# Uninformed Search (Blind Search)

Boris Velichkov

#### Exercises by Week

- 1. Introduction
- 2. Uninformed (Blind) Search
- 3. Informed (Heuristic) Search
- 4. Constraint Satisfaction Problems
- 5. Genetic Algorithms
- 6. Games
- 7. <u>Introduction to Machine Learning</u>
- 8. k-Nearest Neighbors
- 9. Naïve Bayes Classifier
- 10. Decision Tree
- 11. kMeans
- 12. Neural Networks
- 13. Additional Topics, Questions and Homeworks' Presentations
- 14. Additional Topics, Questions and Homeworks' Presentations
- 15. Additional Topics, Questions and Homeworks' Presentations

#### Problem Solving and Search

- Basic Concepts
  - State Space
  - Representation of the State Space
  - Search Strategies Evaluation
  - Global Search vs Local Search
  - Uninformed Search vs Informed Search
  - Graph and Tree Traversal

#### State Space

- State: a representation (formulation) of the task in the process of its solution.
  - Initial State (presented with S)
  - Intermediate State (presented with any capital letter except S and G)
  - Goal State (presented with G, if there are more:  $G_1$ ,  $G_2$ , etc.)
- Successor Function (Operator): obtaining one state from another
- Path Cost: additive; e.g., sum of distances, number of actions executed, etc.
- **State Space**: the totality of all possible states that can be obtained from a given initial state.

<sup>\*</sup> A **solution** is a sequence of actions leading from the initial state to a goal state

# Representation of the State Space

- Graph or Tree where each state is represented with Node and the successor function (operator) is represented with Edge.
- When the state space can be represented as a tree, it is called Search Tree.
  - The initial state is the root of the tree.
  - The terminal states and the goal state are represented with leaves.

#### Search Strategies Evaluation

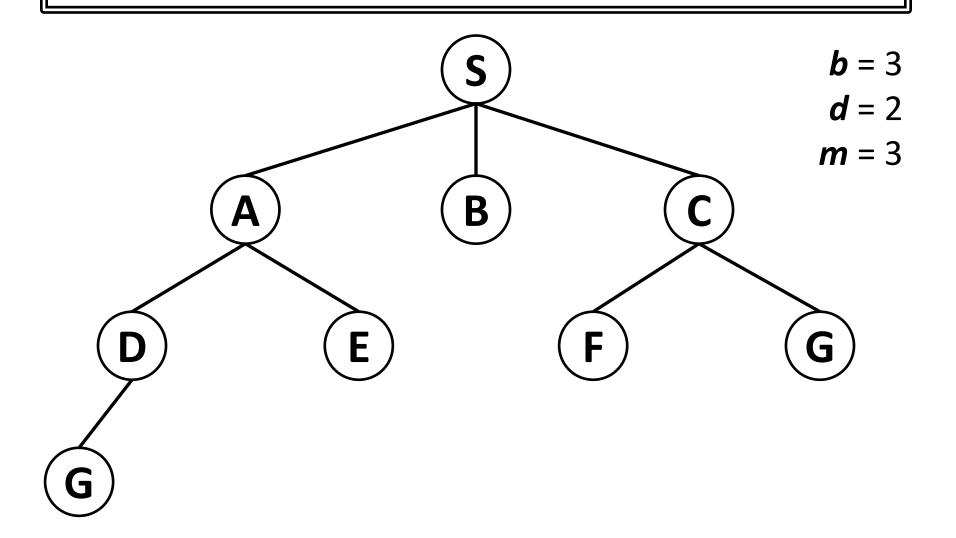
Strategies are evaluated along the following dimensions:

- Completeness: does it always find a solution if one exists?
- Optimality: does it always find a least-cost solution?
- Time Complexity: number of nodes generated/expanded
- Space Complexity: maximum number of nodes in memory
- \* We look at the **worst-case** complexity (Big O notation).

Time and space complexity are measured in terms of:

- b maximum branching factor of the search tree
- d depth of the least-cost solution
- m maximum depth of the state space (may be ∞)

# Search Strategies Evaluation



#### Global Search vs Local Search

- Global Search: They look all over the state space. If necessary, all states will be traversed.
- Local Search: They only look at the local area, so they can only look at states in this area. If the solution is outside of it, they will not be able to find it.

<sup>\*</sup> Generally Local Search is not used for finding a path from state A to state B.

#### Local Search

- Local Search in Artificial Intelligence is an optimizing algorithm to find the optimal solution more quickly.
- Local search algorithms are used when we care only about a solution but not the path to a solution.
- Local search is a heuristic method for solving computationally hard optimization problems.
- Local search can be used on problems that can be formulated as finding a solution maximizing a criterion among a number of candidate solutions.
- Local search algorithms move from solution to solution in the space of candidate solutions (the search space) by applying local changes, until a solution deemed optimal is found or a time bound is elapsed.

#### Uninformed (*Blind*) Search vs Informed (*Heuristic*) Search

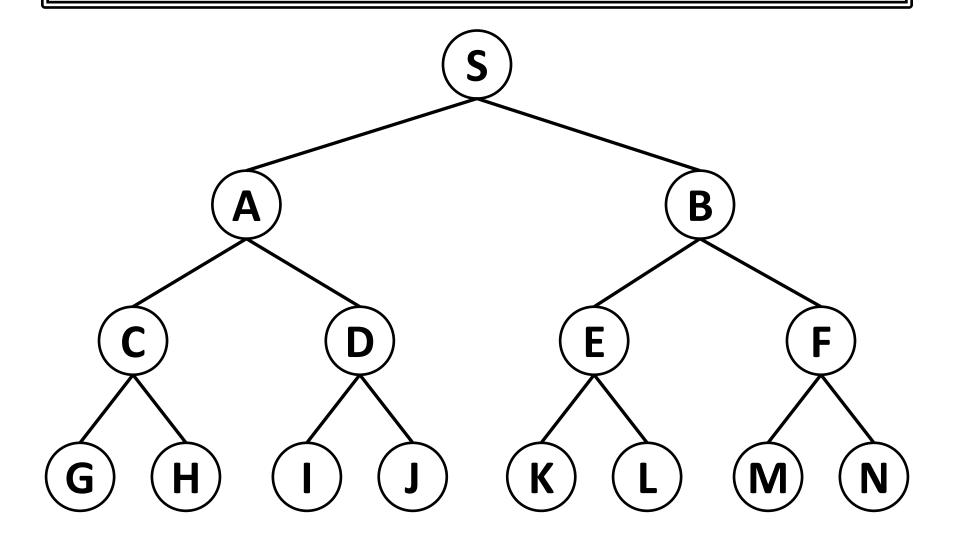
- Uninformed Search: Uninformed strategies use only the information available in the problem definition.
  - Examples: DFS, BFS, UCS, DLS, IDS.
- Informed Search: Informed strategies have information on the *goal state* which helps in more efficient searching. This information is obtained by a function (heuristic) that estimates how close a state is to the *goal state*.
  - Examples: Greedy Best-First Search, A\*, Beam Search, Hill Climbing.

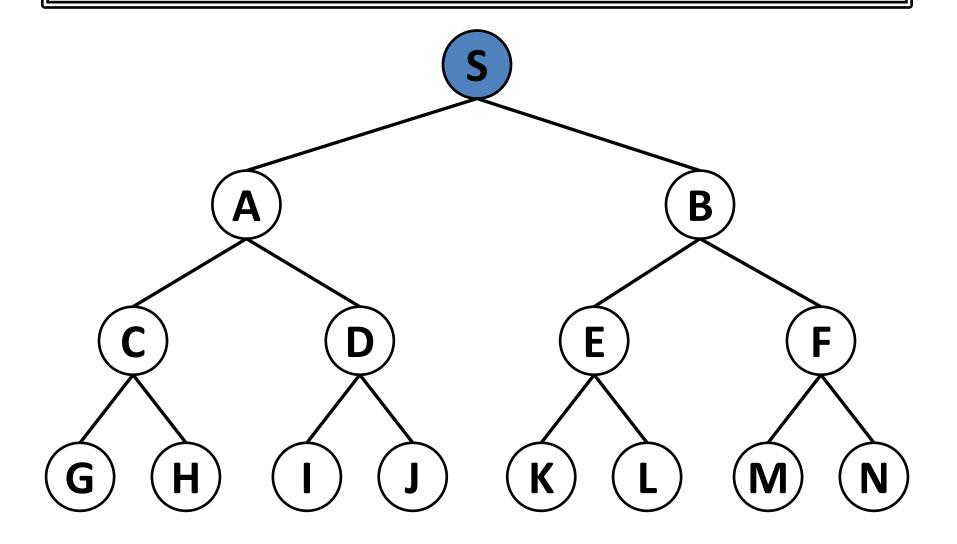
 Node: The graphs and trees consist of nodes and edges. When a node is neither "visited" nor "expanded", we do not mark it in any way (sits uncolored).

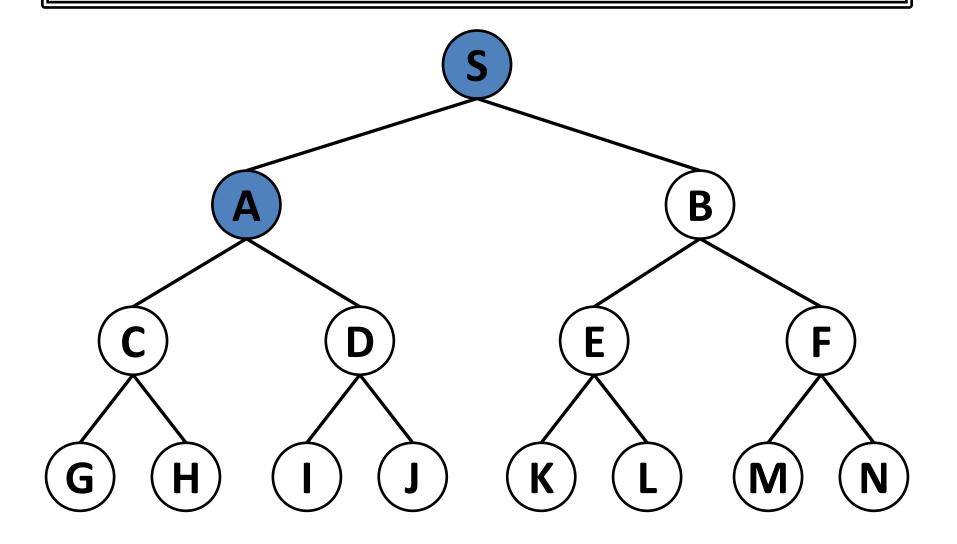
- Example:

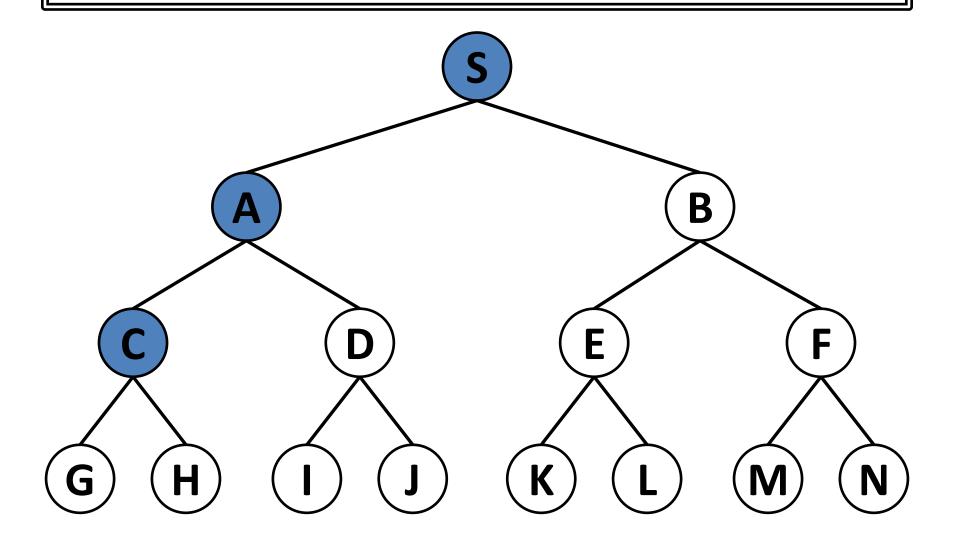
- **Visited Node**: When we visit a *node* or in other words add it to the data structure we use we call it "visited node". We mark it in blue.
  - Example: S
- Expanded Node: When we expand a node or in other words remove it from the data structure we use and add its children in this data structure (of course if the node has children) we call it "expanded node". We mark it in red.
  - Example:

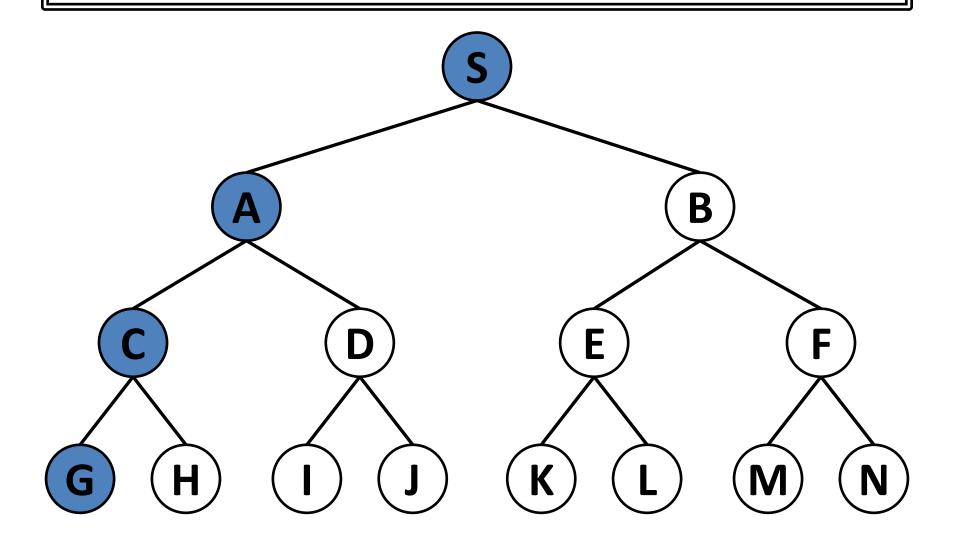


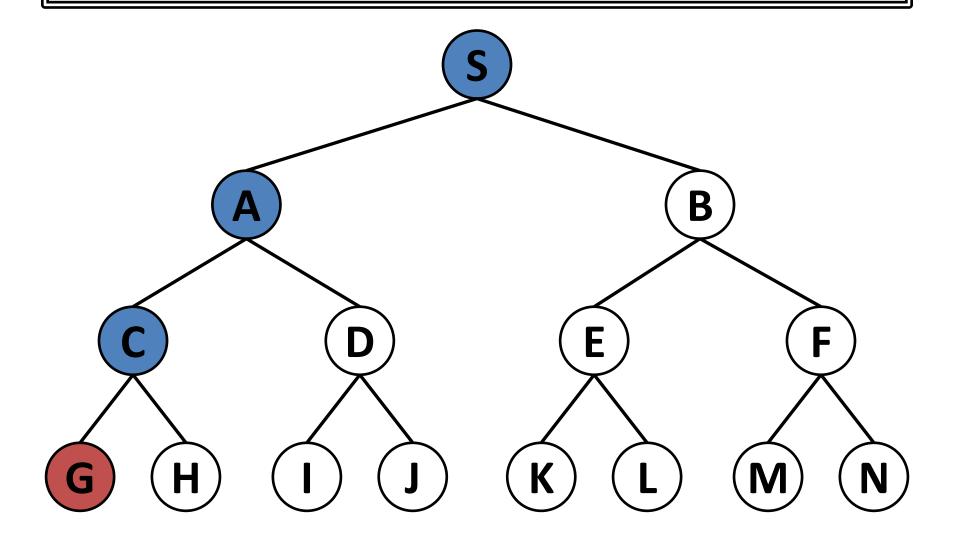


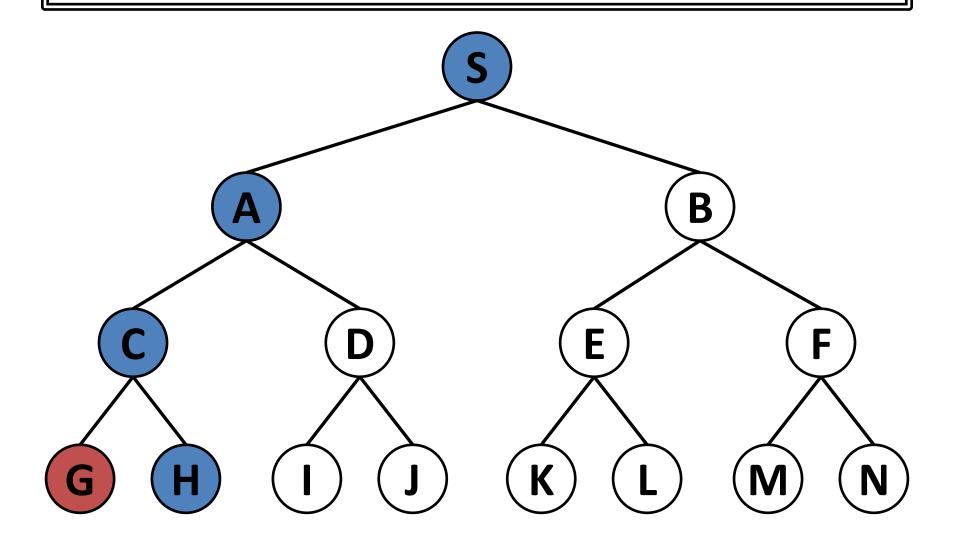


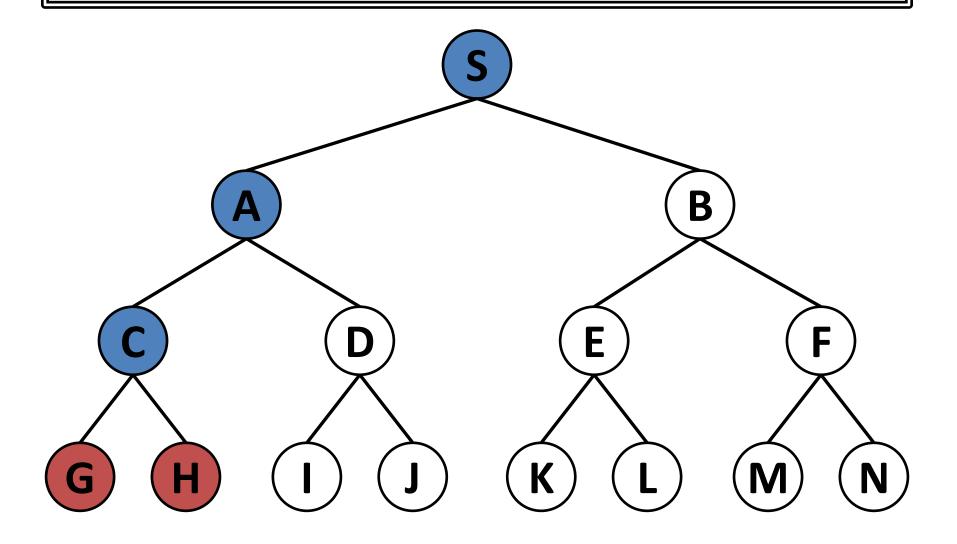


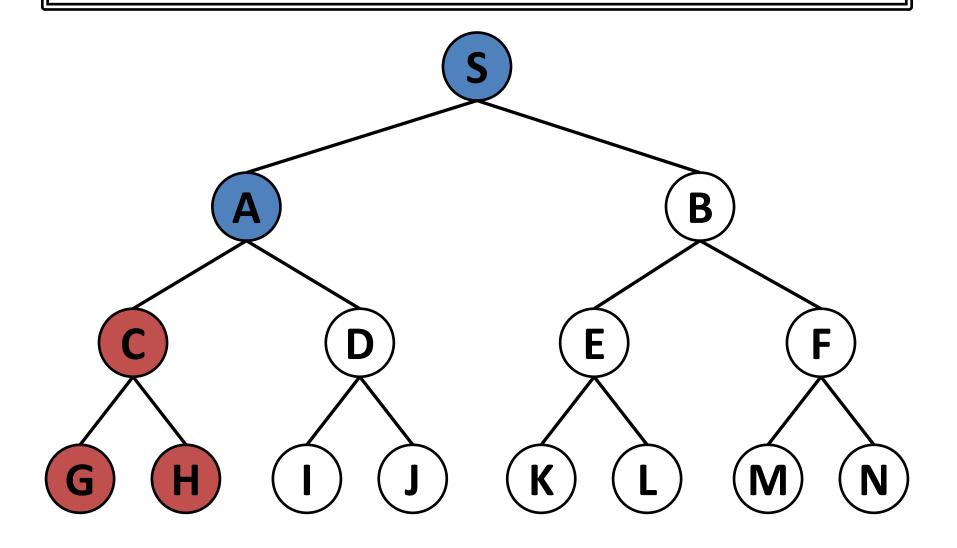


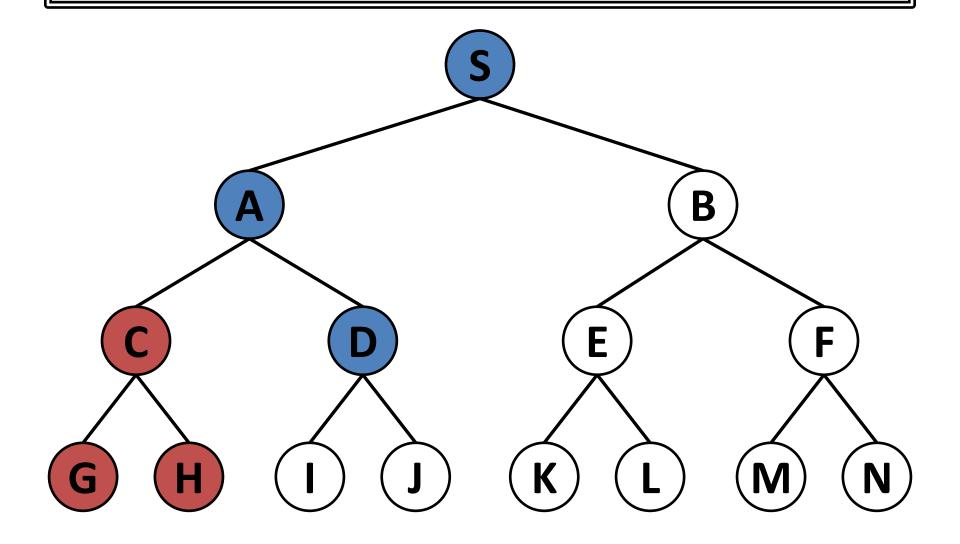


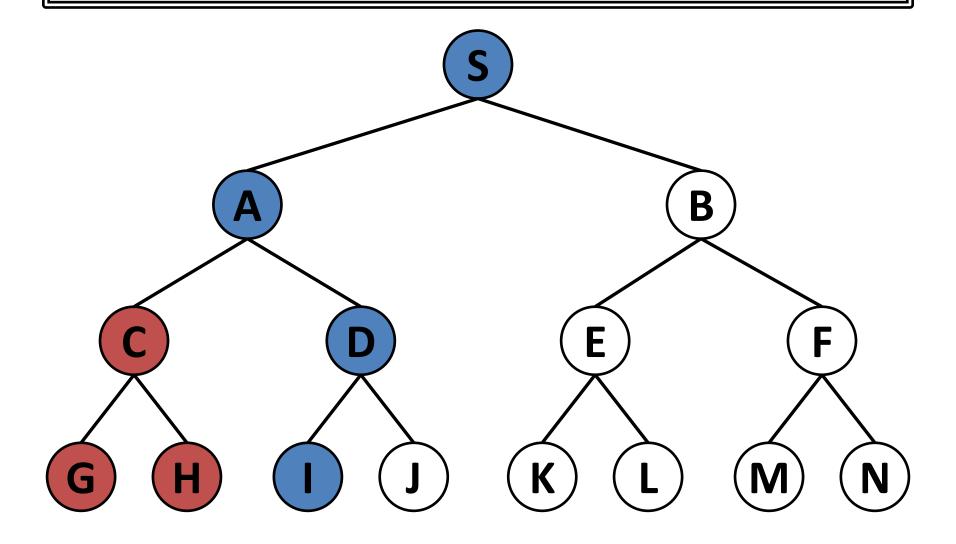


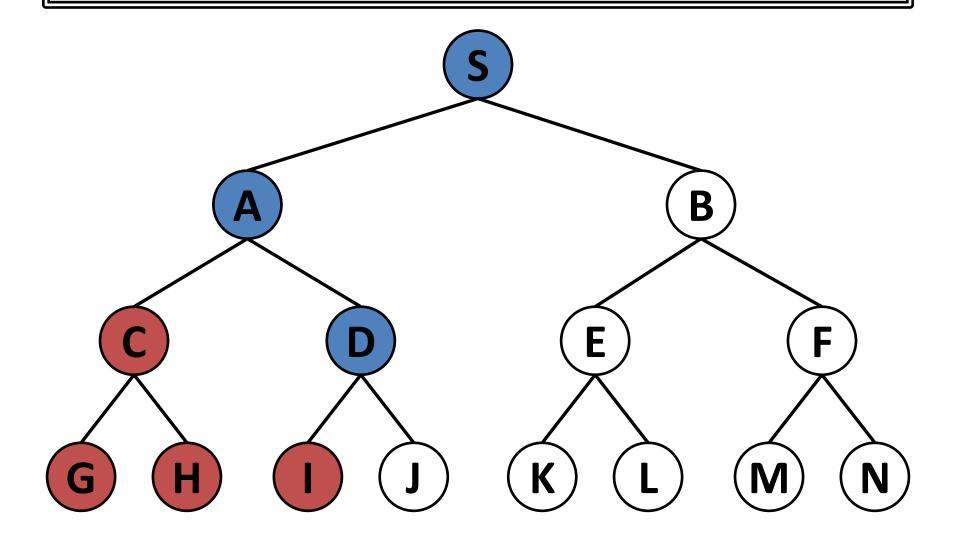


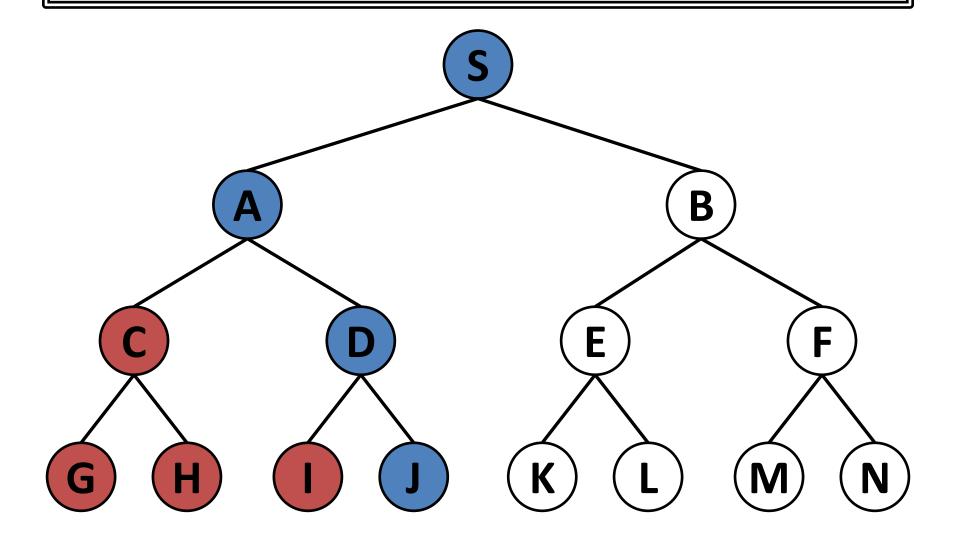


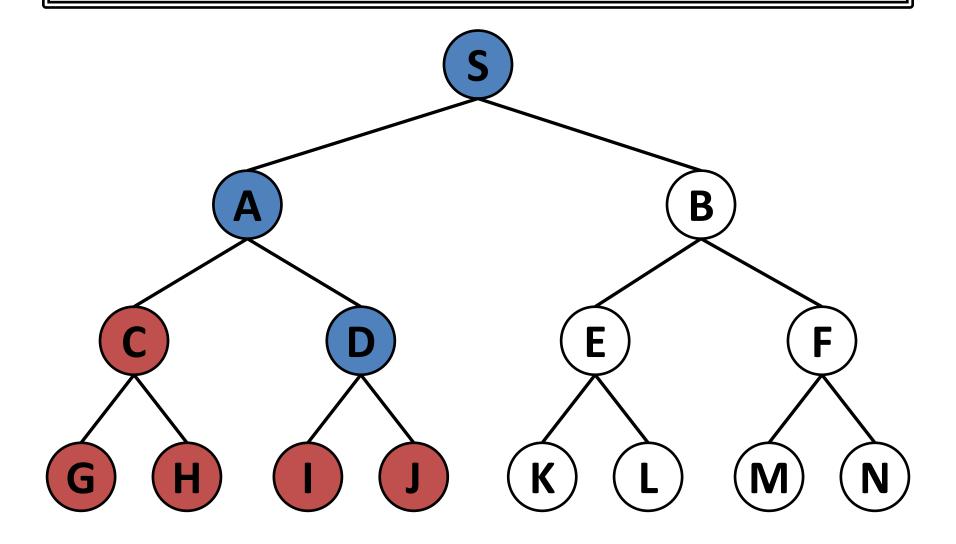


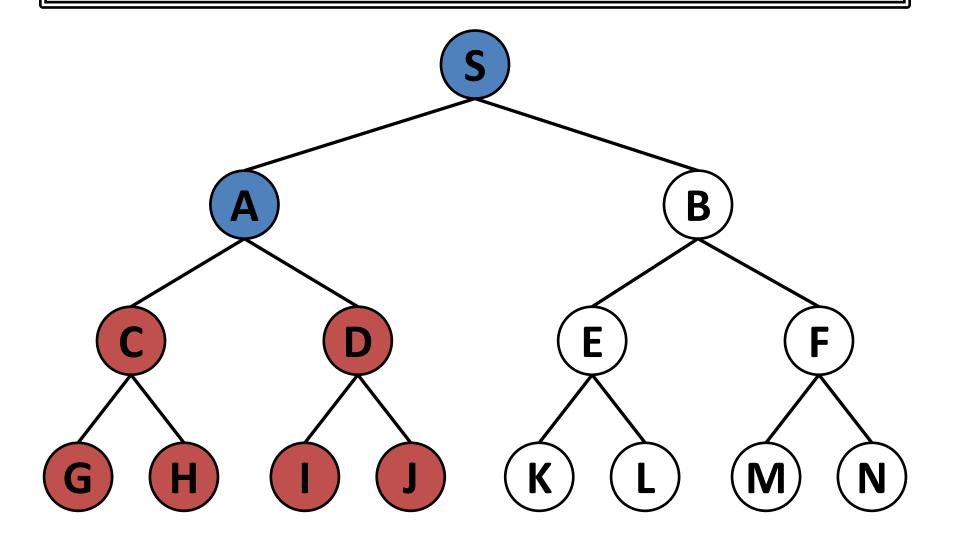


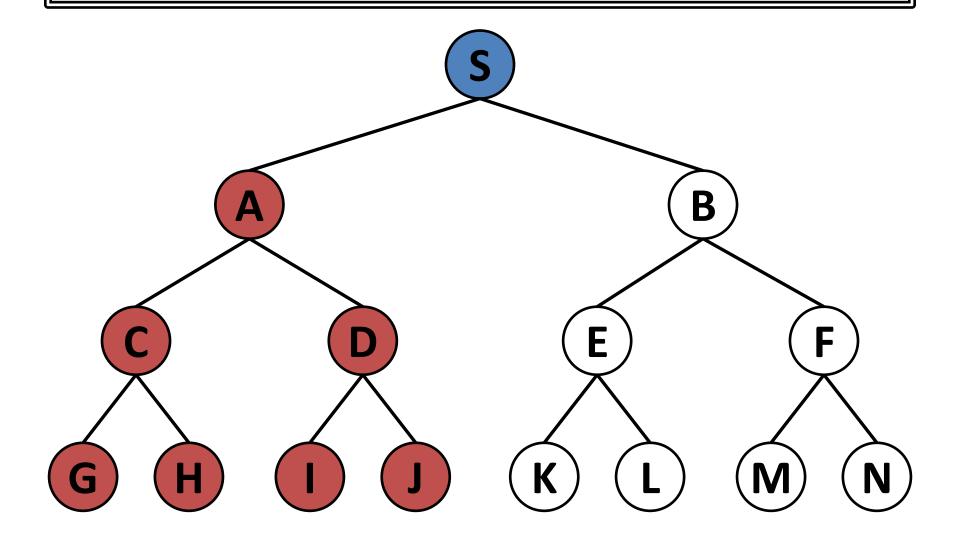


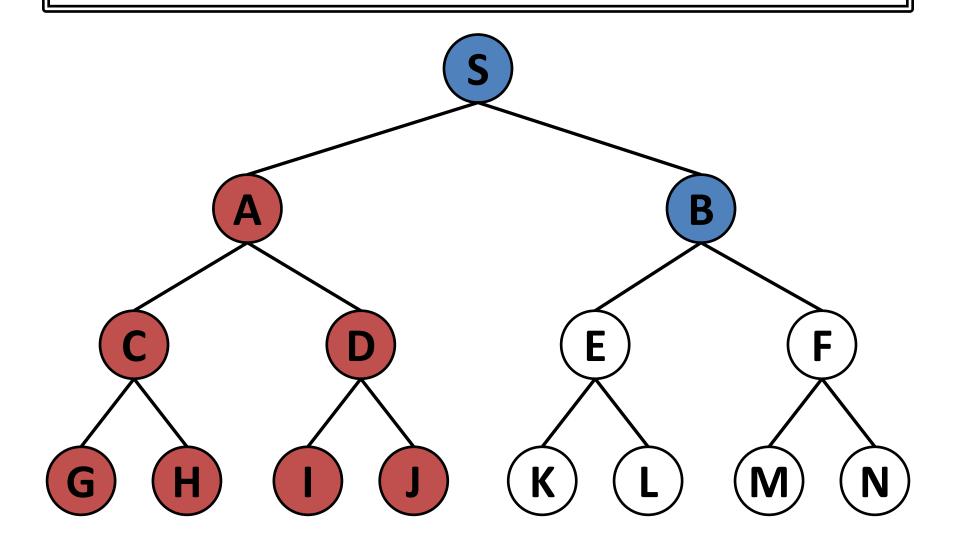


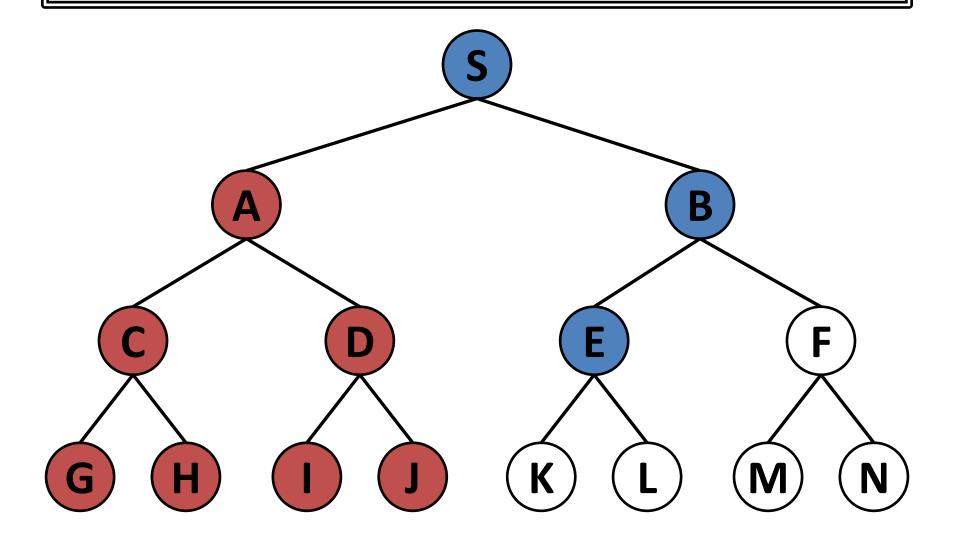


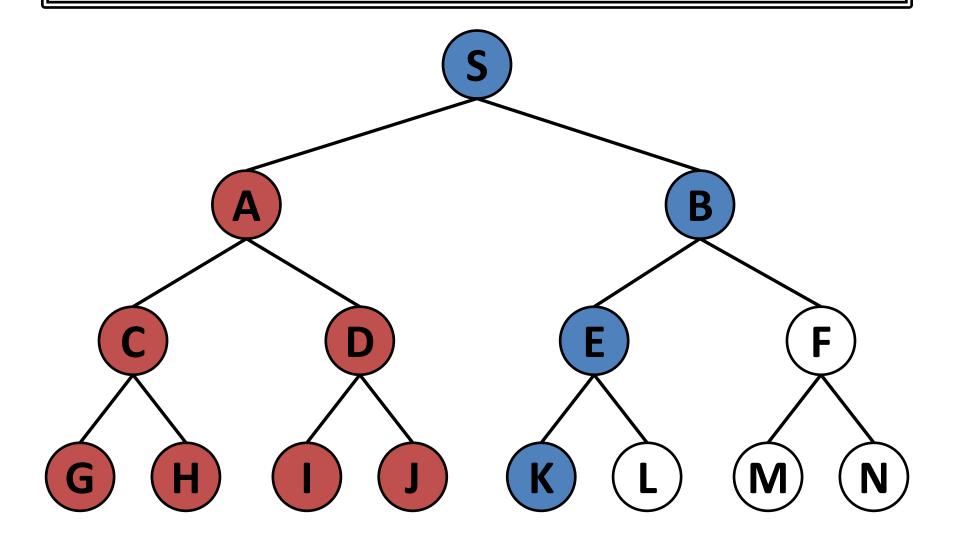


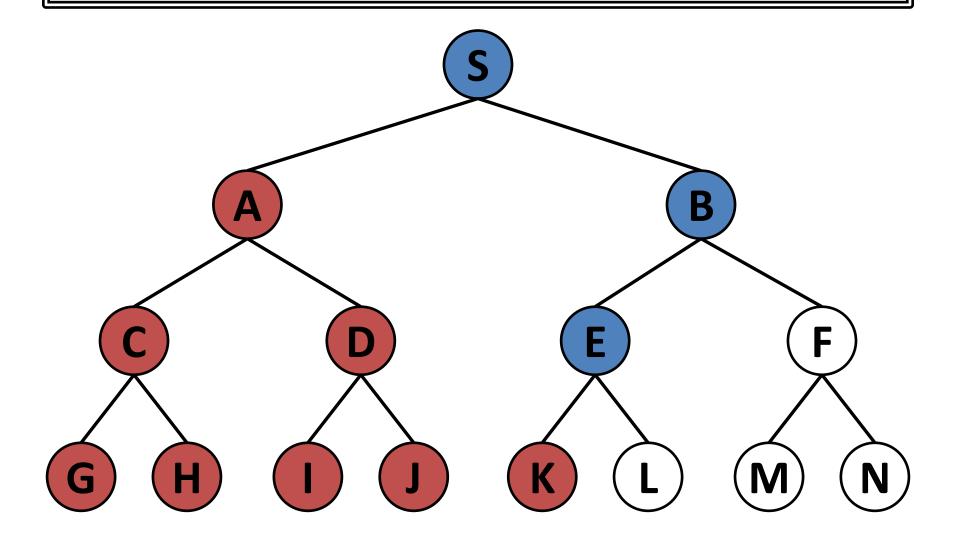


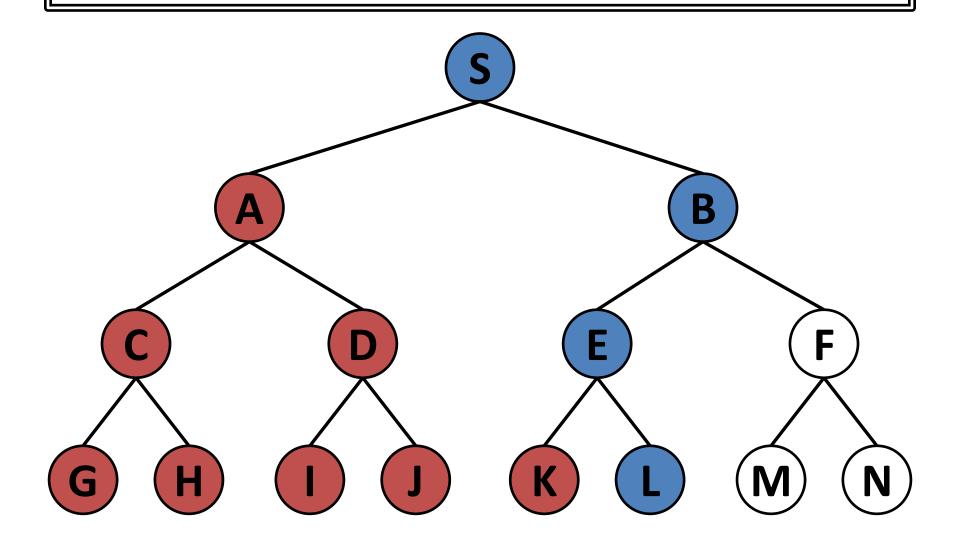


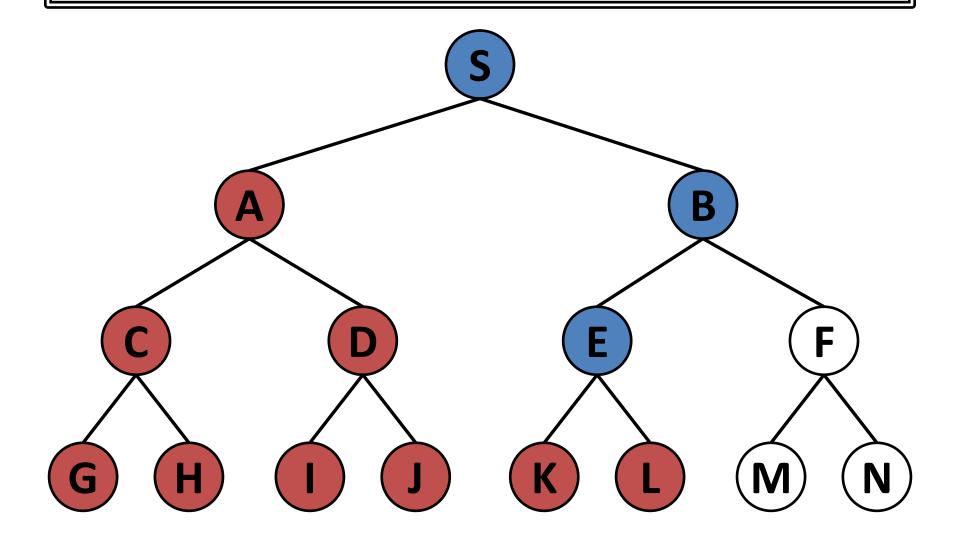


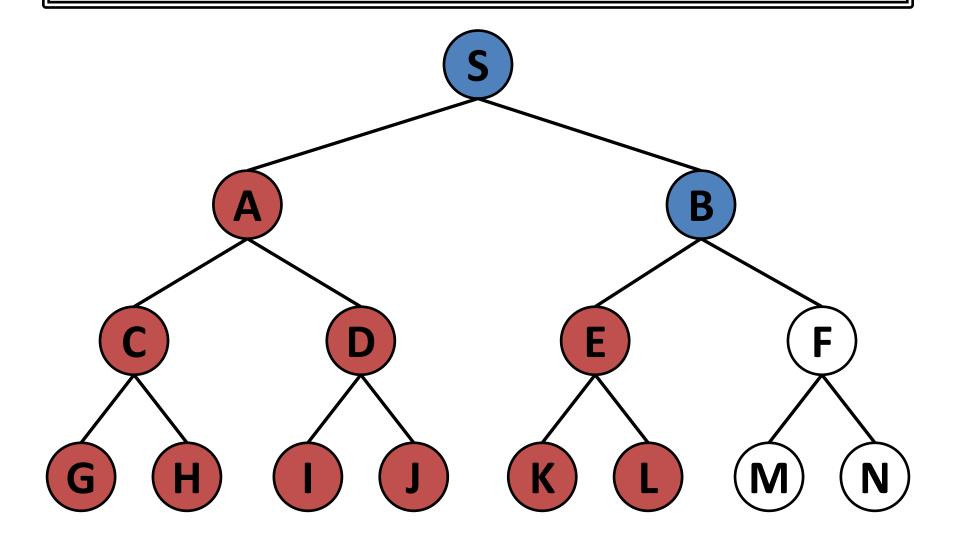


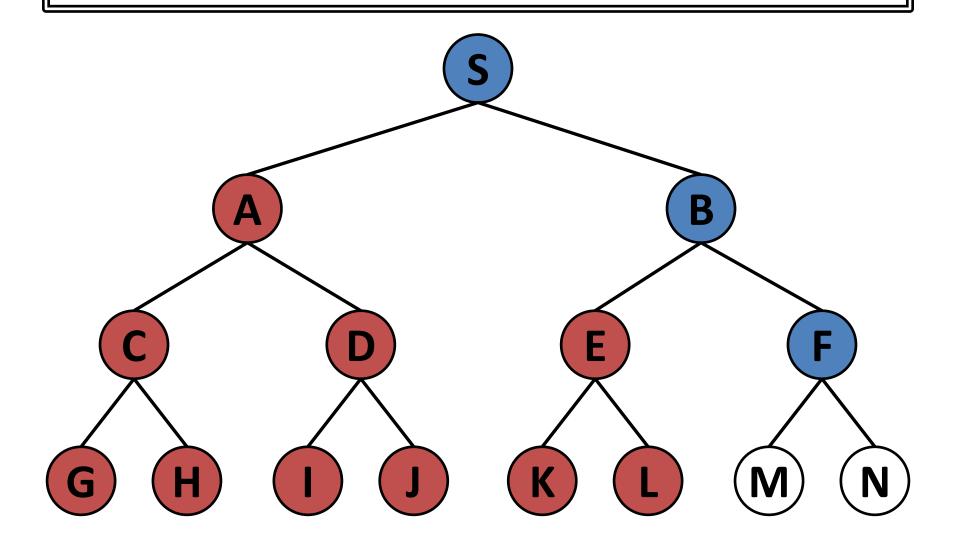


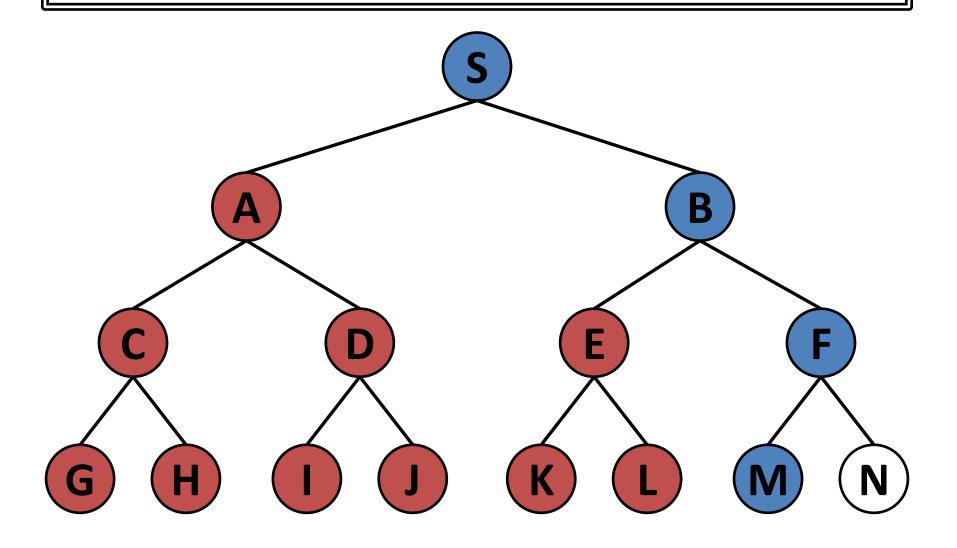


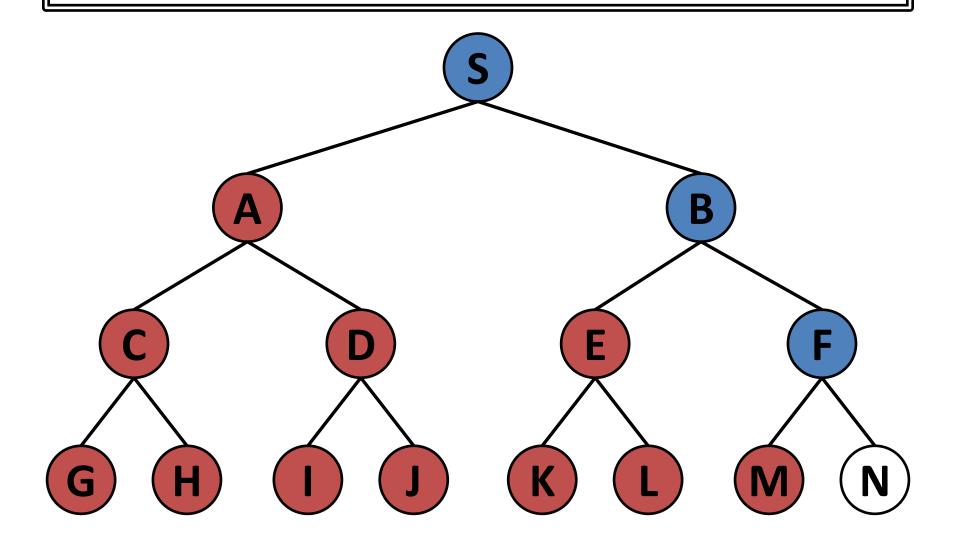


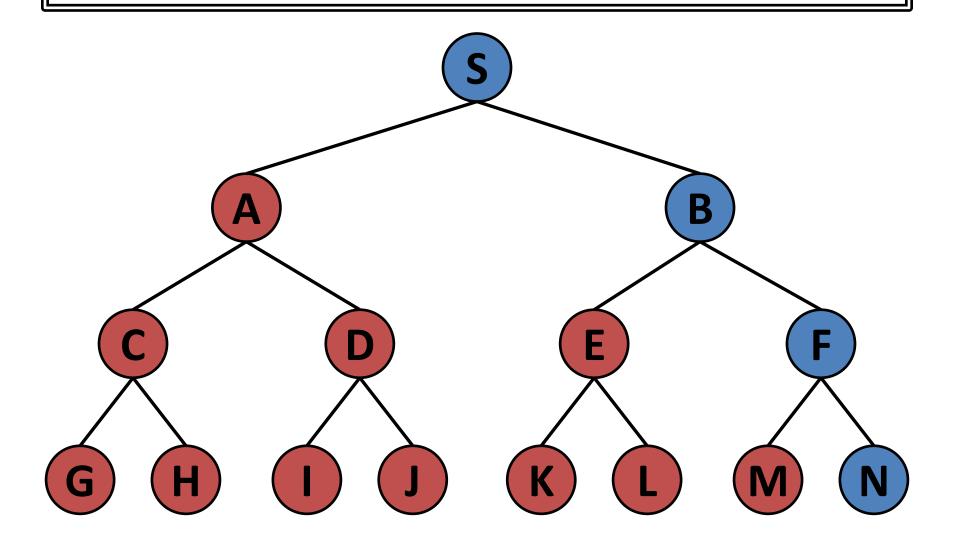


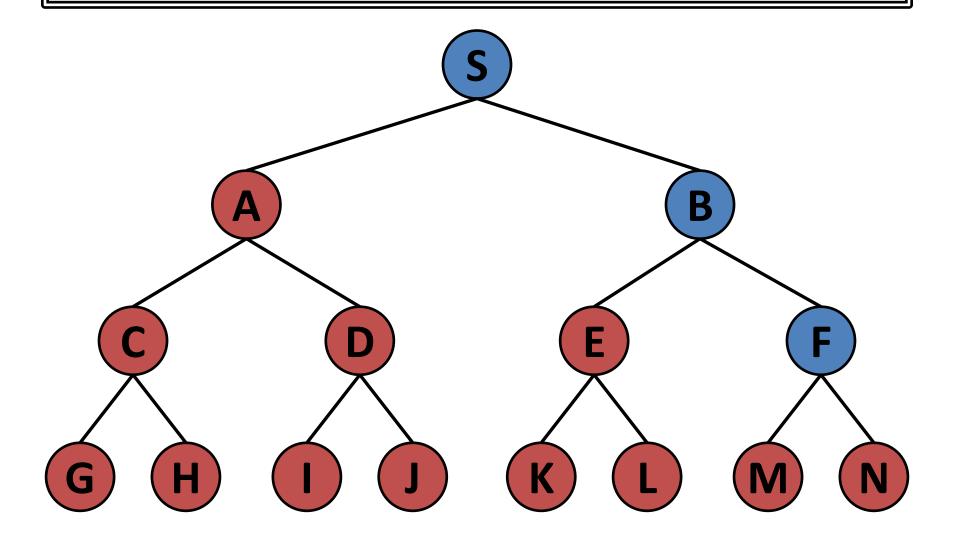


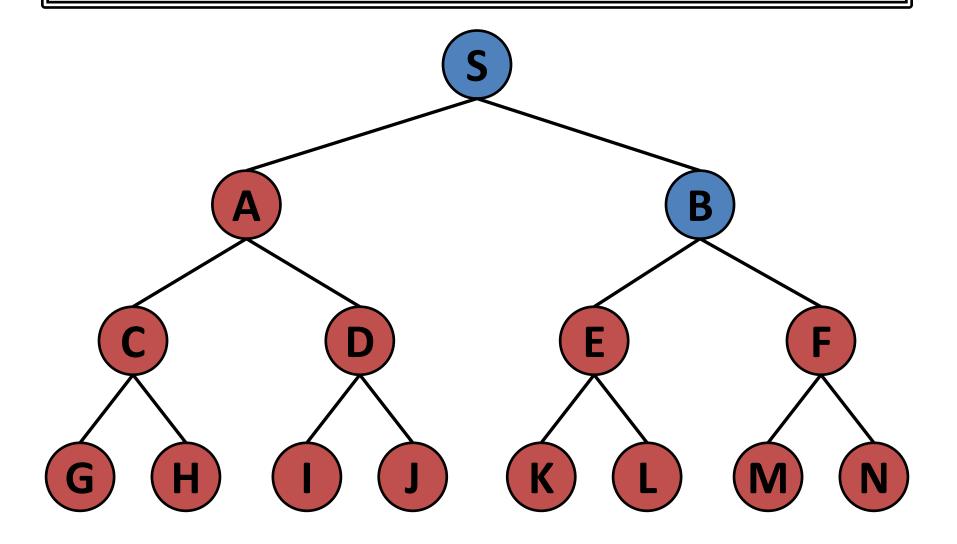


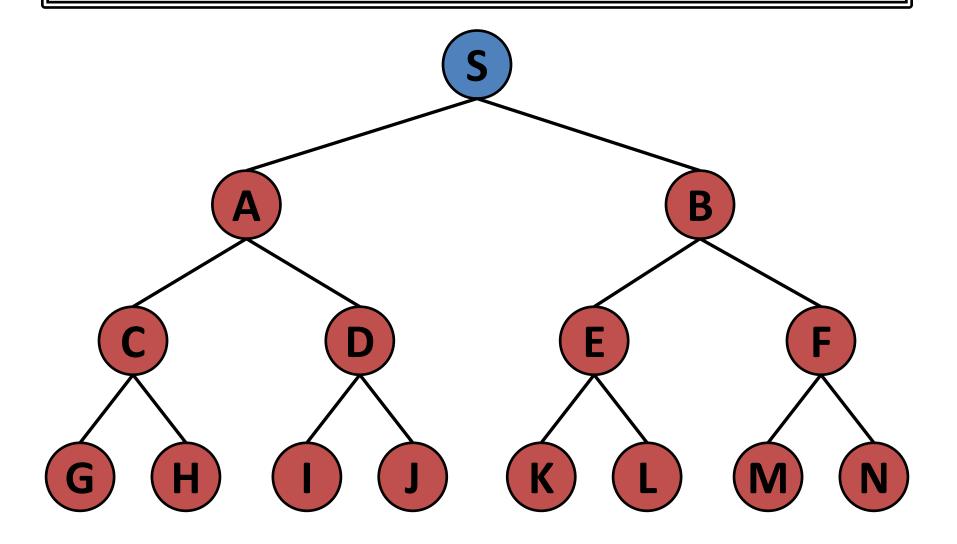


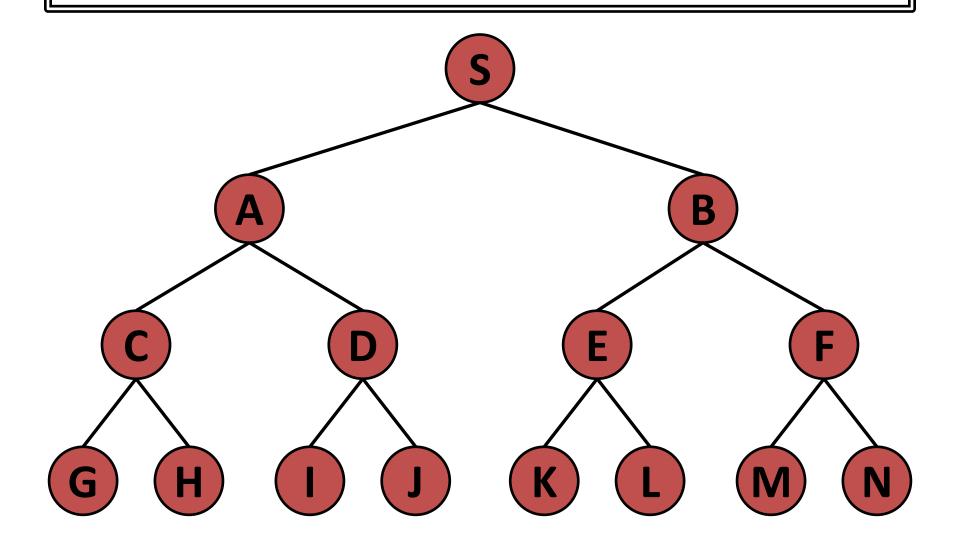


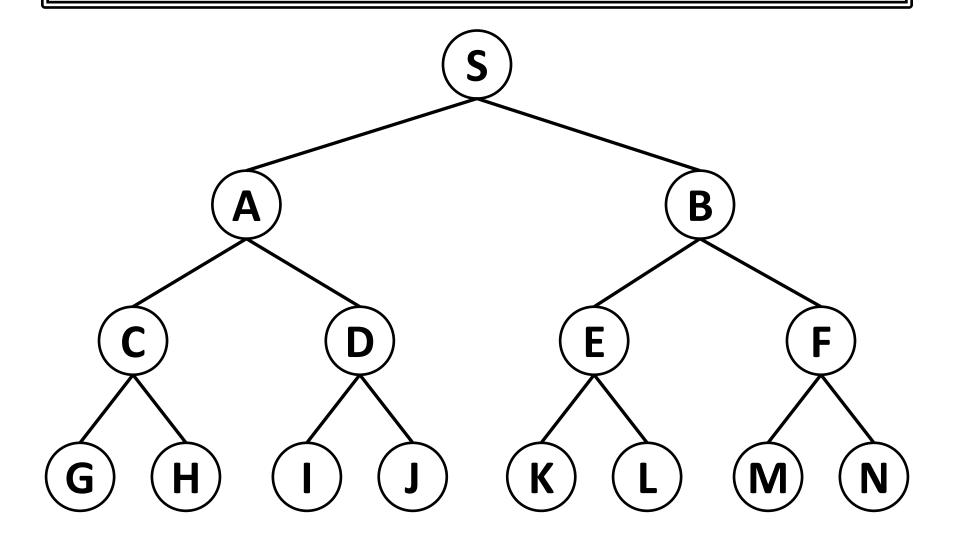


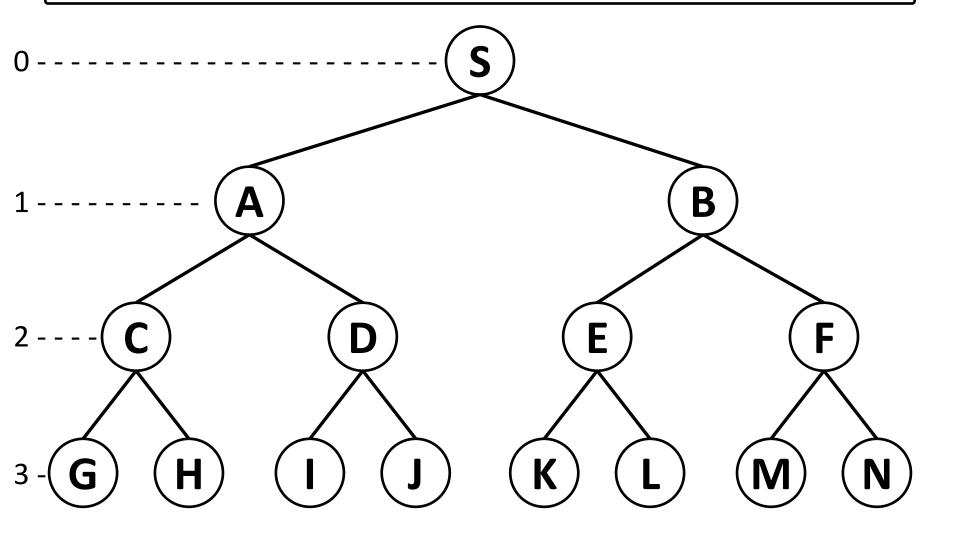


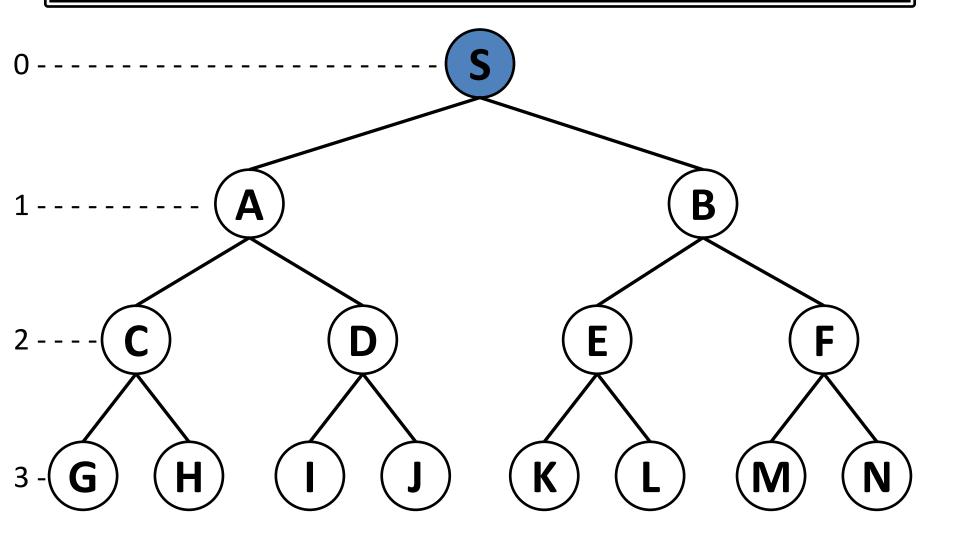


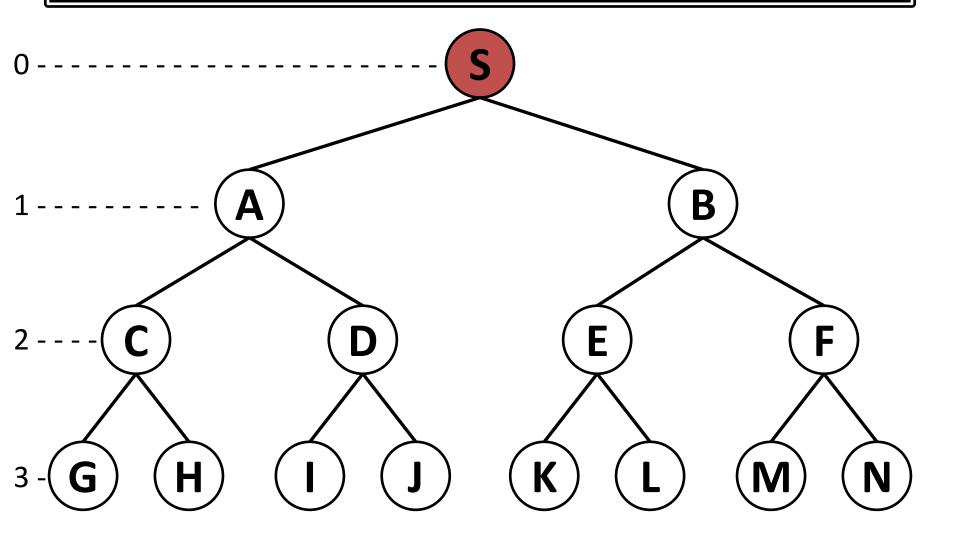


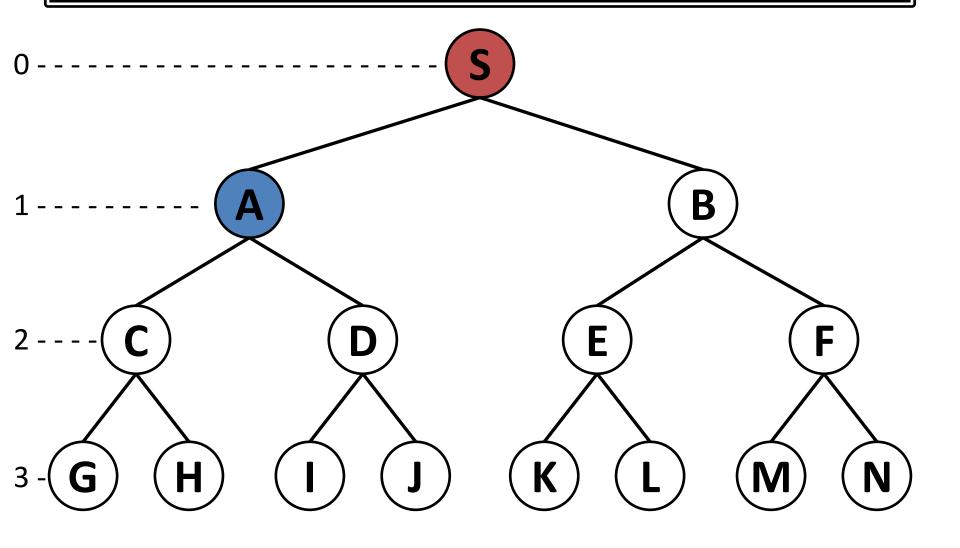


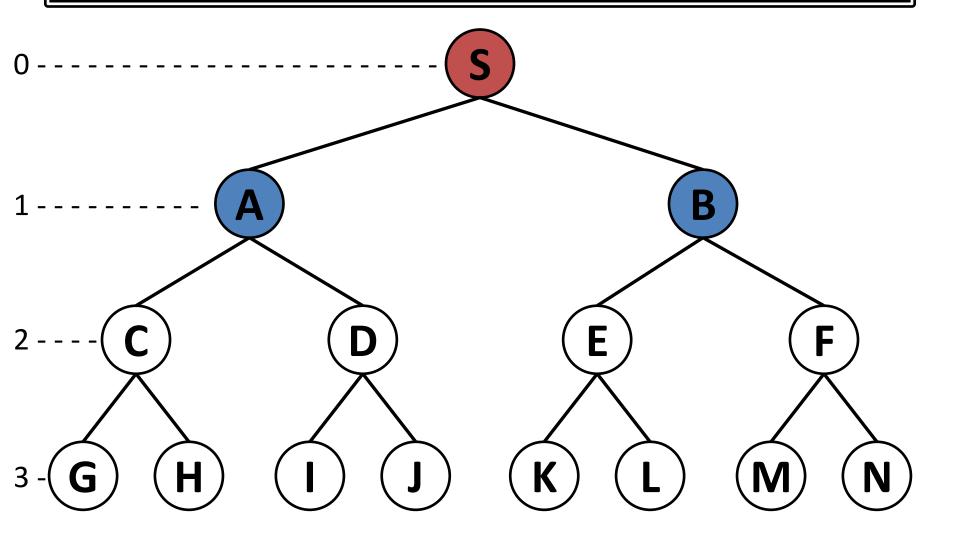


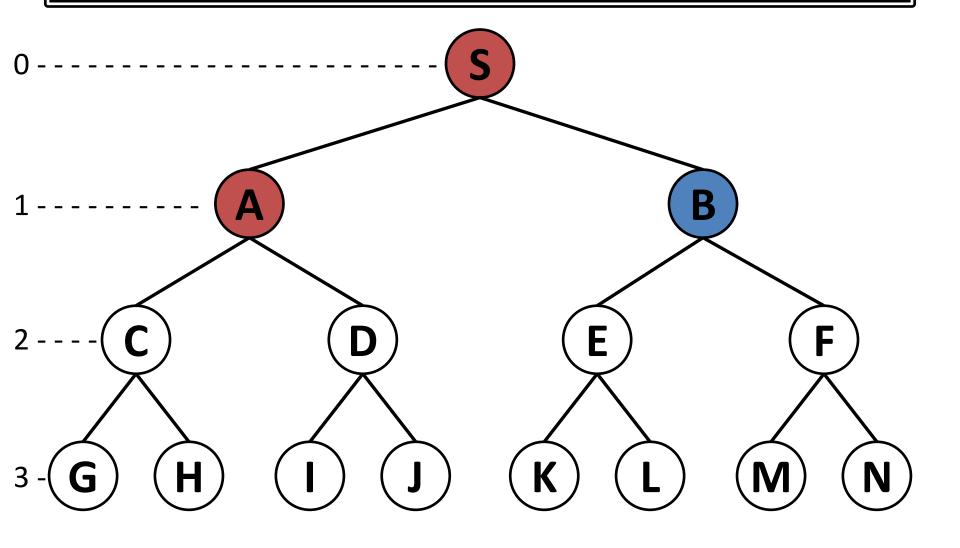


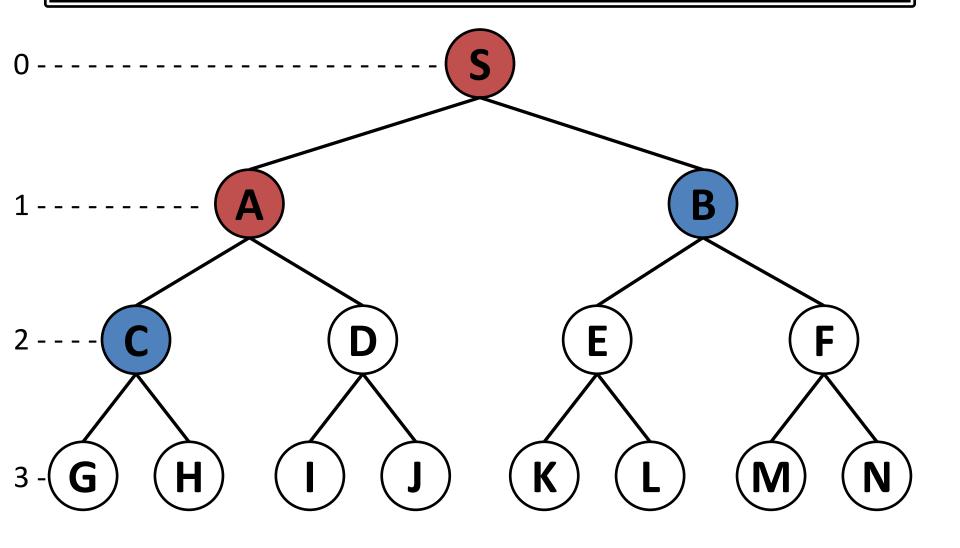


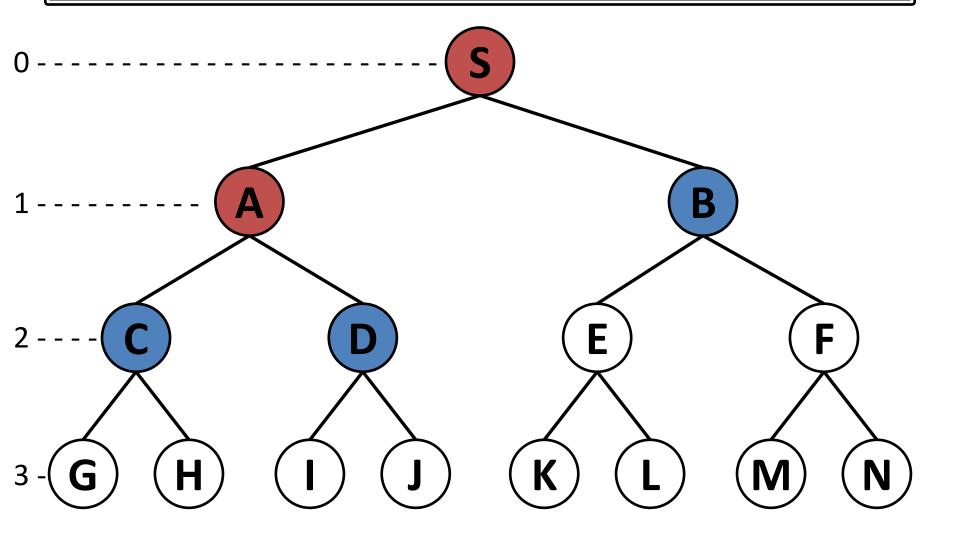


