

# On Traffic Analysis in Tor

*Guest Lecture, ELE 574*

*Communications Security and Privacy*

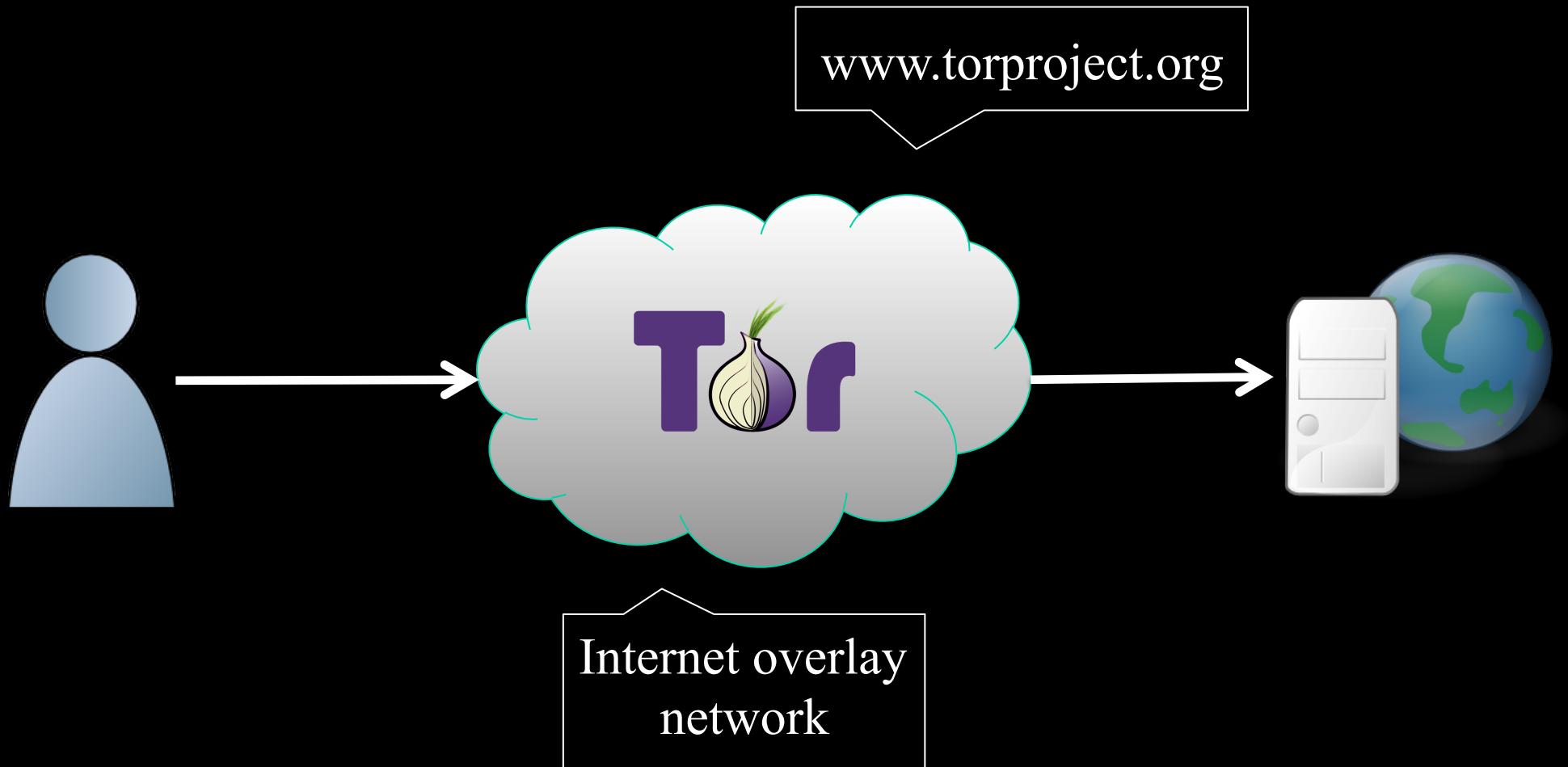
*Princeton University*

*April 3<sup>rd</sup>, 2014*

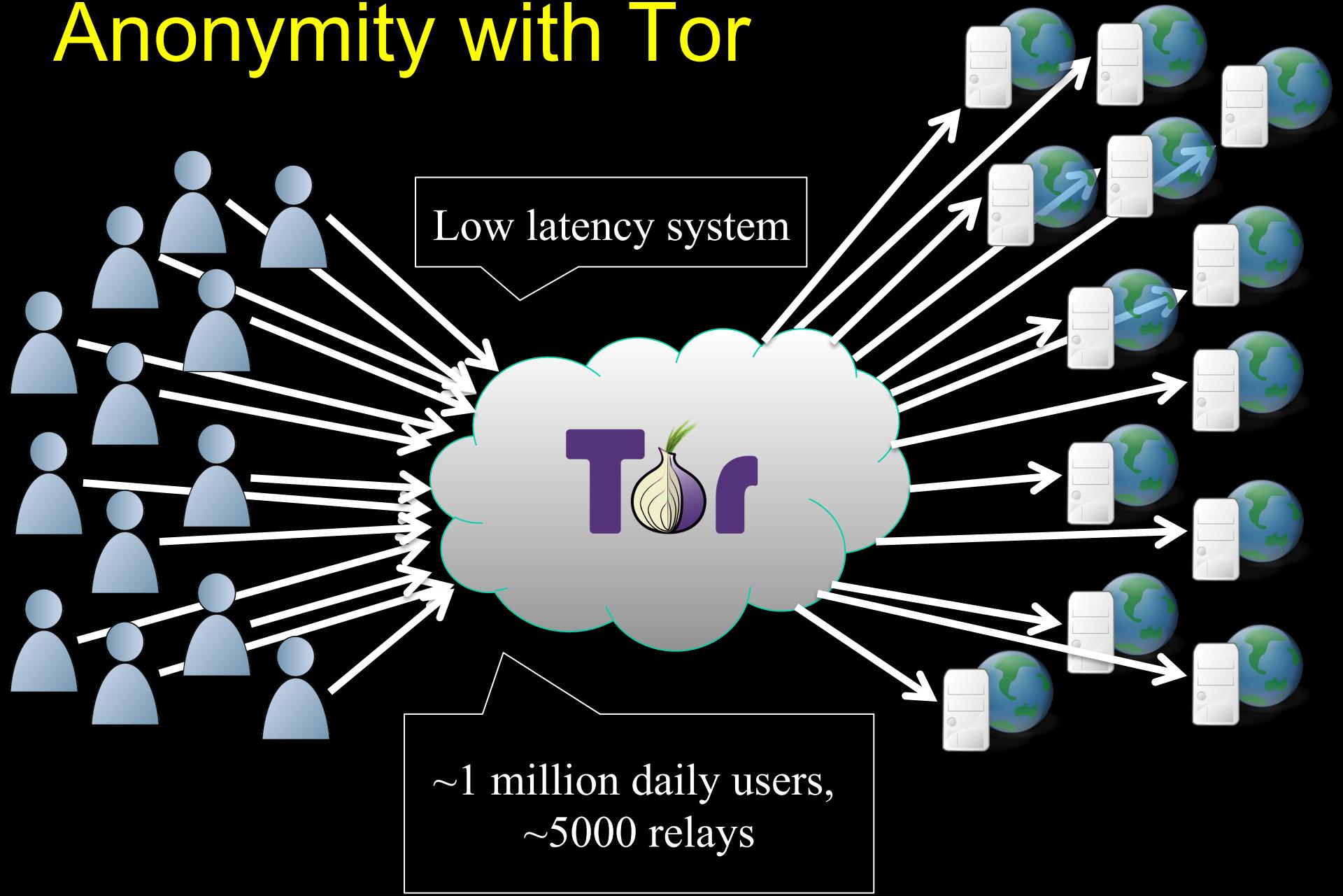


Dr. Rob Jansen  
U.S. Naval Research Laboratory  
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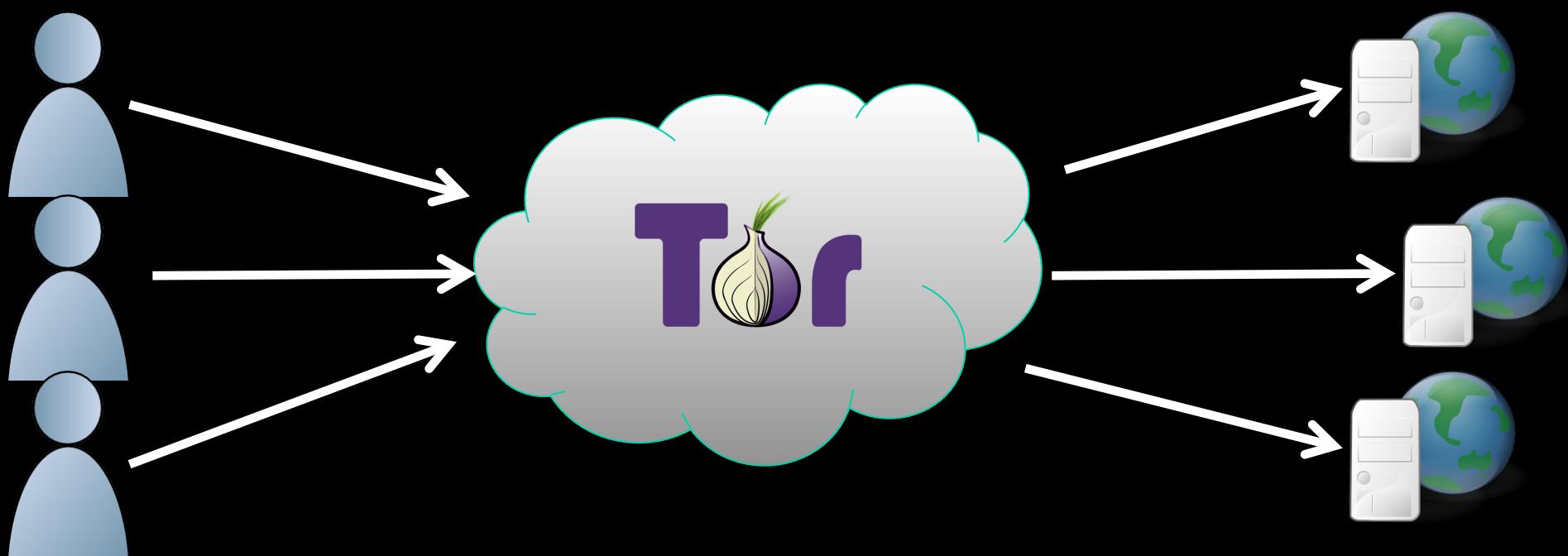
# Anonymity with Tor



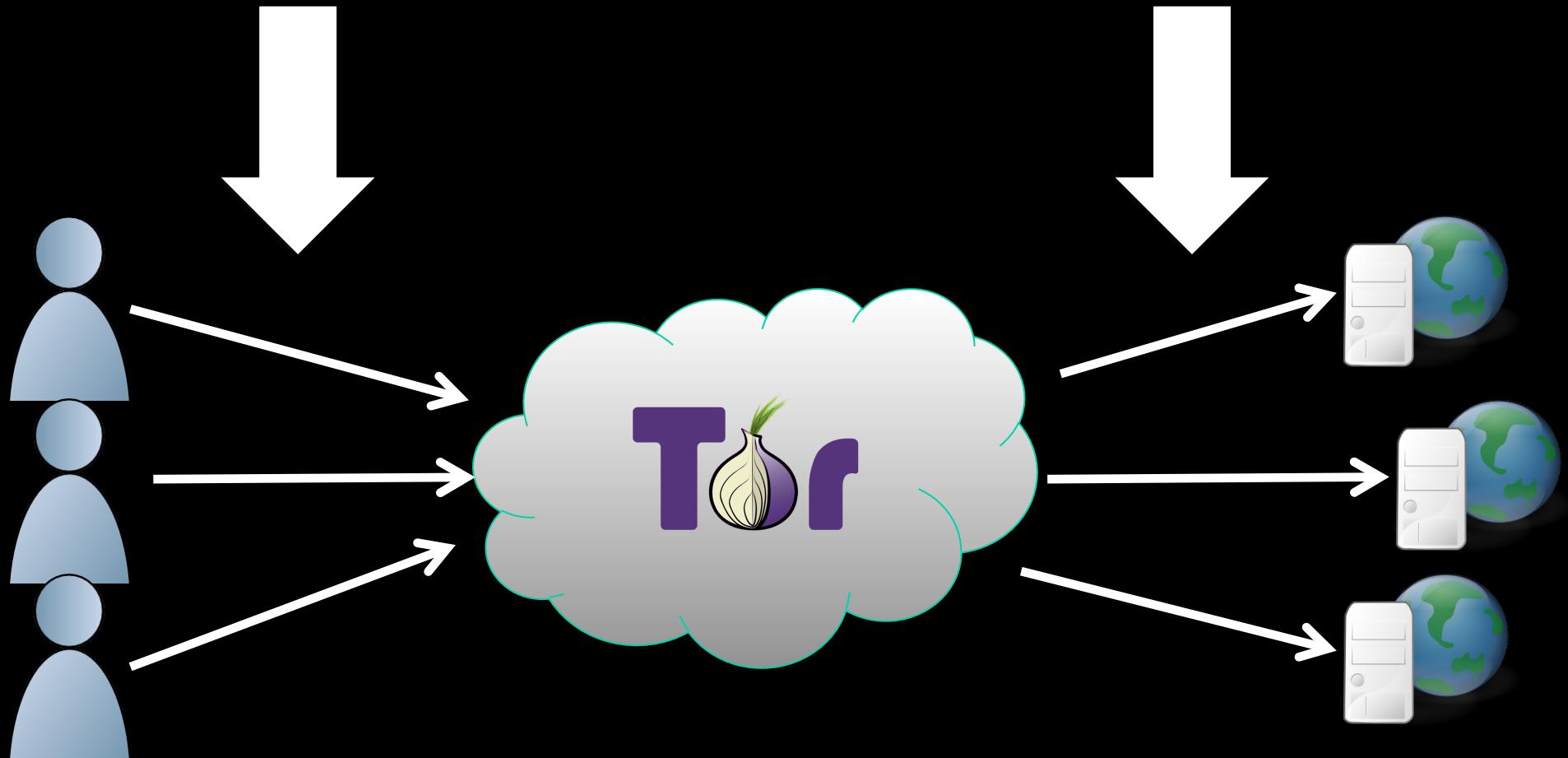
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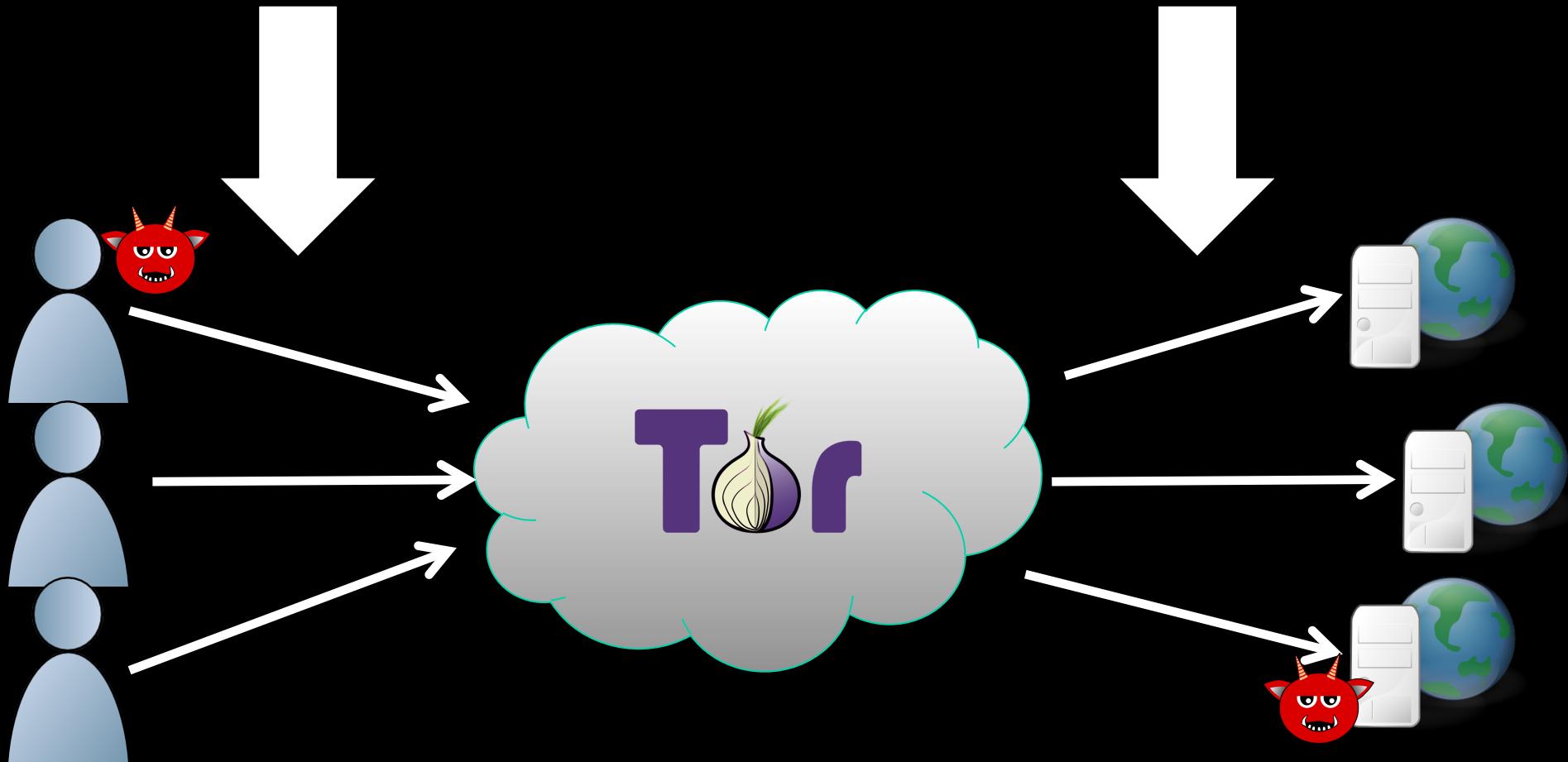
# Traffic Correlation



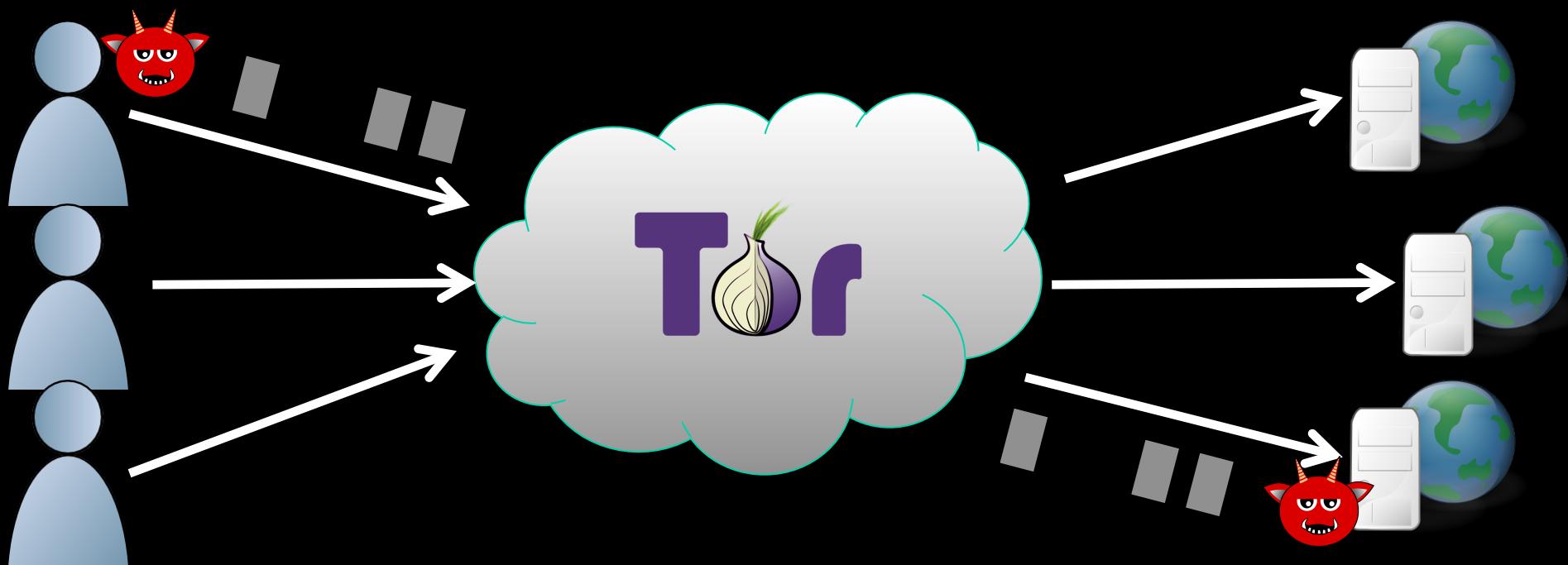
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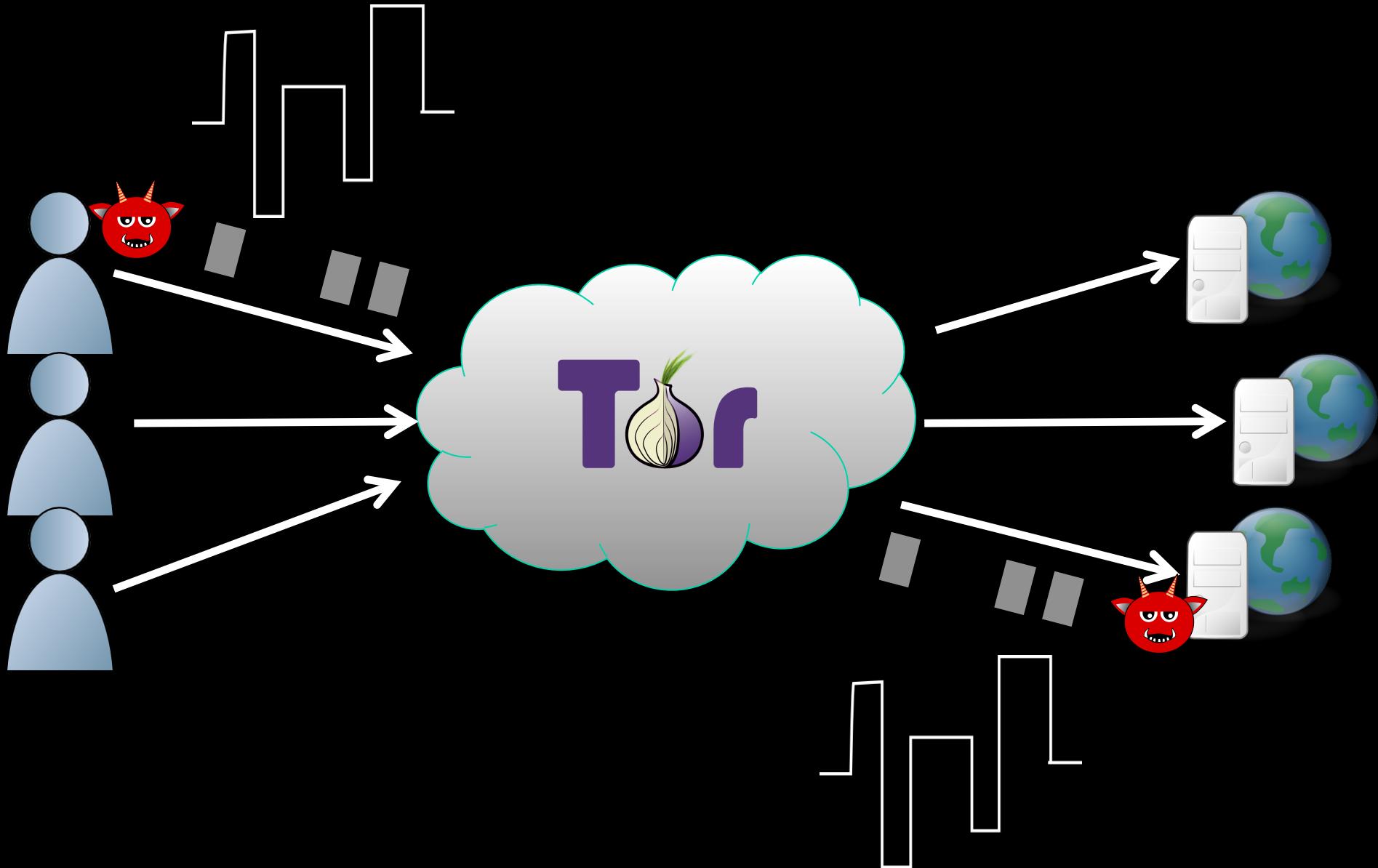
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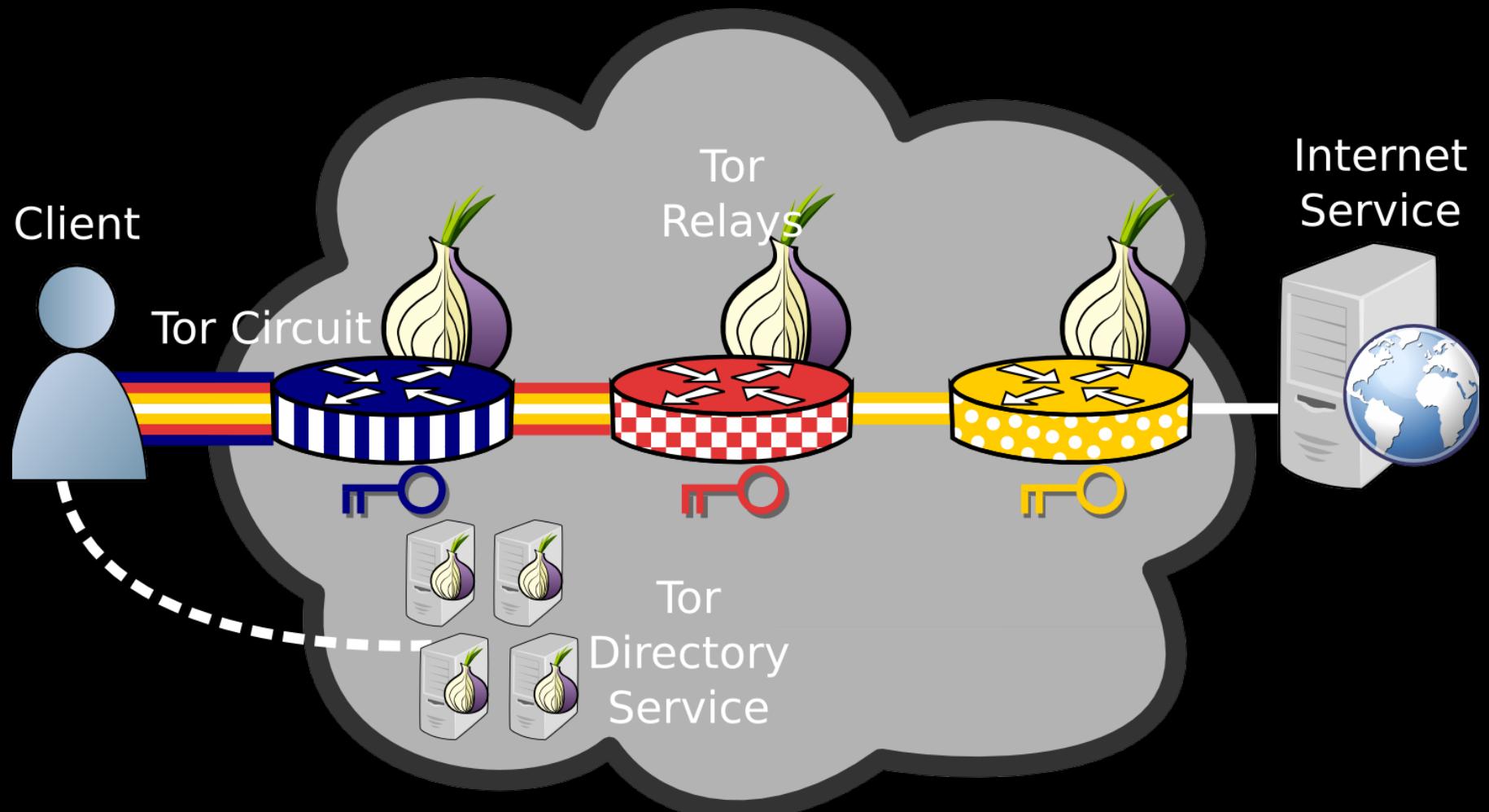


