip scans
- Parallel GIN index creation
- Query optimizer improvements
- EXPLAIN enhancements
- UUID v7 performance
- Real-world Applications

Traditional I/O

Traditional PostgreSQL I/O:

Backend: "I need page 1000"

Kernel: [reads page 1000]

Backend: [waits...]

Kernel: "Here's page 1000"

Backend: "Thanks! Now I need page 1001"

Kernel: [reads page 1001]

Backend: [waits...]

Inefficient because:

- One request at a time
- Backend idle while waiting for I/O
- Can't batch or parallelize requests
- Underutilizes modern storage (NVMe, SSD)

Postgres 17 paved the way with the introduction of read stream and vectored I/O APIs, internal abstractions. See Andres Freund https://youtu.be/qX50xrHwQa4

Async I/O

PostgreSQL 18 Async I/O:

Backend: "I need pages 1000, 1001, 1002, 1003..."

Kernel: [queues all requests]
Backend: [continues other work]

Kernel: [returns pages as ready]
Backend: [processes completed I/O]

Benefits:

- Batch multiple I/O requests
- Kernel can optimize request ordering
- Better utilization of parallel storage
- Backend does useful work while waiting

Async I/O: Configuration

GUC Parameters:

```
# 17+: Control I/O batching (blocks of 8kb)
io_combine_limit = 16  # requests per batch (128kb default)

# 18: I/O method selection
io_method = 'worker'  # default, usually performs better.

# 18: Largest I/O size in operations that combine I/O (blocks of 8kb, default is 16 = 128kB)
# silently limits io_combine_limit. Typically 1MB on Unix and 128kB on Windows.
io_max_combine_limit = 16
```

Operations Most Benefitted:

- Sequential scans of large tables
- Bitmap heap scans (multi-index queries)
- VACUUM operations

Async I/O: Benchmark Results

Async I/O Performance Comparison



Key Findings:

Sequential scans: 15-25% faster

Bitmap heap scans: 10-18% faster

VACUUM: 20-300% faster

Async I/O: Real-World Impact

Example: Analytics Query

```
SELECT category, COUNT(*), AVG(amount)
FROM large_orders
WHERE created_at >= '2025-01-01'
GROUP BY category;
```

PG16: 12.3 seconds (sequential scan)

PG18: 9.8 seconds (sequential scan with async I/O)

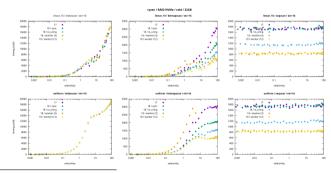
20% improvement without code changes

Async I/O: Tuning for Production

```
TID PRIO USER
                       DISK READ DISK WRITE>
                                               COMMAND
 605556 be/4 999
                        0.00 B/s
                                  16.73 M/s postgres: postgres postgres 10.10.0.208(50254) VACUUM
                        0.00 B/s 63.44 K/s postgres: walwriter
604968 be/4 999
604957 be/4 999
                      66.72 M/s 0.00 B/s postgres: io worker 1
604958 be/4 999
                      102.11 M/s 0.00 B/s postgres: io worker 0
604959 be/4 999
                 46.70 M/s 0.00 B/s postgres: io worker 2
                 23.00 M/s 0.00 B/s postgres: io worker 4
604960 be/4 999
604961 be/4 999
                    33.70 M/s
                                   0.00 B/s postgres: io worker 3
io_method = worker # default: pool of I/O worker processes
                  # io_uring: Linux-specific asvnc I/O queues
                  # sync: traditional synchronous I/O (backwards compatibility)
io workers = 3
                  # Default too low for larger systems
                  # Probably: set to approximately 1/4 of total CPU cores
io max combine limit = 16 # limits io combine limit.
                          # typical max: Unix 128 (1 MB), Windows 16 (128 kB)
                          # requests per batch (128kb default)
io combine limit = 16
                          # raise with io max combine limit to increase the I/O size
```

Async I/O: Performance Comparison

Query Timing by I/O Method¹. Benchmark: Ryzen 9900X (12 cores/24 threads), 4x NVMe SSDs (RAID0)