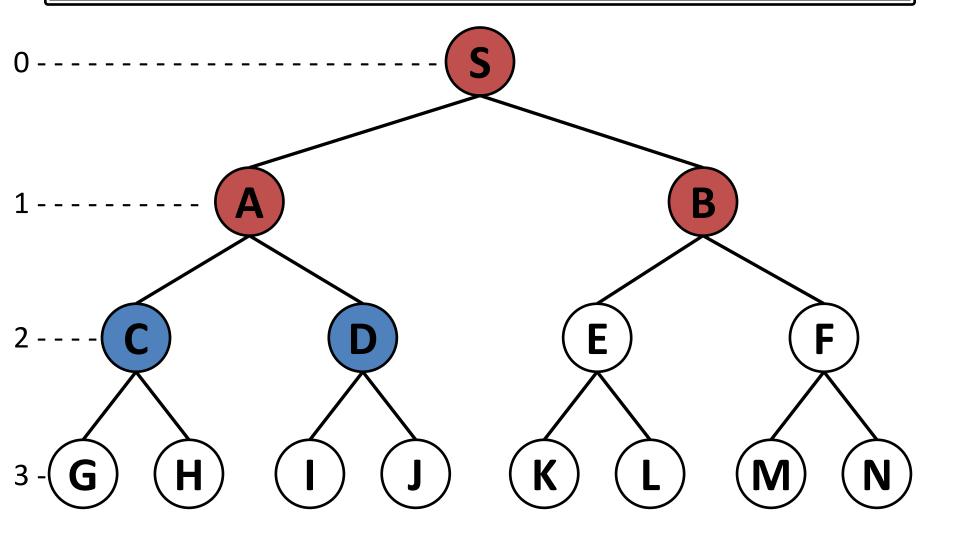


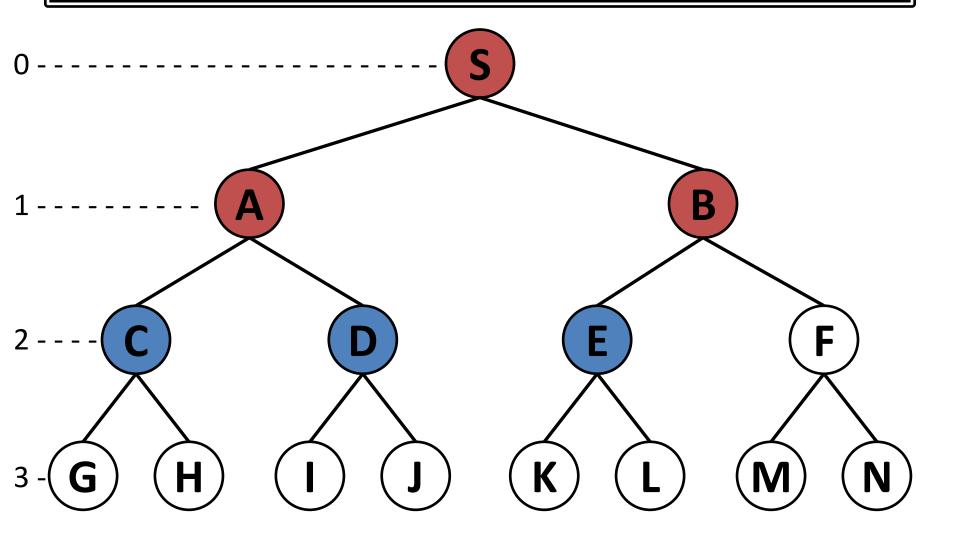


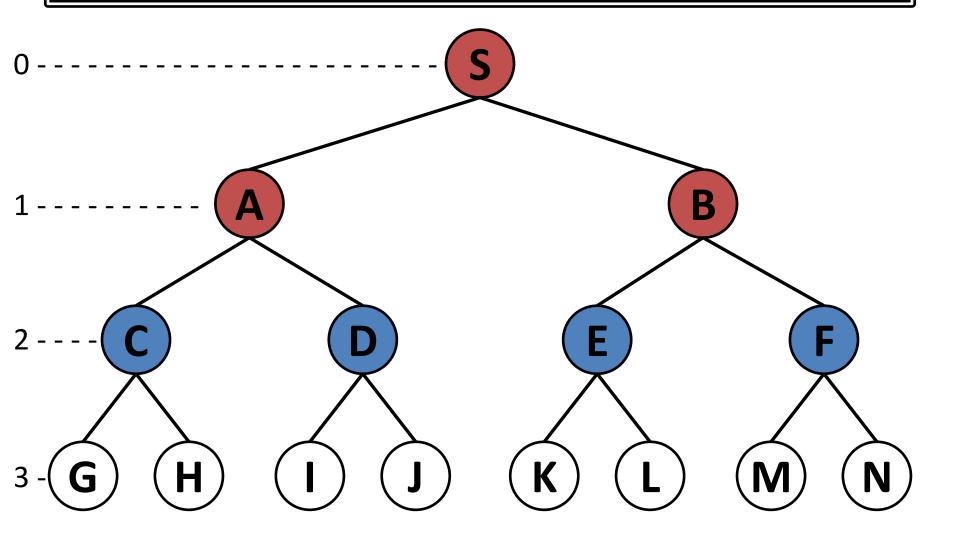


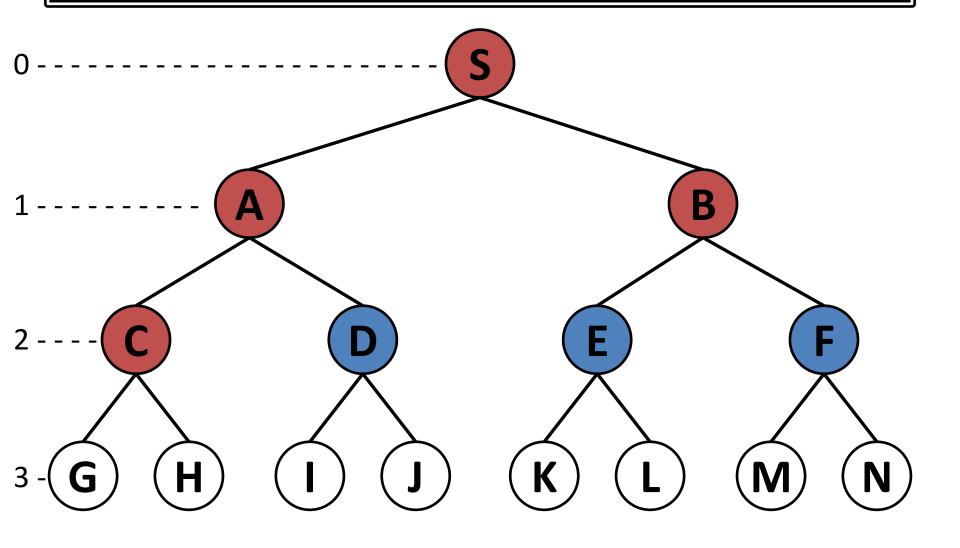


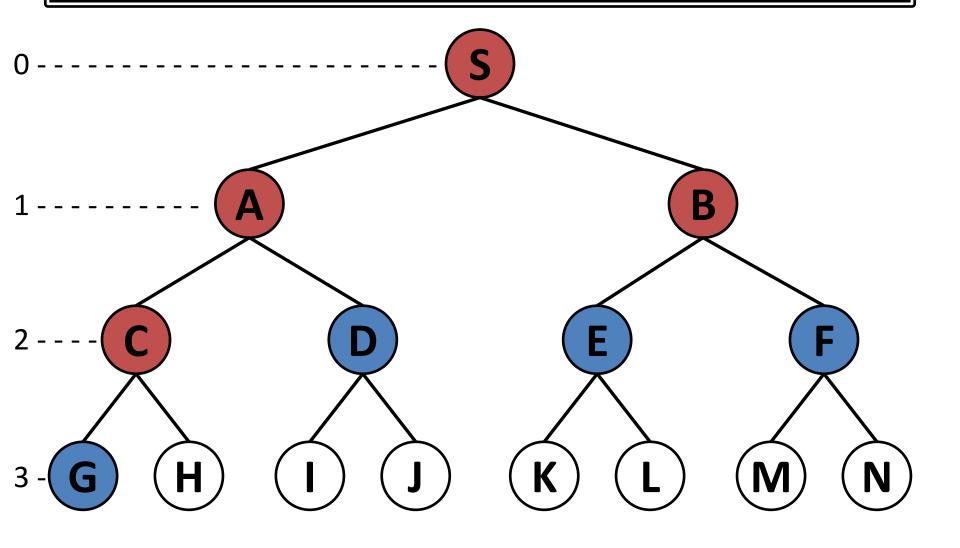


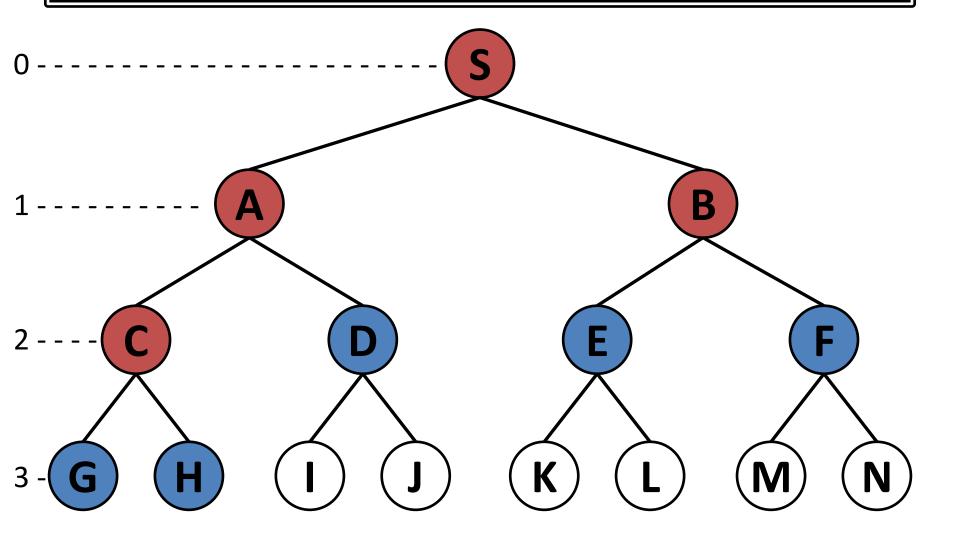


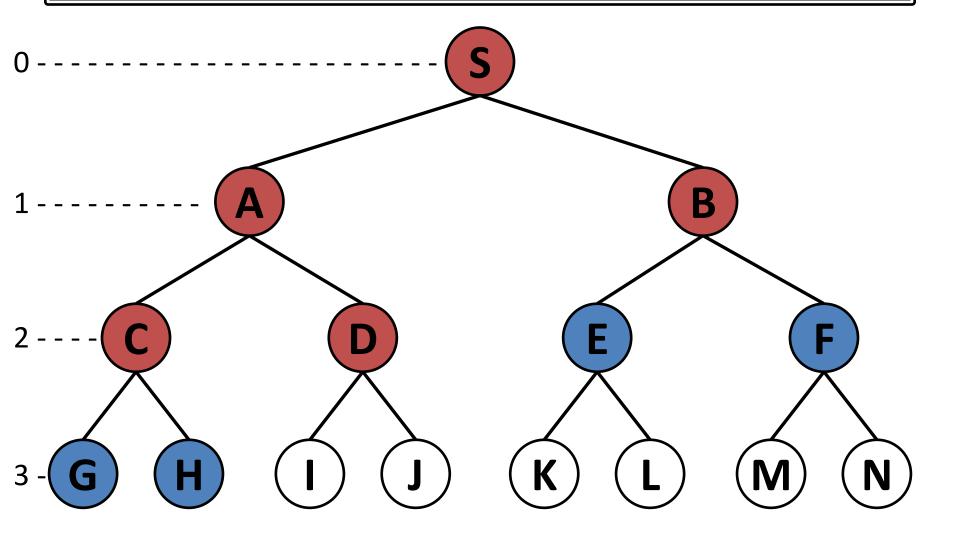


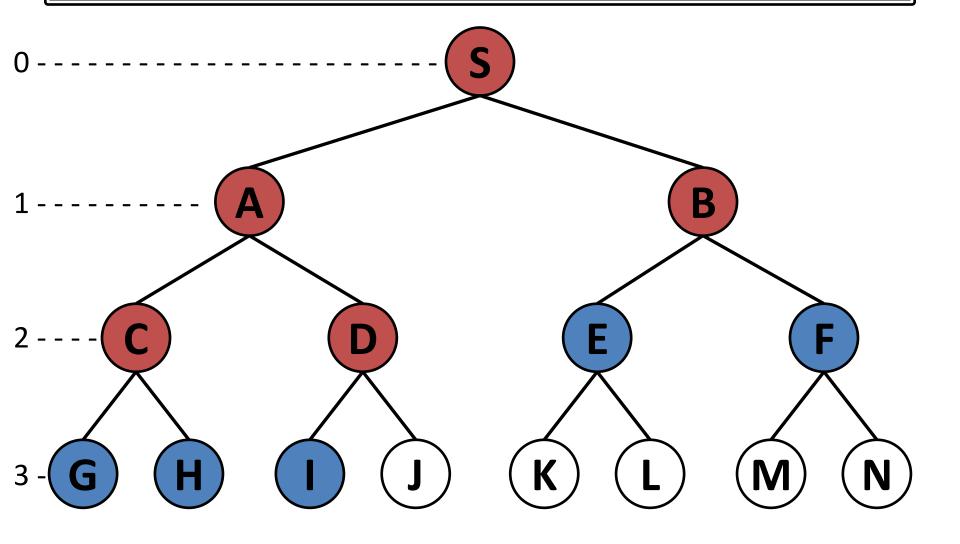


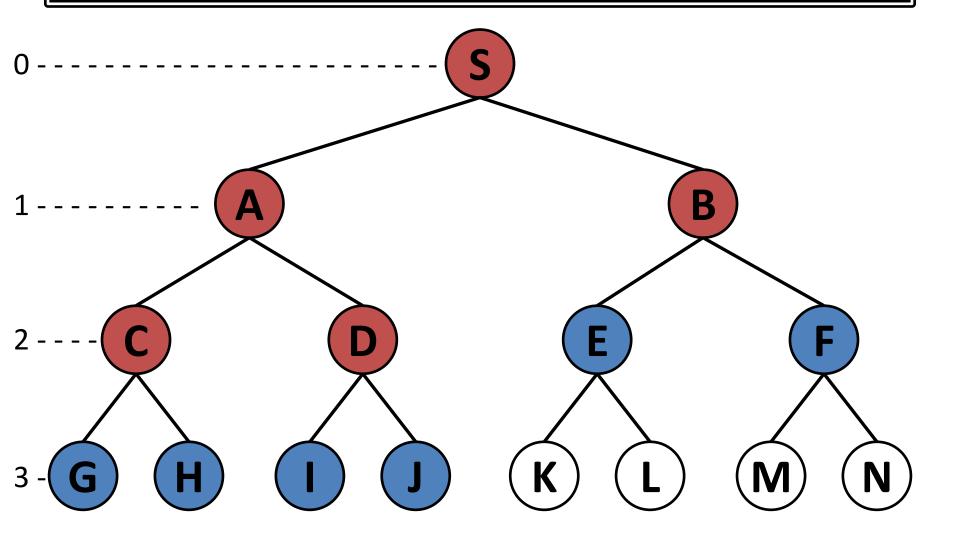


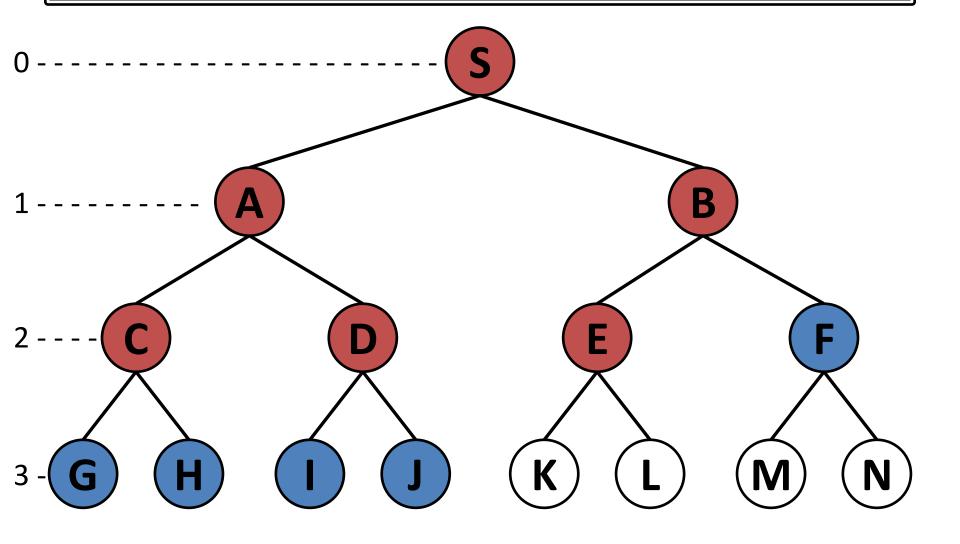


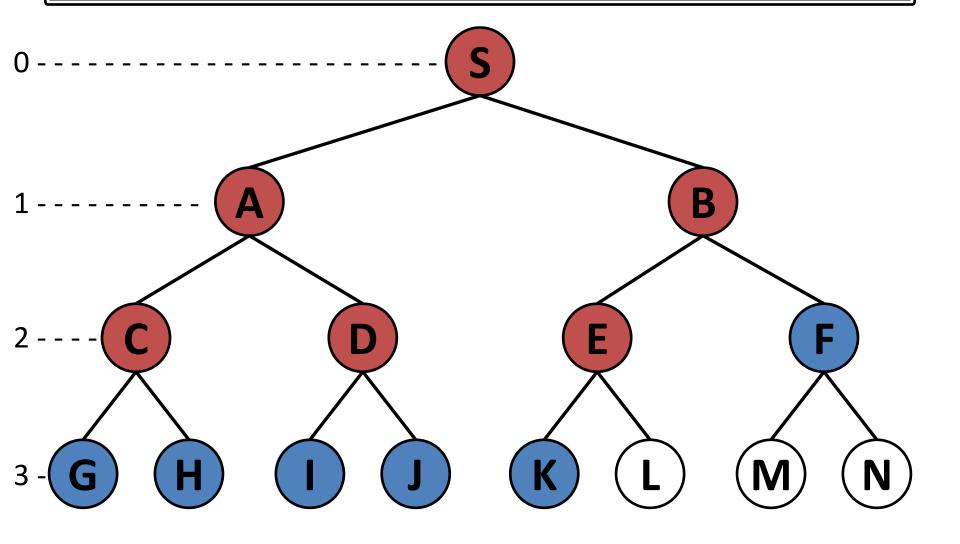


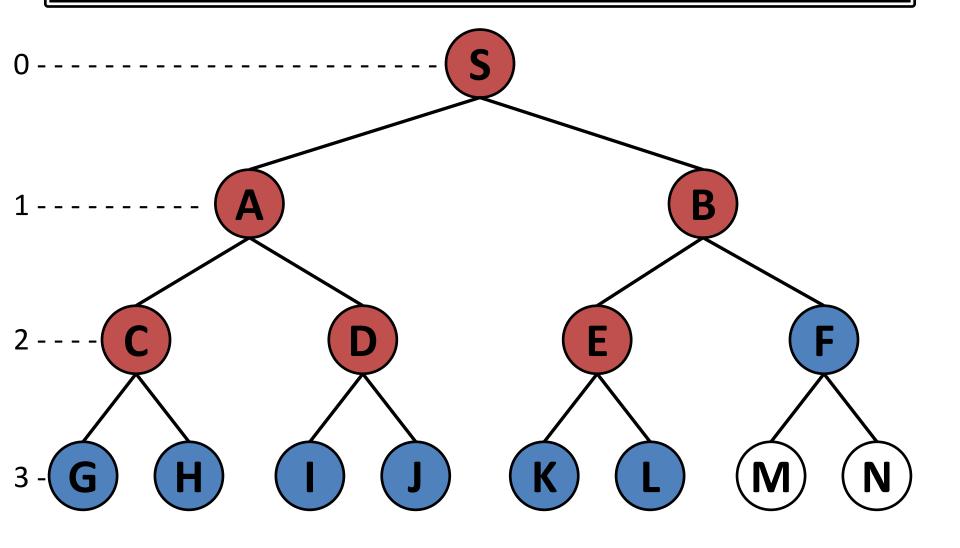


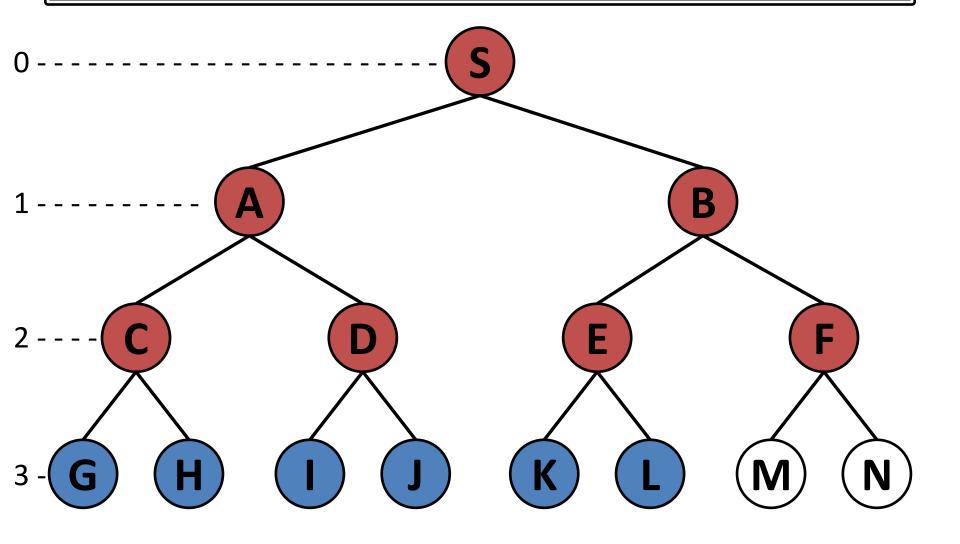


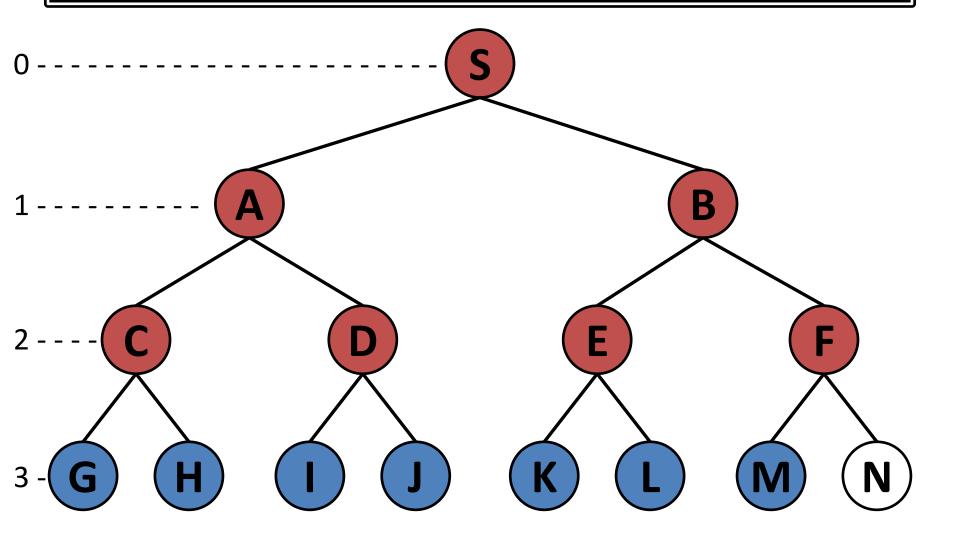


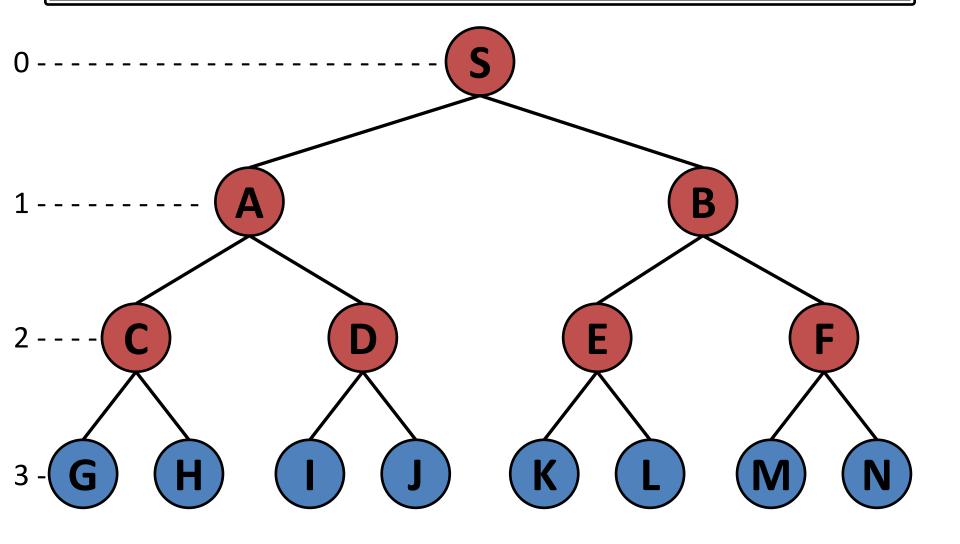


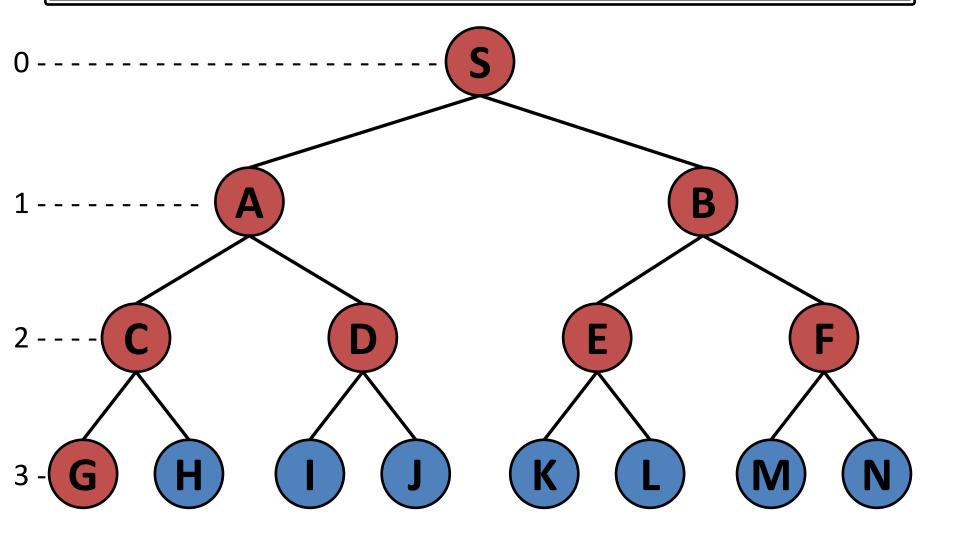


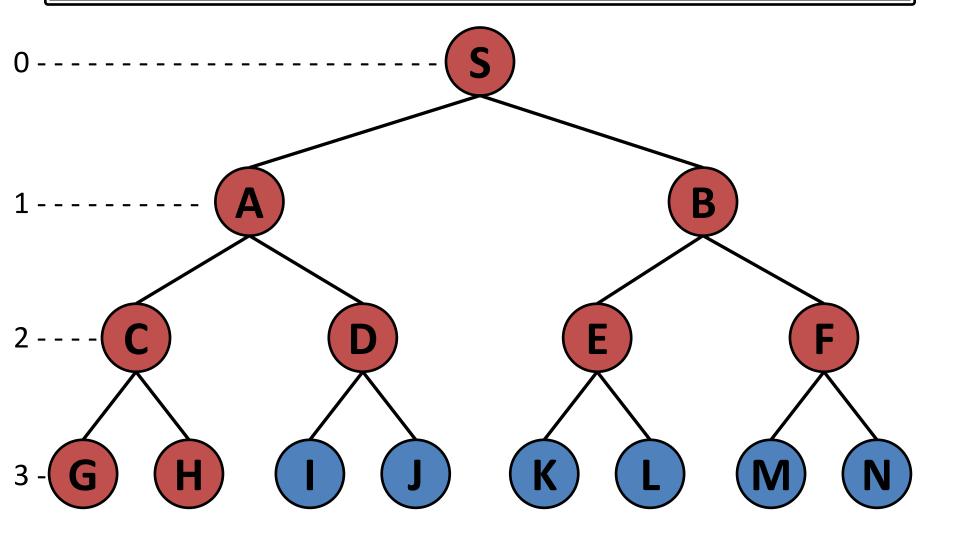


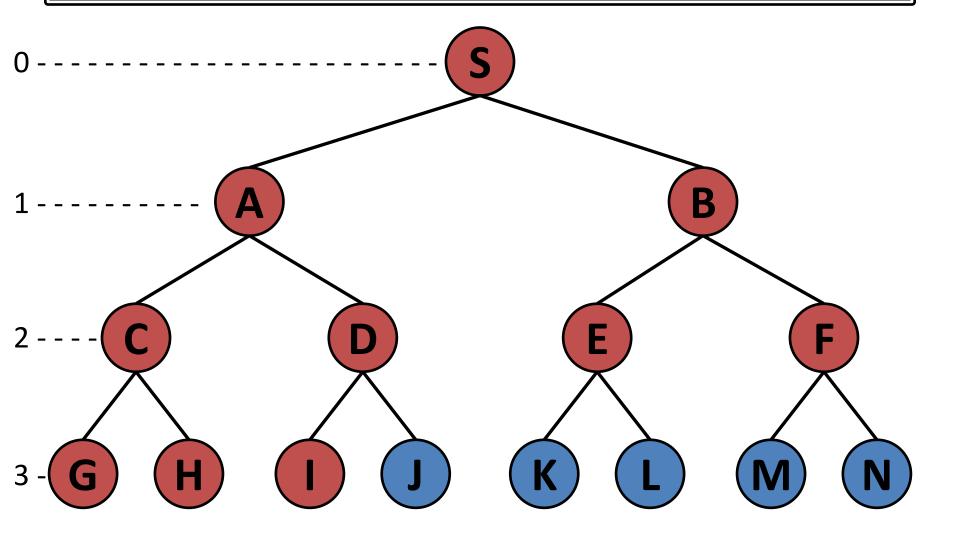


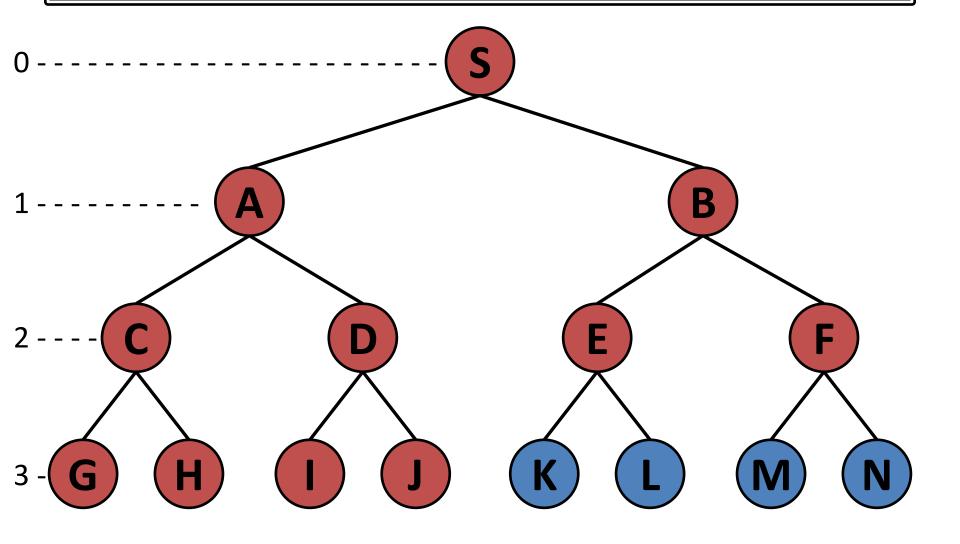


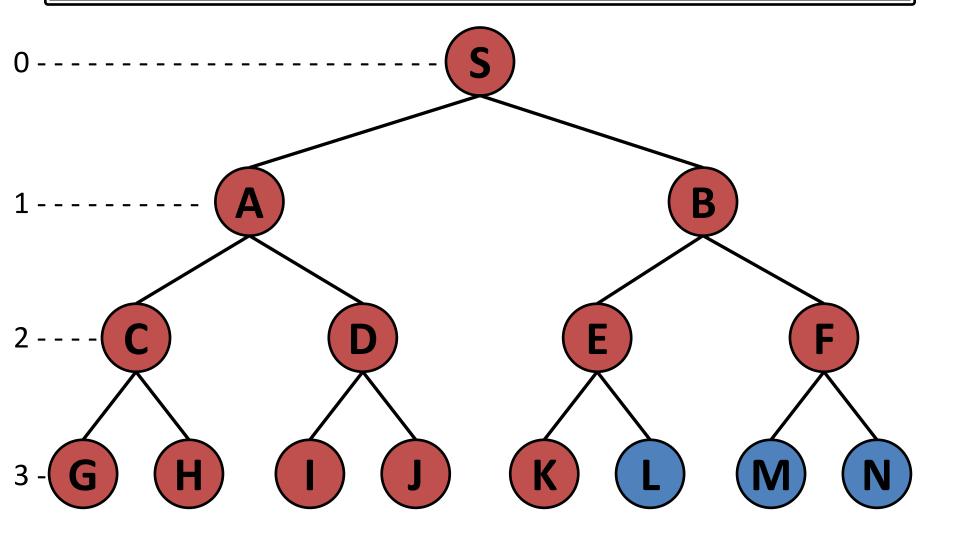


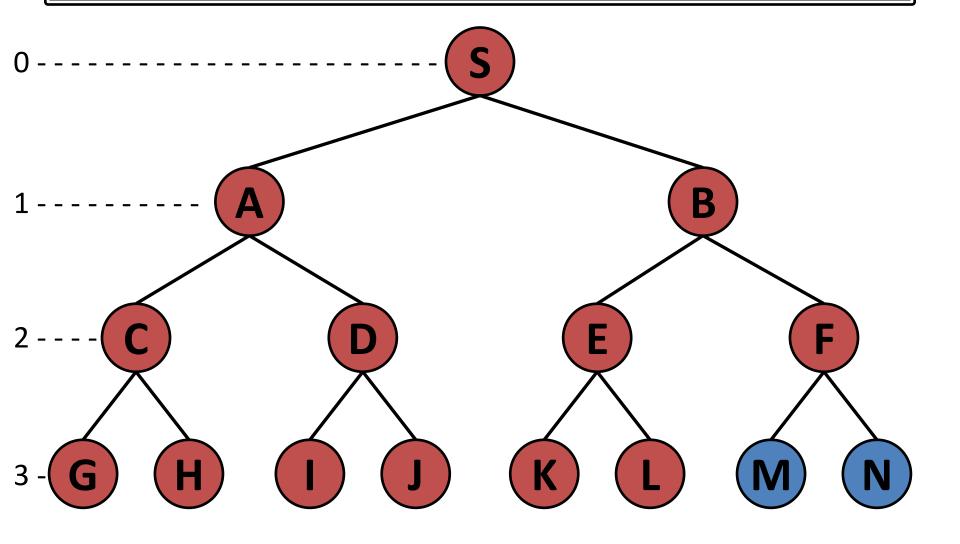




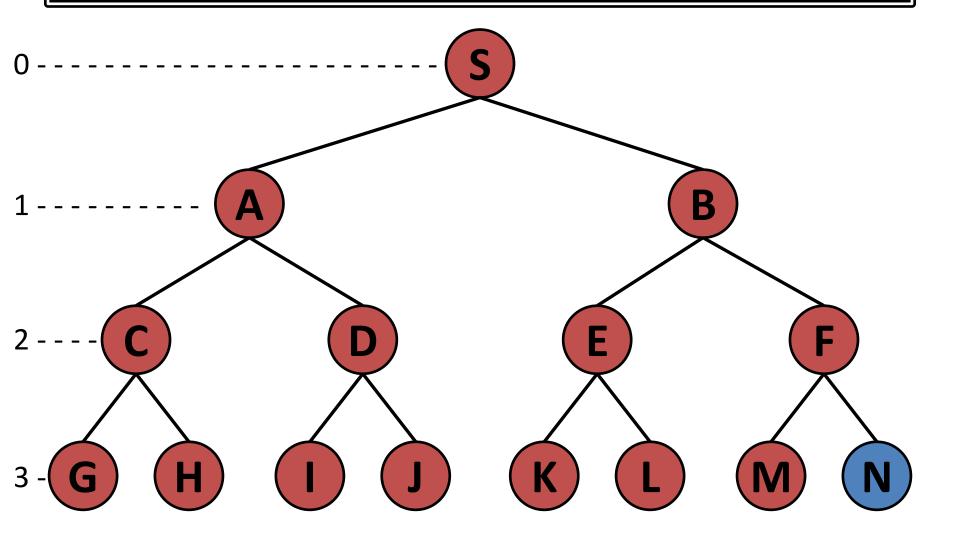




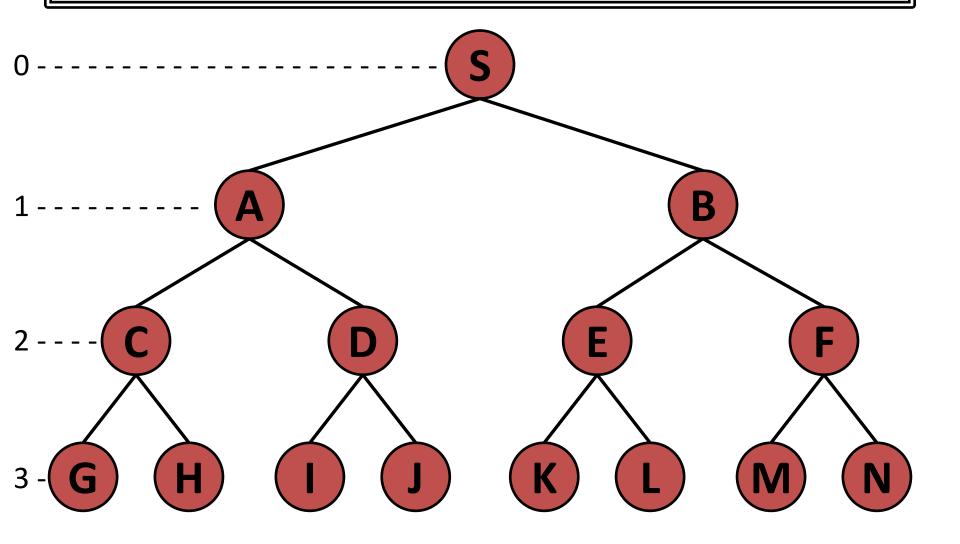




# Graph and Tree Traversal: BFS

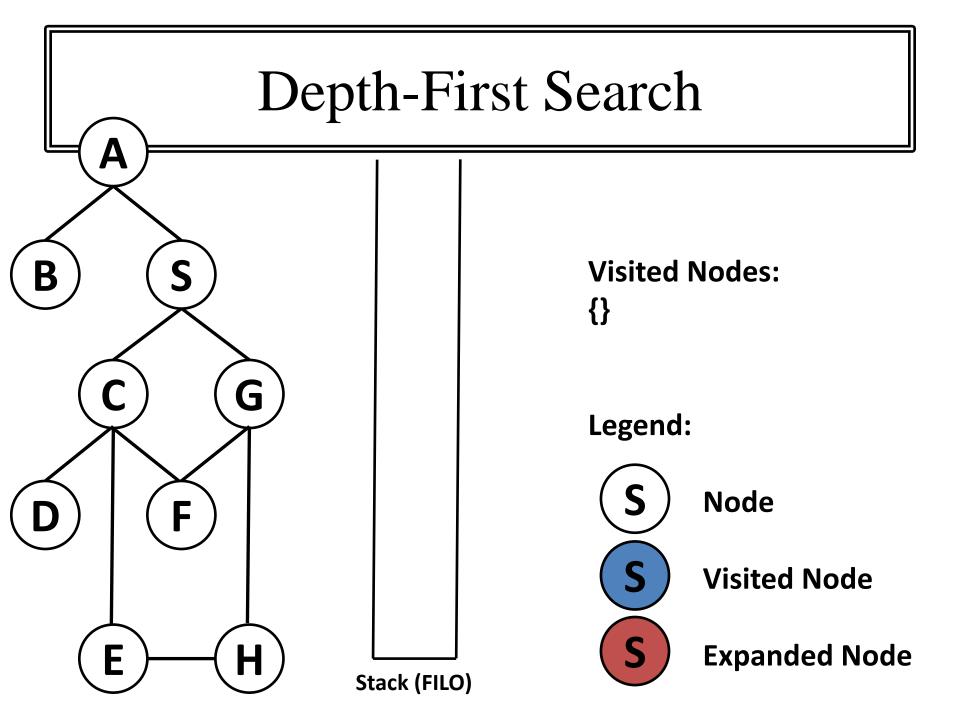


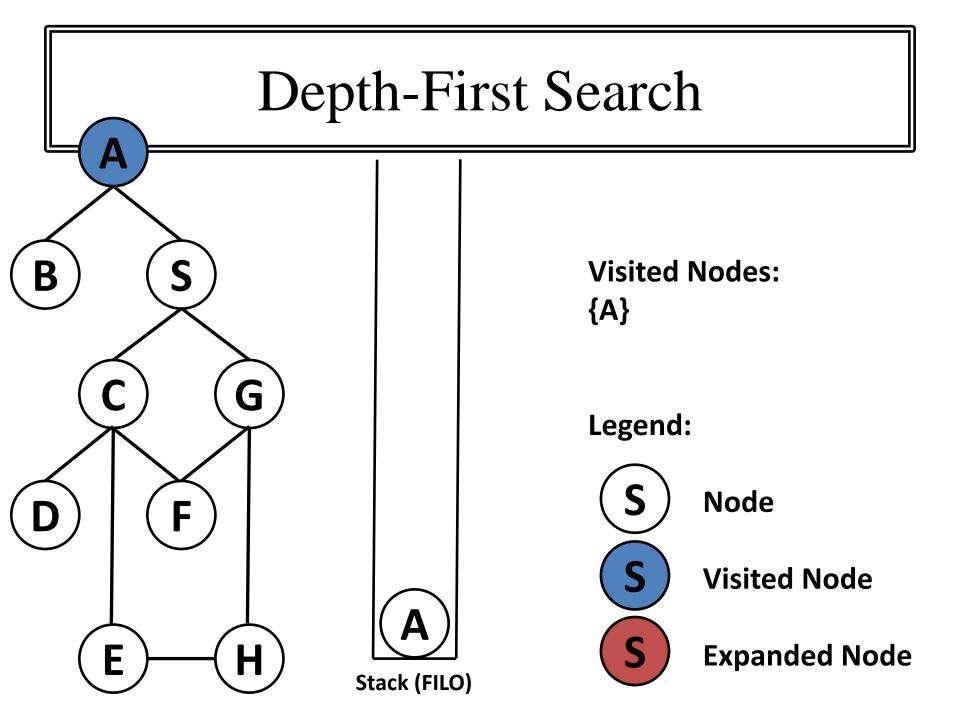
# Graph and Tree Traversal: BFS

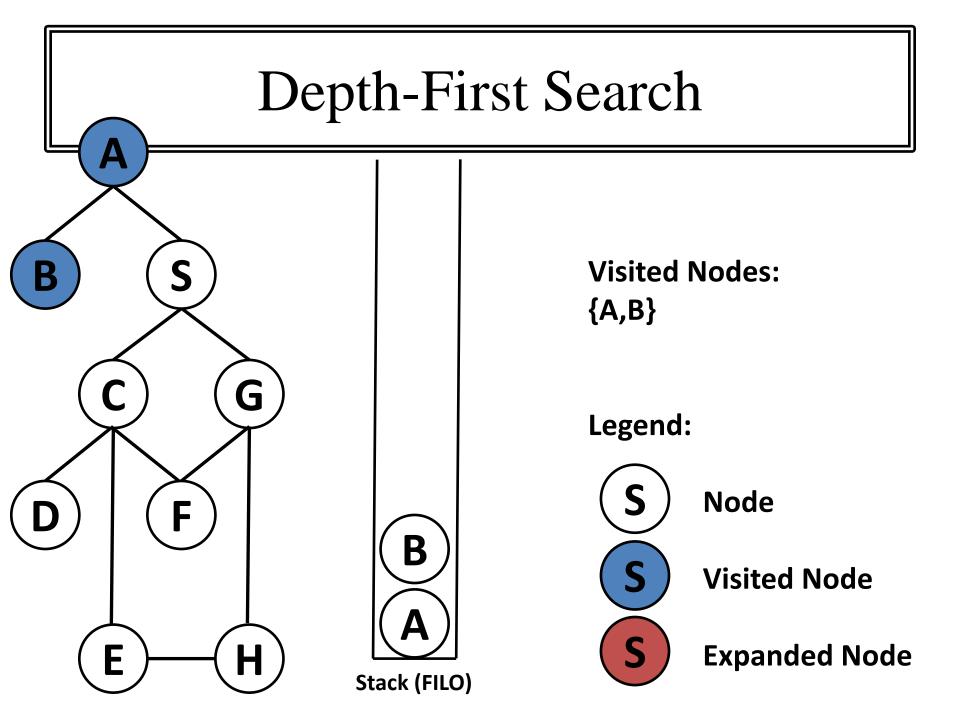


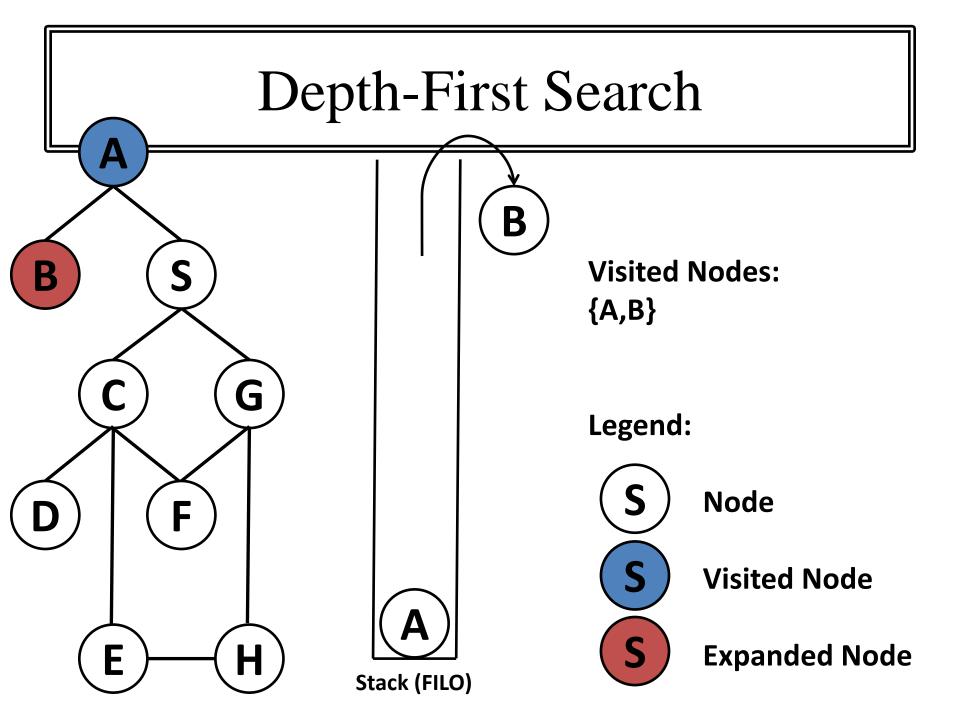
### Uninformed (Blind) Search

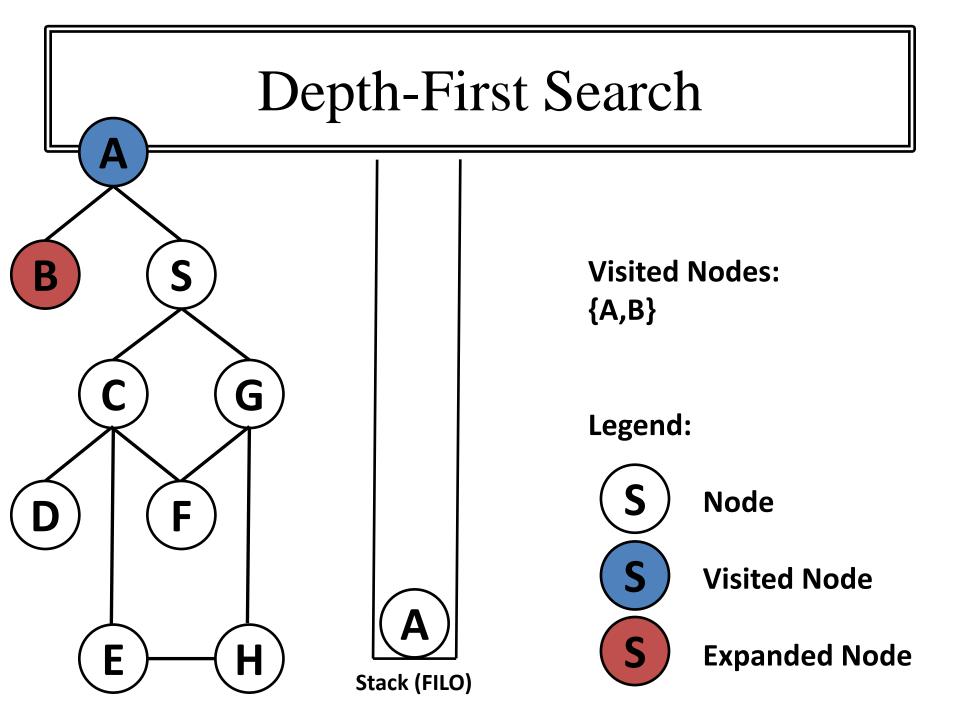
- "Uninformed strategies use only the information available in the problem definition."
- Uninformed Search Algorithms:
  - Depth-First Search (DFS)
  - Breadth-First Search (BFS)
  - Uniform-Cost Search (UCS)
  - Depth-Limited Search (DLS)
  - Iterative Deepening Search (IDS)