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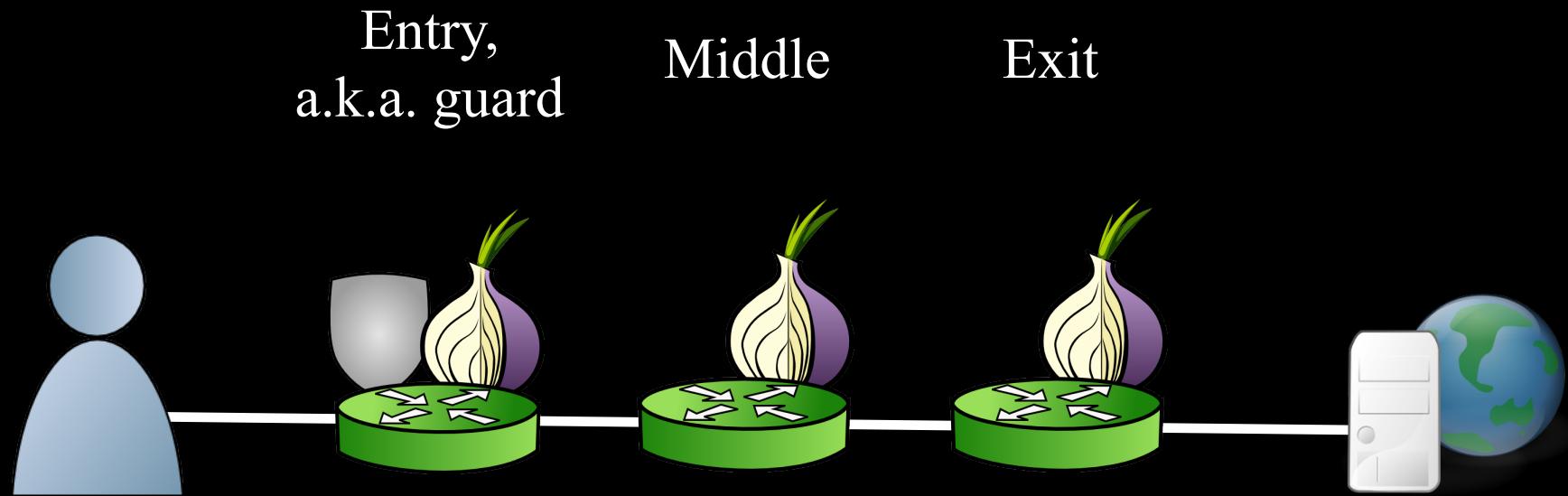
- Is traffic correlation realistic?
- Who might be in these positions?
- Would a nation-state be willing to launch correlation attacks?

The biggest threat  
to Tor's anonymity

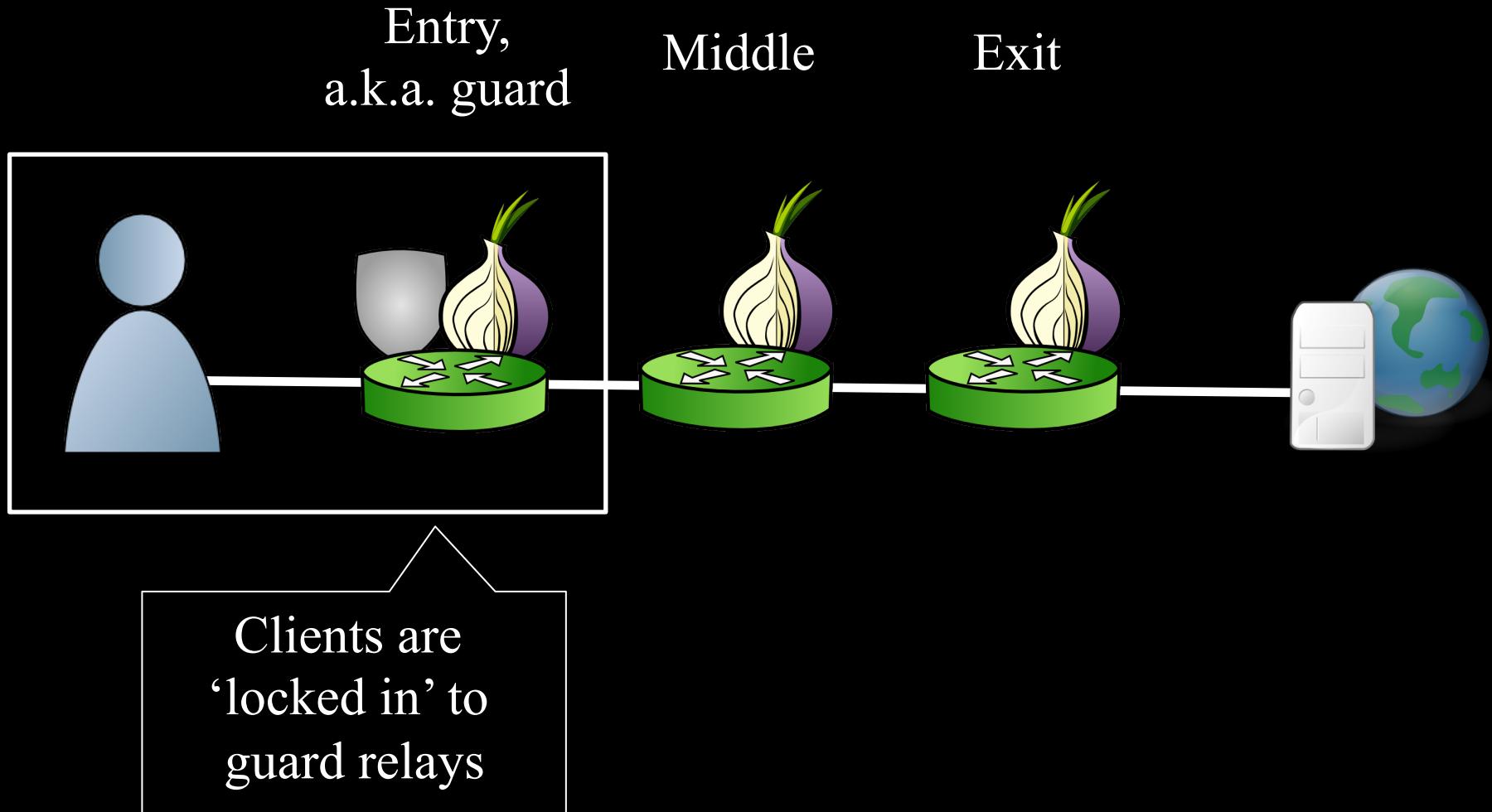
# Anonymity with Onion Routing



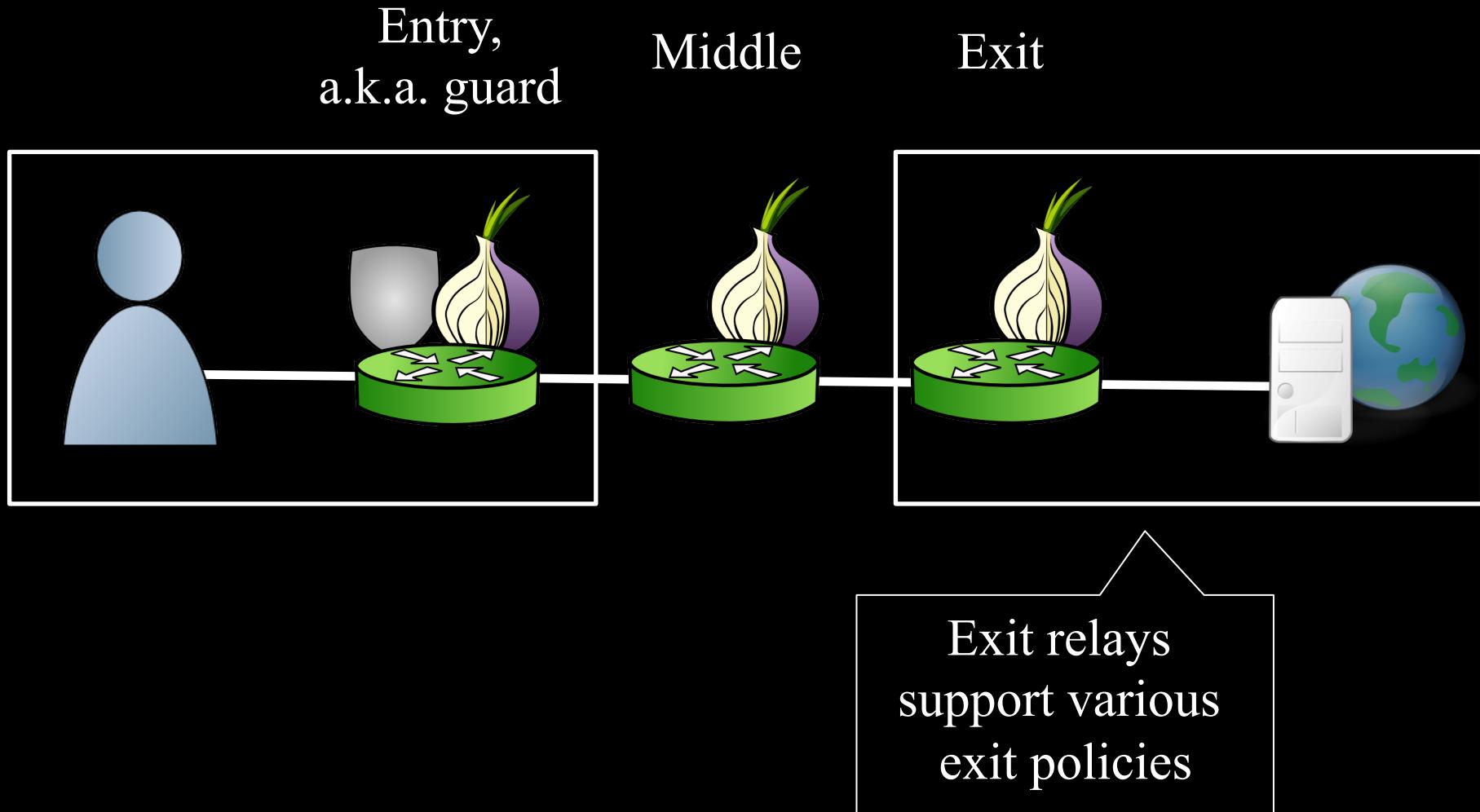
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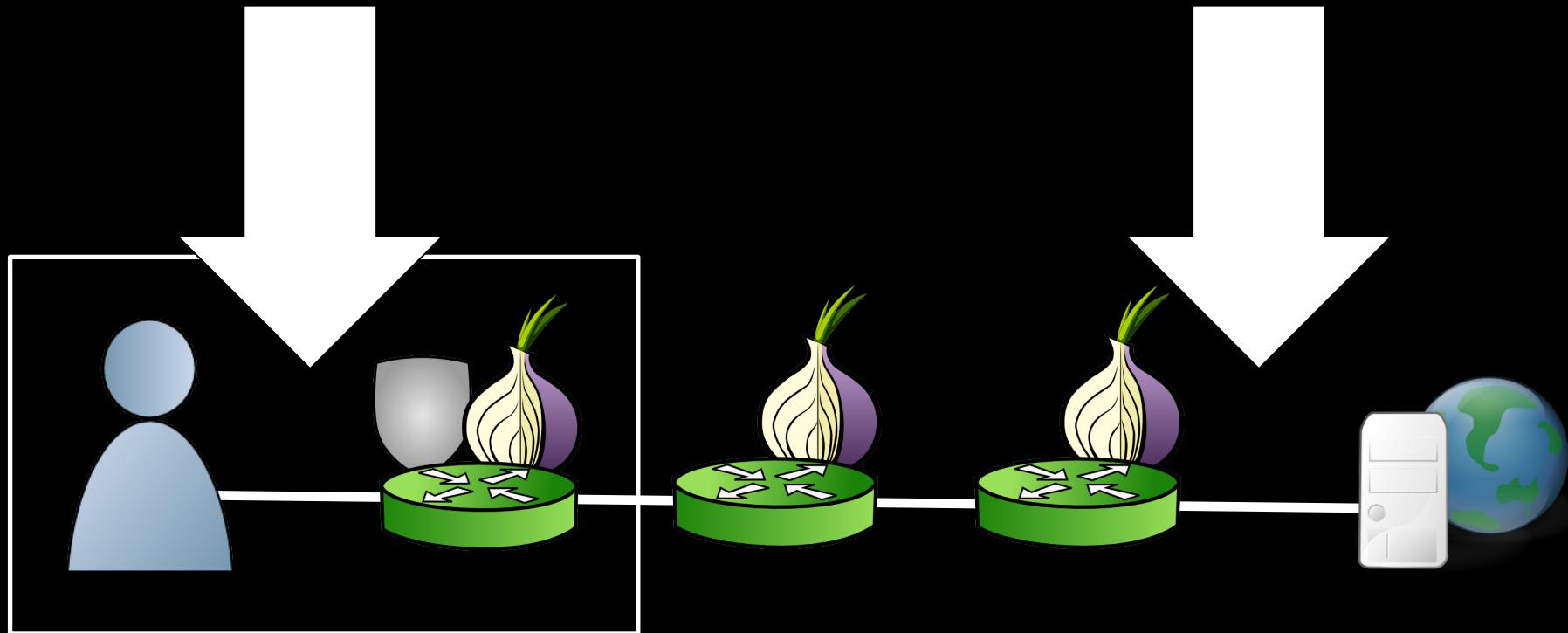
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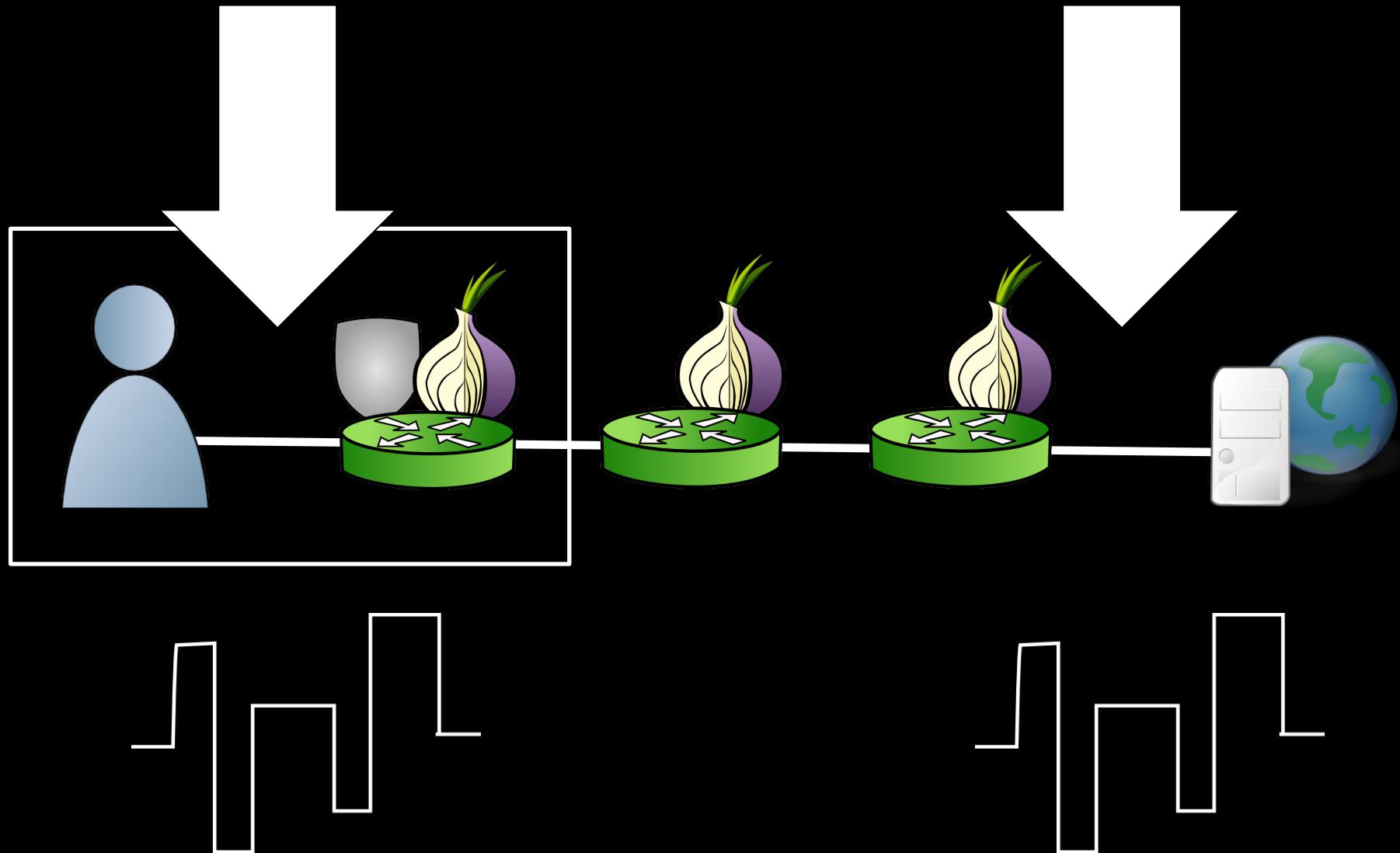
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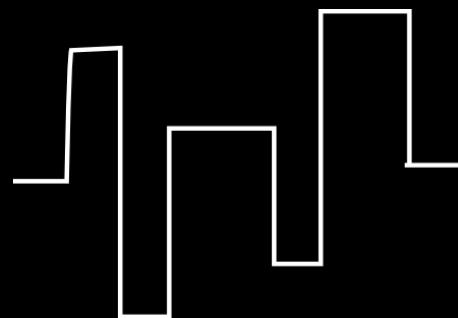
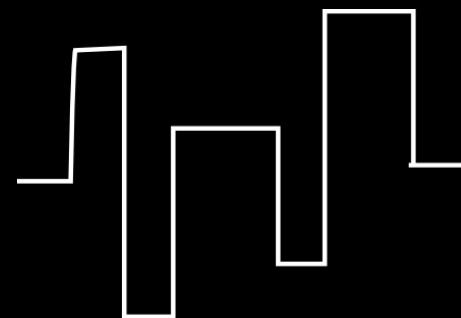


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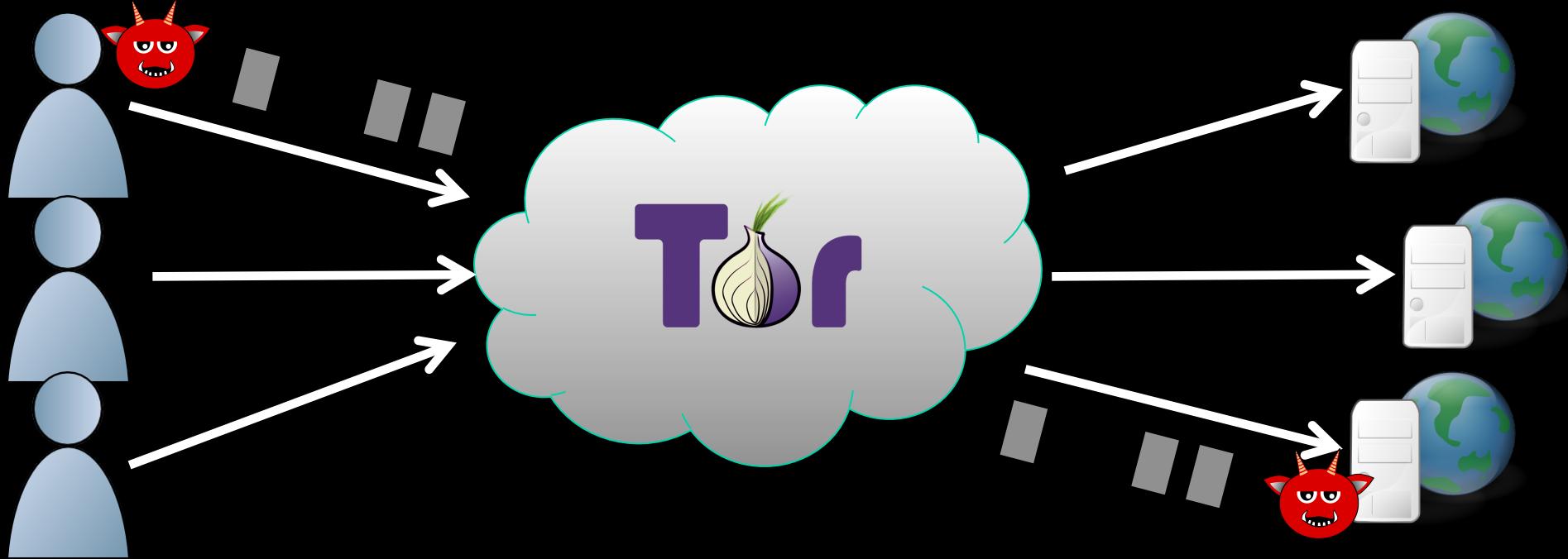
- How does the volunteer resource model affect the vulnerability to correlation attacks?



# Outline

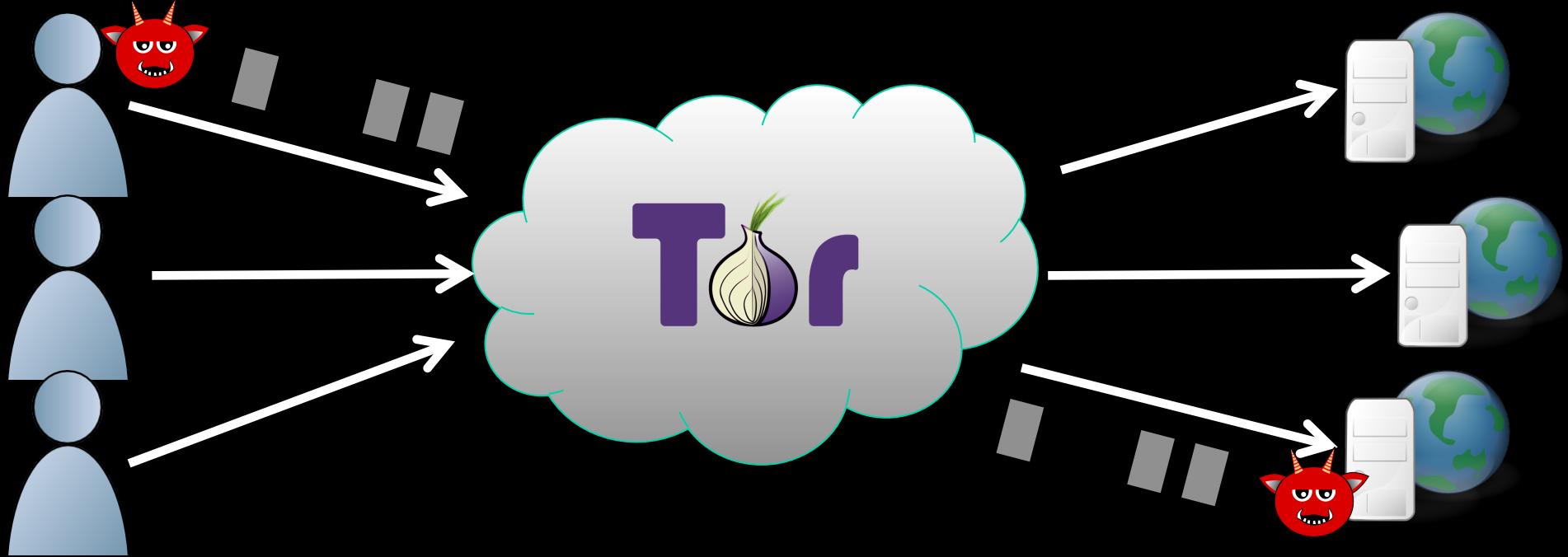
- ~~Background~~
- Security against correlation (end-to-end)
  - Metrics and methodology
  - Node adversaries
  - Link adversaries
- Correlation attacks (partial)
  - Stealthy throughput
  - Induced throttling
    - . Traffic admission control
    - . Congestion control

# Traffic Correlation



- How can one measure how vulnerable real clients on the real network are to traffic correlation?

# Traffic Correlation



- Is there a difference between targeted correlation and general surveillance?

# Security Metrics

## Principles

- Probability distribution
- Measured on human timescales
- Based on real network and adversaries

# Security Metrics

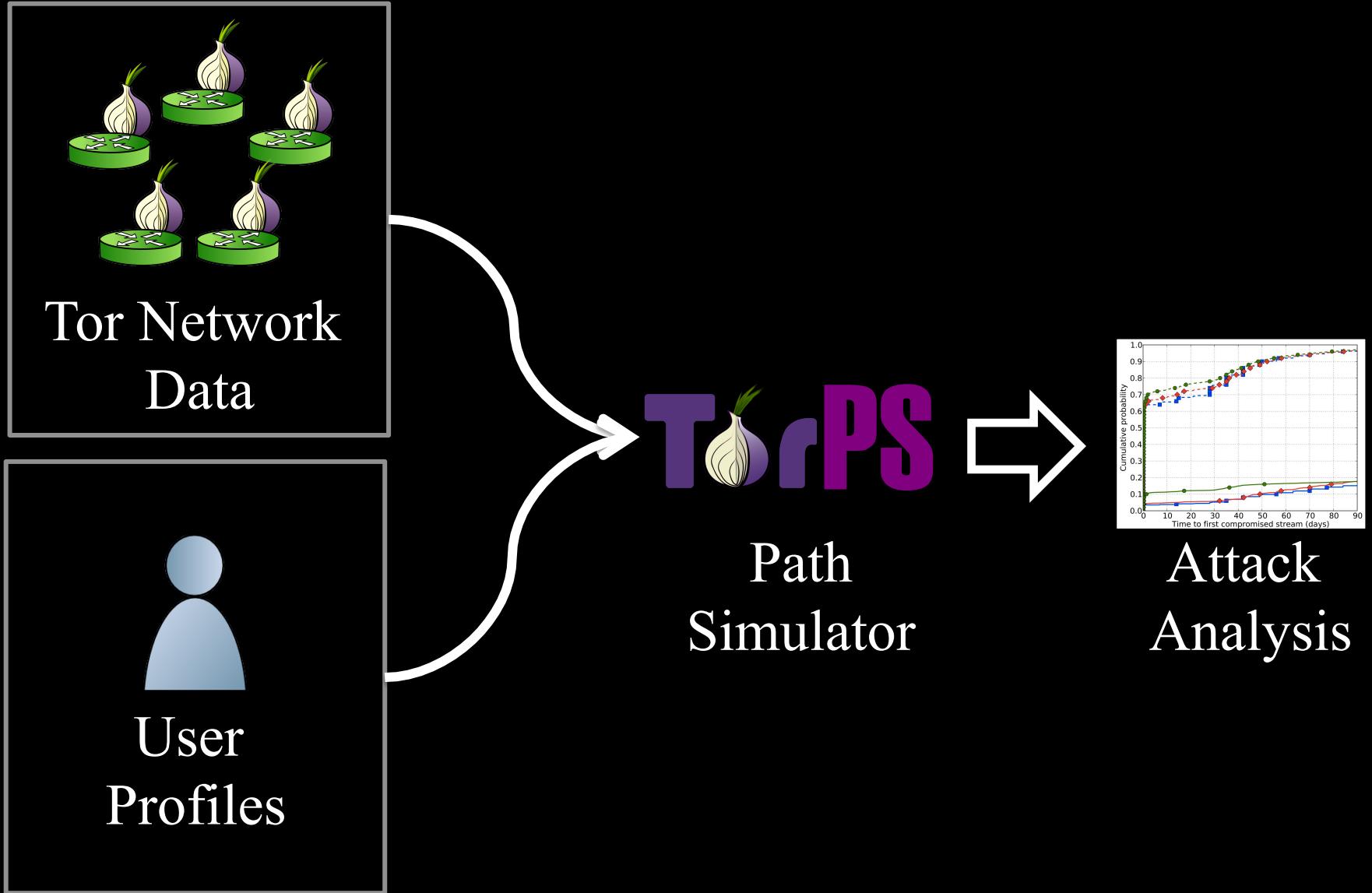
## Principles

- Probability distribution
- Measured on human timescales
- Based on real network and adversaries

## Metrics (Probability distributions)

- Time until first path compromise
- Number of path compromises for a given user over given time period

# Approach: Overview



# Approach: User Profiles

Consider how users actually use Tor

## Typical



Gmail/GChat



GCal/GDocs



Facebook



Web search

## Chat



IRC



BitTorrent

Build a 20-minute trace  
of each activity.  
Capture destinations/  
ports visited

# Approach: User Profiles

“Replay” traces to generate streams based on user behavior

Typical	Chat	File Sharing
<ul style="list-style-type: none"><li>• 2632 traces per week</li><li>• 205 destinations</li><li>• 2 ports</li></ul>	<ul style="list-style-type: none"><li>• 135 traces per week</li><li>• 1 destinations</li><li>• 1 port</li></ul>	<ul style="list-style-type: none"><li>• 6768 traces per week</li><li>• 171 destinations</li><li>• 118 ports</li></ul>

# Approach: User Profiles

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## Typical

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- 2 ports

## Chat

- 135 traces per week
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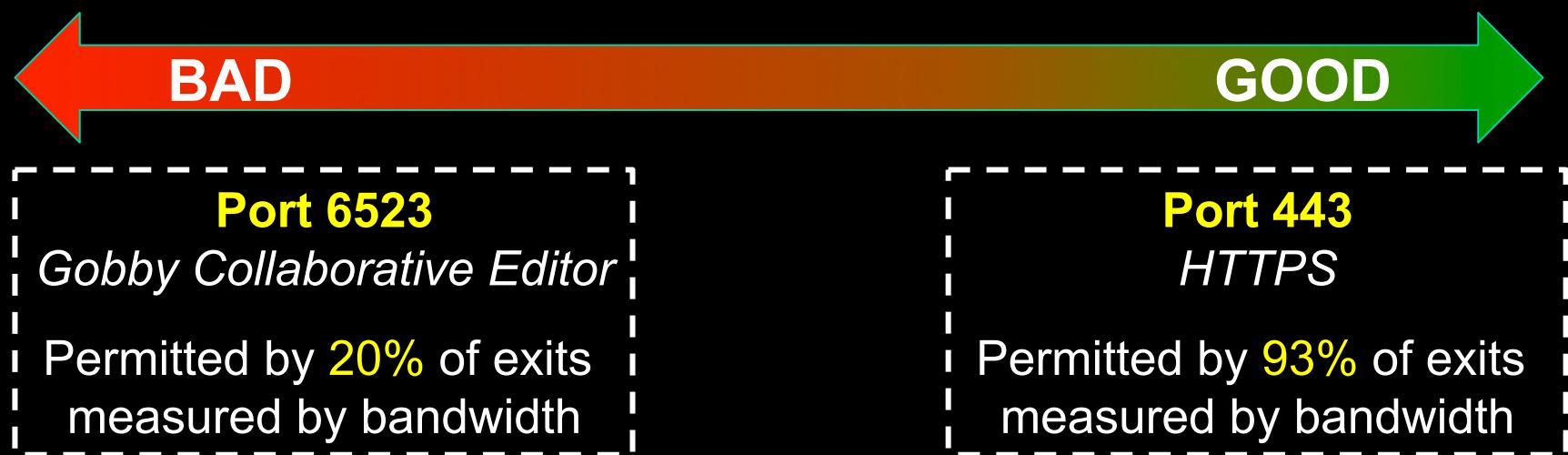
## File Sharing

- 6768 traces per week
- 171 destinations
- 118 ports

Is the user model accurate?  
What are the challenges?

# User Behavior Affects Relay Selection

Some applications are not well-supported by Tor due to exit policies



# Approach: Tor Network Data

Consider the Tor network as it changes over a long period of time:

- Relays join and leave
- Bandwidth changes
- Exit/Guard designations change

Use Tor Project archives to obtain state of network over 3 to 6 months



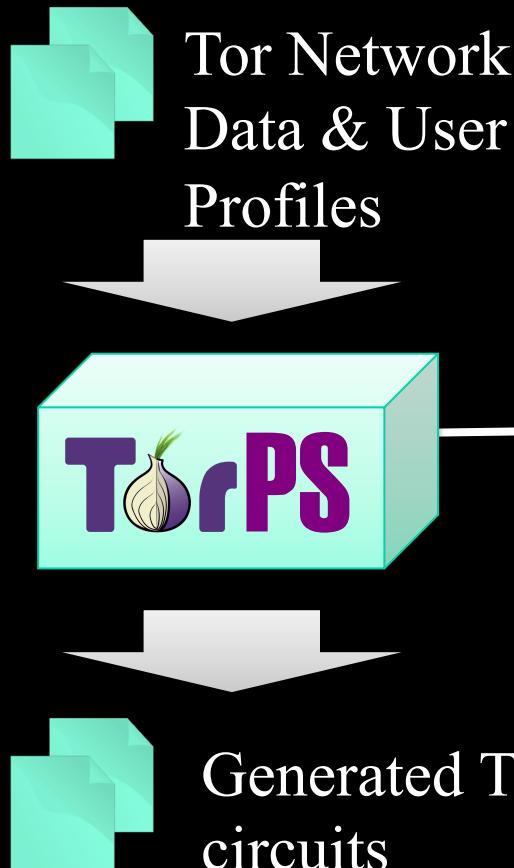
Hourly  
consensuses



Monthly server  
descriptors

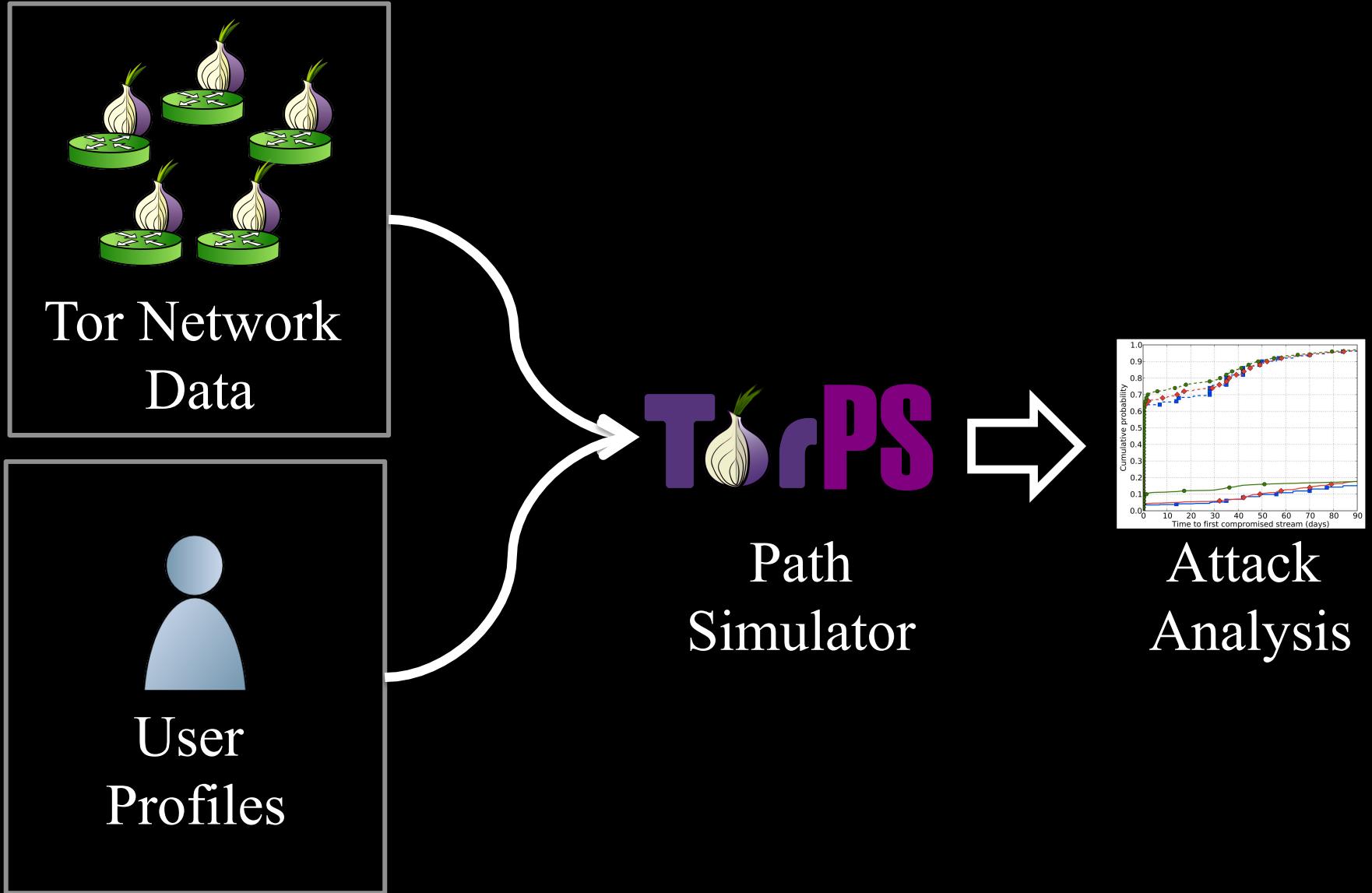
# Approach: Simulate Tor with TorPS

Combine User and Tor Network models using TorPS to produce the circuits Tor would use



- Re-implements path selection
- Based on Tor stable version (0.2.3.25)
- Considers:
  - Bandwidth weighting
  - Exit policies
  - Guards and guard rotation
  - Hibernation
  - /16 and family conflicts
- Omits effects of network performance

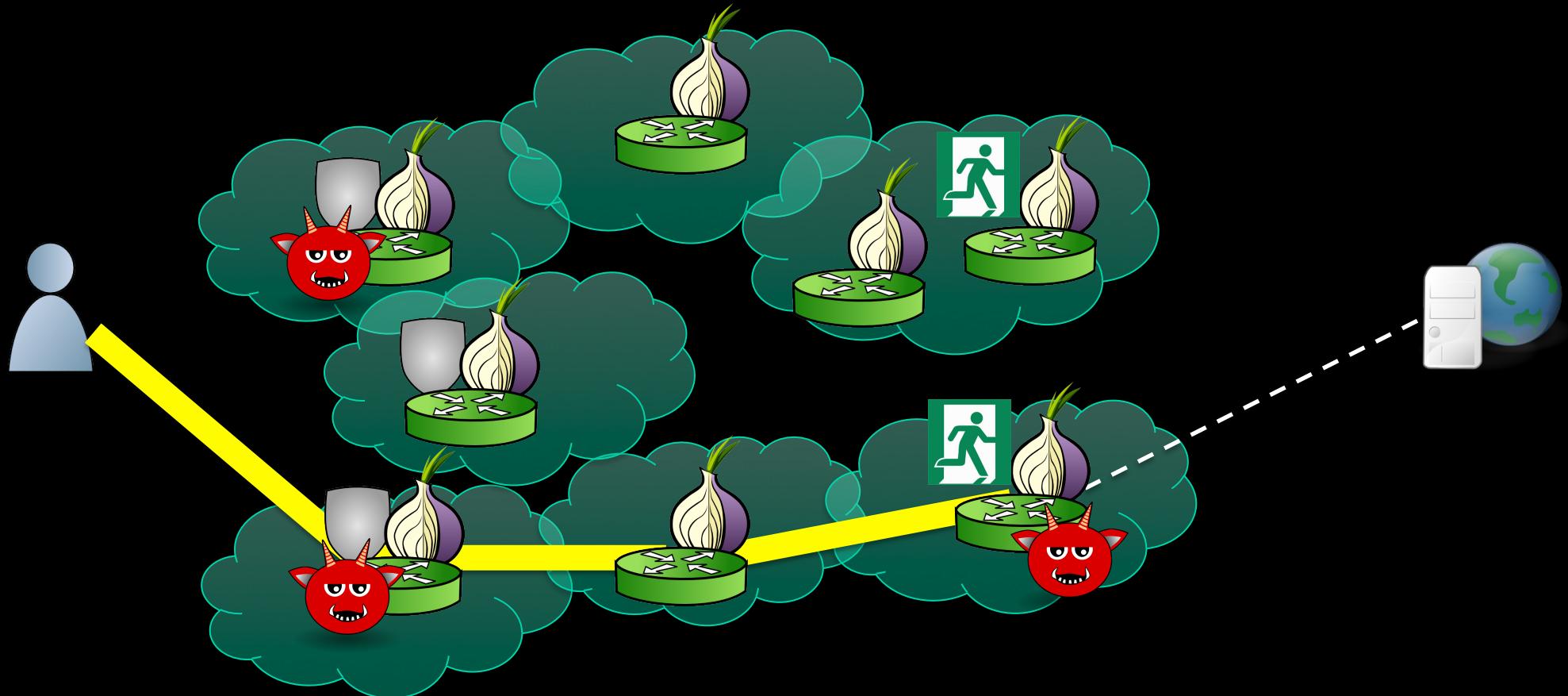
# Approach: Overview



# Outline

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  - ~~Metrics and methodology~~
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# Node Adversary



# Node Adversary

Controls a fixed allotment of relays based on bandwidth budget

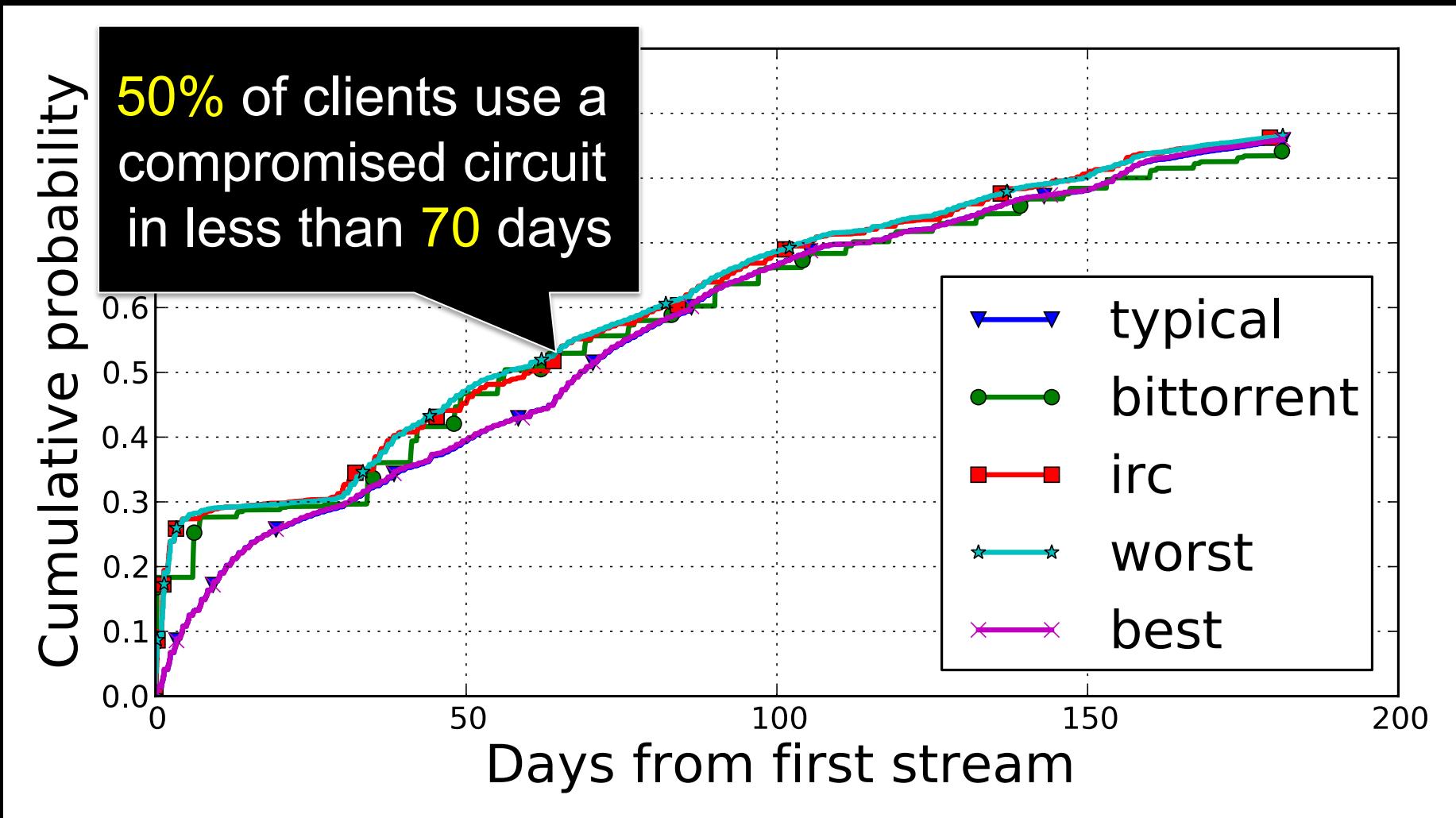
- We assume adversary has 100 MiB/s – comparable to large family of relays
- Adversaries apply 5/6th of bandwidth to **guard** relays and the rest to **exit** relays. (We found this to be the most effective allocation we tested.)

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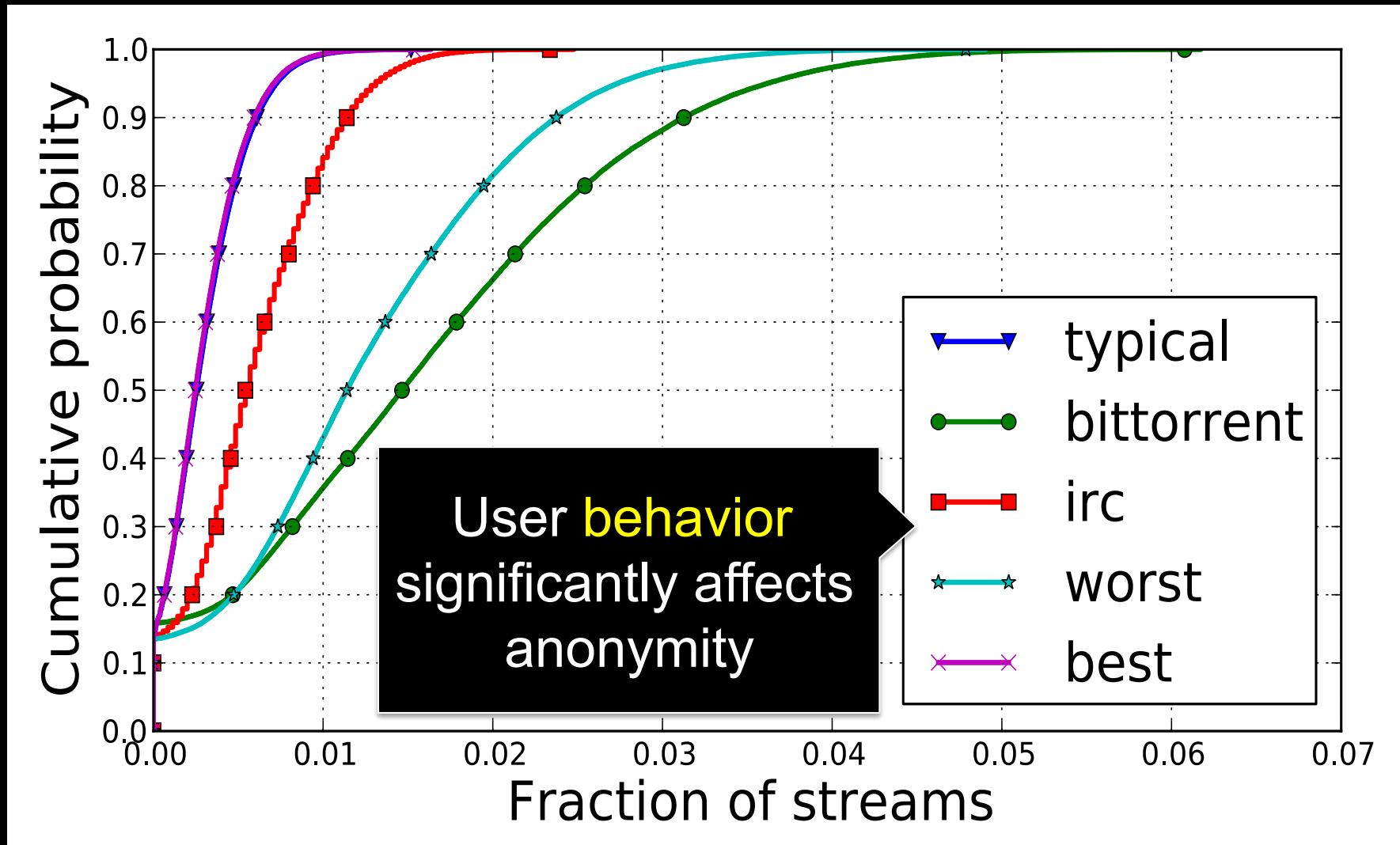
- We assume adversary has 100 MiB/s – comparable to large family of relays
  - Is 100 MiB/s realistic for an adversary?
  - Adversaries apply 5/6th of bandwidth to guard relays and the rest to exit relays. (We found this to be the most effective allocation we tested.)

# Time to First Compromised Circuit



October 2012 – March 2013

# Fraction of Compromised Streams

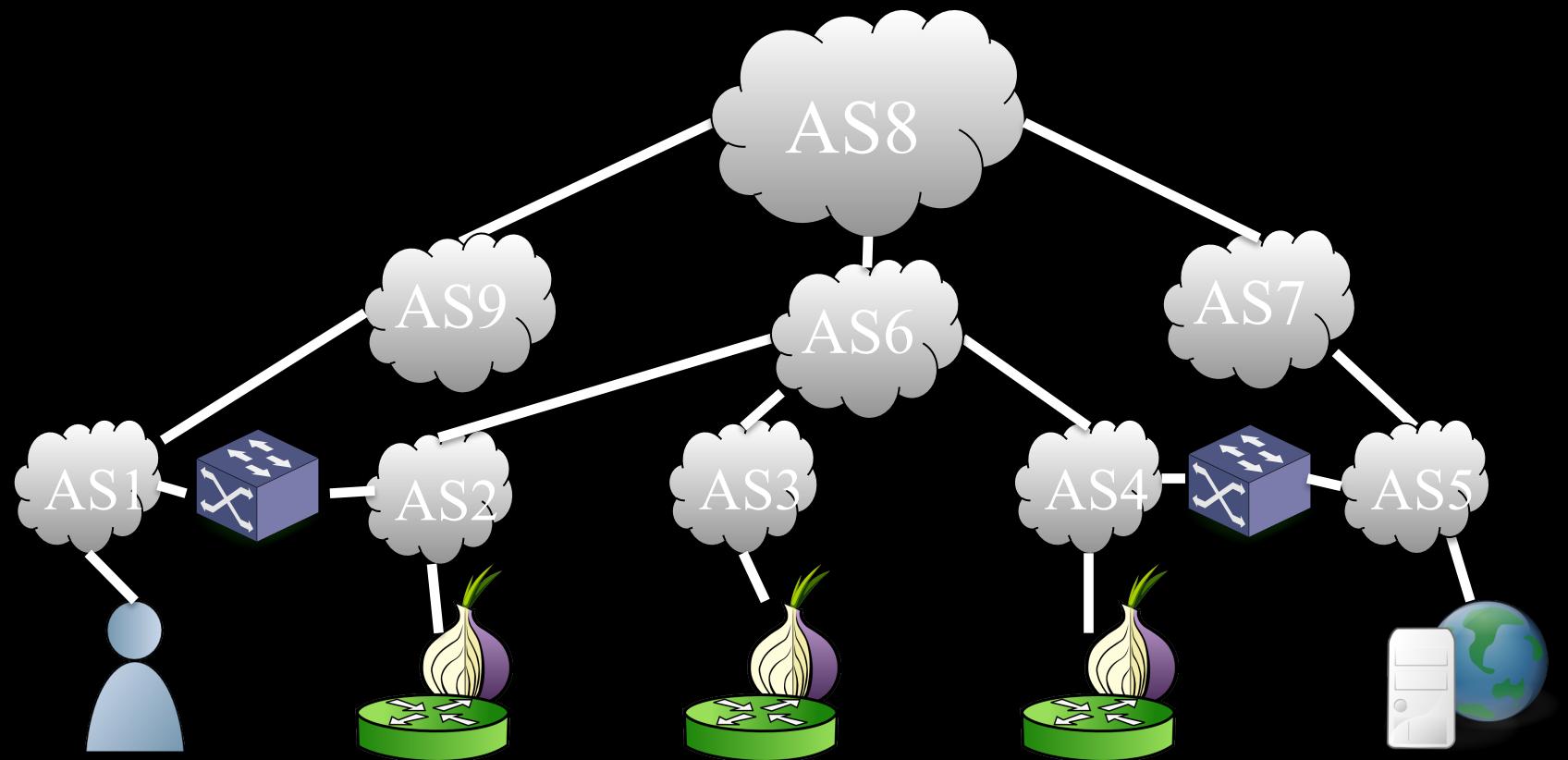


October 2012 – March 2013

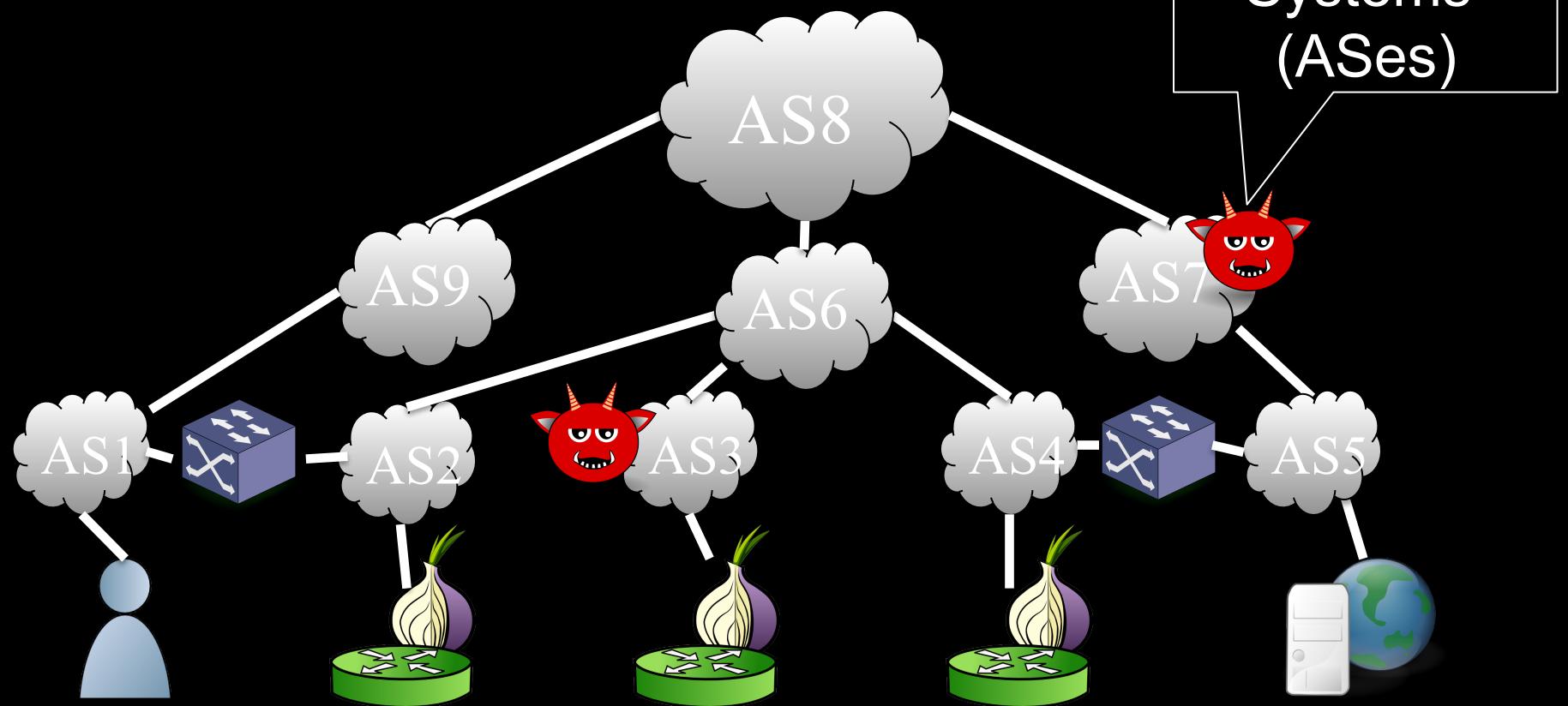
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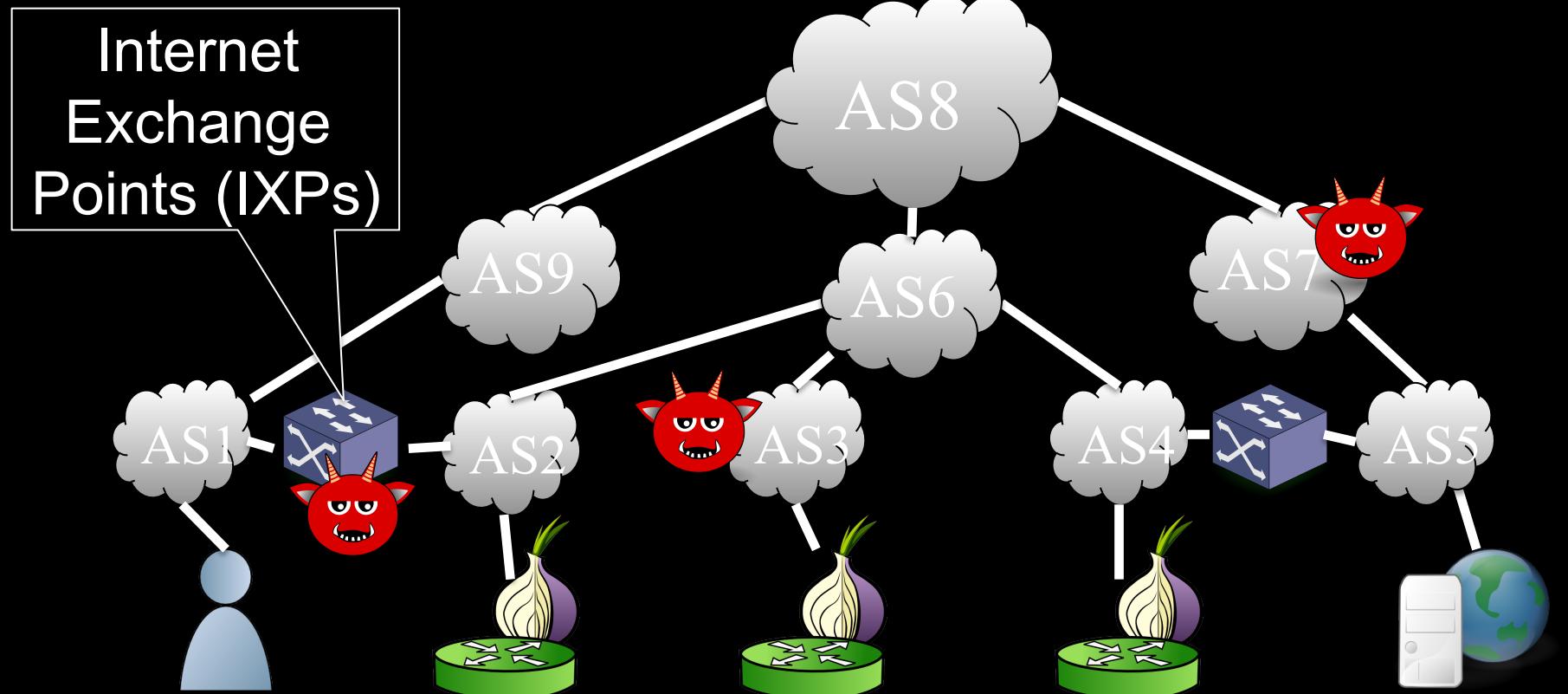
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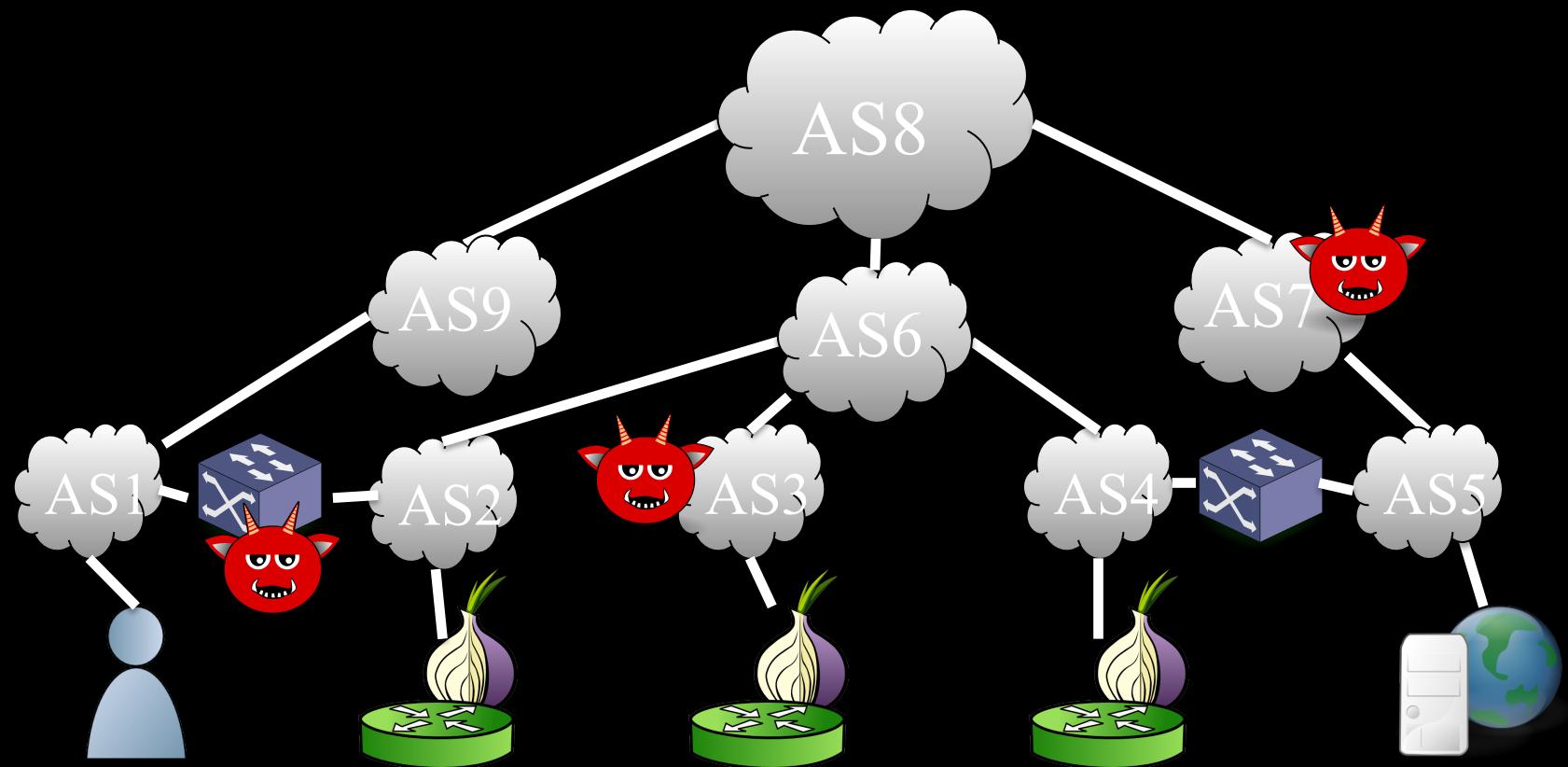
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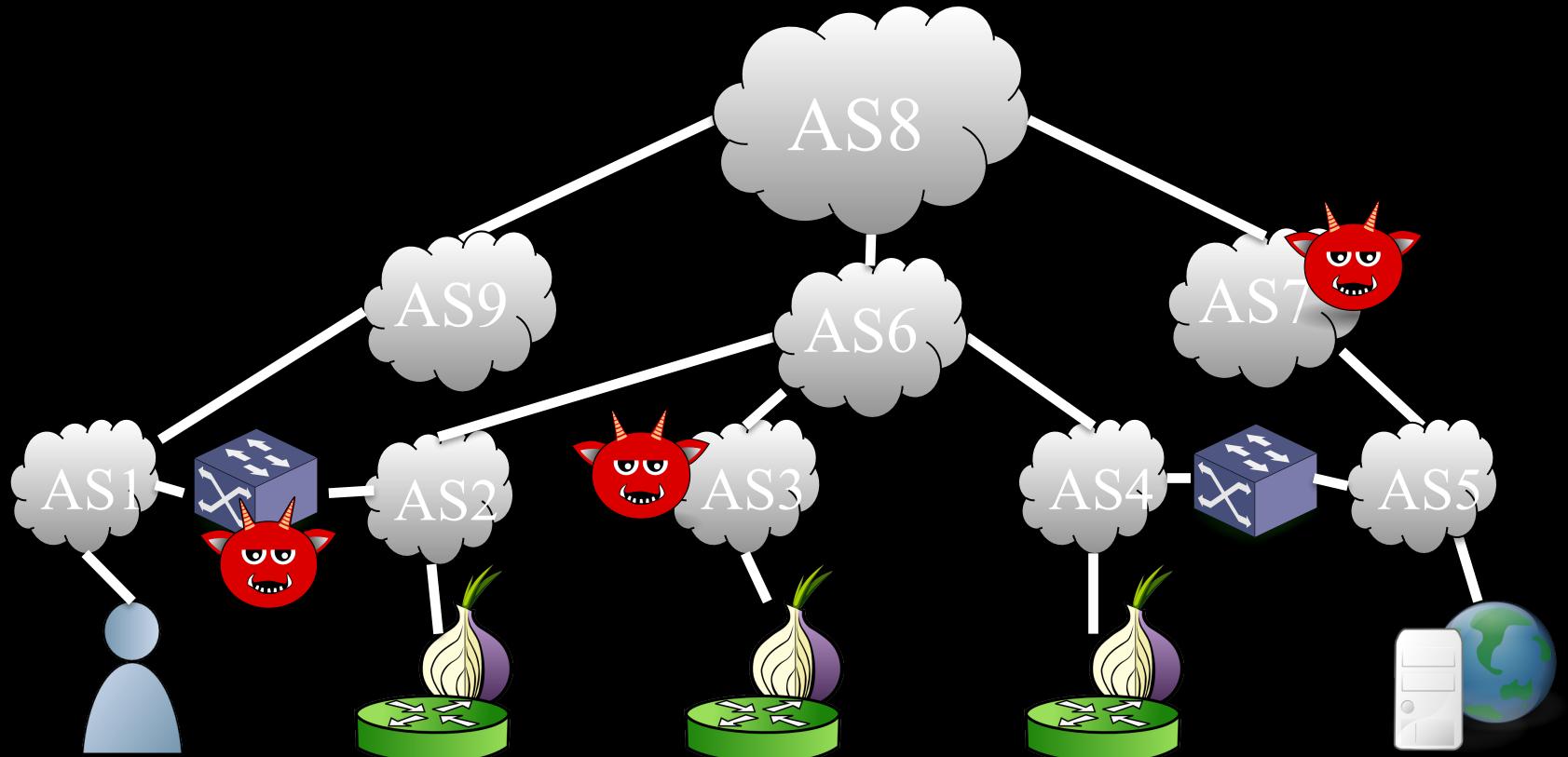


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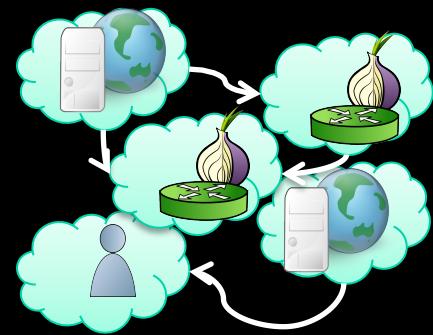
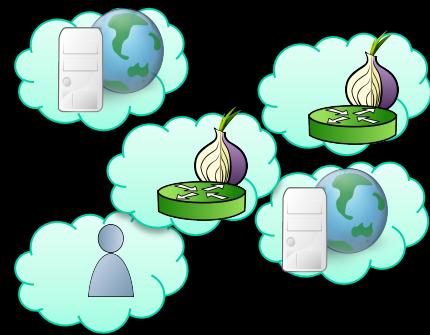
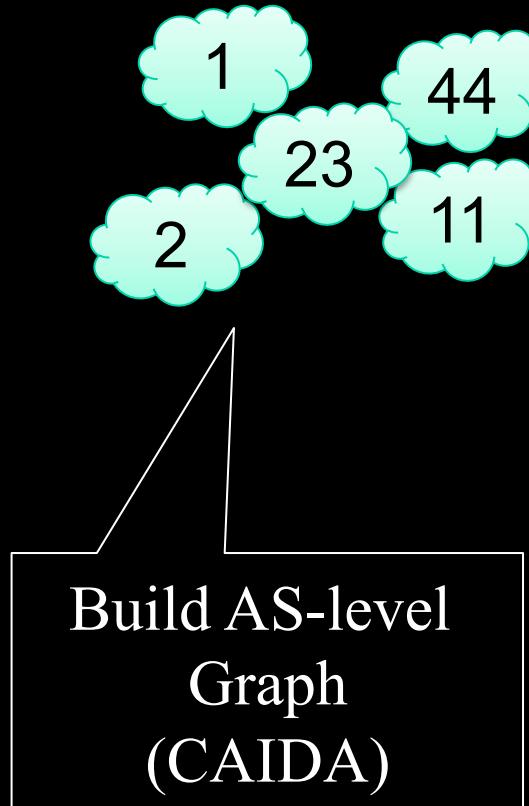
- Adversary has **fixed** location
- Adversary may control **multiple** entities

# Network Adversary

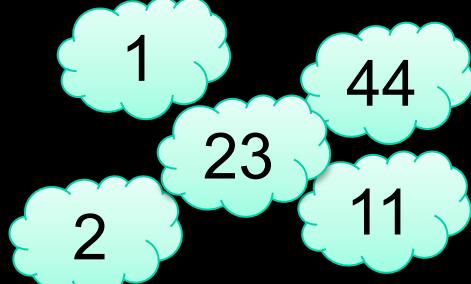


- Should most users be concerned with a network adversary?
- Adversary has **fixed** location
- Adversary may control **multiple** entities

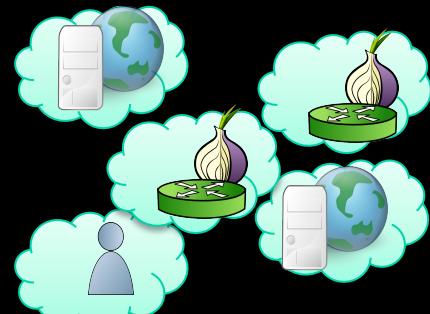
# Simulating a Network Adversary



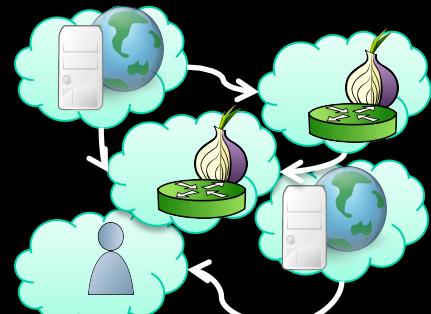
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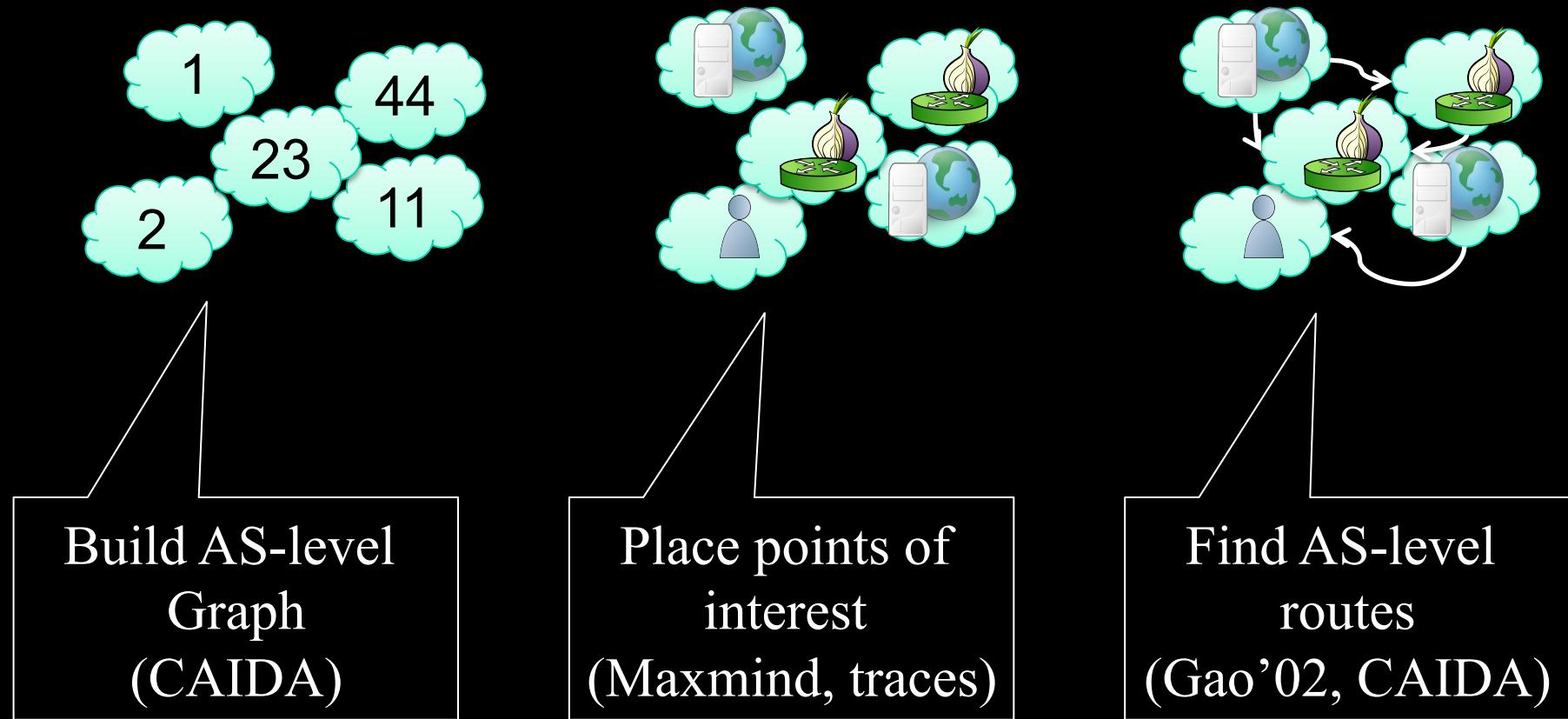
Build AS-level  
Graph  
(CAIDA)



Place points of  
interest  
(Maxmind, traces)



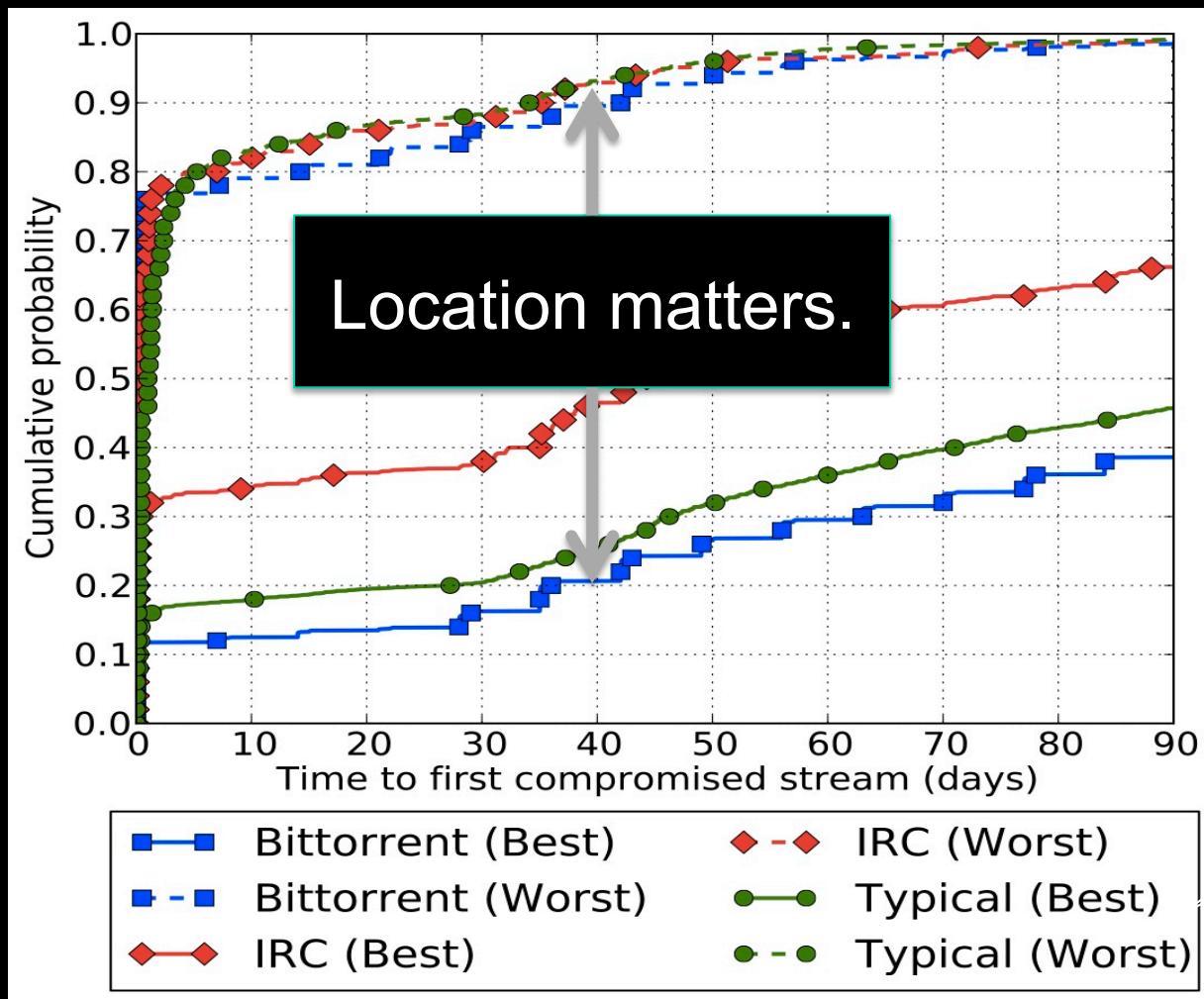
# Simulating a Network Adversary



# Selecting Network Adversaries

1. Rank each AS/IXP for each client location by frequency on entry or exit paths;
2. Exclude src/dst ASes (compromises nearly all paths); and
3. Assign adversary to top  $k$  ASes or IXPs

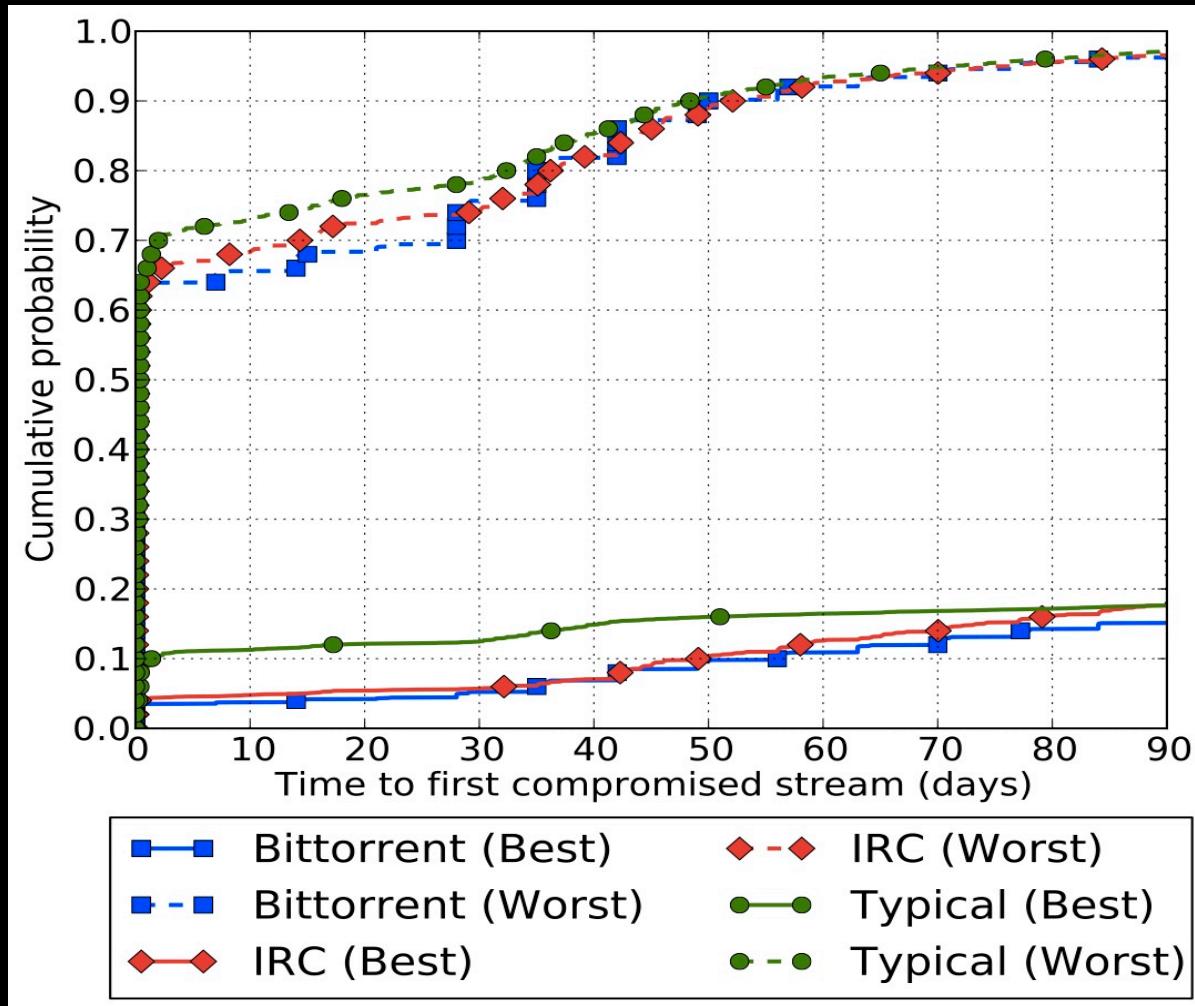
# Adversary Controls One AS



“best”/“worst”  
denote most/least  
secure client

January 2013 – March 2013

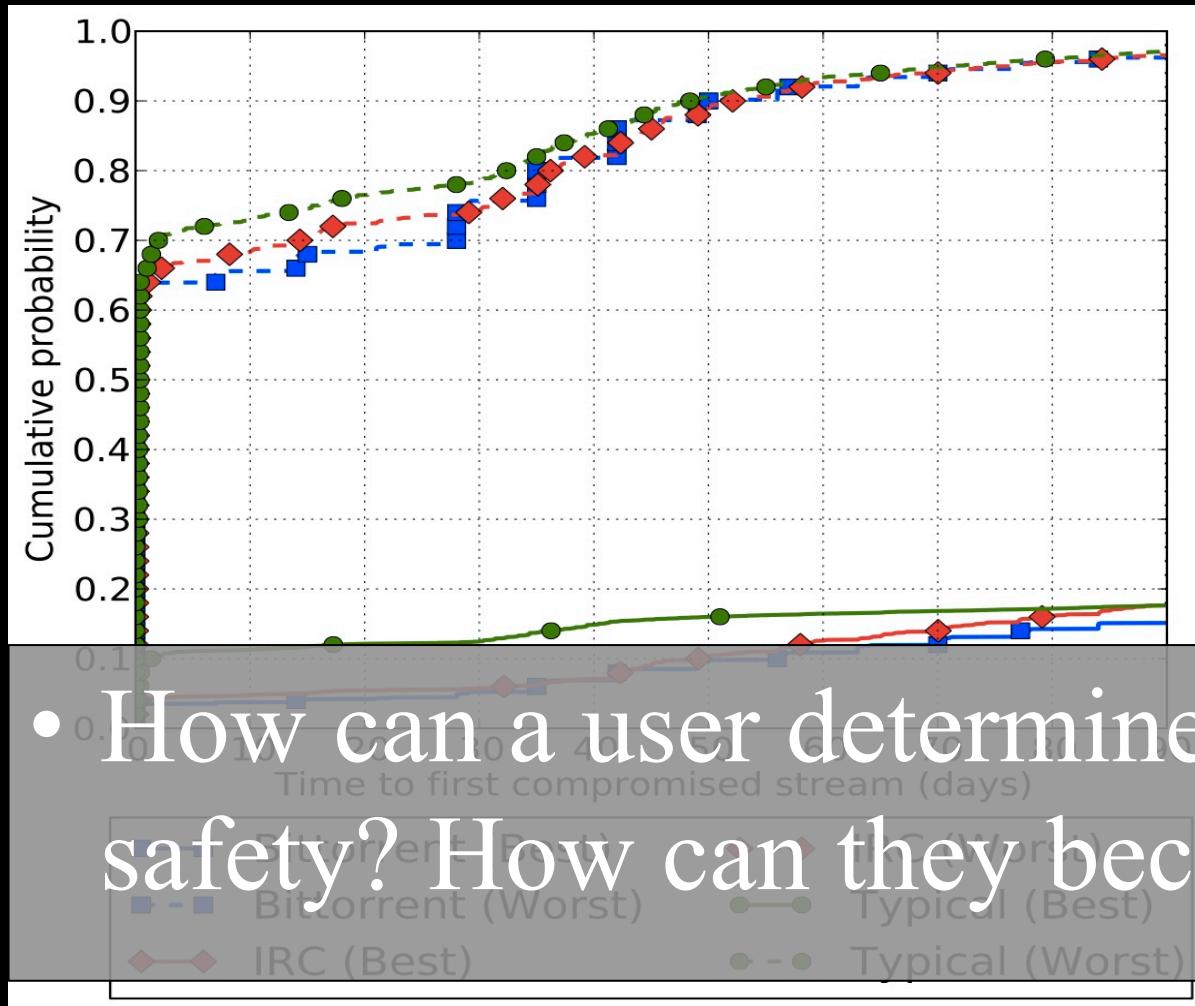
# Adversary Controls One IXP Organization



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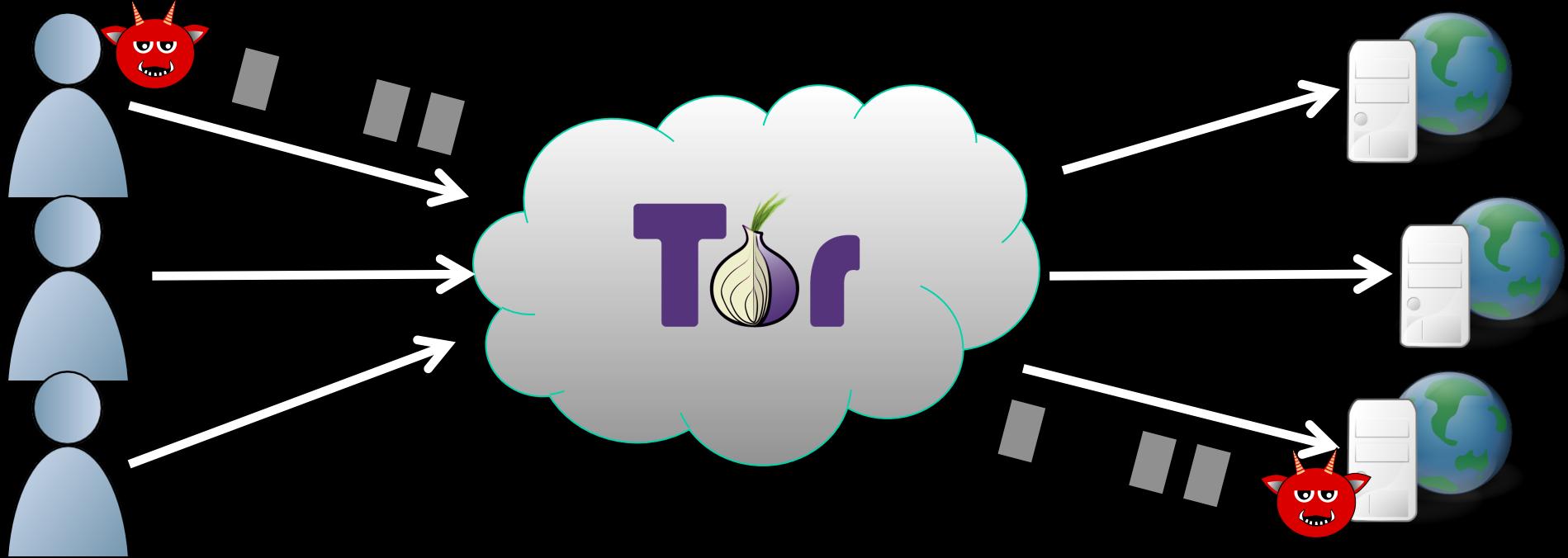


“best”/“worst”  
denote most/least

secure client

- How can a user determine their safety? How can they become safer?

# Traffic Correlation

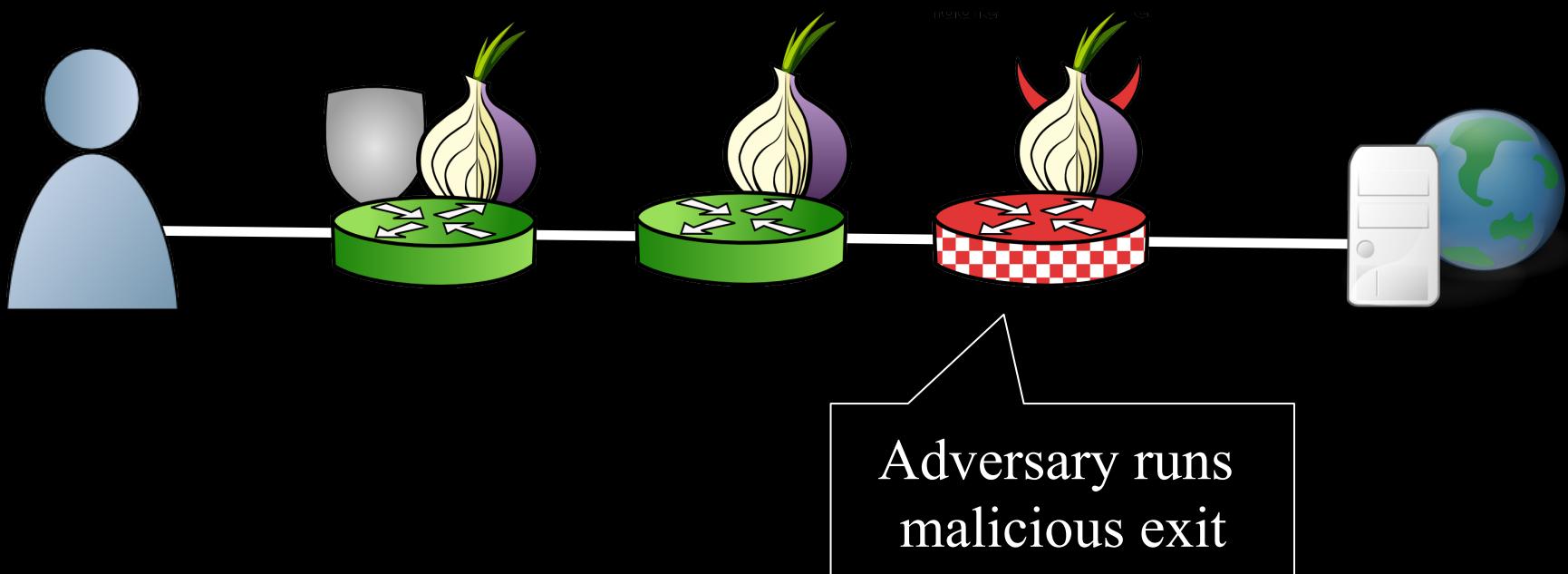


- What if the adversary only controls one of the ends?

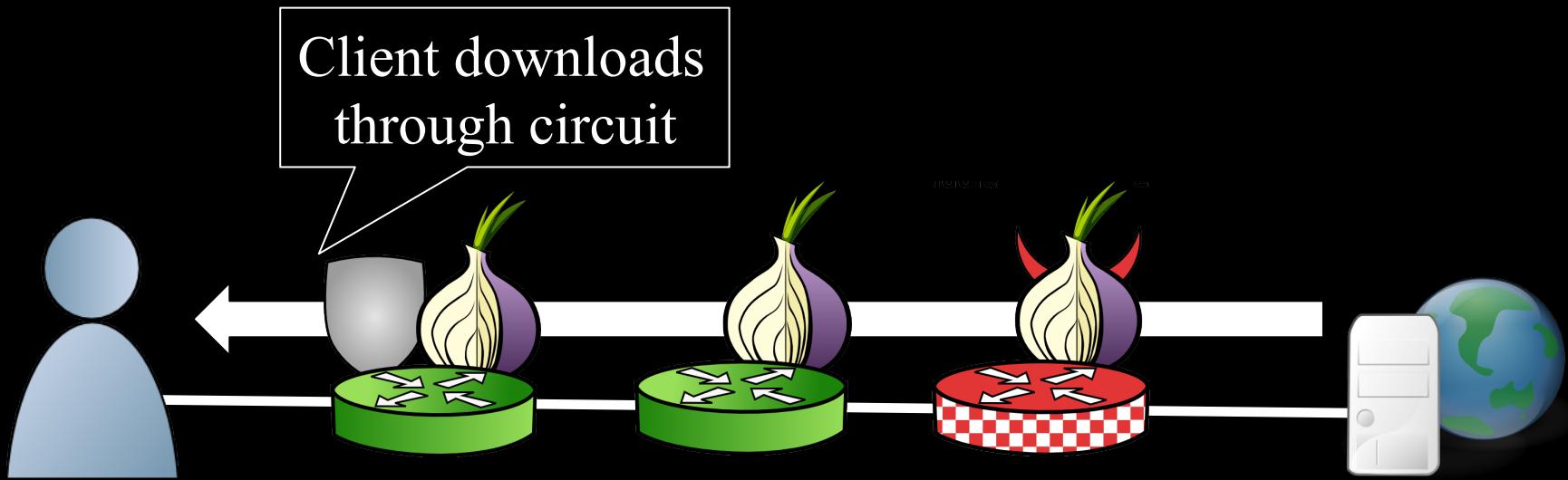
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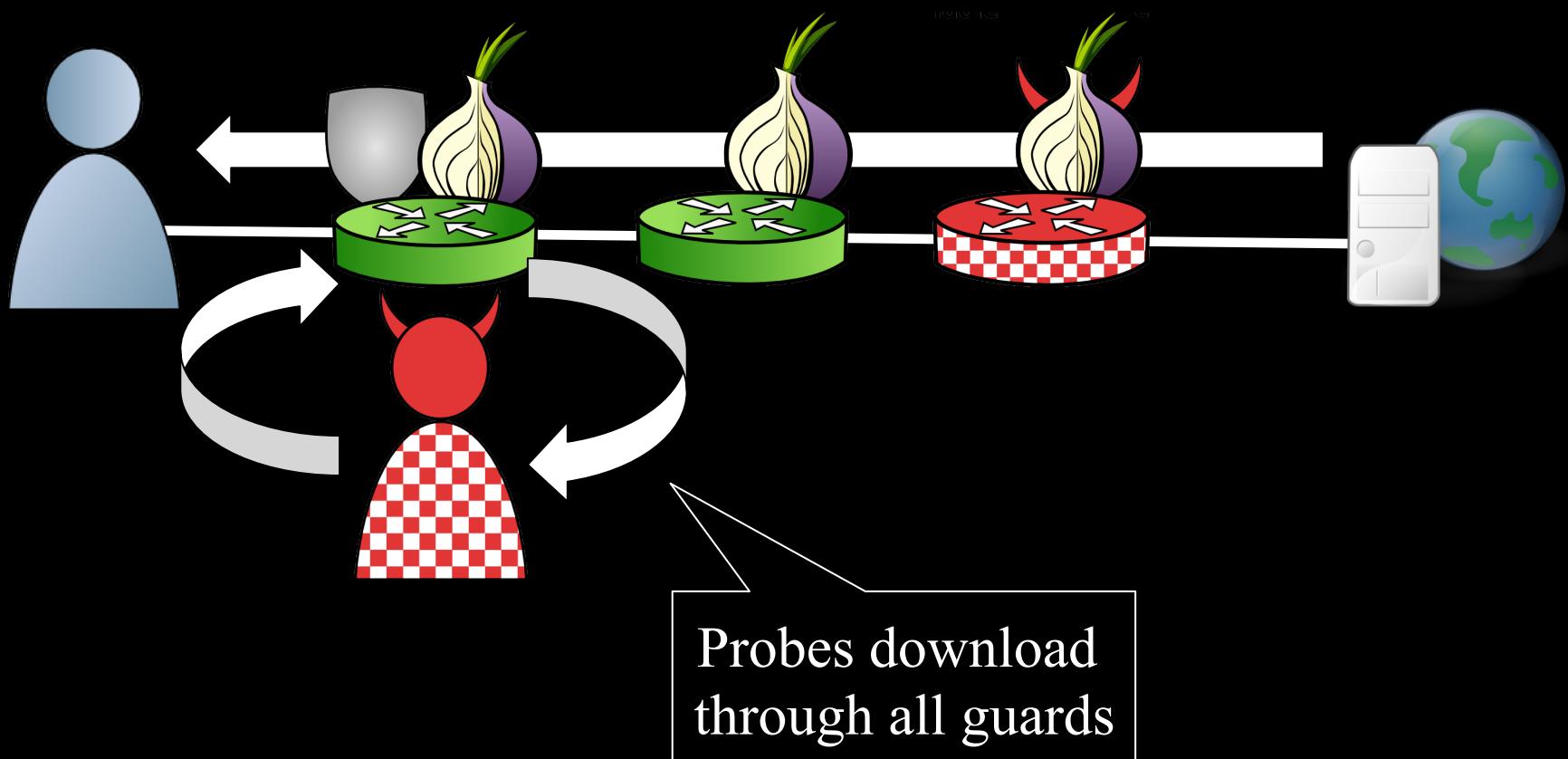
# Traffic Correlation: Throughput



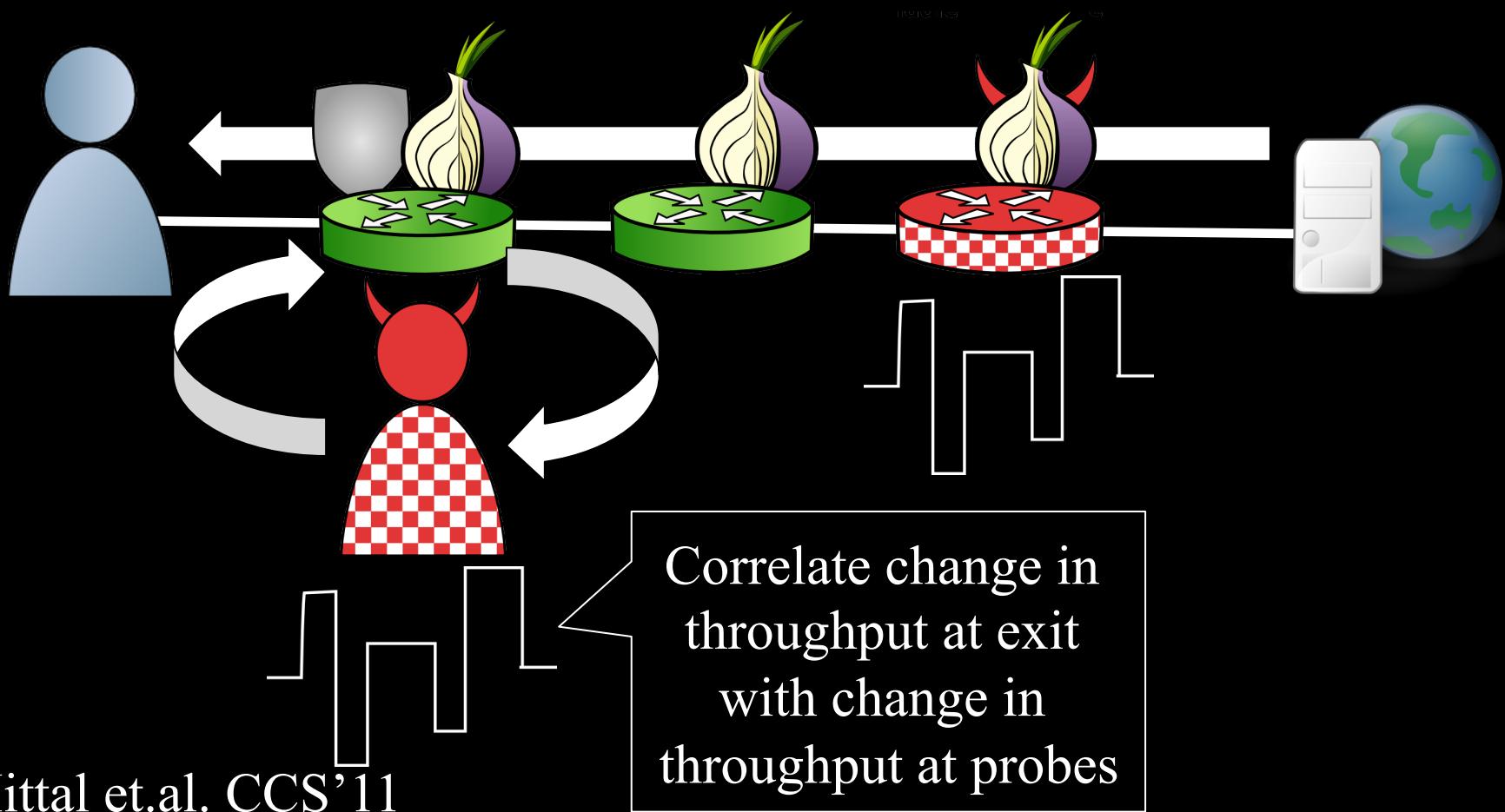
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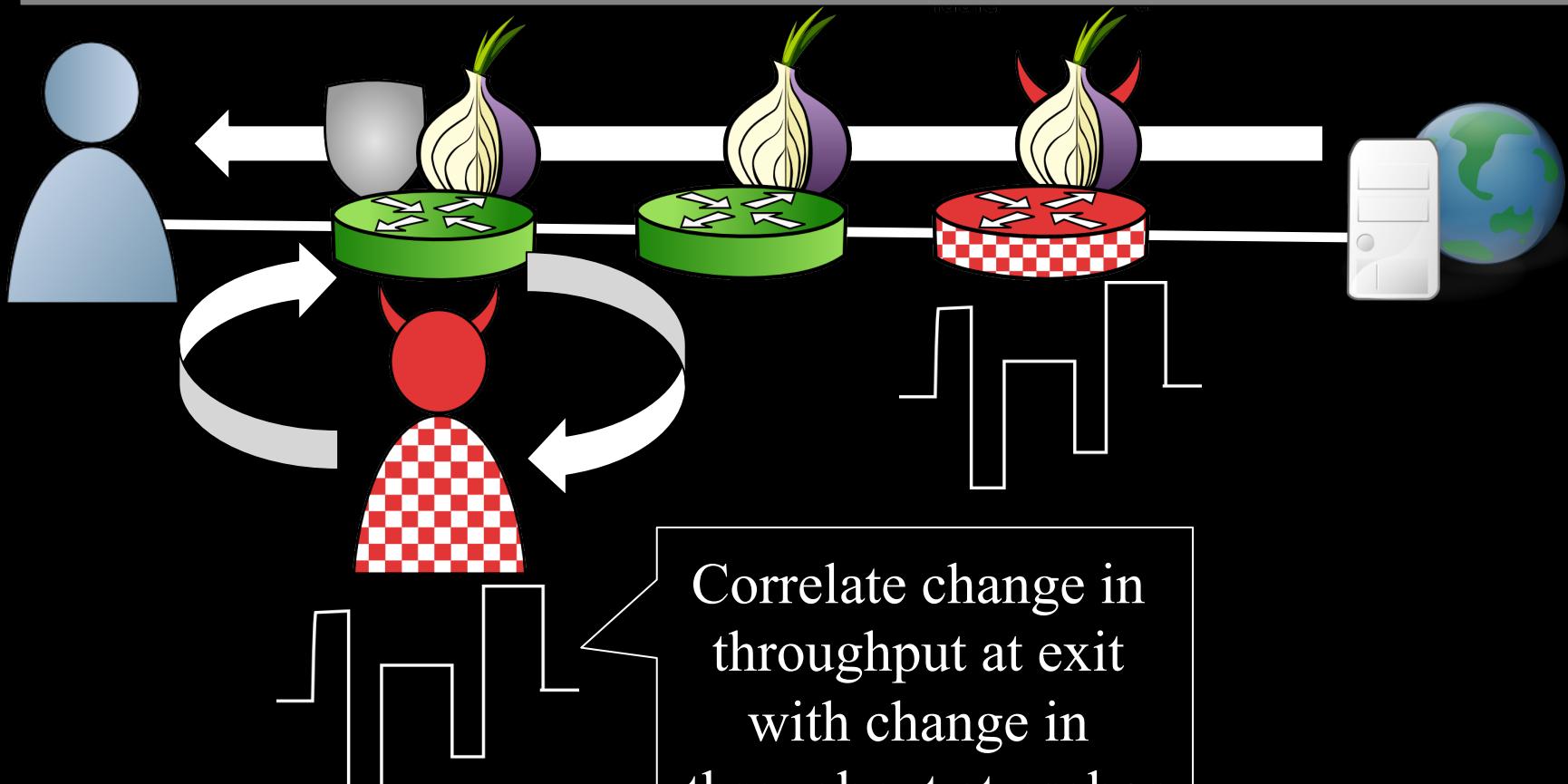


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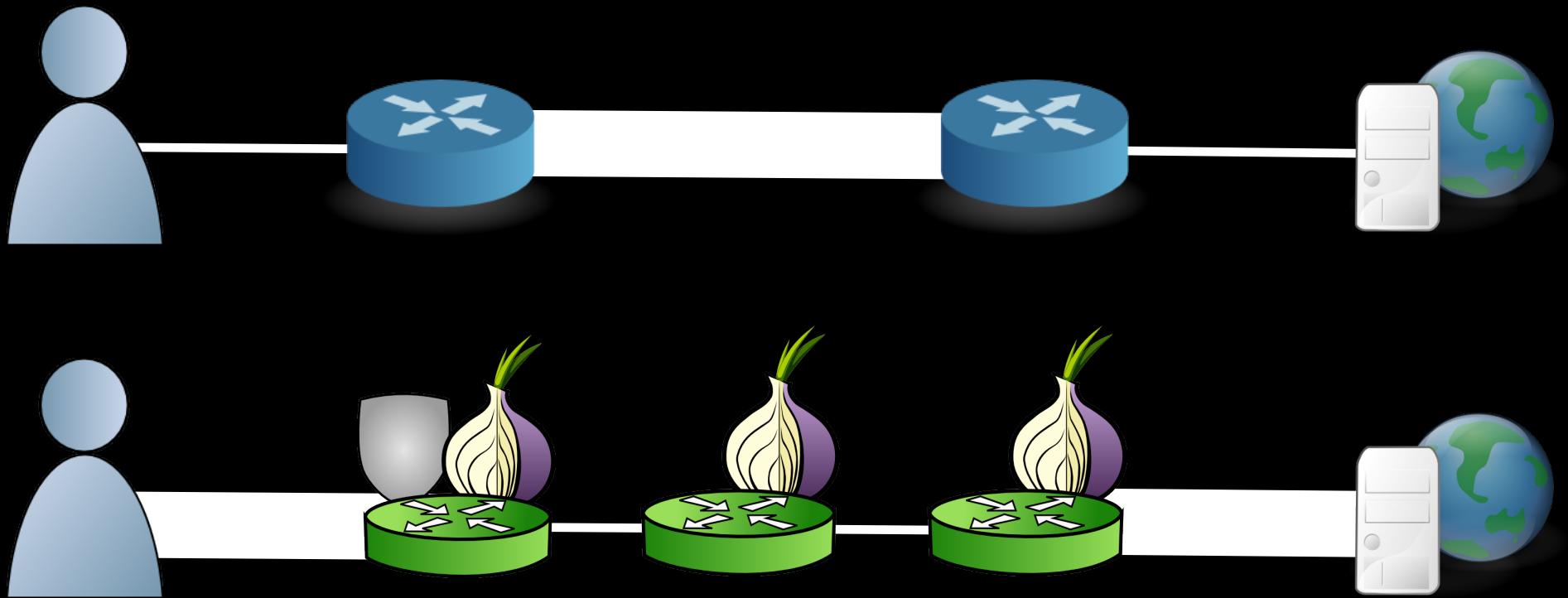
- How is this attack “stealthy”?



# Outline

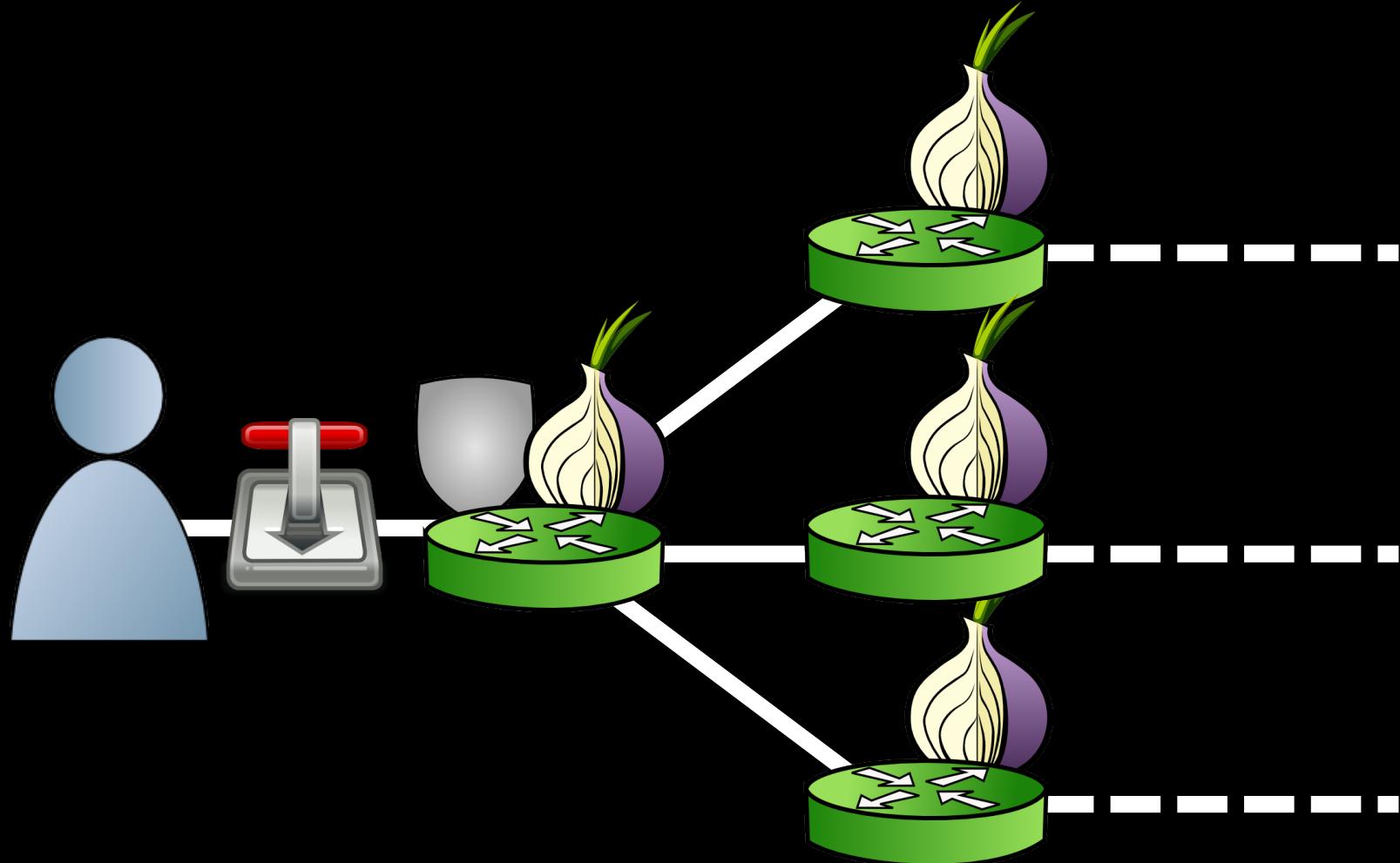
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# Tor != Internet

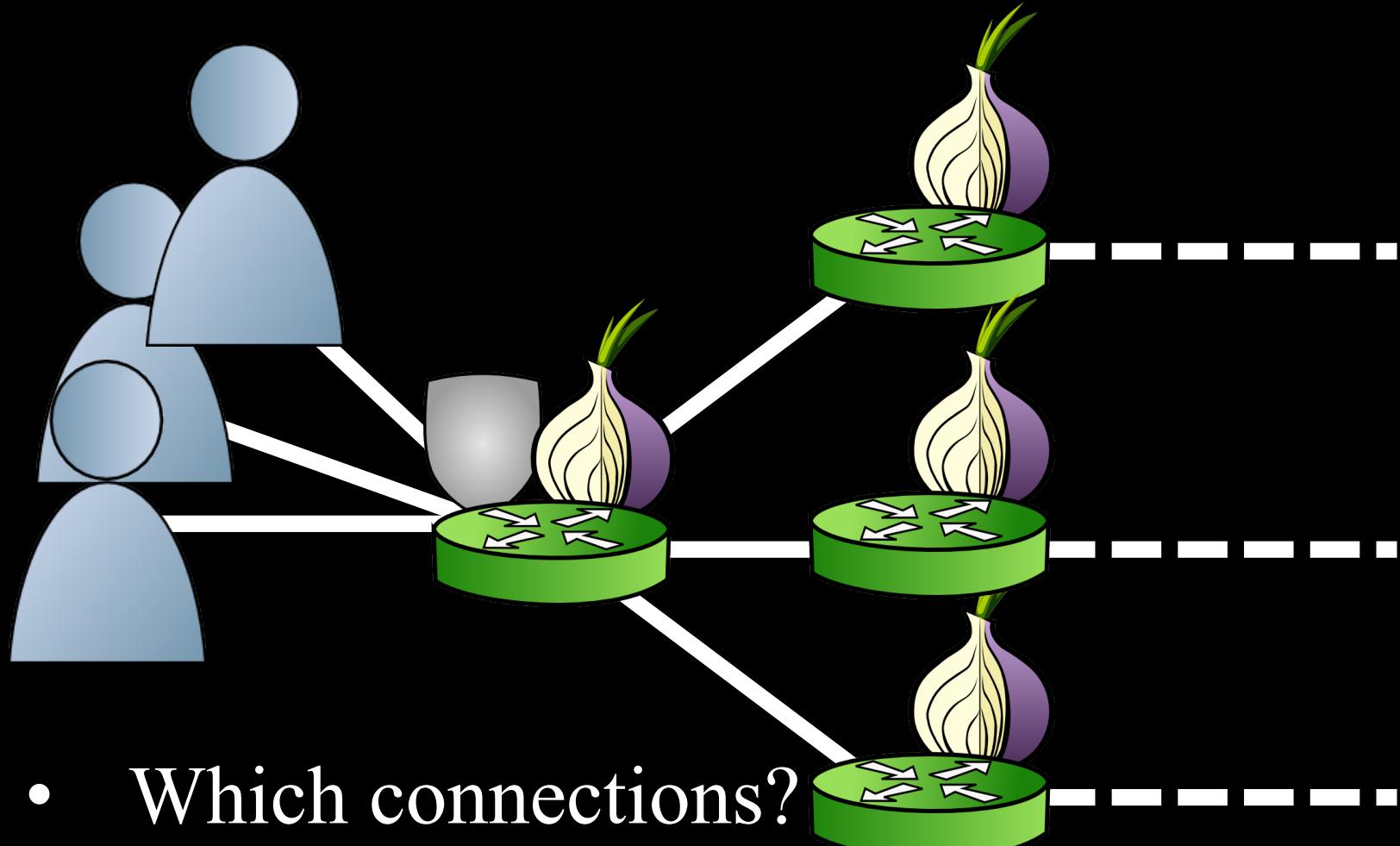


- Specialized Tor performance enhancements
  - Reducing load: **traffic admission control**
  - Reducing load, improving utilization: congestion control

# Traffic Admission Control



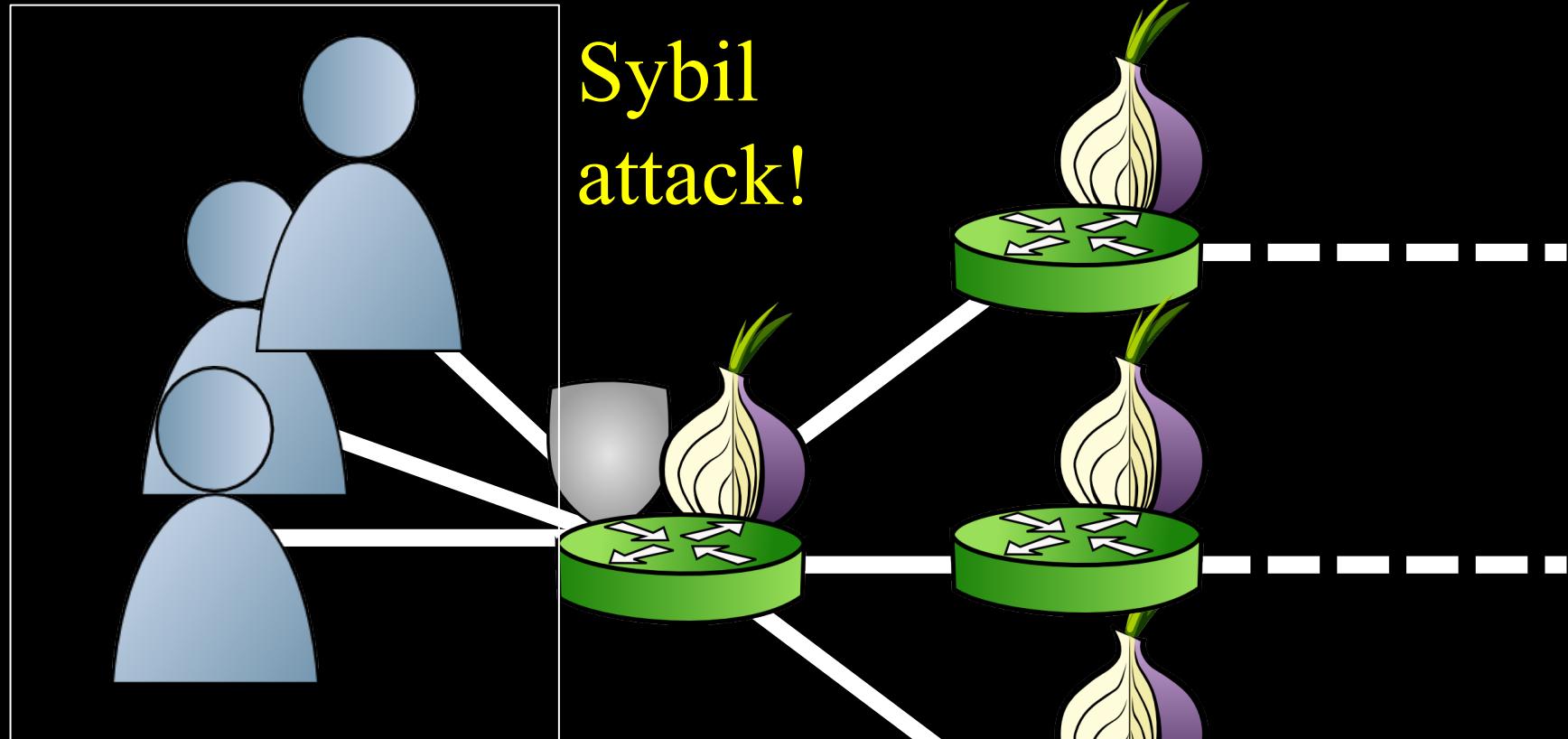
# Traffic Admission Control



- Which connections?
- At what rate?



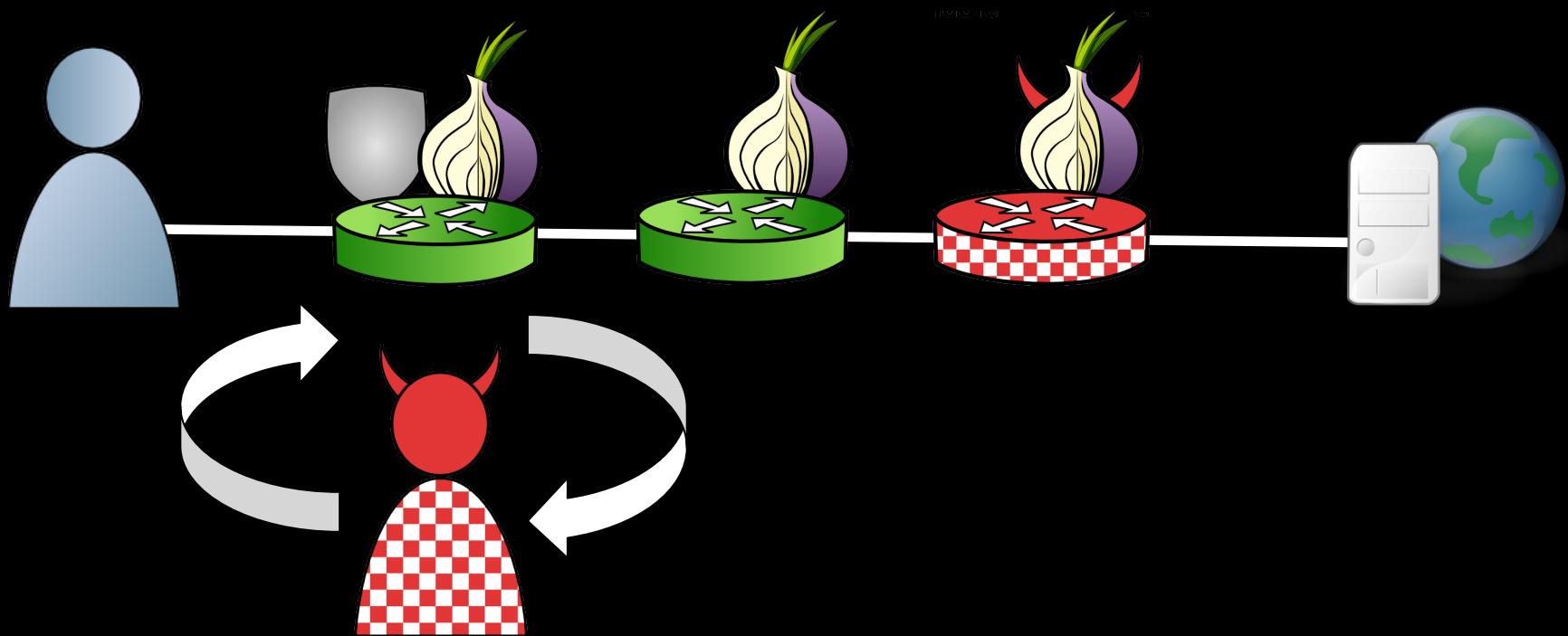
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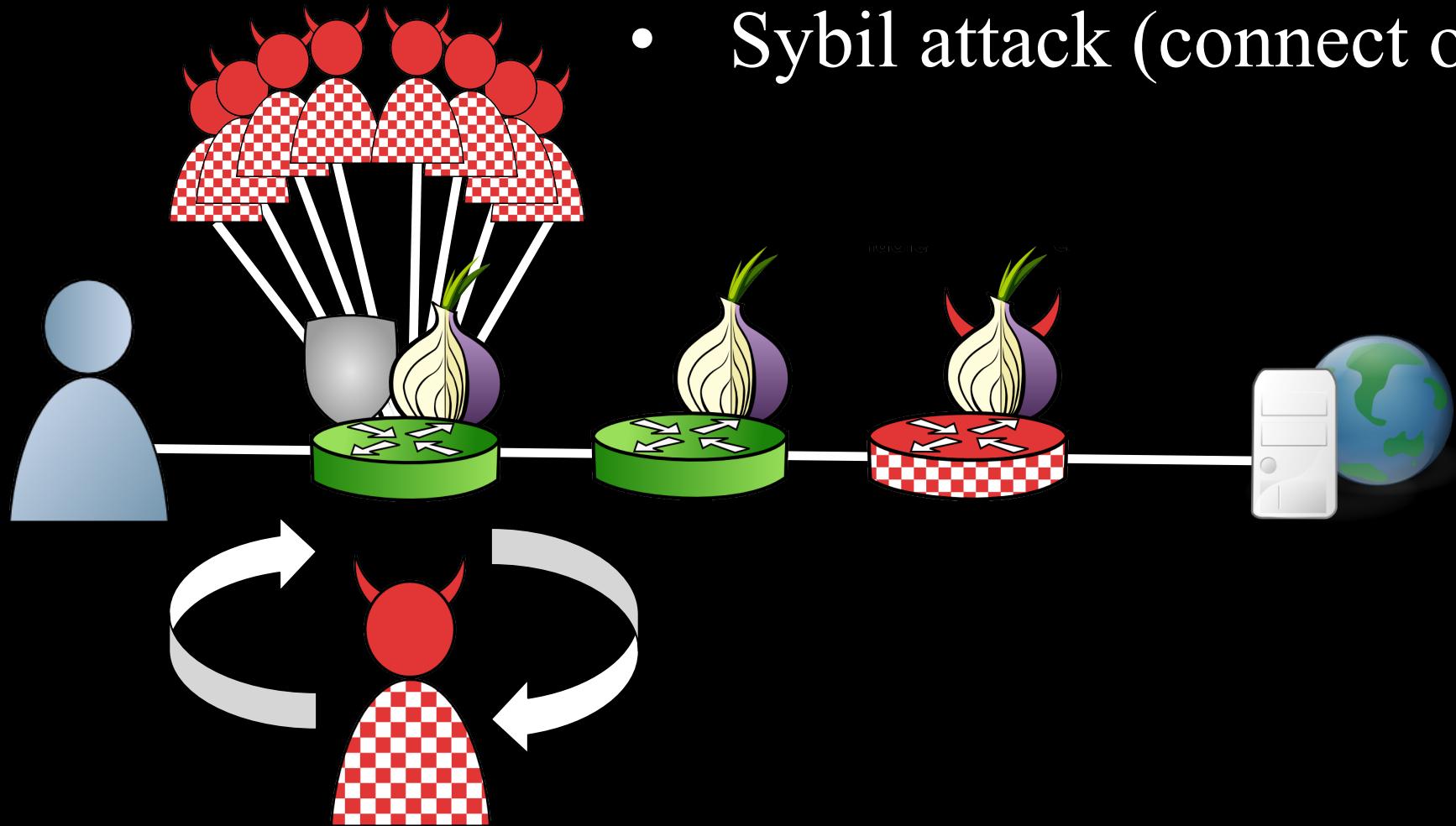


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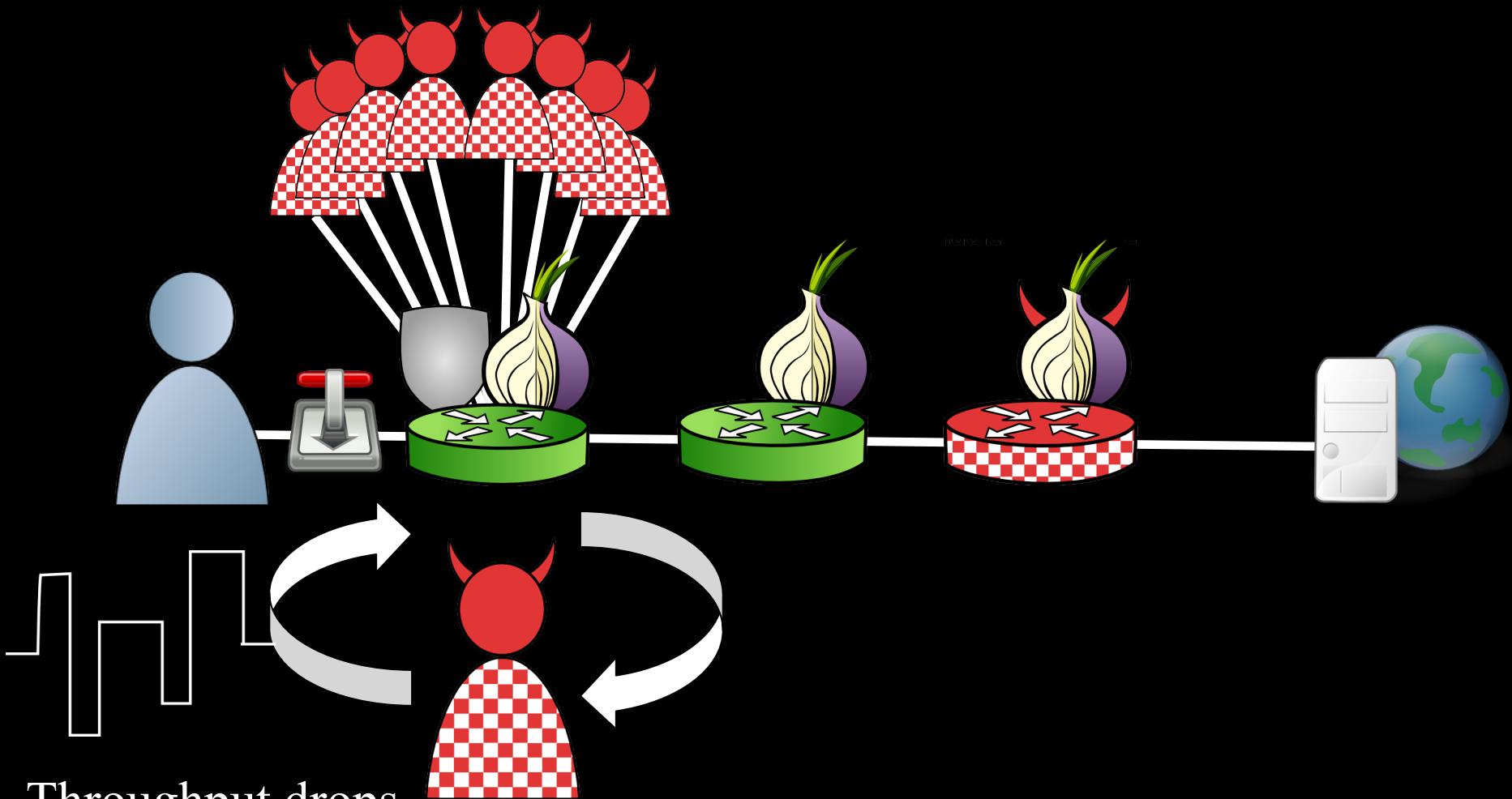


# Traffic Admission Control

- Sybil attack (connect only)



