This is CS50x

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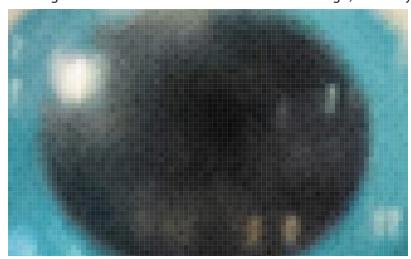
Lecture 4

- Pixels
- Hexadecimal
- Addresses, pointers
- Strings
- Pointer arithmetic
- Compare and copy

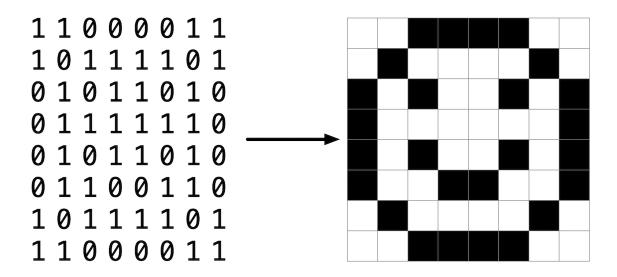
- Memory allocation
- valgrind
- Garbage values
- Swap
- Memory layout
- scanf
- Files
- JPEG

Pixels

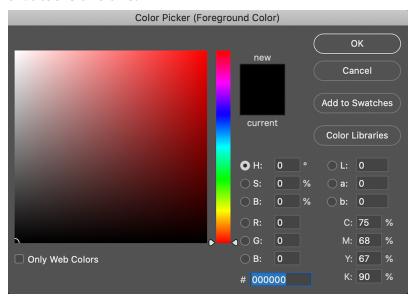
- Last week, we took a look at memory and how we could use arrays to store data.
- We might zoom in further and further in an image, but very soon we'll see individual pixels, like the ones in this puppet's eye:



- Since this image is stored with a finite number of bytes, each perhaps representing a red, green, or blue value for each pixel, there is in turn a finite number of pixels we can see.
- A simpler image of a smiley face might be represented with a single bit per pixel:



- We can visit cs50.ly/art (https://cs50.ly/art) with a Google Account to make a copy of a spreadsheet, which we can then fill with colors to create our own pixel art (https://en.wikipedia.org/wiki/Pixel art).
- Adobe Photoshop (https://en.wikipedia.org/wiki/Adobe_Photoshop), a popular image editing software, includes a color picker that looks like this:



■ Here, the color black is selected, and we see that the values for R, G, and B, or red, green, and blue respectively, are all

- 0. Furthermore, we see another value, #000000 that seems to represent all three values for the color black.
- We take a look at a few more colors:
 - white, with R: 255, G: 255, and B: 255, and #FFFFFF
 - red, with R: 255, G: 0, and B: 0, and #FF0000
 - green, with R: 0, G: 255, and B: 0, and #00FF00
 - blue, with R: 0, G: 0, and B: 255, and #0000FF

Hexadecimal

■ We might notice a pattern for the new notation, where it appears that each value for red, green, and blue are represented with two character. It turns out that there's another base system, **hexadecimal**, or base-16, where there are 16 digits:

```
0 1 2 3 4 5 6 7 8 9 A B C D E F
```

- The digits start with 0 through 9, and continue with A through F as equivalents to the decimal values of 10 through 15.
- Let's consider a two-digit hexadecimal number:

```
16^1 16^0
0 0
```

- Here, the $\boxed{0}$ in the ones place (since $16^0=1$) has a decimal value of 0. We'll count up, and after $\boxed{09}$ we'll use $\boxed{0A}$ to represent 10 in decimal.
- We can keep counting until OF, which is equivalent to 15 in decimal.
- After OF, we need to carry the one, as we would go from 9 to 10 in decimal:

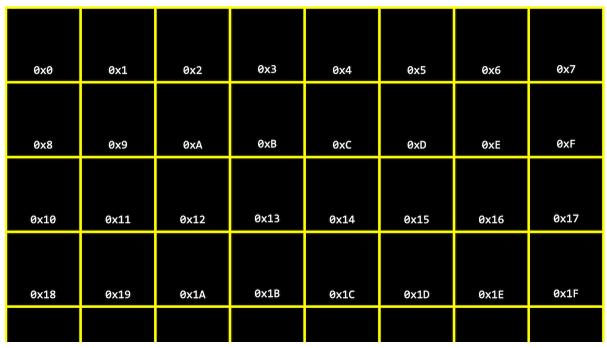
```
16^1 16^0
1 0
```

- lacksquare Here, the lacksquare has a value of $16^1 imes 1 = 16$, so lacksquare in hexadecimal is 16 in decimal.
- lacksquare With two digits, we can have a maximum value of FF , or $16^1 imes 15 + 16^0 imes 15 = 16 imes 15 + 1 imes 15 = 240 + 15 = 255$.

- The values in a computer's memory are still stored as binary, but this way of representation helps us humans represent larger numeric values with fewer digits needed.

Addresses, pointers

• For our computer's memory, too, we'll see hexadecimal used to describe each address or location:

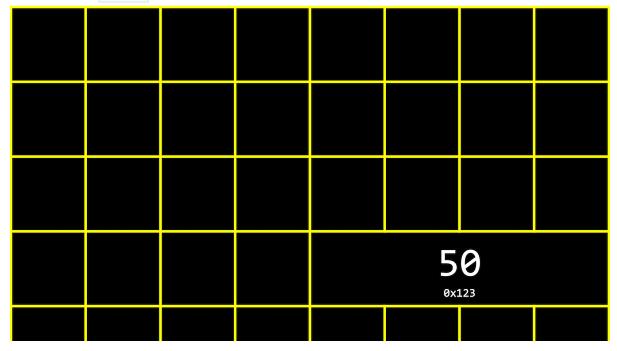


- By writing 0x in front of a hexadecimal value, we can distinguish them from decimal values.
- We might create a value n, and print it out:

```
#include <stdio.h>
int main(void)
```

```
int n = 50;
printf("%i\n", n);
}
```

■ In our computer's memory, there are now 4 bytes somewhere that have the binary value of 50, with some value for its address, like 0x123:



- A **pointer** is a variable that stores an address in memory, where some other variable might be stored.
- The & operator can be used to get the address of some variable, as with &n. And the * operator declares a variable as a pointer, as with int *p, indicating that we have a variable called p that points to an int. So, to store the address of a variable n into a pointer p, we would write:

```
int *p = &n;
```

■ In C, we can actually see the address with the & operator, which means "get the address of this variable":

```
#include <stdio.h>
```

```
int main(void)
{
   int n = 50;
   int *p = &n;
   printf("%p\n", p);
}
```

```
$ make address
$ ./address
0x7ffcb4578e5c
```

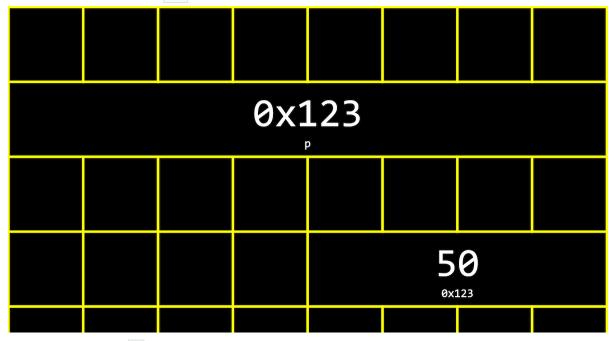
- %p is the format code to print an address with printf. And we only need to use the name of the variable, p, after we've declared it.
- In our instance of VS Code, we see an address with a large value like 0x7ffcb4578e5c. The value of the address in itself is not significant, since it's just some location in memory that the variable is stored in; instead, the important idea is that we can *use* this address later.
- We can run this program a few times, and see that the address of n in memory changes, since different addresses in memory will be available at different times.
- With C, we can also go to specific addresses in memory, which might cause **segmentation faults**, where we've tried to read or write to memory we don't have permission to.
- The * operator is also the **dereference operator**, which goes to an address to get the value stored there. For example, we can say:

```
#include <stdio.h>

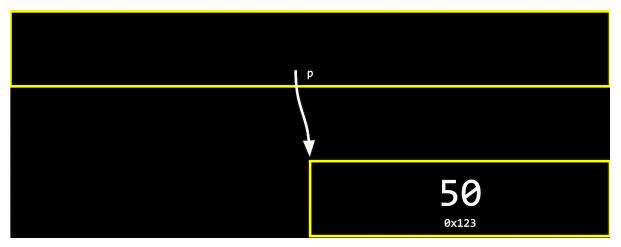
int main(void)
{
    int n = 50;
    int *p = &n;
    printf("%p\n", p);
    printf("%i\n", *p);
```

```
$ ./address
0x7ffda0a4767c
50
```

- Now, we see the value of the pointer itself (an address), and then the value at the address with *p , which is 50.
- Since we declared p to be an int *, the compiler knows that *p is an int, so the right number of bytes are read.
- In memory, we might have one variable, p, with the value of some address, like 0x123, stored, and another variable, an integer with the value 50, at that addresss:



- Notice that p takes up 8 bytes, since in modern computer systems, 64 bits are used in order to address the billions of bytes of memory available. With 32 bits, we can only count up to about 4 billion bytes, or about 4GB of memory.
- We can abstract away the actual value of the addresses, since they'll be different as we declare variables in our programs. We can simply think of p as pointing at some value in memory:



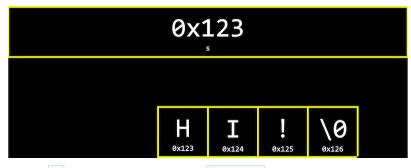
■ In the real world, we might have a mailbox labeled "p", among many mailboxes with addresses. Inside our mailbox, we can put a value like 0x123, which is the address of some other mailbox that's labeled "n".

Strings

• We can declare a string with string s = "HI!"; which will be stored one character at a time in memory. And we can access each character with s[0], s[1], s[2], and s[3]:



■ But it turns out that each character, since it's stored in memory, *also* has some unique address, and s is actually just a pointer with the address of the first character:



• s is a variable of type string, which is a pointer to a character.

- Recall that we can read the entire string by starting at the address in s, and continue reading one character at a time from memory until we reach \0.
- It turns out that string s = "HI!" is the same as char *s = "HI!"; And we can use strings in C in the exact same way without the CS50 Library, by using char *.
- Let's print out a string:

```
#include <cs50.h>
#include <stdio.h>

int main(void)
{
    string s = "HI!";
    printf("%s\n", s);
}

$ make address
$ ./address
HI!
```

■ Now, we can remove the CS50 Library, and say:

```
#include <stdio.h>
int main(void)
{
    char *s = "HI!";
    printf("%s\n", s);
}

$ make address
$ ./address
HI!
```

• We can experiment and see the address of characters:

```
#include <cs50.h>
#include <stdio.h>

int main(void)
{
    string s = "HI!";
    char c = s[0];
    char *p = &c;
    printf("%p\n", s);
    printf("%p\n", p);
}
```

```
$ make address
$ ./address
0x402004
0x7ffd4227fdd7
```

- We store the first character of s into c, and print out its address with p. We also print out s as an address with %p, and we see that the values are difference since we made a copy of the first character with char c = s[0];
- Now, we'll print the address of the first character in s:

```
#include <cs50.h>
#include <stdio.h>

int main(void)
{
    string s = "HI!";
    char *p = &s[0];
    printf("%p\n", p);
    printf("%p\n", s);
}
```

```
$ make address
$ ./address
0x402004
0x402004
```

- With char *p = &s[0]; , we store the address of the first character in s into a pointer called p. And now, when we print p and s as addresses, we see the same value.
- We can see the address of each character in s:

```
#include <stdio.h>

int main(void)
{
    char *s = "HI!";
    printf("%p\n", s);
    printf("%p\n", &s[0]);
    printf("%p\n", &s[1]);
    printf("%p\n", &s[2]);
    printf("%p\n", &s[3]);
}
```

```
$ make address
$ ./address
0x402004
0x402004
0x402005
0x402006
0x402007
```

- Again, the address of the first character, &s[0], is the same as the value of s. And each following character has an address that is one byte higher.
- In the CS50 Library, a string is defined with just typedef char *string; . With typedef, we're creating a custom data

type for the word string, making it equivalent to char *.

Pointer arithmetic

• We can print out each character in a string:

```
#include <cs50.h>
#include <stdio.h>

int main(void)
{
    string s = "HI!";
    printf("%c\n", s[0]);
    printf("%c\n", s[1]);
    printf("%c\n", s[2]);
    printf("%c\n", s[3]);
}
```

```
$ make address
$ ./address
H
I
```

- When we declare a string with double quotes, ", the compiler figures out where to put those characters in memory as an array.
- Let's simplify our code to use char * and show just the printable characters:

```
#include <stdio.h>
```

```
int main(void)
{
    char *s = "HI!";
    printf("%c\n", s[0]);
    printf("%c\n", s[1]);
    printf("%c\n", s[2]);
}

$ make address
$ ./address
H
I
```

■ But we can go to addresses directly:

```
#include <stdio.h>

int main(void)
{
    char *s = "HI!";
    printf("%c\n", *s);
    printf("%c\n", *(s + 1));
    printf("%c\n", *(s + 2));
}
```

- *s goes to the address stored in s, and *(s + 1) goes to the location in memory with the next character, an address that is one byte higher.
- s[1] is **syntactic sugar**, like an abstraction for *(s + 1), equivalent in function but more human-friendly to read and write.
- **Pointer arithmetic** is the process of applying mathematical operations to pointers, using them just like numbers (which they are).

• We can declare an array of numbers, and access them with pointer arithmetic:

```
#include <stdio.h>

int main(void)
{
    int numbers[] = {4, 6, 8, 2, 7, 5, 0};

    printf("%i\n", *numbers);
    printf("%i\n", *(numbers + 1));
    printf("%i\n", *(numbers + 2));
    printf("%i\n", *(numbers + 3));
    printf("%i\n", *(numbers + 4));
    printf("%i\n", *(numbers + 5));
    printf("%i\n", *(numbers + 6));
}
```

```
$ make address
$ ./address
4
6
8
2
7
5
0
```

- It turns out that we only need to add 1 to the address of numbers, instead of 4 (even though int s are 4 bytes in size), since the compiler already knows that the type of each value in numbers is 4 bytes. With + 1, we're telling the compiler to move the next value in the array, not the next byte.
- And notice that numbers is an array, but we can use it as a pointer with *numbers.

Compare and copy

• Let's try to compare two integers from the user:

```
#include <cs50.h>
#include <stdio.h>

int main(void)
{
    int i = get_int("i: ");
    int j = get_int("j: ");

    if (i == j)
        {
            printf("Same\n");
        }
        else
        {
                printf("Different\n");
        }
}
```

```
$ make compare
$ ./compare
i: 50
j: 50
Same
$ ./compare
i: 50
j: 42
Different
```

- We compile and run our program, and it works as we'd expect, with the same values of the two integers giving us "Same" and different values "Different".
- When we try to compare two strings, we see that the same inputs are causing our program to print "Different":

```
$ make compare
$ ./compare
s: HI!
t: BYE!
Different
$ ./compare
s: HI!
t: HI!
```

• Even when our inputs are the same, we see "Different" printed.

- Each "string" is a pointer, char *, to a different location in memory, where the first character of each string is stored.

 So even if the characters in the string are the same, this will always print "Different".
- And get_string, this whole time, has been returning just a char *, or a pointer to the first character of a string from the user. Since we called get_string twice, we got two different pointers back.
- We can fix our program with:

```
#include <cs50.h>
#include <stdio.h>
#include <string.h>

int main(void)
{
    char *s = get_string("s: ");
    char *t = get_string("t: ");

    if (strcmp(s, t) == 0)
    {
        printf("Same\n");
    }
    else
    {
        printf("Different\n");
    }
}
```

```
$ make compare
$ ./compare
s: HI!
t: HI!
Same
```

• We'll take a look at the values of s and t as pointers:

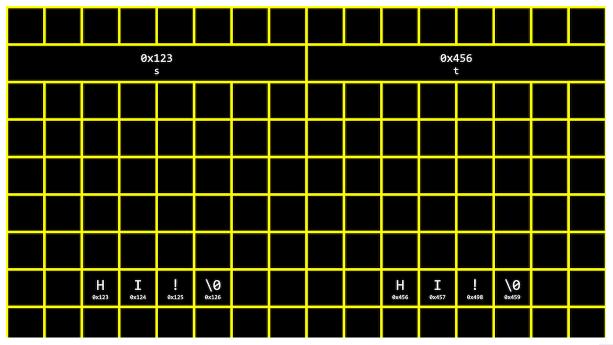
```
#include <cs50.h>
#include <stdio.h>

int main(void)
{
    char *s = get_string("s: ");
    char *t = get_string("t: ");

    printf("%p\n", s);
    printf("%p\n", t);
}
```

```
$ make compare
$ ./compare
s: HI!
t: HI!
0x19e06b0
0x19e06f0
```

- We see that the addresses of our two strings are indeed different.
- Let's visualize how this might look in our computer's memory. Our first string might be at address 0x123, our second might be at 0x456, and s will have the value of 0x123, pointing at that location, and t will have the value of 0x456, pointing at another location:



- Since our computer puts each string at some location in memory for us, we need s and t to point to each of them. And now we see why comparing s and t directly will always print "Different". strcmp, in contrast, will go to each string and compare them character by character.
- In C, we can also get the address of s or t, and store them in a variable of the type char **, a pointer to a pointer.
- Let's try to copy a string:

```
#include <cs50.h>
#include <ctype.h>
#include <stdio.h>
#include <string.h>

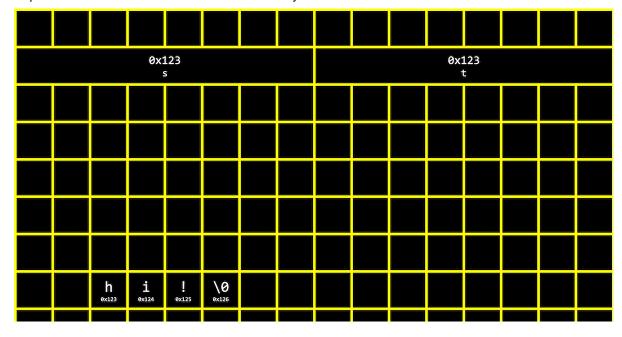
int main(void)
{
    string s = get_string("s: ");
    string t = s;
```

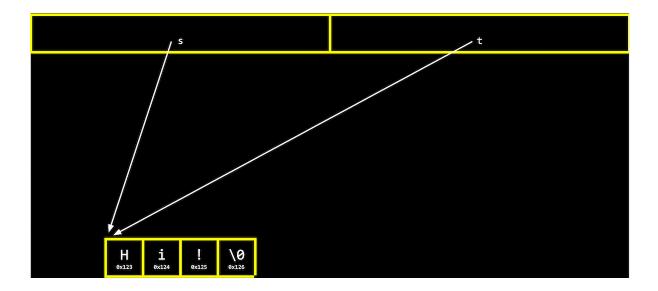
```
t[0] = toupper(t[0]);

printf("s: %s\n", s);
printf("t: %s\n", t);
}
```

```
$ make copy
$ ./copy
s: hi!
s: Hi!
t: Hi!
```

- We get a string s, and copy the value of s into t. Then, we capitalize the first letter in t.
- But when we run our program, we see that both s and t are now capitalized.
- Since we set s and t to the same value, or the same address, they're both pointing to the same character, and so we capitalized the same character in memory:





Memory allocation

- To actually make a copy of a string, we have to do a little more work, and copy each character in s to somewhere else in memory.
- We'll need to use a new function, malloc, to *allocate* some number of bytes in memory. And we'll use free to mark memory as usable when we're done with it, so the operating system can do something else with it.
 - Our computers might slow down if a program we're running has a bug where it allocates more and more memory but never frees it. The operating system will take longer and longer to find enough available memory for our program.
- Let's copy a string now:

```
#include <cs50.h>
#include <ctype.h>
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

int main(void)
{
```

```
char *s = get_string("s: ");

char *t = malloc(strlen(s) + 1);

for (int i = 0, n = strlen(s) + 1; i < n; i++)
{
     t[i] = s[i];
}

t[0] = toupper(t[0]);

printf("s: %s\n", s);
printf("t: %s\n", t);
}</pre>
```

- We create a new variable to point to a new string with char *t . The argument to malloc is the number of bytes we'd like to use. We already know the length of s, but we need to add 1 for the terminating null character.
- Then, we copy each character, one at a time, with a for loop. We use strlen(s) + 1 since we want to copy the null character too to end the string. In the loop, we set t[i] = s[i], copying the characters.
- We could also use a library function, strcpy:

```
#include <ctype.h>
#include <stdio.h>
#include <stdib.h>
#include <string.h>

int main(void)
{
    char *s = get_string("s: ");
    char *t = malloc(strlen(s) + 1);
```

```
strcpy(t, s);

t[0] = toupper(t[0]);

printf("s: %s\n", s);
printf("t: %s\n", t);

free(t);
}
$ make copy
$ /copy
```

```
$ make copy
$ ./copy
s: hi!
s: hi!
t: Hi!
```

- Now, we can capitalize the first letter of t.
- We'll remember to call free on t, since we allocated it ourselves.
- We can add some error-checking to our program:

```
#include <ctype.h>
#include <stdio.h>
#include <stdib.h>
#include <stdlib.h>
#include <string.h>

int main(void)
{
    char *s = get_string("s: ");
    char *t = malloc(strlen(s) + 1);
    if (t == NULL)
    {
}
```

```
return 1;
}

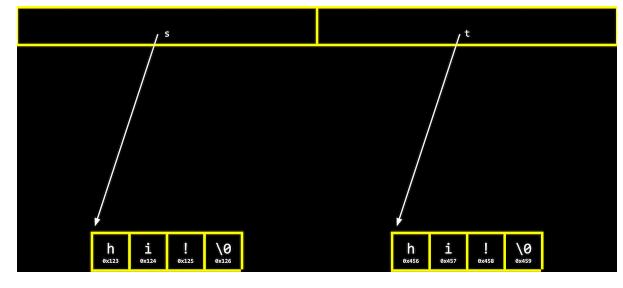
strcpy(t, s);

if (strlen(t) > 0)
{
    t[0] = toupper(t[0]);
}

printf("s: %s\n", s);
printf("t: %s\n", t);

free(t);
}
```

- If our computer is out of memory, malloc will return NULL, the null pointer, or a special value of all 0 bits that indicates there isn't an address to point to. So we should check for that case, and exit if t is NULL.
- We should also check that t has a length, before trying to capitalize the first character.
- We can visualize how this looks in our computer's memory:



We've allocated memory at 0x456 and set t to point to it. Then, we used strcpy to copy each character's value, starting from the address s is pointing to, to the address t is pointing to.

valgrind

• Let's allocate memory for some integers:

```
#include <stdio.h>
#include <stdib.h>

int main(void)
{
    int *x = malloc(3 * sizeof(int));
    x[1] = 72;
    x[2] = 73;
    x[3] = 33;
}

$ make memory
$ ./memory
$
```

- We'll use malloc to get enough memory for 3 times the size of an int, which we can find out with size of.
- We've deliberately made a mistake where we forgot arrays are 0-indexed, and started at x[1] instead. Then, with x[3], we're trying to access memory beyond the bounds of what we have access to.
- We also don't free the memory we've allocated.
- When we compile and run our program, though, nothing seems to happen. It turns out that our mistake wasn't bad enough to cause a segmentation fault this time, though it might next time.
- valgrind is a command-line tool that we can use to run our program and see if it has any memory-related issues.
- We'll run valgrind ./memory after compiling, and we'll see a lot of output:

```
$ valgrind ./memory
==5902== Memcheck, a memory error detector
==5902== Copyright (C) 2002-2017, and GNU GPL'd, by Julian Seward et al.
==5902== Using Valgrind-3.15.0 and LibVEX; rerun with -h for copyright info
==5902== Command: ./memory
==5902==
==5902== Invalid write of size 4
==5902==
                                               at 0x401162: main (memory.c:9)
==5902== Address 0x4bd604c is 0 bytes after a block of size 12 alloc'd
==5902==
                                               at 0x483B7F3: malloc (in /usr/lib/x86_64-linux-qnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/v
                                           by 0x401141: main (memory.c:6)
==5902==
==5902==
==5902==
==5902== HEAP SUMMARY:
==5902==
                                              in use at exit: 12 bytes in 1 blocks
                                           total heap usage: 1 allocs, 0 frees, 12 bytes allocated
==5902==
==5902==
==5902== 12 bytes in 1 blocks are definitely lost in loss record 1 of 1
                                              at 0x483B7F3: malloc (in /usr/lib/x86_64-linux-qnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/v
==5902==
==5902==
                                               by 0x401141: main (memory.c:6)
==5902==
==5902== LEAK SUMMARY:
==5902==
                                               definitely lost: 12 bytes in 1 blocks
                                               indirectly lost: 0 bytes in 0 blocks
==5902==
==5902==
                                                       possibly lost: 0 bytes in 0 blocks
                                              still reachable: 0 bytes in 0 blocks
==5902==
==5902==
                                                                    suppressed: 0 bytes in 0 blocks
==5902==
==5902== For lists of detected and suppressed errors, rerun with: -s
==5902== ERROR SUMMARY: 2 errors from 2 contexts (suppressed: 0 from 0)
```

■ We'll see some snippets like Invalid write of size 4 at ... memory.c:9, which gives us a hint to look at line 9, where we're using x[3].

We'll fix that mistake:

```
#include <stdio.h>
#include <stdib.h>

int main(void)
{
    int *x = malloc(3 * sizeof(int));
    x[0] = 72;
    x[1] = 73;
    x[2] = 33;
}
```

```
$ make memory
$ ./memory
$ valgrind ./memory
==6435== Memcheck, a memory error detector
==6435== Copyright (C) 2002-2017, and GNU GPL'd, by Julian Seward et al.
==6435== Using Valgrind-3.15.0 and LibVEX; rerun with -h for copyright info
==6435== Command: ./memory
==6435==
==6435==
==6435== HEAP SUMMARY:
==6435== in use at exit: 12 bytes in 1 blocks
==6435== total heap usage: 1 allocs, 0 frees, 12 bytes allocated
==6435==
==6435== 12 bytes in 1 blocks are definitely lost in loss record 1 of 1
                                       at 0x483B7F3: malloc (in /usr/lib/x86_64-linux-qnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/vgpreload_memcheck-amd64-linux-gnu/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/valgrind/v
==6435==
                                        by 0x401141: main (memory.c:6)
==6435==
==6435==
==6435== LEAK SUMMARY:
==6435==
                                    definitely lost: 12 bytes in 1 blocks
==6435==
                                   indirectly lost: 0 bytes in 0 blocks
                                      possibly lost: 0 bytes in 0 blocks
==6435==
```

```
==6435== still reachable: 0 bytes in 0 blocks
==6435== suppressed: 0 bytes in 0 blocks
==6435==
==6435== For lists of detected and suppressed errors, rerun with: -s
==6435== ERROR SUMMARY: 1 errors from 1 contexts (suppressed: 0 from 0)
```

- Now, we see less output, with only one error telling us 12 bytes ... are definitely lost, in that we've allocated them, but not freed them.
- Once we free our memory, we'll see no errors:

```
#include <stdio.h>
#include <stdlib.h>

int main(void)
{
    int *x = malloc(3 * sizeof(int));
    x[0] = 72;
    x[1] = 73;
    x[2] = 33;
    free(x);
}
```

```
$ make memory
$ ./memory
$ valgrind ./memory
==6812== Memcheck, a memory error detector
==6812== Copyright (C) 2002-2017, and GNU GPL'd, by Julian Seward et al.
==6812== Using Valgrind-3.15.0 and LibVEX; rerun with -h for copyright info
==6812== Command: ./memory
==6812==
==6812==
==6812== in use at exit: 0 bytes in 0 blocks
```

```
==6812== total heap usage: 1 allocs, 1 frees, 12 bytes allocated
==6812==
==6812== All heap blocks were freed -- no leaks are possible
==6812==
==6812== For lists of detected and suppressed errors, rerun with: -s
==6812== ERROR SUMMARY: 0 errors from 0 contexts (suppressed: 0 from 0)
```

Garbage values

■ Let's take a look at this program:

```
#include <stdio.h>
#include <stdib.h>

int main(void)
{
    int scores[3];
    for (int i = 0; i < 3; i++)
    {
        printf("%i\n", scores[i]);
    }
}</pre>
```

```
$ make garbage
$ ./garbage
68476128
32765
0
```

- We declare an array, scores, but we didn't initialize it with any values.
- The values in the array are **garbage values**, or whatever unknown values that were in memory, from whatever program was running in our computer before.

