

Chapter 1

Proofs of Parser Laws

The Haskell gigaparsec library uses a continuation-passing style (cps) to define its parsers, which makes it easier to reason about their semantics. The approach taken to prove the following parser law for parsley is via equational reasoning on gigaparsec semantics, under the assumption that their semantics are equivalent. While there is no formal proof of this equivalence at the present, gigaparsec was designed to have semantics equivalent to parsley's.

1.1 Left absorption for fmap

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f <$> empty
=   { applicative functor law }
pure f <*> empty
=   { definition of <*> }
liftA2 ($) (pure f) empty
=   { semantics of liftA2 }
Parsec $ \st ok err →
  let ok' x st' = (unParsec empty) st' (ok . (x $)) err
  in (unParsec $ pure f) st ok' err
=   { semantics of empty }
Parsec $ \st ok err →
  let ok' x st' = (unParsec $ raise (`emptyErr` 0)) st' (ok . (x $)) err
  in (unParsec $ pure f) st ok' err
=   { semantics of raise }
Parsec $ \st ok err →
  let ok' x st' = (unParsec $ Parsec $ \st'' _ bad →
    useHints bad (emptyErr st'' 0) st') st' (ok . (x $)) err
  in (unParsec $ pure f) st ok' err
=   {  $\beta$ -reduction }
Parsec $ \st ok err →
  let ok' x st' = useHints err (emptyErr st' 0) st'
  in (unParsec $ pure f) st ok' err
=   { semantics of pure }
Parsec $ \st ok err →
  let ok' x st' = useHints err (emptyErr st' 0) st'
  in (unParsec $ Parsec $ \st'' ok'' _ → ok'' f st'') st ok' err
=   {  $\beta$ -reduction }
Parsec $ \st ok err →
  let ok' x st' = useHints err (emptyErr st' 0) st'
  in ok' f st
=   { inline ok' }
Parsec $ \st ok err → useHints err (emptyErr st 0) st
=   { rearrange and  $\alpha$ -conversion }
Parsec $ \st _ bad → useHints bad ((`emptyErr` 0) st) st
=   { fold definition of raise }
raise (`emptyErr` 0)
=   { fold definition of empty }
empty

```