# Traffic Admission Control

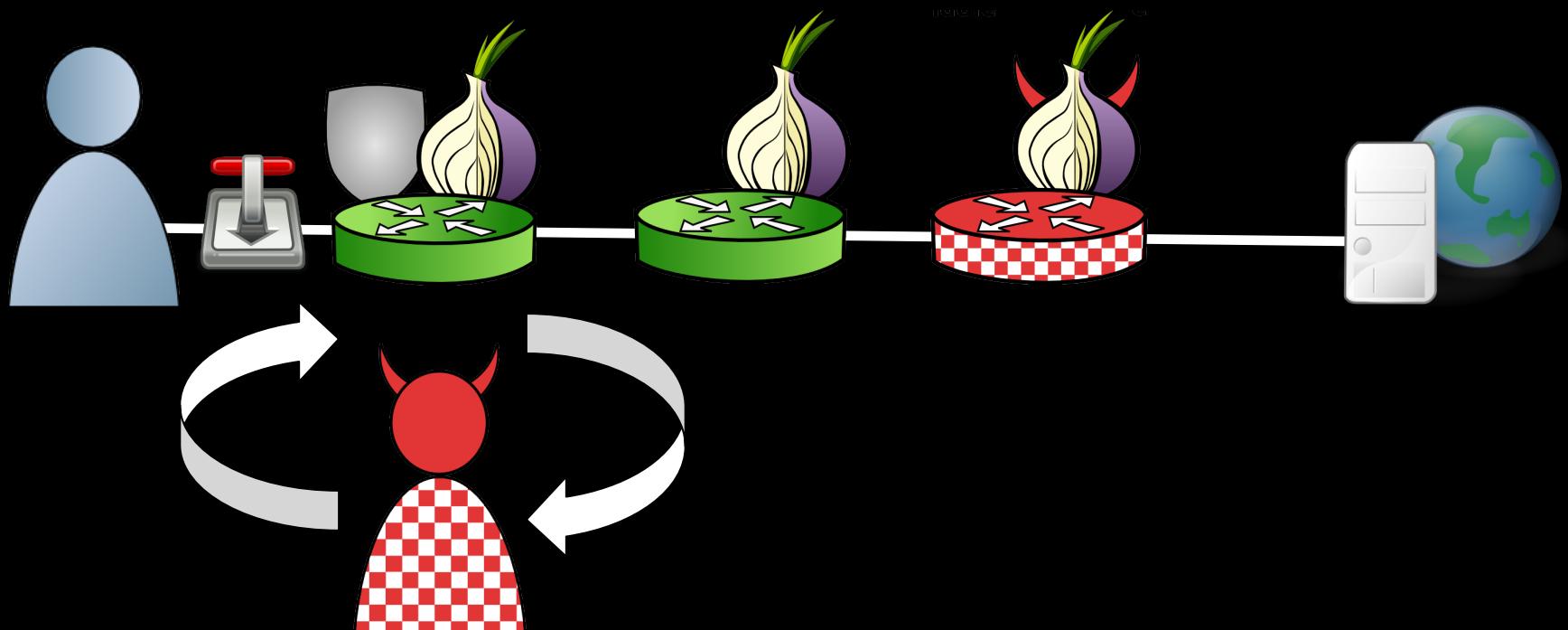


Throughput drops  
to throttle rate

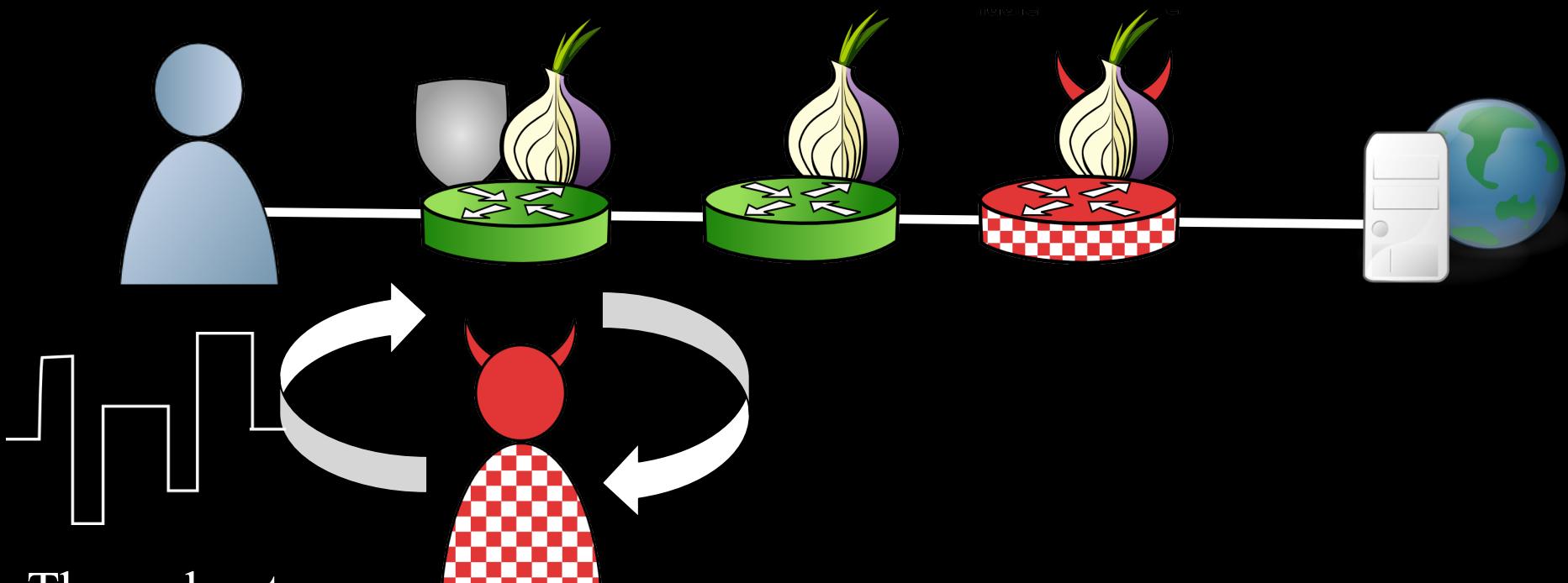
Geddes et.al.  
PETS'13

# Traffic Admission Control

- Disconnect sybils



# Traffic Admission Control

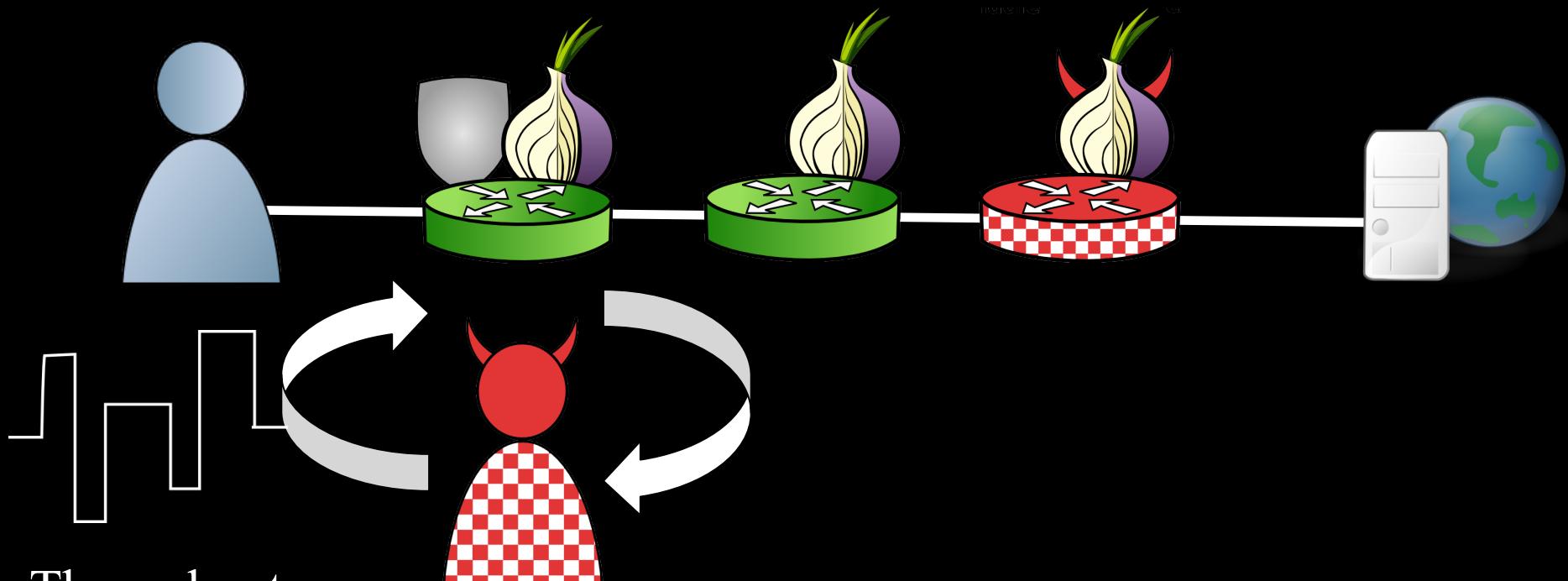


Throughput  
increases

Geddes et.al.  
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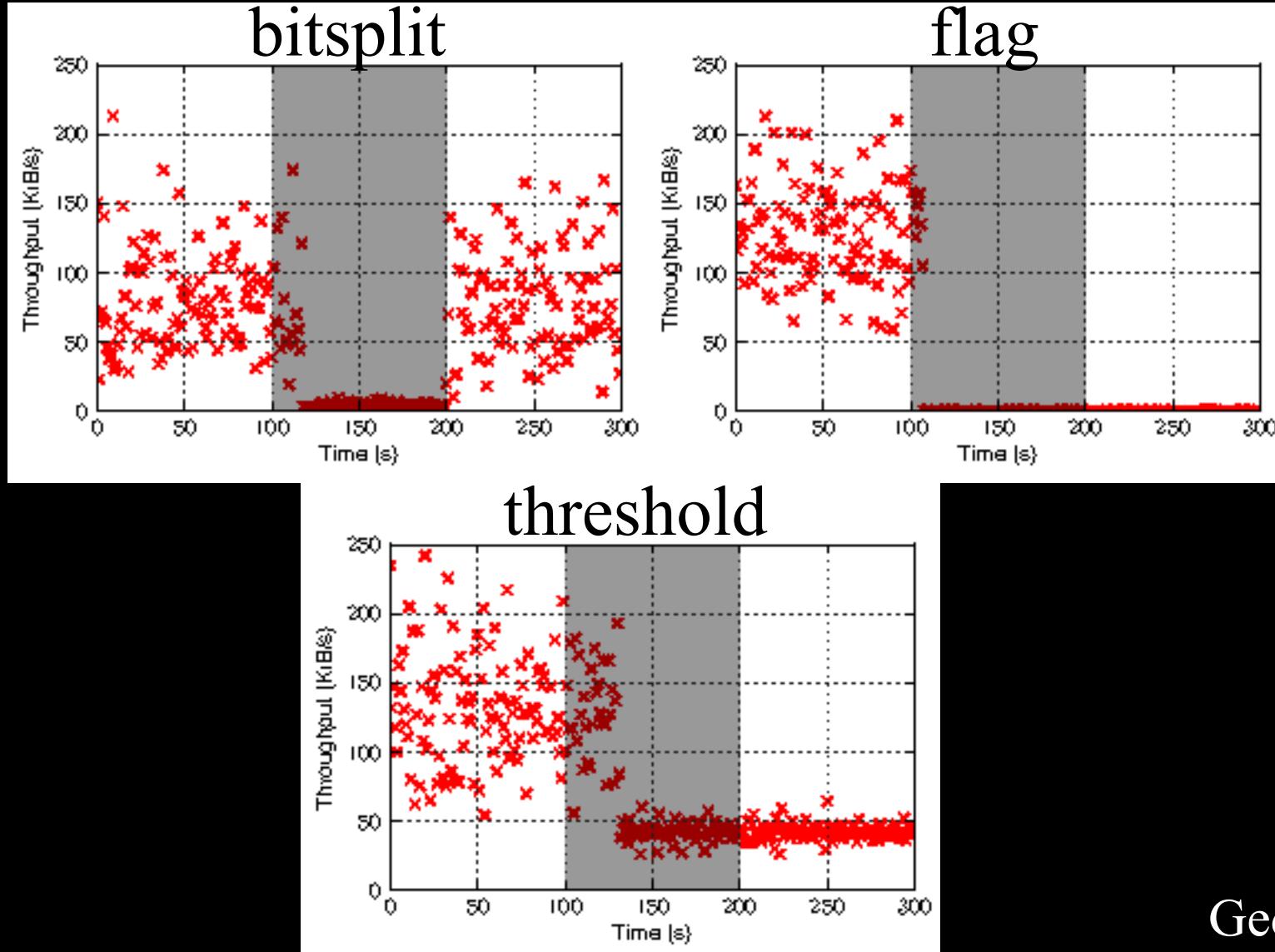
- Is this attack “stealthy”?



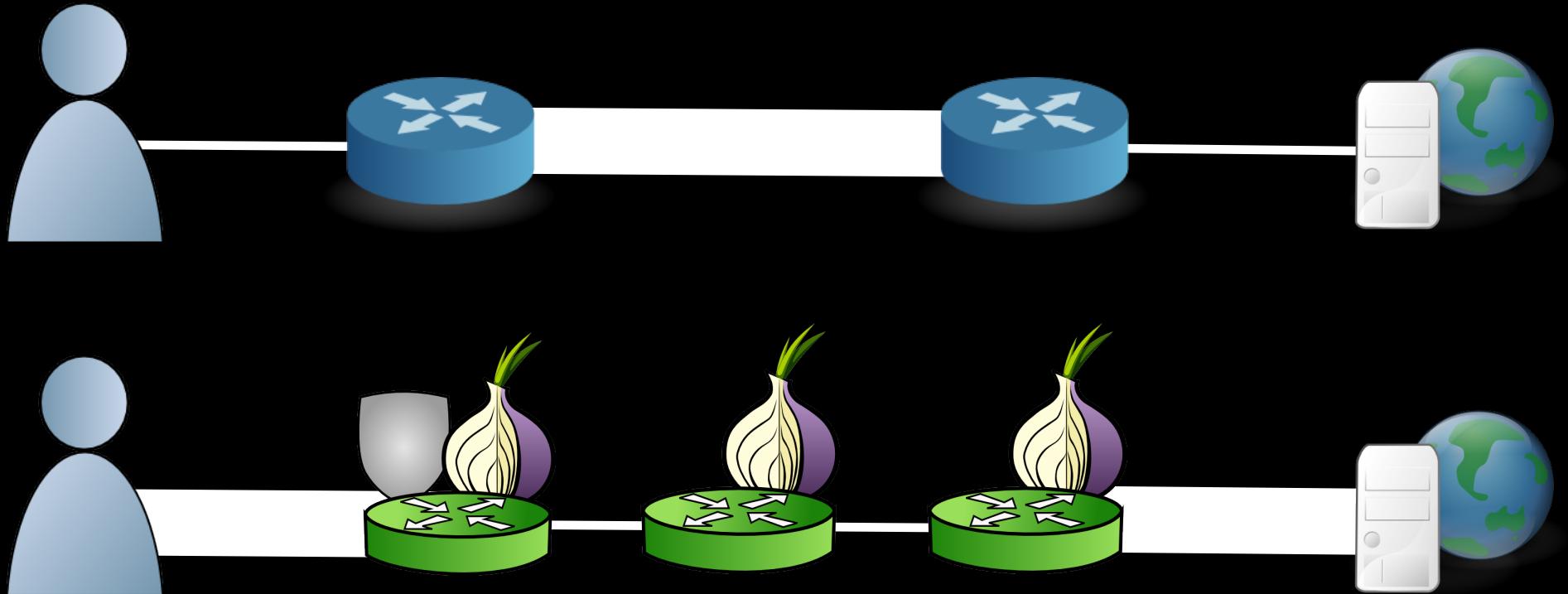
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# Induced Throttling Prototype

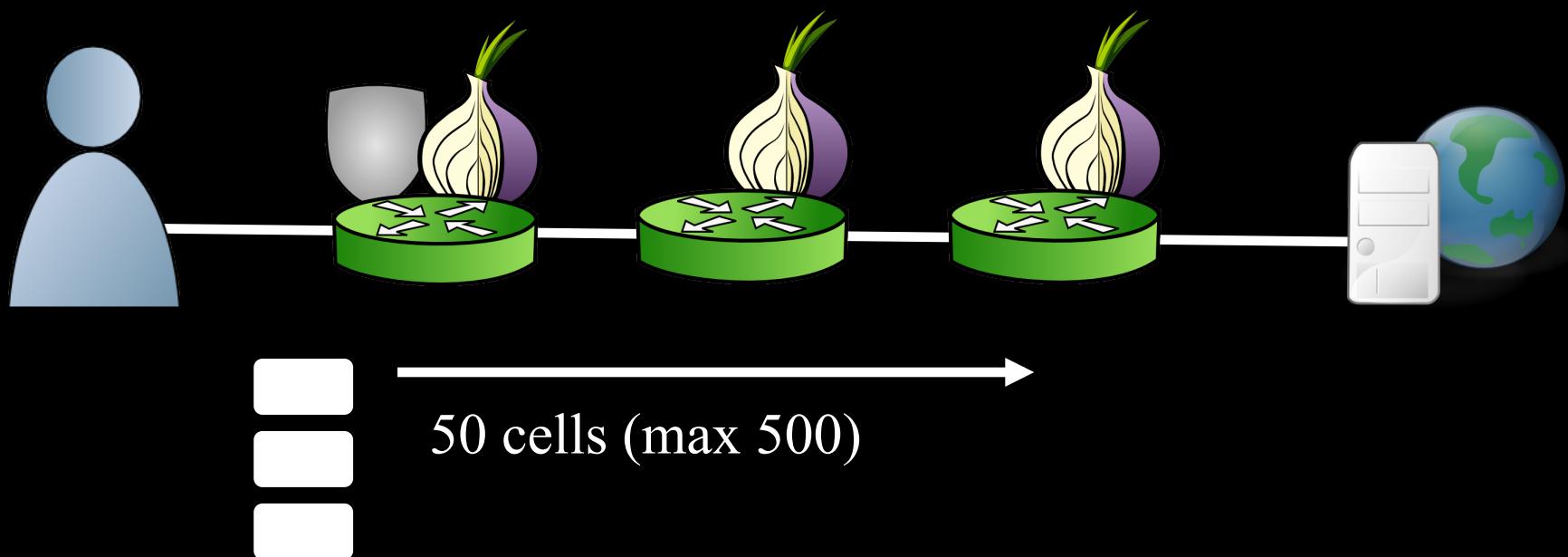


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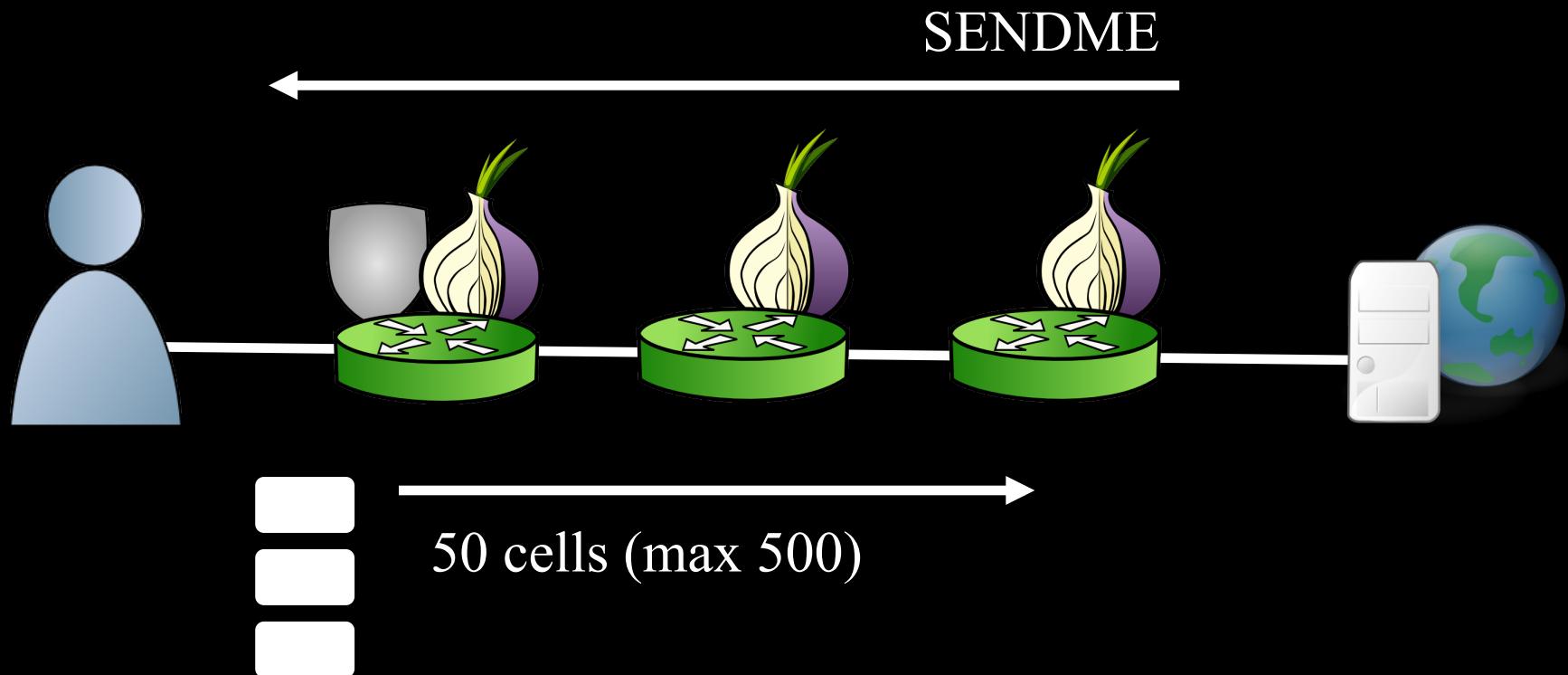


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  - Reducing load, improving utilization: congestion control

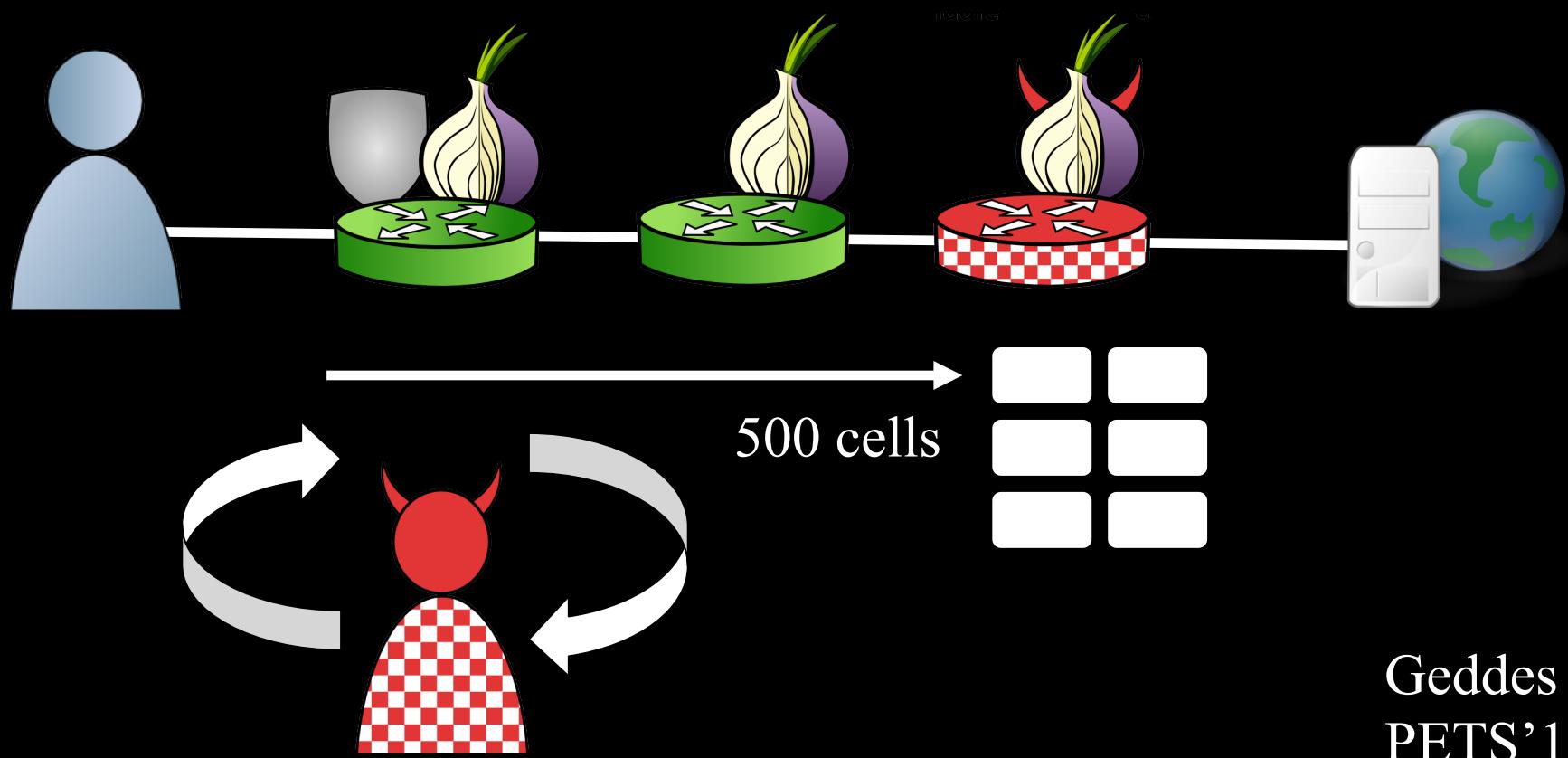
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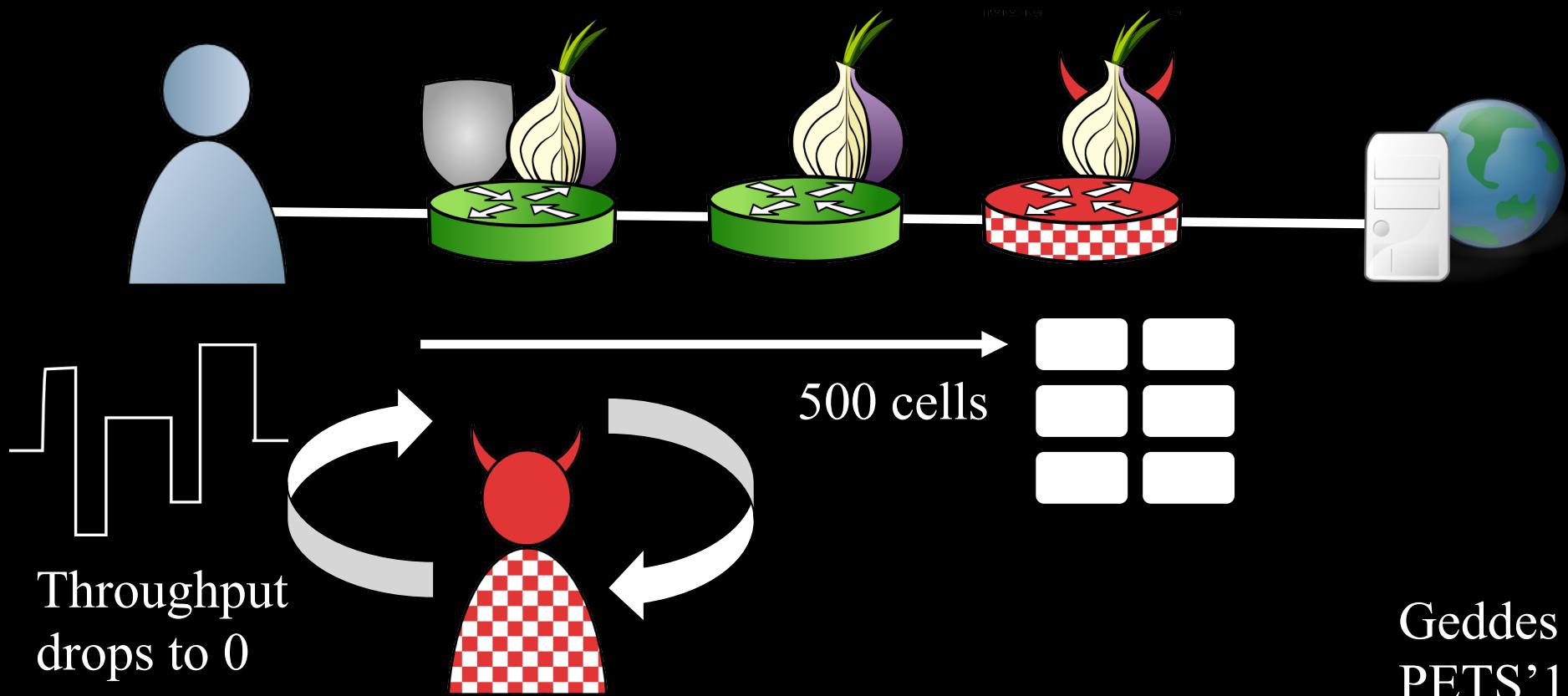