¹See: Tomas Vondra, https://vondra.me/posts/tuning-aio-in-postgresql-18/

Async I/O: Tuning Recommendations

- Keep default io_method = worker
 - Best compatibility across workloads. io_uring is Linux-specific
- Increase io_workers based on cores
 - Start with 1/4 of CPU cores. Monitor and adjust based on workload
- Test with your workload
 - Performance varies by query patterns. Bitmap scans benefit most
- Watch out for
 - Signal overhead between backends and workers
 - File descriptor limits
 - I/O bandwidth saturation

Skip Scan

Acceptable performance with seldom-run queries that might not require a dedicated index.

```
CREATE INDEX idx_country_user
  ON orders(country, user_id);
-- Query on second column only
SELECT * FROM orders WHERE user_id = 12345;
```

Before PG18:

- Index not usable (query doesn't start with country)
- Falls back to sequential scan
- Slow on large tables
- Advice: "Create a single-column index on user_id"

Skip Scan: The Solution

PG18: Skip Scan to the rescue

- Planner recognizes opportunity to use multi-column index
- "Skips" over distinct values of leading column
- For each distinct country, searches for user_id
- Most effective when leading column has low cardinality

Example:

- 10 distinct countries (low cardinality)
- 1M distinct user_ids (high cardinality)
- Skip scan does 10 targeted index searches
- Much faster than sequential scan

B-tree Skip Scan
Parallel GIN Index Creation

Skip Scan

PostgreSQL 16: (without single-column index on created_at)

```
EXPLAIN (ANALYZE, TIMING) SELECT id, project id, created at FROM skip scan test
WHERE created at > NOW() - INTERVAL '30 days' ORDER BY created at DESC LIMIT 100:
Limit (cost=93265.53..93277.50 rows=100 width=20) (actual time=455.743..466.308 rows=100 loops=1)
   -> Gather Merge (cost=93265.53..135927.94 rows=356308 width=20) (actual time=455.741..466.277 rows=100 loops=1)
        Workers Planned: 4
        Workers Launched: 4
         -> Sort (cost=92265.47..92488.16 rows=89077 width=20) (actual time=450.614..450.619 rows=90 loops=5)
              Sort Key: created at DESC
              Sort Method: top-N heapsort Memory: 36kB
              Worker 0: Sort Method: top-N heapsort Memory: 36kB
              Worker 1: Sort Method: top-N heapsort Memory: 37kB
              Worker 2: Sort Method: top-N heapsort Memory: 37kB
              Worker 3: Sort Method: top-N heapsort Memory: 37kB
              -> Parallel Seq Scan on skip_scan_test (cost=0.00..88861.01 rows=89077 width=20) (actual time=0.024..440.729 rows=699
                    Filter: (created at > (now() - '30 days'::interval))
                    Rows Removed by Filter: 930074
Planning Time: 0.116 ms
 Execution Time: 466.354 ms
```

Skip Scan

PostgreSQL 18: (uses composite index with skip scan)

```
QUERY PLAN
```

```
Limit (cost=78397..78409 rows=100) (actual time=151..173 rows=100)

-> Gather Merge (cost=78397..122040 rows=364500)

Workers Planned: 4

Workers Launched: 4

-> Sort (cost=77397..77624 rows=91125)

-> Parallel Bitmap Heap Scan on skip_scan_test

Recheck Cond: (created_at > ...)

-> Bitmap Index Scan on idx_skip_scan_composite

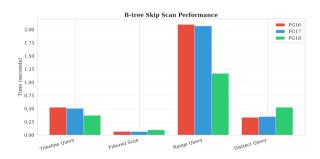
Index Cond: (created_at > ...)

Index Searches: 1001

Execution Time: 173.719 ms
```

Faster using composite index (project_id, created_at) via skip scan

Skip Scan: Benchmark Results



- Biggest improvement with low-cardinality leading column
- Effective for narrow range queries on composite indexes
- Best for queries on non-leading columns with selective filters

