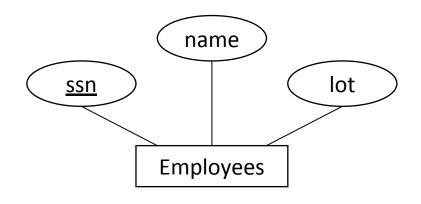
# Introduction to Database System ER to Relational



## Logical DB Design: ER to Relational

Entity sets to tables:



CREATE TABLE Employees
(ssn CHAR(11),
name CHAR(20),
lot INTEGER,
PRIMARY KEY (ssn))

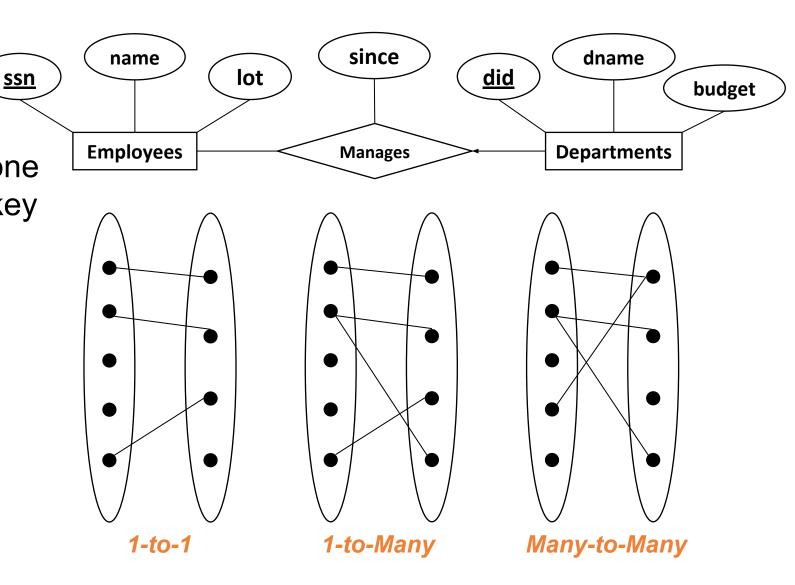
## Relationship Sets to Tables

- In translating a relationship set to a relation, attributes of the relation must include:
  - Keys for each participating entity set (as foreign keys).
    - This set of attributes forms a superkey for the relation.
  - All descriptive attributes.

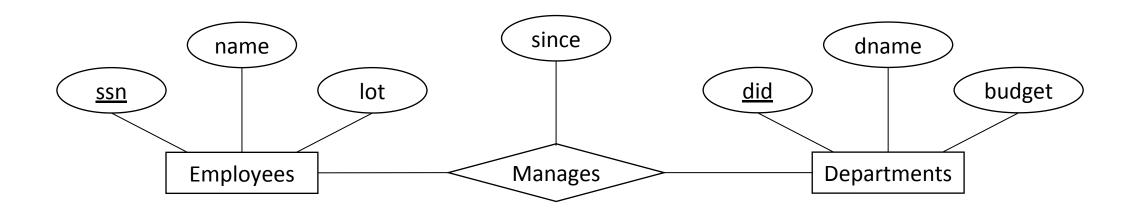
```
CREATE TABLE Works_In(
    ssn CHAR(11),
    did INTEGER,
    since DATE,
    PRIMARY KEY (ssn, did),
    FOREIGN KEY (ssn)
    REFERENCES Emplyees,
    FOREIGN KEY (did)
    REFERENCES Departments)
```

## **Review Key Constraints**

Each dept has at most one manager, according to the key <u>key constraint</u> on Manages.