- If we aren't careful with how our programs access memory, users might end up seeing data from previous programs, like passwords. And if we try to go to an address that's a garbage value, our program is likely to crash from a segmentation fault.
- We watch Pointer Fun with Binky (https://www.youtube.com/watch?v=3uLKjb973HU), an animated video demonstrating pointers, malloc, and dereferencing. The code from the video might look like this in one program:

```
int main(void)
{
    int *x;
    int *y;

    x = malloc(sizeof(int));

    *x = 42;
    *y = 13;

    y = x;

    *y = 13;
}
```

- In the first two lines, we declare two pointers. Then, we allocate memory for x, but not y, so we can assign a value to the memory x is pointing to with *x = 42; But *y = 13; is problematic, since we haven't allocated any memory for y, and the garbage value there points to some area in memory we likely don't have access to.
- We can write y = x; so that y points to the same allocated memory as x, and use *y = 13; to set the value there.

Swap

- We'll have a volunteer on stage try to swap two liquids, one orange liquid in one glass, and one purple liquid in another. But we need a third glass to temporarily pour one liquid into, such as the orange liquid in the first glass. Then, we can pour the purple liquid into the first glass, and finally the orange liquid from the third glass into the second.
- Let's try to swap the values of two integers:

```
#include <stdio.h>
void swap(int a, int b);
int main(void)
    int x = 1;
    int y = 2;
    printf("x is %i, y is %i\n", x, y);
    swap(x, y);
    printf("x is %i, y is %i\n", x, y);
}
void swap(int a, int b)
    int tmp = a;
    a = b;
    b = tmp;
}
```

```
$ make swap
$ ./swap
x is 1, y is 2
x is 1, y is 2
```

- In our swap function, we have a third variable to use as temporary storage space as well. We put a into tmp, and then set a to the value of b, and finally b can be changed to the original value of a, now in tmp.
- But, if we tried to use that function in a program, we don't see any changes.
- It turns out that the swap function gets passed in copies of variables, a and b, which are **local variables** that only the surrounding function can access. Changing those values won't change x and y in the main function:

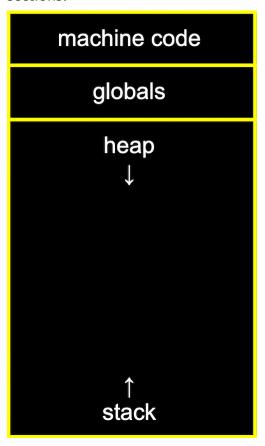
```
#include <stdio.h>
void swap(int a, int b);
int main(void)
    int x = 1;
    int y = 2;
    printf("x is %i, y is %i\n", x, y);
    swap(x, y);
    printf("x is %i, y is %i\n", x, y);
}
void swap(int a, int b)
{
    printf("a is %i, b is %i\n", a, b);
   int tmp = a;
    a = b;
    b = tmp;
    printf("a is %i, b is %i\n", a, b);
}
```

```
$ make swap
$ ./swap
x is 1, y is 2
a is 1, b is 2
a is 2, b is 1
x is 1, y is 2
```

• Our swap function works while we're inside it.

Memory layout

■ Within our computer's memory, different types of data that need to be stored for our program are organized into different sections:



- The **machine code** section is our compiled program's binary code. When we run our program, that code is loaded into memory.
- Just below, or in the next part of memory, are **global variables** we declared in our program.
- The **heap** section is an empty area from where malloc can get free memory for our program to use. As we call malloc, we start allocating memory from the top down.
- The **stack** section is used by functions and local variables in our program as they are called, and grows upwards.

- If we call malloc for too much memory, we will have a **heap overflow**, since we end up going past our heap. Or, if we call too many functions without returning from them, we will have a **stack overflow**, where our stack has too much memory allocated as well.
- Our program for swapping integers might have a stack that looks like this:



- Our main function has two local variables, x and y. Our swap function is created on top of main when it's called, and has three local variables, a, b, and tmp.
- Once swap returns, its memory is freed and its values are now garbage values, and the variables in main haven't been changed.
- By passing in the address of x and y, our swap function will be able to change the original values:

```
void swap(int *a, int *b)
{
   int tmp = *a;
   *a = *b;
   *b = tmp;
```

}

- The addresses of x and y are passed in from main to swap with &x and &y, and we use the int *a syntax to declare that our swap function takes in pointers.
- We save the first value to tmp by following the pointer a, and then set the second value to location pointed to by a by following the second pointer b.
- Finally, we store the value of tmp to the location pointed to by b.
- Our stack might look like this:



- If x is at 0x123, a will contain that value. b will have the address of y, 0x127.
- Our first step will be putting the value of x into tmp by following the pointer a:

| tmp 1 | | | | swap -----| a 0x123 | b 0x127 | main | x 1 | y 2 | ■ Then, we'll follow the pointer b, and store the value there into the location pointed to by a: | tmp 1 | | | | swap -----| a 0x123 | b 0x127 | main | x 2 | y 2 | ■ Finally, we'll go to the location pointed to by b, and put the value of tmp into it:

- Now, swap can return, and the variables in main will still be changed.
- In our program, we'll need to pass in the addresses of x and y to our swap function:

```
#include <stdio.h>

void swap(int *a, int *b);
int main(void)
{
    int x = 1;
    int y = 2;

    printf("x is %i, y is %i\n", x, y);
    swap(&x, &y);
    printf("x is %i, y is %i\n", x, y);
}

void swap(int *a, int *b)
{
    int tmp = *a;
```

```
*a = *b;

*b = tmp;

}

$ make swap

$ ./swap

x is 1, y is 2

x is 2, y is 1
```

• With &x, we can get the address of x to pass in.

scanf

x: 50

• We can get an integer from the user with a C library function, scanf:

```
#include <stdio.h>

int main(void)
{
    int x;
    printf("x: ");
    scanf("%i", &x);
    printf("x: %i\n", x);
}

$ make scanf
$ ./scanf
x: 50
```

scanf takes a format, %i, so the input is "scanned" for that format. We also pass in the address in memory where we want that input to go with &x.

• We can try to get a string the same way:

```
#include <stdio.h>
int main(void)
{
    char *s;
    printf("s: ");
    scanf("%s", s);
    printf("s: %s\n", s);
}

$ clang -o scanf scanf.c
$ ./scanf
$ s: HI!
$ s: (null)
```

- make prevents us from making this mistake, so we'll use clang to demonstrate.
- We haven't actually allocated any memory for s, so scanf is writing our string to an unknown address in memory.
- We can call malloc to allocate memory:

```
#include <stdio.h>

int main(void)
{
    char *s = malloc(4);
    printf("s: ");
    scanf("%s", s);
    printf("s: %s\n", s);
}
```

```
$ clang -o scanf scanf.c
$ ./scanf
```

```
s: HI!
s: HI!
```

• And we can declare an array of 4 characters:

```
#include <stdio.h>

int main(void)
{
    char s[4];
    printf("s: ");
    scanf("%s", s);
    printf("s: %s\n", s);
}
```

- Now, if the user types in a string of length 3 or less, our program will work safely. But if the user types in a longer string, scanf might be trying to write past the end of our array into unknown memory, causing our program to crash.
- get_string from the CS50 library continuously allocates more memory as scanf reads in more characters, so it doesn't have this issue.

Files

■ With the ability to use pointers, we can also open files, like a digital phone book in phonebook.c (https://cdn.cs50.net /2021/fall/lectures/4/src4/phonebook.c?highlight):

```
// Saves names and numbers to a CSV file
```

```
#include <cs50.h>
#include <stdio.h>
#include <string.h>
int main(void)
{
   // Open CSV file
   FILE *file = fopen("phonebook.csv", "a");
    if (!file)
        return 1;
    // Get name and number
    string name = get_string("Name: ");
    string number = get_string("Number: ");
    // Print to file
   fprintf(file, "%s,%s\n", name, number);
   // Close file
   fclose(file);
}
```

- fopen is a new function we can use to open a file with a new type, FILE.
- We can use fprintf to write to a file.
- We'll see more details about working with files in this week's problem set.

JPEG

■ Let's look at a program that opens a file and tells us if it's a JPEG file, a particular format for image files, with jpeg.c

(https://cdn.cs50.net/2021/fall/lectures/4/src4/jpeg.c?highlight):

```
// Detects if a file is a JPEG
#include <stdint.h>
#include <stdio.h>
typedef uint8_t BYTE;
int main(int argc, char *argv[])
{
   // Check usage
   if (argc != 2)
        return 1;
   }
   // Open file
   FILE *file = fopen(argv[1], "r");
    if (!file)
        return 1;
   }
   // Read first three bytes
    BYTE bytes[3];
    fread(bytes, sizeof(BYTE), 3, file);
   // Check first three bytes
    if (bytes[0] == 0xff && bytes[1] == 0xd8 && bytes[2] == 0xff)
        printf("Yes, possibly\n");
    else
```

```
{
    printf("No\n");
}

// Close file
fclose(file);
}
```

```
$ make jpeg
$ ./jpeg .src4/lecture.jpg
Yes, possibly
```

- First, we define a BYTE as 8 bits, so we can refer to a byte as a type more easily in C.
- Then, we'll read from a file with a function called fread.
- We can compare the first three bytes (in hexadecimal) to the three bytes required to begin a JPEG file. If they're the same, then our file is likely to be a JPEG file (though, other types of files may still begin with those bytes). But if they're not the same, we know it's definitely not a JPEG file.
- It turns out that BMP files, another format for images, have even more bytes in its header, or beginning of the file.
- We'll learn more about these in this week's problem set as well, and even implement our own version of image filters, like one that only shows the color red:

```
#include "helpers.h"

// Only let red through
void filter(int height, int width, RGBTRIPLE image[height][width])
{
    // Loop over all pixels
    for (int i = 0; i < height; i++)
    {
        for (int j = 0; j < width; j++)
        {
            image[i][j].rgbtBlue = 0x00;
        }
}</pre>
```

```
image[i][j].rgbtGreen = 0x00;
}
}
```

Here, we have a loop that iterates over all the pixels in a two-dimensional array, and sets the blue and green values to
 0.