

# Type Theory with Records for Natural Language Semantics: Lecture 3

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## Today's Lecture

- ▶ Adding abstract entities to the ontology.
- ▶ Basic Illocutionary Interaction in KoS
- ▶ Negation

# Outline

Adding abstract entities to the semantic ontology

Basic Illocutionary Interaction in KoS

## Adding abstract entities to the semantic ontology

Based on Chapter 3 from Jonathan Ginzburg. 2011. *The Interactive Stance: Meaning for conversation*. Oxford University Press.

## Propositions

- ▶ Propositional-like entities, more intensional than events/situations, are a necessary ingredient for accounts of illocutionary acts, as well as of attitude reports.
- ▶ Sets of situations, although somewhat more fine grained than sets of worlds, also succumb to sophisticated variants of logical omniscience (see e.g. Soames' puzzle (Soames (1985))).).
- ▶ Building on a conception articulated 30 years earlier by Austin (Austin (1961)), Barwise and Etchemendy (1987) developed a theory of propositions in which a proposition is a structured object  $prop(s, \sigma)$ , individuated in terms of a situation  $s$  and a situation type  $\sigma$ .

- (1) a.  $prop(s, \sigma)$  is true iff  $s : \sigma$  ( $s$  is of type  $\sigma$ ).  
 b.  $prop(s, \sigma)$  is false iff  $s \not:\sigma$  ( $s$  is not of type  $\sigma$ ).

## Propositions

- ▶ There are two ways to maintain the insights of an Austinian approach in TTR, implicitly Austinian or explicitly so.
- ▶ Cooper (2005) develops the former in which a proposition  $p$  is taken to be a record type.
- ▶ A witness for this type is a situation.
- ▶ On this strategy, a witness is not directly included in the semantic representation.

## Propositions

- ▶ record types *are* competitive in such a role: they are sufficiently fine-grained to distinguish identity statements that involve distinct constituents.
- ▶ In this set up substitutivity of co-referentials (2e) and cross-linguistic equivalents ((2e), the Hebrew equivalent of (2a)) can be enforced:

- (2)
- a. Enescu is identical with himself.
  - b. Poulenc is identical with himself.
  - c.  $[c : \text{Identical}(\text{enescu}, \text{enescu})]$
  - d.  $[c : \text{Identical}(\text{poulenc}, \text{poulenc})]$
  - e. He is identical with himself.
  - f. Enesku zehe leacmo.

## Lappin Concerns

TTR satisfies:

- ▶ Fine grain.
- ▶ Ultra fine grain, but refine grain via equivalence classes.
- ▶ Subtyping: a notion of entailment.



## Propositions

- ▶ A situational witness for the record type could also be deduced to explicate cases of event anaphora, as in(3); indeed, a similar strategy will be invoked when we discuss nominal anaphora.

- (3) a. A: Jo and Mo got married yesterday. It was a wonderful occasion.
- b. A: Jo's arriving next week. B: No, that's happening in about a month.

# Propositions I

- ▶ Here we develop an explicitly Austinian approach, where the situational witness is directly included in the semantic representation.
- ▶ The original Austinian conception was that  $s$  is a situation deictically indicated by a speaker making an assertion—teasing the semantic difference between implicit and explicit witnesses is a difficult semantic task.
- ▶ Some other motivation from negation later on today.

## Propositions II

- ▶ propositions can also play a role in characterizing the communicative process: in subsequent lectures we will show that (*locutionary propositions*) individuated in terms of an utterance event  $u_0$  as well as to its grammatical type  $T_{u_0}$  allows one to simultaneously define update and clarification potential for utterances.
- ▶ In this case, there are potentially many instances of distinct locutionary propositions, which need to be differentiated on the basis of the utterance token—minimally any two utterances classified as being of the same type by the grammar.

## Propositions

- ▶ TTR offers a straightforward way for us to model propositions using records. A proposition is a record of the form in (4a). The type of propositions is the record type (4b):

$$(4) \text{ a. } \left[ \begin{array}{l} \text{sit} = r_0 \\ \text{sit-type} = p_0 \end{array} \right]$$

$$\text{b. Prop} = \left[ \begin{array}{l} \text{sit} : \text{Record} \\ \text{sit-type} : \text{RecType} \end{array} \right]$$

# Propositions

- ▶ TTR offers a straightforward way for us to model propositions using records. A proposition is a record of the form in (6a). The type of propositions is the record type (6b):

$$(6) \text{ a. } \left[ \begin{array}{l} \text{sit} = r_0 \\ \text{sit-type} = p_0 \end{array} \right]$$

$$\text{b. Prop} = \left[ \begin{array}{l} \text{sit} : \text{Record} \\ \text{sit-type} : \text{RecType} \end{array} \right]$$

- ▶ Truth:

$$(7) \quad \text{A proposition } \left[ \begin{array}{l} \text{sit} = r_0 \\ \text{sit-type} = p_0 \end{array} \right] \text{ is true iff } r_0 : p_0$$

## Meets, Joins, and Compound Propositions

- ▶ Meet and join types are useful for various applications.
- ▶ Here defined in the traditional Boolean way:
  - (8) a. **Meet/Join types**: if  $T_1$  and  $T_2$  are types, then so are  $T_1 \wedge T_2$  and  $T_1 \vee T_2$ .
  - b.  $a : T_1 \wedge T_2$  ( $T_1 \vee T_2$ ) iff  $a : T_1$  and (or)  $a : T_2$ .

# Disjunctive Propositions

- ▶ a disjunctive witness—one of the situational witnesses, depending which of these is indeed a witness:

$$\begin{aligned}
 (9) \quad & \text{Given } s_1, s_2 : \text{Rec} \text{ and } T_1, T_2 : \text{RecType}, \\
 & \text{Disjrec}((s_1, T_1), (s_2, T_2)) =_{\text{def}} \begin{aligned} & s_1 \text{ if } s_1 : T_1 \\ & \text{otherwise: } s_2 \text{ if } s_2 : T_2 \\ & \quad \llbracket \text{(the empty record)} \rrbracket \\ & \text{otherwise} \end{aligned}
 \end{aligned}$$

## Disjunctive Propositions

- define the disjoined proposition (the operation  $\vee_{prop}$ ) as follows:

$$(10) \quad \begin{bmatrix} \text{sit} = s_1 \\ \text{sit-type} = T_1 \end{bmatrix} \vee_{prop} \begin{bmatrix} \text{sit} = s_2 \\ \text{sit-type} = T_2 \end{bmatrix} =_{def} \begin{bmatrix} \text{sit} = \text{Disjrec}((s_1, T_1), (s_2, T_2)) \\ \text{sit-type} = T_1 \vee T_2 \end{bmatrix}$$

- It is fairly straightforward to show that the disjunctive proposition is true iff at least one of its disjuncts is true.



## Disjunctive Propositions

- ▶ What of the somewhat contingent identity conditions of this disjunctive proposition (or rather its situational component)?
- ▶ In (11) the situational anaphor 'that' seems able to refer to this split eventuality:

(11)     A: Bill left town or his bicycle has been stolen.  
          B: Yeah, that must have happened more than five  
          hours ago.

## Conjunctive Propositions

- ▶ Either: introduce a mereological structure, as commonly assumed in the literature on plurals Landman (1996)
- ▶ Or: define the conjoined proposition (the operation  $\wedge_{prop}$ ) as follows:

$$(12) \quad \left[ \begin{array}{l} \text{sit} = s_1 \\ \text{sit-type} = T_1 \end{array} \right] \wedge_{prop} \left[ \begin{array}{l} \text{sit} = s_2 \\ \text{sit-type} = T_2 \end{array} \right] =_{def} \left[ \begin{array}{l} \text{sit} = s_1 \cup s_2 \\ \text{sit-type} = T_1 \wedge T_2 \end{array} \right]$$

- ▶ Or:  $=_{def} \left[ \begin{array}{l} \text{sit} = \langle s_1, s_2 \rangle \\ \text{sit-type} = T_1 \wedge T_2 \end{array} \right]$

assuming:  $\langle s_1, s_2 \rangle$  is a witness for  $T_1 \wedge T_2$  (classical  
Constructive Type Theory)

## Questions in TTR

- ▶ Given the existence of Austinian-like propositions and a theory of  $\lambda$ -abstraction given to us by existence of functional types, it is relatively straightforward to develop a theory of questions as propositional abstracts in TTR, building on earlier work in situation theory in Ginzburg (1995); Ginzburg and Sag (2000).

## Questions as Propositional Abstracts

- ▶ any theory of questions needs to associate an abstract-like object with interrogatives to explicate the resolution of short answers.

(13) a. A: Who attended the meeting?

B: Mo/No students./A friend of Jo's.

b. A: When did Bo leave?

B: Yesterday./At two.

c. A: Why did Maire cross the road?

B: Because she thought no cars were passing.

- ▶ explicating answerhood involves a characterization that utilizes in some form an abstract-like object, as will be demonstrated below.
- ▶ If questions simply *are* abstracts, then the fact that abstracts need to be associated with interrogative utterances follows automatically.

## Questions: some simple examples

- A question will be a function from records into propositions:

(14) a. who ran

b. TTR representation— $(r : \begin{bmatrix} x : \text{Ind} \\ c1 : \text{person}(x) \end{bmatrix})$

$$\begin{bmatrix} \text{sit} = r_1 \\ \text{sit-type} = \begin{bmatrix} c : \text{run}(r.x) \end{bmatrix} \end{bmatrix}$$

That is, a function that maps records  $r : T_{\text{who}} =$

$\begin{bmatrix} x : \text{Ind} \\ c1 : \text{person}(x) \end{bmatrix}$  into propositions of the form

$$\begin{bmatrix} \text{sit} = r_1 \\ \text{sit-type} = \begin{bmatrix} c : \text{run}(r.x) \end{bmatrix} \end{bmatrix}$$

## Questions: some simple examples

(15) a. who saw what

b. TTR representation— $(r : \begin{bmatrix} x : \text{Ind} \\ c1 : \text{person}(x) \\ y : \text{Ind} \\ c2 : \text{thing}(y) \end{bmatrix})$

$\begin{bmatrix} \text{sit} = r_1 \\ \text{sit-type} = [c : \text{saw}(r.x, r.y)] \end{bmatrix}$

## Questions: some simple examples

- And, by extension, a question will be a 0-ary propositional abstract.

(16) a. Did Bo run

b. TTR representation— $(r : \left[ \begin{array}{c} \end{array} \right]) \left[ \begin{array}{c} \text{sit} = r_1 \\ \text{sit-type} = [c : \text{run}(b)] \end{array} \right]$ .

That is, a function that maps records  $r : T_0 = \left[ \begin{array}{c} \end{array} \right]$

into propositions of the form  $\left[ \begin{array}{c} \text{sit} = r_1 \\ \text{sit-type} = [c : \text{run}(b)] \end{array} \right]$

## Answerhood

- ▶ There are a number of notions of answerhood that are of importance to dialogue.
- ▶ We concern ourself here with one that relates to coherence: any speaker of a given language can recognize, independently of domain knowledge and of the goals underlying an interaction, that certain propositions are *about* or *directly concern* a given question.



## Simple Answerhood in TTR

(17) Given a question  $q : (A)B$ :

- a.  $\text{AtomAns}(q) =_{\text{def}} \left[ \begin{array}{l} \text{shortans} : A \\ \text{propans} = q(\text{shortans}) : \text{Prop} \end{array} \right]$
- b.  $\text{NegAtomAns}(q) =_{\text{def}} \left[ \begin{array}{l} \text{shortans} : A \\ \text{propans} = \neg q(\text{shortans}) : \text{Prop} \end{array} \right]$
- c.  $\text{SimpleAns}(q) =_{\text{def}} \text{AtomAns}(q) \vee \text{NegAtomAns}(q)$

## Simple Answerhood

- ▶ Simple answerhood covers a fair amount of ground. But it clearly underdetermines aboutness.
- ▶ On the polar front, it leaves out the whole gamut of answers to polar questions that are weaker than  $p$  or  $\neg p$  such as conditional answers 'If  $r$ , then  $p$ ' (e.g. 18a) or weakly modalized answers 'probably/possibly/maybe/possibly not  $p$ ' (e.g. 18b).
- ▶ As far as wh-questions go, it leaves out quantificational answers (18c-g), as well as disjunctive answers.

## Simple Answerhood v. Aboutness

- (18) a. Christopher: Can I have some ice-cream then?  
Dorothy: you can do if there is any. (BNC, KBW)
- b. Anon: Are you voting for Tory?  
Denise: I might. (BNC, KB?, slightly modified)
- c. Dorothy: What did grandma have to catch?  
Christopher: A bus. (BNC, KBW, slightly modified)
- d. Rhiannon: How much tape have you used up?  
Chris: About half of one side. (BNC, KB?)
- e. Dorothy: What do you want on this?  
Andrew: I would like some yogurt please. (BNC, KBW, slightly modified)
- f. Elinor: Where are you going to hide it?  
Tim: Somewhere you can't have it. (BNC, KBW)
- g. Christopher: Where is the box?  
Dorothy: Near the window. (BNC, KBW)

## Aboutness via disjunction

- ▶ One straightforward way to enrich simple answerhood is to consider the relation that emerges by closing simple answerhood under disjunction. See Ginzburg and Sag (2000) for more discussion.

(19)  $p$  is about  $q$  iff  $p$  entails a disjunction of simple answers.

## Aboutness via disjunction

- If  $r : \text{SimpleAns}(q)$ , then since  $r : \text{AtomAns}(q)$  or  $r : \text{NegAtomAns}(q)$ ,  $r$  is a record of the form in (i),  $p1$  then is the value  $r$  gets on the `propans` field:

$$(i) \quad r = \begin{bmatrix} \text{shortans} = a \\ \text{propans} = p \end{bmatrix}$$

(20) Given a question  $q : (A)B$ :

a.  $\text{Aboutness}(q) =_{\text{def}}$

$$\begin{bmatrix} r1 : \text{SimpleAns}(q) \\ p1 = r1.\text{propans} : \text{Prop} \\ r2 : \text{SimpleAns}(q) \\ p2 = r2.\text{propans} : \text{Prop} \\ \text{propans} : \text{Prop} \\ c : \text{Entails}(\text{propans}, p1 \vee_{\text{prop}} p2) \end{bmatrix}$$

## Answering a question with a question

- (21) a. A: Who murdered Smith? B: Who was in town?
- b. A: Who is going to win the race? B: Who is going to participate?
- c. Carol: Right, what do you want for your dinner?  
Chris: What do you (pause) suggest? (BNC, Kbj)
- d. Chris: Where's mummy?  
Emma: Mm?  
Chris: Mummy?  
Emma: What do you want her for? (BNC, Kbj)

## Answering a question with a question I

- ▶ There has been much work on relations among questions within the framework of *Inferential Erotetic Logic* (IEL) (see e.g. Wiśniewski (2001, 2003)), yielding notions of *q(uestion)–implication*.
- ▶ From this a natural hypothesis can be made about such query responses, as in (22a);
- ▶ A related proposal, first articulated by Carlson (1983), is that they are constrained by the semantic relations of *dependence*, or its converse *influence*.

(22)  $q_2$  influences  $q_1$  iff any proposition  $p$  such that  $p$  Resolves  $q_2$ , also satisfies  $p$  entails  $r$  such that  $r$  is About  $q_1$ .

## Answering a question with a question II

- ▶ Its intuitive rationale is this: discussion of  $q_2$  will necessarily bring about the provision of information about  $q_1$



# Outline

Adding abstract entities to the semantic ontology

Basic Illocutionary Interaction in KoS

## Basic Illocutionary Interaction in KoS

Based on Chapter 4 from Jonathan Ginzburg. 2011. *The Interactive Stance: Meaning for conversation*. Oxford University Press.

## Dialogue Analysis

- ▶ Dialogue analyst's task: describe conventionally acceptable patterns of interaction (*protocols*), in terms of sequences of information states.
- ▶ Methodological constraint: compositionality (but as with the sentential level not obsessively [cf. the need for constructionism]).

## Contexts in KoS

- ▶ in KoS, there is actually no single context, for reasons that will soon emerge.
- ▶ Instead of a single context, analysis is formulated at a level of information states, one per conversational participant.
- ▶ What the typed entities depicted in (23),(24) amount to is explained in the next few lectures.
- ▶ The type of information states is given in (23):

(23)    TotalInformationState (TIS) =  
          $\left[ \begin{array}{l} \text{dialoguegameboard} : \text{DGBType} \\ \text{private} : \text{Private} \end{array} \right]$

## Contexts in KoS

- ▶ The dialogue gameboard (DGB) represents information that arises from publicized interactions and, for now, we can identify it as the public context.

$$(24) \quad \text{DGBType (initial definition)} = \left[ \begin{array}{l} \text{spkr: Ind} \\ \text{addr: Ind} \\ \text{c-utt : addressing(spkr,addr)} \\ \text{Facts : Set(Prop)} \\ \text{Moves : list(IllocProp)} \\ \text{QUD : poset(Question)} \end{array} \right]$$

## Basics of Interaction

- ▶ The basic units of change are mappings between DGBs that specify how one DGB configuration can be modified into another.  $\therefore$  *conversational rule*.
- ▶ The types specifying its domain and its range respectively the *preconditions* and the *effects*.

## Basics of Interaction

- Notationwise a conversational rule will be specified as in (25a). We will often notate such a mapping as in (25b):

$$(25) \text{ a. } r : \begin{bmatrix} \dots \\ \text{dgb1} : \text{DGBType} \\ \dots \end{bmatrix} \mapsto \begin{bmatrix} \dots \\ \text{dgb2} : \text{DGBType} \\ \dots \end{bmatrix}$$

$$\text{b. } \begin{bmatrix} \text{pre(conds)} : \text{RType} \\ \text{effects} : \text{RType} \end{bmatrix}$$

## Greeting and Parting

- The conversational rule associated with greeting:

$$\left[ \begin{array}{l} \text{pre : } \left[ \begin{array}{l} \text{spkr : Ind} \\ \text{addr : Ind} \\ \text{moves = elist : list(IllocProp)} \\ \text{qud = elist : list(Question)} \\ \text{facts = commonground1 : Prop} \end{array} \right] \\ \text{effects : } \left[ \begin{array}{l} \text{spkr = pre.spkr : Ind} \\ \text{addr = pre.addr : Ind} \\ \text{LatestMove = Greet(spkr,addr):IllocProp} \\ \text{qud = pre.qud : list(Question)} \\ \text{facts = pre.facts : Prop} \end{array} \right] \end{array} \right]$$



## Greeting and Parting

- So we can *abbreviate* conversational rules as in (26a), which for the rule for greeting yields (26b).

$$\begin{aligned}
 (26) \quad & \text{a. } \left[ \begin{array}{l} \text{pre: PreCondSpec} \\ \text{effects : ChangePrecondSpec} \end{array} \right] \\
 & \text{b. } \left[ \begin{array}{l} \text{pre : } \left[ \begin{array}{l} \text{moves = elist : list(IllocProp)} \\ \text{qud = elist : list(Question)} \end{array} \right] \\ \text{effects : } \left[ \text{LatestMove = Greet(spr,addr):IllocProp} \right] \end{array} \right]
 \end{aligned}$$

## Greeting and Parting

$$\left[ \begin{array}{lcl} & & \left[ \begin{array}{l} \text{qud} = \text{eset} : \text{poset}(\text{Question}) \\ \text{facts} : \text{Set}(\text{Prop}) \\ \text{f} = \text{MinInteraction}(\{\text{spkr}, \text{addr}\}) : \text{Prop} \\ \text{c1} : \text{member}(\text{f}, \text{facts}) \end{array} \right] \\ \text{preconds} & : & \\ \text{effects} & : & \left[ \text{LatestMove} = \text{Part}(\text{spkr}, \text{addr}) : \text{IllocProp} \right] \end{array} \right]$$

## Greeting and Parting

► Counterparting:



$$\left[ \begin{array}{l} \text{preconds} : \left[ \begin{array}{l} \text{LatestMove} = \text{Part}(\text{spkr}, \text{addr}) : \text{IllocProp} \\ \text{qud} = \text{eset} : \text{poset}(\text{Question}) \end{array} \right] \\ \text{effects} : \left[ \begin{array}{l} \text{spkr} = \text{preconds.addr} : \text{Ind} \\ \text{addr} = \text{preconds.spkr} : \text{Ind} \\ \text{LatestMove} = \text{CounterPart}(\text{spkr}, \text{addr}) : \text{IllocProp} \end{array} \right] \end{array} \right]$$

# Greeting and Parting

$$\left[ \begin{array}{lcl} \text{preconds} & : & \left[ \begin{array}{l} \text{LatestMove} = \text{CounterPart}(\text{spkr}, \text{addr}) : \text{IllocProp} \\ \text{qud} = \text{eset} : \text{poset}(\text{Question}) \end{array} \right] \\ \text{effects} & : & \left[ \text{LatestMove} = \text{Disengaged}(\{\text{spkr}, \text{addr}\}) : \text{IllocProp} \right] \end{array} \right]$$

## Asking, Asserting, Answering, and Accepting

- ▶ Broadly speaking queries and assertions are either *issue initiating*—they introduce an issue unrelated to those currently under discussion— or they are *reactive*—they involve a reaction to a previously raised issue.
- ▶ Accounting for the the reactive ones using DGB-based conversational rules is simple. These can also be used to explicate the effects *issue initiating* moves have. More on that a bit later.

## Simple assertion and querying: ingredients

querying	assertion
LatestMove = Ask(A,q)	LatestMove = Assert(A,p)
A: push q onto QUD; release turn;	A: push p? onto QUD; release turn
B: push q onto QUD; take turn; make q—specific utterance take turn.	B: push p? onto QUD; take turn; Option 1: Discuss p?  Option 2: Accept p
	LatestMove = Accept(B,p)
	B: increment FACTS with p; pop p? from QUD;
	A: increment FACTS with p; pop p? from QUD;

## Simple assertion and querying: ingredients

- ▶ q-specific utterance: an utterance whose content is either a proposition  $p$  About max-qud (*partial answer*) or a question  $q_1$  on which max-qud Depends (*sub-question*).

## Asking, Asserting, Answering, and Accepting

- ▶ Two aspects of this protocol are not query specific:
  1. The protocol is like the ones we have seen for greeting and parting a 2-person turn exchange protocol (2-PTEP).
  2. The specification make  $q$ -specific utterance is an instance of a general constraint that characterizes the contextual background of reactive queries and assertions.



## Asking, Asserting, Answering, and Accepting

- QSPEC: if  $q$  is QUD-maximal, then subsequent to this either conversational participant may make a move constrained to be  $q$ -specific (i.e. either About or Influencing  $q$ ).

(27) QSPEC

$$\left[ \begin{array}{l} \text{pre} : \left[ \text{qud} = \langle q, Q \rangle : \text{poset}(\text{Question}) \right] \\ \text{effects} : \text{TurnUnderspec} \wedge_{\text{merge}} \\ \left[ \begin{array}{l} r : \text{AbSemObj} \\ R : \text{IllocRel} \\ \text{LatestMove} = R(\text{spkr}, \text{addr}, r) : \text{IllocProp} \\ c1 : \text{Qspecific}(r, q) \end{array} \right] \end{array} \right]$$

## Turn Change and abbreviation

- Turnholder-Underspecified =

$$\left[ \begin{array}{l} \text{pre : } \left[ \begin{array}{l} \text{spkr : Ind} \\ \text{addr : Ind} \end{array} \right] \\ \\ \text{effects : } \left[ \begin{array}{l} \text{PrevAud} = \{ \text{pre.spkr}, \text{pre.addr} \} : \text{Set(Ind)} \\ \text{spkr : Ind} \\ \text{c1 : member(spkr, PrevAud)} \\ \text{addr : Ind} \\ \text{c2 : member(addr, PrevAud) } \wedge \text{ addr } \neq \text{spkr} \end{array} \right] \end{array} \right]$$

# Asking, Asserting, Answering, and Accepting

- The only query specific aspect of the querying protocol is:

(28) Ask QUD–incrementation:

$$\left[ \begin{array}{l} \text{pre : } \left[ \begin{array}{l} q : \text{Question} \\ \text{LatestMove} = \text{Ask}(\text{spkr}, \text{addr}, q) : \text{IllocProp} \end{array} \right] \\ \text{effects : } \left[ \text{qud} = \langle q, \text{pre.qud} \rangle : \text{poset}(\text{Question}) \right] \end{array} \right]$$

## Asking, Asserting, Answering, and Accepting

- ▶ What are the components of the assertion protocol? Not specific to assertion is the fact that it is a 2-PTEP; similarly, the discussion option is simply an instance of QSPEC.
- ▶ This leaves two novel components: QUD incrementation with  $p?$  and acceptance.