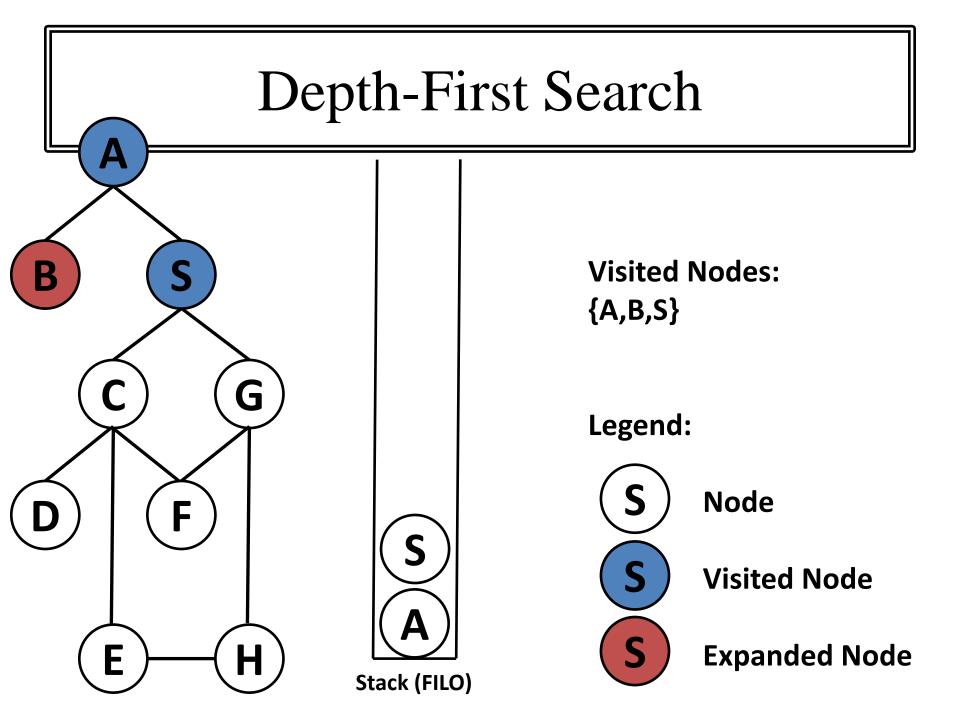


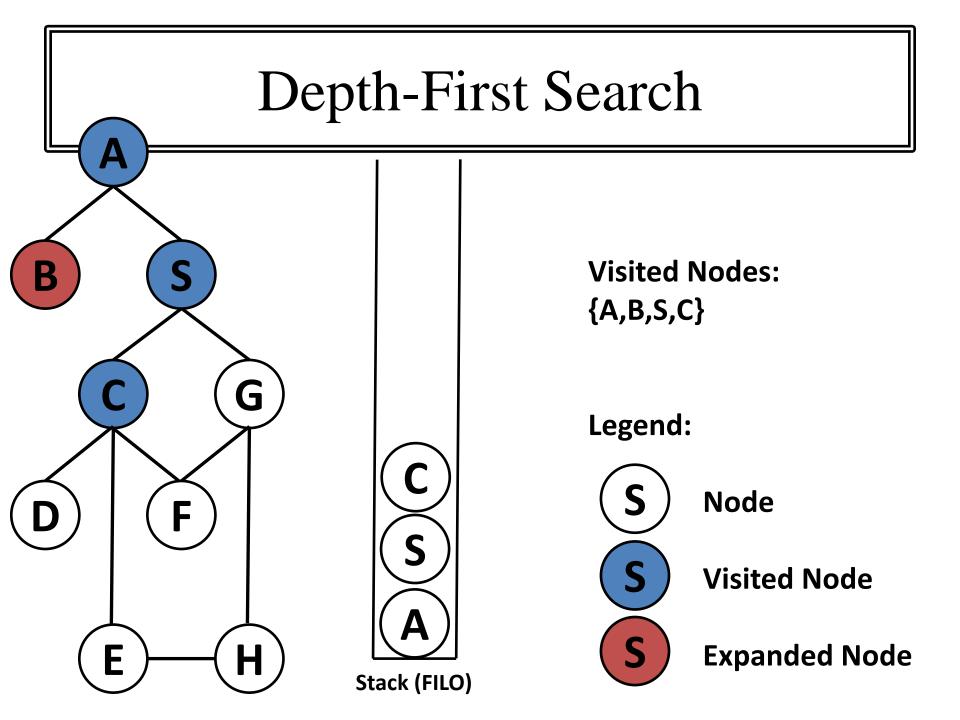


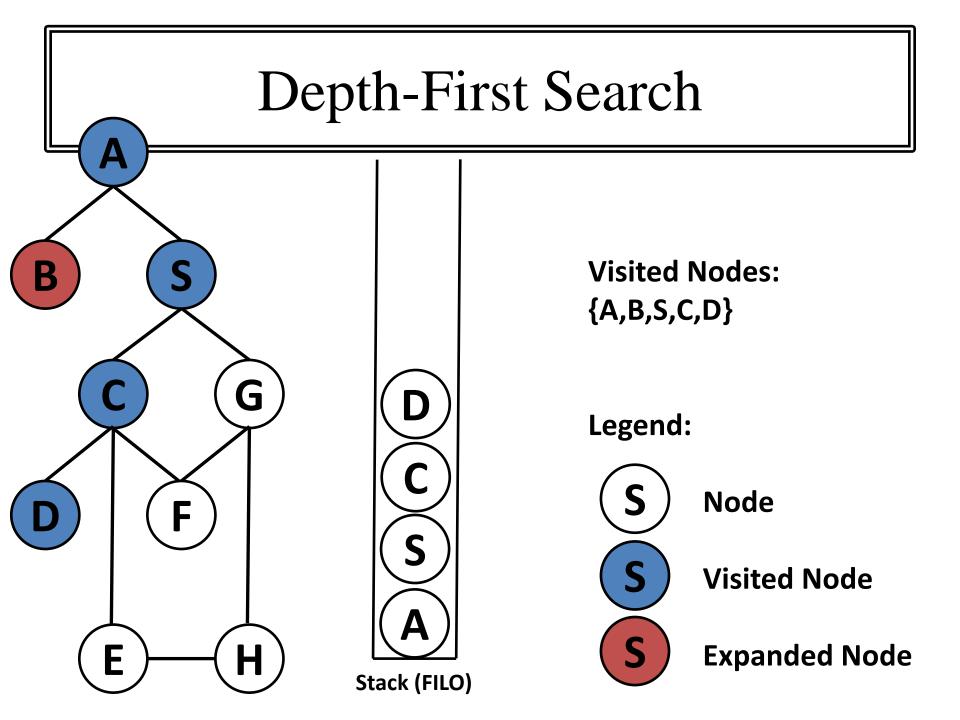


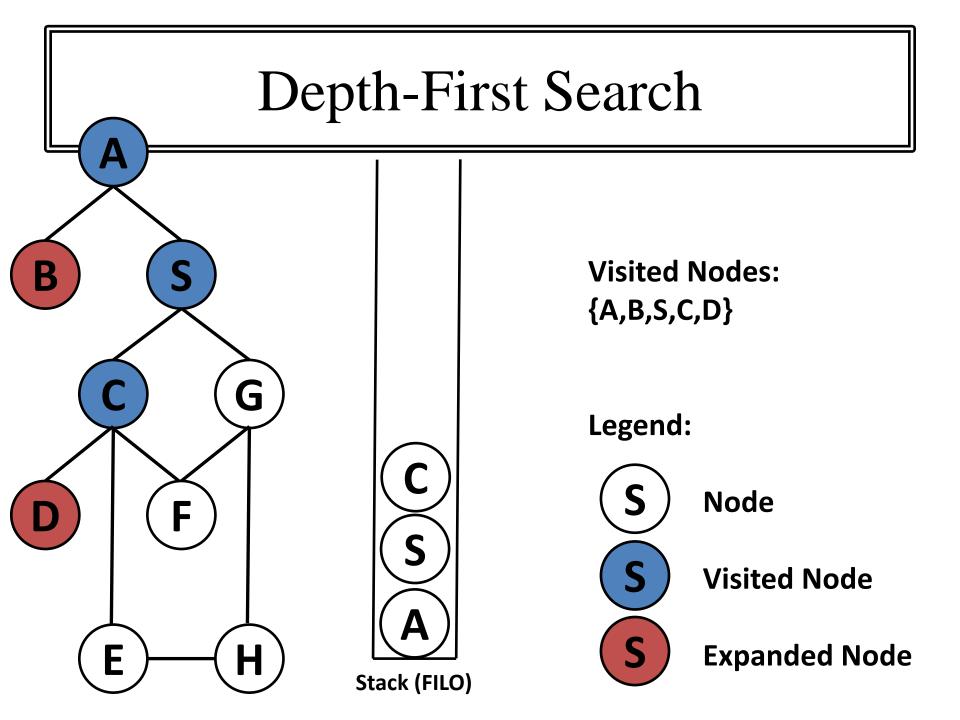


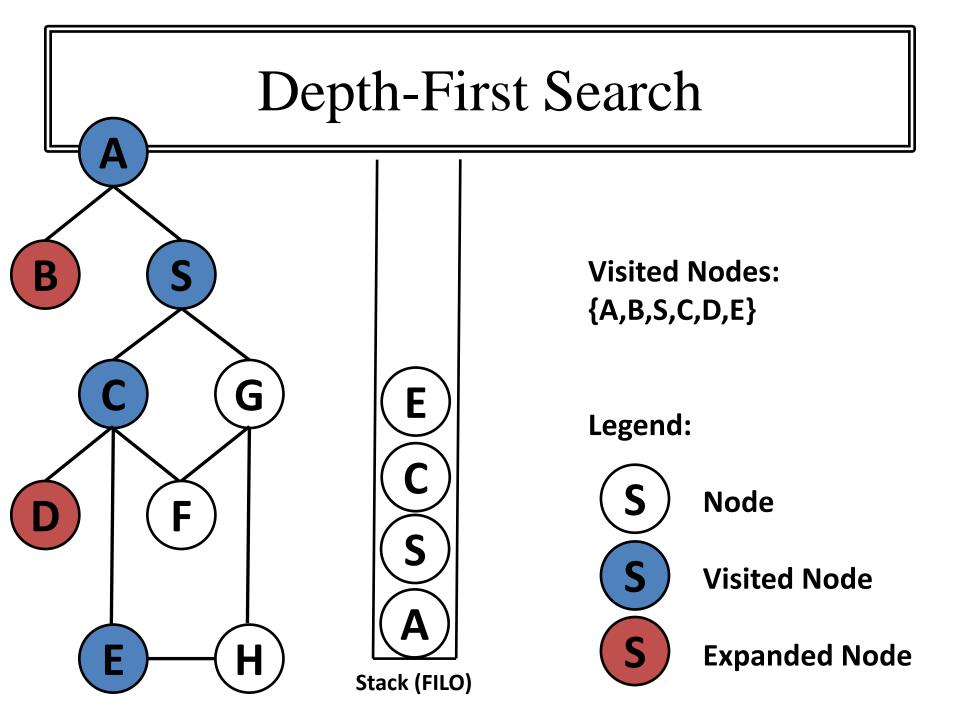


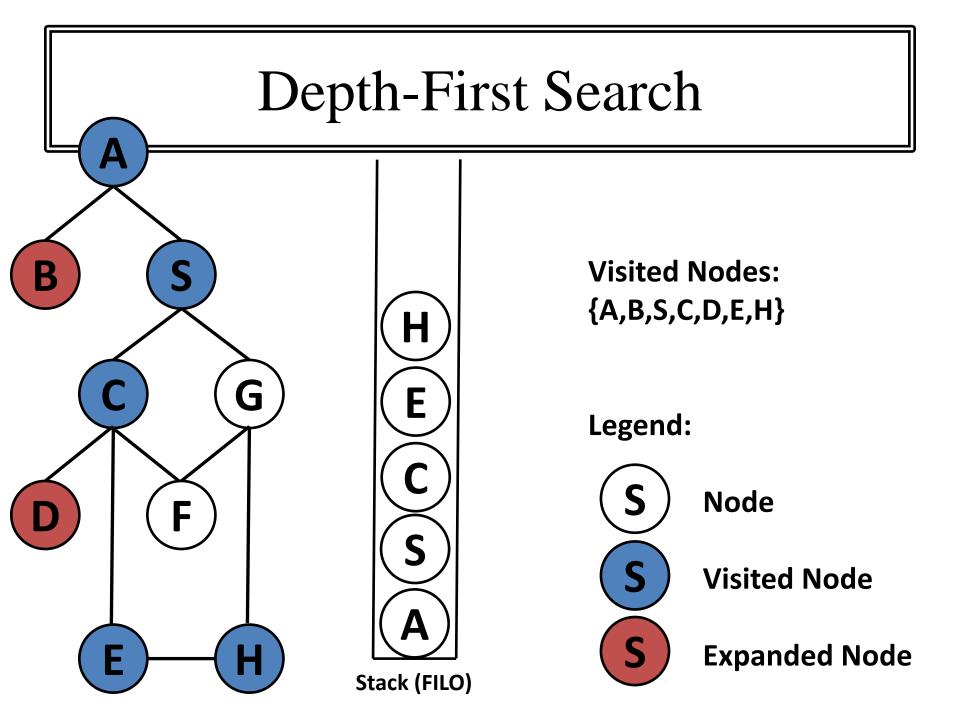


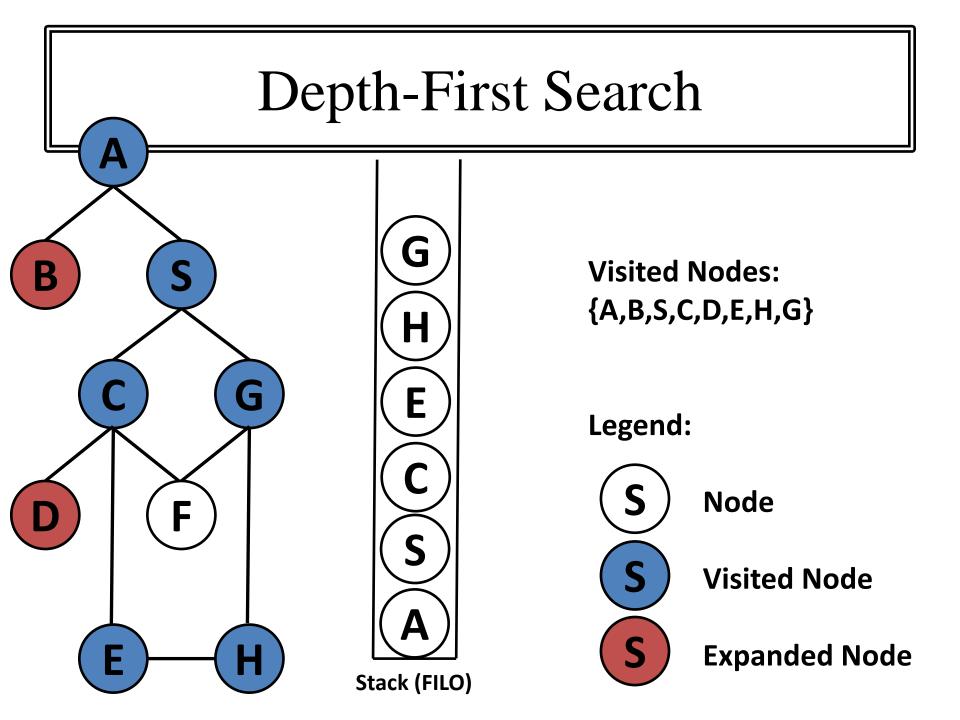


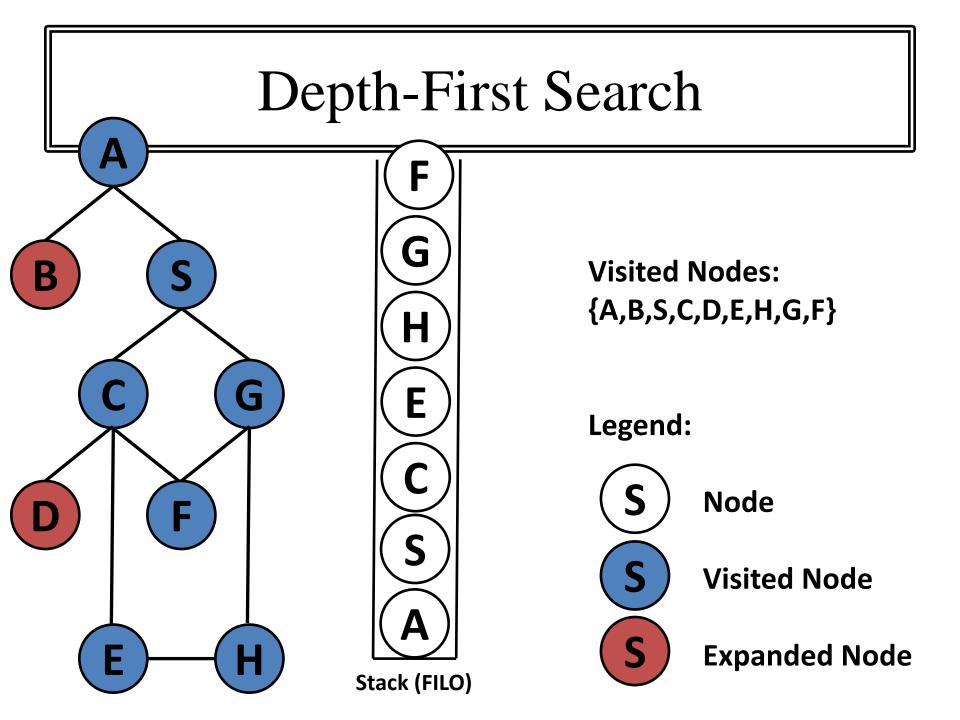


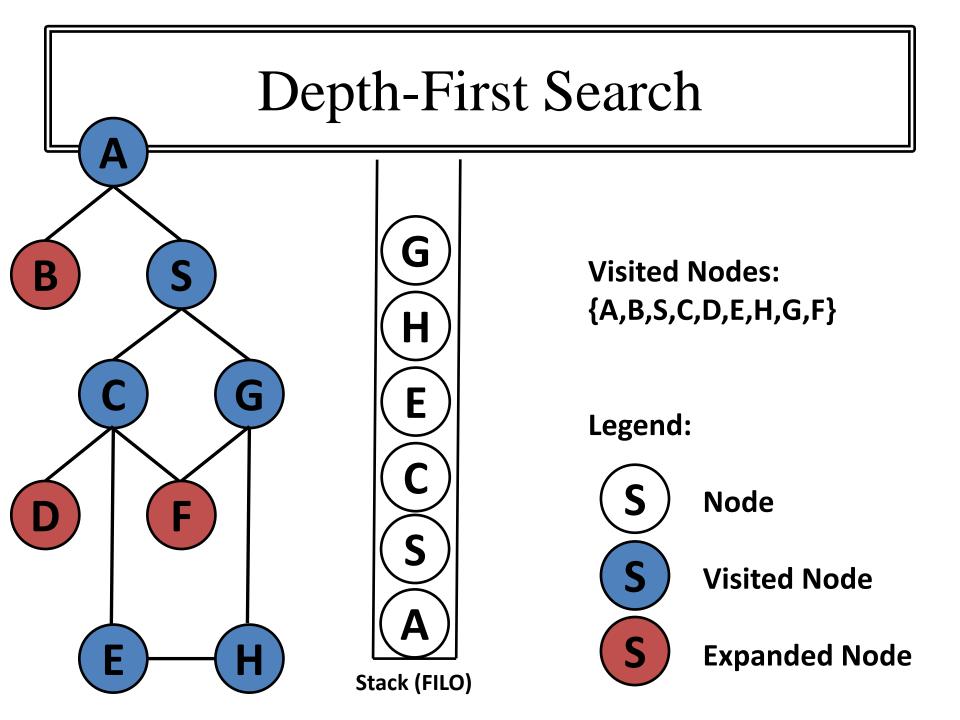


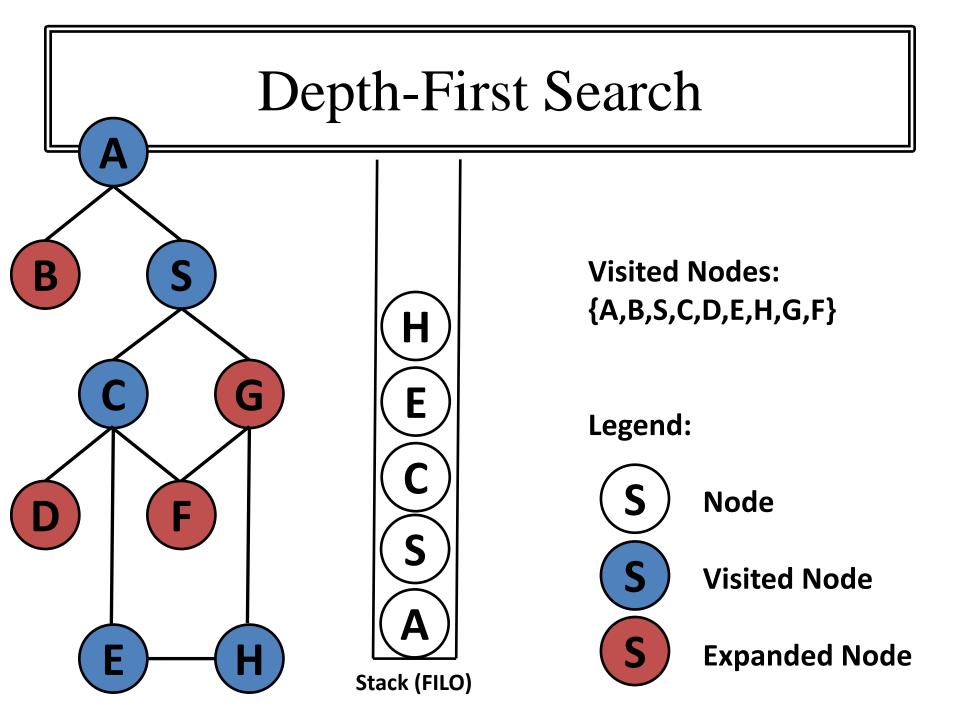


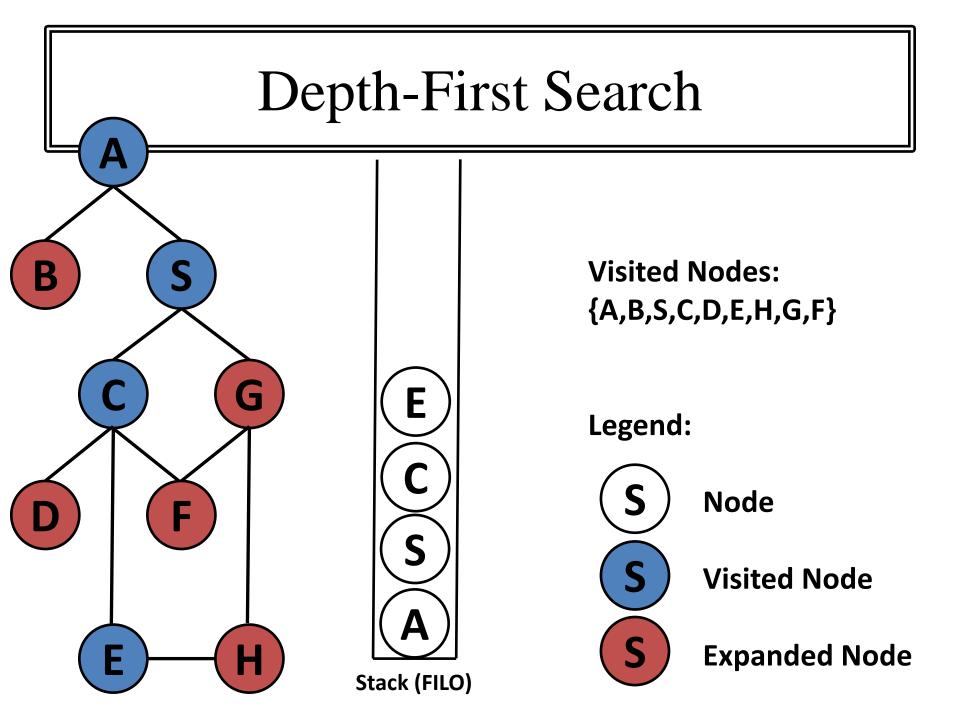


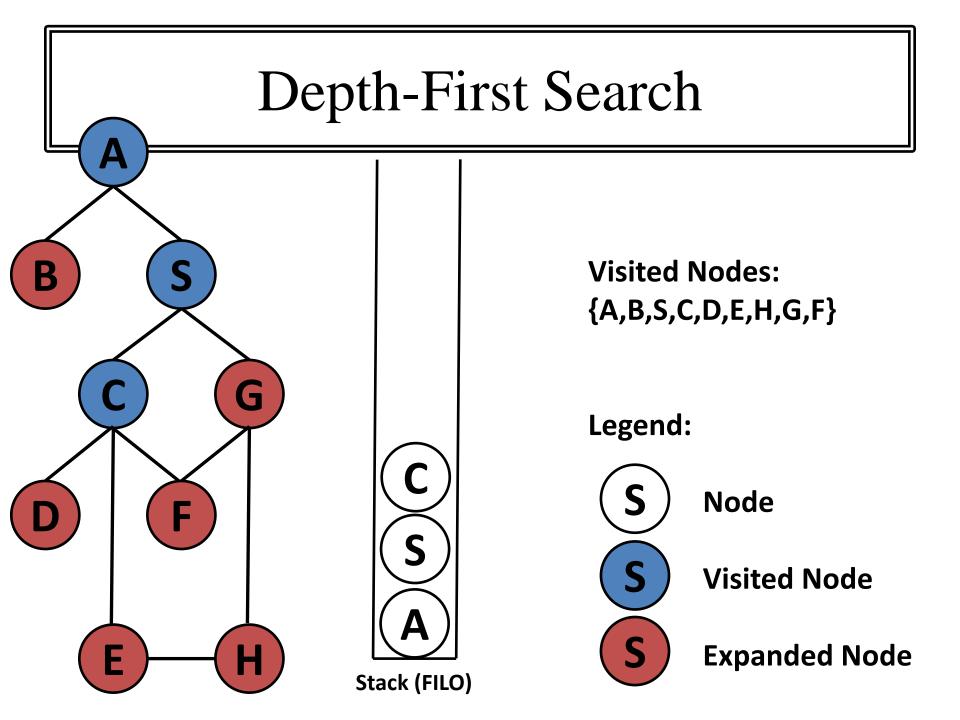


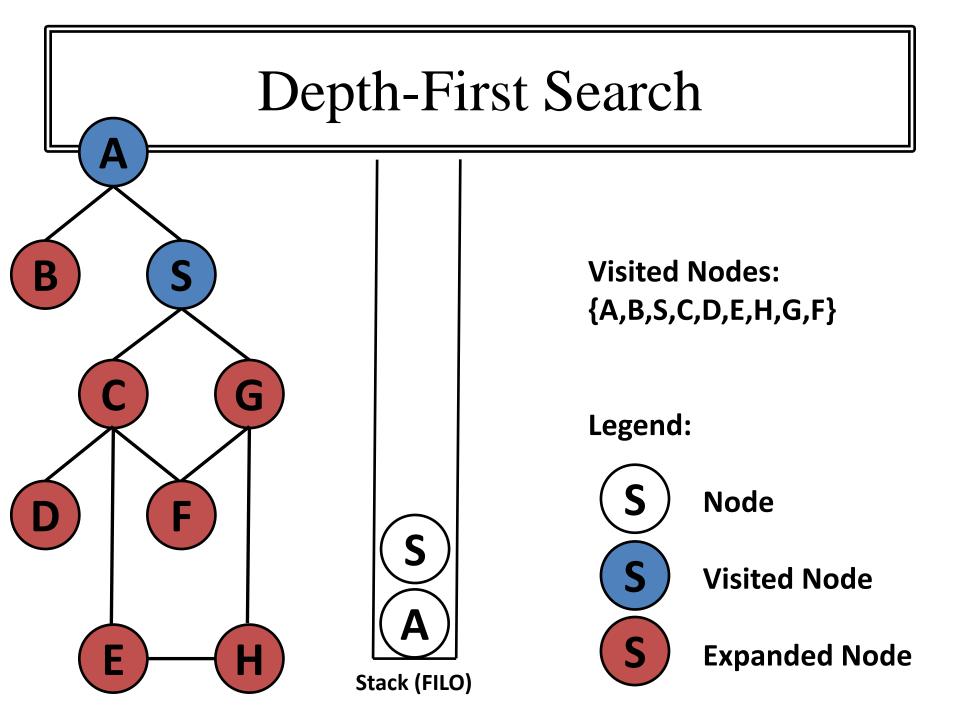


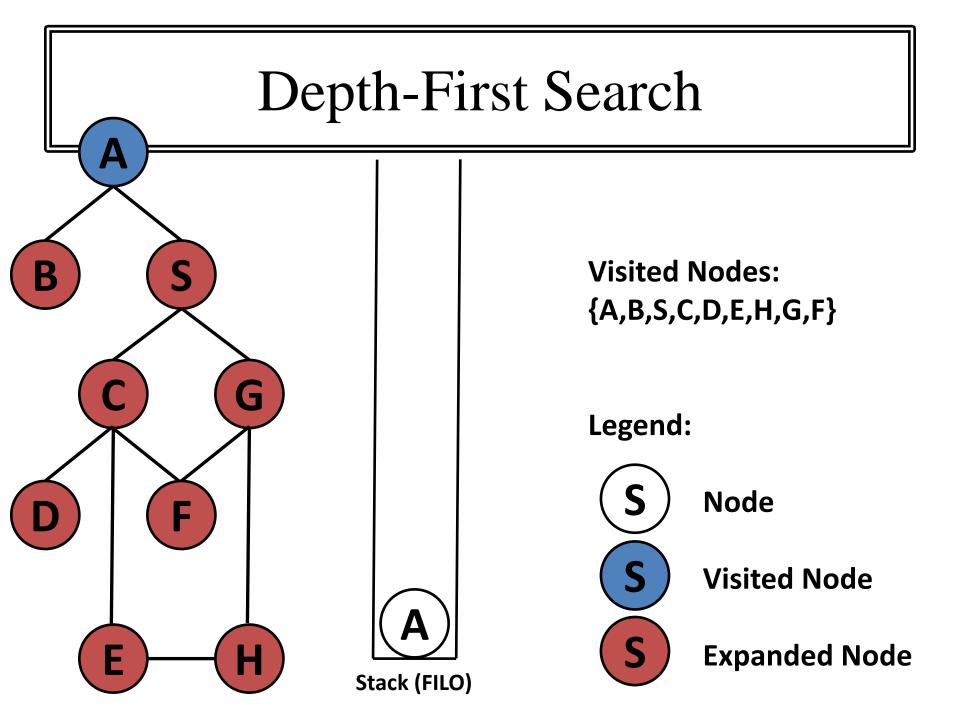


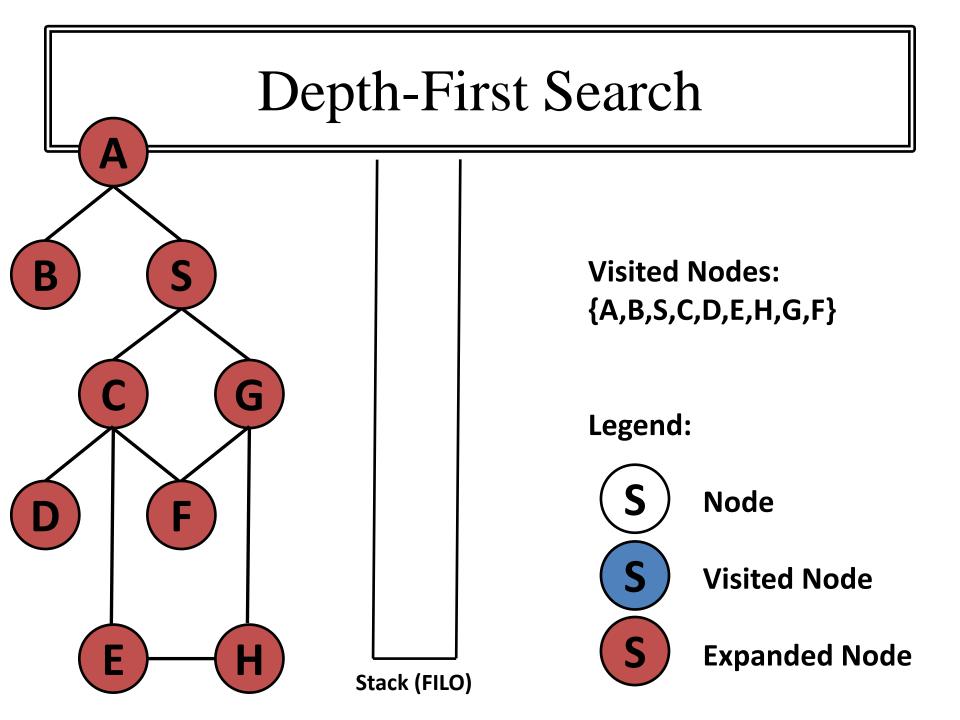












### Depth-First Search (DFS)

#### Complete?

- No: fails in infinite-depth spaces, spaces with loops
   Modify to avoid repeated states along path
  - *⇒* complete in finite spaces

#### Optimal?

- No

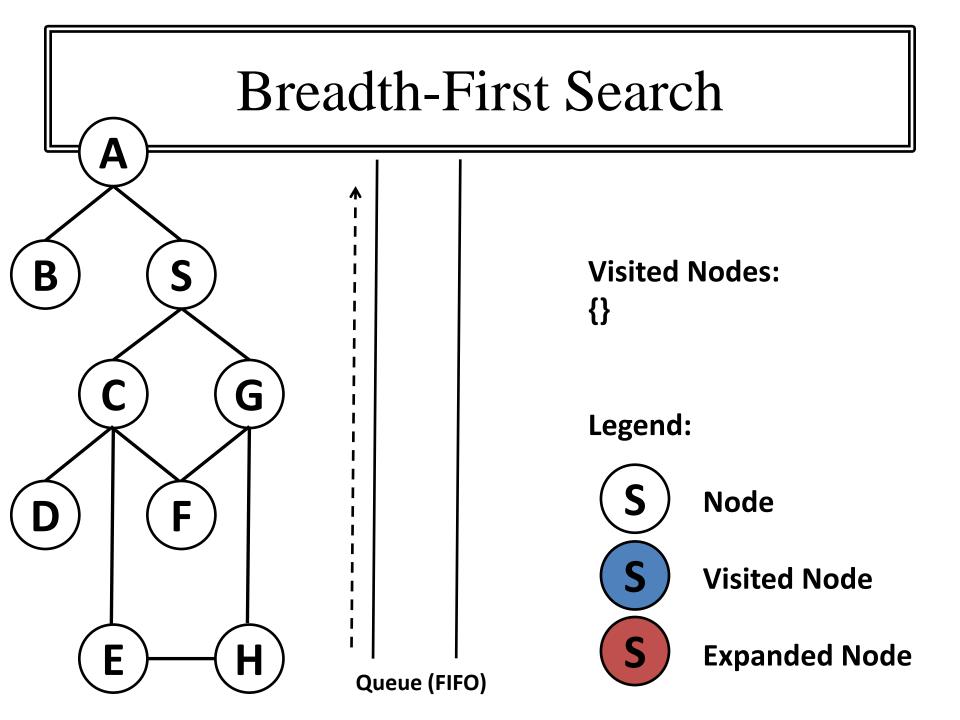
#### Time?

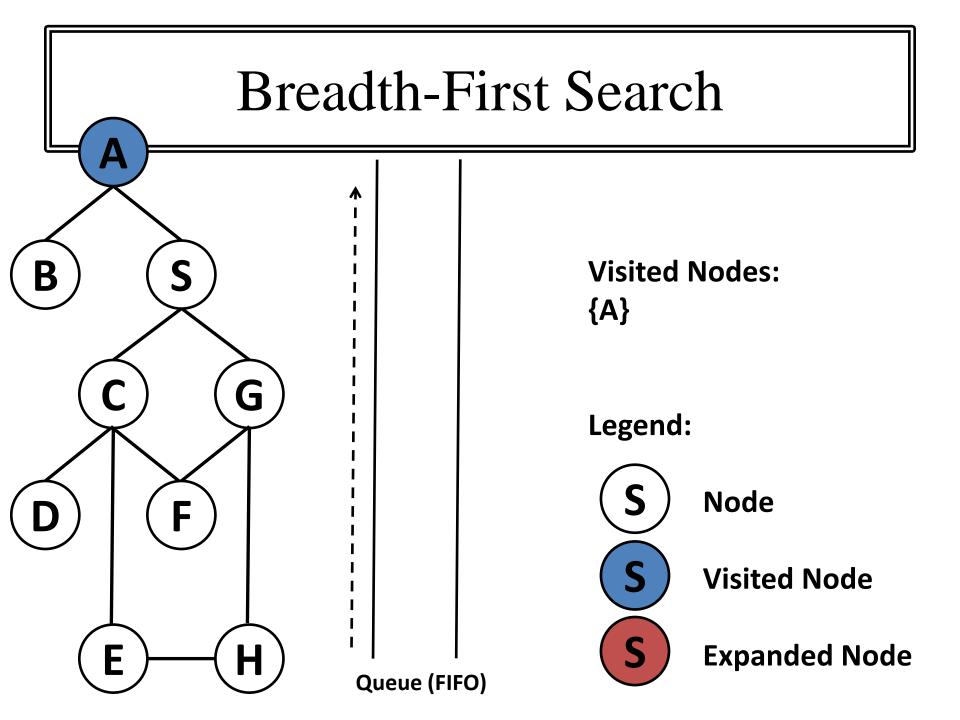
O(b<sup>m</sup>): terrible if m is much larger than d
 but if solutions are dense, may be much faster than breadth-first

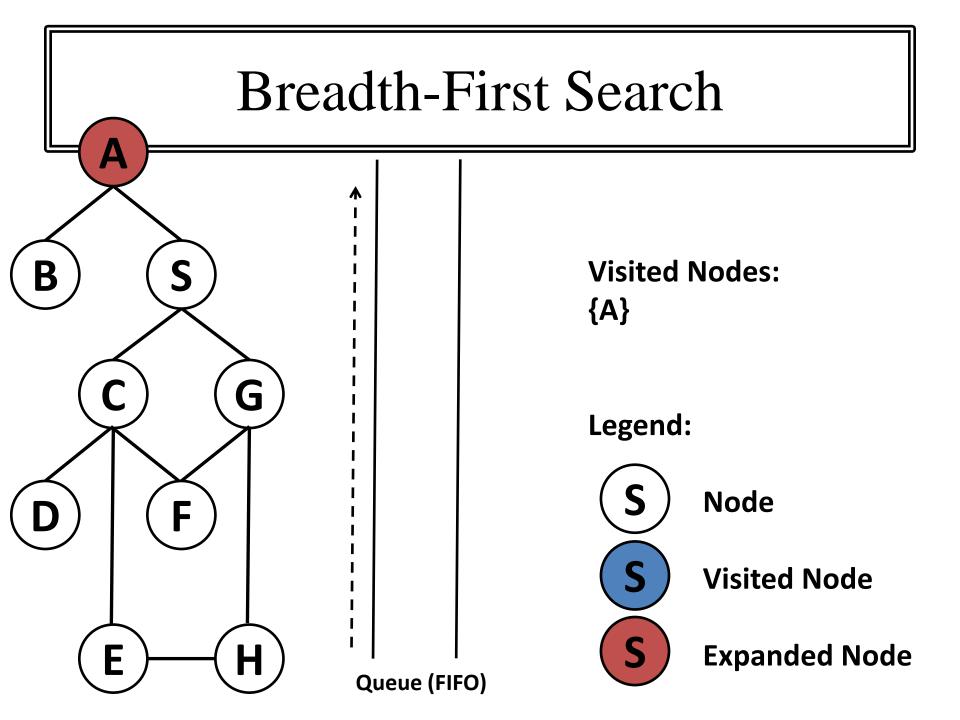
#### • Space?

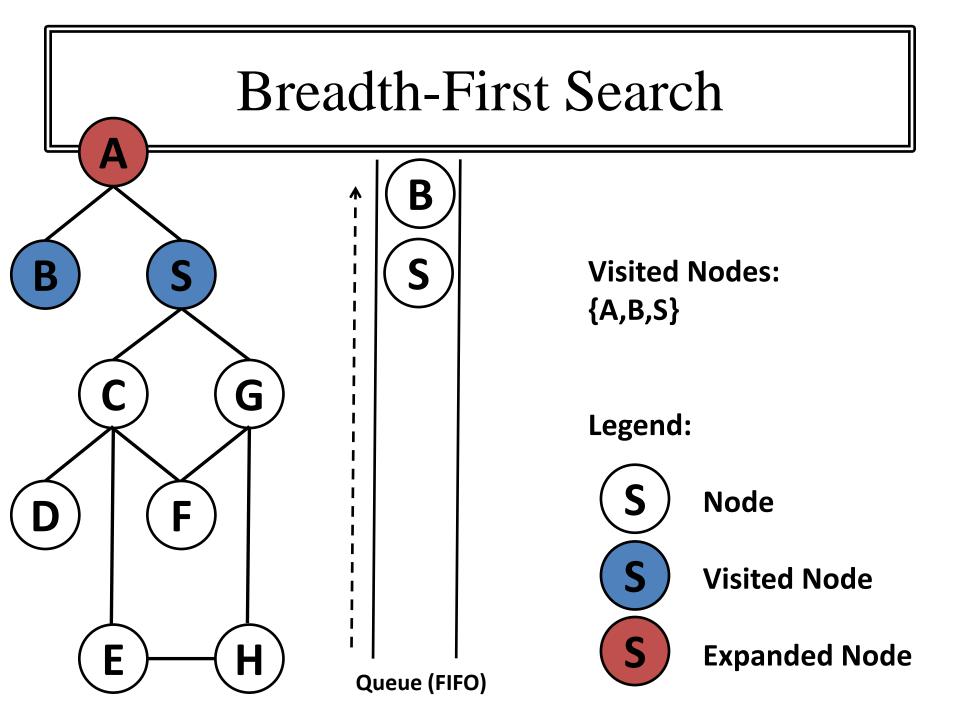
— O(bm), i.e., linear space!

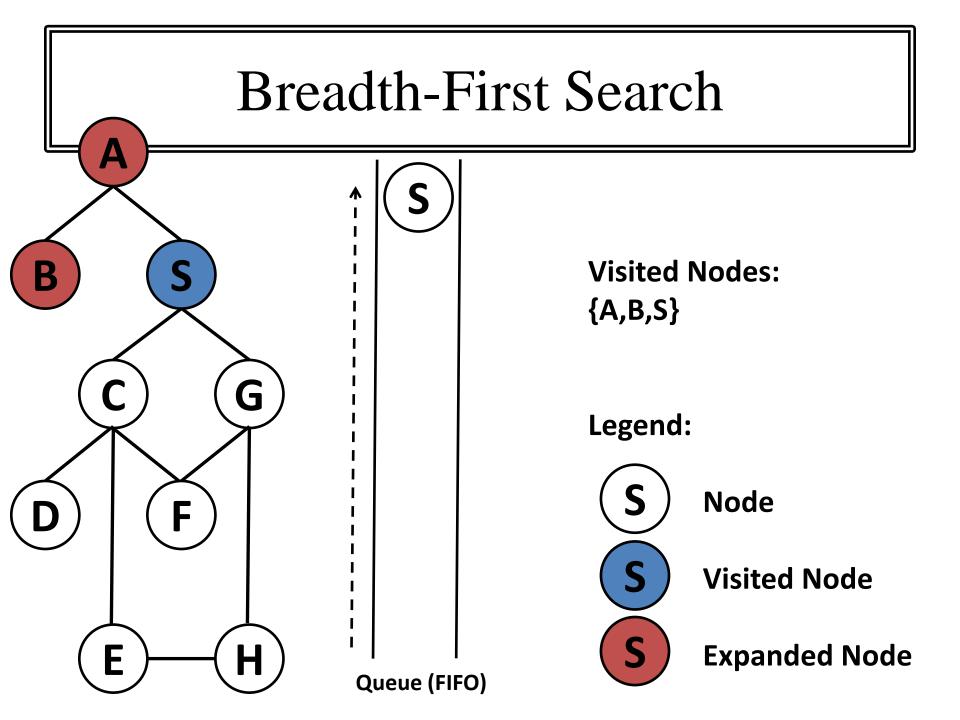
<sup>\*</sup> A <u>dense graph</u> is a graph in which the number of edges is close to the maximal number of edges.

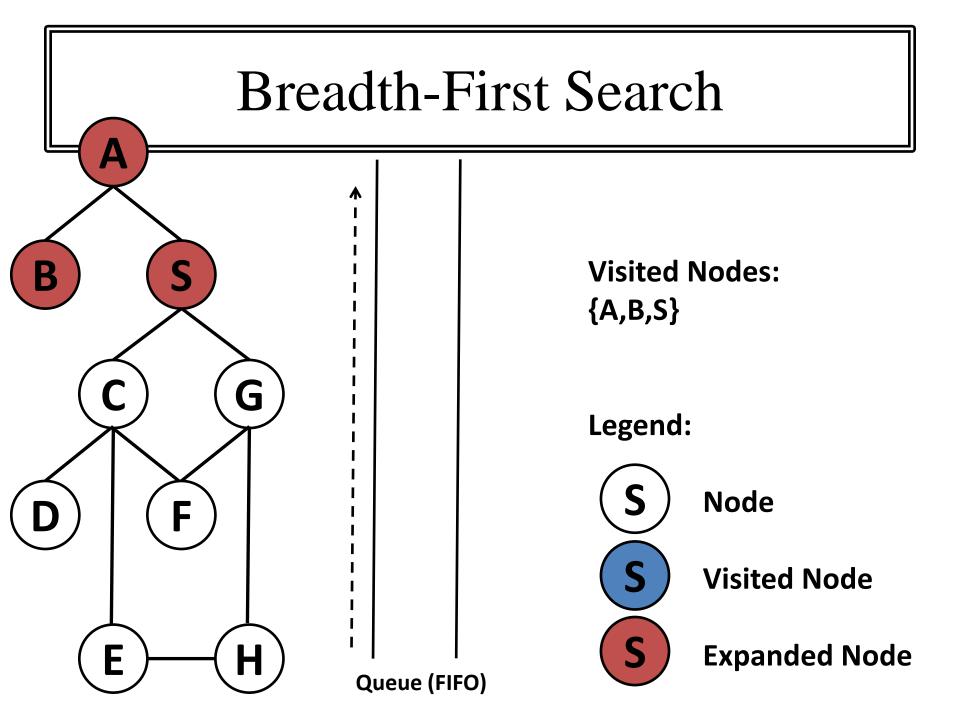


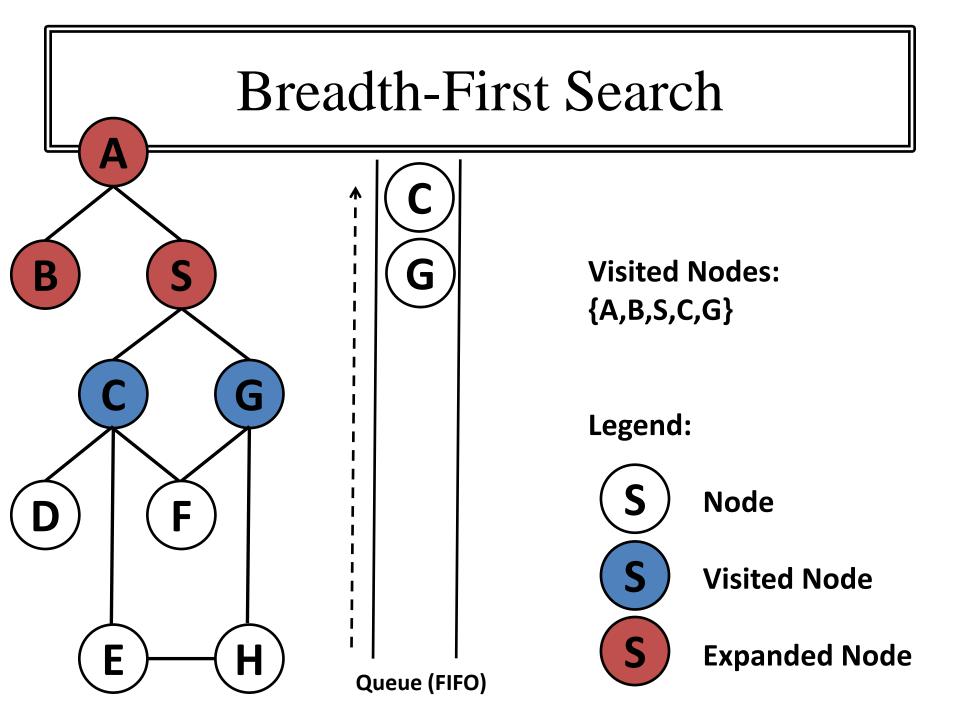


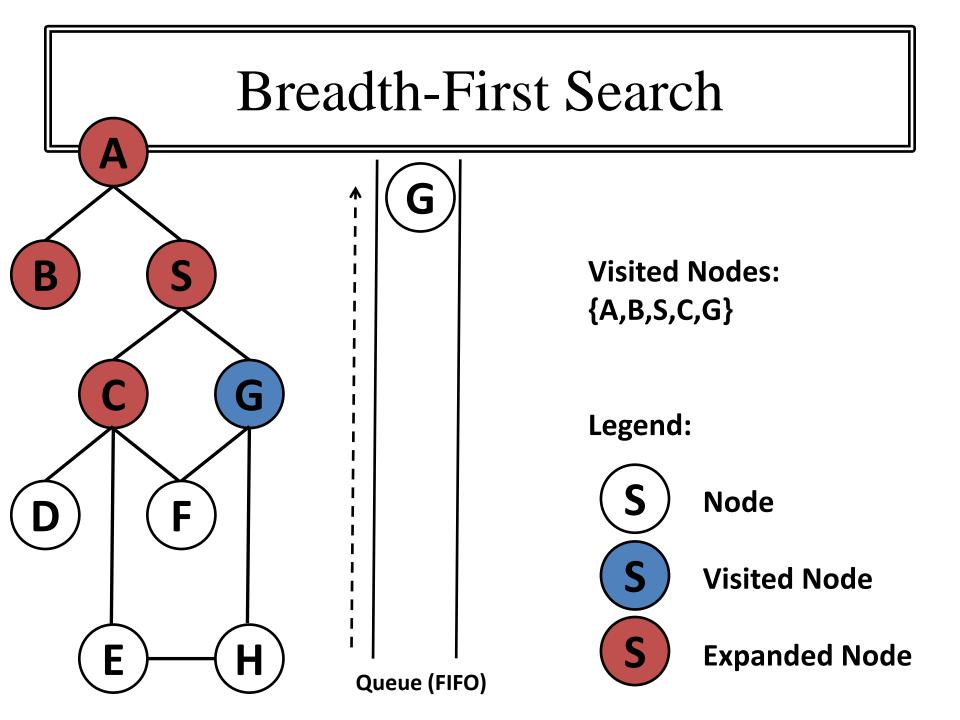


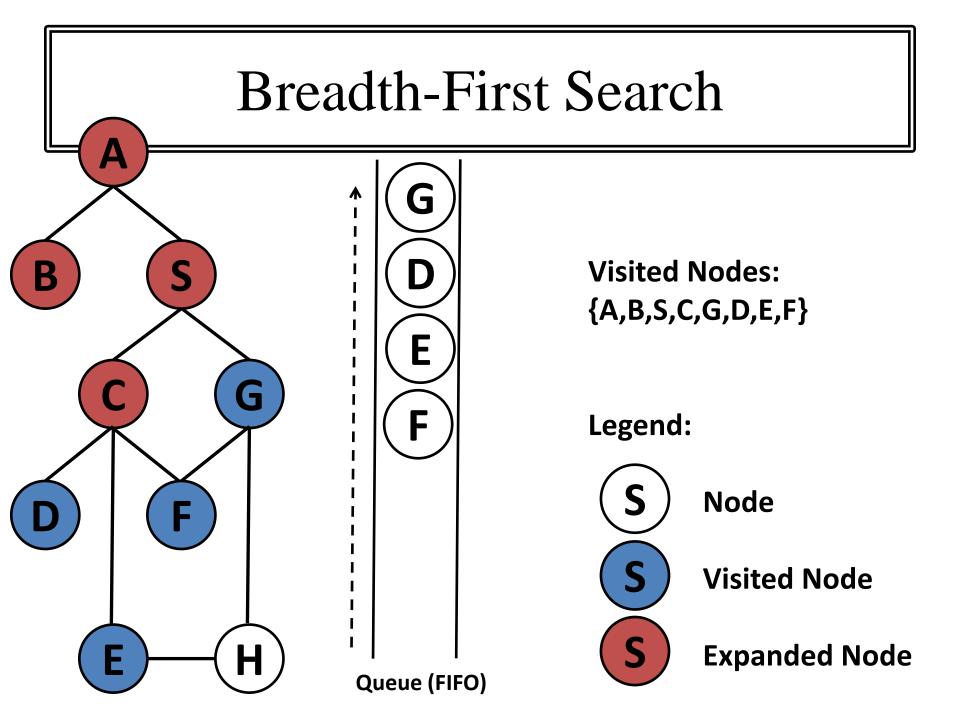


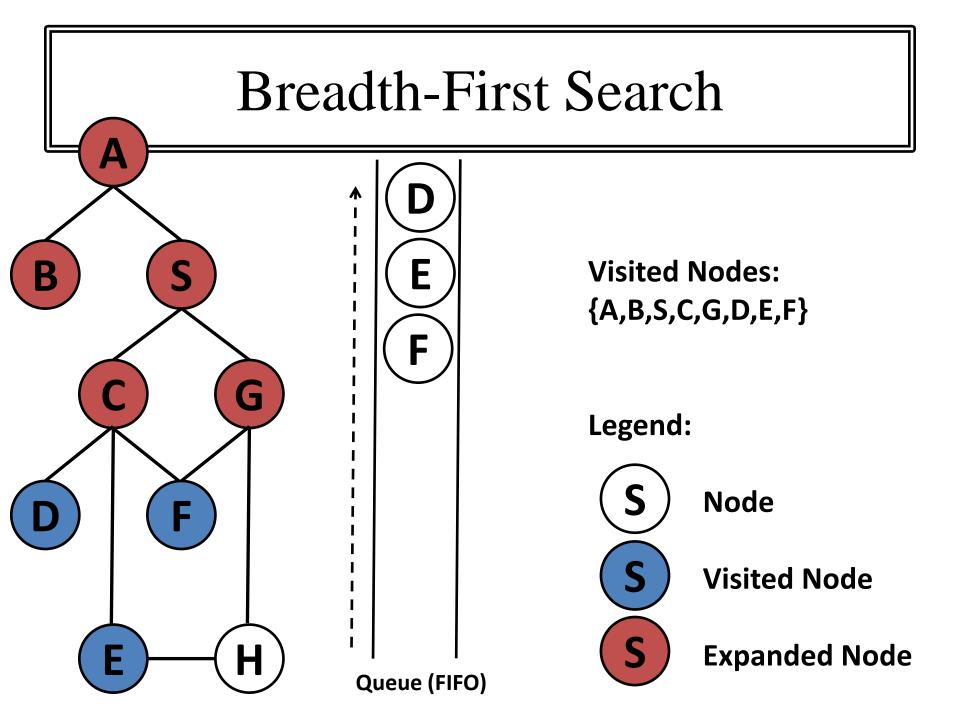


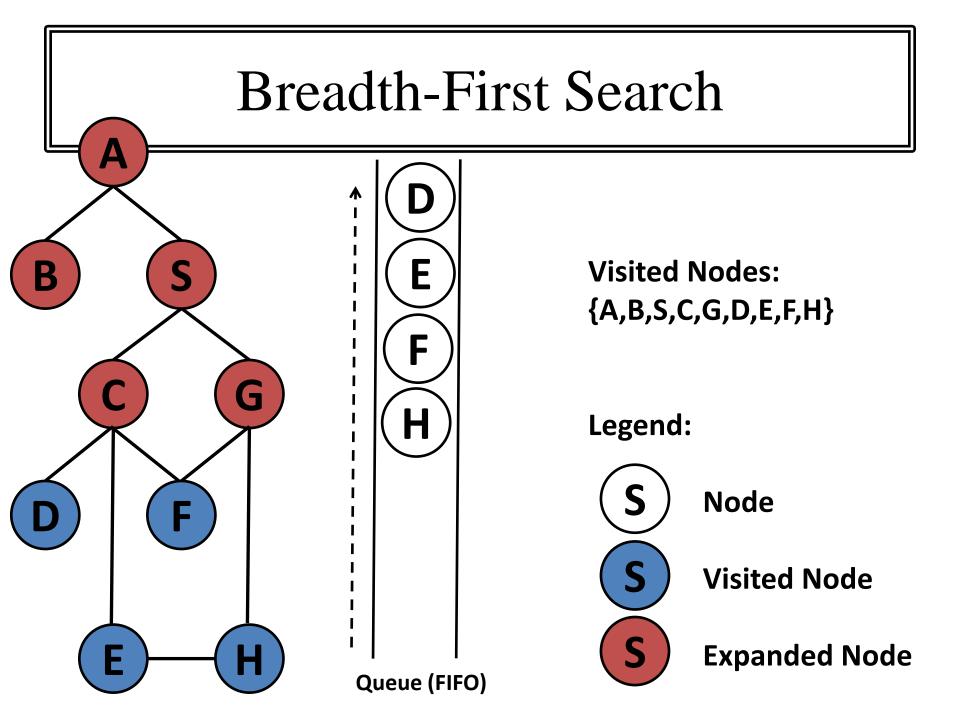


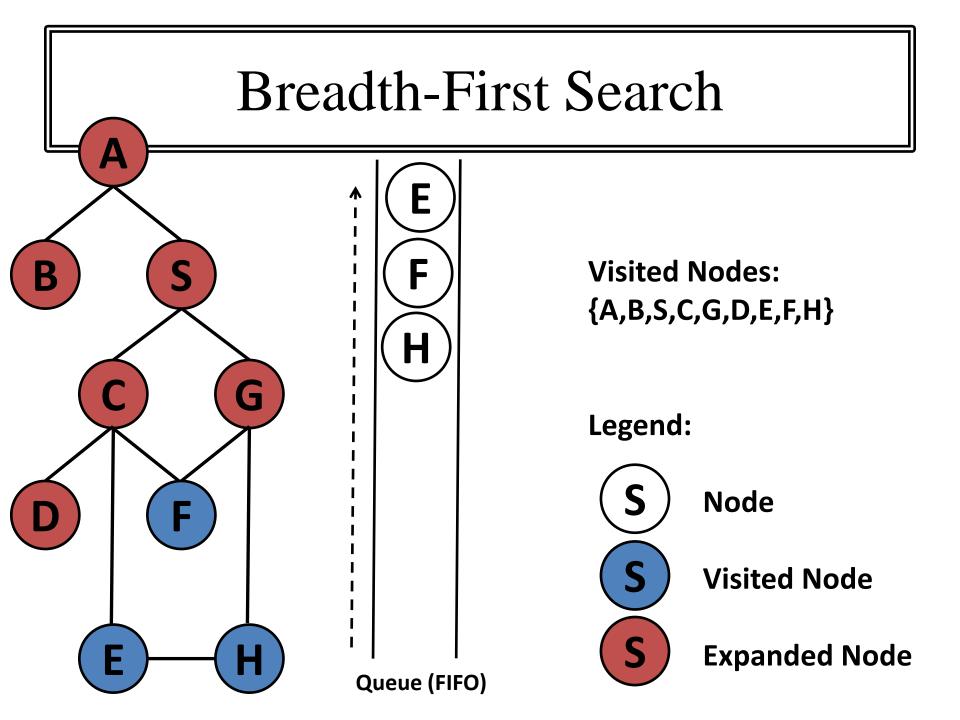


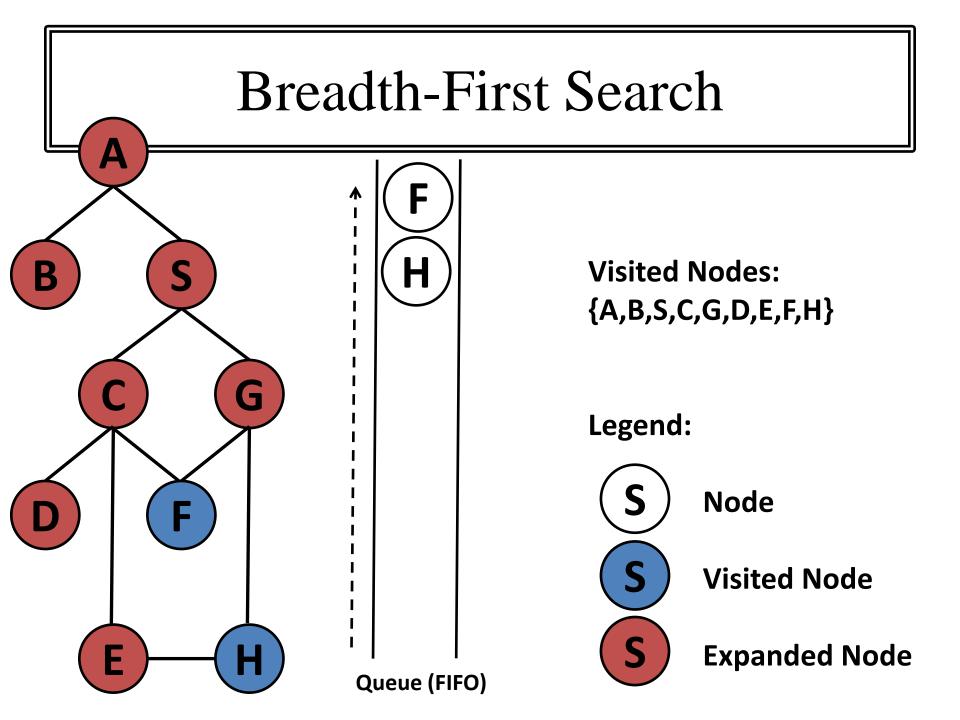


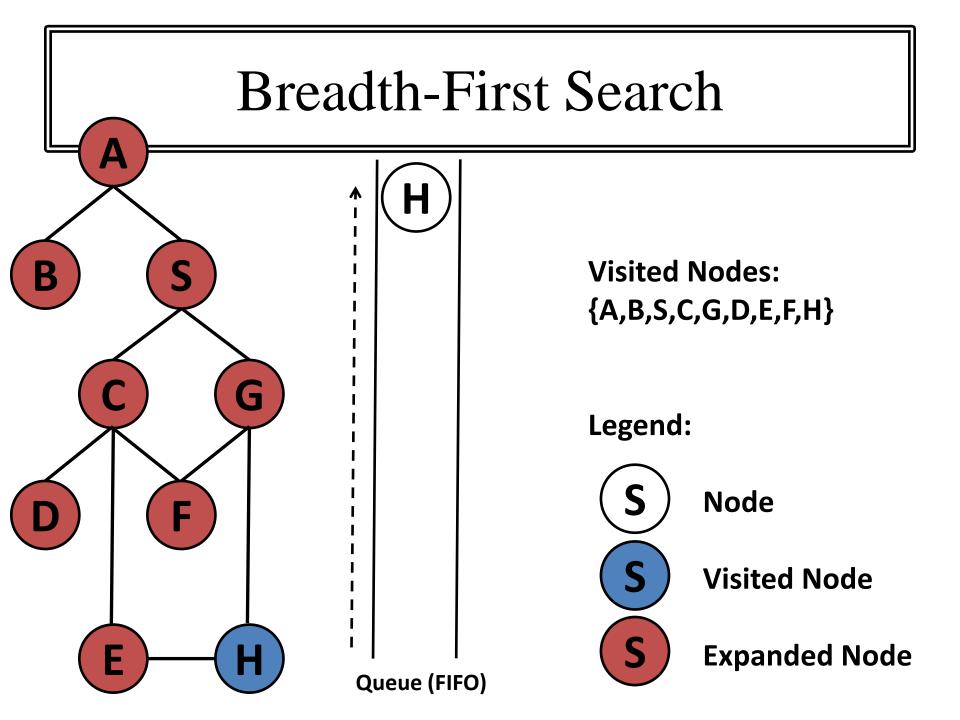


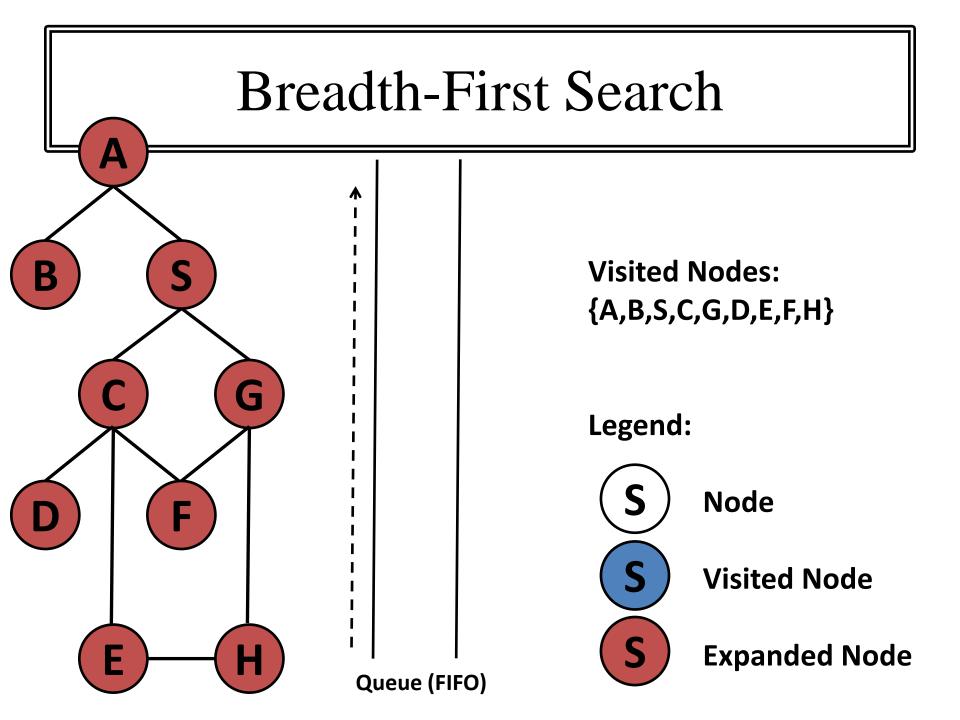












#### Breadth-First Search

```
Node {
    Node parent
    T value
    List<Node> children
}
```

## Breadth-First Search (BFS)

#### Complete?

– Yes (if b is finite)

#### Optimal?

— Yes (if cost = 1 per step); not optimal in general

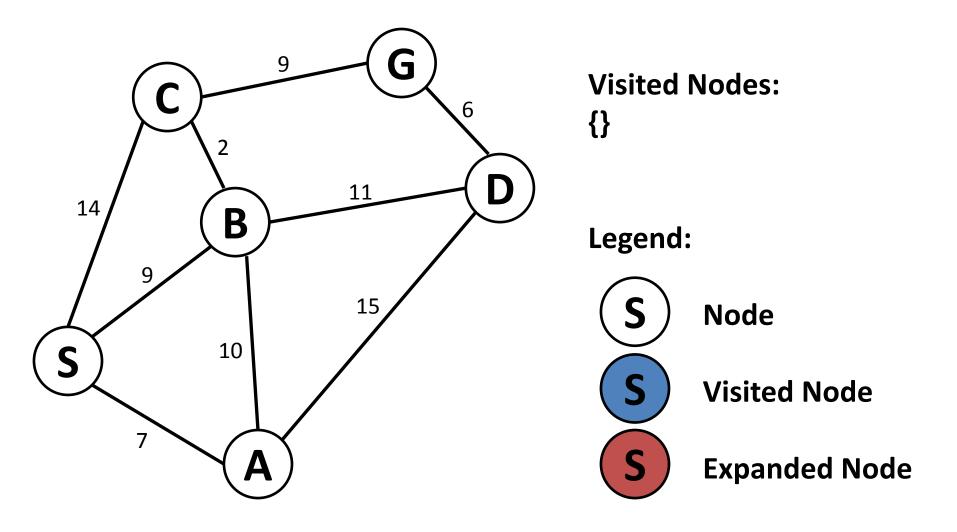
#### Time?

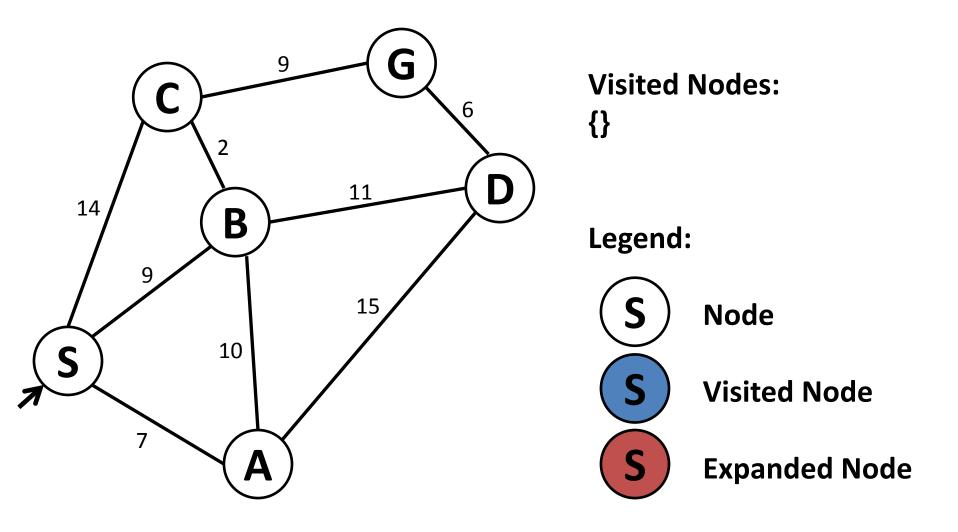
 $-1+b+b^2+b^3+...+b^d+b(b^d-1)=O(b^{d+1})$ , i.e., exp. in **d** 

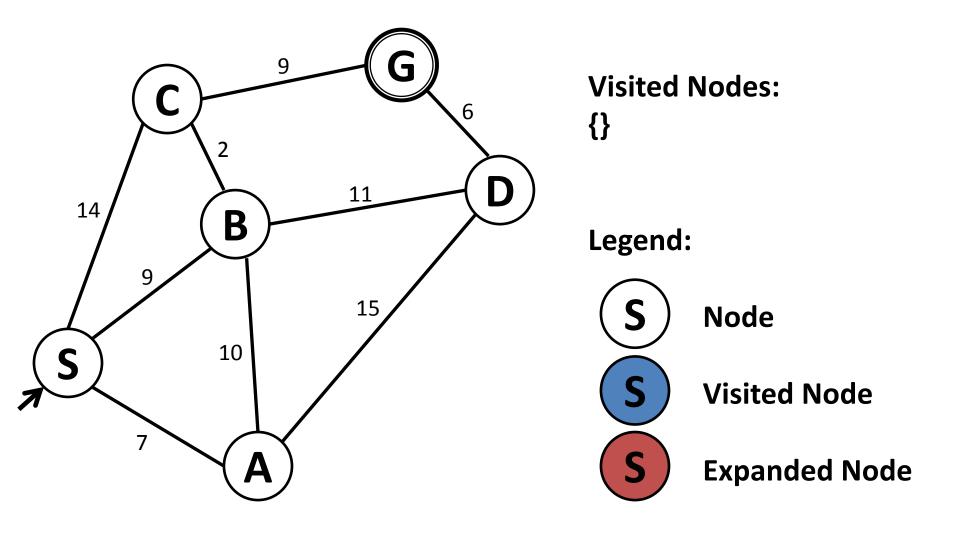
#### • Space?

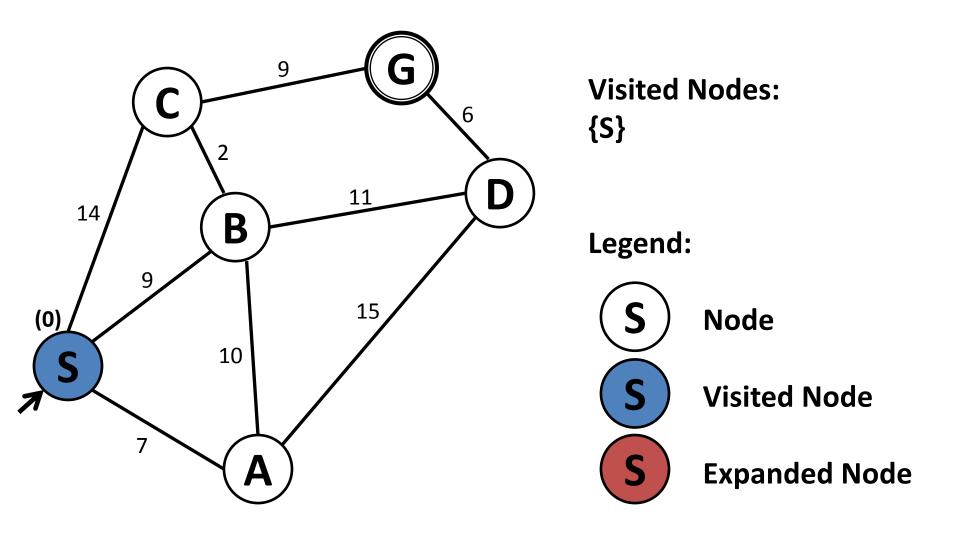
 $-O(b^{d+1})$  (keeps every node in memory)

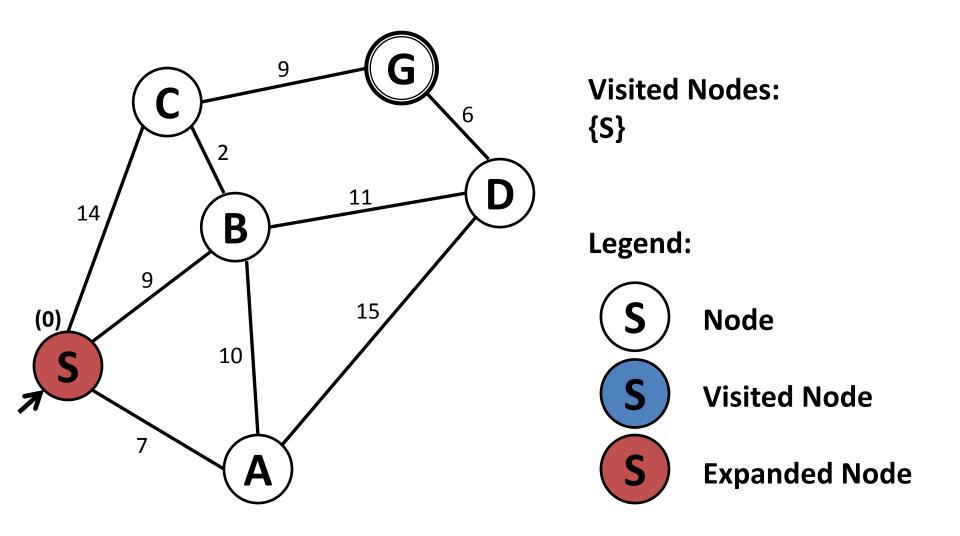
<sup>\*</sup> **Space** is the big problem; can easily generate nodes at 100MB/sec so 24hrs = 8640GB.

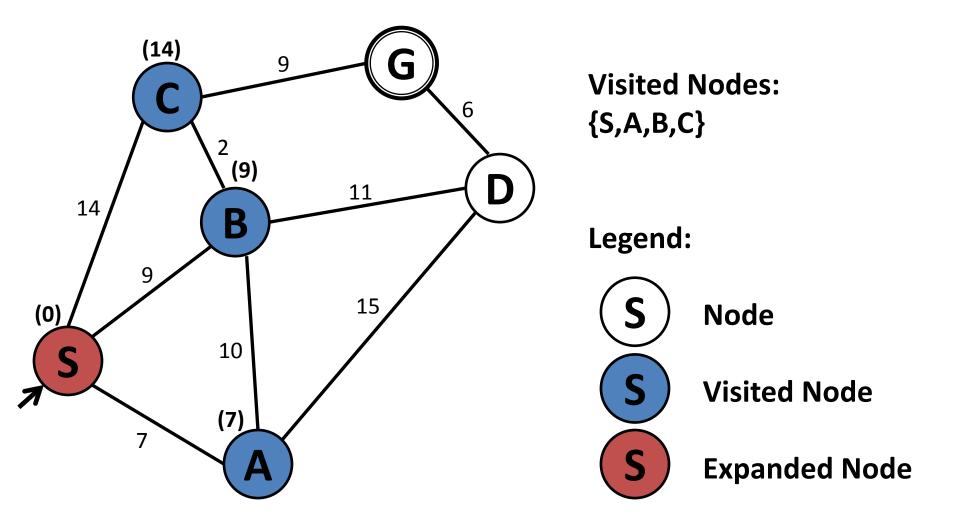


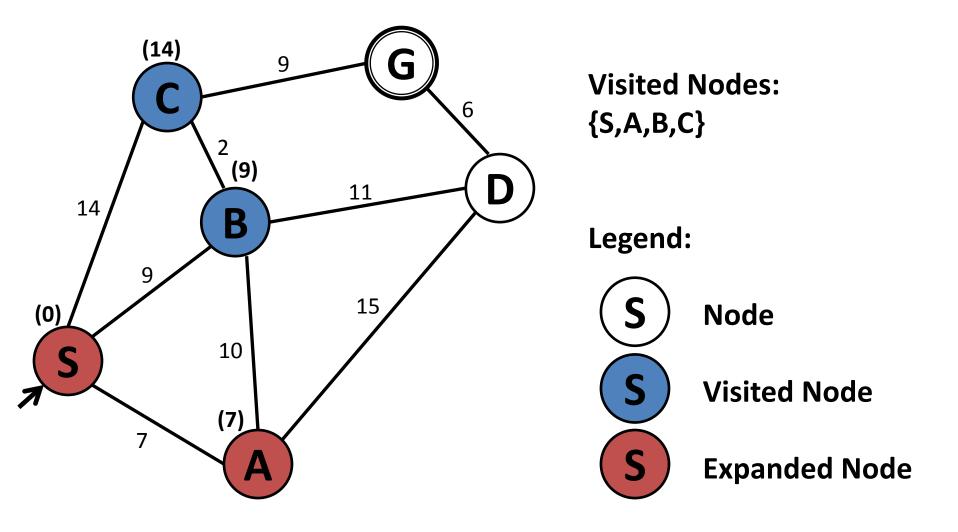


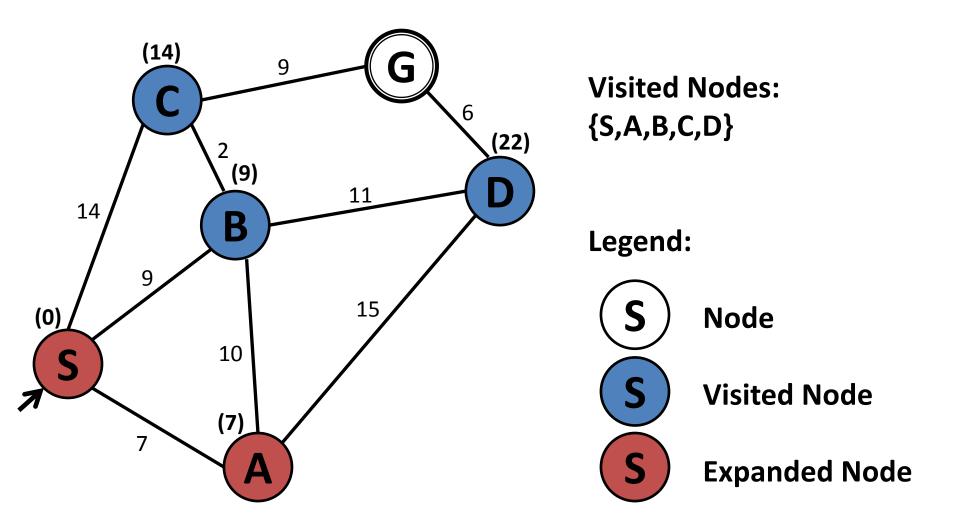


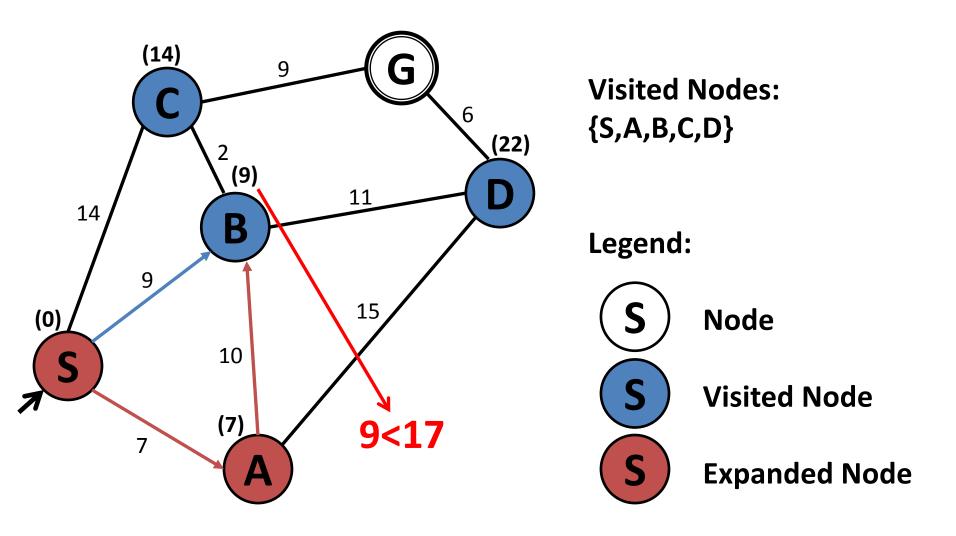


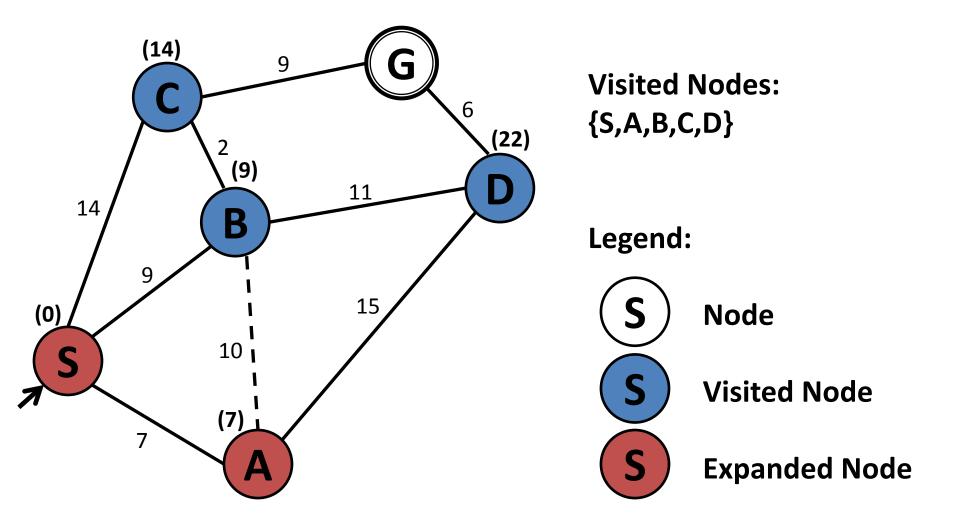


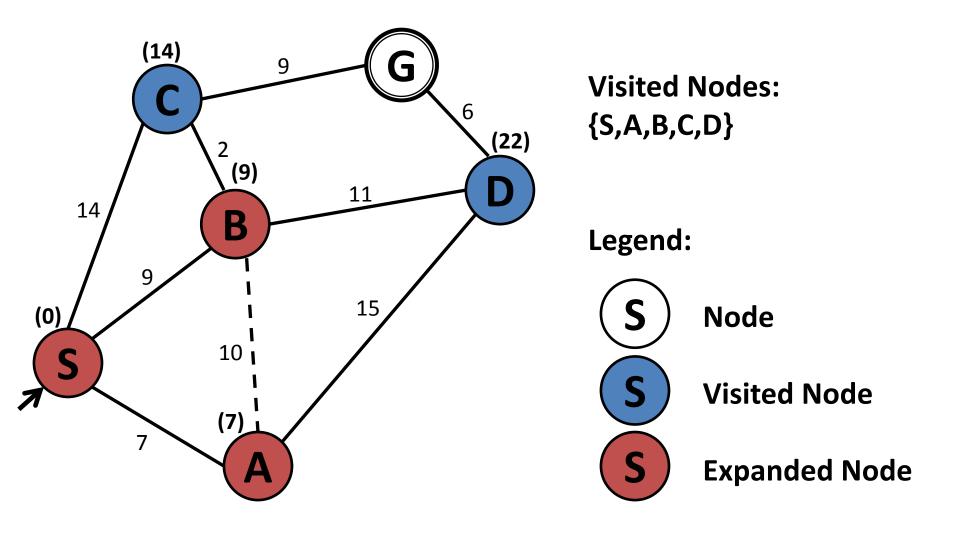


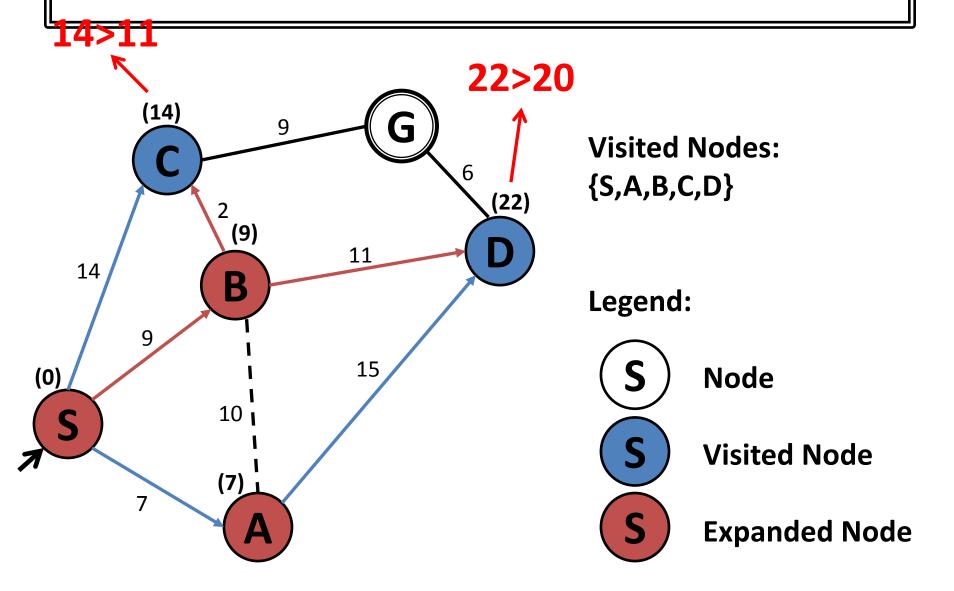


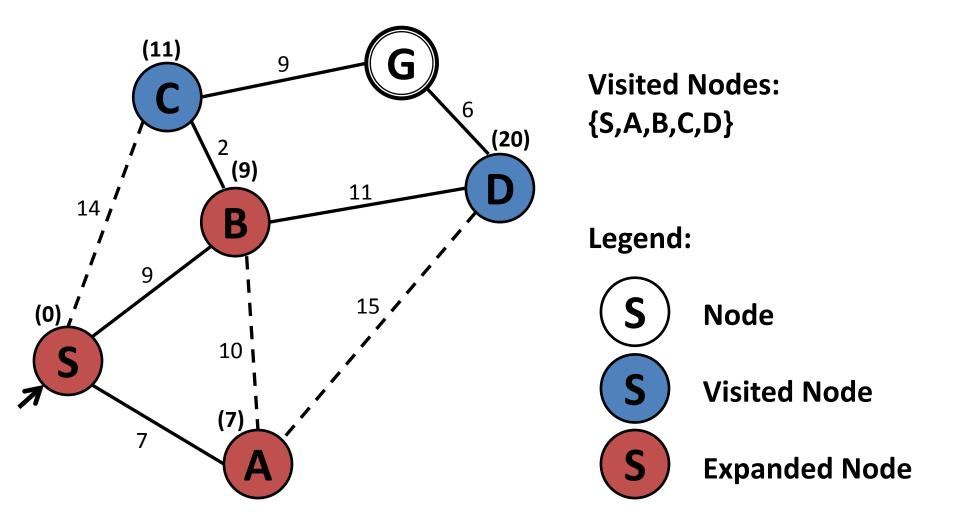


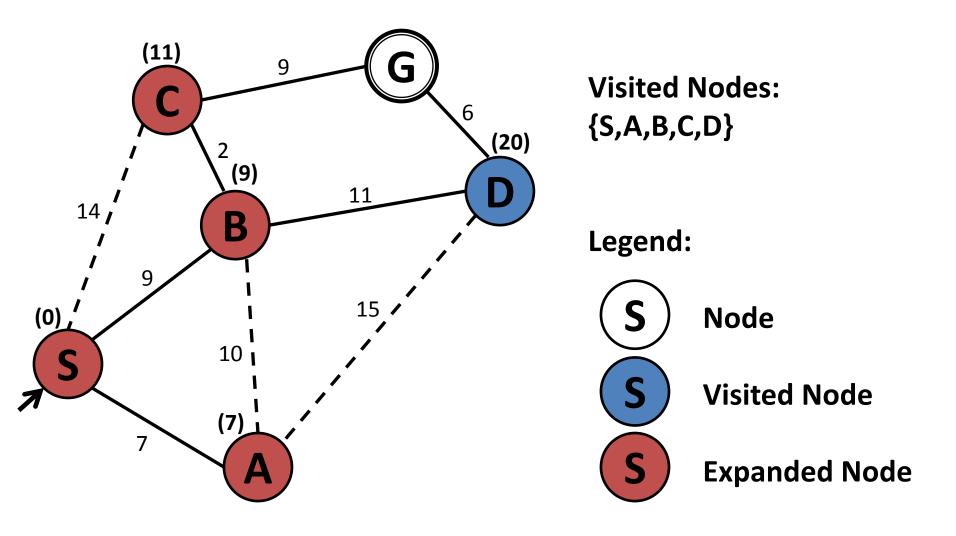


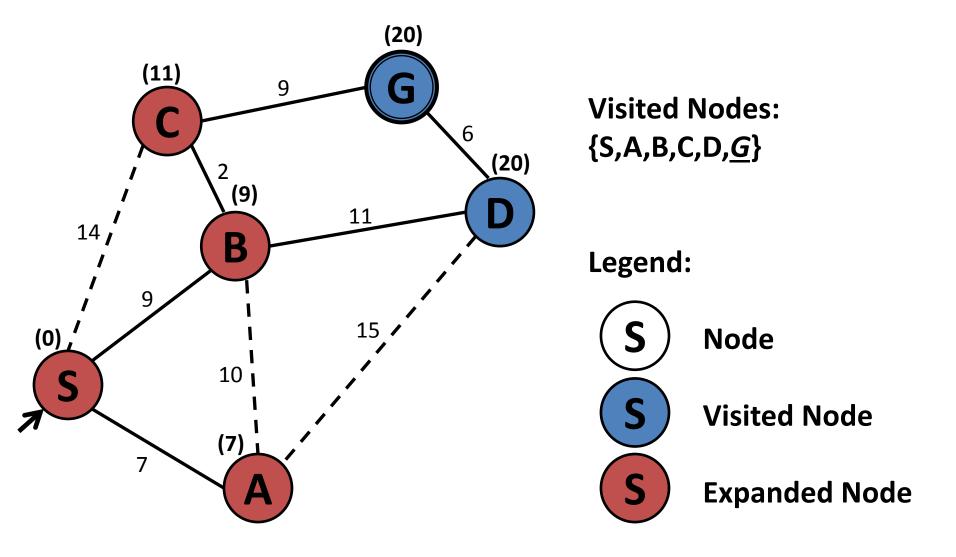


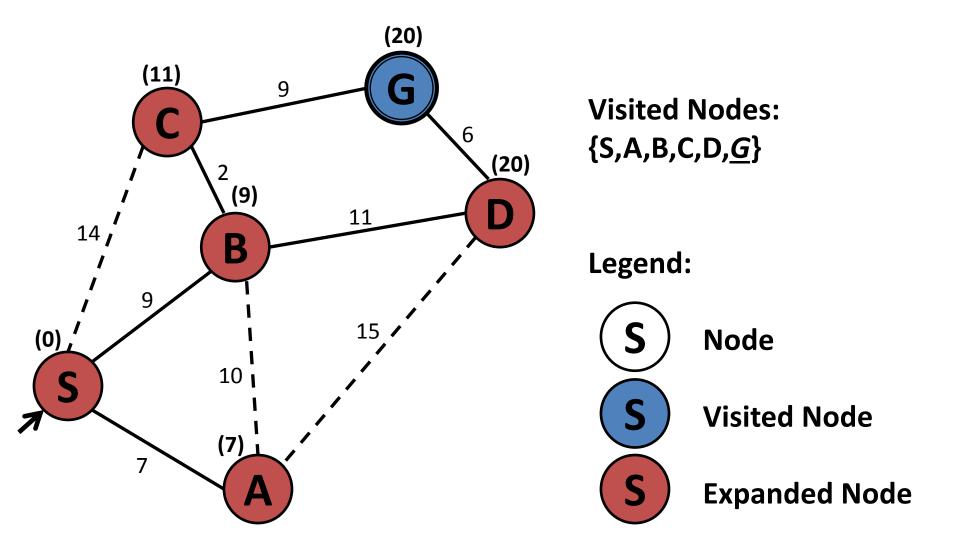


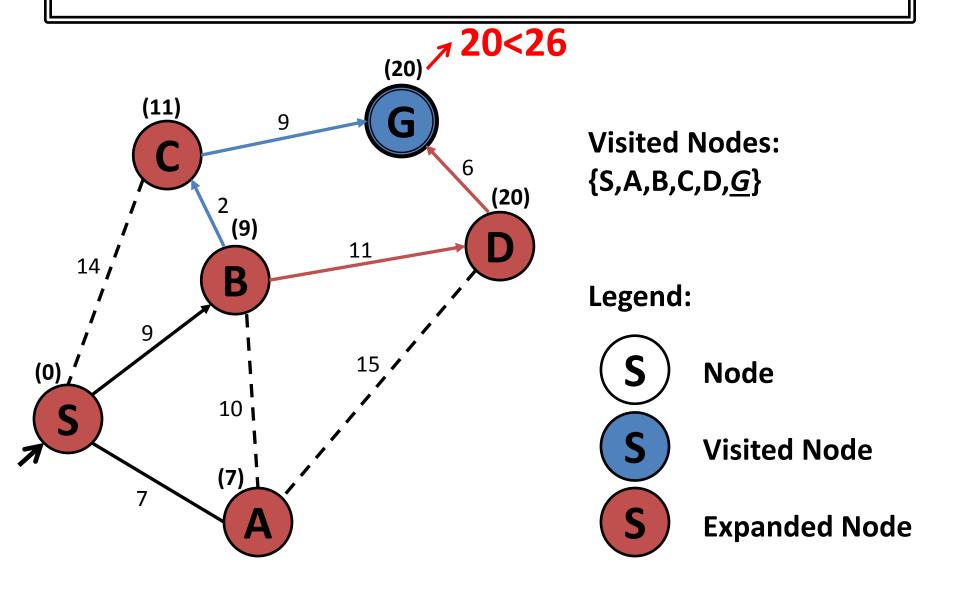


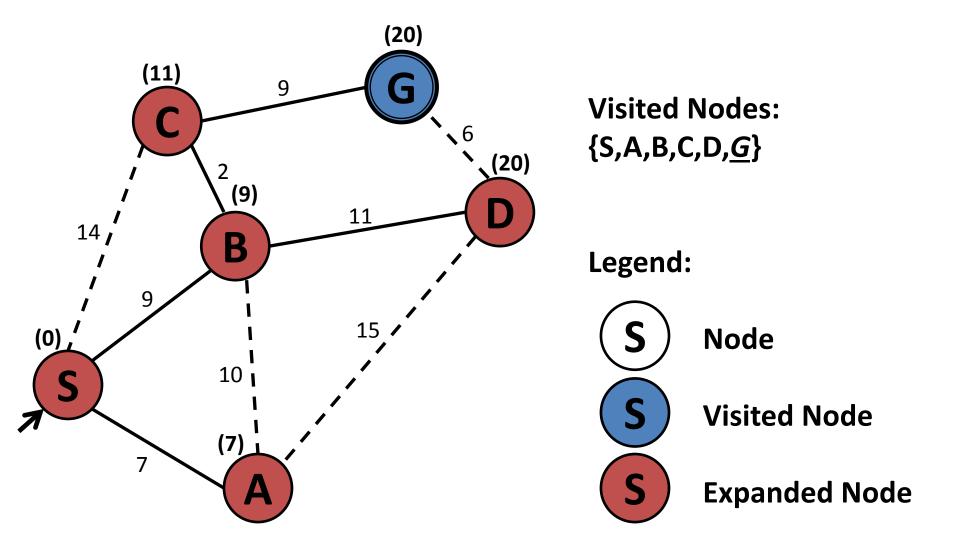


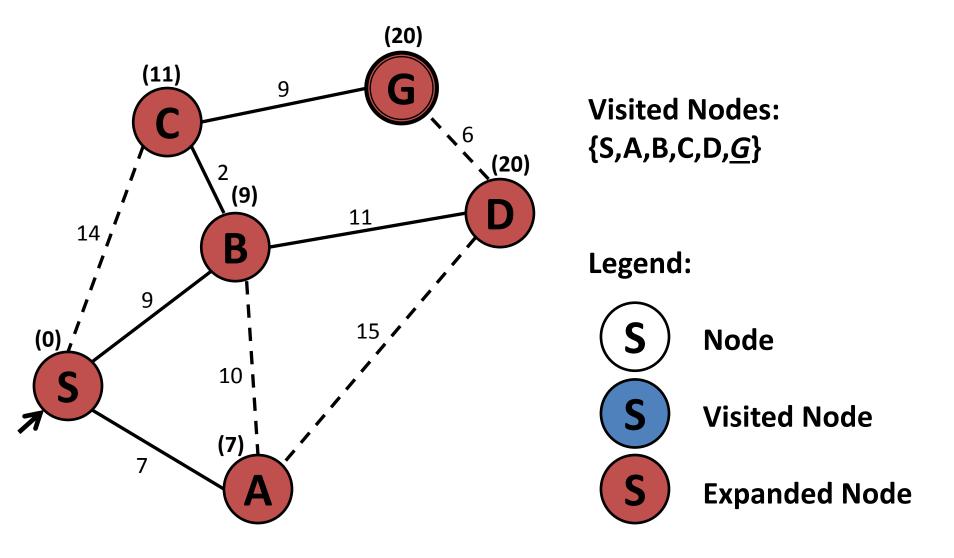


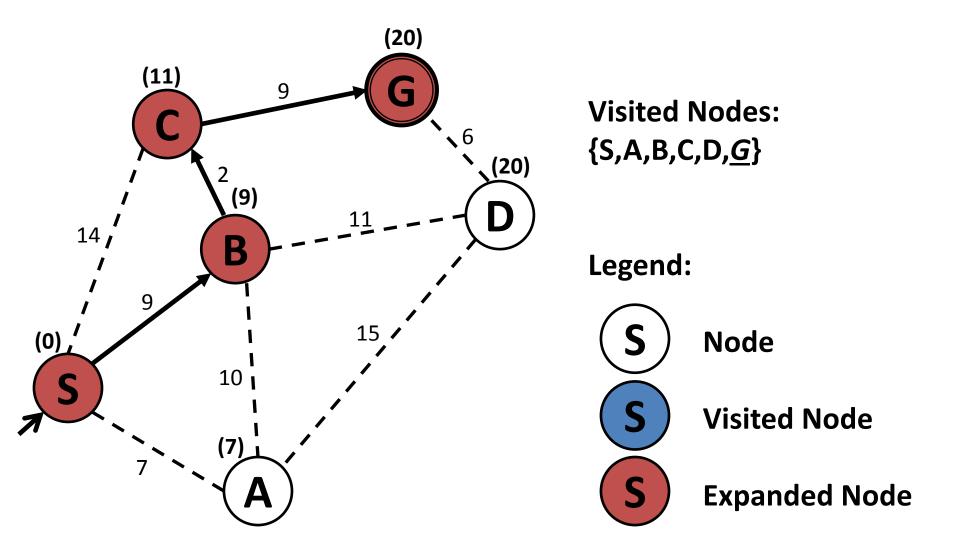












## Uniform-Cost Search (UCS)

#### Complete?

- **Yes**, if step cost ≥ ε

#### Optimal?

Yes - nodes expanded in increasing order of g(n)

#### • <u>Time?</u>

- # of nodes with  $g \le cost$  of optimal solution,  $O(b^{\lceil C^*/\epsilon \rceil})$  where  $C^*$  is the cost of the optimal solution

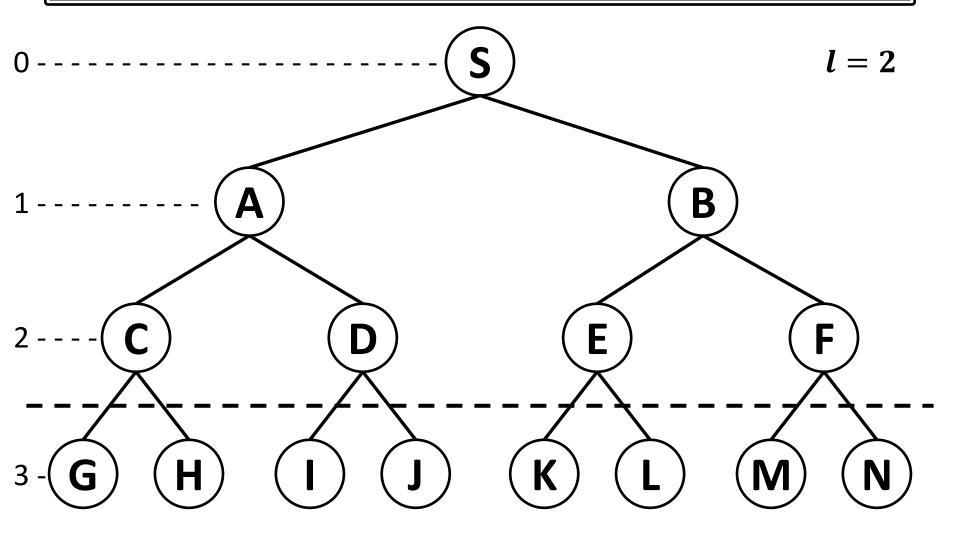
#### • Space?

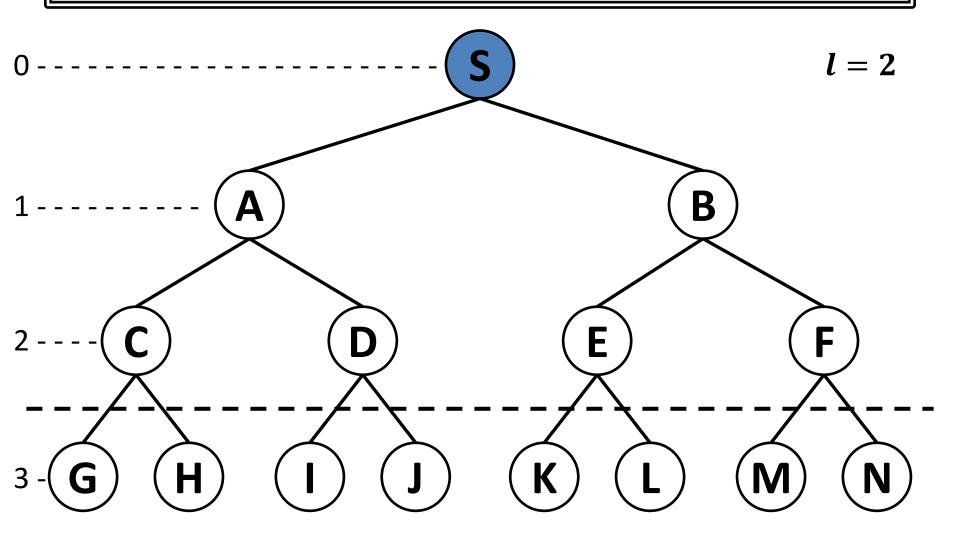
- # of nodes with  $g \le cost$  of optimal solution,  $O(b^{\lceil C^*/\epsilon \rceil})$ 

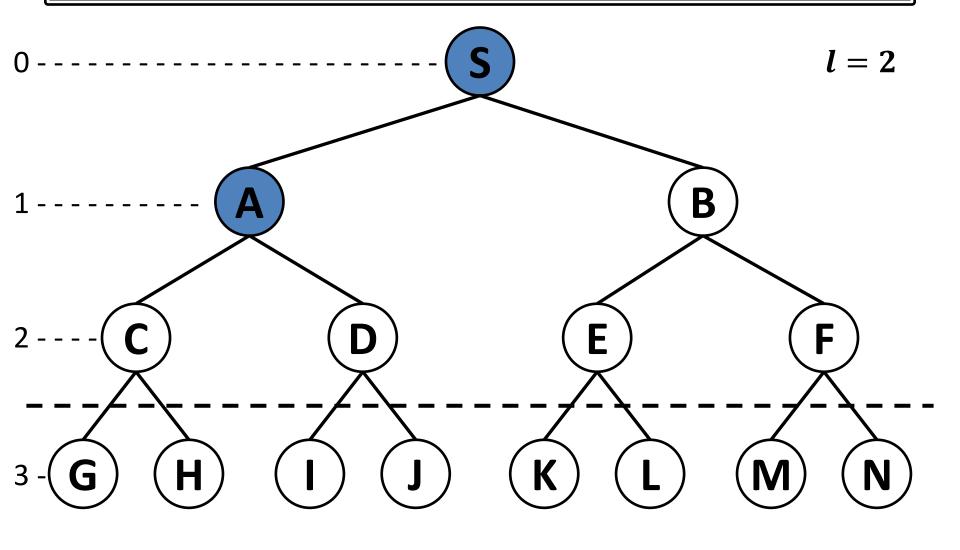
<sup>\*</sup> Expand least-cost unexpanded node using queue ordered by path cost, lowest first

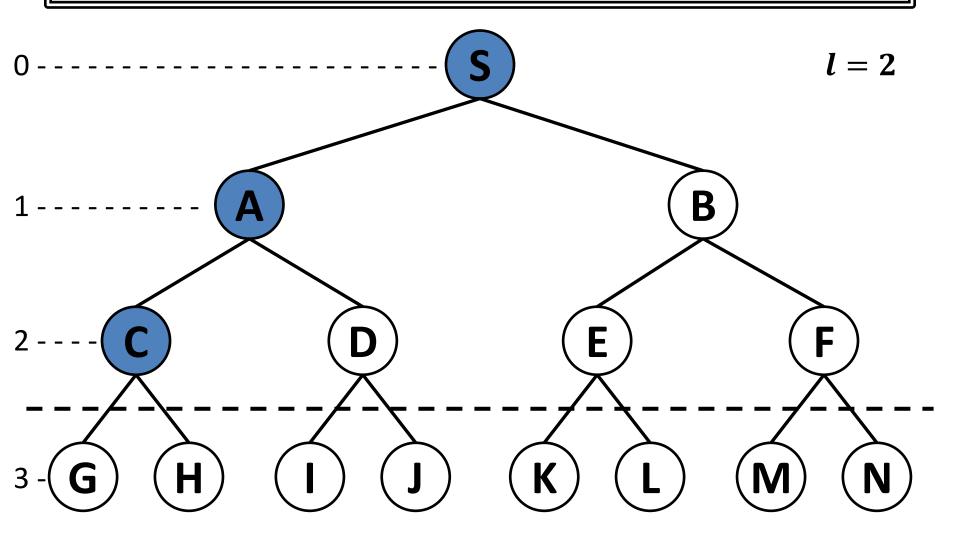
<sup>\*</sup> Equivalent to breadth-first if step costs all equal

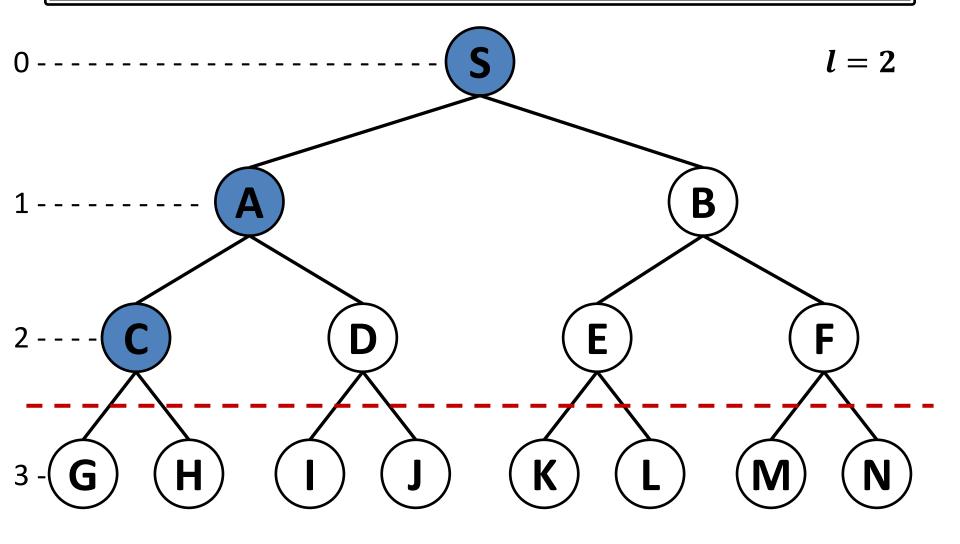
- = depth-first search with depth limit I,
   i.e., nodes at depth I have no successors
- Also known as Depth-Bounded Search

