CSc 134 Database Management and File Organization

7. Functional Dependencies and Normalization for Relational Databases

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Introduction

- What is relational database design?

  The grouping of attributes to form relation schemas

- What are good relational design?
- Formal measures

Functional Dependencies

- FDs are **constraints** that are derived from

  meaning and interrelationships of the data attributes

- A functional dependency is a property of the semantics or meaning of the attributes.

Definition of functional dependency

- A **functional dependency**, denoted by $X \rightarrow Y$, between two sets of attributes X and Y that are subsets of R specifies a constraint on the possible tuples that can form a relation state r of R. The constraint is that, for any two tuples t1 and t2 in r that have $t1[X] = t2[X]$, they must also have $t1[Y]=t2[Y]$.

FD example

- A set of attributes X *functionally determines* a set of attributes Y if the value of X determines a unique value for Y.
- Social security number functionally determines employee name

  SSN $\rightarrow$ ENAME

Notation of Functional Dependencies

- $X \rightarrow Y$
  - function dependency from X to Y
  - Y is functionally dependent on X
  - X: left hand side FD. Y: right hand side FD
- $X \rightarrow Y$ holds if whenever two tuples have the same value for X, they *must have the same value for Y*
- A FD is a property of the relation schema R, not of a particular legal relation state r of R.
- $X \rightarrow Y$ in R specifies a *constraint* on **all** relation instances r(R)

Examples of FD

- Social security number determines employee name

  SSN $\rightarrow$ ENAME
- Project number determines project name and location

  PNUMBER $\rightarrow$ {PNAME, PLOCATION}
- Employee ssn and project number determines the hours per week that the employee works on the project

  {SSN, PNUMBER} $\rightarrow$ HOURS

Infer additional FDs

- Given a set of FDs F, we can *infer* additional FDs that hold whenever the FDs in F hold.

- Given a set of functional dependencies F
  - F= {SSN $\rightarrow$ ENAME
    PNUMBER $\rightarrow$ {PNAME, PLOCATION}
    {SSN, PNUMBER} $\rightarrow$ HOURS }

Infer?

- {ssn,bdata} $\rightarrow$ {ename,bdata}
- Pnumber $\rightarrow$ pname
- ssn $\rightarrow$ hours

Inference Rules for FDs

- Notation: XZ stands for {X,Z}

Armstrong's inference rules:

IR1. (Reflexive)

If $Y \subseteq X$, then $X \rightarrow Y$

IR2. (Augmentation)

If $X \rightarrow Y$, then $XZ \rightarrow YZ$

IR3. (Transitive)

If $X \rightarrow Y$ and $Y \rightarrow Z$, then $X \rightarrow Z$

Additional Inference Rules

IR 4:(**Decomposition**)

If $X \rightarrow YZ$, then $X \rightarrow Y$ and $X \rightarrow Z$

IR 5: (**Union**)

If $X \rightarrow Y$ and $X \rightarrow Z$, then $X \rightarrow YZ$

IR6: (**Psuedotransitivity**)

If $X \rightarrow Y$ and $WY \rightarrow Z$, then $WX \rightarrow Z$

- Deduced from IR1, IR2, and IR3

Closure

- **F^+: Closure** of F. The set of all dependencies that include F as well as all dependencies that can be inferred from f is called the closure of F.
- **X^+: Closure** of X under F. The set of attributes that are functionally determined by X based on F.
- X^+ can be calculated by repeatedly applying IR1, IR2, IR3 using the FDs in F

Algorithm to calculate X^+

Determining X^+, the closure of x under F

$X^+ := X;$

Repeat

    $oldX^+ := X^+;$

    for each functional dependency Y->Z in F do

    if $Y \subseteq X^+$ then $X^+ := X^+ \cup Z;$

Until $(X^+ = oldX^+);$

Example of calculate X^+

- F= {SSN-> ENAME,
  PNUMBER -> {PNAME, PLOCATION},
  {SSN,PNUMBER} ->HOURS}

- $\{SSN\}^+ = \{SSN, ENAME\}$
- $\{SSN, PNUMBER\}^+ = \{SSN, PNUMBER, ENAME, PNAME, PLOCATION, HOURS\}$
- $\{PNUMBER\}^+ =$ ____?