Parallel GIN: Background

GIN Indexes: Generalized Inverted Index

- Used for arrays, JSONB, full-text search
- Can be slow to build on large tables (higher maintenance_work_mem helps)

Postgres 18

Parallel index builds available for GIN, in addition to B-tree, BRIN

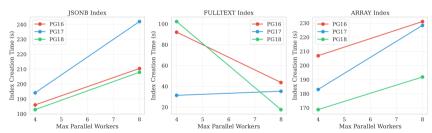
Parallel GIN: Configuration

Configuration:

```
SET max_parallel_maintenance_workers = 4;
-- Create index (automatically uses parallel workers)
CREATE INDEX idx_tags ON posts USING GIN(tags);
  TID PRIO< USER
                    DISK READ DISK WRITE
                                          COMMAND
604957 be/4 999
                   377.75 K/s 0.00 B/s postgres: io worker 1
604958 be/4 999
                  47.22 K/s 0.00 B/s postgres: io worker 0
               39.35 K/s 0.00 B/s postgres: io worker 2
604959 be/4 999
604962 be/4 999
                    31.48 K/s 31.48 K/s postgres; checkpointer
604965 be/4 999
                15.74 K/s
                              7.87 K/s postgres: checkpointer
610510 be/4 999
                15.74 K/s 25.95 M/s postgres: parallel worker for PID 2032
610511 be/4 999
                   129.85 K/s 0.00 B/s postgres: parallel worker for PID 2032
610512 be/4 999
                    7.87 K/s 0.00 B/s postgres: parallel worker for PID 2032
```

Parallel GIN: Benchmark Results

Parallel GIN Index Creation Performance



Results (8-core test system):

- 4 workers: Best performance for JSONB & Array indexes
- 8 workers: Performance degradation (12-25% slower)

 $\textbf{Important:} \ \texttt{Set} \ \texttt{max_parallel_maintenance_workers} \leq \texttt{CPU} \ \texttt{cores}$

Optimizer Improvements Hash Join & GROUP BY Self-join Removal EXPLAIN Improvements

Optimizer: Multiple Enhancements

- Hash Join & GROUP BY improvements
- **IN (VALUES)** \rightarrow **= ANY** transformation
- OR clauses → array operations for indexable queries
- Unnecessary self-join removal
- Speed up of INTERSECT/EXCEPT, window aggregates, view column aliases
- SELECT DISTINCT internal reordering to avoid sorting

Mostly no query changes needed.

Optimizer: IN (VALUES) Performance

Postgres 18 adds nbtree skip scan building on Postgres 17 work on IN() / = ANY() condition index scans

- items that are close together (1,2,3) or far apart (10_000, 20_000)
- Supports complex combinations of IN() conditions, = conditions, as well as <, >,
 <=, => conditions
- Only reads index leaf pages that might have matches

Optimizer: OR to Array

Rewrite OR conditions to better use indexes

```
SELECT * FROM products
WHERE category = 'electronics'
OR category = 'clothing'
OR category = 'food';
```

PG18 Optimization:

```
SELECT * FROM products
WHERE category = ANY (ARRAY['electronics', 'clothing', 'food']);
```

Can use bitmap index scans more efficiently - reports of 100x improvement.

Optimizer: Hash Join & GROUP BY Improvements

- Hash Right Semi Join support
 - Planner can now choose which table to hash based on size
 - Previously constrained to hashing inner table only
 - 40% reduction in memory usage for large datasets
- 3 JIT-compiled hash value generation
 - Hashing for GROUP BY and hashed subplans
 - Enables JIT compilation of hash values
 - Faster hash value computation during execution

Improved performance and reduced memory for hash joins, GROUP BY, EXCEPT, and subplan hash lookups

Hash Join & GROUP BY: Real-World Example

Query Pattern: Semi-join with GROUP BY aggregation

```
-- Find flights with at least one ticket sold
SELECT f.flight_id, f.flight_no, COUNT(DISTINCT tf.ticket_no)
FROM flights f
WHERE EXISTS (
    SELECT 1 FROM ticket_flights tf WHERE tf.flight_id = f.flight_id
)
GROUP BY f.flight_id, f.flight_no;
```

