## COMP 230: Computer Architecture and Organization Fall 2018 EXAM 2 (PRACTICE)

## Instructions:

- This practice exam is meant to give you the flavor of questions that will be asked on the exam. Do not expect the real exam to be the same questions with only the numbers or MIPS instructions changed.
- Your real exam will be open book and notes. However, you are not allowed to use laptops, cell phones, calculators, or other electronics unless you have previously discussed this with the instructor.
- If you do not show your work, do not expect partial credit for incorrect answers.
- If you believe a problem is incorrectly or incompletely specified, make a reasonable assumption and solve the problem. The assumption should not result in a trivial solution.
- In all cases, clearly state any assumptions you make in your answers.

Question	Points Possible	Grade
1	10	
2	20	
3	20	
4	20	
5	30	
Total	100	

1. (a) Draw a table showing the process of dividing 12 by 5. Assume a machine that represents integers using only 4 bits. You may use either version of the division hardware we have seen.

Simple hardware:

Iteration	Step	Quotient	Divisor	Remainder
0	initial	0000	0101 0000	00001100
1	sub? and set Q0	000 <u>0</u>		(no change)
1	shift Div right		00101000	
2	sub? and set Q0	00 <u>00</u>		(no change)
2	shift Div right		00010100	
3	sub? and set Q0	0 <u>000</u>		(no change)
9	shift Div right		00001010	
4	sub? and set Q0	<u>0001</u>		00000010
4	shift Div right		00000101	
5	sub? and set Q0	<u>0010</u>		(no change)
J	shift Div right		_	

Concise hardware:

Iteration	Step	8 Bits
0	initial	0000 1100
1	shift	0001 1000
1	sub? and set Q0	0001 1000
2	$\operatorname{shift}$	0011 0000
2	sub? and set Q0	0011 0000
3	$\operatorname{shift}$	0110 0000
3	sub? and set Q0	0001 0001
4	$\operatorname{shift}$	0010 0010
4	sub? and set Q0	0010 0010

(b) If our MIPS machine operates on 4-bit words (so that we only have 4-bits to represent our integers), after a divide instruction, what instruction puts the quotient into \$t0?

mflo \$t0

2. Add 1.75 and 0.390625 using binary floating-point addition, assuming one round bit, one guard bit, one sticky bit, and using 5 fraction bits (5 bits of precision). Show your work. Also, what would the result be without a sticky bit?

$$1.75 = 1.11 \times 2^{0}$$
 =1.11000 | 00  
 $0.390625 = 0.011001 \times 2^{0}$  =0.01100 | 10  
 $0.00100 = 10$ 

The result is  $1.00010 \times 2^1$  no matter how you round. The sticky bit actually does not matter for this problem.

3. Write MIPS assembly code for a procedure divides that takes in two integers a and b as parameters. If a is divisible by b divides should return 1; otherwise 0.

```
divides:
           add
                  $v0.
                        $zero,
                                $zero
           div
                  $a0,
                        $a1
           mfhi
                  $t0
                  $t0,
                        $zero, exit
           bne
           addi
                  $v0,
                        $zero,
                                1
exit:
           jr
                  $ra
```

4. You have just written a new program for your company, which exhibits no pipeline stalls and has the following breakdown of executed instructions:

add	addi	not	beq	lw	sw
20%	20%	0%	25%	25%	10%

(a) In what fraction of all cycles is the data memory used?

$$25\% + 10\% = 35\%$$

(b) In what fraction of all cycles is the input of the sign-extend circuit needed? What is this circuit doing in cycles in which its input is not needed?

$$20\% + 25\% + 25\% + 10\% = 80\%$$

5. Consider the following sequence of instructions, executing using the pipeline shown in figure 1:

add \$s1, \$s2, \$s3 add \$s2, \$s1, \$s4 add \$s1, \$s1, \$s2

(a) Assuming no forwarding, what are the hazards?

3 data hazards: \$s1 RAW (read-after-write) between instructions 1 and 2, \$s2 RAW between instructions 2 and 3, and \$s1 RAW between instructions 1 and 3.

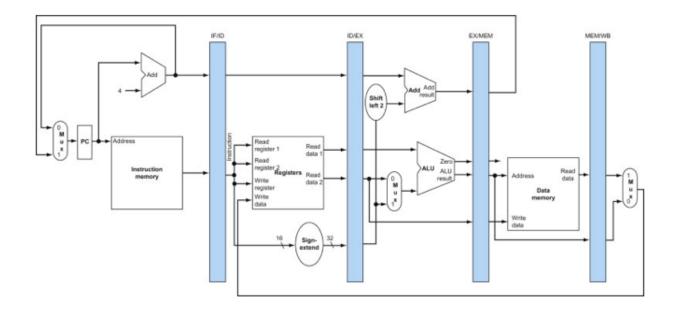


Figure 1: A simple MIPS pipeline.

(b) Assuming no forwarding, how many cycles does this sequence take to execute?

There are two bubbles after each add, so we can only start the last add on cycle 7. Then we need 4 more cycles to finish it, so 11 cycles total.

(c) Assuming full forwarding, what hazards remain?

None.

(d) Assuming full forwarding, how many cycles does this sequence take to execute?

Start the final add on the third cycle, then four more to finish it. In total, 7.