# **DS** sheating-sheet

hupxer-7meqqe-hUgzik

## 工具

int(str,n) 将字符串str转换为n进制的整数。

for key,value in dict.items() 遍历字典的键值对。

for index,value in enumerate(list) 枚举列表,提供元素及其索引。

dict.get(key,default) 从字典中获取键对应的值,如果键不存在,则返回默认值 default。

list(zip(a,b)) 将两个列表元素一一配对,生成元组的列表。

math.pow(m,n) 计算m的n次幂。

math.log(m,n) 计算以n为底的m的对数。

### Irucache

```
from functools import lru_cache
@lru_cache(maxsize=None)
```

#### bisect

```
import bisect
# 创建一个有序列表
sorted_list = [1, 3, 4, 4, 5, 7]
# 使用bisect_left查找插入点
position = bisect.bisect_left(sorted_list, 4)
print(position) # 输出: 2
# 使用bisect_right查找插入点
position = bisect.bisect_right(sorted_list, 4)
print(position) # 输出: 4
# 使用insort_left插入元素
bisect.insort_left(sorted_list, 4)
print(sorted_list) # 输出: [1, 3, 4, 4, 4, 5, 7]
# 使用insort_right插入元素
```

```
bisect.insort_right(sorted_list, 4)
print(sorted_list) # 输出: [1, 3, 4, 4, 4, 5, 7]
```

## 字符串

- 1. str\_lstrip() / str\_rstrip(): 移除字符串左侧/右侧的空白字符。
- 2. str.find(sub): 返回子字符串sub在字符串中首次出现的索引,如果未找到,则返回-1。
- 3. str.replace(old, new): 将字符串中的old子字符串替换为new。
- 4. str.startswith(prefix) / str.endswith(suffix): 检查字符串是否以 prefix开头或以suffix结尾。
- 5. str.isalpha() / str.isdigit() / str.isalnum(): 检查字符串是否全部由字母/数字/字母和数字组成。
  - 6.str<sub>title()</sub>:每个单词首字母大写。

## counter: 计数

## permutations: 全排列

```
from itertools import permutations
# 创建一个可迭代对象的排列
perm = permutations([1, 2, 3])
# 打印所有排列
for p in perm:
    print(p)
# 输出: (1, 2, 3), (1, 3, 2), (2, 1, 3), (2, 3, 1), (3, 1, 2), (3, 2, 1)
```

## combinations: 组合

```
from itertools import combinations
# 创建一个可迭代对象的组合
comb = combinations([1, 2, 3], 2)
# 打印所有组合
for c in comb:
    print(c)
# 输出: (1, 2), (1, 3), (2, 3)
```

## reduce: 累次运算

```
from functools import reduce

# 使用reduce计算列表元素的乘积

product = reduce(lambda x, y: x * y, [1, 2, 3, 4])

print(product) # 输出: 24
```

## product:笛卡尔积

```
from itertools import product
# 创建两个可迭代对象的笛卡尔积
prod = product([1, 2], ['a', 'b'])
# 打印所有笛卡尔积对
for p in prod:
    print(p)
# 输出: (1, 'a'), (1, 'b'), (2, 'a'), (2, 'b')
```

### defaultdict

defaultdict 是另一种字典子类,它提供了一个默认值,用于字典所尝试访问的键不存在时返回。

```
from collections import defaultdict

# 使用 lambda 来指定默认值为 0
d = defaultdict(lambda: 0)

d['key1'] = 5
print(d['key1']) # 输出: 5
print(d['key2']) # 输出: 0, 因为 key2 不存在, 返回默认值 0
```

## namedtuple

namedtuple生成可以使用名字来访问元素内容的元组子类。

```
from collections import namedtuple

Point = namedtuple('Point', ['x', 'y'])
p = Point(11, y=22)

print(p.x + p.y) # 输出: 33
print(p[0] + p[1]) # 输出: 33 # 还可以像普通元组那样用索引访问
```

#### **OrderedDict**

OrderedDict 是一个字典子类,它保持了元素被添加的顺序,这在某些情况下非常有用。

```
from collections import OrderedDict

od = OrderedDict()
od['z'] = 1
od['y'] = 2
od['x'] = 3

for key in od:
    print(key, od[key])
# 输出:
# z 1
# y 2
# x 3
```

## heapq

- 1. heapify(x)
- 用途: 将列表 x 原地转换为堆。
- 示例

```
import heapq
data = [3, 1, 4, 1, 5, 9, 2, 6, 5]
heapq.heapify(data)
print(data) # 输出将是堆,但可能不是完全排序的
```

- 2. heappush(heap, item)
- 用途:将 item 加入到堆 heap中,并保持堆的不变性。

• 示例

```
heap = []
heapq.heappush(heap, 3)
heapq.heappush(heap, 1)
heapq.heappush(heap, 4)
print(heap) # 输出最小元素总是在索引0
```

- 3. heappop(heap)
- 用途: 弹出并返回 heap 中最小的元素, 保持堆的不变性。
- 示例

```
print(heapq.heappop(heap)) # 返回1
print(heap) # 剩余的堆
```

- 4. heapreplace(heap, item)
- **用途**: 弹出堆中最小的元素,并将新的 item 插入堆中,效率高于先 heappop() 后 heappush()。
- 示例

```
heapq.heapreplace(heap, 7)
print(heap)
```

- 5. heappushpop(heap, item)
- 用途: 先将 item 压入堆中, 然后弹出并返回堆中最小的元素。
- 示例

```
result = heapq.heappushpop(heap, 0)
print(result) # 输出0
print(heap) # 剩余的堆
```

- 6. nlargest(n, iterable, key=None) 和 nsmallest(n, iterable, key=None)
- 用途: 从 iterable 数据中找出最大的或最小的 n 个元素。
- 示例

```
data = [3, 1, 4, 1, 5, 9, 2, 6, 5]
print(heapq.nlargest(3, data)) # 输出[9, 6, 5]
print(heapq.nsmallest(3, data)) # 输出[1, 1, 2]
```

#### 注意事项

- 如需实现最大堆功能,可以通过对元素取反来实现。将所有元素取负后使用 heapq,然 后再取负回来即可。
- 堆操作的时间复杂度一般为 O(log n), 适合处理大数据集。
- heapq 只能保证列表中的第一个元素是最小的,其他元素的排序并不严格。

```
import queue

# 创建一个 LIFO 队列
lifo_queue = queue.LifoQueue()

# 添加元素
lifo_queue.put('a')
lifo_queue.put('b')
lifo_queue.put('c')

# 依次取出元素
print(lifo_queue.get()) # 输出 'c'
print(lifo_queue.get()) # 输出 'b'
print(lifo_queue.get()) # 输出 'a'
```

#### math

gcd包, 计算最大公因式

```
from math import gcd
x = gcd(15,20,25)
print(x)
## 5
```

math.pow(m,n) 计算m的n次幂。

math.log(m,n) 计算以n为底的m的对数。

#### eval

eval()是 python 中功能非常强大的一个函数

将字符串当成有效的表达式来求值,并返回计算结果

所谓表达式就是: eval 这个函数会把里面的字符串参数的引号去掉,把中间的内容当成 Python的代码, eval 函数会执行这段代码并且返回执行结果

也可以这样来理解: eval() 函数就是实现 list、dict、tuple、与str 之间的转化

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```
result = eval("1 + 1")
print(result) # 2

result = eval("'+' * 5")
print(result) # +++++

# 3. 将字符串转换成列表
a = "[1, 2, 3, 4]"
result = type(eval(a))
print(result) # <class 'list'>

input_number = input("请输入一个加减乘除运算公式: ")
print(eval(input_number))
## 1*2 +3
## 5
```

## 埃氏筛法,得到质数表

```
def judge(number):
    nlist = list(range(1,number+1))
    nlist[0] = 0
    k = 2
    while k * k <= number:
        if nlist[k-1] != 0:
            for i in range(2*k,number+1,k):
                nlist[i-1] = 0
        k += 1
    result = []
    for num in nlist:
        if num != 0:
            result.append(num)
    return result</pre>
```

## print保留小数

```
print("%.6f" % x)
print("{:.6f}".format(result))
# 当输出内容很多时:
print('\n'.join(map(str, ans)))
```

## 动态规划

## 背包问题

```
# 0-1
dp = [0 \text{ for i in } range(m+1)]
for i in range(1,n+1):
    for j in range(m, w[i]-1,-1):
         dp[j] = \max(dp[j], dp[j-w[i]]+v[i])
# complete
dp = [0 \text{ for i in } range(m+1)]
for i in range(1,n+1):
    for j in range(w[i],m+1):
         dp[j] = max(dp[j], dp[j-w[i]]+v[i])
# multi
W = [None]
V = [None]
for i in range(1,n+1):
    w,v,num = map(int,input().strip().split())
    k = 1
    while k <= num:
         num = num - k
         W<sub>∗</sub> append (k*w)
         V<sub>■</sub> append (k*v)
         k = k*2
    if num > 0:
         W<sub>■</sub> append (num*w)
         V<sub>■</sub>append(num*v)
dp = [0 \text{ for i in } range(m+1)]
for i in range(1,len(W)):
    for j in range(m,W[i]-1,-1):
         dp[j] = max(dp[j], dp[j-W[i]]+V[i])
```

## 最长下降子列

```
# 二分版
from bisect import *

n = int(input())
a = list(map(int, input().split()))
ans = 0
b = [-a[0]]
for i in range(1, n):
    if -a[i] > b[-1]:
        b.append(-a[i])
    else:
        pos = bisect_left(b, -a[i])
        b[pos] = -a[i]
print(len(b))
```

## 单调栈

维护所有前缀的后缀最值

```
import bisect
n = int(input())
a = [0] * n
for i in range(n):
    a[i] = int(input())
st1 = []
st2 = []
ans = 0
for i in range(n):
    while st1 and a[i] \le a[st1[-1]]:
        st1.pop()
    while st2 and a[i] > a[st2[-1]]:
        st2.pop()
    if st2:
        k = bisect_bisect_right(st1, st2[-1])
        if k < len(st1):
            ans = max(ans, i - st1[k] + 1)
    elif st1:
```

```
ans = max(ans, i - st1[0] + 1)
st1.append(i)
st2.append(i)
print(ans)
```

## 排序

归并排序

```
def merge_sort(lst):
    # The list is already sorted if it contains a single element.
    if len(lst) <= 1:</pre>
        return lst, 0
    # Divide the input into two halves.
    middle = len(lst) // 2
    left, inv_left = merge_sort(lst[:middle])
    right, inv_right = merge_sort(lst[middle:])
    merged, inv_merge = merge(left, right)
    # The total number of inversions is the sum of inversions in the
recursion and the merge process.
    return merged, inv_left + inv_right + inv_merge
def merge(left, right):
    merged = []
    inv_count = 0
    i = j = 0
    # Merge smaller elements first.
    while i < len(left) and j < len(right):
        if left[i] <= right[j]:</pre>
            merged.append(left[i])
            i += 1
        else:
            merged.append(right[j])
            inv_count += len(left) - i #left[i~mid)都比right[j]要大, 他们都
会与right[j]构成逆序对,将他们加入答案
```

```
# If there are remaining elements in the left or right half, append
them to the result.
    merged += left[i:]
    merged += right[j:]

    return merged, inv_count

while True:
    n = int(input())
    if n == 0:
        break

lst = []
    for _ in range(n):
        lst.append(int(input()))

_, inversions = merge_sort(lst)
    print(inversions)
```

## 树算法

## 四种遍历

前(中、后) 序遍历

```
def __show(self,root):
    if root!=None:
        print(root.data,end=',')
        self.__show(root.left)
        self.__show(root.right)
    else:
        return 0
```

层次遍历

```
def level_order(self,root):
    queue = deque()
    queue.append(root) # 将根节点放到队列之中
    while len(queue) > 0: # 只要队列不空
        node = queue.popleft() # 队头出队
        print(node.data, end=',')
        if node.left: # 出队的节点存在左孩子,就把左孩子入队
            queue.append(node.left)
        if node.right: # 如果存在右孩子,就把右孩子入队
            queue.append(node.right)
```

## 建树

current移动法

```
class Node:
    def __init__(self, _v):
        self_v = v
        self.left = None
        self right = None
        self.father = None
n = int(input())
for _ in range(n):
    root = Node(input())
    cu = root
    cu h = 1
    while True:
        s = input()
        if s == '0':
            break
        nn = Node(s[-1])
        while cu_h != len(s) - 1:
            cu = cu.father
            cu h -= 1
        if cu.left:
             cu₁right = nn
        else:
             cu_{\bullet}left = nn
        if s[-1] != '*':
            nn_{\bullet}father = cu
```

```
cu = nn
cu_h += 1
```

#### 递归建树

```
class Node:
    def __init__(self, _v):
        self_v = v
        self.left = None
        self.right = None
index = 0
def build():
    global index
    root = Node(s[index])
    index += 1
    if root.v != '#':
        root left = build()
        root.right = build()
    return root
while True:
    try:
        n = int(input())
        if n == 0:
            break
        s = input().split()
        index = 0
        build()
        print('T' if index == n else 'F')
    except:
        print('F')
```

#### 前中建树 (转后)

```
def build_tree(preorder, inorder):
   if not preorder:
     return ''
```

```
root = preorder[0]
root_index = inorder.index(root)

left_preorder = preorder[1:1 + root_index]
right_preorder = preorder[1 + root_index:]

left_inorder = inorder[:root_index]
right_inorder = inorder[root_index + 1:]

left_tree = build_tree(left_preorder, left_inorder)
right_tree = build_tree(right_preorder, right_inorder)
return left_tree + right_tree + root
```

#### 中后建树

```
class Node:
    def __init__(self, data):
        self.data = data
        self.left = None
        self.right = None

def build_tree(inorder, postorder):
    if inorder:
        root = Node(postorder.pop())
        root_index = inorder.index(root.data)
        root.right = build_tree(inorder[root_index+1:], postorder)
        root.left = build_tree(inorder[:root_index], postorder)
        return root
```

## 二叉堆

```
class tree_node:
    def __init__(self):
        self.heap = []

def parent(self, i):
        return (i - 1) // 2

def left_child(self, i):
        return 2 * i + 1

def right_child(self, i):
```

```
return 2 * i + 2
    def swap(self, i, j):
        self.heap[i], self.heap[j] = self.heap[j], self.heap[i]
    def insert(self, item):
        self.heap.append(item)
        self_heapify up(len(self_heap) - 1)
    def delete(self):
        if len(self.heap) == 0:
            raise IndexError("Heap is empty")
        self.swap(0, len(self.heap) - 1)
        min_value = self.heap.pop()
        self heapify down(0)
        return min_value
    def heapify_up(self, i):
        while i > 0 and self_heap[i] < self_heap[self_parent(i)]:
            self.swap(i, self.parent(i))
            i = self.parent(i)
    def heapify_down(self, i):
        min_index = i
        left = self.left_child(i)
        right = self.right_child(i)
        if left < len(self_heap) and self_heap[left] <</pre>
self.heap[min_index]:
            min_index = left
        if right < len(self_heap) and self_heap[right] <</pre>
self.heap[min_index]:
            min_index = right
        if i != min index:
            self.swap(i, min_index)
            self.heapify_down(min_index)
n = int(input())
lst = tree node()
for _ in range(n):
    s = input()
    if s[0] == '1':
```

```
lst.insert(int(s[2:]))
if s[0] == '2':
    print(lst.delete())
```

## 搜索树

### 二叉搜索树

```
class TreeNode:
    def __init__(self, value):
        self.value = value
        self.left = None
        self.right = None

def insert(node, value):
    if node is None:
        return TreeNode(value)
    if value < node.value:
        node.left = insert(node.left, value)
    elif value > node.value:
        node.right = insert(node.right, value)
    return node
```

## AVL树

```
from collections import deque # 在遍历的时候使用而已,其实并不需要
class TreeNode():
   def __init__(self): # 这是没有设置父连接的版本
      self_data=0
      self_left=None
                       # 左孩子
      self.right=None # 右孩子
      self.height=0
                       # 节点高度(这里设置所有节点的初始高度为0,后续会
根据情况调整,根节点的高度会随之增加)
class BTree():
   def init (self):
      self.root=None
   def __Max(self, h1, h2): # 比较左右孩子的高度,返回最大的那个
      if h1>h2:
         return h1
```

```
elif h1<=h2:
          return h2
   def __LL(self,r):#左左情况,向右旋转 r是第一个高度差不满足的节点
      node=r.left
      r.left=node.right
      node right=r
r.height=self.__Max(self.getHeight(r.right),self.getHeight(r.left))+1
node.height=self.__Max(self.getHeight(node.right),self.getHeight(node.l
eft))+1
      return node # 返回的是最开始r的左孩子,也是现在子树中的根节点
   def ___RR(self,r):#右右,左旋
                                 r是第一个高度差不满足条件的节点
      node = r.right
      r.right = node.left
      node_left = r
      r.height = self.__Max(self.getHeight(r.right),
self.getHeight(r.left)) + 1
      node.height = self.__Max(self.getHeight(node.right),
self.getHeight(node.left)) + 1
                         # 返回的是最开始r的右孩子,也是现在子树中的根节点
      return node
   def _LR(self,r):#左右,先左旋再右旋 r是第一个高度差不满足条件的节点
      r.left=self.__RR(r.left)
                                   # 对r的左孩子进行左旋
      return self.__LL(r)
                                   # 对r进行右旋,随即返回
   def __RL(self,r):#右左, 先右旋再左旋
                                   r是第一个高度差不满足条件的节点
      r.right=self.__LL(r.right) # 对r的右孩子进行右旋
      return self<sub>----</sub>RR(r)
                                   # 对r进行左旋, 随即返回
   def insert(self,data,r):
      if r==None:
                           # 假入没有根节点,就新创建一个二叉树,插入值作
为根节点的值
         node=TreeNode()
         node.data=data
         return node
      elif data==r.data: # 假设插入值和根节点相同,那么就直接返回根节点
          return r
      elif data<r.data:
                           # 假设插入值比根节点的值要小,使用递归迭代,
          r.left=self.__insert(data,r.left)
          if self.getHeight(r.left)-self.getHeight(r.right)>=2: # 如
果左孩子的高度 比 右孩子的高度 大于等于2;简单的说其实是在判定是在左孩子的子树进行插入
             if data<r.left.data: # 插入数值比左孩子要小(即在左子树中插
入),使用右旋
                r=self<sub>*</sub>__LL(r)
                                # 插入数值比右孩子要大(即在右子树中插
             else:
入),使用左旋-右旋
```

```
r=self<sub>*</sub>__LR(r)
      else:
          r.right=self.__insert(data,r.right)
          if self.getHeight(r.right)-self.getHeight(r.left)>=2: # 如
果右孩子的高度 比 左孩子的高度 大于等于2;简单的说其实是在判定是在右孩子的子树进行插入
             if data>r.right.data: # 插入数值比右孩子要大(即在右子树中
插入),使用左旋
                 r=self RR(r)
                                 # 插入数值比左孩子要小(即在左子树中
             else:
插入),使用右旋-左旋
                 r=self<sub>-</sub>__RL(r)
r.height=self.__Max(self.getHeight(r.left),self.getHeight(r.right))+1
  # 修正树的根节点的深度
      return r
   # 删除data节点
   def __delete(self,data,r):
      if r==None:
                              # 假如节点为空, 就返回
          print("don't have %d"%data)
          return r
      elif r.data==data:
                             # 当节点的值和删除值相同时
          if r.left==None: #如果被删除的节点只有右子树,直接将右子树赋
值到此节点
             return r.right
          elif r.right==None: #如果被删除的节点只有左子树,直接将左子树赋
值到此节点
             return r.left
          else:#如果同时有左右子树
             if self.getHeight(r.left)>self.getHeight(r.right): #左子
树高度大于右子树
                 # 找到左子树的最右节点(最大值) 返回节点值 并删除该节点
                 node=r_left
                 while(node.right!=None):
                    node=node right
                 r=self.__delete(node.data,r) # 调用自身删除node
                 r.data=node.data
                 return r
             else:
                                        # 右子树高度大于左子树
                 node=r₁right
                 # 找到右子树的最小节点(最小值) 返回节点值 并删除该节点
                 while node.left!=None:
                    node=node left
                 r=self.__delete(node.data,r) # 调用自身删除node
                 r.data=node.data
```

```
return r
                                   # 当删除值小于根节点的值时
      elif data<r.data:
          r.left=self.__delete(data,r.left)
                                               # 在左子树中删除,使
用递归删除
          if self.getHeight(r.right)-self.getHeight(r.left)>=2: # 删
除后,如果右子树高度与左子树高度相差超过1
             if
self.getHeight(r.right.left)>self.getHeight(r.right.right):
                 r=self<sub>-</sub>__RL(r)
                                 # 第一个错误点在右孩子的左子树
中,使用右旋-左旋
             else:
                 r=self<sub>*</sub>__RR(r)
                                      # 第一个错误点在右孩子的右子树
中,使用左旋
      elif data>r.data:
                                 # 当删除值大于根节点的值时
          r.right=self.__delete(data,r.right)
                                               # 右子树中删除
          if self.getHeight(r.left)-self.getHeight(r.right)>=2:
# 左子树与右子树高度相差超过1
             if
self.getHeight(r.left.right)>self.getHeight(r.left.left):
                 r=self<sub>*</sub>__LR(r)
                                      # 第一个错误点在左孩子的右子树
中,使用左旋-右旋
             else:
                 r=self.__LL(r) # 第一个错误点在左孩子的左子树
中,使用右旋
r.height=self.__Max(self.getHeight(r.left),self.getHeight(r.right))+1
        # 根节点的高度为 左右孩子节点最大的那个+1
      return r
   def Insert(self,data):
                                  # 插入操作
      self.root=self.__insert(data, self.root) # 把新插入的节点
从根节点开始比较
      return self.root # 返回根节点
   def Delete(self,data):
                                  # 删除操作
      self.root=self.__delete(data,self.root)
                                         # 从根节点开始寻
找,找到要删除的节点位置
      return self root
```

## 哈夫曼编码

```
import heapq
class Node:
```

```
def __init__(self, weight, char=None):
        self.weight = weight
        self.char = char
        self.left = None
        self right = None
    def __lt__(self, other):
        if self.weight == other.weight:
            return self.char < other.char
        return self_weight < other_weight
def build_huffman_tree(characters):
    heap = []
    for char, weight in characters.items():
        heapq heappush (heap, Node (weight, char))
    while len(heap) > 1:
        left = heapq.heappop(heap)
        right = heapq heappop(heap)
        merged = Node(left.weight + right.weight) #note: 合并后, char 字段
默认值是空
        merged.left = left
        merged right = right
        heapq.heappush(heap, merged)
    return heap[0]
def encode_huffman_tree(root):
    codes = {}
    def traverse(node, code):
        if node.char:
            codes[node_char] = code
        else:
            traverse(node left, code + '0')
            traverse(node right, code + '1')
    traverse(root, '')
    return codes
def huffman_encoding(codes, string):
    encoded = ''
    for char in string:
        encoded += codes[char]
    return encoded
```

```
def huffman_decoding(root, encoded_string):
    decoded = ''
    node = root
    for bit in encoded_string:
        if bit == '0':
            node = node.left
        else:
            node = node.right

    if node.char:
        decoded += node.char
        node = root
    return decoded
```

## 并查集

### 普诵并查集

```
def find set(x): #有路径压缩优化的查询
   if x != s[x]:
                           # 不等于自己的集
       s[x] = find_set(s[x]) # 把集改成根节点的集
   return s[x]
def merge_set(x,y): #合并
   x = find_set(x)
   y = find_set(y)
   if x != y : s[x] = s[y]
n,m = map(int,input().split())
s = list (range(10**6)) # 大小看题目要求,最大规模10**6,或者n+10也可以
for i in range(m):
   op,x,y = map(int,input().split())
   if op == 1:
       merge_set(x,y)
   else:
       if find_set(x) == find_set(y): # 两个人的组织相同,则是朋友
          print("YES")
       else:
          print("NO")
```

## 带权并查集

```
def find(_x):
    if _x != pre[_x]:
        t = find(pre[_x])
        h[_x] = (h[pre[_x]] + h[_x]) % multi_num
        pre[_x] = t
    return pre[_x]
def merge(_a, _b, sign):
    _{fa}, _{fb} = find(_{a}), find(_{b})
    pre[_fa] = _fb
    h[_fa] = h[_b] - h[_a] + sign
def same(_a, _b):
    return (h[a] - h[b]) % multi_num == 0 and find(_a) == find(_b)
T = int(input())
multi_num = 2
for in range(T):
    n, k = map(int, input().split())
    pre = [i for i in range(n)]
    h = [0 \text{ for } \underline{\ } \text{ in } range(n)]
    for __ in range(k):
        op, a, b = input().split()
        a, b = int(a) - 1, int(b) - 1
        fa, fb = find(a), find(b)
        if op == 'A':
             if find(a) != find(b):
                 print('Not sure yet.')
             elif same(a, b):
                 print('In the same gang.')
             else:
                 print('In different gangs.')
        elif op == 'D':
             merge(a, b, 1)
```

## 中序转后序

