ANY-1 Instruction Set

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Instruction Formats

Immediate Format:

Extended Immediate

31 Constant_{31..8} 8 7 0 50_8 31 Constant_{55..32} 8 7 0 51_8

Constant_{79..56} 52₈

Register Format:

SR1 (one source register)

SR2 (two source register)

58 Instruction Modifier

59 Bitfield Modifier

Branch

5A Branch Modifier

Instruction Modifier

5C Stride Modifier

z: 1 = zero vector element if mask bit clear, 0 = vector element unchanged (ignored for scalar ops) m_3 : vector mask register (ignored for scalar operations).

Rm₃: rounding mode

If any of Rt, Ra, Rb, Rc are vector registers, then the instruction is a vector instruction.

 Rn_8

0 to 63 scalar registers 64 to 127 vector registers

128 to 255 Rn is a seven-bit constant

Execution Unit Integer Floating-point Decimal floating-point Posit	Qualifier .int .fp .dfp .pos	
Size	Oualifier	Alt Qualifier
byte	.b	
wyde	.W	
tetra	.t	.s (single)
octa	.0	.d (double)
hexi	.h	.q (quad)
SIMD byte	.bp	
SIMD wyde	.wp	
SIMD tetra	.tp	.sp
SIMD octa	.op	.dp
SIMD hexi	.hp	.qp
	Integer Floating-point Decimal floating-point Posit Size byte wyde tetra octa hexi SIMD byte SIMD wyde SIMD tetra SIMD octa	Integer .int Floating-point .fp Decimal floating-point .dfp Posit .pos Size Qualifier byte .b wyde .w tetra .t octa .o hexi .h SIMD byte .bp SIMD wyde .wp SIMD tetra .tp SIMD octa .op

Example Instruction

add.int.o x1,x2,x3,x0 ; scalar add of integers x2,x3

add.int.o v1,v2,v3,v0 ; vector add of integers v2,v3

add.int.o v1,v2,v0,x4 ; vector add scalar integers v2,x4

add.fp.o v1,v2,v3,v0 ; vector add float-point double v2,v3

Instructions

Arithmetic / Logical

ABS – Absolute Value

Description:

This instruction takes the absolute value of a register and places the result in a target register.

Instruction Format: R1

Operation:

$$If Ra < 0$$

$$Rt = -Ra$$

$$else$$

$$Rt = Ra$$

Vector Operation

for
$$x=0$$
 to VL - 1
$$if\left(Vm[x]\right)Rt[x]=Ra[x]<0\ ?\ -Ra[x]:Ra[x]$$

Execution Units: I, F, D, P

Exceptions: none

Notes:

For sign-magnitude formats this instruction simply clears the MSB of the number. No rounding occurs.

ADD - Addition

Description:

Add two values. The first operand must be in a register. The second operand may be in a register or may be an immediate value specified in the instruction.

Operation:

$$Rt = Ra + Imm$$

or

$$Rt = Ra + Rb$$

Vector Operation

for
$$x = 0$$
 to $VL - 1$

if
$$(Vm[x]) Vt[x] = Va[x] + Vb[x]$$

else if
$$(z)$$
 $Vt[x] = 0$

Instruction Format: RI

Instruction Format: R2

Vector Instruction Format: RI

Vector Instruction Format: R2

AND – Bitwise And

Description:

Perform a bitwise 'and' operation between operands. The first operand must be in a register. The second operand may be in a register of may be an immediate value specified in the instruction. A third source operand must be in a register. The immediate constant is one extended before use.

Instruction Format: RI

Instruction Format: R2

Operation:

$$Rt = Ra \& Imm$$

or

Rt = Ra & Rb & Rc

Vector Operation

for
$$x=0$$
 to VL - 1
$$if \ (Vm[x]) \ Vt[x] = Va[x] \ \& \ Vb[x] \ \& \ Vc[x]$$
 else if (z) $Vt[x] = 0$

AISIP - Add Immediate Shifted to IP

Description:

This instruction forms the sum of the instruction pointer and an immediate value shifted left a multiple of 32 times. The result is then placed in the target register. The low order 32 bits of the target register are zeroed out.

Instruction Format

63		32	3128	2724	23 16	15 8	7 0
	Constant ₃₂		F_4	Sh_4	63_{8}	Rt_8	Opcode ₈

BLEND – Blend Colors

Description:

This instruction blends two colors whose values are in Ra and Rb according to an alpha value in Rc. The resulting color is placed in register Rt. The alpha value is an eight-bit value assumed to be a binary fraction less than one. The color values in Ra and Rb are assumed to be RGB888 format colors. The result is a RGB888 format color. The high order eight bits of the result register are set to the high order eight bits of Ra. Note that a close approximation to 1.0 – alpha is used. Each component of the color is blended.

Instruction Format: R3

Operation:

$$Rt.R = (Ra.R * alpha) + (Rb.R * \sim alpha)$$

$$Rt.G = (Ra.G * alpha) + (Rb.G * \sim alpha)$$

$$Rt.B = (Ra.B * alpha) + (Rb.B * \sim alpha)$$

Clock Cycles: 2

BMM – Bit Matrix Multiply

BMM Rt, Ra, Rb

Description:

The BMM instruction treats the bits of register Ra and register Rb as an 8x8 matrix and performs a bit matrix multiply of the two registers and stores the result in the target register. An alternate mnemonic for this instruction is MOR.

Instruction Format: S2

63 61	60 58	57	50	4948	47 44	4341	40	39	32	31	24	23	16	15	8	7	0
Fn_3	Rm_3	03	h_8	U_2	Sz_4	m_3	Z	~	8	Rb) 8	R	a_8	R	8	($03h_8$

Fn ₃	Function
0	MOR
1	MXOR
2	MORT (MOR transpose)
3	MXORT (MXOR transpose)
4 to 7	reserved

Operation:

```
for I = 0 to 7  for \ j = 0 \ to \ 7 \\ Rt.bit[i][j] = (Ra[i][0]\&Rb[0][j]) \ | \ (Ra[i][1]\&Rb[1][j]) \ | \ \dots \ | \ (Ra[i][15]\&Rb[15][j])
```

Clock Cycles: 1

Execution Units: Integer ALU

Exceptions: none

Notes:

The bits are numbered with bit 63 of a register representing I,j = 0,0 and bit 0 of the register representing I,j = 7,7.

BYTNDX – **Byte Index**

Description:

This instruction searches Ra, which is treated as an array of eight bytes, for a byte value specified by Rb or an immediate value and places the index of the byte into the target register Rt. If the byte is not found -1 is placed in the target register. A common use would be to search for a null byte. The index result may vary from -1 to +7. The index of the first found byte is returned (closest to zero).

The result tag type is set to integer.

Instruction Format: SR2

0_{6}^{31}	6 25 20 Rb ₆	19 14 Ra ₆	Rt_6	7 0 $^{1}Ah_{8}$
31 28	27 20	19 14	13 8	7 0
1	Imm_8	Ra_6	Rt_6	$1Ah_8$

R2 Supported Formats: .0

Clock Cycles: 1

Execution Units: Integer ALU

Operation:

Rt = Index of (Rb in Ra)

CMP – Compare

Description

Compare two registers or a register and an immediate value and return the relationship between them. Both values are treated as signed numbers.

Instruction Format: R1

Operation:

$$Rt = Ra < Rb ? -1 : Ra = Rb ? 0 : 1$$

Vector Operation

for
$$x=0$$
 to VL - 1
$$if\left(Vm[x]\right)\,Vt[x] = Va[x] < Vb[x]\,\,?\,-1:Va[x] = Vb[x]\,\,?\,0:1$$

CNTPOP – Count Population

Description:

Count the number of ones and place the count in the target register.

Vector Operation

for
$$x = 0$$
 to $VL - 1$
$$if (Vm[x]) Vt[x] = popcnt(Va[x])$$

Instruction Format: R1

Execution Units: integer ALU

CNTLZ – Count Leading Zeros

Description:

Count the number of leading zeros (starting at the MSB) in Ra and place the count in the target register.

Instruction Format: R1

R1 Supported Formats: .o

Clock Cycles: 1

Execution Units: Integer ALU

COM – Ones Complement

Description:

Complement all the bits in the register. 1's become 0's and 0's become 1's.

Instruction Format: RI

Operation

$$Rt = Ra$$

OR

$$Rt = Ra$$

Vector Operation

for
$$x = 0$$
 to VL-1
$$if (Vm[x]) \ Vt[x] = \sim Va[x]$$

$$else \ if (z) \ Vt[x] = 0$$

$$else \ Vt[x] = Vt[x]$$

DEP – **Deposit**

Description:

Insert to a bitfield. Rc specifies the bitfield offset, Rd specifies the width of the bitfield. Rb specifies the data to insert. Ra contains the original source data. The least significant Rd minus one bits of Rb are inserted into Ra at the position specified by Rc. The final result is placed into Rt.

This instruction may also be used to perform a left shift of a single register by specifying x0 for Ra.