Equivalence of Sets of FDs

- Two sets of FDs F and G are **equivalent** if:
  - every FD in F can be inferred from G, *and*
  - every FD in G can be inferred from F
- F and G are **equivalent** if $F^{+} = G^{+}$

Definition: F **covers** G if every FD in G can be inferred from F (i.e., if $G^{+} \subseteq F^{+}$)

- F and G are **equivalent** if F covers G and G covers F

Minimal Sets of FDs

A set of FDs is **minimal** if it satisfies the following conditions:

(1) Every dependency in F has a single attribute for its right hand side.

(2) We cannot replace any dependency X -> A in F with a dependency Y -> A, where $Y \subseteq X$, and still have a set of dependencies that is equivalent to F.

(3) We cannot remove any dependency from F and have a set of dependencies that is equivalent to F.

Minimal Sets of FDs

- Every set of FDs has an equivalent minimal set
- There can be several equivalent minimal sets
- We can always find at least one minimal set using Algorithm 10.2

Algorithm 10.2 Finding a Minimal Cover F for a set of functional Dependencies E

1. Set F: = E;
2. Replace each functional dependency X → {A1, A2,..., An} in F by the n functional dependencies X → A1, X → A2,..., X → An
3. For each functional dependency X →A in F

   for each attribute B that is an element of X

   if {{F-{X →A}} U {{x-{B}} →A}} is equivalent to F

   then replace X → A with (X-{B}) → A in F
4. For each remaining functional dependency X→A in F

   if (F-{X →A}) is equivalent to F,

   then remove X →A from F.

What is normalization?

- **Normalization**: The process of decomposing unsatisfactory relations by breaking up their attributes into smaller relations
  - Use
    - keys
    - FDs

  to certify whether a relation schema is in a particular normal form

Practical Use of Normal Forms

- **Normalization** is carried out so that the resulting designs are of high quality and meet the desirable properties
- The practical utility of these normal forms becomes questionable when the constraints on which they are based are **hard to understand** or to **detect**
- The database designers ***need not*** normalize to the highest possible normal form.
- **Denormalization:** the process of storing the join of higher normal form relations as a base relation—which is in a lower normal form

Definitions of Keys and Attributes Participating in Keys

- A **superkey** of a relation schema $R = \{A_1, A_2, ..., A_n\}$ is a set of attributes $S \subseteq R$ with the property that no two tuples t_1 and t_2 in any legal relation state r of R will have $t_1[S] = t_2[S]$

$(k)^+ =$ ______ ?

- A **key** K is a superkey with the *additional property* that removal of any attribute from K will cause K not to be a superkey any more.

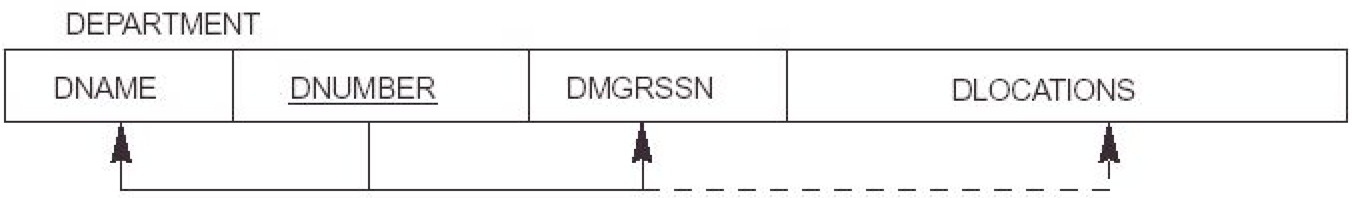
Definitions of Keys and Attributes Participating in Keys (Cont.)

- If a relation schema has more than one key, each is called a **candidate key.** One of the candidate keys is *arbitrarily* designated to be the **primary key,** and the others are called *secondary keys*.

First Normal Form

- Disallows composite attributes, multivalued attributes
- Disallows attributes whose values *for an individual tuple* are non-atomic
- Considered to be part of the definition of relation

Figure 10.8 Normalization into 1NF



DEPARTMENT

DNAME	DNUMBER	DMGRSSN	DLOCATIONS
Research	5	333445555	{Bellaire, Sugarland, Houston}
Administration	4	987654321	{Stafford}
Headquarters	1	888665555	{Houston}

Normalization into 1 NF

- Solution 1 (best)
  - Department(dname,<u>dnumber</u>, dmgrssn)
  - dept_loc(<u>dnumber, dlocation</u>)
- Solution 2
  - department(<u>dnumber,dlocation</u>,dname,dmgrssn)
- Solution 3
  - department(<u>dnumber</u>,dname,dmgrssn, dlocation1, dlocation2,dlocation3)

Full functional dependency

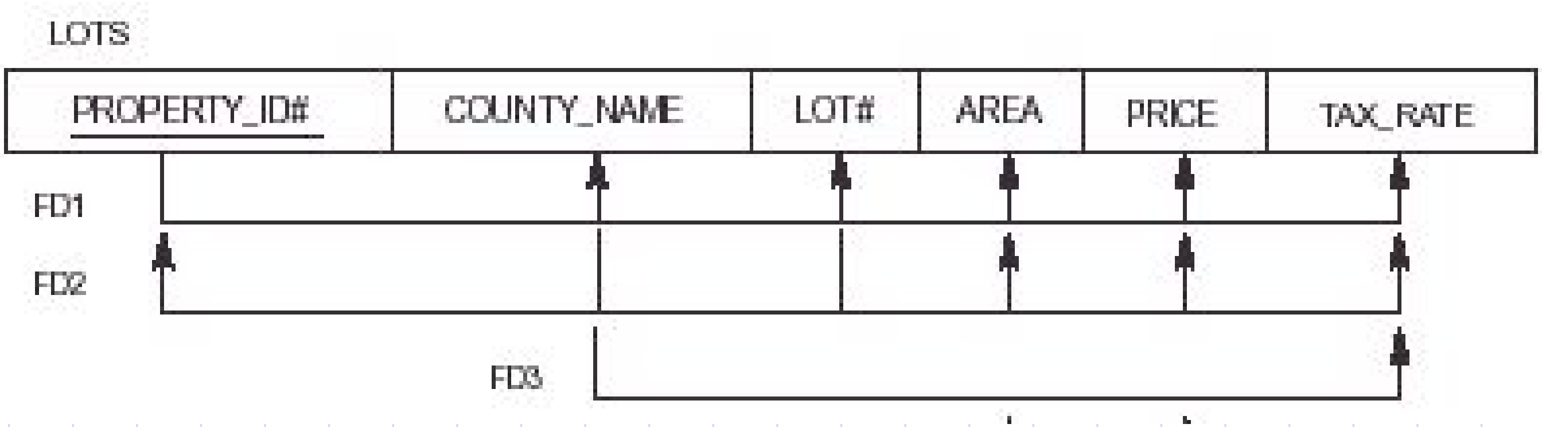
- Full functional dependency
  - a FD $Y \rightarrow Z$, where removal of any attribute from Y means the FD does not hold any more
  - e.g. {SSN, PNUMBER} $\rightarrow$ HOURS
- Partial dependency
  - e.g. {SSN, PNUMBER} $\rightarrow$ ENAME

2 NF

- General definition
- Take into account relations with *multiple candidate keys*
- **Prime attribute**: An attribute that is part of any candidate key
- A relation schema R is in **second normal form (2NF)** if <u>every</u> non-prime attribute A in R is fully functionally dependent on <u>every</u> key of R.