```
def infix_to_postfix(expression):
```

```
precedence = {'+':1, '-':1, '*':2, '/':2}
    stack = []
    postfix = []
    number = ''
    for char in expression:
        if char.isnumeric() or char == '.':
            number += char
        else:
            if number:
                num = float(number)
                postfix append(int(num) if num is_integer() else num)
                number = ''
            if char in '+-*/':
                while stack and stack[-1] in '+-*/' and precedence[char]
<= precedence[stack[-1]]:</pre>
                    postfix.append(stack.pop())
                stack_append(char)
            elif char == '(':
                stack append(char)
            elif char == ')':
                while stack and stack[-1] != '(':
                    postfix.append(stack.pop())
                stack.pop()
    if number:
        num = float(number)
        postfix.append(int(num) if num.is_integer() else num)
    while stack:
        postfix.append(stack.pop())
    return ' '.join(str(x) for x in postfix)
n = int(input())
for _ in range(n):
    expression = input()
    print(infix_to_postfix(expression))
```

## 图算法

```
q = deque()
q append((0, 0, 0))
v = set()
steps = -1
while q:
    x0, y0, h = q.popleft()
    if pct[x0][y0] == 1:
        steps = h
        break
    v<sub>a</sub>add((x0, y0))
    for dx, dy in [(1, 0), (-1, 0), (0, 1), (0, -1)]:
        x1 = x0 + dx
        y1 = y0 + dy
        if 0 \le x1 \le m and 0 \le y1 \le n and pct[x1][y1] != 2 and (x1, y1)
not in v:
             q_aappend((x1, y1, h + 1))
```

## Dijkstra

```
def dijkstra(graph: List[List[int]], n: int, k: int):
 # 稀疏图用邻接表graph, 其中graph[node]存储以(nb, w)形式存储邻居和边权
 expanded = [False for _ in range(n)]
 curDist = [float('inf') for _ in range(n)]
 curDist[k] = 0
 h = [(0, k)]
 while h:
   # 1. 找到没有扩展过的点中到起点距离最短的点node
   node = heappop(h)
   if not expanded[node]: # 确保不重复扩展
     # 2. 扩展
     for nb, w in graph[node]:
       # 扩展意味着搜索树中nb(或暂时成为, 取决于搜索树类型)node的子节点
       newDist = curDist[node] + w
       if newDist < curDist[nb]: # 剪枝
         curDist[nb] = newDist
         heappush(h, (newDist, nb))
     # 3. 扩展完标记该点
     expanded[node] = True
```

```
import heapq
k, n, r = int(input()), int(input()), int(input())
def dij(g, s, e):
    dis = {v: float('inf') for v in range(1, n + 1)}
    dis[s] = 0
    q = [(0, s, 0)]
    heapq heapify(q)
    while q:
        d, now, fee = heapq.heappop(q)
        if now == n:
            return d
        for neighbor, distance, c in g[now]:
            if fee + c \leq k:
                 dis[neighbor] = distance + d
                 heapq_heappush(q, (distance + d, neighbor, fee + c))
    return -1
g = \{v: [] \text{ for } v \text{ in } range(1, n + 1)\}
for _ in range(r):
    s, e, m, j = map(int, input().split())
    g[s]append((e, m, j))
p = dij(g, 1, n)
print(p)
```

#### 记录路径

```
import heapq

def dijkstra(adjacency, start):
    # 初始化, 将其余所有顶点到起始点的距离都设为inf (无穷大)
    distances = {vertex: float('inf') for vertex in adjacency}
    # 初始化, 所有点的前一步都是None
    previous = {vertex: None for vertex in adjacency}
    # 起点到自身的距离为0
    distances[start] = 0
    # 优先队列
```

```
pq = [(0, start)]
   while pq:
       # 取出优先队列中,目前距离最小的
       current_distance, current_vertex = heapq.heappop(pq)
       # 剪枝, 如果优先队列里保存的距离大于目前更新后的距离, 则可以跳过
       if current distance > distances[current vertex]:
           continue
       # 对当前节点的所有邻居,如果距离更优,将他们放入优先队列中
       for neighbor, weight in adjacency[current_vertex].items():
           distance = current distance + weight
           if distance < distances[neighbor]:</pre>
               distances[neighbor] = distance
               # 这一步用来记录每个节点的前一步
               previous[neighbor] = current_vertex
               heapq.heappush(pq, (distance, neighbor))
    return distances, previous
def shortest_path_to(adjacency, start, end):
   # 逐步访问每个节点上一步
   distances, previous = dijkstra(adjacency, start)
   path = []
   current = end
   while previous[current] is not None:
       path insert(0, current)
       current = previous[current]
   path insert(0, start)
    return path, distances[end]
#Read the input data
P = int(input())
places = {input().strip() for _ in range(P)}
Q = int(input())
graph = {place: {} for place in places}
for _ in range(Q):
   src, dest, dist = input().split()
   dist = int(dist)
   graph[src][dest] = dist
   graph[dest][src] = dist # Assuming the graph is bidirectional
R = int(input())
requests = [input().split() for _ in range(R)]
```

```
#Process each request
for start, end in requests:
    if start == end:
        print(start)
        continue

path, total_dist = shortest_path_to(graph, start, end)
    output = ""
    for i in range(len(path) - 1):
        output += f"{path[i]}->({graph[path[i]][path[i+1]]})->"
    output += f"{end}"
    print(output)
```

### **Floyd**

```
def floyd_warshall(graph):
    n = len(graph)
    dist = [[float('inf')] * n for _ in range(n)]

for i in range(n):
    for j in range(n):
        if i == j:
            dist[i][j] = 0
        elif j in graph[i]:
            dist[i][j] = graph[i][j]

for k in range(n):
    for i in range(n):
        for j in range(n):
            dist[i][j] = min(dist[i][j], dist[i][k] + dist[k][j])

return dist
```

## 最小生成树

#### Prim

```
def prim(graph: List[List[int]], n: int):
# 稠密图用邻接矩阵graph, 其中存边权, 无边存无穷大
curDist = [float('inf') for _ in range(n)] # 点到当前树的最小距离, 是边权
inMST = [False for _ in range(n)] # 标记是否已加入到MST中
```

```
totalWeight = 0
for _ in range(n): # 每次加一个点一条边到树中(第一次只加点不加边)
 # 1. 通过枚举点找到连接树和树外一点的最短边
 minNode = None
 for node in range(n):
   if not inMST[node] and (minNode is None
       \ or curDist[node] < curDist[minNode]):</pre>
     minNode = nodei
 # 2. 把最短边及其连接的树外点加入到MST中(第一次循环只加点不加边)
 if i != 0: # 当然也可将起点的 curDist 初始化为 0, 则此处无需判断
   totalWeight += curDist[minNode]
   # 如果这条最短边为inf, 就代表该树外点与树中任一点都不连通, 即原图是不连通的
 inMST[minNode] = True
 # 3. 更新树外节点到树的最小距离
 for nb in graph[minNode]:
   curDist[nb] = min(curDist[nb], graph[minNode][nb])
return totalWeight
```

```
from heapq import *
while True:
    n = int(input())
    if n == 0:
        break
    trucks = [input() for _ in range(n)]
    trucks_sort()
    sd = [[0 for _ in range(n)] for _ in range(n)]
    for i in range(n):
        for j in range(i + 1, n):
            sd[i][j] = sd[j][i] = sum(a != b for a, b in zip(trucks[i],
trucks[j]))
    # Prim
    v = [False for _ in range(n)]
    dis = [float('inf') for _ in range(n)]
    dis[0] = 0
    q = [(0, 0)]
    total weight = 0
    while q:
```

```
weight, node = heappop(q)
if v[node]:
    continue
v[node] = True
total_weight += weight
for nb in range(n):
    if nb != node and not v[nb] and dis[nb] > sd[nb][node]:
        dis[nb] = sd[nb][node]
        heappush(q, (dis[nb], nb))
print(f'The highest possible quality is 1/{total_weight}.')
```

### 拓扑排序

#### **DFS**

```
def dfs(num):
    v[num] = 1
    for neighbour in g[num]:
        if v[neighbour] == 0 and dfs(neighbour):
            return True
        if v[neighbour] == 1:
            return True
        v[num] = 2
        ans.append(num) # 最后要反转
        return False
```

#### Kahn

```
def topo_sort(graph):
    in_degree = {u:0 for u in graph}
    for u in graph:
        for v in graph[u]:
            in_degree[v] += 1
    q = deque([u for u in in_degree if in_degree[u] == 0])
    topo_order = [];flag = True
    while q:
        if len(q) > 1:
            flag = False#topo_sort不唯一确定
        u = q.popleft()
        topo_order.append(u)
        for v in graph[u]:
```

```
in_degree[v] -= 1
    if in_degree[v] == 0:
        q.append(v)
if len(topo_order) != len(graph): return 0
return topo_order if flag else None
```

### 强连通分量

```
def dfs1(graph, node, visited, stack):
    visited[node] = True
    for neighbor in graph[node]:
        if not visited[neighbor]:
            dfs1(graph, neighbor, visited, stack)
    stack.append(node)
def dfs2(graph, node, visited, component):
    visited[node] = True
    component_append(node)
    for neighbor in graph[node]:
        if not visited[neighbor]:
            dfs2(graph, neighbor, visited, component)
def kosaraju(graph):
    # Step 1: Perform first DFS to get finishing times
    stack = []
    visited = [False] * len(graph)
    for node in range(len(graph)):
        if not visited[node]:
            dfs1(graph, node, visited, stack)
    # Step 2: Transpose the graph
    transposed_graph = [[] for _ in range(len(graph))]
    for node in range(len(graph)):
        for neighbor in graph[node]:
            transposed_graph[neighbor] append(node)
    # Step 3: Perform second DFS on the transposed graph to find SCCs
    visited = [False] * len(graph)
    sccs = []
    while stack:
        node = stack pop()
        if not visited[node]:
            scc = []
            dfs2(transposed_graph, node, visited, scc)
```