PG17: Uses Hash Semi Join, must hash larger ticket_flights table (2.3s)

PG18: Uses Hash Right Semi Join, hashes smaller flights table (<1s)

50%+ faster with 40% less memory usage

Optimizer: Self-Join Removal

```
SELECT t1.*
FROM orders t1 JOIN orders t2 ON t1.id = t2.id
WHERE t1.status = 'completed';
```

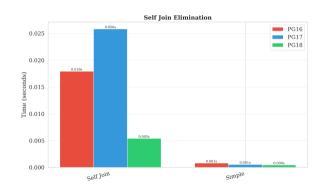
Postgres 18:

- Detects redundant self-join and removes t2 table
- Executes as simple SELECT * FROM orders WHERE status = 'completed'

Common in:

- ORM-generated queries
- View definitions
- Query builder tools

Optimizer: Self-Join Removal



Self-Join Removal: Up to 5x faster in PG18

EXPLAIN: What's New

- Automatic BUFFERS in EXPLAIN ANALYZE
- EXPLAIN ANALYZE VERBOSE shows hardware stats (CPU, Memory, I/O)
- Per connection stats on I/O and WAL utilization
- Better observability by default
- Easier performance troubleshooting

```
EXPLAIN ANALYZE SELECT * FROM orders WHERE value < 100;

Seq Scan on orders (cost=0.00..1834.00 rows=98.3 ...)

Filter: (value < 100)

Rows Removed by Filter: 9902

Buffers: shared hit=834 read=0
```

Optimizer Improvements Hash Join & GROUP BY Self-join Removal EXPLAIN Improvements

EXPLAIN: What's New

```
=# explain (analyze, wal, timing, memory, verbose)
-# select * from uuid v4 test where id <= '00036a6a-6218-4f4b-8bfa-fd2dce5e1443'
-\# and id >= '00021dd6-79de-426e-809d-440a0160504d' order by id:
                                           QUERY PLAN
Index Scan using unid v4 test pkey on public unid v4 test (cost=0.43..2.65 rows=1 width=90)
                                              (actual time=0.023..0.278 rows=182.00 loops=1)
   Output: id. user_id. event_type, data, created_at
   Index Cond: ((uuid v4 test.id <= '00036a6a-6218-4f4b-8bfa-fd2dce5e1443'::uuid) AND
                (uuid v4 test.id \ge '00021dd6-79de-426e-809d-440a0160504d'::uuid))
   Index Searches: 1
   Buffers: shared hit=186
Query Identifier: 5337184330030628219
Planning:
   Buffers: shared hit=8
  Memory: used=15kB allocated=32kB
Planning Time: 0.122 ms
 Execution Time: 0 299 ms
```

Optimizer Improvements Hash Join & GROUP BY Self-join Removal EXPLAIN Improvements

EXPLAIN: What's New

UUID v7: Time-Ordered UUIDs

Problems with UUID v4:

- Completely random values
- Poor index locality
- Index bloat and fragmentation
- Slower inserts as table grows

UUID v7 (RFC 9562):

- First 48 bits: timestamp (millisecond precision)
- Remaining bits: random
- Time-ordered like SERIAL, but globally unique
- Better B-tree performance, with less disk use

UUID v7: Usage

```
-- New function for UUID v7

CREATE TABLE users (
  id UUID PRIMARY KEY DEFAULT uuidv7(),
  email TEXT,
  created_at TIMESTAMP DEFAULT NOW()
);

-- Inserts are naturally ordered by time
INSERT INTO users (email)

VALUES ('user@example.com');
```

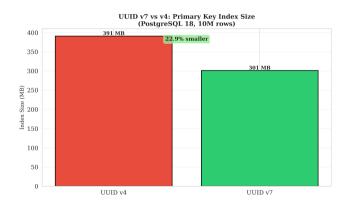
Benefits:

- Smaller indexes (better locality)
- Can infer creation time from UUID

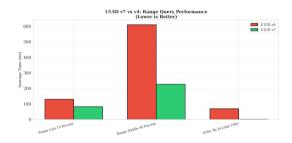
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UUID v7 Statistics Retention GitLab Discourse