Formats Supported: R4

DT_3	Meaning
00	Rc,Rd are both regs
01	Rc is a six bit immediate, Rd is a reg
10	Rd is a six bit immediate, Rc is a reg
11	Both Rc, Rd are six bit immediates

Operation Size: .o

Execution Units: integer ALU

Exceptions: none

Example:

DIF – **Difference**

Description:

This instruction computes the difference between two signed values in registers Ra and Rb and places the result in a target Rt register. The difference is calculated as the absolute value of Ra minus Rb.

Instruction Format: R2, R2S

Supported Formats: .b .w, .t, .o, .h, .bv, .wv, .tv, .ov, .hv

Clock Cycles: 1

Execution Units: Integer

Operation:

Rt = Abs(Ra - Rb)

DIV – Division

Description:

Divide two operand values and place the result in the target register. The first operand must be in a register specified by the Ra field of the instruction. The second operand may be a register specified by the Rb field of the instruction or an immediate value. Both operands are treated as signed values.

Instruction Format: RI

Instruction Format: R2

Execution Units: ALU

Clock Cycles: 67

DIVR – **Division**

Description:

This instruction is supplied as division is not commutative. Divide two operand values and place the result in the target register. The first operand must be an immediate value. The second operand must be a register specified by the Ra field of the instruction. Both operands are treated as signed values. This instruction allows a constant to be divided by a register value "reverse" to how the DIV instruction works.

Formats Supported: RI

Execution Units: ALU

Clock Cycles: 67

Exceptions: none

DIVU – Division Unsigned

Description:

Divide two operand values and place the result in the target register. The first operand must be in a register specified by the Ra field of the instruction. The second operand may be either a register specified by the Rb field of the instruction, an immediate value. Both operands are treated as unsigned values.

Instruction Format: RI

Instruction Format: R2

Execution Units: ALU

Clock Cycles: 67

EXI0,EXI1,EXI2 – Extended Immediate

Description:

These instructions are used to extend the constant field of the following instruction. The constant is extended from bit eight. Multiple constant extensions may be present to extend a constant up to 64 bits. When multiple extensions are present they should be placed in order least significant to most significant. (EXIO first, EXI1 second, EXI2 third). The constant extensions sign extend to the width of the machine.

Constant extensions may be applied for most instructions with a constant field.

Interrupts are locked out between the modifier and the following instruction.

Instruction Format: EXI

31	Constant ₃₁₈	8	7 0 50 ₈
31	Constant ₅₅₃₂	8	7 0 51 ₈
31	Constant ₇₉₅₆	8	7 0 52 ₈

EXT – Extract Bitfield

Description:

A bitfield is extracted from the source by shifting the source to the right and 'and' masking. The result is sign extended to the width of the machine. This instruction may be used to sign extend a value from an arbitrary bit position. The width specified should be one less than the desired width. The source is value is contained in the register pair Ra, Rb. The field width is specified by Rc and field offset by Rd.

Instruction Format: R4

DT_3	Meaning
00	Rc,Rd are both regs
01	Rc is a six bit immediate, Rd is a reg
10	Rd is a six bit immediate, Rc is a reg
11	Both Rc, Rd are six bit immediates

Execution Units: Integer ALU

Exceptions: none

Notes:

EXTU – Extract Bitfield Unsigned

Description:

A bitfield is extracted from the source by shifting the source to the right and 'and' masking. The result is zero extended to the width of the machine. This instruction may be used to zero extend a value from an arbitrary bit position. The width specified should be one less than the desired width. The source is a 128-bit value which is the concatenation of Rb and Ra. Rc contains the field offset, Rd the width.

Instruction Format: R4

DT_3	Meaning
00	Rc,Rd are both regs
01	Rc is a six bit immediate, Rd is a reg
10	Rd is a six bit immediate, Rc is a reg
11	Both Rc, Rd are six bit immediates

Execution Units: Integer ALU

Exceptions: none

Notes:

FDP - Fused Dot Product

Description:

Calculate the dot product x = (a * b) + (c * d). The operations are fused together meaning no rounding occurs until the final product is produced.

Instruction Format: R4

FFO -Find First One

Description:

A bitfield contained in Ra is searched beginning at the most significant bit to the least significant bit for a bit that is set. The index into the bitfield of the bit that is set is stored in Rt. If no bits are set, then Rt is set equal to -1. The field offset is specified by Rc, the field width by Rd.

Instruction Format: R4

DT_3	Meaning
00	Rc,Rd are both regs
01	Rc is a six bit immediate, Rd is a reg
10	Rd is a six bit immediate, Rc is a reg
11	Both Rc, Rd are six bit immediates

Clock Cycles:

Execution Units: Integer

MAX – Maximum Value

Description:

Determines the maximum of two values in registers Ra, Rb and places the result in the target register Rt.

Instruction Format: R2

Operation:

$$IF\ Ra > Rb$$

$$Rt = Ra$$

$$else$$

$$Rt = Rb$$

MIN – Minimum Value

Description:

Determines the minimum of two values in registers Ra, Rb and places the result in the target register Rt.

Instruction Format: R2

Operation:

$$IF\ Ra < Rb$$

$$Rt = Ra$$

$$else$$

$$Rt = Rb$$

MOD – Instruction Modifier

Description:

Used to modify the operation of the following instruction.

Interrupts are locked out between the modifier and the following instruction.

Instruction Format: F8

```
31\ 29  28\ 26  25\ 20  19\ 14  1312  11\ 9  8  7  0  DT_3  Rm_3  Rd_6  Rc_6  A  m_3  z  58h_8  A[0]: 1 = apply vector mask, 0=ignore mask spec  A[1]: 1 = apply rounding mode. <math>0 = ignored rounding mode spec
```

There are four basic additional elements supplied for the following instruction.

- 1) A vector mask specification, used only by vector instructions.
- 2) Two additional source registers
- 3) A rounding mode specification, useful only to applicable instructions
- 4) A data type to help identify the operation required.

Two additional register fields allow up to four source operands for the following instruction. If these registers are not required they should be specified as x0.

Application of the vector mask and rounding mode are optional. Two bits in the 'A' field indicate which of these modifiers is applied.

DT_3	Interpretation: Instruction is operating on:
0	Integer
1	Floating point (double precision)
2	Posit (64-bit)
3-7	reserved

Instruction Format: F9

The F9 modifier is almost the same as the F8 modifier, except that the data type DT field is interpreted differently. The data type field uses two bits to indicate whether Rc and Rd are register specs or six-bit unsigned integer values. This modifier is used primarily for the bitfield extract / deposit instructions although it is also applicable to the SLLP instruction.

Instruction Format: FA

DT ₃	Consta	20	Rc_6	Rt_6	$5Ah_8$	
31 29	28	20	19 14	13 8	7 0	

The FA modifier applies to branch instructions to both extend the range of a branch and allow branch-to-register, and branch-and-link capability. When the FA modifier is present, the Rc register overrides the use of the IP in calculating the branch target address. The target address is then the sum of register Rc and a constant supplied in the instruction.

The constant field of the FA modifier adds an additional nine bits to the branch displacement. This allows branching extended to $\pm 4MB$.

The Rt field may be set to the address of the instruction following the branch, to allow conditional branch to subroutine capability.

Finally, the data type field identifies the type of data being compared by the branch instruction.

MUL – Signed Multiply

Description:

Multiply two values. The first operand must be in a register. The second operand may be in a register or may be an immediate value specified in the instruction. Both the operands are treated as signed values, the result is a signed result.

Instruction Format: RI

Instruction Format: R2

Execution Units: ALU

Vector Operation

for
$$x = 0$$
 to $VL - 1$
$$if (Vm[x]) Vt[x] = Va[x] * Vb[x]$$

MULF – Fast Unsigned Multiply

Description:

Multiply two values. The first operand must be in a register. The second operand may be in a register or may be an immediate value specified in the instruction. Both the operands are treated as unsigned values. The result is an unsigned result. The fast multiply multiplies only the low order 24 bits of the first operand times the low order 16 bits of the second. The result is a 40-bit unsigned product.

Instruction Format: R2

Instruction Format: RI

Execution Units: ALU

Clock Cycles: 1

MULU – Unsigned Multiply

Description:

Multiply two values. The first operand must be in a register. The second operand may be in a register or may be an immediate value specified in the instruction. Both the operands are treated as unsigned values, the result is a unsigned result.

Instruction Format: RI

Instruction Format: R2

Vector Operation

for
$$x = 0$$
 to $VL - 1$
$$if (Vm[x]) Vt[x] = Va[x] * Vb[x]$$

MUX – Multiplex

Description:

The MUX instruction performs a bit-by-bit copy of a bit of Rb to the target register if the corresponding bit in Ra is set, or a copy of a bit from Rc if the corresponding bit in Ra is clear.

Instruction Format

Exceptions: none

Execution Units: integer ALU

NEG - Negate

Description:

This is an alternate mnemonic for the SUBF instruction where the constant is zero.