## Translation to relational model (SQL)?



## Translating ER Diagrams with Key Constraints

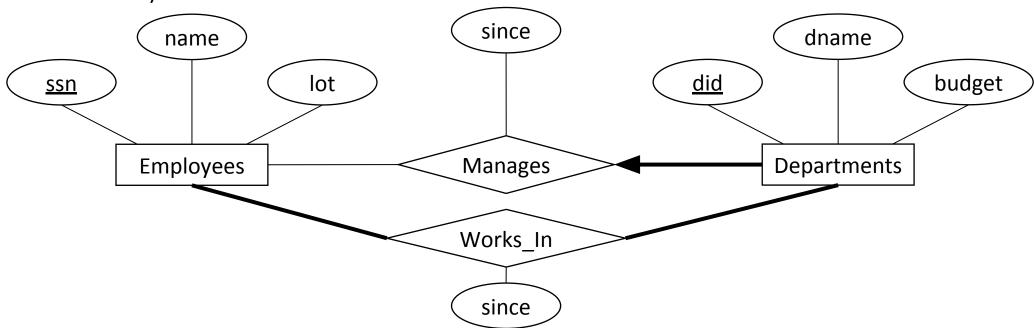
- Map relationship to a table:
  - Note that did is the key now!
  - Separate tables for Employees and Departments.
- Since each department has a unique manager, we could instead combine Manages and Departments.

```
CREATE TABLE Manages(
    ssn CHAR(11),
    did INTEGER,
    since DATE,
    PRIMARY KEY (did),
    FOREIGN KEY (ssn) REFERENCES Employees,
    FOREIGN KEY (did) REFERENCES Departments)
```

```
CREATE TABLE Dept_Mgr(
    did INTEGER,
    dname CHAR(20)
    budget REAL,
    ssn CHAR(11),
    since DATE,
    PRIMARY KEY (did),
    FOREIGN KEY (ssn) REFERENCES Employees)
```

## **Review: Participation Constraints**

- Does every department have a manager?
  - If so, this is a <u>participation constraint:</u> the participation of Departments in Manages is said to be total (vs. partial).
    - Every did value in Departments table must appear in a row of the Manages table (with a non-null ssn value)



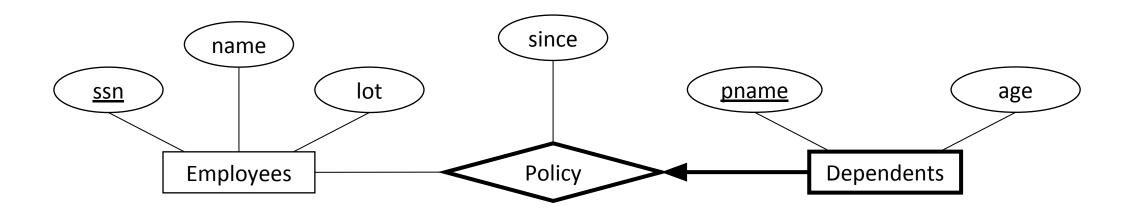
## Participation Constraints in SQL

• We can capture participation constraints involving one entity set in a binary relationship, but little else (without resorting to CHECK constraints).

```
CREATE TABLE Dept_Mgr(
    did INTEGER,
    dname CHAR(20),
    budget REAL,
    ssn CHAR(11) NOT NULL,
    since DATE,
    PRIMARY KEY (did),
    FOREIGN KEY (ssn) REFERENCES Employees,
    ON DELETE NO ACTION)
```

### **Review: Weak Entities**

- A weak entity can be identified uniquely only by considering the primary key of another (owner) entity.
  - Owner entity set and weak entity set must participate in a one-to-many relationship set (1 owner, many weak entities).
  - Weak entities set must have total participation in this identifying relationship set.

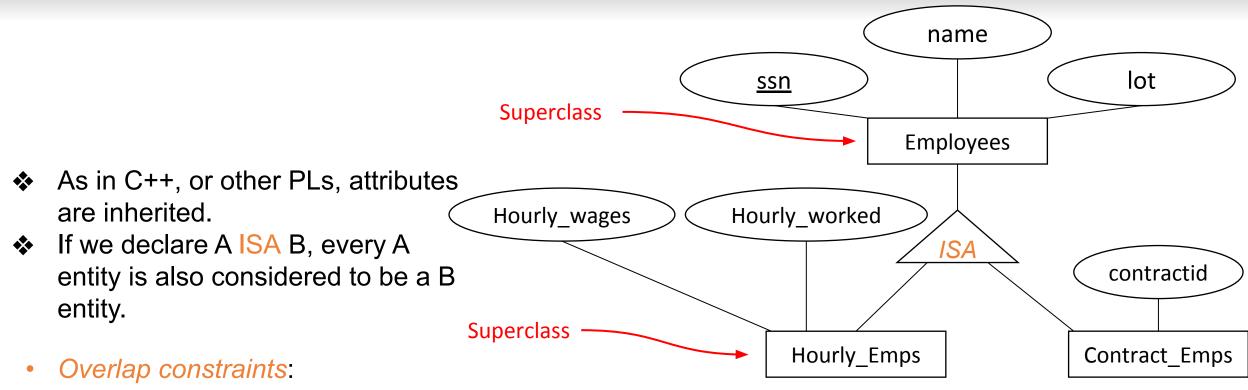


## **Translating Weak Entity Sets**

- Weak entity set and identifying relationship set are translated into a single table.
  - When the owner entity is deleted, all owned weak entities must also be deleted.

```
CREATE TABLE Dept_Mgr(
    pname CHAR(20),
    age INTEGER,
    cost REAL,
    ssn CHAR(11) NOT NULL,
    PRIMARY KEY (did),
    FOREIGN KEY (ssn) REFERENCES Employees,
    ON DELETE NO ACTION)
```