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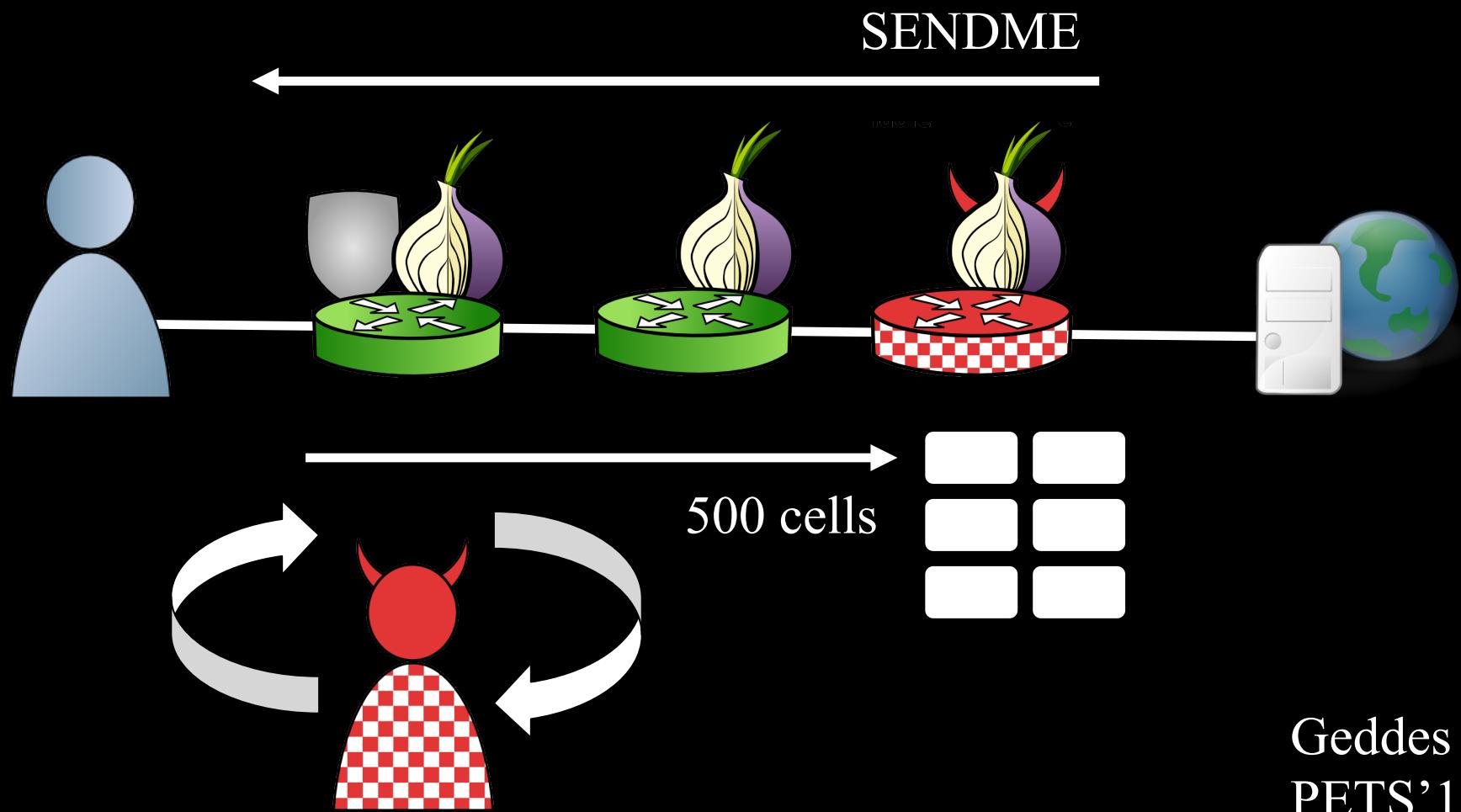


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PETS'13

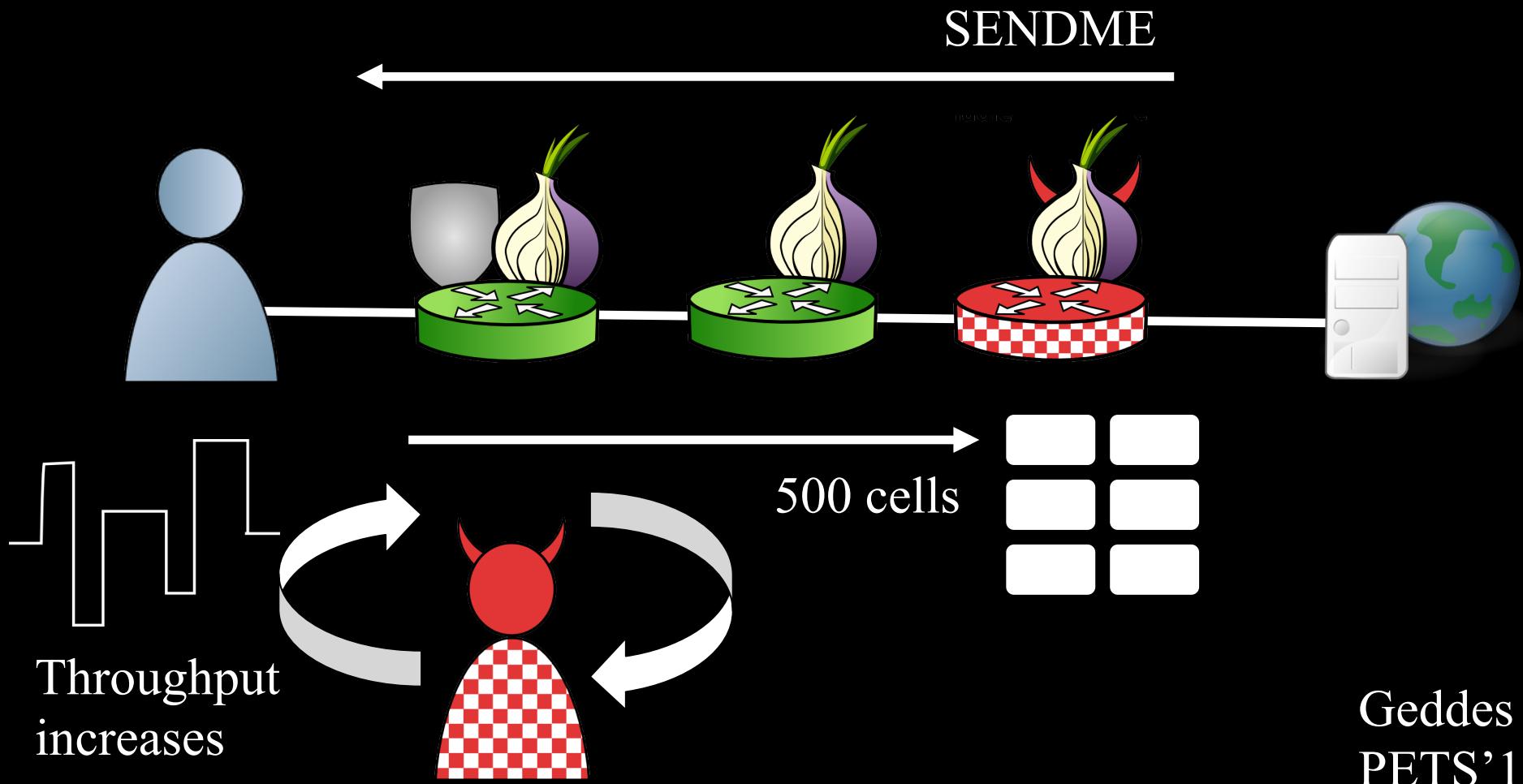
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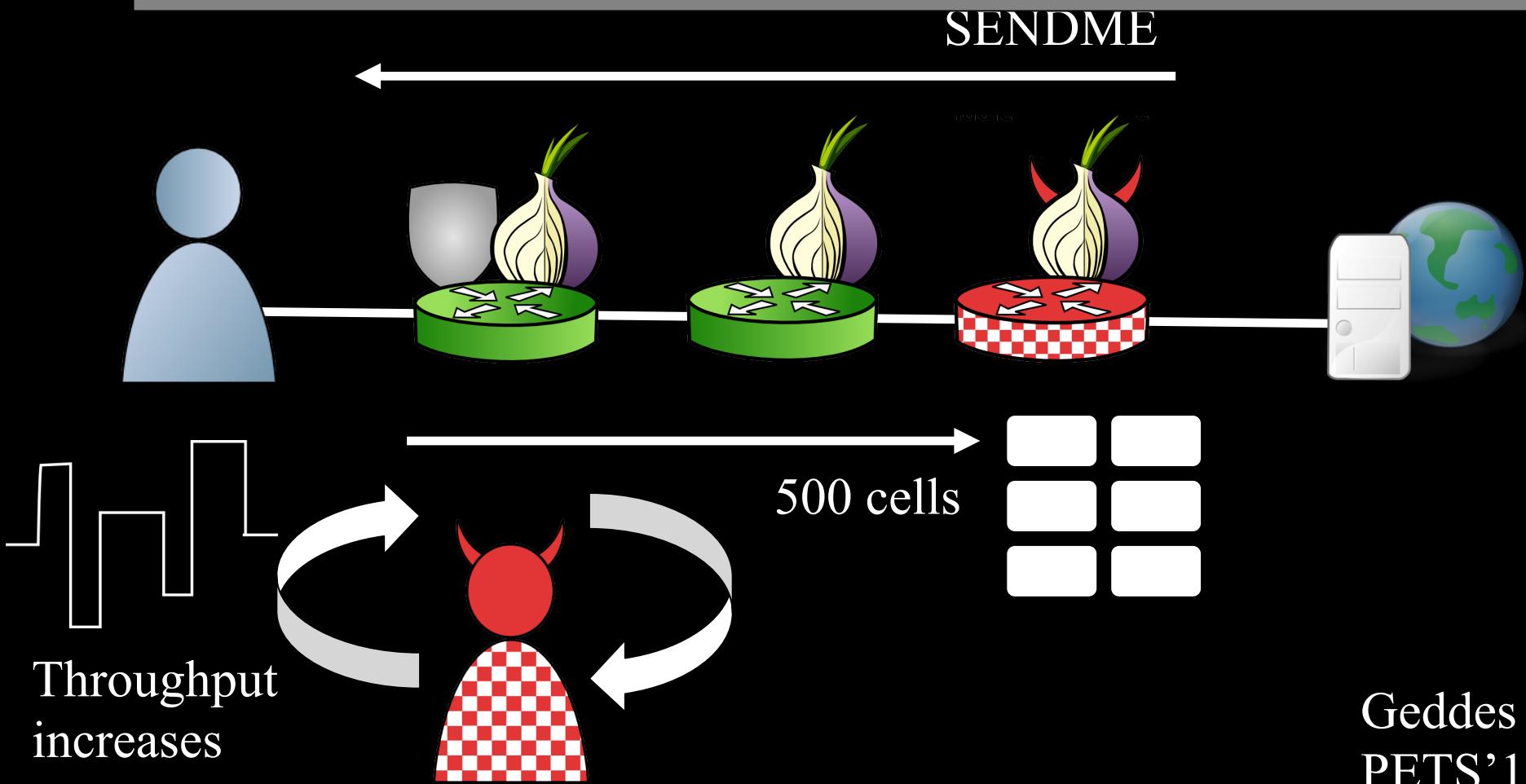


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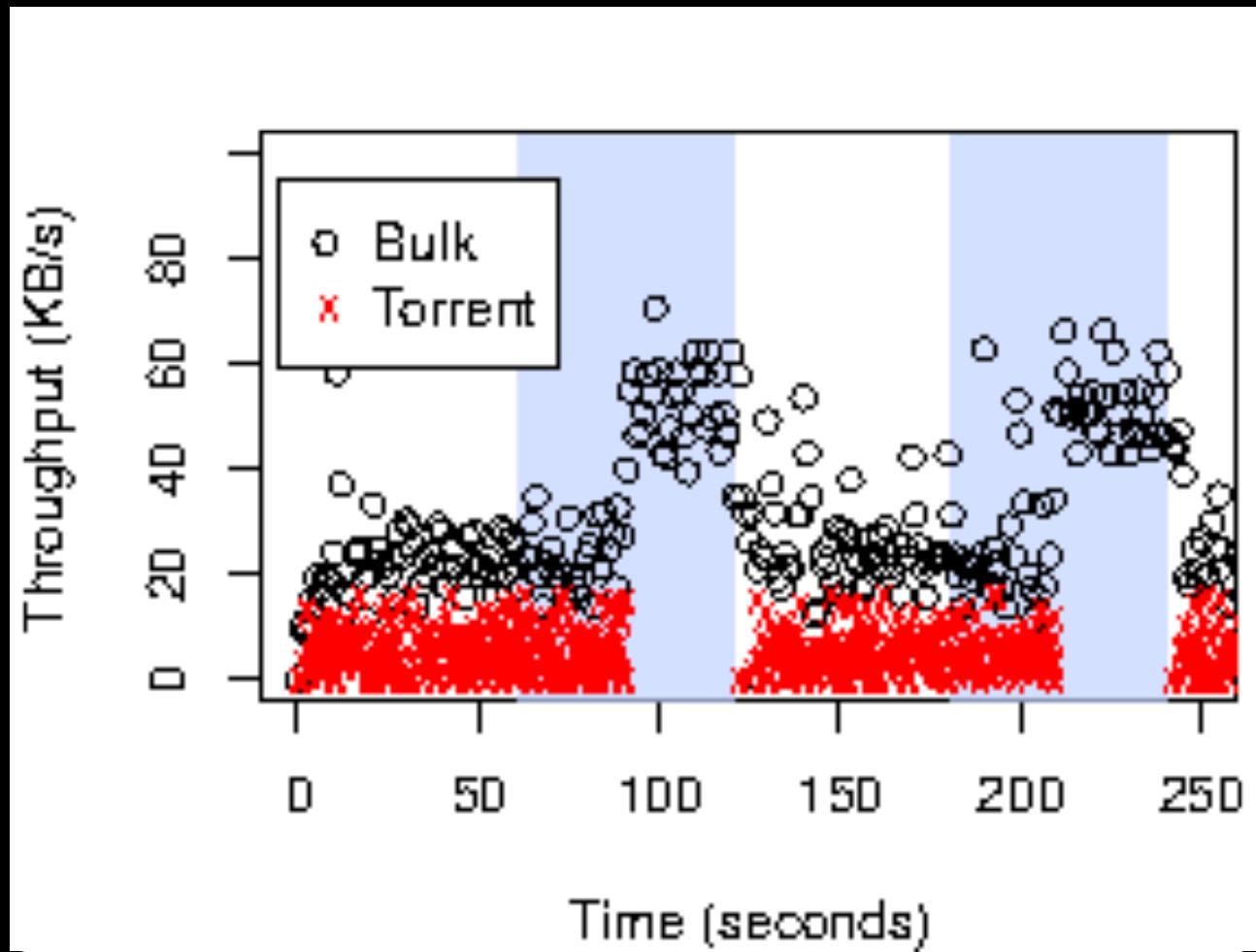


# Congestion Control

- Is this attack “stealthy”?

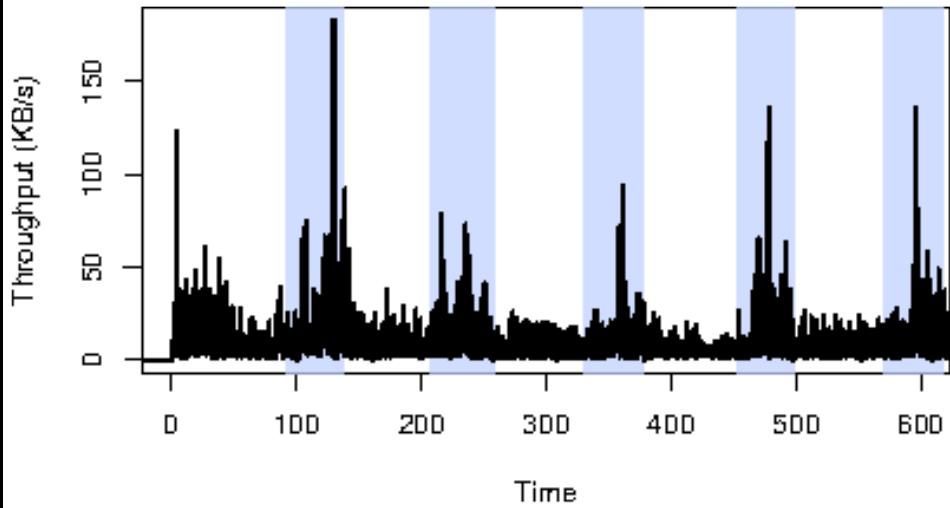


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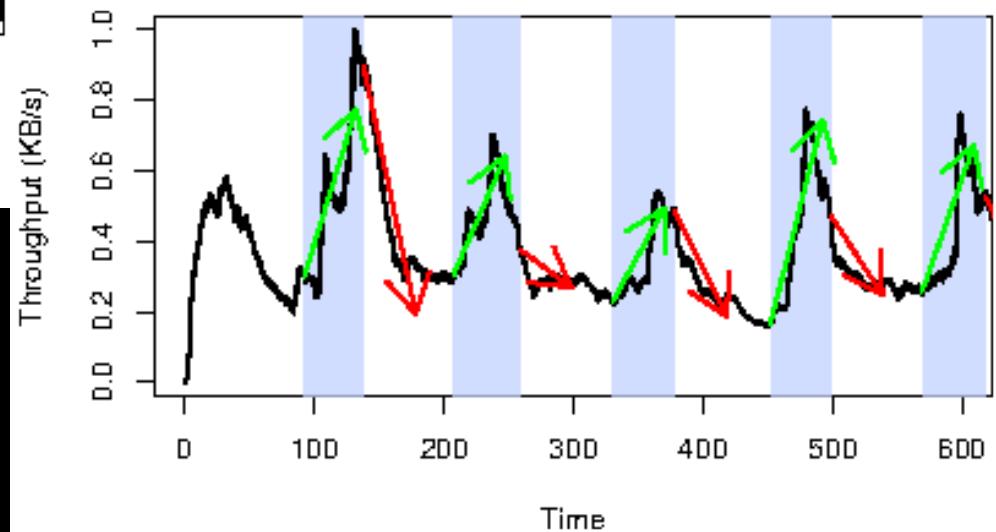


# Induced Throttling Results

Raw throughput



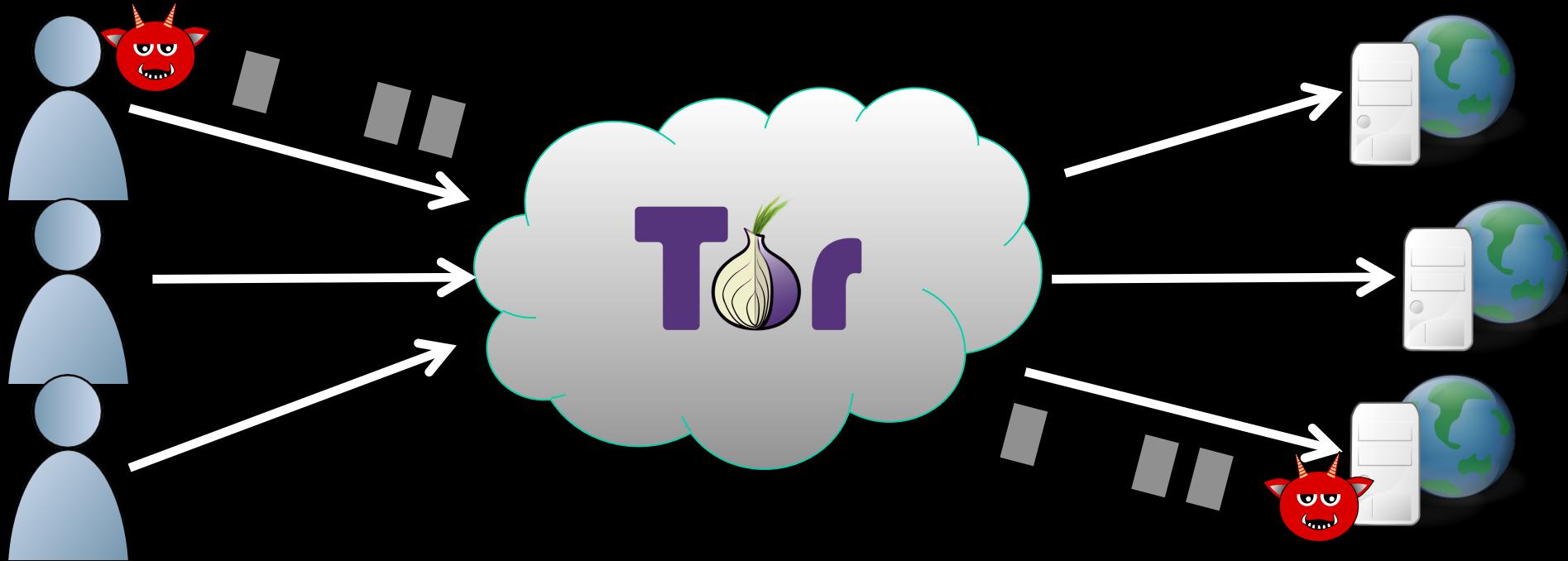
Smoothed throughput



# Outline

- ~~Background~~
- ~~Security against correlation (end-to-end)~~
  - ~~Metrics and methodology~~
  - ~~Node adversaries~~
  - ~~Link adversaries~~
- ~~Correlation attacks (partial)~~
  - ~~Stealthy throughput~~
  - ~~Induced throttling~~
    - . ~~Traffic admission control~~
    - . ~~Congestion control~~

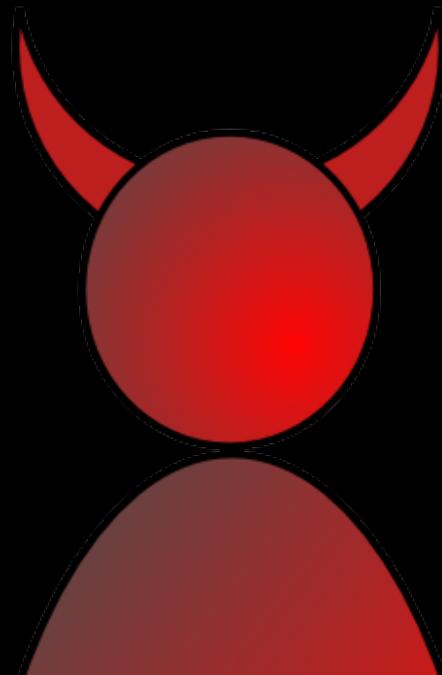
# Traffic Correlation



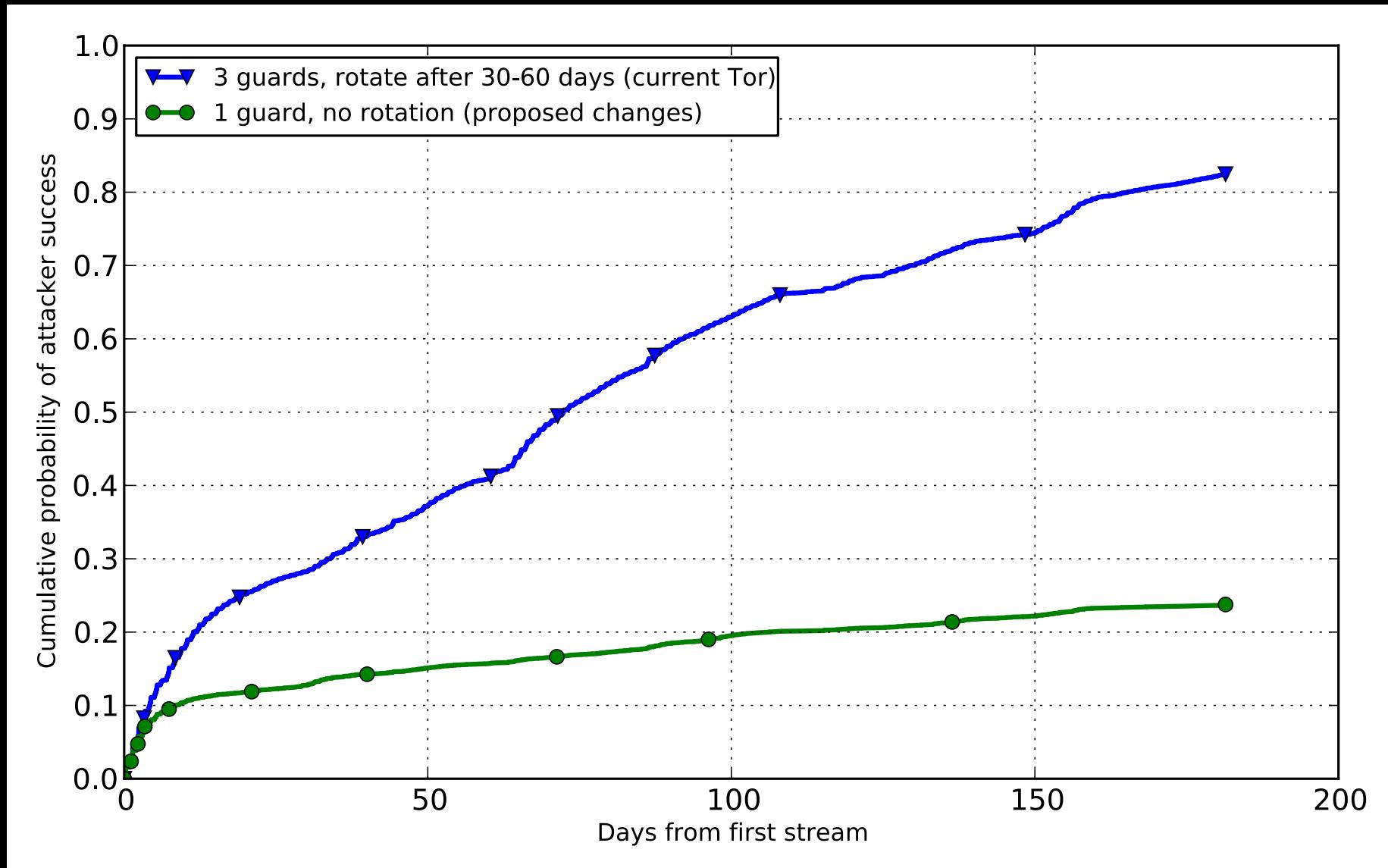
- How might we defend against ALL traffic correlation attacks?

# Questions?

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# Conclusion



# Tor is Efficient: ~65% Utilization