2NF - Example

- Keys:
  - property_id#
  - {county_name,lot#}
- Violate/satisfy? 2NF



3NF

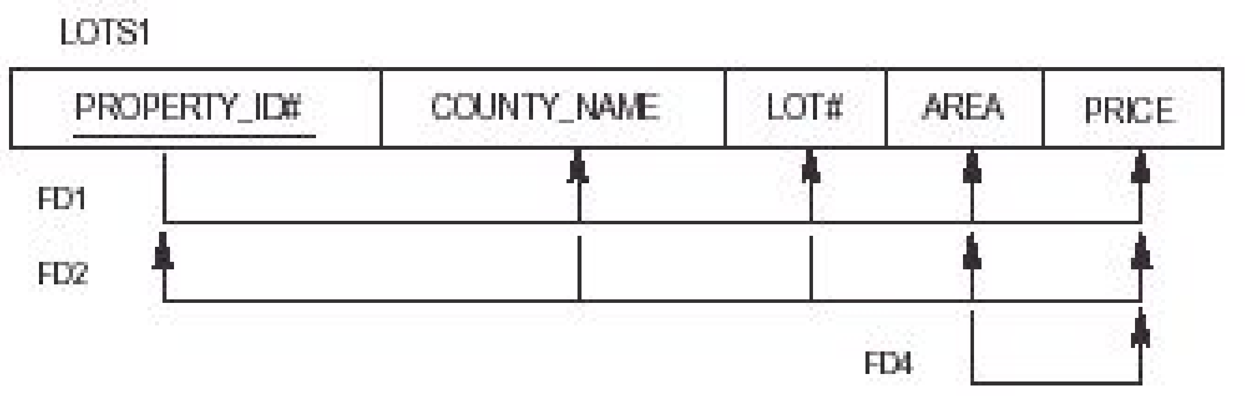
- $X \rightarrow Y$ is **trivial** if $Y \subset X$, otherwise, it is nontrival.
- A relation schema R is in **third normal form** (**3NF**) if, whenever a non-trivial FD $X \rightarrow A$ holds in R, then either:

  (1) X is a superkey of R, or

  (2) A is a prime attribute of R

3NF - example

- Keys:
  - property_id#
  - {county_name,lot#}
- Violate/satisfy? 3NF

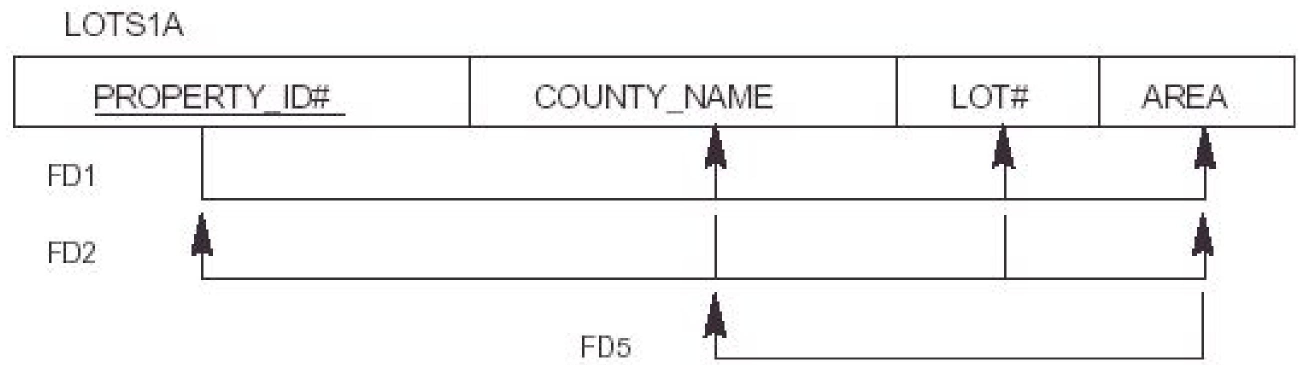


BCNF (Boyce-Codd Normal Form)

- A relation schema R is in **Boyce-Codd Normal Form** (**BCNF**) if whenever an nontrivial FD $X \rightarrow A$ holds in R, then X is a superkey of R
- Each normal form is strictly stronger than the previous one
  - Every 2NF relation is in 1NF
  - Every 3NF relation is in 2NF
  - Every BCNF relation is in 3NF
- There exist relations that are in 3NF but not in BCNF

BCNF - example

- Keys:
  - property_id#
  - {county_name,lot#}
- Violate/satisfy? BCNF



These slides are based on the textbook of:

R. Elmaseri and S. Navathe, *Fundamentals of Database Systems*, 6th Edition, Addison-Wesley. Chapter 15.