Instruction Format: R2

Scalar Operation

$$Rt = 0 - Rb$$

Vector Operation

for
$$x = 0$$
 to $VL - 1$
$$if (Vm[x]) \ Vt[x] = 0 - Vb[x]$$

$$else \ if (z) \ Vt[x] = 0$$

$$else \ Vt[x] = Vt[x]$$

Notes

For sign-magnitude operations the sign bit is inverted, no subtract occurs. The result is not rounded.

NOT – Logical Not

Description:

This instruction takes the logical 'not' value of a register and places the result in a target register. If the source register contains a non-zero value, then a zero is loaded into the target. Otherwise, if the source register contains a zero a one is loaded into the target register.

NOT reduces the value to a single bit Boolean.

Instruction Format: R1

Operation:

$$Rt = !Ra$$

OR – Bitwise Or

Description:

Perform a bitwise or operation between operands. The immediate constant is zero extended before use.

Instruction Format: RI

Instruction Format: R2

Operation

$$Rt = Ra \mid Immediate$$
 OR $Rt = Ra \mid Rb$

Vector Operation

for
$$x = 0$$
 to VL-1
$$if \ (Vm[x]) \ Vt[x] = Va[x] \ | \ Vb[x] \ | \ Vc[x]$$

PERM – Permute Bytes

Description:

This instruction allows any combination of bytes in a source register to be copied to a target register. The low order twenty-four bits of register Rb or a 12-bit immediate constant are used to identify which source bytes are copied to the destination. The twenty-four-bit value is composed of eight three-bit fields. Field S0 indicates the source byte for target byte position 0. S1 indicates the source byte for target byte position 1. S2 to S7 work similarly for the remaining target bytes. There are many interesting possibilities with this instruction. A single source byte could be copied to all target byte positions for instance. Or the order of bytes in a word could be reversed.

Instruction Format: SR2, PERM

63 61	60 58	57 50	4948	47 44	4341	40	39 32	31 24	23 16	15 8	7 0
0_3	Rm_3	~8	O_2	Sz_4	m_3	Z	∼ 8	Rb_8	Ra_8	Rt_8	$17h_8$
63 61	60.50	57 50	4049	47 44	1211	40	20	24	22 16	15 0	7 0
05 01	00.38	37 30	4940	4/44	4341	40	39	24	23 10	13 6	/ 0
1_3	Rm_3	Imm ₂₃₁₆	0_2	Sz_4	m_3	Z	Imr	n_{150}	Ra_8	Rt_8	$17h_8$

Execution Units: integer ALU

Clock Cycles: 1

PTRDIF – Difference Between Pointers

Description:

Subtract two values then shift the result right. Both operands must be in a register. The right shift is provided to accommodate common object sizes. It may still be necessary to perform a divide operation after the PTRDIF to obtain an index into odd sized or large objects. Rc may vary from zero to thirty-one.

The result tag is forced to integer.

Instruction Format: R3

$$Rt = Abs(Ra - Rb) >> Rc$$

Clock Cycles: 1

Execution Units: Integer

Exceptions:

None

SEQ – Set if Equal

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is equal to a second operand in register (Rb) or an immediate constant then the target register is set to a one, otherwise the target register is set to a zero.

For floating-point operations positive and negative zero are considered equal.

If a vector operation is taking place then the target register is one of the vector mask registers.

Instruction Format: RI

SGE – Set if Greater Than or Equal

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is greater than or equal to a second operand in register (Rb) then the target register is set to a one, otherwise the target register is set to a zero. The operands are treated as signed values.

There is no immediate form to this instruction. An immediate equivalent may be achieved using the SGT instruction and adjusting the constant by one.

SGEU – Set if Greater Than or Equal Unsigned

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is greater than or equal to a second operand in register (Rb) then the target register is set to a one, otherwise the target register is set to a zero. The operands are treated as signed values.

There is no immediate form to this instruction. An immediate equivalent may be achieved using the SGTU instruction and adjusting the constant by one.

SGT - Set if Greater Than

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is greater than a second operand which is a constant supplied in the instruction, then the target register is set to a one, otherwise the target register is set to a zero. The operands are treated as signed values.

There is no register form of this instruction. The register equivalent operation may be performed using the SLT instruction and swapping the registers.

Instruction Format: RI

SGTU – Set if Greater Than Unsigned

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is greater than a second operand which is a constant supplied in the instruction, then the target register is set to a one, otherwise the target register is set to a zero. The operands are treated as signed values.

There is no register form of this instruction. The register equivalent operation may be performed using the SLTU instruction and swapping the registers.

Instruction Format: RI

SIGN - Sign

Synopsis

Take sign of value. This is an extended Mnemonic for the CMP instruction.

Description

The sign of a register is placed in the target register Rt.

Instruction Format: R1

Operation:

$$Rt = Ra < 0 ? -1 : Ra = 0 ? 0 : 1$$

Vector Operation

for
$$x = 0$$
 to VL - 1
$$if \ (Vm[x]) \ Vt[x] = Va[x] < 0 \ ? -1 : Va[x] = 0 \ ? \ 0 : 1$$

SLL –Shift Left Logical

Description:

Left shift an operand value by an operand value and place the result in the target register. Zeros are shifted into the least significant bits. The first operand must be in a register specified by the Ra. The second operand may be either a register specified by the Rb field of the instruction, or an immediate value.

Instruction Formats: R2

Instruction Formats: R2I

Operation Size: .o, .t, .w, .b

Execution Units: integer ALU

Exceptions: none

Example:

SLLP –Shift Left Logical Pair

Description:

Left shift a pair of operand values by an operand value and place the result in the target register. The upper 64 bits of the result are placed in the target register. Zeros are shifted into the least significant bits. The operand pair must be in registers specified by the Ra and Rc field of the instruction. The third operand may be either a register specified by the Rb field of the instruction, or an immediate value.

This instruction may also be used to perform a left rotate of a single register by specifying the same register for Ra and Rc.

Instruction Formats: R3

Operation Size: .o

Execution Units: integer ALU

Exceptions: none

Example:

SLT – Set if Less Than

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is less than a second operand in either a register (Rb) or a constant supplied in the instruction, then the target register is set to a one, otherwise the target register is set to a zero. The operands are treated as signed values.

The register form of the instruction may also be used to test for greater than by swapping the operands around.

Instruction Format: R2

47 44 58 56 55 48 39 32 31 24 23 16 15 8 4341 40 $03h_8$ Rm_3 $0Ch_8$ Sz_4 $2Ch_8$ Rb_8 Ra_8 Rt_8 m_3 Z

SLE – Set if Less Than or Equal

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is less than or equal to a second operand in register (Rb) then the target register is set to a one, otherwise the target register is set to a zero. The operands are treated as signed values.

There is no immediate form to this instruction. An immediate equivalent may be achieved using the SLT instruction and adjusting the constant by one.

SLEU – Set if Less Than or Equal

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is less than or equal to a second operand in register (Rb) then the target register is set to a one, otherwise the target register is set to a zero. The operands are treated as unsigned values.

There is no immediate form to this instruction. An immediate equivalent may be achieved using the SLTU instruction and adjusting the constant by one.

SLTU – Set if Less Than Unsigned

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is less than a second operand in either a register (Rb) or a constant supplied in the instruction, then the target register is set to a one, otherwise the target register is set to a zero. The operands are treated as unsigned values.

The register form of the instruction may also be used to test for greater than by swapping the operands around.

59 58 56 55 48 47 44 4341 40 39 32 31 24 23 16 15 8 7 0
$$\sim$$
 Rm₃ 0Ch₈ Sz₄ m₃ z 2Eh₈ Rb₈ Ra₈ Rt₈ 03h₈

SNE – Set if Not Equal

Description:

The set instruction places a 1 or 0 in the target register based on the relationship between the two source operands. If operand Ra is not equal to a second operand in register (Rb) or an immediate constant then the target register is set to a one, otherwise the target register is set to a zero.

For floating-point operations positive and negative zero are considered equal.

Instruction Format: RI

SRA –Shift Right Arithmetic Pair

Description:

This is an alternate mnemonic for the signed field extract **EXT** instruction.

Right shift a pair of operand values by an operand value and place the result in the target register. The lower 64 bits of the result are placed in the target register. The sign bit is shifted into the most significant bits. The operand pair must be in registers specified by the Ra and Rb field of the instruction. The third operand may be either a register specified by the Rc field of the instruction, or an immediate value.

Instruction Format: R4

DT_3	Meaning			
10	Rd is a six bit immediate, Rc is a reg			
11	Both Rc, Rd are six bit immediates			

Operation Size: .o

Execution Units: integer ALU

Exceptions: none

Example:

SRL -Shift Right Logical Pair

Description:

This is an alternate mnemonic for the unsigned field extract **EXTU** instruction.

Right shift a pair of operand values by an operand value and place the result in the target register. The lower 64 bits of the result are placed in the target register. Zeros are shifted into the most significant bits. The operand pair must be in registers specified by the Ra and Rb field of the instruction. The third operand may be either a register specified by the Rc field of the instruction, or an immediate value.

This instruction may also be used to perform a right rotate of a single register by specifying the same register for Ra and Rb.

Instruction Format: R4

DT_3	Meaning
10	Rd is a six bit immediate, Rc is a reg
11	Both Rc, Rd are six bit immediates

Instruction Format: SR4

58 56 47 44 32 31 24 23 16 15 8 0 4341 40 39 $2Ch_8$ Rm_3 BFh₈ 54 Rc_8 Rb_8 Rt_8 m_3 \mathbf{Z} Ra_8

Operation Size: .w

Execution Units: integer ALU

Exceptions: none

Example:

SUB - Subtract

Description:

Subtract two values. Both operands must be in a register.

Instruction Format: R2

Scalar Operation

$$Rt = Ra - Rb$$

Vector Operation

for
$$x = 0$$
 to $VL - 1$
$$if (Vm[x]) \ Vt[x] = Va[x] - Vb[x]$$

$$else \ if (z) \ Vt[x] = 0$$

$$else \ Vt[x] = Vt[x]$$

SUBF – Subtract From

Description:

Subtract two values. The first operand must be in a register. The second operand must be an immediate value specified in the instruction. There is no register form for this instruction.

Instruction Format: RI

Operation:

$$Rt = Imm - Ra$$

U21NDX - UTF21 Index

Description:

This instruction searches Ra, which is treated as an array of three UTF21 values, for a value specified by Rb and places the index of the value into the target register Rt. If the UTF21 value is not found -1 is placed in the target register. A common use would be to search for a null. The index result may vary from -1 to +2. The index of the first found value is returned (closest to zero).

The result is tagged as an integer.

Instruction Format: R2

Supported Formats: .0

Clock Cycles: 1

Execution Units: Integer ALU

Operation:

Rt = Index of (Rb in Ra)

WYDNDX – Wyde Index

Description:

This instruction searches Ra, which is treated as an array of four wydes, for a wyde value specified by Rb and places the index of the wyde into the target register Rt. If the wyde is not found -1 is placed in the target register. A common use would be to search for a null wyde. The index result may vary from -1 to +3. The index of the first found wyde is returned (closest to zero).

Instruction Format: R2

R2 Supported Formats: .0

Clock Cycles: 1

Execution Units: Integer ALU

Operation:

Rt = Index of (Rb in Ra)

XOR – Bitwise Exclusive Or

Description:

Perform a bitwise exclusive or operation between operands. The first operand must be in a register. The second operand may be a register or immediate value. A third operand must be in a register. The immediate constant is zero extended before use.

Instruction Format: RI

Instruction Format: R2

Operation

OR

$$Rt = Ra \wedge Rb$$

Vector Operation

for
$$x=0$$
 to VL-1
$$if \ (Vm[x]) \ Vt[x] = Va[x] \land Vb[x] \land Vc[x]$$

$$else \ if \ (z) \ Vt[x] = 0$$

$$else \ Vt[x] = Vt[x]$$

ZXB –**Zero** Extend Byte

Description:

This is an alternate mnemonic for the bitfield extract (EXTU) operation.

Instruction Format: EXT

A bitfield in the source specified by Ra is extracted, the result is copied to the target register. Rc specifies the bit offset. Rd specifies the bit width.

Clock Cycles: 1

Execution Units: Integer ALU

Exceptions: none

Notes:

ZXW -Zero Extend Wyde

Description:

This is an alternate mnemonic for the bitfield extract (EXTU) operation.

Instruction Format: BFI

A bitfield in the source specified by Ra is extracted, the result is copied to the target register. Rc specifies the bit offset. Rd specifies the bit width.

Clock Cycles: 1

Execution Units: Integer ALU

Exceptions: none

Notes:

ZXT –Zero Extend Tetra

Description:

This is an alternate mnemonic for the bitfield extract (EXTU) operation.

Instruction Format: EXT

A bitfield in the source specified by Ra is extracted, the result is copied to the target register. Rc specifies the bit offset. Rd specifies the bit width.

Clock Cycles: 1

Execution Units: Integer ALU

Exceptions: none

Notes:

Memory Operations

CACHE - Cache Command

CACHE Cmd, d[Rn]

Description:

This instruction commands the cache controller to perform an operation. Commands are summarized in the command table below. Commands may be issued to both the instruction and data cache at the same time. The address of the cache line to be invalidated is passed in Ra if needed.

Instruction Formats: CACHE

Commands:

IC_3	Mne.	Operation
0	NOP	no operation
1	invline	invalidate line associated with given address
2	invall	invalidate the entire cache (address is ignored)
3 to 7		reserved

DC_3	Mne.	Operation	
0	NOP	no operation	
1	enable	enable cache (instruction cache is always enabled)	
2	disable	not valid for the instruction cache	
3	invline	invalidate line associated with given address	
4	invall	invalidate the entire cache (address is ignored)	
5 to 7		reserved	

Notes:

LDx - Load

Description:

Load a value from memory into a register.

Formats Supported:

Register Indirect with Displacement

This mode may make use of immediate prefixes to extend the range.

Scalar Indexed Form (LD)

The effective address (EA) is calculated as the sum of Ra plus Rb multiplied by a scale.

z: 1 = zero extend, 0 = sign extend

$$\begin{array}{ccc} Sc_3 & Multiplier \\ 0 & 1 \\ 1 & 2 \\ 2 & 4 \\ 3 & 8 \\ 4 & 16 \end{array}$$

Operation:

Rt = Memory[d + Ra + Rb * Sc]

Vector forms

Stridden Form (LDS)

Data is loaded from memory addresses separated by the stride amount specified by register field Rc, beginning with the sum of Ra and an immediate value. If the vector mask bit is clear and the 'z' bit is set in the instruction then the corresponding element of the vector register is loaded with zero. If the vector mask bit is clear and the 'z' bit is clear in the instruction then the corresponding element of the vector register is left unchanged (no value is loaded from memory).

Elements are loaded only up to the length specified in the vector length register.

```
Vm[x]
          Z
               Result
  0
          0
               Vt[x] = Vt[x] (unchanged)
  0
          1
               Vt[x] = 0 (set to zero)
  1
          0
               Vt[x] = memory, sign extended
  1
          1
               Vt[x] = memory, zero extended
  U_2
         Unit
  0
         integer
   1
         floating-point
   2
         decimal-float
   3
         posit
 Sz_4
         Operation Size
  0
         byte
  1
         wyde
  2
         tetra
  3
         octa
  4
         hexi (double octa)
  5
         quad octa
  6
         reserved
  7
         pointer
         reserved
  . . .
  15
         cache cmd
```

Operation:

Indexed Form

```
\label{eq:continuous_section} \begin{split} & \text{for } x = 0 \text{ to vector length} \\ & \text{if } (Vm[x]) \\ & Vt[x] = Memory[d + Ra + Rb * x] \\ & \text{else} \\ & Vt[x] = z ? 0 : Vt[x] \end{split}
```

Data is loaded from memory addresses beginning with the sum of Ra and a vector element from Vc.

Operation:

$$\begin{split} n &= 0 \\ \text{for } x &= 0 \text{ to vector length} \\ &\quad \text{if } (Vm[x]) \\ &\quad Vt[x] &= Memory[d + Ra + Vb[x]] \\ &\quad \text{else} \\ &\quad Vt[x] &= z ? \ 0 : Vt[x] \end{split}$$

LDB – Load Byte (8 bits)

Description:

Data is loaded from the memory address which is the sum of an immediate value and the sum of Ra and Rb times a scale. The value loaded is sign extended from bit 7 to the machine width.

Formats Supported: LD

Operation:

 $Rd = Memory_8[d + Ra + Rb*Sc]$

Exceptions: none

LDBZ – Load Byte, Zero Extend (8 bits)

Description:

Data is loaded from the memory address which is the sum of an immediate value and the sum of Ra and Rb times a scale. The value loaded is zero extended from bit 8 to the machine width.

Formats Supported: LD

Operation:

 $Rd = Memory_8[d + Ra + Rb*Sc]$

LDO – Load Octa (64 bits)

Description:

Data is loaded into Rt from the memory address which is the sum of an immediate value and the sum of Ra and Rb scaled.

Formats Supported: RR,RI

Operation:

 $Rt = Memory_{64}[d + Ra + Rb*Sc]$

Execution Units: Mem

LDT – Load Tetra (32 bits)

Description:

Data is loaded from the memory address which is the sum of Ra and an immediate value or the sum of Ra and Rb scaled. The value loaded is sign extended from bit 31 to the machine width.

Formats Supported: RR,RI

Operation:

 $Rt = Memory_{32}[d + Ra + Rb*Sc]$

Execution Units: Mem

Exceptions: none

LDTZ – Load Tetra, Zero Extend (32 bits)

Description:

Data is loaded from the memory address which is the sum of Ra and an immediate value or the sum of Ra and Rb scaled. The value loaded is zero extended from bit 8 to the machine width.

Formats Supported: RR,RI

Operation:

 $Rt = Memory_{32}[d + Ra + Rb*Sc]$

Execution Units: Mem

LDW – Load Wyde (16 bits)

Description:

Data is loaded from the memory address which is the sum of Ra and an immediate value or the sum of Ra and Rb scaled. The value loaded is sign extended from bit 15 to the machine width.

Formats Supported: LD

Operation:

 $Rt = Memory_{16}[d + Ra + Rb*Sc]$

Execution Units: Mem

Exceptions: none

LDWZ – Load Wyde, Zero Extend (16 bits)

Description:

Data is loaded from the memory address which is the sum of Ra and an immediate value or the sum of Ra and Rb scaled. The value loaded is zero extended from bit 16 to the machine width.

Formats Supported: LD

Operation:

 $Rt = Memory_{16}[d + Ra + Rb*Sc]$

Execution Units: Mem

LEA – Load Effective Address

Description:

This instruction computes the effective address for a load/store operation. The data type tag for the target register is set to indicate it contains a pointer.

Formats Supported:

Scalar Indexed Form (LD)

Operation:

$$Rt = d + Ra + Rb * Sc$$

Vector forms

Stridden Form (LDS)

```
4948
                              47 44
                                        4341
                                                40
                                                       39
                                                            32
                                                                                   23 16
                                                                                                15 8
                                                                      Rb_{8} \\
                                                                                                              69h<sub>8</sub>
Const_{21..8}
                      U_2
                               Sz_4
                                        m_3
                                                z Const<sub>7..0</sub>
                                                                                   Ra_8
                                                                                                 Rt_8
```

- Vm[x] z Result
 - 0 Vt[x] = Vt[x] (unchanged)
 - 0 1 Vt[x] = 0 (set to zero)
 - 1 0 Vt[x] = memory address
 - 1 Vt[x] = memory address
 - U₂ Unit
 - 0 integer
 - 1 floating-point
 - 2 decimal-float
 - 3 posit
 - Sz₄ Operation Size
 - 0 byte
 - 1 wyde
 - 2 tetra
 - 3 octa
 - 4 hexi

Operation:

$$\label{eq:continuous_section} \begin{split} &for \ x=0 \ to \ vector \ length \\ &if \ (Vm[x]) \\ &Vt[x] = d + Ra + Rb * x \\ &else \\ &Vt[x] = z \ ? \ 0 : Vt[x] \end{split}$$

Indexed Form

Operation:

$$\begin{split} n &= 0 \\ \text{for } x &= 0 \text{ to vector length} \\ &\quad \text{if } (Vm[x]) \\ &\quad Vt[x] &= d + Ra + Vb[x] \\ &\quad \text{else} \\ &\quad Vt[x] &= z \ ? \ 0 : Vt[x] \end{split}$$

LSM – Load or Store Multiple

Description:

The LSM prefix instruction allows multiple registers or values to be loaded or stored using the following load / store instruction. Register x0 cannot be stored using this prefix. If the register spec field is zero then no load or store takes place at that position. Up to seven registers may be specified.

Formats Supported: LSM

Execution Units: Mem

STx - Store

Description:

Store values to memory. Either the contents of a scalar or vector register or a six-bit immediate constant may be stored. Both scalar and vector store operations are possible.

Formats Supported:

Register Indirect with Displacement

Scalar Indexed Form (ST)

The effective address (EA) is calculated as the sum of Ra plus Rc multiplied by a scale.

Operation:

Memory[d+Ra + Rb * Sc] = Rs

Vector forms

Stridden Form (STS)

Data is stored to memory addresses separated by the stride amount specified by register field Rb, beginning with the sum of Ra and an immediate value. If the vector mask bit is clear and the 'z' bit is set in the instruction then memory for the corresponding element of the vector register is

stored with zero. If the vector mask bit is clear and the 'z' bit is clear in the instruction then memory corresponding to the element of the vector register is left unchanged (no value is stored to memory).

Elements are loaded only up to the length specified in the vector length register.

```
Vm[x]
               Result
          \mathbf{z}
  0
          0
               Memory = Memory (unchanged)
  0
               Memory = 0 (set to zero)
          1
  1
          0
               memory = Vt[x]
  1
               memory = Vt[x]
          1
  U_2
         Unit
  0
         integer
  1
         floating-point
  2
         decimal-float
  3
         posit
 Sz_4
         Operation Size
  0
         byte
  1
         wyde
  2
         tetra
  3
         octa
  4
         hexi
 5,6
         reserved
  7
         pointer
```

Operation:

```
\label{eq:for x = 0 to vector length} $if (Vm[x])$ $Memory[d+Ra+Rb*x] = Vt[x]$ $else$ $Memory[d+Ra+Rb*x] = z?0: Memory[d+Ra+Rb*x]$
```

Indexed Form

Data is stored to memory addresses beginning with the sum of Ra and a vector element from Vb.

Operation:

```
\begin{split} n &= 0 \\ for \ x &= 0 \ to \ vector \ length \\ & if \ (Vm[x]) \\ & Memory[d + Ra + Vb[x]] = Vt[x] \\ & else \\ & Memory = z \ ? \ 0 \ : Memory \end{split}
```

STB – Store Byte (8 bits)

Description:

This instruction stores a byte (8 bit) value to memory. The memory address is calculated as the sum of an immediate constant and the sum of Ra and Rb scaled.

Instruction Format: ST

Operation:

$$Memory_8[d + Ra + Rb*Sc] = Rs$$

STBZ – Store Byte and Zero (8 bits)

Description:

This instruction stores a byte (8 bit) value to memory. The memory address is calculated as the sum of an immediate constant and the sum of Ra and Rb scaled. After the byte is stored to memory the register is zeroed out.

Instruction Format: ST

Operation:

$$\begin{aligned} & Memory_8[d+Ra+Rb*Sc] = Rs \\ & Rs = 0 \end{aligned}$$

STO – Store Octa (64 bits)

Description:

This instruction stores an octa-byte (64 bit) value to memory. The memory address is calculated as the sum of an immediate constant and the sum of Ra and Rb scaled.

Instruction Format: ST

Operation:

 $Memory_{64}[d + Ra + Rb*Sc] = Rs$

STOZ – Store Octa and Zero (64 bits)

Description:

This instruction stores an octa-byte (64 bit) value to memory. The memory address is calculated as the sum of an immediate constant and the sum of Ra and Rb scaled. After the tetra is stored to memory the register is zeroed out.

Instruction Format: ST

Operation:

 $Memory_{64}[d + Ra + Rb*Sc] = Rs$

Rs = 0

STPTR – Store Pointer (64 bits)

Description:

This instruction stores an octa-byte (64 bit) value to memory. The memory address is calculated as the sum of an immediate constant and the sum of Ra and Rb scaled. STPTR begins a series of stores to memory addresses scaled by eight bits, until the address zero is reached. The first store proceeds normally, for the second and subsequent stores a byte store operation takes place with the value zero being to memory.

The purpose of the STPTR instruction is to allow a code dense implementation of a write barrier that indicates where in memory a pointer is stored with increasing resolution.

This instruction assumes that card memory used to record pointer locations is located at the low end of the memory system.

Instruction Format: ST

Operation:

```
\begin{aligned} ea &= d + Ra + Rb*Sc \\ Memory_{64}[ea] &= Rs \\ while &ea <> 0 \\ &ea = ea >> 8 \\ Memory_{8}[ea] &= 0 \end{aligned}
```

STT – Store Tetra (32 bits)

Description:

This instruction stores a tetra-byte (32 bit) value to memory. The memory address is calculated as the sum of an immediate constant and the sum of Ra and Rb scaled.

Instruction Format: ST

Operation:

 $Memory_{32}[d + Ra + Rb*Sc] = Rs$

STTZ – Store Tetra and Zero (32 bits)

Description:

This instruction stores a tetra-byte (32 bit) value to memory. The memory address is calculated as the sum of an immediate constant and the sum of Ra and Rb scaled. After the tetra is stored to memory the register is zeroed out.

Instruction Format: ST

Operation:

 $Memory_{32}[d + Ra + Rb*Sc] = Rs$ Rs = 0

STW – Store Wyde (16 bits)

Description:

This instruction stores a byte (16 bit) value to memory. The memory address is calculated as the sum of an immediate constant and the sum of Ra and Rb scaled.

Instruction Format: ST

Operation:

 $Memory_{16}[d + Ra + Rb*Sc] = Rs$

STWZ – Store Wyde and Zero (16 bits)

Description:

This instruction stores a byte (16 bit) value to memory. The memory address is calculated as the sum of an immediate constant and the sum of Ra and Rb scaled. After the wyde is stored to memory the register is zeroed out.

Instruction Format: ST

Operation:

$$Memory_{16}[d + Ra + Rb*Sc] = Rs$$

$$Rs = 0$$

Flow Control (Branch Unit) Operations

Branches

The branch modifier may be used to make it possible to branch to a target address contained in a register, and to store the return address in a register. Simultaneously the branch displacement is increased to 24 bits allowing a ± 32 MB branch range.

BAL – Branch and Link

Description:

This instruction may be used to call a subroutine. The address of the instruction after the BAL is stored in the specified return address register (Rt) then a jump to the address specified in the instruction is made. The address range is 24 bits or $\pm 8MB$. The constant is shifted left twice before use.

The return address register is assumed to be x1 if not otherwise specified. The BAL instruction does not require space in branch predictor tables.

Formats Supported: BAL

Flags Affected: none

Operation:

Rt = IP + 4

IP = IP + displacement

Execution Units: Branch

Exceptions: none

BEQ – Branch if Equal

Description:

This instruction branches to the target address if the contents of Ra and Rb are equal, otherwise program execution continues with the next instruction. With a branch modifier instruction, the target address is formed as the sum of Rc and a displacement. If Rc is r63 then the instruction pointer value is used. Otherwise, the target address is the sum of the instruction pointer value and the displacement specified in the instruction.

Formats Supported: BR

Operation:

```
If (Ra = Rb)
IP = IP + Displacement 14
With Modifier
Rt = IP + 4
If (Ra = Rb)
IP = Rc + Displacement 23
```

Execution Units: Branch

Exceptions: none

Notes:

For a floating-point comparison positive and negative zero are considered equal.

BGE – Branch if Greater Than or Equal

Description:

This instruction branches to the target address if the contents of Ra is greater than or equal to Rb, otherwise program execution continues with the next instruction. The values in Ra and Rb are treated as signed values. The target address is formed as the sum of Rc and a displacement. If Rc is x63 then the instruction pointer value is used.

Formats Supported: BR

Operation:

If
$$(Ra \ge Rb)$$

 $IP = IP + Displacement$

Execution Units: Branch

BGEU – Branch if Greater Than or Equal Unsigned

Description:

This instruction branches to the target address if the contents of Ra is greater than or equal to Rb, otherwise program execution continues with the next instruction. The values in Ra and Rb are treated as unsigned values. The target address is formed as the sum of Rc and a displacement. If Rc is r63 then the program counter value is used.

Formats Supported: BR

Operation:

$$Rt = IP + 8$$

$$If (Ra >= Rb)$$

$$PC = Rc + Displacement$$

Execution Units: Branch

BGT – Branch if Greater Than

Description:

This instruction is an alternate mnemonic for the <u>BLT</u> instruction where the register operands have been swapped.

This instruction branches to the target address if the contents of Ra is less than Rb, otherwise program execution continues with the next instruction. The values in Ra and Rb are treated as signed values. The target address is formed as the sum of Rc and a displacement. If Rc is x63 then the program counter value is used.

Formats Supported: BR

Operation:

$$Rt = IP + 8$$

$$If (Ra < Rb)$$

$$PC = Rc + Displacement$$

Execution Units: Branch

BGTU – Branch if Greater Than Unsigned

Description:

This instruction is an alternate mnemonic for the <u>BLTU</u> instruction where the register operands have been swapped.

This instruction branches to the target address if the contents of Ra is less than Rb, otherwise program execution continues with the next instruction. The values in Ra and Rb are treated as unsigned values. The target address is formed as the sum of Rc and a displacement. If Rc is x63 then the program counter value is used.

Formats Supported: BR

Operation:

$$Rt = IP + 8$$

$$If (Ra < Rb)$$

$$PC = Rc + Displacement$$

Execution Units: Branch

BNE – Branch if Not Equal

Description:

This instruction branches to the target address if the contents of Ra and Rb are not equal, otherwise program execution continues with the next instruction. The target address is formed as the sum of Rc and a displacement. If Rc is x63 then the program counter value is used.

Formats Supported: BR

Operation:

$$Rt = IP + 8$$

$$If (Ra <> Rb)$$

$$PC = Rc + Displacement$$

Execution Units: Branch

BLE – Branch if Less Than or Equal

Description:

This is an alternate mnemonic for the BGE instruction, where the register operands have been swapped.

This instruction branches to the target address if the contents of Ra is greater than or equal to Rb, otherwise program execution continues with the next instruction. The values in Ra and Rb are treated as signed values. The target address is formed as the sum of Rc and a displacement. If Rc is x63 then the program counter value is used.

Formats Supported: BR

Operation:

```
If (Ra \ge Rb)

PC = Rc + Displacement
```

Execution Units: Branch

Exceptions: none

BLEU – Branch if Less Than or Equal Unsigned

Description:

This is an alternate mnemonic for the BGEU instruction, where the register operands have been swapped.

This instruction branches to the target address if the contents of Ra is greater than or equal to Rb, otherwise program execution continues with the next instruction. The values in Ra and Rb are treated as unsigned values. The target address is formed as the sum of Rc and a displacement. If Rc is x63 then the program counter value is used.

Formats Supported: BR

Operation:

```
If (Ra \ge Rb)

PC = Rc + Displacement
```

Execution Units: Branch

BLT – Branch if Less Than

Description:

This instruction branches to the target address if the contents of Ra is less than Rb, otherwise program execution continues with the next instruction. The values in Ra and Rb are treated as signed values. The target address is formed as the sum of Rc and a displacement. If Rc is x63 then the program counter value is used.

Formats Supported: BR

63 50	4948	47 44	43 40	39 32	31 24	23 16	15 8	7 0
Displacement ₂₀₇	U_2	Sz_4	$Disp_{63}$	Rc_8	Rb_8	Ra_8	Rt_8	$48h_8$

Operation:

```
Rt = IP + 8 If (Ra < Rb) PC = Rc + Displacement
```

Execution Units: Branch

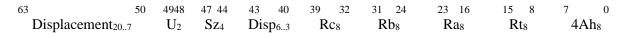
Exceptions: none

BLTU – Branch if Less Than Unsigned

Description:

This instruction branches to the target address if the contents of Ra is less than Rb, otherwise program execution continues with the next instruction. The values in Ra and Rb are treated as unsigned values. The target address is formed as the sum of Rc and a displacement. If Rc is x63 then the program counter value is used.

Formats Supported: BR



Operation:

$$Rt = IP + 8$$

$$If (Ra < Rb)$$

$$PC = Rc + Displacement$$

Execution Units: Branch

BRA – Unconditional Branch

Description:

This instruction is an alternate mnemonic for the \underline{BAL} instruction. The address range is 24 bits or $\pm 8MB$. The constant field is shifted left twice before use.

Formats Supported: JAL

Flags Affected: none

Operation:

IP = IP + Displacement

Execution Units: Branch

Exceptions: none

CHK – Check Register Against Bounds

Description:

A register is compared to two values. If the register is outside of the bounds then an exception will occur.

Instruction Format: RI

Cn₂ Interpretation

- 0 Ra \leq Rc \leq Constant
- 1 Ra < Rc <= Constant
- $2 Ra \le Rc < Constant$
- $3 ext{Ra} < Rc < Constant$

Instruction Format: R3

Cn₂ Interpretation

- $0 ext{Ra} \ll Rb \ll Rc$
- 1 Ra < Rb <= Rc
- $2 ext{Ra} <= Rb < Rc$
- 3 Ra < Rb < Rc

Supported Formats: .0

Clock Cycles: 2

Execution Units: Integer ALU, Float, Decimal Float, Posit

Exceptions: bounds check

Notes:

The system exception handler will typically transfer processing back to a local exception handler.

JAL – Jump and Link

Description:

This instruction may be used to both call a subroutine and return from it. The address of the instruction after the JAL is stored in the specified return address register (Rt) then a jump to the address specified in the instruction is made. The address range is 24 bits or ± 8 MB. The constant field is shifted left twice before use.

The return address register is assumed to be x1 if not otherwise specified. The JAL instruction does not require space in branch predictor tables.

Formats Supported: JAL

Flags Affected: none

Operation:

Rt = IP + 4 $IP = IP + displacement \quad (BAL)$ $OR \qquad IP = displacement \qquad (JAL)$

Execution Units: Branch

Exceptions: none

JALR – Jump and Link to Register

Description:

This instruction may be used to both call a subroutine and return from it. The sum of the current IP and a small constant is stored in the specified return address register (Rt) then a jump to the address specified in the instruction plus an index register value is made.

The return address register is assumed to be x1 if not otherwise specified. The JALR instruction does not require space in branch predictor tables.

If x63 is specified for Ra then the current instruction pointer value is used.

Formats Supported: JALR

Flags Affected: none

Operation:

```
Rt = IP + Cnst_4*2 If Ra=63 IP = IP + displacement Else IP = Ra + Displacement
```

Execution Units: Branch

Exceptions: none

JMP – Jump

Description:

This instruction is an alternate mnemonic for the <u>JAL</u> instruction. It may be used to jump directly to a specific address. The address range is 24 bits or ± 8 MB. The constant field is shifted left twice before use.

The return address register is assumed to be x0 (discarding the return address). The JMP instruction does not require space in branch predictor tables.

Formats Supported: JAL

Flags Affected: none

Operation:

IP = displacement

Execution Units: Branch

Exceptions: none

RET – Return from Subroutine

Description:

This instruction is an alternate mnemonic for the \underline{JALR} instruction. Register Ra is assumed to be x1 and register Rt is assumed to be x0. The constant is assumed to be zero.

Formats Supported: JALR

31	20	19 14	13 10	98	7	0
Const	ant ₁₂	01_{6}	O_4	0_2	42	h_8

Flags Affected: none

Operation:

Execution Units: Branch

Exceptions: an unimplemented instruction exception may occur if a vector register is specified.

Notes:

Return address prediction hardware may make use of the RET instruction.

System Instructions

BRK – Break

Description:

This instruction initiates the processor debug routine. The processor enters debug mode. The cause code register is set to the value specified in the instruction. Interrupts are disabled. The instruction pointer is reset to the contents of tvec[4] and instructions begin executing. There should be a jump instruction placed at the break vector location. The address of the BRK instruction is stored in the EIP register.

Instruction Format: BRK

31 26 25 22 21 14 13 8 7 0
$$\sim_6$$
 \sim_4 Cause₈ 0_6 $00h_8$

Operation:

```
\begin{split} PMSTACK &= (PMSTACK << 4) \mid 10 \\ CAUSE &= Const_8 \\ EIP &= IP \\ IP &= tvec[4] \end{split}
```

Execution Units: Branch

Clock Cycles:

Exceptions: none

CSRx – Control and Special / Status Access

Description:

The CSR instruction group provides access to control and special or status registers in the core. For the read operation the current value of the CSR is placed in the target register Rt.

This instruction is usually used with an extended immediate modifier, however it may be used without the modifier in which case only a read of user CSRs is possible.

Instruction Format: CSR

Op ₃		Operation
0	CSRR	Only read the CSR, no update takes place, Ra should be R0.
1	CSRW	Write to CSR
2	CSRS	Set CSR bits
3	CSRC	Clear CSR bits
4 to 7		reserved

CSRS and CSRC operations are only valid on registers that support the capability.

The Regno_[15..12] field is reserved to specify the operating mode. Note that registers cannot be accessed by a lower operating mode.

Execution Units: Integer, the instruction may be available on only a single execution unit (not supported on all available integer units).

Clock Cycles: 1

Exceptions: privilege violation attempting to access registers outside of those allowed for the operating mode.

PFI – Poll for Interrupt

Description:

The poll for interrupt instruction polls the interrupt status lines and performs an interrupt service if an interrupt is present. Otherwise, the PFI instruction is treated as a NOP operation. Polling for interrupts is performed by managed code. PFI provides a means to process interrupts at specific points in running software.

Instruction Format: SYS

Clock Cycles: 1 (if no exception present)

Execution Units: Branch

REX – Redirect Exception

Description:

This instruction redirects an exception from an operating mode to a lower operating mode. This instruction if successful jumps to the target exception handler and does not return. If this instruction fails execution will continue with the next instruction.

This instruction may fail if exceptions are not enabled at the target level.

The location of the target exception handler is found in the trap vector register for that operating mode (tvec[xx]).

The cause (cause) and bad address (badaddr) registers of the originating mode are copied to the corresponding registers in the target mode.

Instruction Format: REX

```
^{59} ^{58} ^{56} ^{55} ^{48} ^{47} ^{44} ^{4341} ^{40} ^{39} ^{32} ^{31} ^{24} ^{23} ^{16} ^{15} ^{8} ^{7} ^{0} ^{\sim} ^{8} ^{16} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15} ^{15}
```

Tm_3	
0	redirect to user mode
1	redirect to supervisor mode
2	redirect to hypervisor mode
3	redirect to machine mode
4 to 7	not used

Clock Cycles: 4

Execution Units: Branch

Example:

```
REX 1 ; redirect to supervisor handler
; If the redirection failed, exceptions were likely disabled at the target level.
; Continue processing so the target level may complete its operation.

RTE ; redirection failed (exceptions disabled ?)
```

Notes:

Since all exceptions are initially handled in debug mode the debug handler must check for disabled lower mode exceptions.

RTE – Return from Exception

Description:

Restore the previous interrupt enable setting and operating level and transfer program execution back to the address in the exception address register (EIP). One of sixty-four semaphore registers specified by the Rb field of the instruction may also be cleared. Semaphore register zero is always cleared by this instruction.

This instruction may be encoded to return a short distance past the exception address point. This may be useful to return to the next instruction or return to a point past inline parameters. The Ra field specifies a return offset in terms of instruction words.

There is really only a single instruction to return from any mode for an exception. Although there are several additional mnemonics.

Instruction Format: SYS

Flags Affected: none

Operation:

```
PMSTACK = PMSTACK >> 4
Semaphore[0] = 0
Semaphore[Rb] = 0
IP = EIP + Ra
```

Execution Units: Branch

Clock Cycles:

Exceptions: none

SYNC -Synchronize

Description:

All instructions for a particular unit before the SYNC are completed and committed to the architectural state before instructions of the unit type after the SYNC are issued. This instruction is used to ensure that the machine state is valid before subsequent instructions are executed.

Instruction Format:

									23 16 Ra ₈		
~3	UD3	? ?N8	U2	SZ4	III3	Z	KC ₈	KD ₈	Kas	Klջ	44N ₈

TLBRW - Read / Write TLB

Description:

This instruction both reads and writes the TLB. Which translation entry to update comes from the value in Ra. The update value comes from the value in Rb. Rb contains the virtual page number, ASID, and physical page number. The current value of the entry selected by Ra is copied to Rt. The TLB will be written only if bit 63 of Ra is set.

The entry number for Ra comes from virtual address bits 14 to 23.

Page numbers are in terms of a 16kB page size.

Instruction Format: SYS

Clock Cycles: 5

Execution Units: Memory

Ra Value Format

63	62	11 10	9 0
W	~	way	entry no

Rb/Rt Value Format

63	56	55	54	53	52	48	47		32	31	20	19		0
AS	ID	G	D	Α	UCR	RWX		VPN			~		PPN	

Bits		Meaning						
0 to 19	PPN	Physical page number						
20 to 31	~	reserved (expansion of physical p	page number)					
32 to 49	VPN	Virtual page number high address	Firtual page number high address order bits 24 to 39					
48	X	1 = page is executable	These three combined indicate					
49	W	1 = page is writeable	page present (P) $0 = \text{not present}$					
50	R	1 = page is readable						
51	C	1 = page is cachable	1 = page is cachable					
52	U	reserved for system usage						
53	A	Accessed, set if translation was u	sed					
54	D	Dirty, set if a write occurred to the	ne page					
55	G	Global, global translation indicate	or					
56 to 63	ASID	ASID address space identifier						

WFI – Wait for Interrupt

Description:

The WFI instruction waits for an external interrupt to occur before proceeding. While waiting for the interrupt, the processor clock is stopped placing the processor in a lower power mode.

Instruction Format: SYS

Clock Cycles: 1 (if no exception present)

Execution Units: Branch

Vector Specific Instructions

Arithmetic / Logical

V2BITS

Synopsis

Convert Boolean vector to bits.

Description

The least significant bit of each vector element is copied to the corresponding bit in the target register. The target register is a scalar register.

Instruction Format: R1

Operation

For
$$x = 0$$
 to VL-1
$$if (Vm[x])$$

$$Rt[x] = Va[x].LSB$$

$$else \ if \ (z)$$

$$Rt[x] = 0$$

VBITS2V

Synopsis

Convert bits to Boolean vector.

Description

Bits from a general register are copied to the corresponding vector target register.

Operation

For
$$x = 0$$
 to VL-1
$$if (Vm[x]) Vt[x] = Ra[x]$$

VCIDX – Compress Index

Synopsis

Vector compression.

Description

A value in a register Ra is multiplied by the element number and copied to elements of vector register Vt guided by a vector mask register.

Operation

$$y = 0$$
 for $x = 0$ to $VL - 1$
$$if (Vm[x])$$

$$Vt[y] = Ra * x$$

$$y = y + 1$$

VCMPRSS – Compress Vector

Synopsis

Vector compression.

Description

Selected elements from vector register Va are copied to elements of vector register Vt guided by a vector mask register.

Operation

$$y = 0$$
 for $x = 0$ to $VL - 1$
$$if (Vm[x])$$

$$Vt[y] = Va[x]$$

$$y = y + 1$$

VEINS / VMOVSV – Vector Element Insert

Synopsis

Vector element insert.

Description

A general-purpose register Rb is transferred into one element of a vector register Vt. The element to insert is identified by Ra.

Operation

Vt[Ra] = Rb

VEX / VMOVS – Vector Element Extract

Synopsis

Vector element extract.

Description

A vector register element from Vb is transferred into a general-purpose register Rt. The element to extract is identified by Ra.

Operation

Rt = Vb[Ra]

VSCAN

Synopsis

.

Description

Elements of Vt are set to the cumulative sum of a value in register Ra. The summation is guided by a vector mask register.

Operation

```
sum = 0 for x = 0 to VL - 1 Vt[x] = sum if (Vm[x]) sum = sum + Ra
```

VSLLV – Shift Vector Left Logical

Description

Elements of the vector are transferred upwards to the next element position. The first is loaded with the value zero. This is also called a slide operation.