UUID v7: Index Size



UUID v7: Range Query Performance



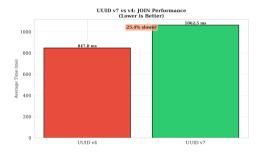
Faster for larger result sets

- Better index locality improves sequential scans
- Most beneficial for queries returning many rows
- Excellent for time-based queries

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UUID v7: JOIN Performance



SELECT e.id, e.event_type, d.score
FROM uuid_v7_test e JOIN uuid_v7_test_details d
 ON e.id = d.event_id LIMIT 1000;

- Best for insert-heavy, time-series data
- Less ideal for JOIN-heavy OLTP workloads

UUID v7 Statistics Retention GitLab Discourse

Statistics Retention Across Upgrades

The Problem:

- Major version upgrades lose optimizer statistics
- First queries after upgrade are slow
- Must wait for autovacuum to collect stats
- Production vs staging plan differences

PG18

- New pg_dump --statistics-only
- Functions to restore statistics
- Preserve query plans across upgrades
- Copy production stats to dev/test

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UUID v7 Statistics Retention GitLab Discourse

Statistics Retention: Usage

pg_dump --statistics-only mydb > stats.sql

```
SELECT * FROM pg_catalog.pg_restore_relation_stats(
        'version', '180000'::integer,
        'schemaname', 'public'.
        'relname', 'async io test'.
        'relpages', '80777'::integer,
        'reltuples', '1.004389e+06'::real,
        'relallvisible', '80777'::integer,
        'relallfrozen', '3768'::integer );
SELECT * FROM pg_catalog.pg_restore_attribute_stats(
        'version', '180000'::integer,
        'schemaname', 'public',
        'relname'. 'asvnc io test'.
        'attname', 'created at'.
        'inherited', 'f'::boolean,
        'null_frac', '0'::real,
        'avg_width', '8'::integer,
        'n_distinct', '116'::real,
        'most_common_vals', '{"2025-10-15 02:54:23.838692","2025-10-15 02:54:29.491851", ...
        'most_common_freqs', '{0.103533335.0.1018.0.10146666.0.1003.0.09996667.0.09893333....
        'histogram bounds', '{"2025-10-15 02:55:00.91417","2025-10-15 02:55:15.533473"....
```

GitLab



- Total execution time: 54% faster (PG18: 10.7s vs PG16: 21.5s)
- Query count: Consistent across versions (~14,400 queries)

Discourse

2000

PG16

PG17

PG18



PG17

- Total execution time: 14% faster (PG18: 1.43s vs PG16: 1.62s)
- **Query count:** Consistent across versions (\sim 9,000 queries)

500

250

PG16

PG16

PG17

PG18

PG18

Real-World Use Cases

Who Benefits Most?

- Analytics Workloads: Async I/O, optimizer improvements
- SaaS Applications: UUID v7, skip scans
- E-commerce: Parallel GIN, optimizer
- Content Platforms: Full-text search (parallel GIN)
- Multi-tenant Apps: Skip scans on tenant_id indexes

Key Takeaways

- PostgreSQL 18 brings measurable performance gains
- Improvements are mostly automatic
- Async I/O: foundational infrastructure-level improvement
- Skip Scan: acceptable performance for multi-column indexes and less-frequent queries
- Parallel GIN: faster index builds
- Optimizer: smarter query planning

Additional Resources

- PostgreSQL 18 Release Notes: https://www.postgresql.org/docs/current/release-18.html
- Tomas Vondra at https://vondra.me/posts/tuning-aio-in-postgresql-18/
- Async I/O Deep Dive: https://pganalyze.com/blog/postgres-18-async-io
- Crunchy Data Blog: https://crunchydata.com/blog/get-excited-about-postgres-18
- Jonathan Katz, Peter Geoghegan: https://www.slideshare.net/slideshow/postgresql-18-a-whirlwind-tour-of-features/283259854

Real-World Use Cases Key Takeaways Resources

Questions

Feedback:



These slides:

