Lagrange Duality and KKT Conditions

ECE 4200 02/21/20 Recitation

Mathematical Optimization

All learning is some optimization problem
 Stick to canonical form

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minimize f_0(x)
subject to f_i(x) \leq b_i, \quad i = 1, \dots, m
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- $x = (x_1, x_2, ..., x_n) opt. variables ; x*$
- $f_0: R^n \rightarrow R objective function$
- $f_i : R^n \rightarrow R constraint function$

Optimization Example

- Well familiar with: regression
 - Least squares

min.
$$||y - X\beta||_2^2$$

Add some constraints

min.
$$||y - X\beta||_2^2$$

s.t. $||\beta||_2^2 \le \lambda$

Why convex optimization?

- Can't solve most OPs
 - E.g. NP Hard, even high polynomial time too slow
- Convex OPs
 - (Generally) No analytic solution
 - Efficient algorithms to find (global) solution

What is Convex Optimization?

• OP with *convex* objective and constraint functions

minimize
$$f_0(x)$$

subject to $f_i(x) \leq b_i, \quad i = 1, \dots, m$

• f₀, ..., f_m are *convex* = convex OP that has an efficient solution!

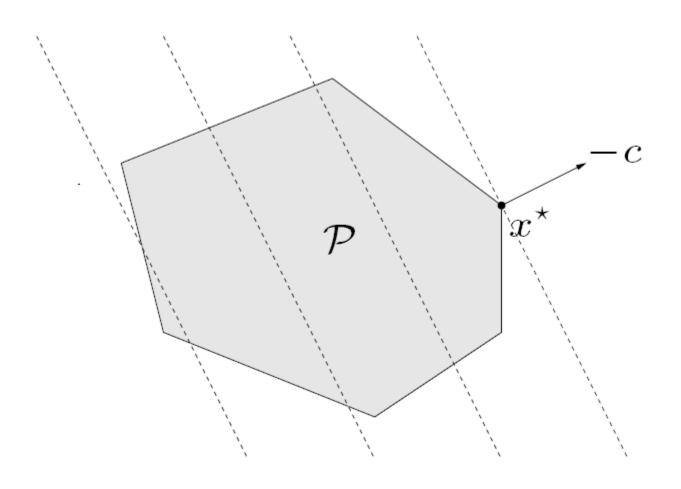
Some common convex OPs

• LP: affine objective function, affine constraints

minimize
$$c^T x + d$$

subject to $Gx \leq h$
 $Ax = b$

LP Visualization



Quadratic Program

• QP: Quadratic objective, affine constraints

minimize
$$(1/2)x^TPx + q^Tx + r$$
 subject to $Gx \leq h$ $Ax = b$

- LP is special case
- Many SVM problems result in QP, regression

SVM Optimization

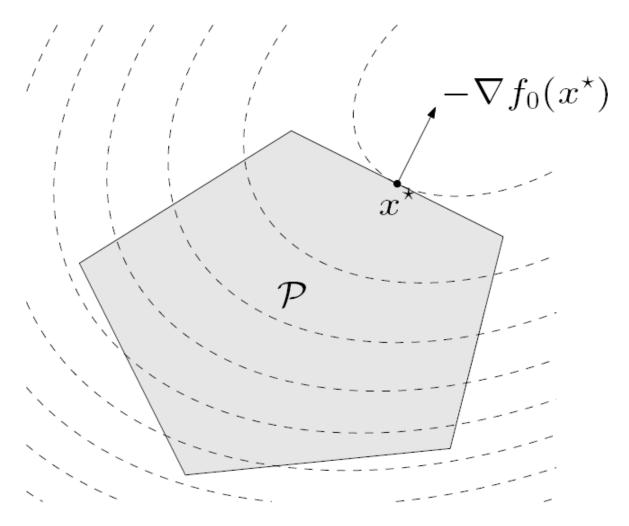
Minimize over \overline{w} , t:

$$\| \overline{w} \|_2^2$$

Subject to:

$$y_i(\overline{w}\cdot \overline{X}_i - t) \ge 1$$

QP Visualization



Lagrange

Standard form:

minimize
$$f_0(x)$$

subject to $f_i(x) \leq 0, \quad i = 1, \dots, m$
 $h_i(x) = 0, \quad i = 1, \dots, p$