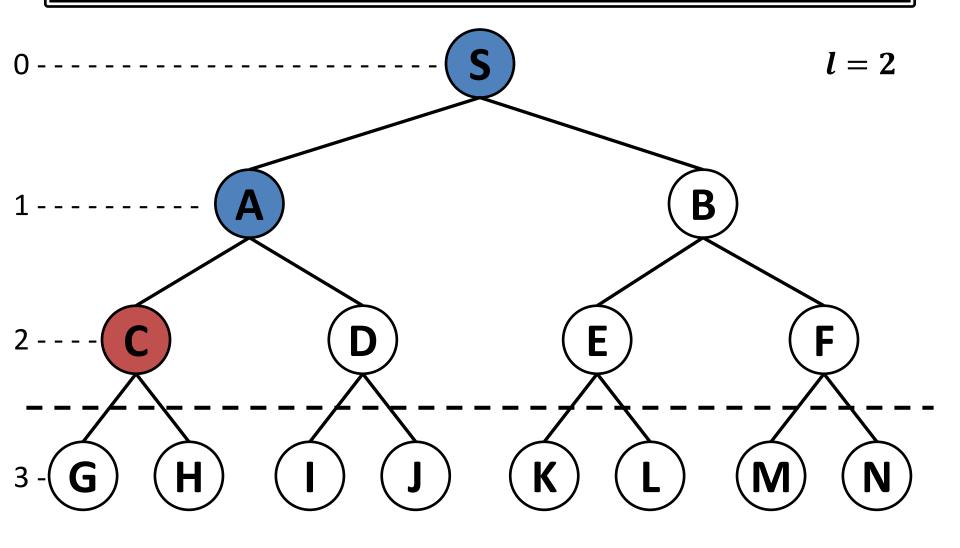


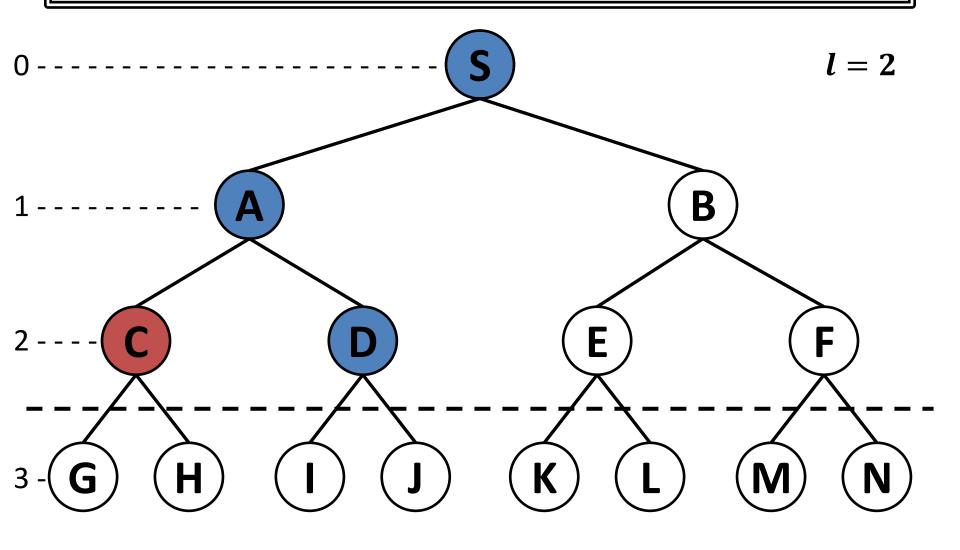


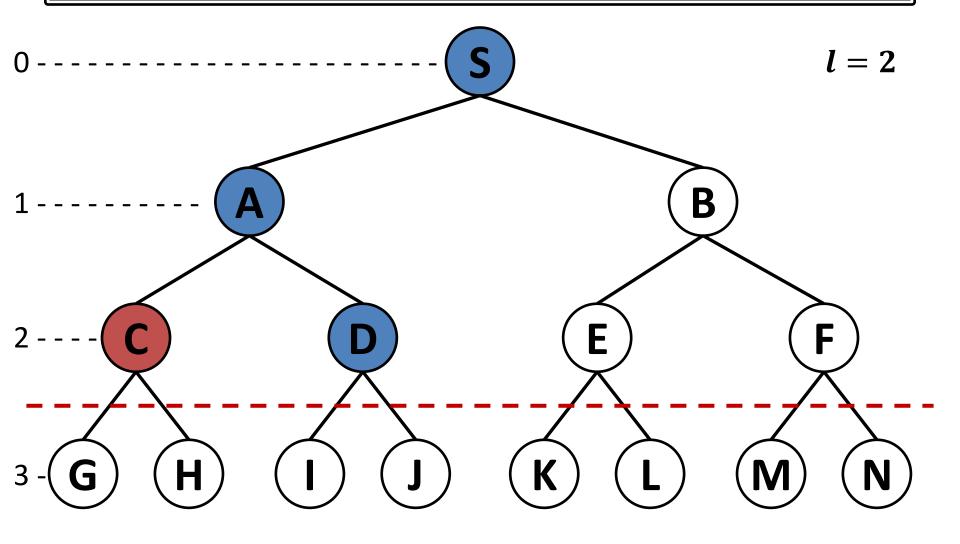


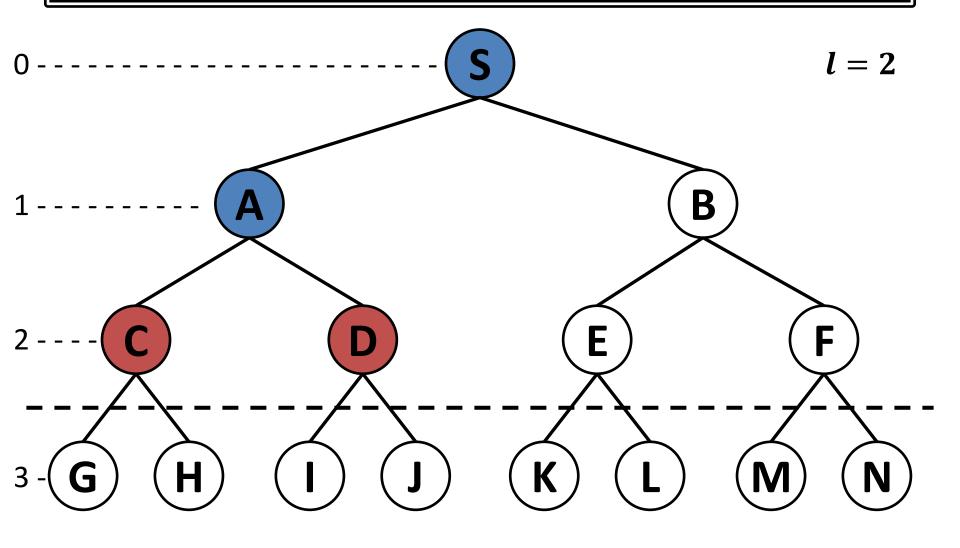


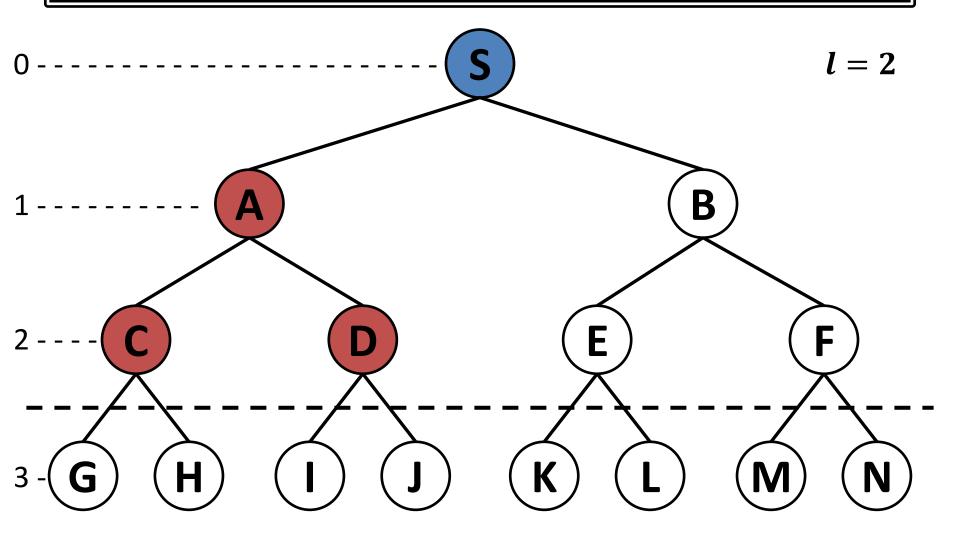


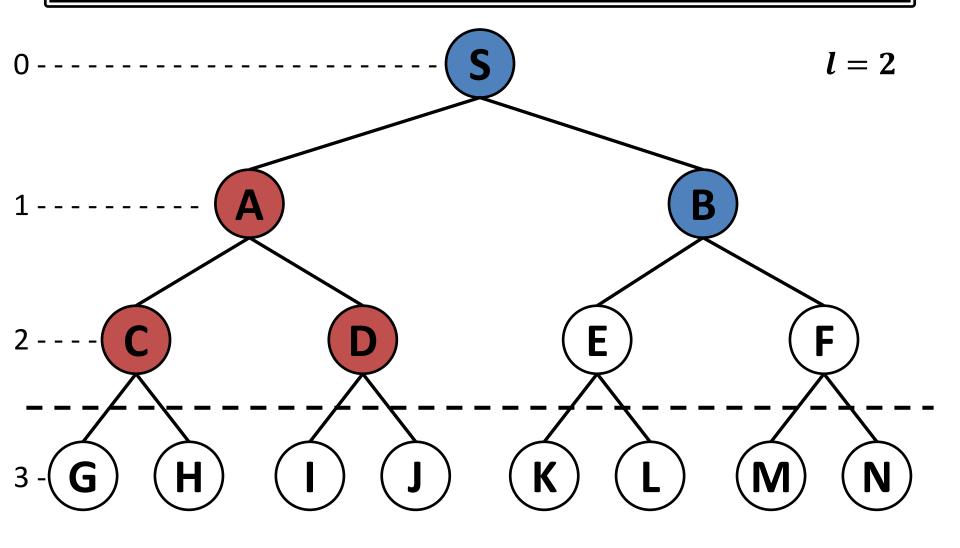


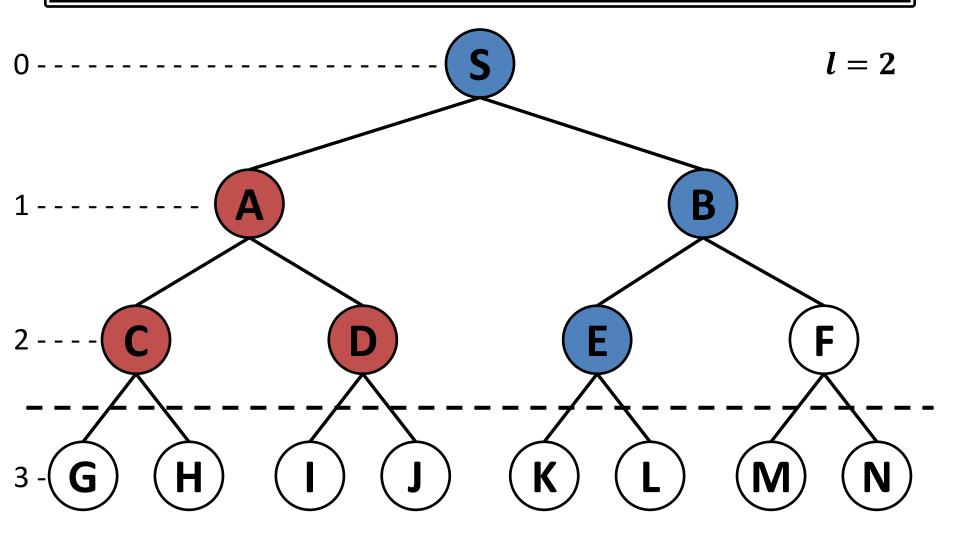


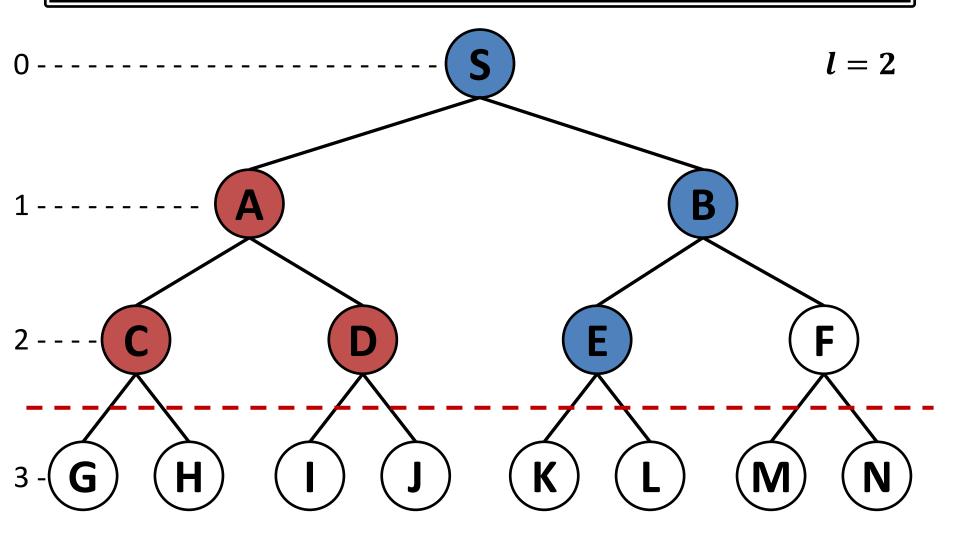


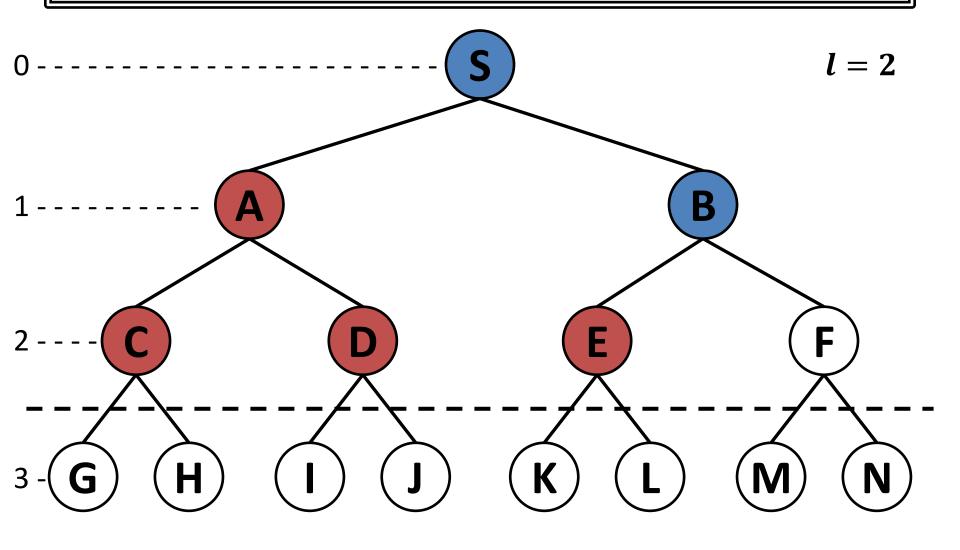


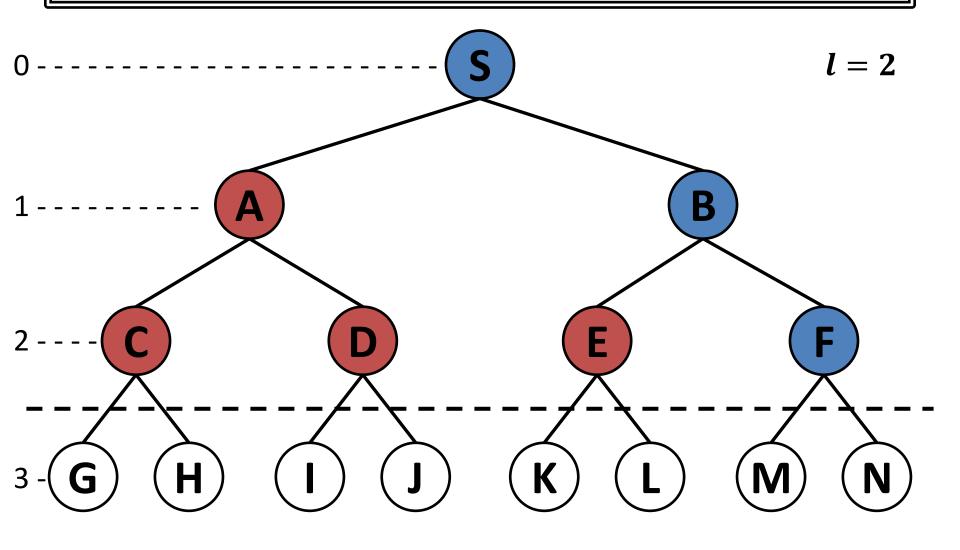


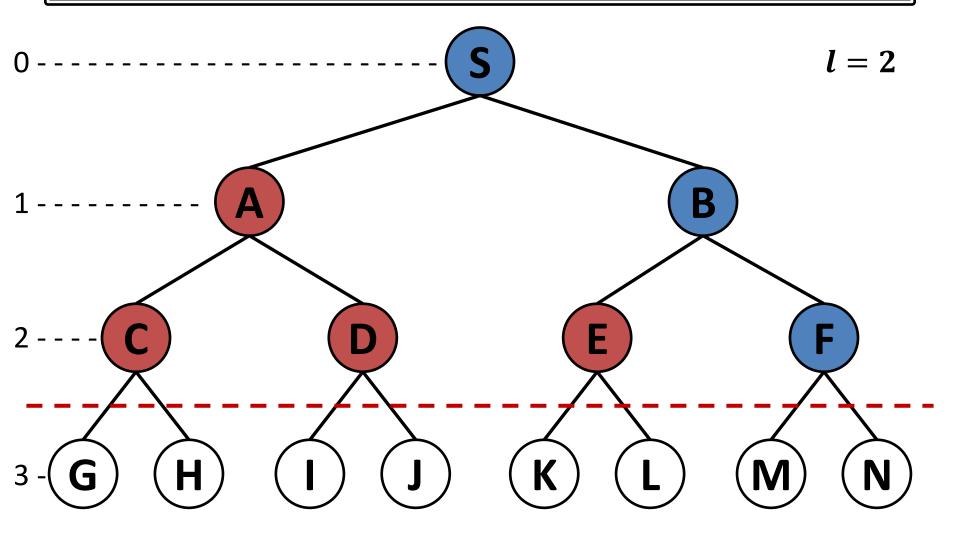


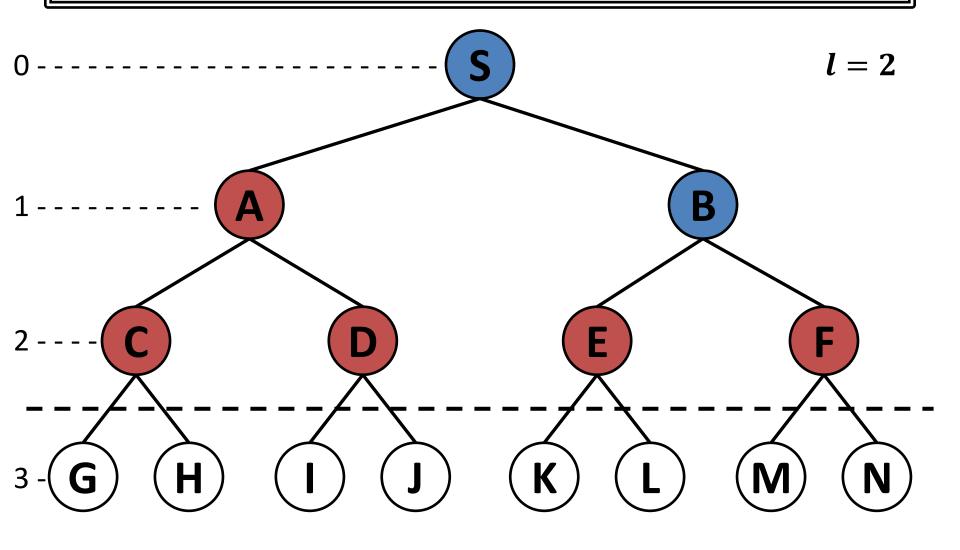


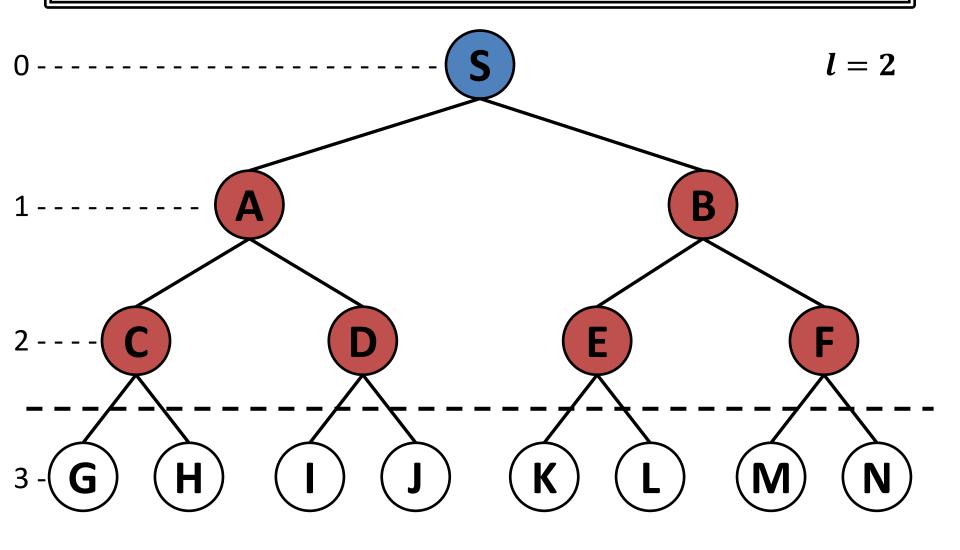


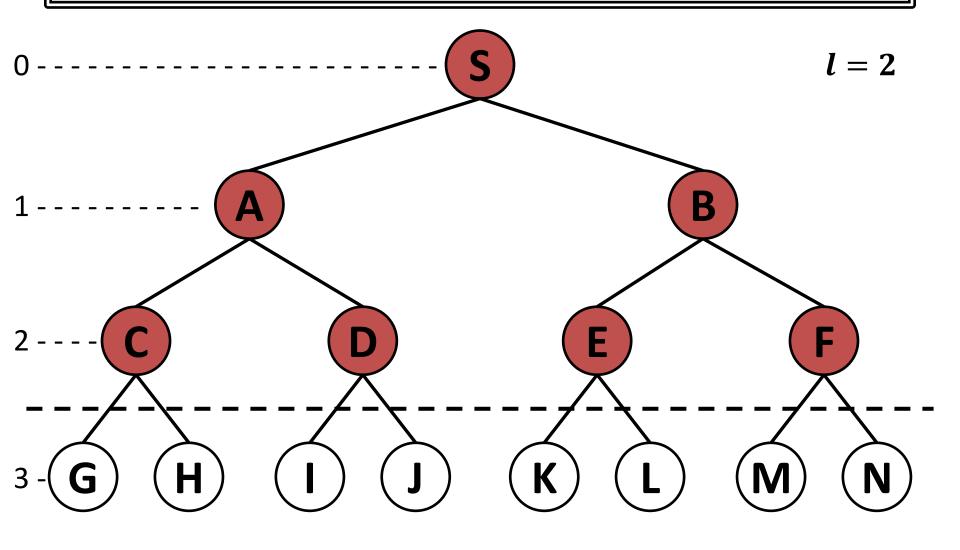












#### Depth-Limited Search (DLS)

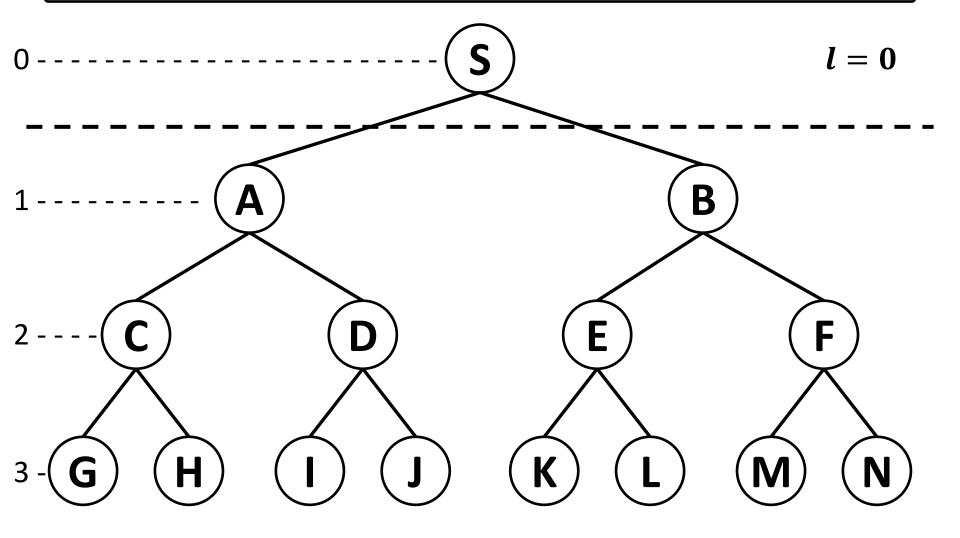
"DLS = depth-first search with depth limit I, i.e., nodes at depth I have no successors"

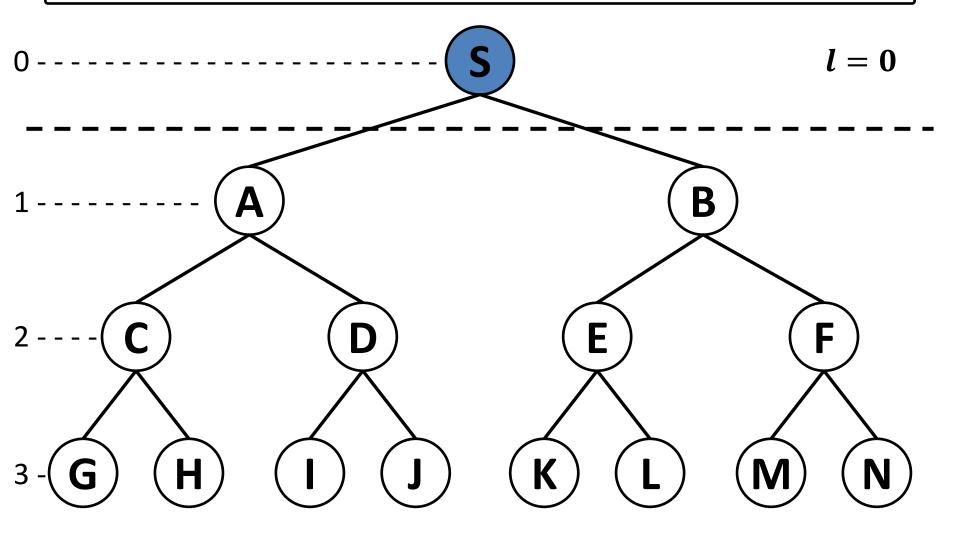
- Complete?
  - Yes, if l ≥ d
- Optimal?
  - No
- Time?
  - -O(b')
- Space?
  - O(bl)

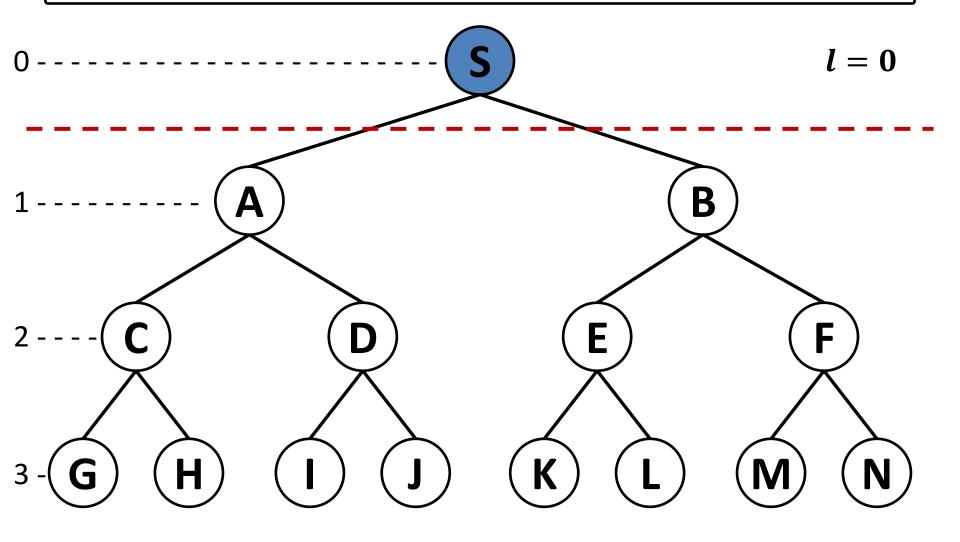
<sup>\*</sup> DLS searches in a subspace (local area), but it is NOT a local search algorithm!

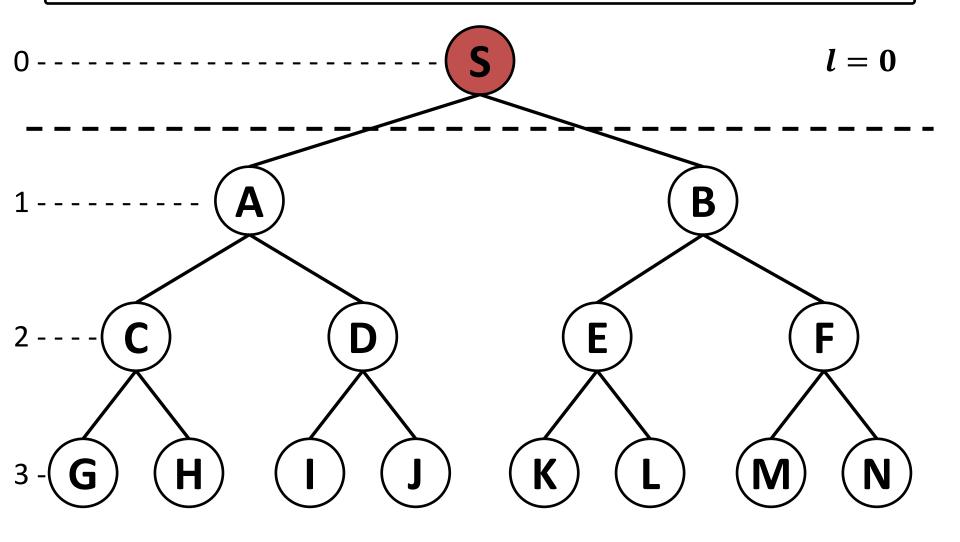
#### Iterative Deepening Search (IDS)

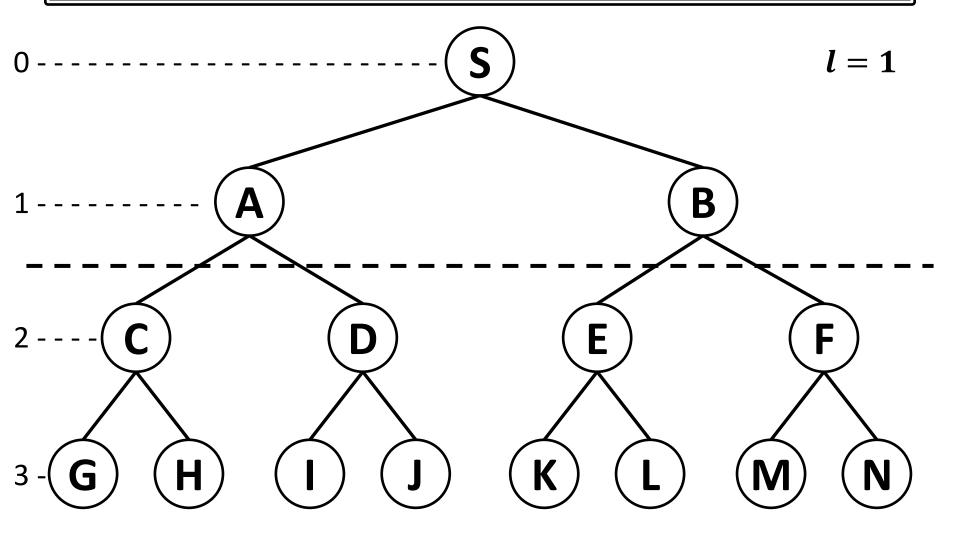
```
function ITERATIVE-DEEPENING-SEARCH (problem) returns a solution inputs: problem, a problem for depth \leftarrow 0 to \infty do result \leftarrow Depth-Limited-Search (problem, depth) if <math>result \neq \text{cutoff then return } result end
```