# Asking, Asserting, Answering, and Accepting

(29) Assert QUD-incrementation:

$$\left[ \begin{array}{ll} \text{preconds} & : \left[ \begin{array}{l} p : \text{Prop} \\ \text{LatestMove} = \text{Assert}(\text{spkr}, \text{addr}, p) : \text{IllocProp} \end{array} \right] \\ \text{effects} & : \left[ \text{qud} = \langle p?, \text{pre.qud} \rangle : \text{poset}(\text{Question}) \right] \end{array} \right]$$

## Asking, Asserting, Answering, and Accepting

- ▶ The second component of acceptance is the incrementation of FACTS by  $p$ .
- ▶ This is not quite as straightforward as it might seem: when FACTS gets incremented, we also need to ensure that  $p?$  gets downdated from QUD—only Nonresolved questions can be in QUD (see notes for discussion, leads to account of rhetorical questions.).
- ▶ In order to ensure that this is the case, we need to check for all existing elements of QUD that they are not resolved by the new value of FACTS.
- ▶ Hence, accepting  $p$  involves both an update of FACTS and a downdate of QUD—minimally just removing  $p?$ , potentially removing other questions as well.

## Fact Update/ QUD Downdate

- NonResolve is a function that maps a partially ordered set of questions  $poset(q)$  and a set of propositions  $P$  to a partially ordered set of questions  $poset'(q)$  which is identical to  $poset(q)$  modulo those questions in  $poset(q)$  resolved by members of  $P$ .

(30) Fact Update/ QUD Downdate

$$\left[ \begin{array}{l} \text{pre : } \left[ \begin{array}{l} p : \text{Prop} \\ \text{LatestMove} = \text{Accept}(\text{spkr}, \text{addr}, p) \vee \\ \text{Confirm}(\text{spkr}, \text{addr}, p) \\ \text{qud} = \langle p?, Q \rangle : \text{poset}(\text{Question}) \end{array} \right] \\ \text{effects : } \left[ \begin{array}{l} \text{facts} = \text{pre.facts} \cup \{p\} : \text{Set}(\text{Prop}) \\ \text{qud} = \text{NonResolve}(Q, \text{facts}) : \text{poset}(\text{Question}) \end{array} \right] \end{array} \right]$$

## What drives the dialogue?

- ▶ Baseline rule: Free Speech; when QUD is empty one can say whatever one likes.
- ▶ This is true when one stands on the proverbial soap box or (various caveats) when sitting in a restaurant with friends ...
- ▶ 
$$\left[ \begin{array}{l} \text{pre} : [\text{qud} = \langle \rangle : \text{eset}(\text{question})] \\ \text{effects} : \text{TurnUnderspec} \wedge_{\text{merge}} \\ \left[ \begin{array}{l} r : \text{AbSemObj} \\ R : \text{IllocRel} \\ \text{LatestMove} = R(\text{spkr}, \text{addr}, r) \text{ IllocProp} \end{array} \right] \end{array} \right]$$



## A simple example

(31) a. A: Hi

B: Hi

A: Who's coming tomorrow?

B: Several colleagues of mine (are coming).

A: I see.

B: Mike (is coming) too.