Lagrange L:

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^{m} \lambda_i f_i(x) + \sum_{i=1}^{p} \nu_i h_i(x)$$

Lagrange Dual Function

$$g(\lambda, \nu) = \inf_{x \in \mathcal{D}} L(x, \lambda, \nu)$$

$$= \inf_{x \in \mathcal{D}} \left(f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \right)$$

 Lagrange Dual found by minimizing L with respect to primal variables

Lagrange Dual Function

Lagrange Dual provides lower bound on objective value at solution

lower bound property: if $\lambda \succeq 0$, then $g(\lambda, \nu) \leq p^*$

proof: if \tilde{x} is feasible and $\lambda \succeq 0$, then

$$f_0(\tilde{x}) \ge L(\tilde{x}, \lambda, \nu) \ge \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) = g(\lambda, \nu)$$

minimizing over all feasible \tilde{x} gives $p^* \geq g(\lambda, \nu)$

Lagrange Dual Problem

- Why not make the lower bound best possible?
- Dual problem: maximize $g(\lambda, \nu)$ subject to $\lambda \succ 0$
- Always convex opt. problem (even when primal is non-convex)
- Weak Duality: d* <= p* (have already seen this)

Lagrange Dual Function

- Note: Lagrange dual function is the point-wise infimum of family of affine functions of (lambda, nu)
- Thus, g is concave even if problem is not convex

Example: SVM

Minimize over \overrightarrow{w} , t:

$$\| \overrightarrow{w} \|_2^2$$

Subject to:

$$\forall i, 1 - y_i(\overrightarrow{w} \cdot \overrightarrow{x_i} - t) \leq 0$$

Lagrange L:

$$\mathcal{L}(\overrightarrow{w}, t, \boldsymbol{\alpha}) = \frac{1}{2} \cdot ||\overrightarrow{w}||_2^2 + \sum_{1 \leq i \leq m} \alpha_j \left(1 - y_i(\overrightarrow{w} \cdot \overrightarrow{x_i} - t)\right)$$

Lagrange Dual function:

$$g(\boldsymbol{\alpha}) = \operatorname{Inf}_{\overrightarrow{w},t} \mathcal{L}(\overrightarrow{w},t,\boldsymbol{\alpha})$$

Dual Problem:

maximize $g(\alpha)$

subject to $\alpha \geq 0$

Optimization problem

$$\mathcal{L}(\overrightarrow{w}, t, \boldsymbol{\alpha}) = \frac{1}{2} \cdot \| \overrightarrow{w} \|_2^2 + \sum_{1 \le i \le n} \alpha_i (1 - y_i (\overrightarrow{w} \cdot \overrightarrow{x_i} - t))$$

Primal: $\min_{\overrightarrow{w},t} \max_{\alpha \geq 0} \mathcal{L}(\overrightarrow{w},t,\alpha) \text{ (p*)}$

Dual: $\max_{\alpha \geq 0} \min_{\overrightarrow{w},t} \mathcal{L}(\overrightarrow{w},t,\alpha) \quad (d^*)$

Strong Duality

- If d* = p*, strong duality holds
- Does not hold in general
- Slater's Theorem: If convex problem, and strictly feasible point exists, then strong duality holds! (proof too involved, omit here)
- => For convex problems, can use dual problem to find solution

When strong duality holds

$$f_{0}(x^{\star}) = g(\lambda^{\star}, \nu^{\star})$$

$$= \inf_{x} \left(f_{0}(x) + \sum_{i=1}^{m} \lambda_{i}^{\star} f_{i}(x) + \sum_{i=1}^{p} \nu_{i}^{\star} h_{i}(x) \right) \quad \text{(definition)}$$

$$\leq f_{0}(x^{\star}) + \sum_{i=1}^{m} \lambda_{i}^{\star} f_{i}(x^{\star}) + \sum_{i=1}^{p} \nu_{i}^{\star} h_{i}(x^{\star})$$

$$\leq f_{0}(x^{\star}). \quad \text{(since constraints satisfied at } x^{\star})$$

 Sandwiched between f₀(x), last 2 inequalities are equalities!

$$f_0(x^*) = f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^p \nu_i^* h_i(x^*)$$

- Which means: $\sum_{i=1}^{n} \lambda_i^{\star} f_i(x^{\star}) = 0$
- Since each term is non-positive, we have complementary slackness:

$$\lambda_i^{\star} f_i(x^{\star}) = 0, \quad i = 1, \dots, m.$$

 Whenever constraint is non-active, corresponding multiplier is zero

$$\lambda_i^* f_i(x^*) = 0, \quad i = 1, \dots, m.$$

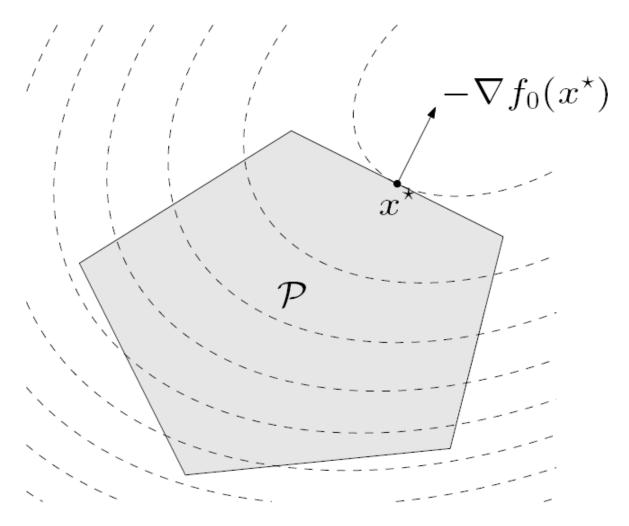
This can also be described by

$$\lambda_i^* > 0 \implies f_i(x^*) = 0,$$

 $f_i(x^*) < 0 \implies \lambda_i^* = 0.$

 Since usually only a few active constraints at solution (see geometry), the dual variable lambda is often sparse

QP Visualization

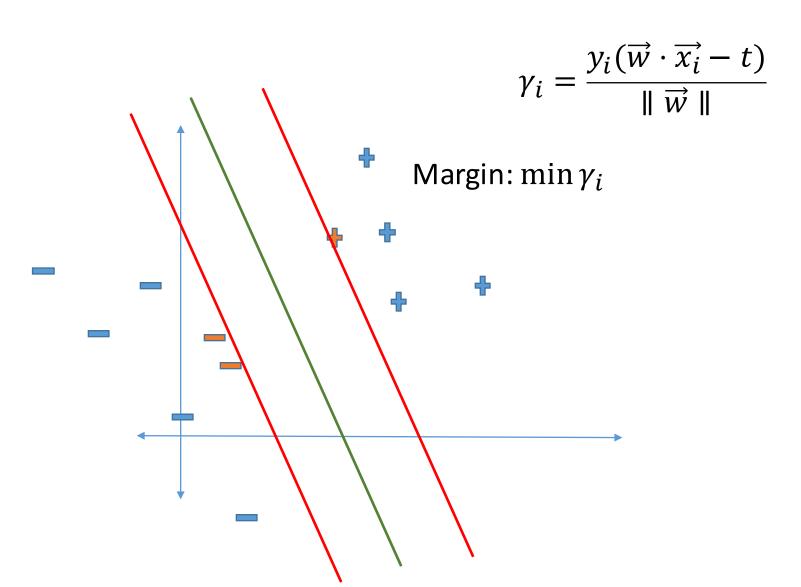


$$\lambda_i^{\star} f_i(x^{\star}) = 0, \quad i = 1, \dots, m.$$

- As we will see, this is why support vector machines result in solution with only key support vectors
 - These come from the dual problem, constraints correspond to points, and complementary slackness ensures only the "active" points are kept

Assume data linearly separable

Distance of the point $\overrightarrow{x_i}$:



KKT Conditions

- The KKT conditions are then just what we call that set of conditions required at the solution (basically list what we know)
- KKT conditions play important role
 - Can sometimes be used to find solution analytically
 - Otherwise can think of many methods as ways of solving KKT conditions