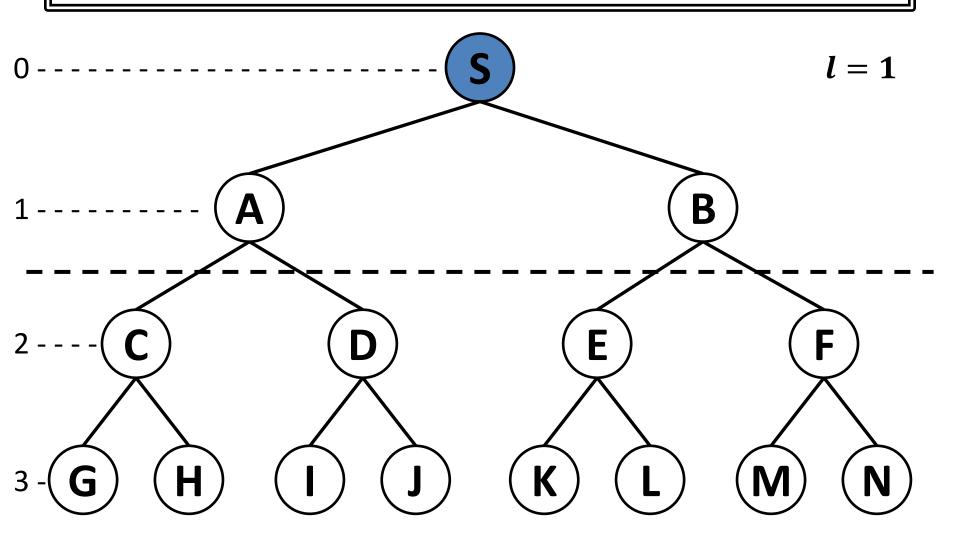


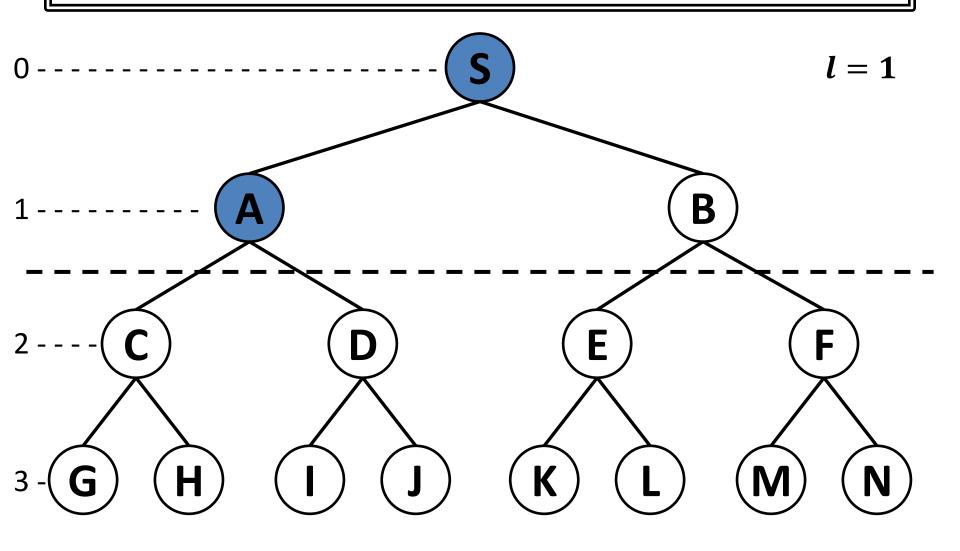


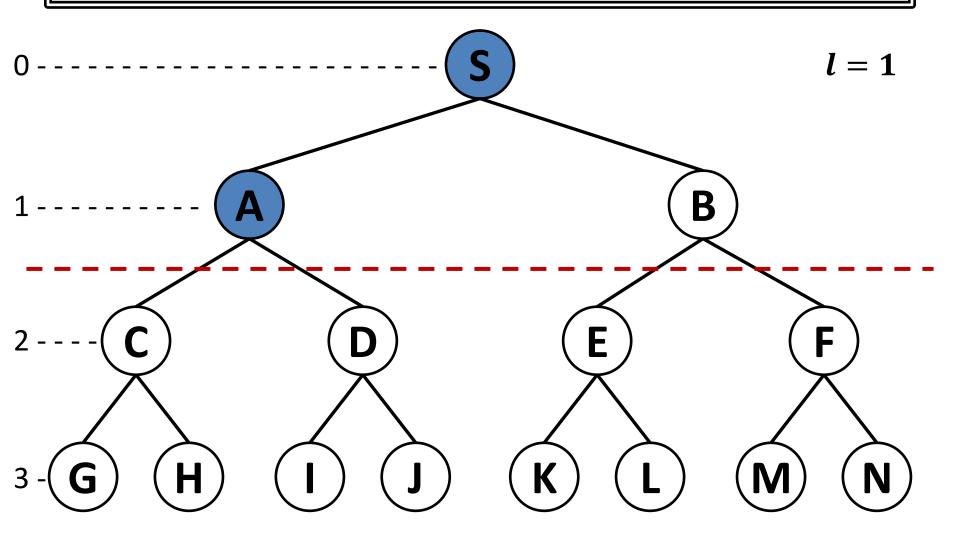


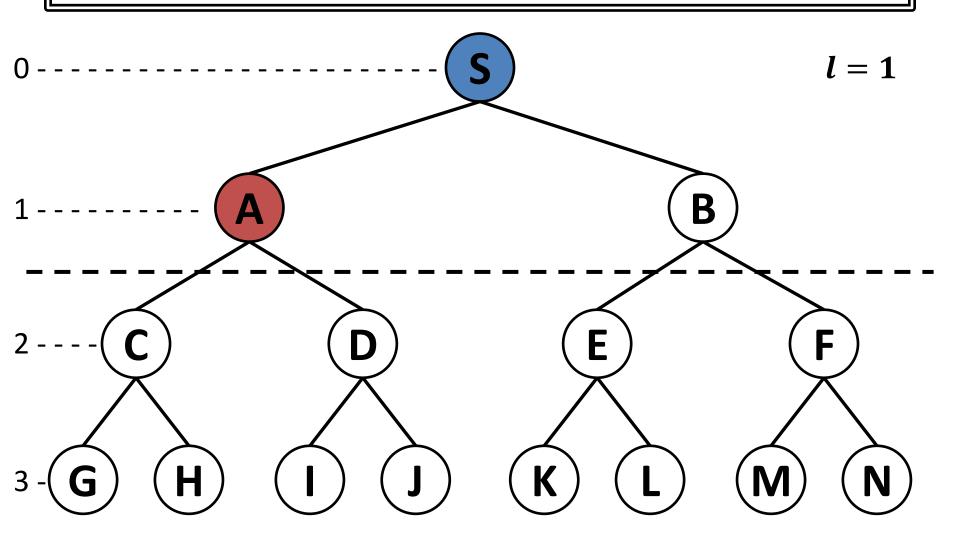


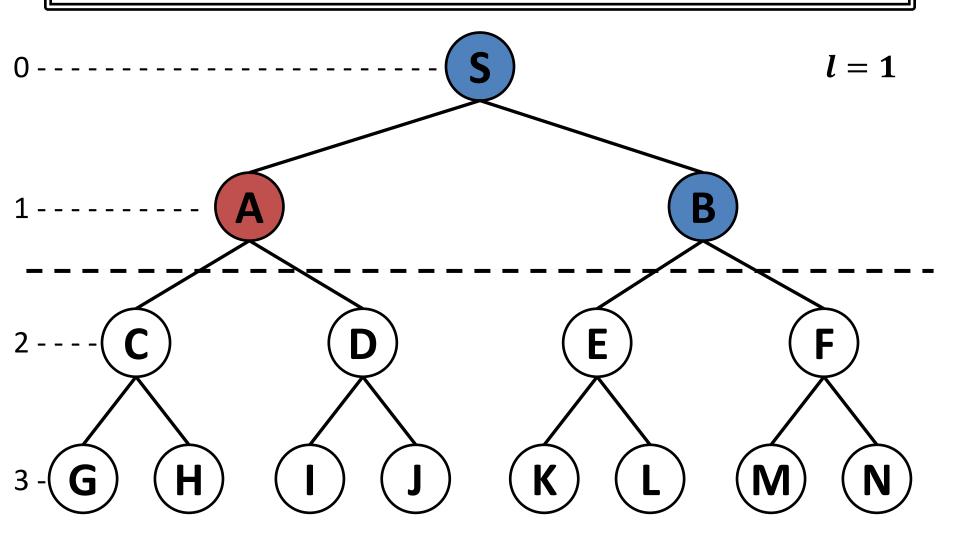


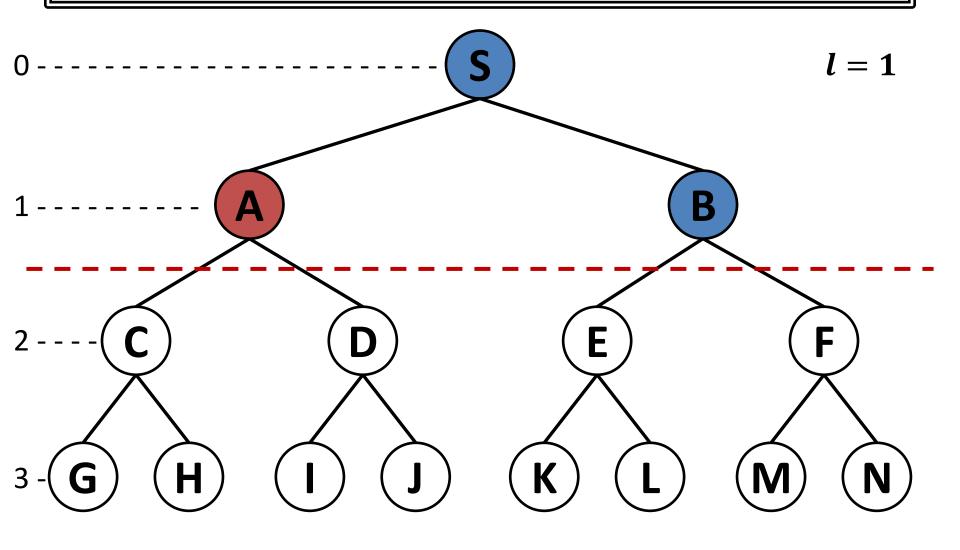


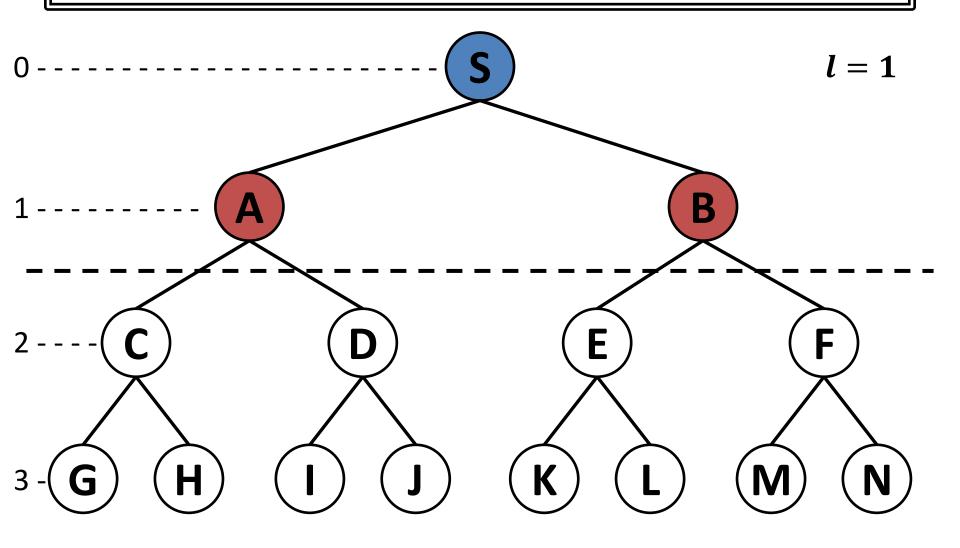


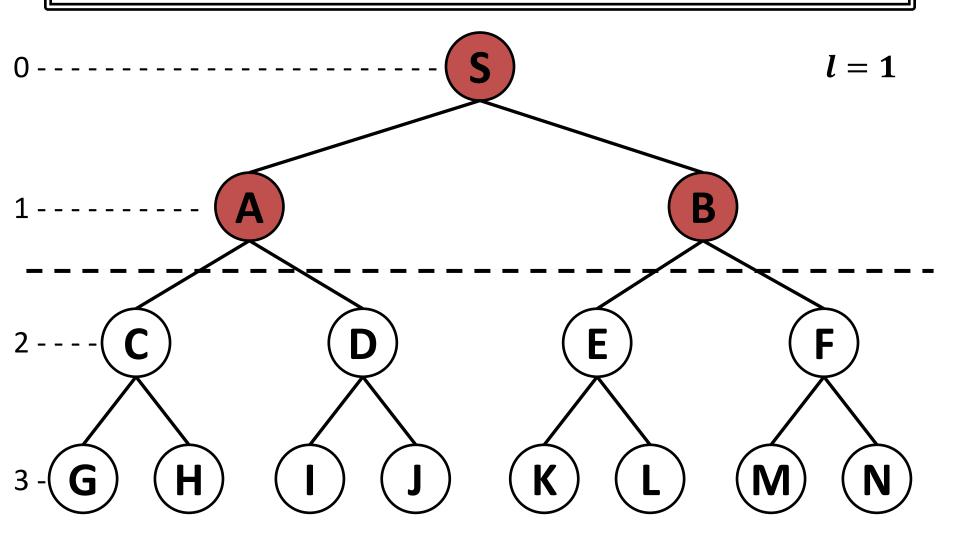


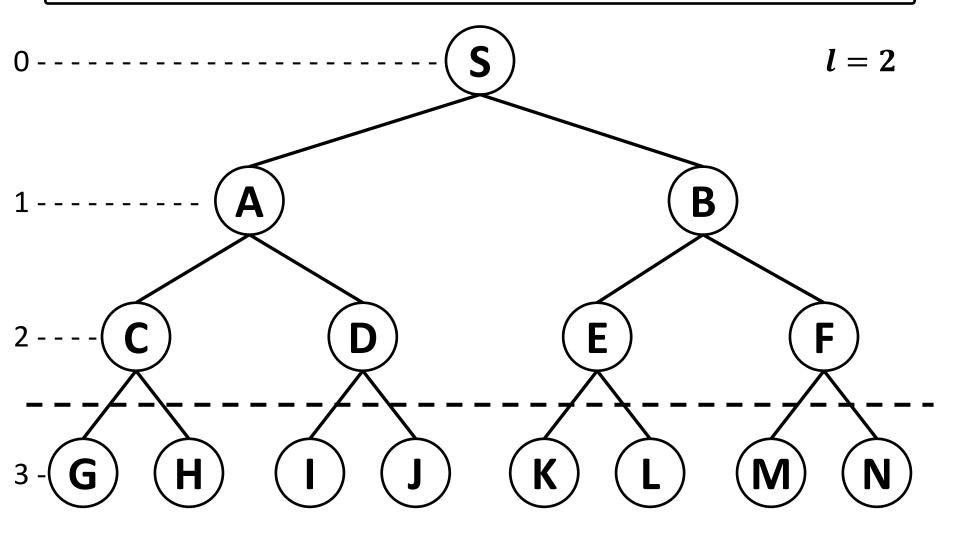


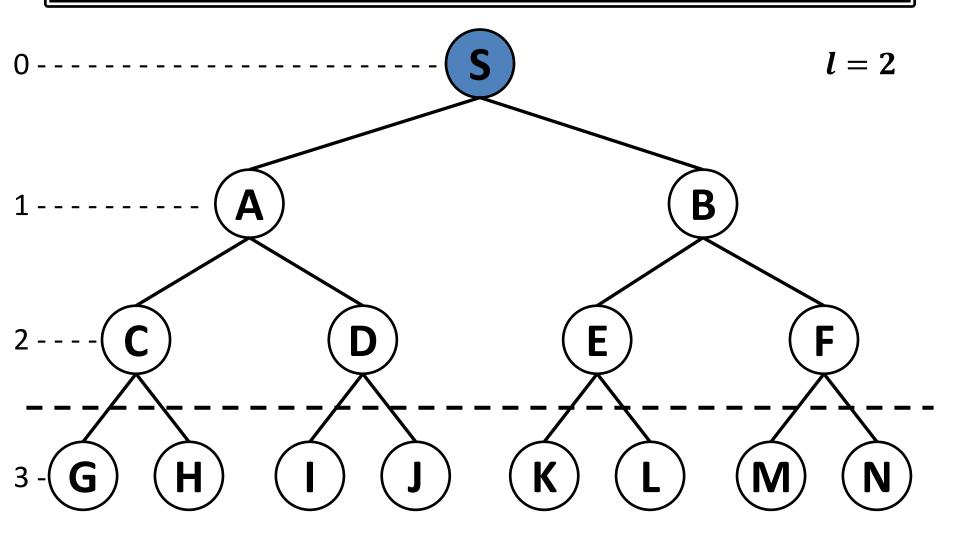


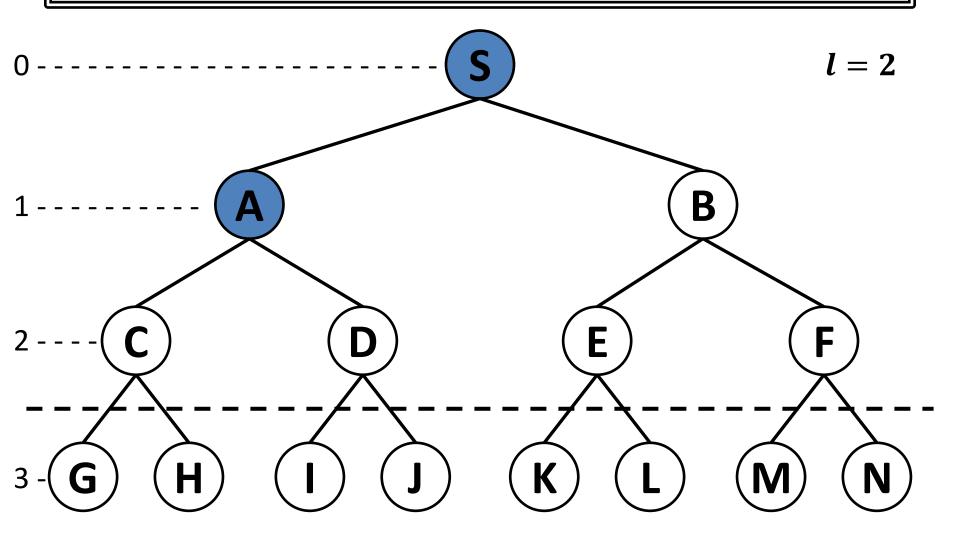


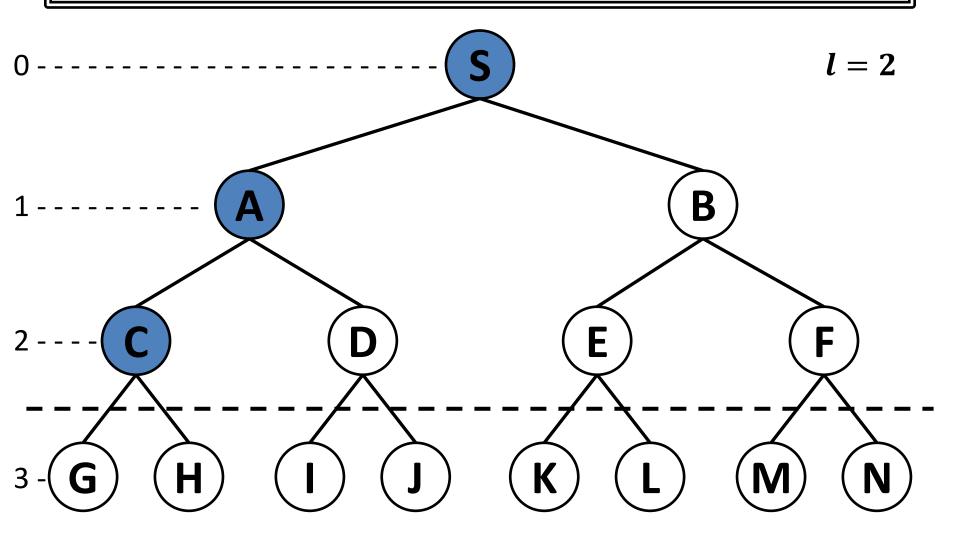


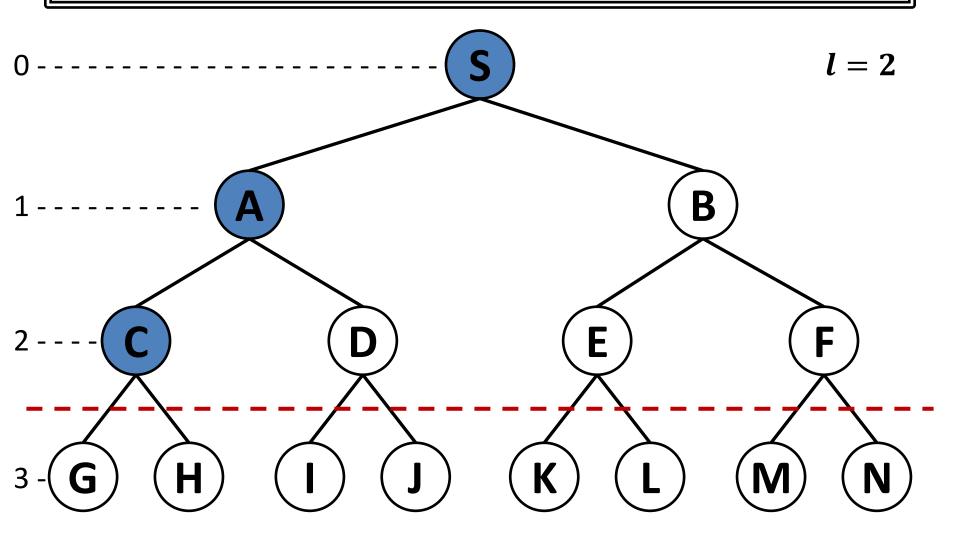


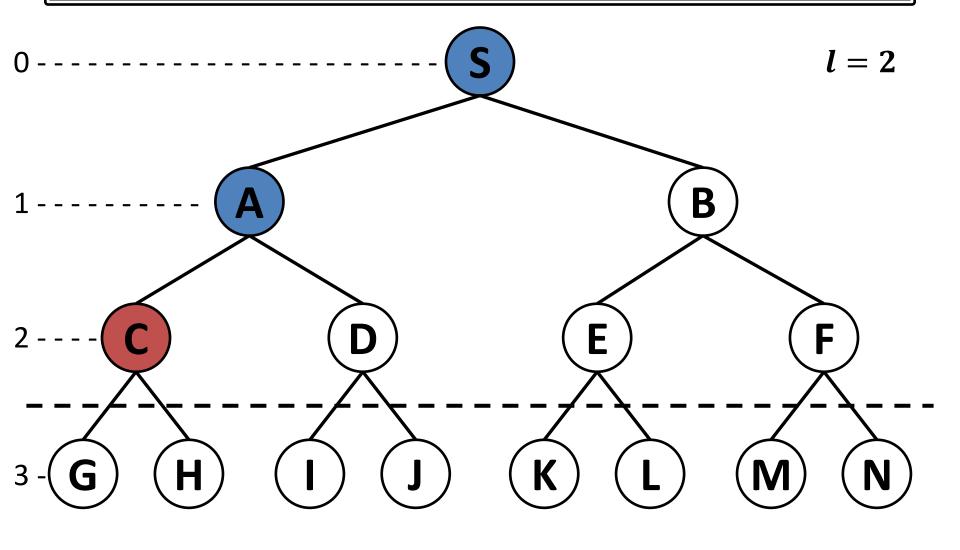


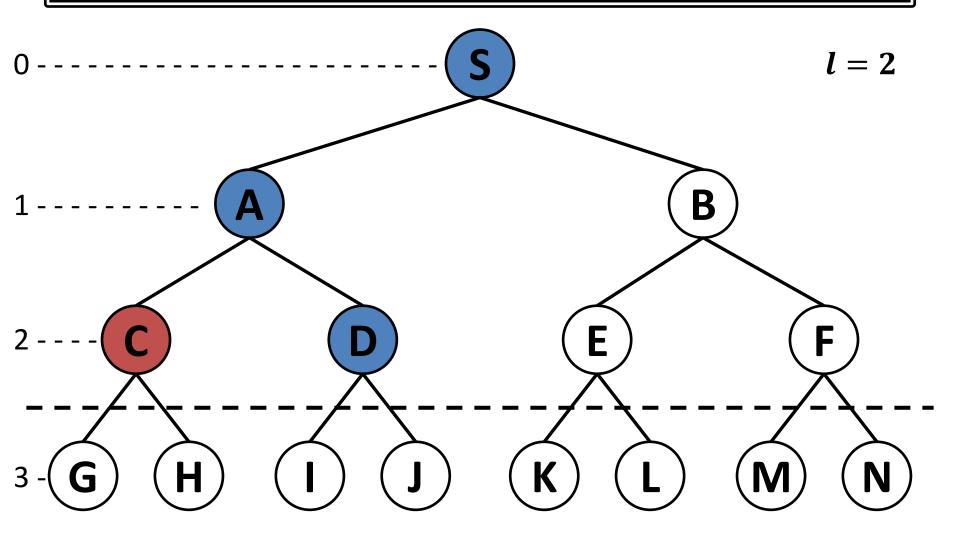


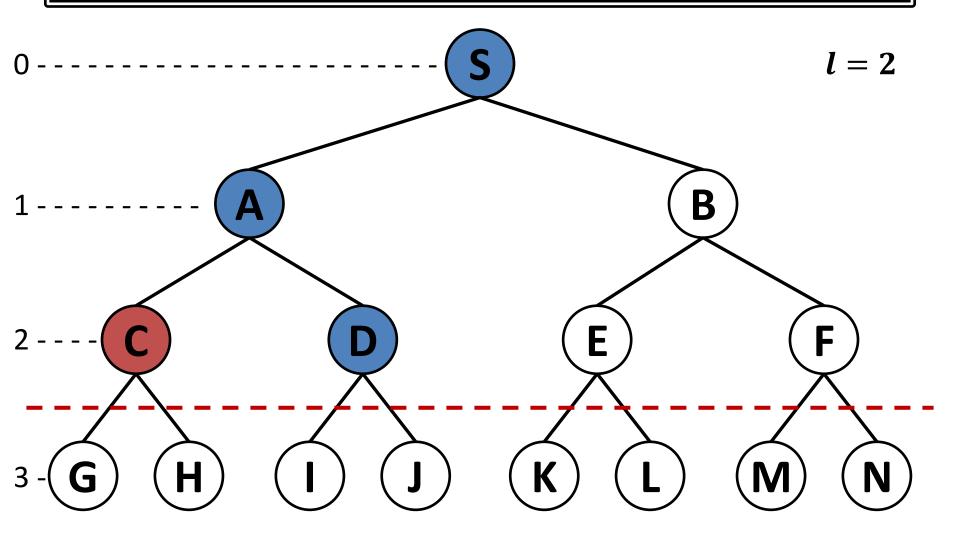


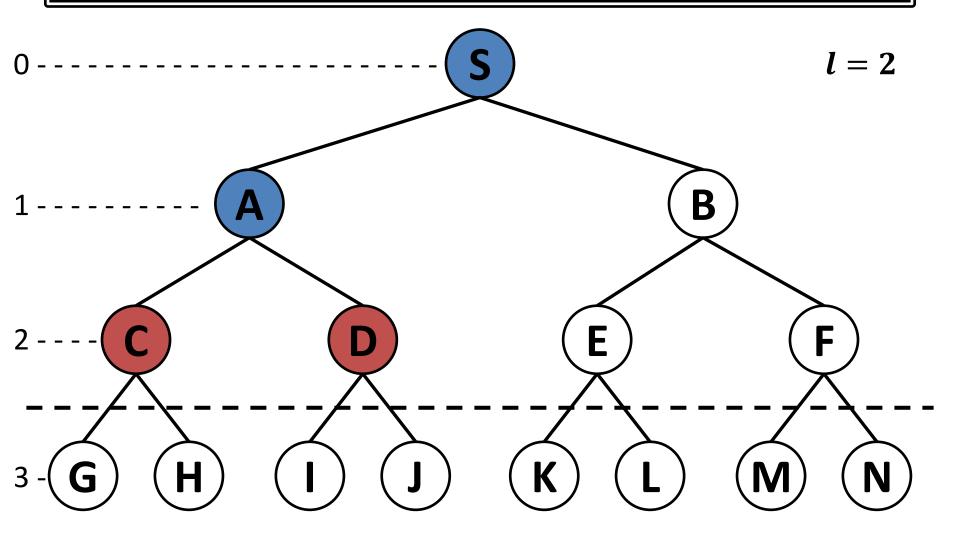


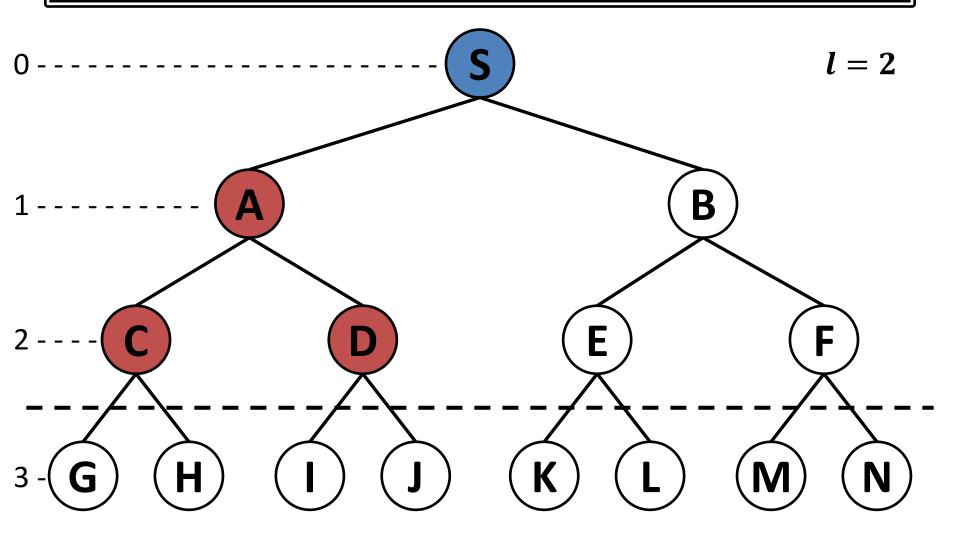


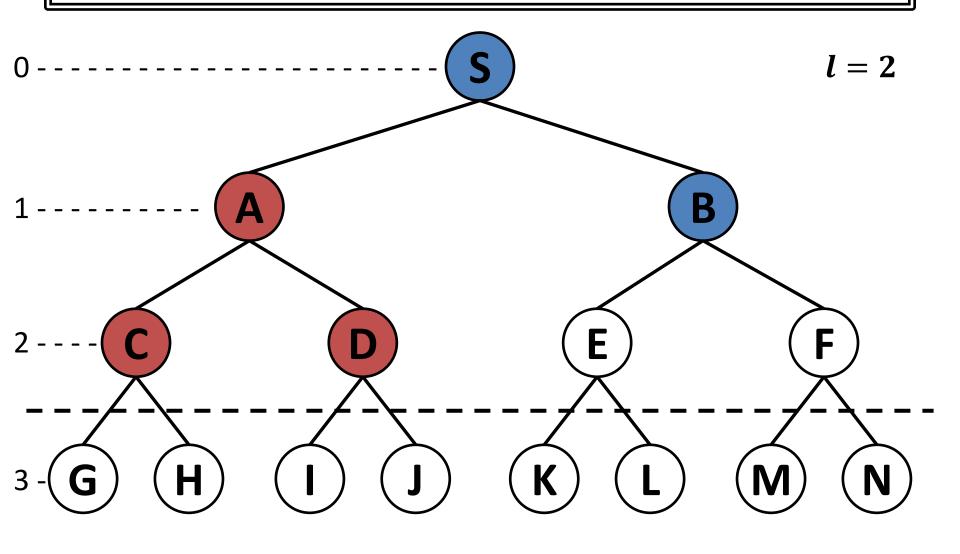


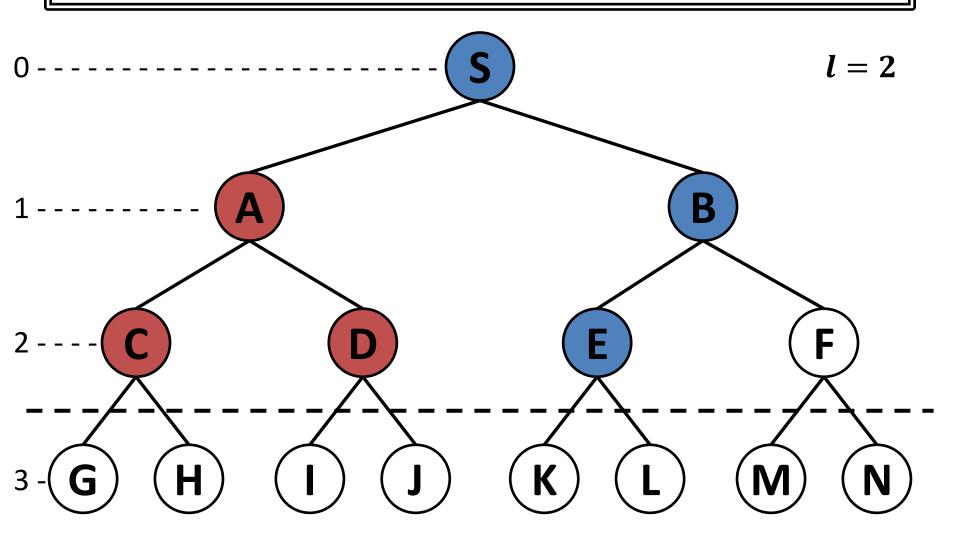


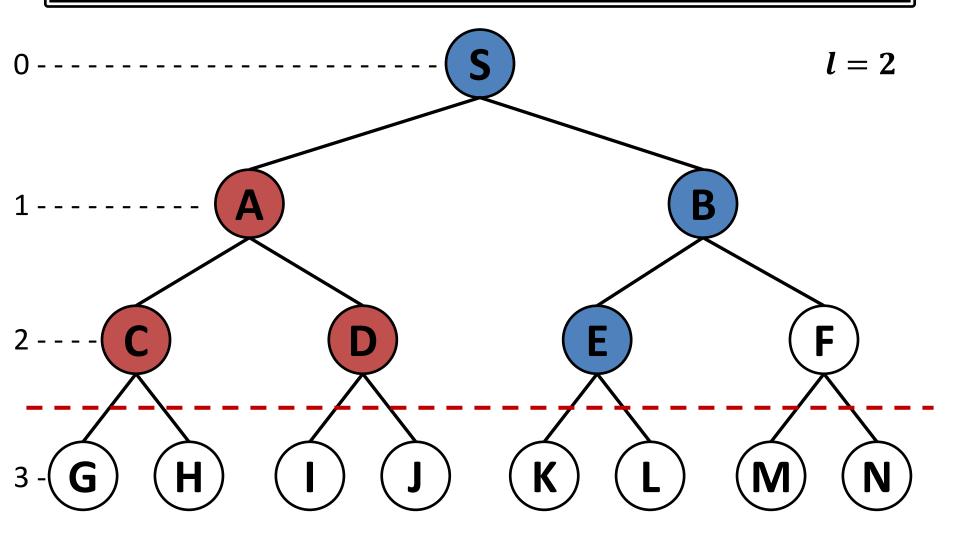


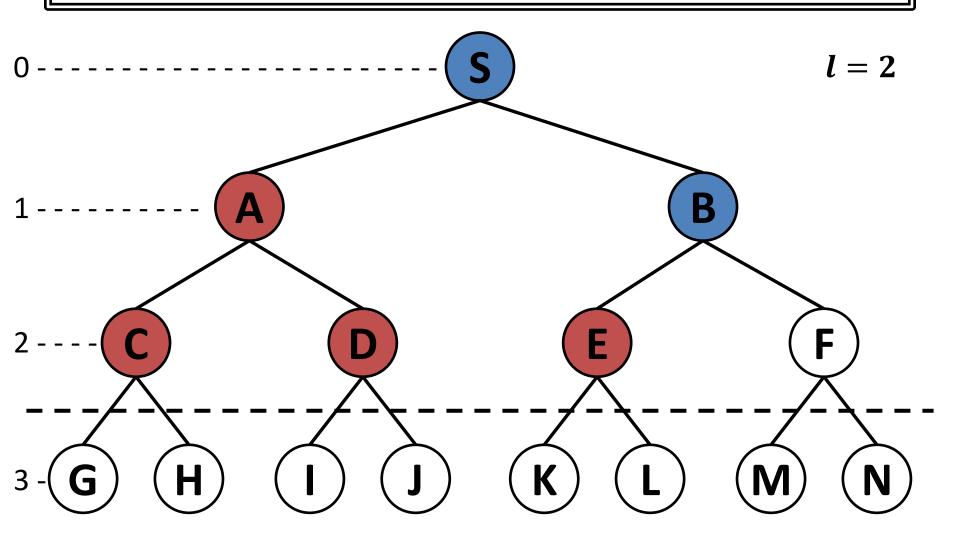


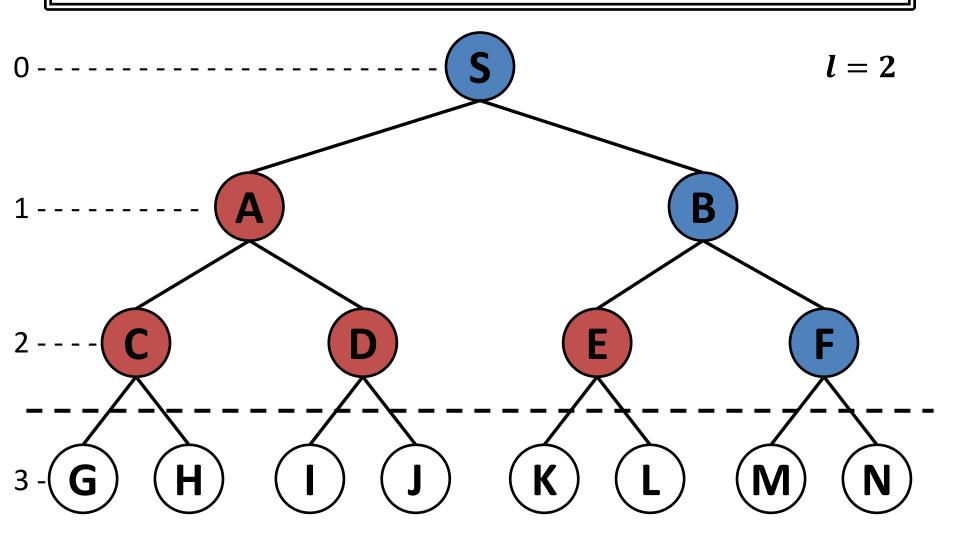


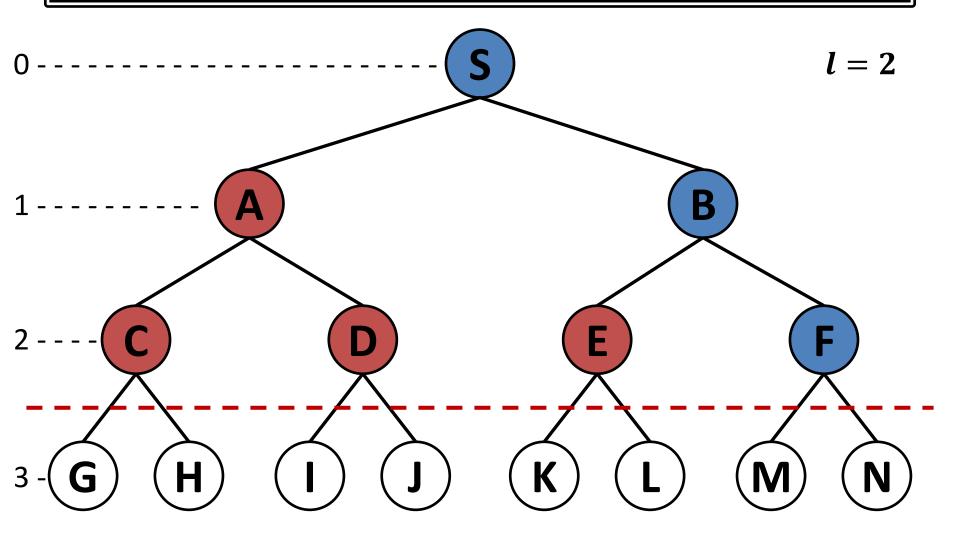


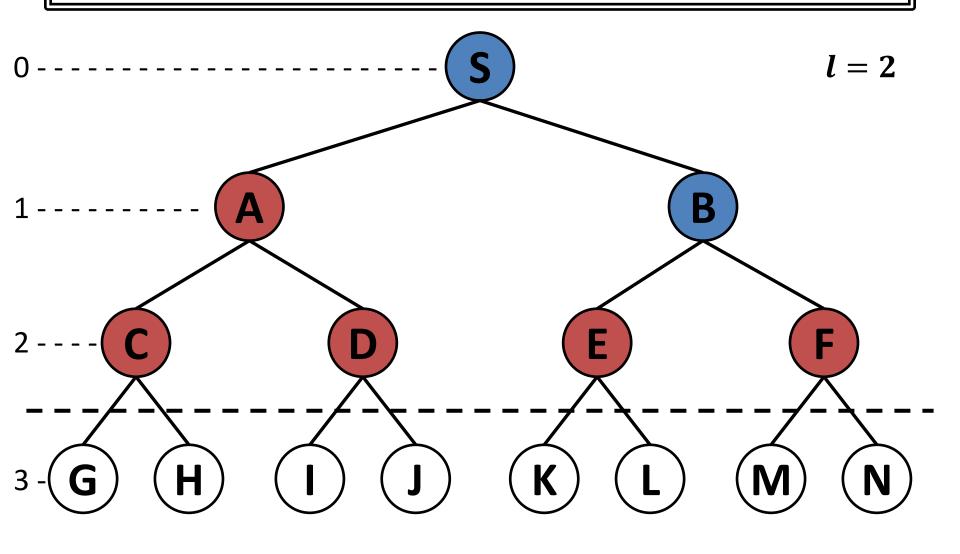


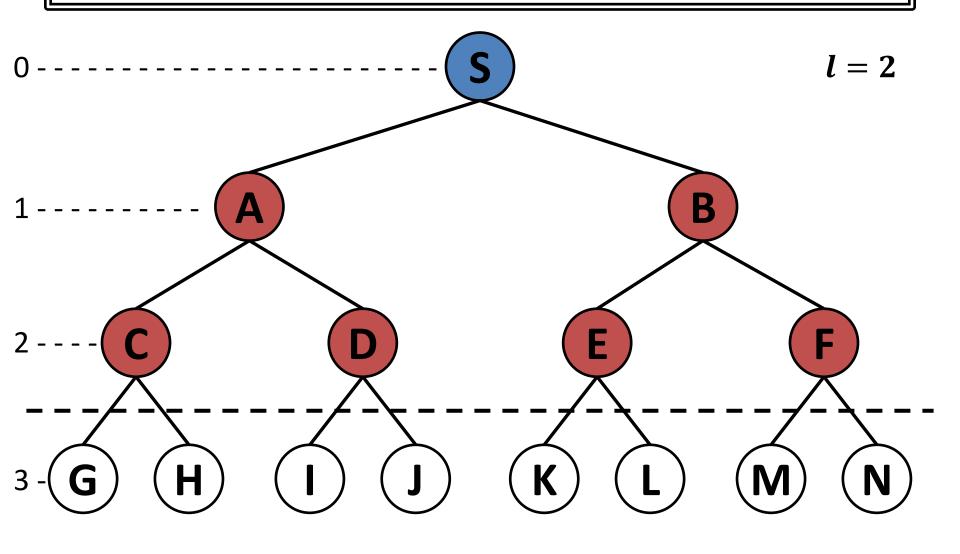


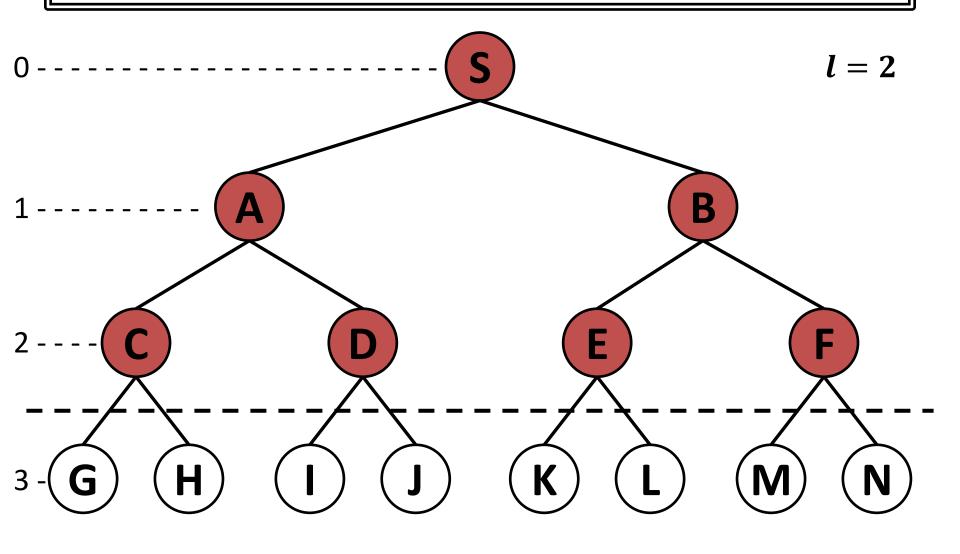


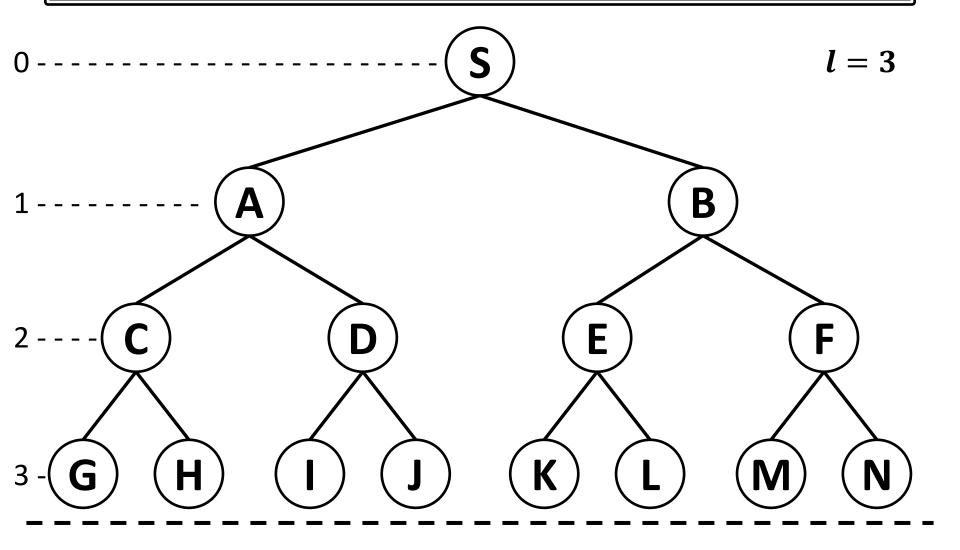


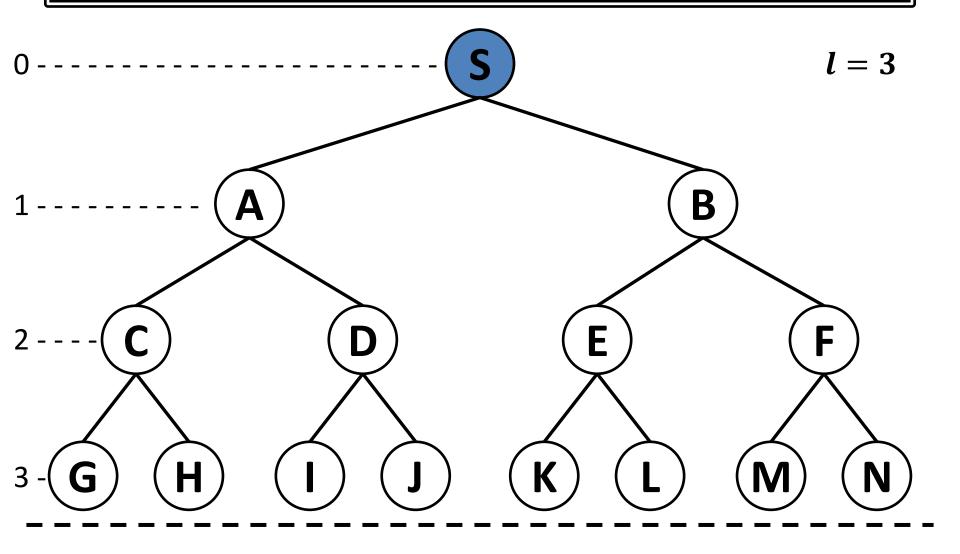


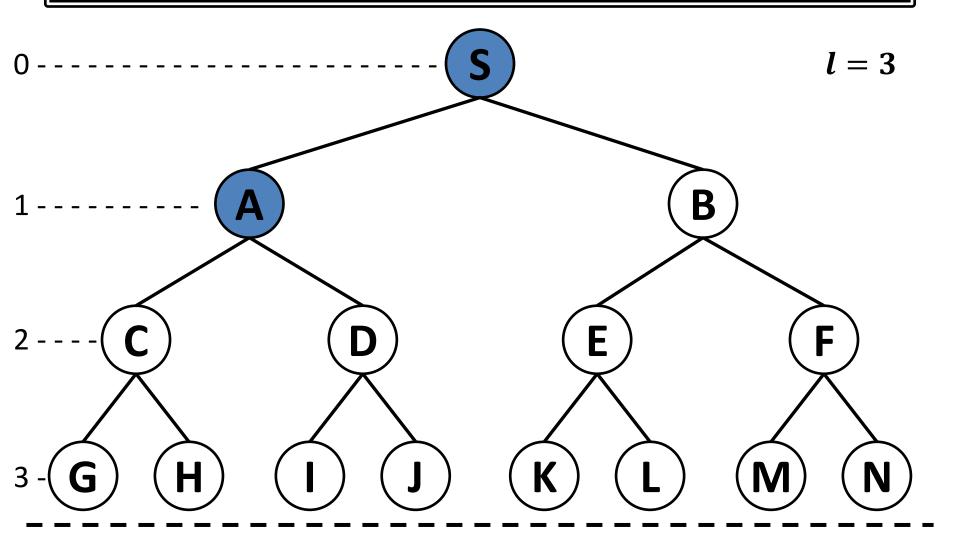


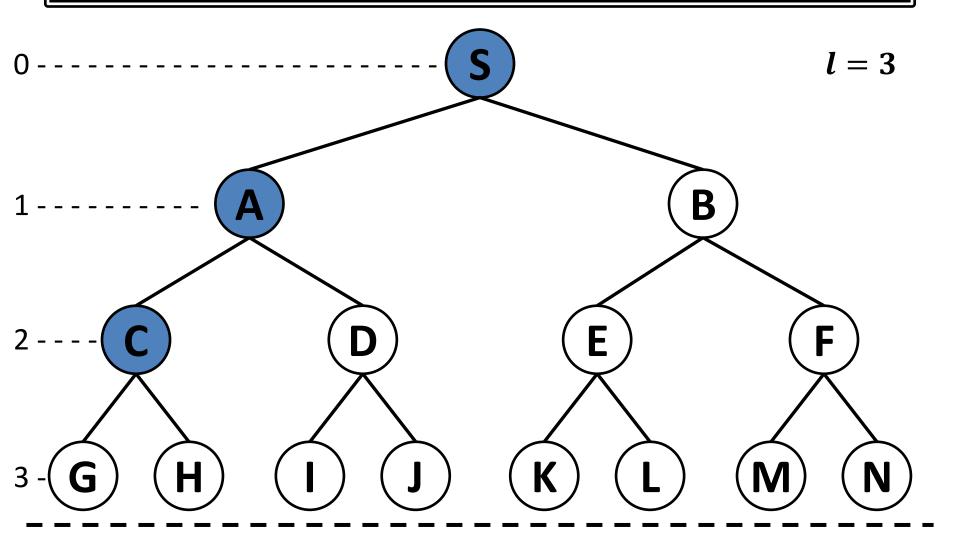


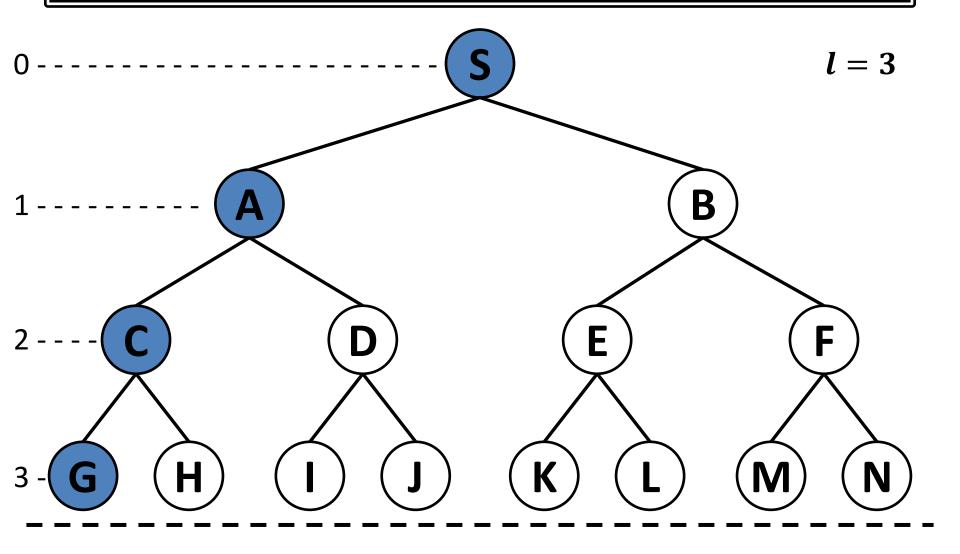


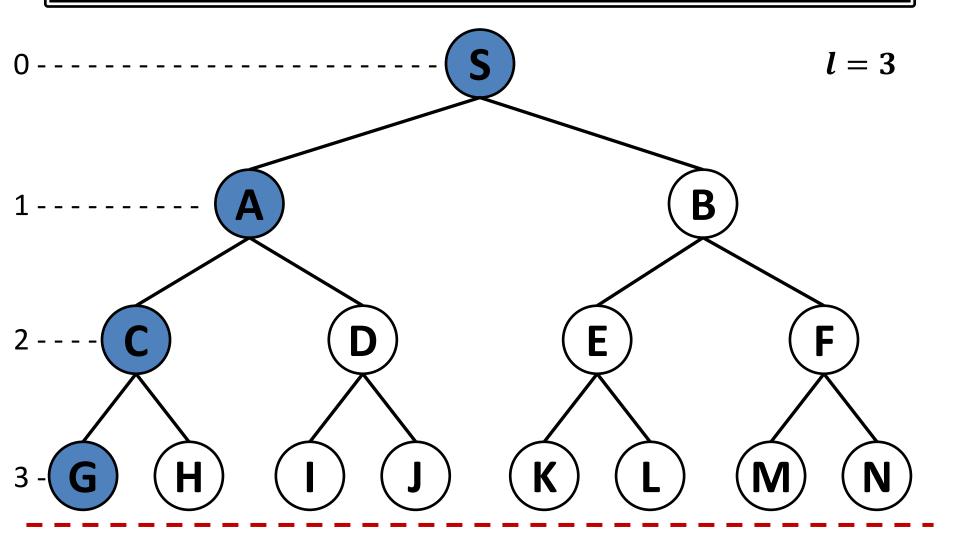


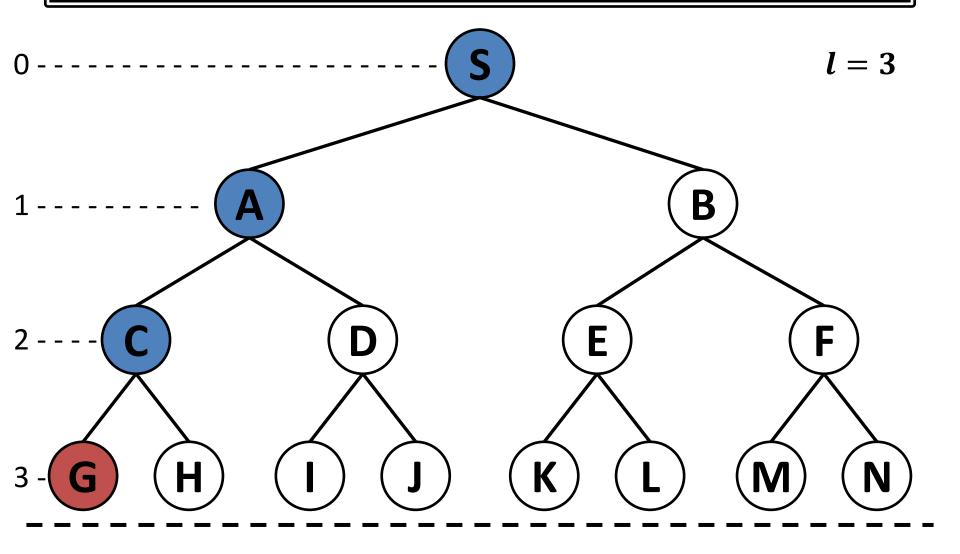


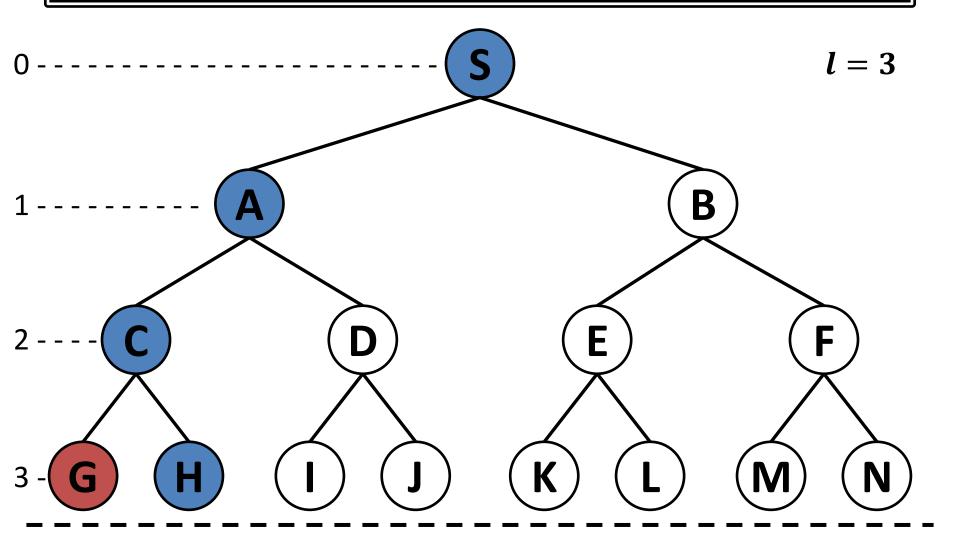


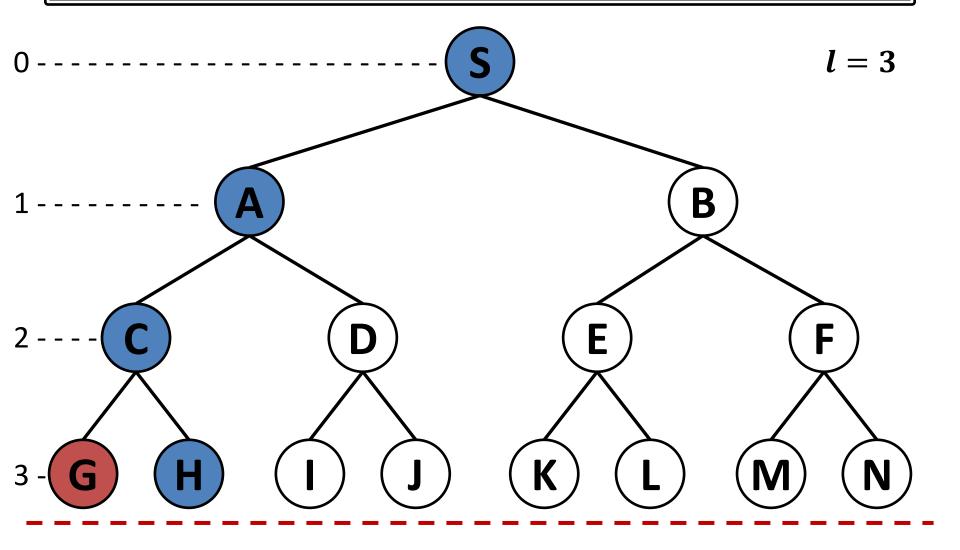


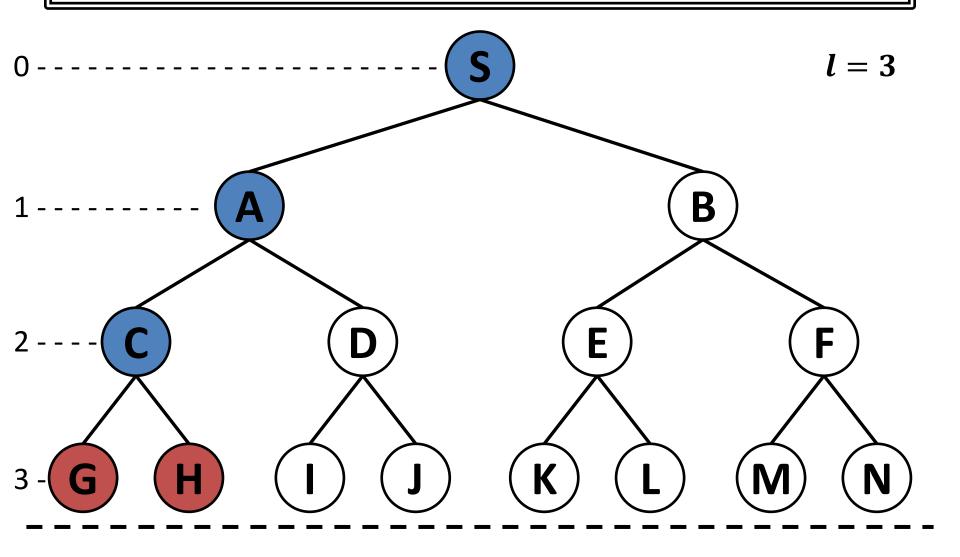


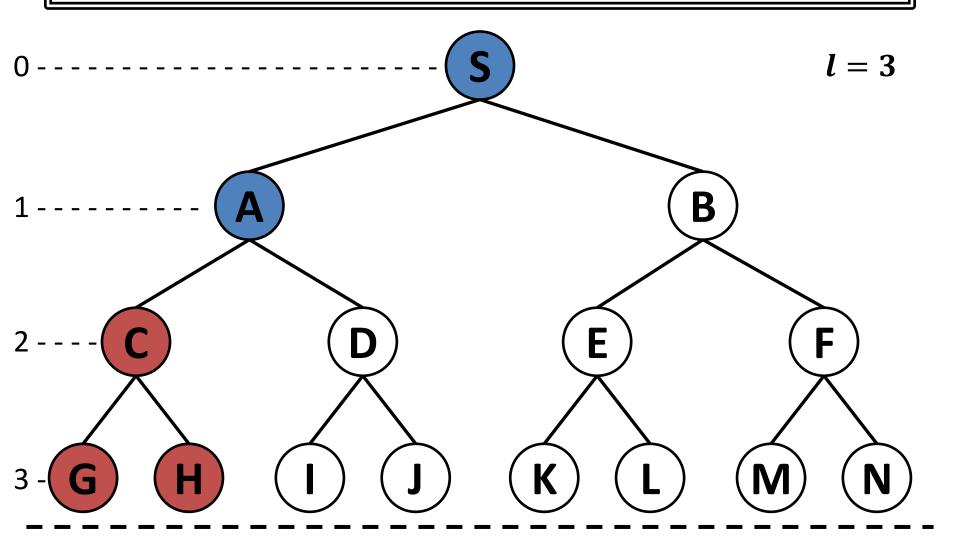


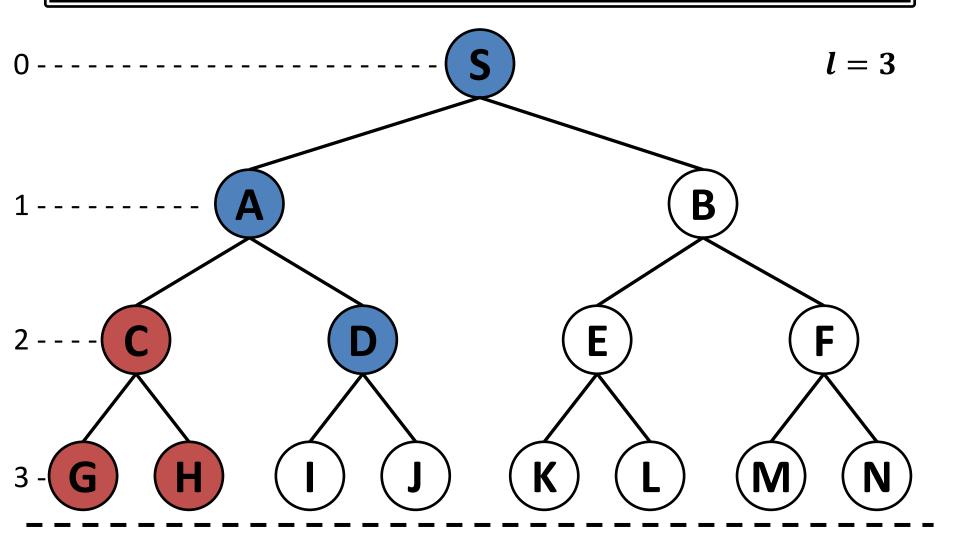


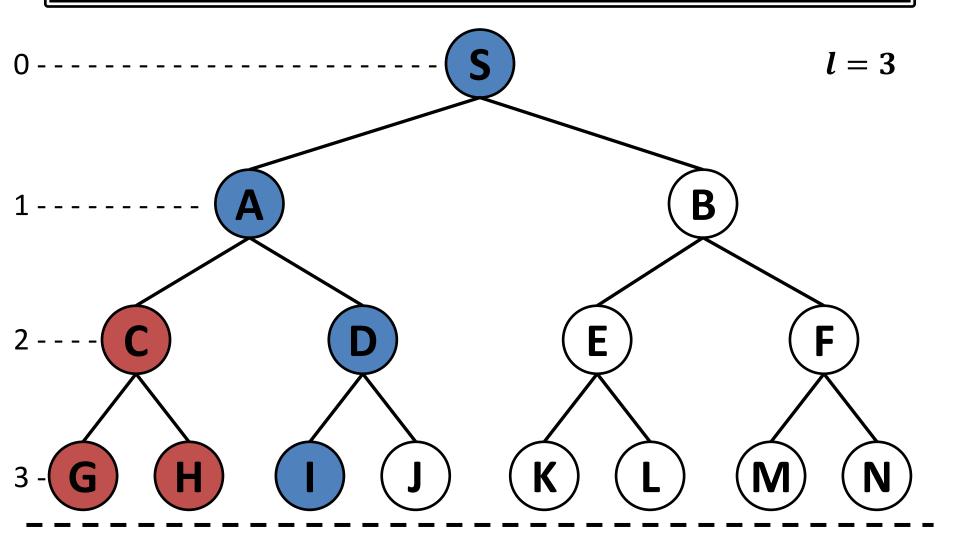


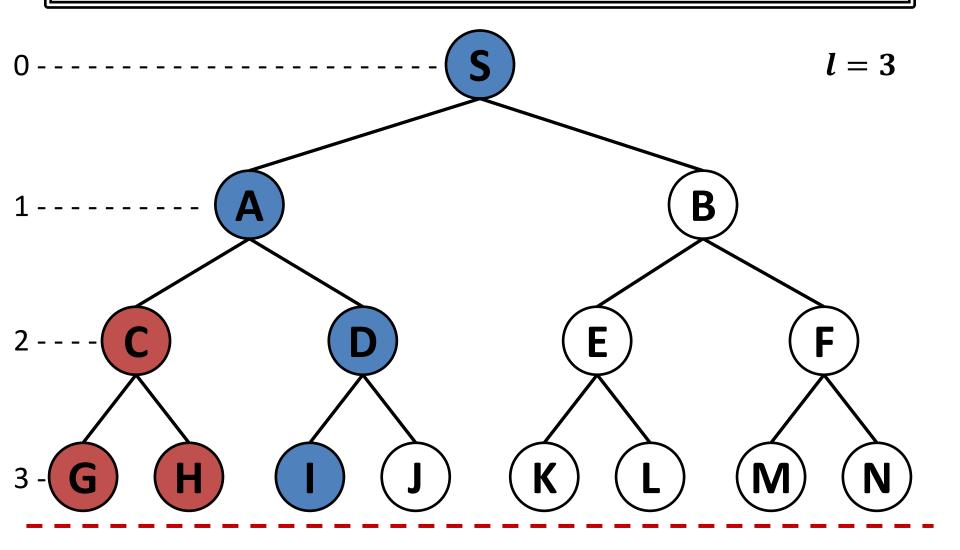


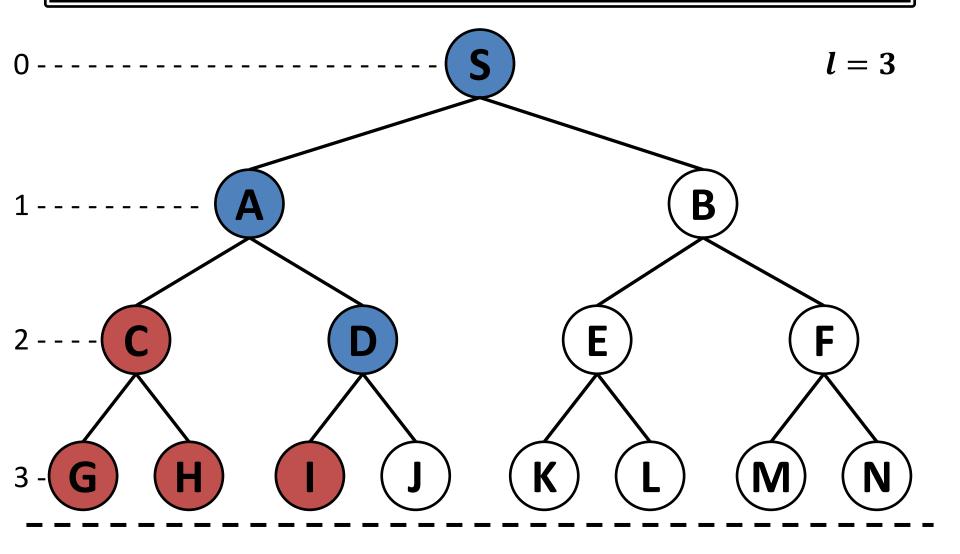


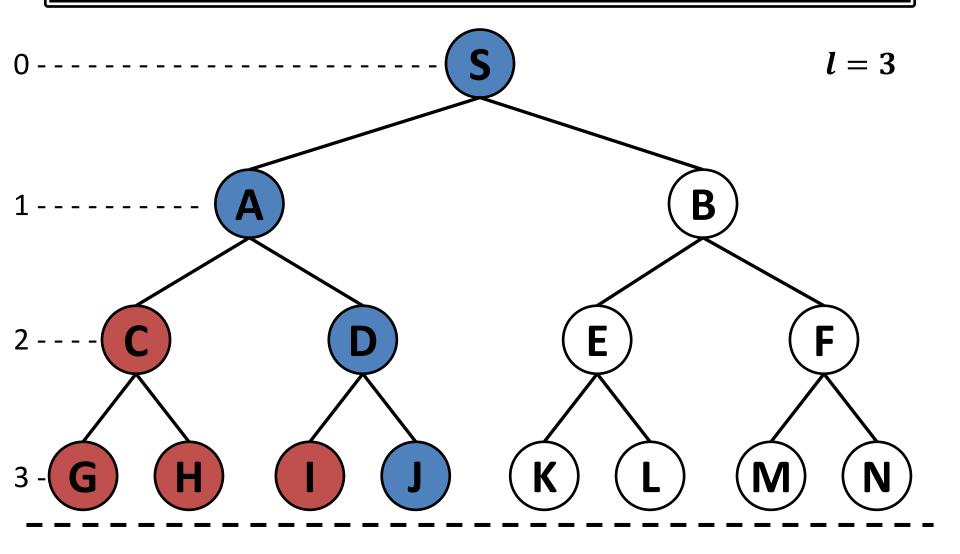


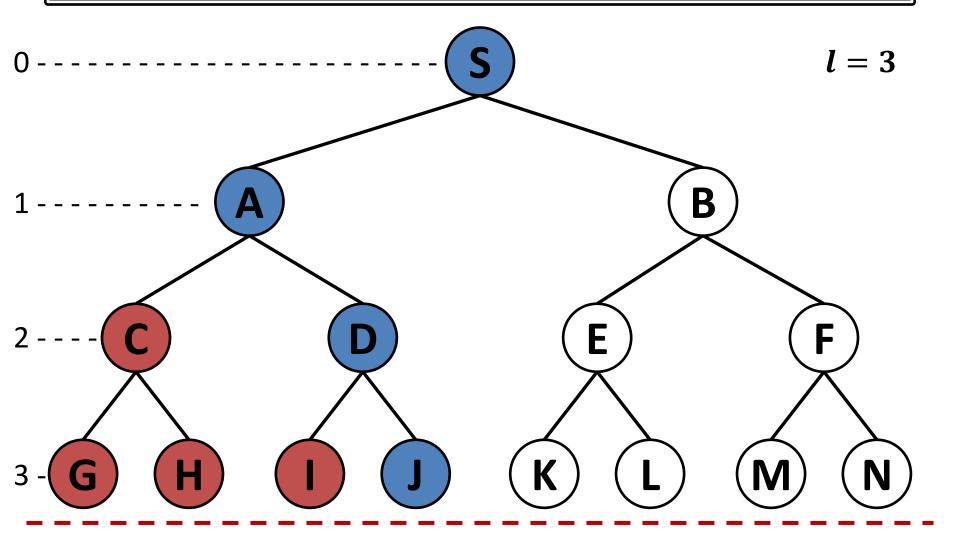


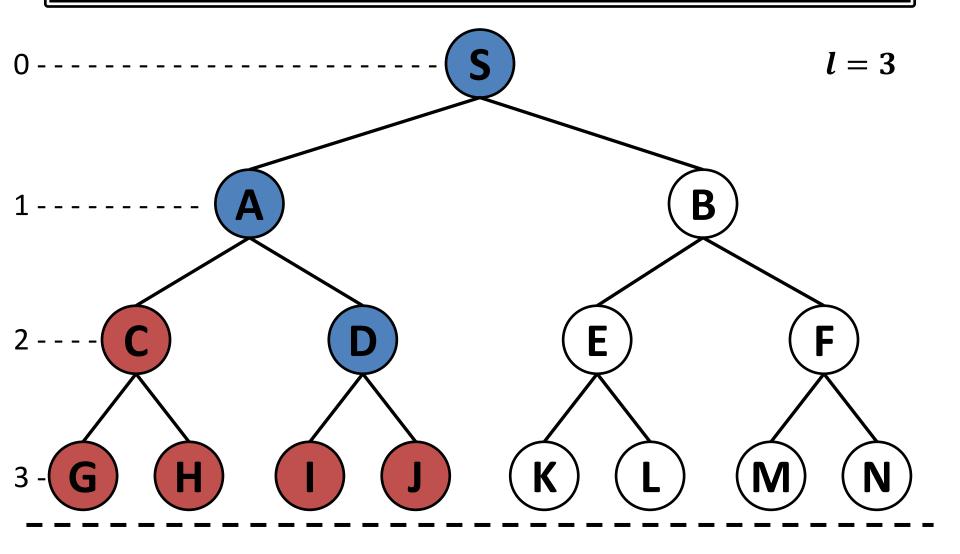


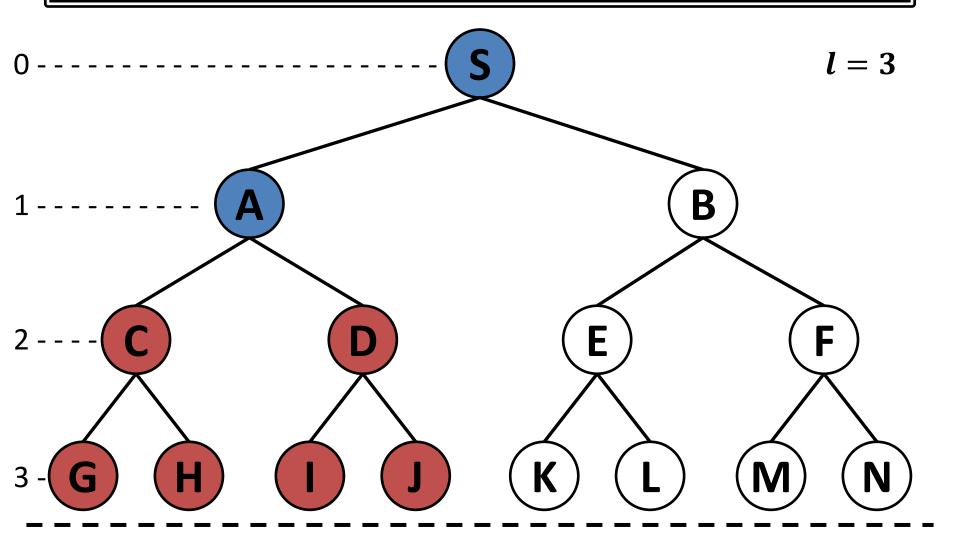


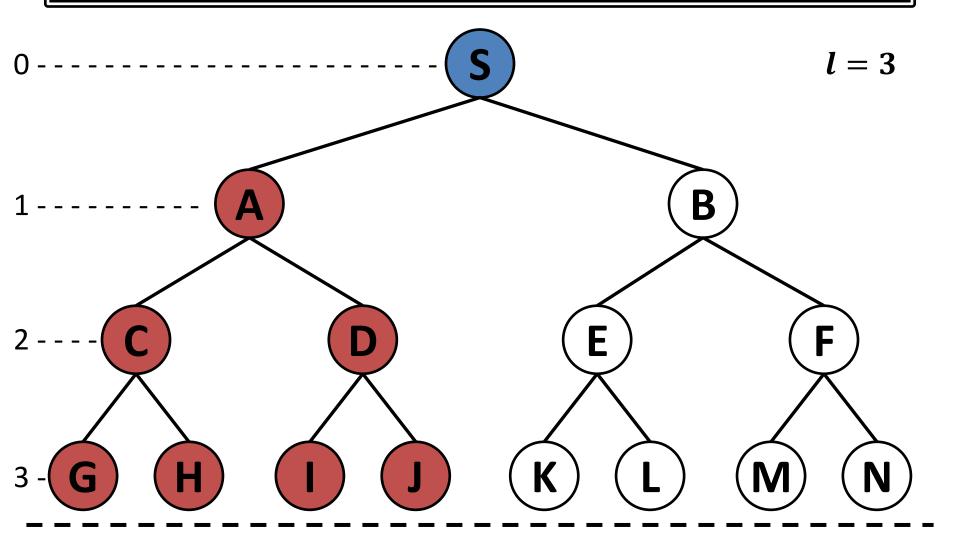


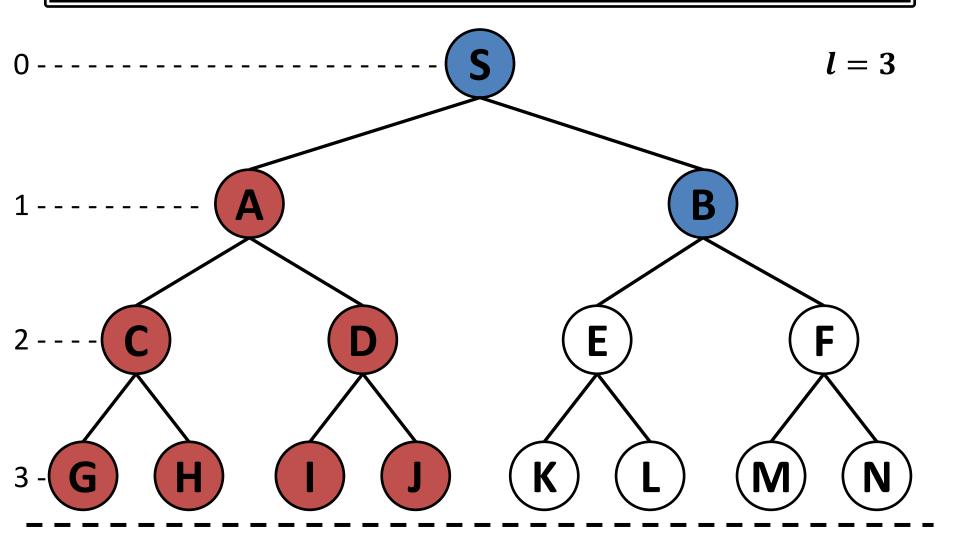


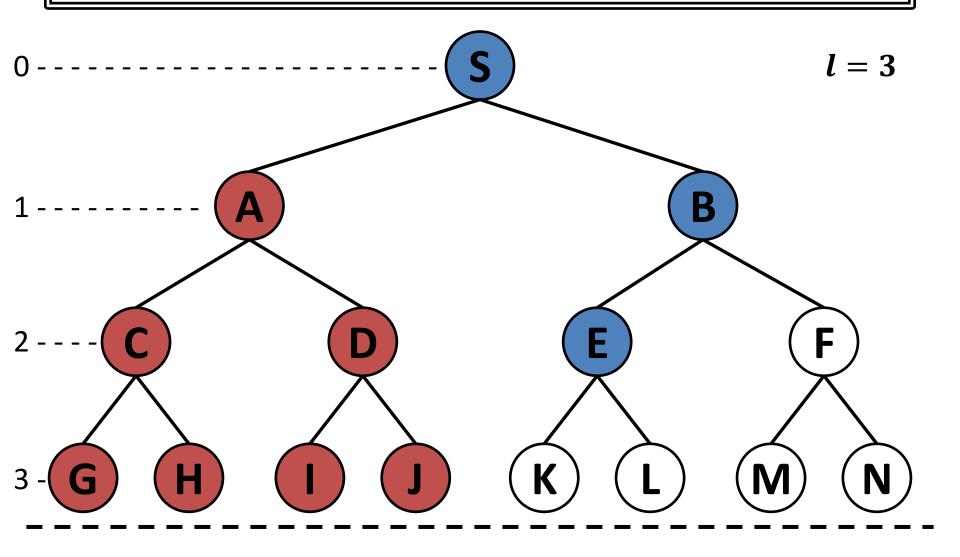


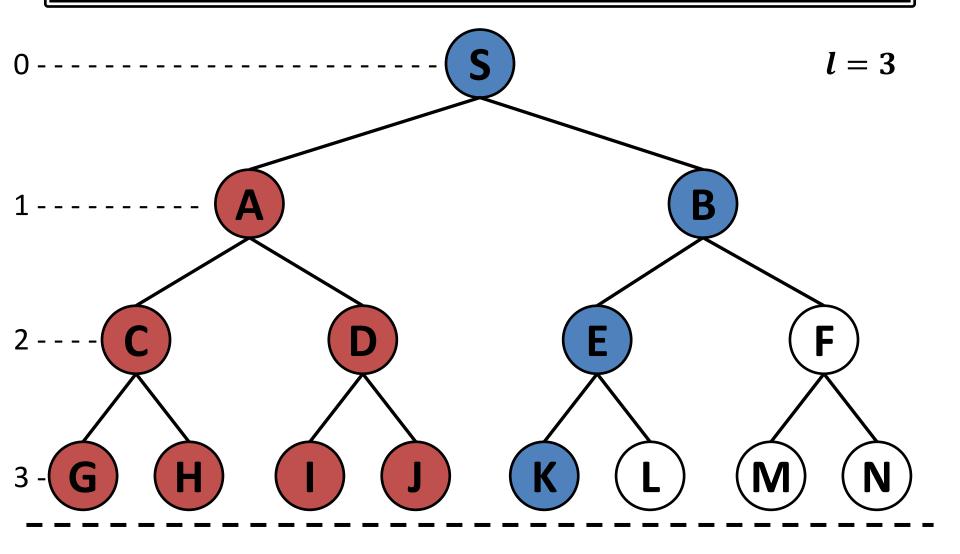


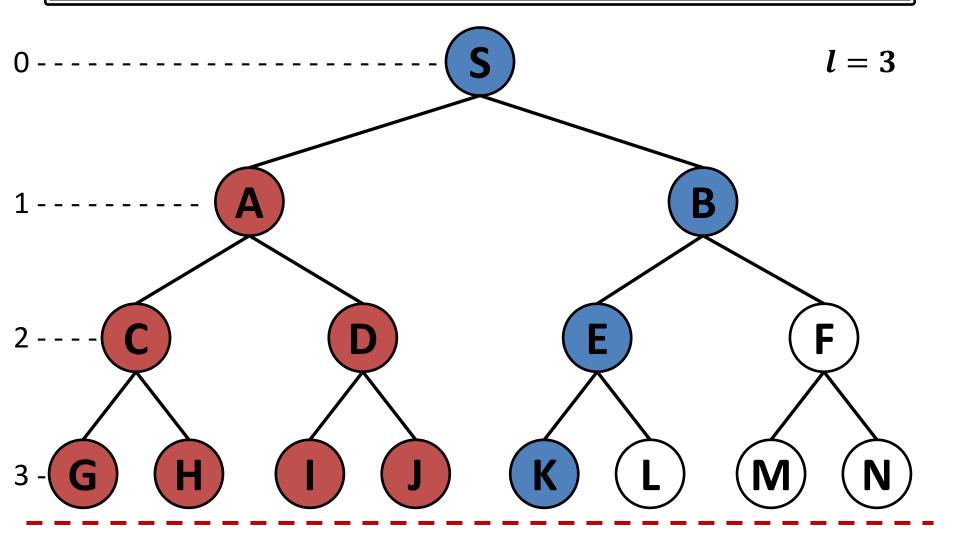


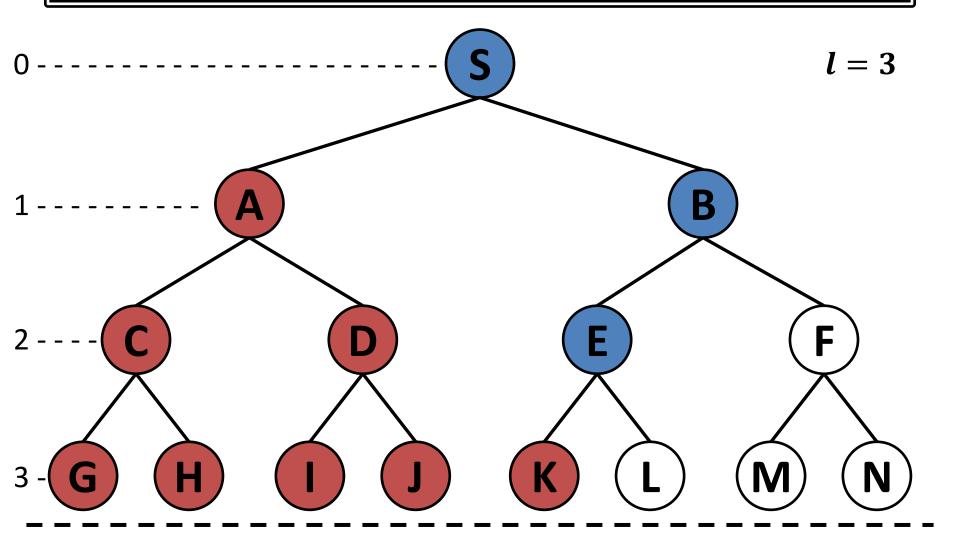


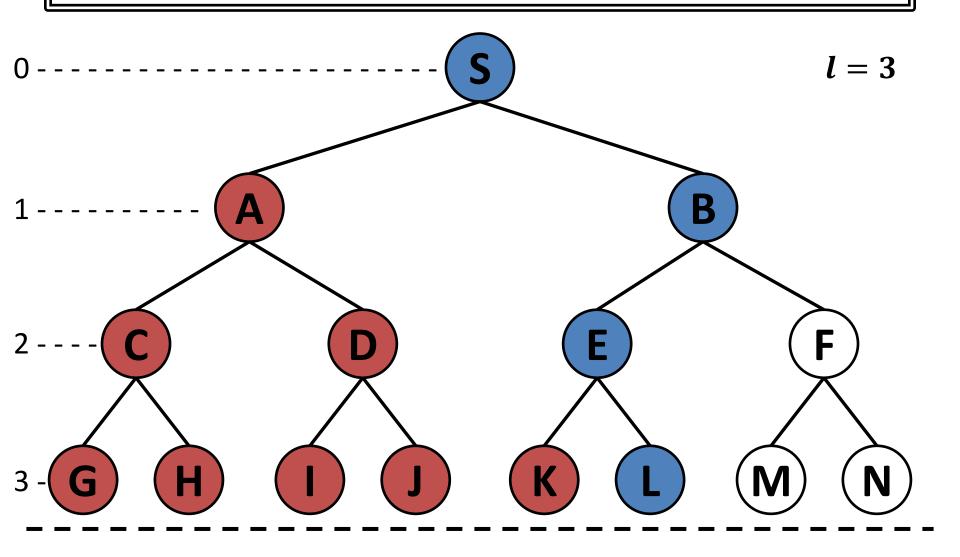


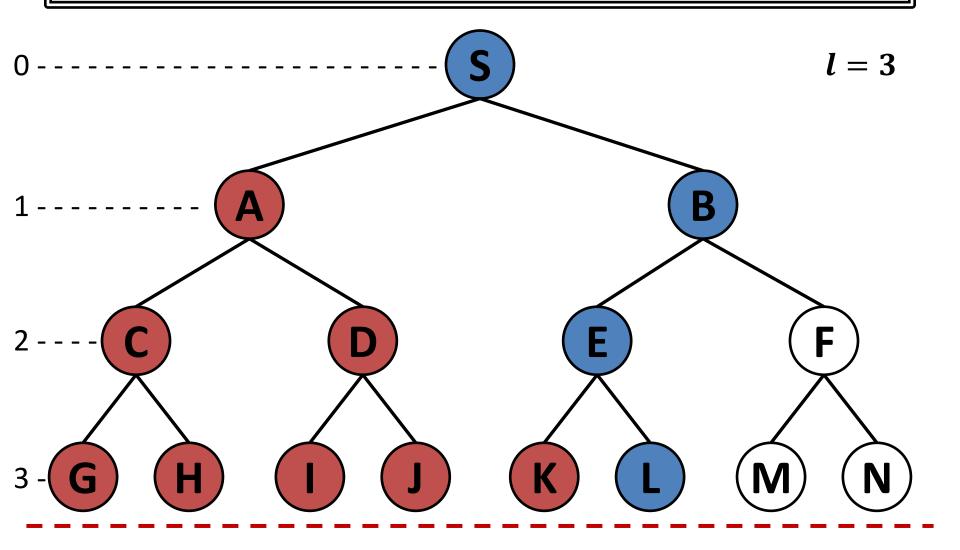


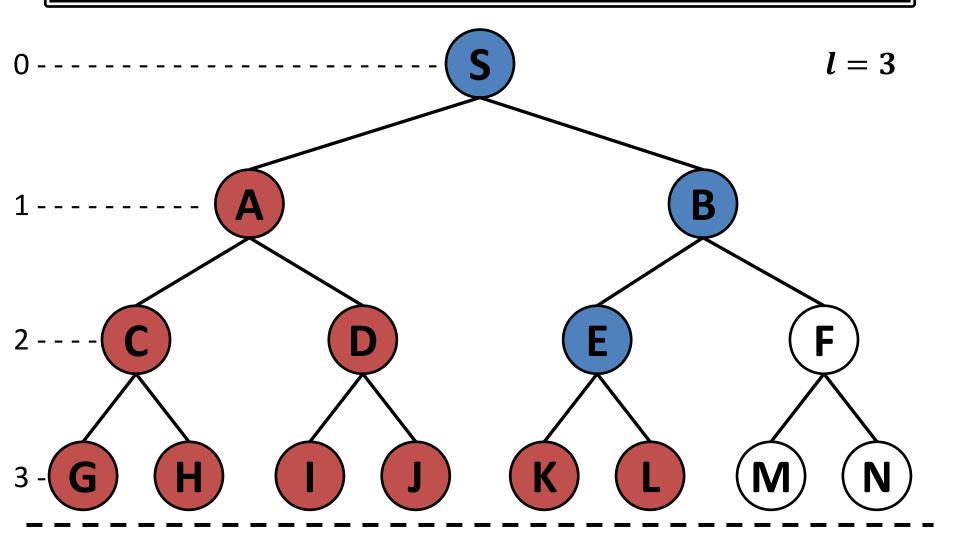


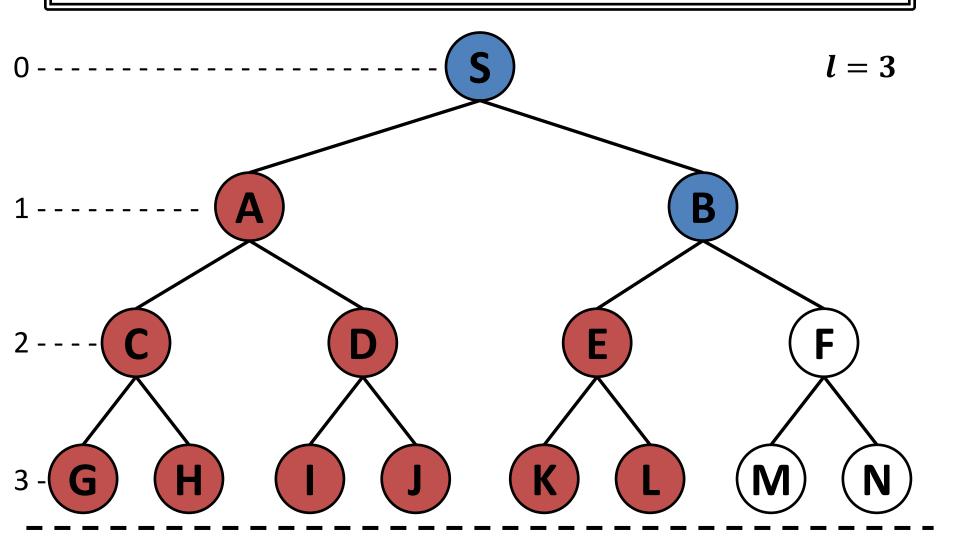


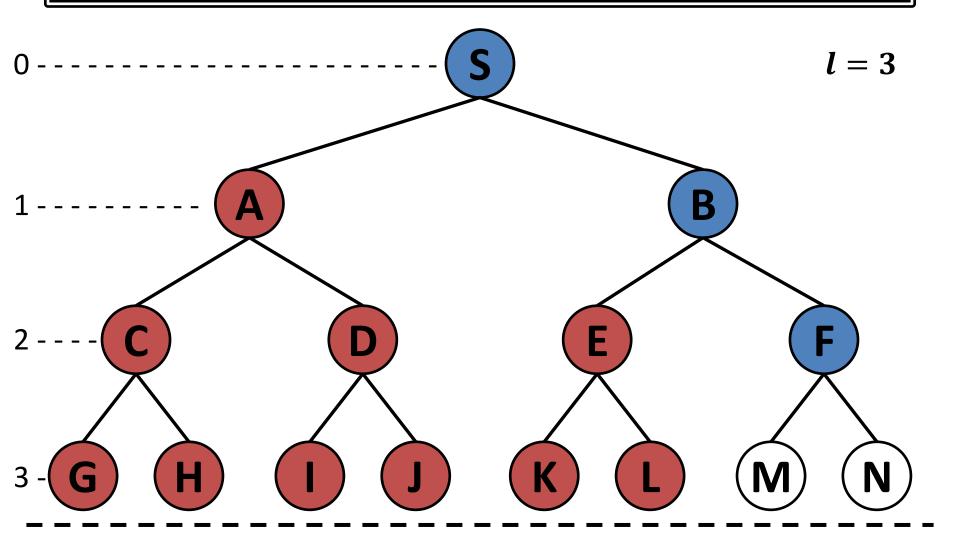


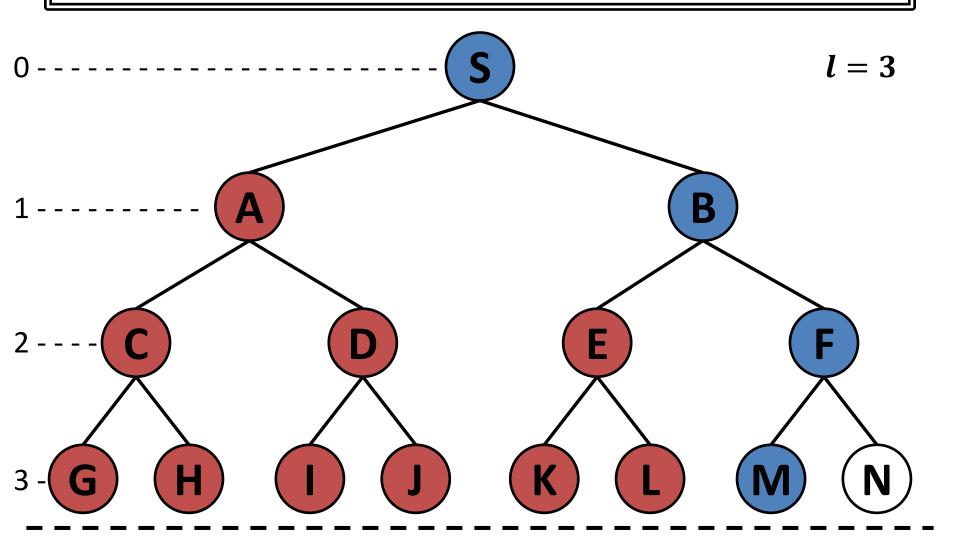


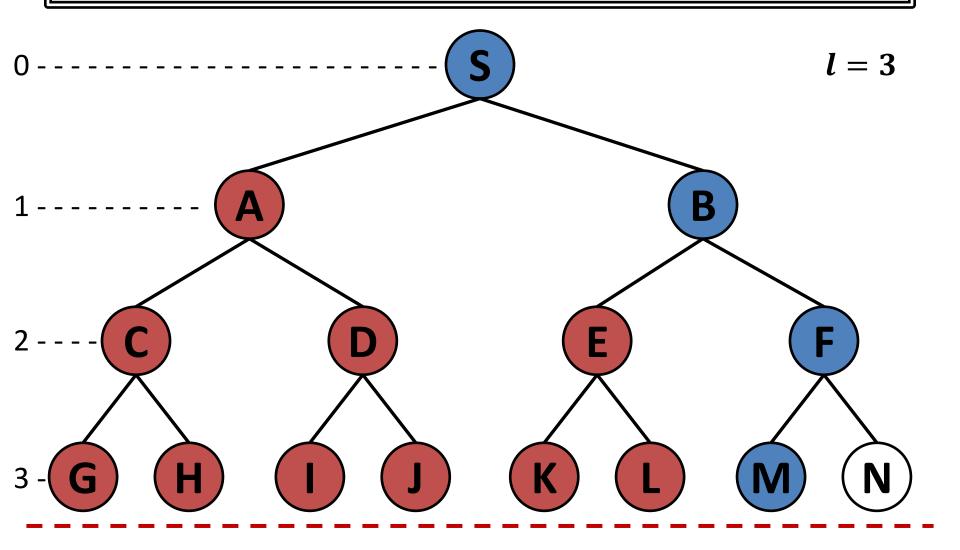


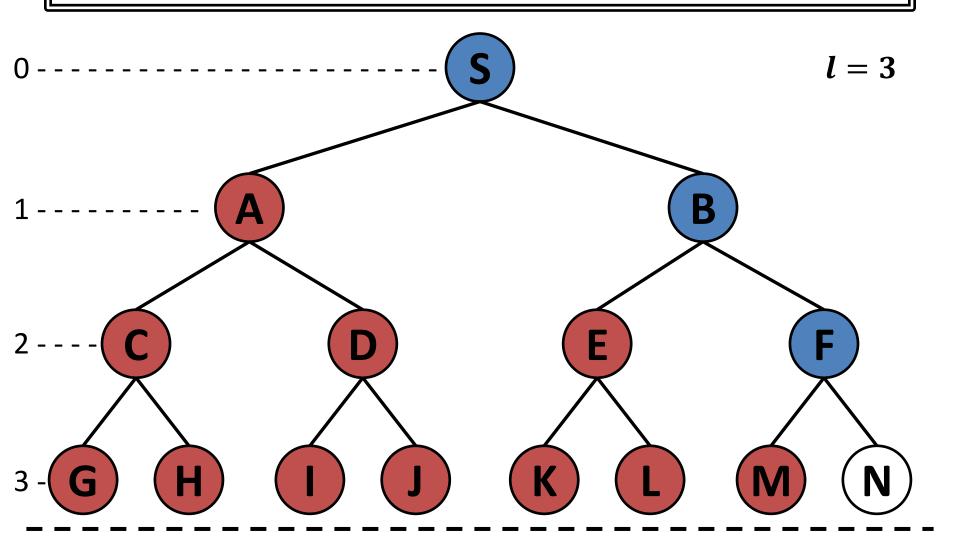


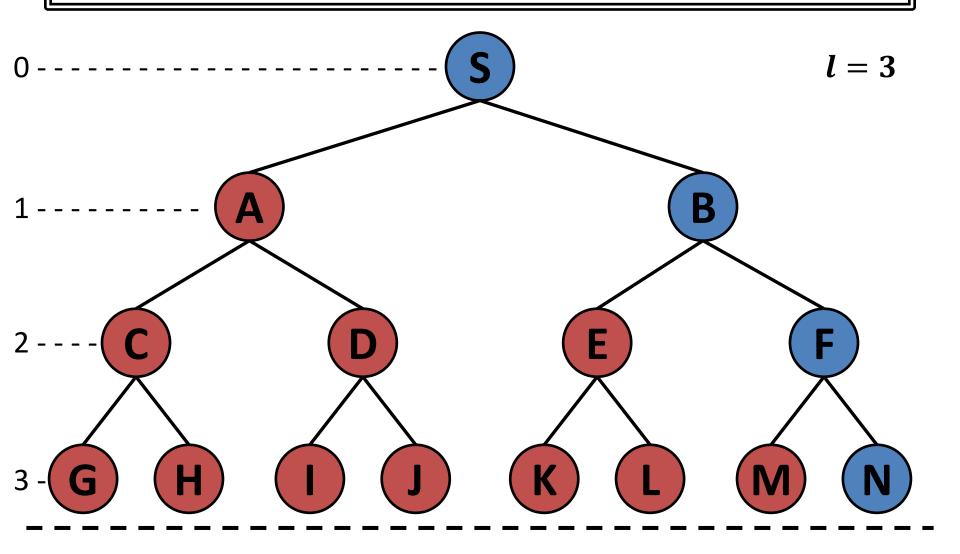


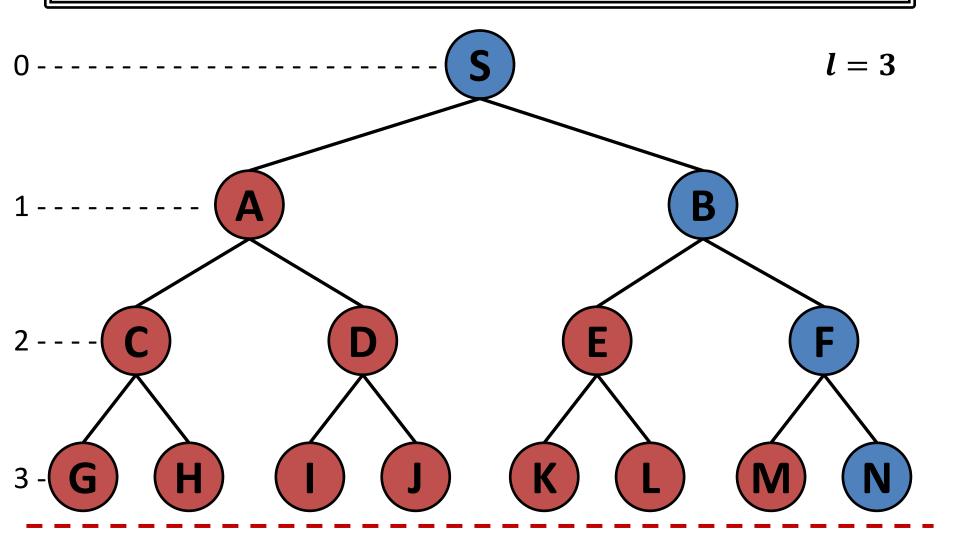


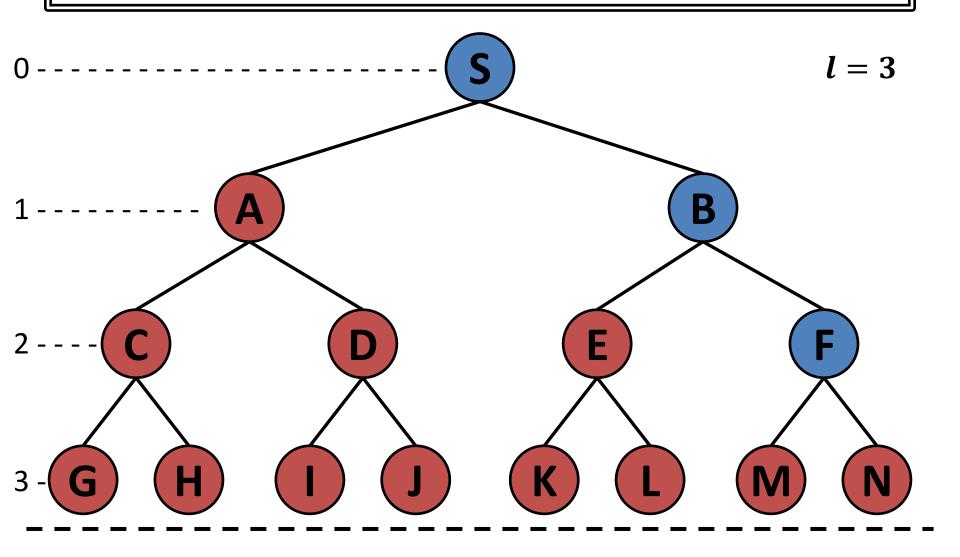


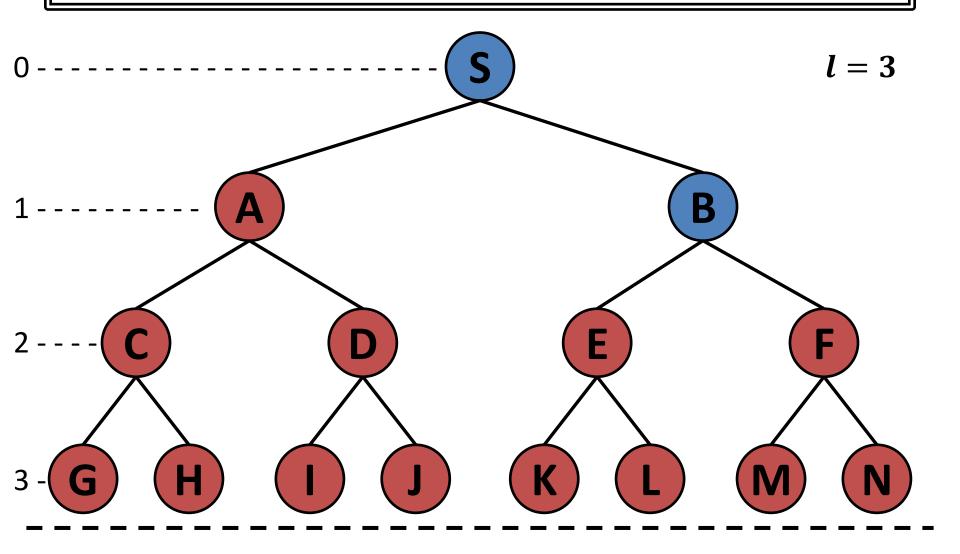


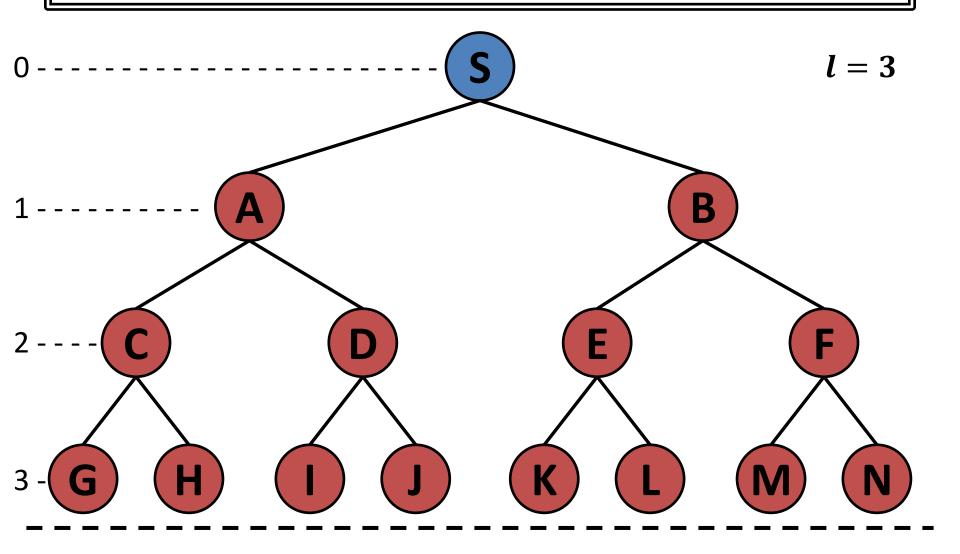


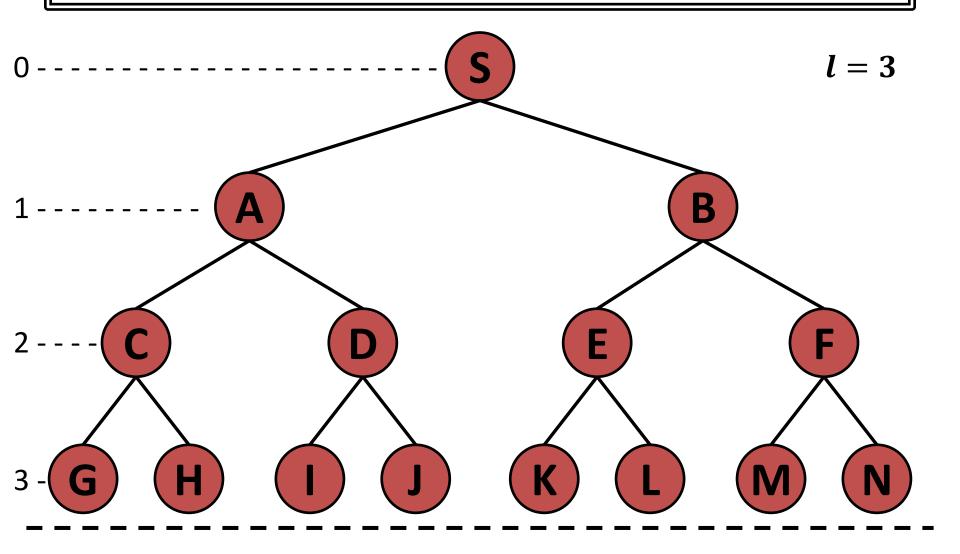