## A simple example

Utt.	DGB Update (Conditions)	Rule
initial	$\text{MOVES} = \langle \rangle$ $\text{QUD} = \langle \rangle$ $\text{FACTS} = \text{cg1}$	
1	$\text{LatestMove} := \text{Greet}(A, B)$	greeting
2	$\text{LatestMove} := \text{CounterGreet}(B, A)$	countergreeting
3	$\text{LatestMove} := \text{Ask}(A, B, q_0)$ $\text{QUD} := \langle q_0 \rangle$	Free Speech Ask QUD-incrementation
4	$\text{LatestMove} := \text{Assert}(B, A, p_1)$ $(\text{About}(p_1, q_0))$ $\text{QUD} := \langle p_1?, q_0 \rangle$	QSPEC Assert QUD-incrementation
5	$\text{LatestMove} := \text{Accept}(A, B, p_1)$ $\text{QUD} := \langle q_0 \rangle$ $\text{FACTS} := \text{cg1} \wedge p_1$	Accept Fact update/QUD downdate
6	$\text{LatestMove} := \text{Assert}(B, A, p_2)$ $(\text{About}(p_2, q_0))$ $\text{QUD} := \langle p_2?, q_0 \rangle$	QSPEC Assert QUD-incrementation

## Assertoric Benchmark 2

### ► Assertion benchmark 2: Accommodate disagreement

- (32)     A(1): Who will agree to come?  
          B(2): Helen and Jelle.  
          A(3): I doubt Helen will want to come after last time.  
          B(4): I think she's forgiven and forgotten.  
          A(5): OK.

## Assertoric Benchmark 2

Utt.	DGB Update (Conditions)	Rule
initial	MOVES = $\langle \rangle$ QUD = $\langle \rangle$ FACTS = cg1	
1	LatestMove := Ask(A,B,q0) QUD := $\langle q0 \rangle$	Free Speech Ask QUD-incrementation
2	LatestMove := Assert(B,A,p1) (About(p1,q0)) QUD := $\langle p1?, q0 \rangle$	QSPEC Assert QUD-incrementation
3	LatestMove := Assert(A,B,p2) (About(p2,p1?)) QUD := $\langle p2?, p1?, q0 \rangle$	QSPEC Assert QUD-incrementation
4	LatestMove := Assert(B,A,p3) (About(p3,p2?)) QUD := $\langle p3?, p2?, p1?, q0? \rangle$	QSPEC Assert QUD-incrementation
5	LatestMove := Accept(A,B,p3) QUD := $\langle p1?, q0 \rangle$	Accept Fact update/QUD downdate

## Query benchmark 4: query responses

► Query benchmark4: accommodate sub-questions.

- (33)     A(1): Who should we invite for tomorrow?  
         B(2): Who will agree to come?  
         A(3): Helen and Jelle and Fran and maybe Sunil.  
         B(4) : (a) I see. (b) So, Jelle I think.  
         A(5): OK.

## Query benchmark 4: query responses

Utt.	DGB Update (Conditions)	Rule
initial	MOVES = $\langle \rangle$ QUD = $\langle \rangle$ FACTS = cg1	
1	LatestMove := Ask(A,B,q0) QUD := $\langle q0 \rangle$	Free Speech Ask QUD-incrementation
2	LatestMove := Ask(B,A,q1) Influence(q1,q0) QUD := $\langle q1, q0 \rangle$	QSPEC Assert QUD-incrementation
3	LatestMove := Assert(A,B,p1) (About(p1,q1)) QUD := $\langle p1?, q1, q0 \rangle$	QSPEC Assert QUD-incrementation
4a	LatestMove := Accept(B,A,p1) FACTS := $cg1 \cup \{p1\}$ QUD := $\langle q0 \rangle$	Accept Fact update/QUD downdate
4b	LatestMove := Assert(B,A,p2) (About(p2,q0)) QUD := $\langle p2?, q0 \rangle$	QSPEC Assert QUD-incrementation
5	LatestMove := Accept(A,B,p2) FACTS := $cg1 \cup \{p1, p2\}$ QUD := $\langle q0 \rangle$	Accept Fact update/QUD downdate

## Query Benchmark 6 (partially)

- ▶ Same basic mechanism seems to regulate queries/assertions, across varying sizes of participant sets:
  - (34) a. Monologue: self answering (*A: Who should we invite? Perhaps Noam.*)
  - b. Dialogue: querier/responder (*A: Who should we invite? B: Perhaps Noam.*)
  - c. Multilogue: discussed later, maybe.
- ▶ Query benchmark 6: ensure approach scales down to monologue

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