#### **Class Hierarchies**



- Can Joe be an Hourly\_Emps as well as a Contract\_Emps entity? (Allowed/disallowed)
- Covering constraints: Does every Employees entity also have to be an Hourly\_Emps or a Contract\_Emps entity? (Yes/no)
- Reasons for using ISA:
  - To add descriptive attributes specific to a subclass.
  - To identify that participate in a relationship.

## Translating ISA Hierarchies to Relations

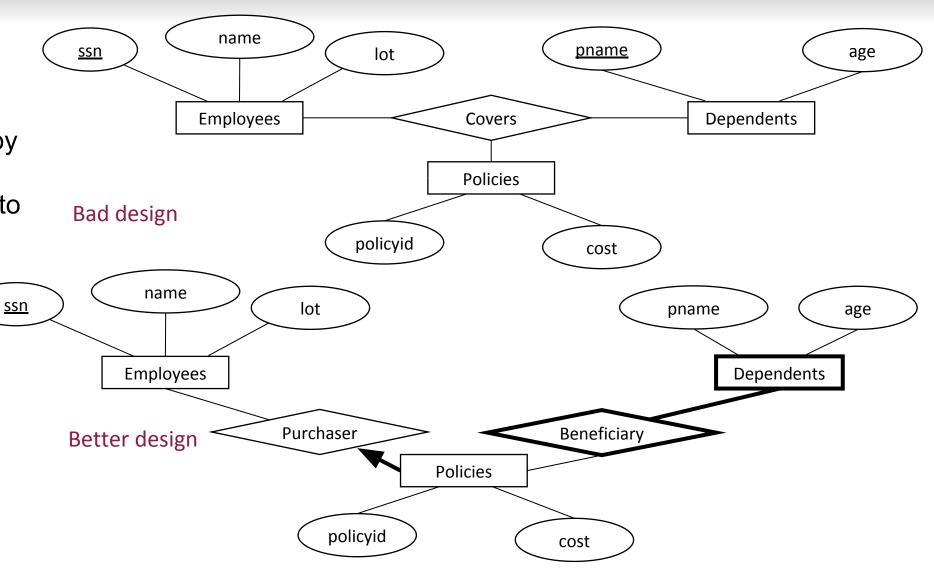
#### General approach:

- 3 relations: Employees, Hourly\_Emps and Contract\_Emps.
  - Hourly\_Emps: Every employee is recorded in Employees. For hourly emps,
    extra info recorded in Hourly\_Emps(hourly\_wages, hours\_worked, ssn); must
    delete Hourly\_Emps tuple if referenced Employees tuple is deleted).
  - Queries involving all employees easy, those involving just Hourly\_Emps require a join to get some attributes.
- Alternative: Just Hourly\_Emps and Contract\_Emps.
  - Hourly\_Emps: <u>ssn</u>, name, lot, hourly\_wages, hours\_worked.
  - Each employee must be in one of these two subclasses.

## Binary vs. Ternary Relationships

 If each policy is owned by just 1 employee, and each department is tied to the covering policy, first diagram is inaccurate.

 What are the additional constraints in the 2<sup>nd</sup> diagram?



## Binary vs. Ternary Relationships (Contd.)

- The key constraints allow us to combine Purchaser with Policies and Beneficiary with Dependents.
- Participation constraints lead to NOT NULL constraints.

```
CREATE TABLE Policies(
         policyid INTEGER,
         cost REAL,
         ssn CHAR(11) NOT NULL,
         PRIMARY KEY (policyid),
         FOREIGN KEY (ssn) REFERENCES Employees,
             ON DELETE CASCADE)
CREATE TABLE Dependents(
         pname CHAR(20),
         age INTEGER,
         policyid INTEGER,
         PRIMARY KEY (pname, policyid),
         FOREIGN KEY (policyid) REFERENCES Employees,
             ON DELETE CASCADE)
```