```
sccs.append(scc)
return sccs

# Example
graph = [[1], [2, 4], [3, 5], [0, 6], [5], [4], [7], [5, 6]]
sccs = kosaraju(graph)
print("Strongly Connected Components:")
for scc in sccs:
    print(scc)

"""
Strongly Connected Components:
[0, 3, 2, 1]
[6, 7]
[5, 4]
"""
```

### 网络最大流

```
import numpy as np
class Node:
    def __init__(self, name, arc_dict):
        self name = name
        self.arc_dict = arc_dict
def create_node(name, next_list, flow_list):
    arc_dict = {}
    for i in range(len(next_list)):
        arc_dict[next_list[i]] = flow_list[i]
    return Node(name, arc_dict)
def create_level_graph(s, e, node_list, name_index_dict):
    level_graph = np.zeros((len(node_list), len(node_list))).tolist()
    cur_layer = [s]
    all_node = set()
    all_node add(s)
    next_layer = set()
    while len(cur_layer) > 0:
        for node_name in cur_layer:
```

```
node = node_list[name_index_dict[node_name]]
            for key in node.arc_dict.keys():
                if key not in all_node and (node.arc_dict[key] is None
or node arc_dict[key] > 0):
                    level_graph[name_index_dict[node_name]]
[name_index_dict[key]] = node_arc_dict[key]
                    next_layer.add(key)
                    all node add(key)
        cur_layer = list(next_layer)
        next_layer = set()
    return level_graph if e in all_node else None
def Dinic_Solve(s, e, node_list, name_index_dict):
    routes = []
    s index = name index dict[s]
   e_index = name_index_dict[e]
    level_graph = create_level_graph(s, e, node_list, name_index_dict)
   while level_graph is not None:
        res list = []
       while True:
            res = dfs(e_index, [s_index], None, level_graph)
            if res is None:
                break
            # 更新 level graph
            route, flow = res
            for i in range(len(route) - 1):
                if level_graph[route[i]][route[i + 1]] is not None:
                    level graph[route[i]][route[i + 1]] -= flow
            # 追加记录增广路径
            res_list_append(res)
            routes.append([[node_list[n].name for n in res[0]], res[1]])
       # 更新残存网络
        for res in res_list:
            update(res, node_list)
        # 重新构造 level graph
        level_graph = create_level_graph(s, e, node_list,
name_index_dict)
    return routes
def update(res, node_list):
    route, flow = res
   for i in range(len(route) - 1):
        n1 = node_list[route[i]]
```

```
n2 = node_list[route[i + 1]]
       # 正向更新 n1 -> n2 剩余流量减少
       if n2.name in n1.arc_dict.keys() and n1.arc_dict[n2.name] is not
None:
           n1.arc_dict[n2.name] = n1.arc_dict[n2.name] - flow
       # 反向更新 n2 -> n1 剩余流量增加
        if n1.name in n2.arc_dict.keys() and n2.arc_dict[n1.name] is not
None:
           n2.arc_dict[n1.name] = n2.arc_dict[n1.name] + flow
def dfs(e_index, cur_route, last_flow, level_graph):
    if cur_route[-1] == e_index:
        return cur_route, last_flow
   for next_node in range(len(level_graph)):
        if next_node not in cur_route:
            if level_graph[cur_route[-1]][next_node] is None or
level_graph[cur_route[-1]][next_node] > 0:
               flow = min_flow(level_graph[cur_route[-1]][next_node],
last_flow)
               cur_route append(next_node)
               res = dfs(e_index, cur_route, flow, level_graph)
                if res is not None:
                    return res
               cur_route pop(-1)
def min_flow(f1, f2):
   求两个流量的较小者
    if f1 is None:
        return f2
   elif f2 is None:
       return f1
   else:
        return min(f1, f2)
if __name__ == '__main__':
   # 格式: [节点名, 后继节点的名称, 当前节点到各个后继的流量] (None 代表流量无穷
大)
   graph = [
        ["S", ["1", "2", "3"], [None, None, None]],
        ["1", ["4"], [1]],
```

```
["2", ["4", "6"], [1, 1]],
        ["3", ["5"], [1]],
        ["4", ["1", "2", "E"], [0, 0, 1]],
        ["5", ["3", "E"], [0, 1]],
       ["6", ["2", "E"], [0, 1]],
        ["E", [], []]
   name_index_dict = dict()
   node_list = []
   for i in range(len(graph)):
       node_list.append(create_node(graph[i][0], graph[i][1], graph[i]
[2]))
       name_index_dict[graph[i][0]] = i
   # 调用算法求解最大流
   routes = Dinic_Solve("S", "E", node_list, name_index_dict)
   for i, (route, flow) in enumerate(routes):
       print(f"Route-{i + 1}: {route} , flow: {flow}")
```

## 波兰表达式

```
s = input().split()
def cal():
    cur = s.pop(0)
    if cur in "+-*/":
        return str(eval(cal() + cur + cal()))
    else:
        return cur
print("%.6f" % float(cal()))
```