### Iterative Deepening Search (IDS)

#### Complete?

- Yes

#### Optimal?

- Yes, if step cost = 1 (Can be modified to explore uniform-cost tree)

#### • <u>Time?</u>

$$-(d+1)b^0+db^1+(d-1)b^2+\ldots+b^d=O(b^d)$$

#### Space?

- O(bd), i.e., linear space!

# Iterative Deepening Search (IDS)

Numerical comparison for b = 10 and d = 5, solution at far right leaf:

$$N(IDS) = 50 + 400 + 3,000 + 20,000 + 100,000 = 123,450$$

$$N(BFS) = 10 + 100 + 1,000 + 10,000 + 100,000 + 999,990 = 1,111,100$$

- \* IDS does better because other nodes at depth d are not expanded
- \* **BFS** can be modified to apply goal test when a node is generated

## Summary of Algorithms

Criterion	DFS	BFS	UCS	DLS	IDS
Complete	No	Yes*	Yes*	Yes*	Yes
		(if <b>b</b> is finite)	(if step cost $\geq \varepsilon$ )	(if $l \ge d$ )	
Time	$b^m$	$b^{d+1}$	$b^{\lceil C^*/\mathcal{E} \rceil}$	b'	b <sup>d</sup>
Space	bm	b <sup>d+1</sup>	b <sup>[C*</sup> / <i>E</i> ]	bl	bd
Optimal	No	<b>Yes*</b> (if step cost = 1)	Yes	No	Yes* (if step cost = 1)
Data Structure	Stack	Queue	Priority Queue	Stack	Stack

# Homework: "Frog Leap Puzzle"

- The game has 2N + 1 fields of play. In the beginning of the game there are N frogs which are looking to the left and placed on the rightmost N fields, and N frogs which are looking to the right and placed on the leftmost fields. The aim of the game is to swap the places of the frogs and reach the opposite configuration.
- The rules of play are as follows: Each frog can only move in the direction it looks at. Every frog can jump in a free field in front of it or jump over a frog to go to an empty field in front of it.
- Input: N number of frogs looking in one direction
- **Output**: All the configurations that needed to go through to get from the start to the final state.

## Homework: "Frog Leap Puzzle"

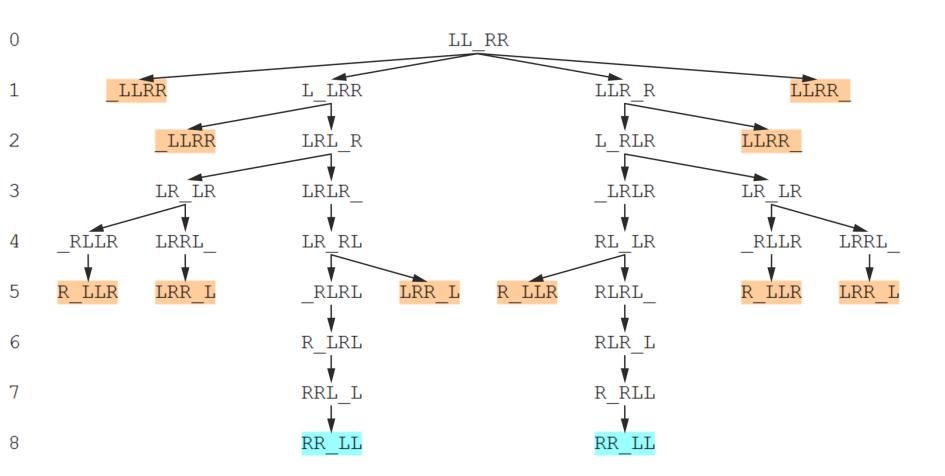
Sample Input:

2

Sample Output:

- How can we solve the problem?
- Can we solve the problem using uninformed search algorithm?
  - If "Yes", how can we represent the puzzle as a tree or a graph?
  - Which uninformed search algorithm is best? Why?
- Can we solve the problem in a better way?
- What about a linear time solution?

- Initial State: LL\_RR
- Successor Function (moves):
  - Left moves (only *R*)
    - Jump:  $LL_RR \rightarrow LLR_R$
    - Double Jump: LL\_RR → LLRR\_
  - Right moves (only L)
    - Jump: *LL\_RR* → *L\_LRR*
    - Double Jump: LL\_RR → \_LLRR
- Goal State: RR\_LL



```
n = readInput();
size = (n * 2) + 1;
board = new board[size];
function dfs(board, zeroState) {
      if isGoalState(board, zeroState)
            return TRUE;
      for move in moves(board, zeroState) {
            if dfs(move.board, move.zeroState) {
                  print(move.board);
                  return TRUE;
      return FALSE;
```