Operation

For
$$x = VL-1$$
 to Amt
$$Vt[x] = Va[x-amt]$$
 For $x = Amt-1$ to 0
$$Vt[x] = 0$$

VSRLV – Shift Vector Right Logical

Description

Elements of the vector are transferred downwards to the next element position. The last is loaded with the value zero. This is also called a slide operation.

Operation

For
$$x = 0$$
 to VL-Amt
$$Vt[x] = Va[x+amt]$$
 For $x = VL-Amt+1$ to VL-1
$$Vt[x] = 0$$

Memory Operations

CVLDx – Compressed Vector Load

Description:

Formats Supported:

Stridden Form

```
63
                           4948
                                 47 44
                                                          39
                                                                 32
                                                                        31 24
                                                                                     23 16
                                                                                                 15 8
      Const<sub>21..8</sub>
                            U_2
                                    Sz_4
                                             m_3
                                                     Z
                                                          Const<sub>7..0</sub>
                                                                         Rb_8
                                                                                      Ras
                                                                                                   Rt_8
                                                                                                               65h_8
```

Data is loaded from memory locations beginning at the sum of Ra and a constant and separated by the stride amount in the stride register Rb. Rb may also be a constant in the range -62 to 63. If Rb = -63 then the Sz₄ field is used to determine the stride.

Operation:

```
y = 0
for x = 0 to vector length
        if Rb is a constant
                if Rb = -63
                         stride = Sz4
                else
                         stride = Rb
        else
                stride = [Rb]
        n = stride * y
        if (Vm[x])
                Vt[y] = Memory[d+Ra + n]
                y = y + 1
for y = y to vector length
        Vt[y] = z ? 0 : Vt[y]
n = 0
```

If the vector mask bit is clear and the 'z' bit is set in the instruction then the corresponding element of the vector register is loaded with zero. If the vector mask bit is clear and the 'z' bit is clear in the instruction then the corresponding element of the vector register is left unchanged (no value is loaded from memory).

Elements are loaded only up to the length specified in the vector length register.

```
Vm[x] z Result
```

Operation:

```
\begin{split} n &= 0 \\ y &= 0 \\ \text{for } x &= 0 \text{ to vector length} \\ &\quad \text{if } (Vm[x]) \\ &\quad Vt[y] = Memory[d+Ra+n] \\ &\quad n = n + size of \text{ precision} \\ &\quad y = y+1 \\ \text{for } y = y \text{ to vector length} \\ &\quad Vt[y] = z ? 0 : Vt[y] \end{split}
```

Indexed Form

Data is loaded from memory addresses beginning with the sum of Ra and a vector element from Vb.

Operation:

```
\begin{aligned} y &= 0 \\ \text{for } x &= 0 \text{ to vector length} \\ &\quad \text{if } (Vm[x]) \\ &\quad Vt[y] &= Memory[d + Ra + Vb[x]] \\ &\quad y &= y + 1 \\ \text{for } y &= y \text{ to vector length} \\ &\quad Vt[y] &= z ? 0 : Vt[y] \end{aligned}
```

CVSTx – Compressed Vector Store

Description:

Formats Supported:

Register Indirect with Displacement

Data is stored to consecutive memory addresses beginning with the sum of Ra and an immediate

Elements are stored only up to the length specified in the vector length register.

```
4140
                    39 36
                             35 33
                                                                              19 14
                                     32
                                            31
                                                                               Ra_6
                                                                                                    1
                                                                                                           74h_7
Const<sub>6</sub>
             U_2
                     Sz_4
                                                    Constant<sub>12</sub>
                                                                                          V_{S_6}
                              m_3
Vm[x]
                     Result
              \mathbf{Z}
   1
              0
                     memory = Vs[x]
   1
                     memory = Vs[x], Vs[x] = 0
```

Operation:

```
\begin{split} n &= 0 \\ \text{for } x &= 0 \text{ to vector length} \\ &\quad \text{if } (Vm[x]) \\ &\quad \text{Memory}[d + Ra + n] = Vs[x] \\ &\quad \text{if } (z) \ Vs[x] = 0 \\ &\quad n = n + \text{sizeof precision} \end{split}
```

Stridden Form

The stridden form works much the same as the register indirect form except that data is stored to memory locations separated by the stride amount in the stride register.

```
4140
                                                                 25 20
                                                                              19 14
                   39 36
                              35 33
                                       32
                                                31
                                                        26
                                                                                           13 8
Const<sub>6</sub>
            U_2
                    Sz_4
                               m_3
                                        Z
                                                 Const<sub>6</sub>
                                                                   Rb_6
                                                                                Ra_6
                                                                                            V_{S_6}
                                                                                                     1
```

Operation:

```
y = 0 for x = 0 to vector length n = Rb * y if (Vm[x]) Memory[d+Ra + n] = Vs[x] if (z) Vs[x] = 0 y = y + 1
```

Indexed Form

Data is stored to memory addresses beginning with the sum of Ra and a vector element from Vb.

Operation:

```
\begin{aligned} y &= 0 \\ \text{for } x &= 0 \text{ to vector length} \\ &\quad \text{if } (Vm[x]) \\ &\quad \text{Memory}[d + Ra + Vb[y]] = Vs[x] \\ &\quad \text{if } (z) \ Vs[x] = 0 \\ &\quad y &= y + 1 \end{aligned}
```

Root Opcode Map

rtoot O _I			010	011	100	101	110	111			
	000	001	010	011	100	101	110	111			
00000	BRK	{R1}	{R2}	ALU {R3}	ADD	SUBF	MUL				
00000	AND	OR	XOR	{K5}	ADD	{SET}	MULU	CSR			
00001	DIV	DIVU	DIVSU			MULF	MULSU	PERM			
00010			BYTNDX	WWDNDV	(DTELD)	MULF	MULSU	PERM			
00011	REM	REMU		WYDNDX	{BTFLD}	a a a	ano	G) III			
00100	REMSU	DIVR	CHK	U21NDX	SAND	SOR	SEQ	SNE			
00101	SLT	SGT	SLTU	SGTU							
00110	MADD	MSUB	NMADD	NMSUB				FDP			
00111								NOP			
				Branch Unit							
01000	JAL	BAL	JALR		{SYS}						
01001	BLT	BGE	BLTU	BGEU		BBS	BEQ	BNE			
				on Modifiers	(Prefixes)						
01010	EXI	EXI	EXI								
01011	IMOD	BTFLD	BRMOD	FB	STRIDE						
Memory Unit											
01100	LDx	LDxX					CVLDS	CVLDVX			
01101								LSM			
01110	STx	STxX					CVSTS	CVSTVX			
01111											
				Vector ALU							
10000		{R1}	{R2}	{R3}	ADD	SUBF	MUL				
10001	AND	OR	XOR				MULU				
10010	DIV	DIVU	DIVSU			MULF	MULSU	PERM			
10011	REM	REMU	BYTNDX	WYDNDX	{BTFLD}						
10100	REMSU	DIVR	CHK	U21NDX			SEQ	SNE			
10101	SLT	SGT	SLTU	SGTU							
10110	MADD	MSUB	NMADD	NMSUB				FDP			
10111								NOP			
11000											
11001											
11010											
11011	IMOD	BTFLD	BRMOD	FB							
	1	·		1			1				
11100			LDSx	LDxVX							
11101											
11110			STSx	STxVX							
11111											

(aba)	, m . 1.	D ' /	
{ SR 3 }	1 r1ad1c	Register	Uns
	IIIaaic	110815101	

	000	001	010	011	100	101	110	111
000					MUX			
001								
010	SLLP	SLLPI						
011	PTRDIF							
100			CHK					
101								
110	BLEND							
111								

{SR2} Dyadic Register Ops

	000	001	010	011	100	101	110	111
000	AND	OR	XOR	BMM	ADD	SUB	MUL	
001	NAND	NOR	XNOR			MULF	MULU	MULH
010	DIV	DIVU	DIVSU	REM	REMU	REMSU	MULSU	PERM
011	DIF	SLL	SLLI		MULF	MULSUH	MULUH	RGF
100							SEQ	SNE
101	MIN	MAX	CMP		SLT	SGE	SLTU	SGEU
110								
111			VEX	VEINS				

{SR1} Monadic Register Ops

	000	001	010	011	100	101	110	111
00	CNTLZ	CNTLO	CNTPOP	COM	NOT	NEG	ABS	
01				TST				
10	PTRINC							
11	V2BITS	BITS2V	VSLLV	VSLRV	VCMPRSS	VCIDX	VSCAN	

{OSR2} System Ops

	000	001	010	011	100	101	110	111
00	LLAL	LLAH			LPAL	LPAH		
01	PUSHQ	POPQ	PEEKQ	STATQ	SETKEY	GCCLR		
10	REX	PFI	WAI	RTE				
11	SETTO	GETTO	GETZL		MVMAP	MVSEG	TLBRW	SYNC