KKT Conditions

• Again given strong duality and assuming differentiable, since x^* minimizes $L(x, \lambda^*, \nu^*)$ gradient must be 0 at x^*

$$\nabla f_0(x^*) + \sum_{i=1}^m \lambda_i^* \nabla f_i(x^*) + \sum_{i=1}^p \nu_i^* \nabla h_i(x^*) = 0.$$

 Thus, putting it all together, for non-convex problems we have

KKT Conditions — non-convex

$$\begin{aligned}
f_{i}(x^{\star}) &\leq 0, & i = 1, \dots, m \\
h_{i}(x^{\star}) &= 0, & i = 1, \dots, p \\
\lambda_{i}^{\star} &\geq 0, & i = 1, \dots, m \\
\lambda_{i}^{\star} f_{i}(x^{\star}) &= 0, & i = 1, \dots, m \\
\nabla f_{0}(x^{\star}) + \sum_{i=1}^{m} \lambda_{i}^{\star} \nabla f_{i}(x^{\star}) + \sum_{i=1}^{p} \nu_{i}^{\star} \nabla h_{i}(x^{\star}) &= 0,
\end{aligned}$$

Necessary conditions

KKT Conditions – convex

$$f_{i}(\tilde{x}) \leq 0, \quad i = 1, \dots, m$$

$$h_{i}(\tilde{x}) = 0, \quad i = 1, \dots, p$$

$$\tilde{\lambda}_{i} \geq 0, \quad i = 1, \dots, m$$

$$\tilde{\lambda}_{i} f_{i}(\tilde{x}) = 0, \quad i = 1, \dots, m$$

$$\nabla f_{0}(\tilde{x}) + \sum_{i=1}^{m} \tilde{\lambda}_{i} \nabla f_{i}(\tilde{x}) + \sum_{i=1}^{p} \tilde{\nu}_{i} \nabla h_{i}(\tilde{x}) = 0,$$

then \tilde{x} and $(\tilde{\lambda}, \tilde{\nu})$ are primal and dual optimal, with zero duality gap.

Also sufficient conditions:

$$\begin{split} g(\tilde{\lambda},\tilde{\nu}) &= L(\tilde{x},\tilde{\lambda},\tilde{\nu}) & \text{(according to the last equation)} \\ &= f_0(\tilde{x}) + \sum_{i=1}^m \tilde{\lambda}_i f_i(\tilde{x}) + \sum_{i=1}^p \tilde{\nu}_i h_i(\tilde{x}) \\ &= f_0(\tilde{x}), \end{split}$$

Optimization problem

$$\mathcal{L}(\overrightarrow{w}, t, \boldsymbol{\alpha}) = \frac{1}{2} \cdot \|\overrightarrow{w}\|_{2}^{2} + \sum_{1 \leq i \leq m} \alpha_{i} \left(1 - y_{i}(\overrightarrow{w} \cdot \overrightarrow{x_{i}} - t)\right)$$

Take the gradient equal to zero:

$$\nabla_{\overrightarrow{w}} \mathcal{L}(\overrightarrow{w}, t, \boldsymbol{\alpha}) = 0 \Rightarrow \overrightarrow{w} = \sum_{1 \leq i \leq m} \alpha_i y_i \overrightarrow{x_i}.$$

Linear combinations

Optimization problem

Take the gradient equal to zero:

$$\frac{\partial \mathcal{L}(\vec{w}, t, \boldsymbol{\alpha})}{\partial t} = 0$$

gives,

$$\sum_{1 \le i \le m} \alpha_i y_i = 0.$$

Dual Problem

We substitute these values of \vec{w} , t.

Maximize:

$$\sum_{1 \leq i \leq n} \alpha_i - \sum_{1 \leq i, j \leq n} \alpha_i \alpha_j y_i y_j \left(\overrightarrow{x_i} \cdot \overrightarrow{x_j} \right)$$

Subject to

$$\sum \alpha_i y_i = 0, \alpha_i \geq 0.$$

Reference

Slides are credited to:

http://www.ittc.ku.edu/bioinfo_seminar/ppt/Summer09/ConvexOpt_July09.ppt

More resources:

1. Stanford Convex Optimization (Boyd & Vandenberghe):

http://ee364a.stanford.edu/lectures/duality.pdf

2. Slater's condition (Wiki):

https://en.wikipedia.org/wiki/Slater%27s_condition