

parsley-garnish

A linter for the parsley parser combinator library

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Statement of Originality

Write about how this is your own work etc etc (check what's needed)

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Chapter 1

Introduction

1.1 Motivation

Parser combinators [Hutton 1992] are an elegant approach for writing parsers in a manner that remains close to the original grammar specification. `parsley` [Willis and Wu 2018] is a parser combinator library implemented as an embedded domain-specific language (DSL) [Hudak 1996] in Scala, with an API inspired by the `parsec` [Leijen and Meijer 2001] family of libraries in Haskell. However, as with many libraries, there exists a learning curve to utilising `parsley` and parser combinator libraries in an idiomatic manner.

While well-documented, the wealth of information to get started with `parsley` can be overwhelming for users, particularly those new to parser combinators. Furthermore, there exists a number of design patterns [Willis and Wu 2022] for writing maintainable parsers, which even experienced users may be unaware of. A potential solution to this problem is tooling to provide automated code hints, which a user can use during the development cycle to evaluate if their code adheres to best practices.

A number of modern integrated development environments (IDEs) provide code hints to warn programmers about problems in their source code, highlighting offending snippets and suggesting actions to improve suboptimal or incorrect code [Kurbatova et al. 2021]. Many of these code analysis tools are designed to detect general issues for the host language, rather than specifically for libraries. However, tools may also utilise domain-specific code analyses in order to detect issues specific to a particular system or problem domain [Renggli et al. 2010; Gregor and Schupp 2006].

1.2 Project Goals

This project aims to explore the potential of harnessing static code analysis techniques to develop a new tool, `parsley-garnish`, that offers code hints aimed at assisting programmers in writing idiomatic and correct `parsley` code. Additionally, for certain issues that can be automatically fixed, `parsley-garnish` will provide automated actions to resolve the issue. The goal of `parsley-garnish` is to be used as a companion library to `parsley`, in order to improve its ease of adoption and to help users enforce best practices.

Chapter 2

Background

2.1 Linting and Static Analysis

Linting is the process of analysing source code to identify and report issues related to coding style and potential logical errors. The term originates from the `Lint` program [Johnson 1978], which examined C source code for bugs, as well as wasteful code patterns that may be legal but error-prone. The tool was also utilised to enforce portability restrictions which aided users in writing portable code that could be compiled on multiple platforms. Since the release of `Lint`, many linting tools, known as **linters**, have been developed for a wide range of programming languages.

Nowadays, many linters can be integrated into IDEs, where code analysis performed by the linter is run incrementally in the background. Any violations found by the linter are displayed directly in the editor as warnings or errors at the relevant locations in the source code. This brings early, real-time feedback to the programmer, allowing them to address issues as they write code, with minimal interruption to their development workflow. Linters can also be integrated as part of the code review process, or into continuous integration (CI) pipelines to ensure that code adheres to a set of standards before being merged into the main codebase.

Although the traditional definition for linting is concerned only with *detecting* issues in code, modern linters have broadened their scope significantly. In addition to detecting issues, many linters provide *auto-fix* capabilities to correct issues by automatically rewriting the offending code snippets. This feature is often integrated into IDEs as well: the popular Language Server Protocol for defining IDE features enables these auto-fix features via *code actions* [Gunasinghe and Marcus 2022]. When a section of code is highlighted by a linter warning, a user can apply a code action to automatically fix the issue with a single click.

Linters and related static analysis tools are increasingly becoming more important in modern software development, as modern code continues to become more complex and difficult to reason about. Industry leaders, such as Google [Sadowski et al. 2018] and Meta/Facebook [Calcagno et al. 2015], have embraced these tools as integral components of their software development workflows. The use of automated tools to detect potential issues is not only faster but in some cases more effective than human review, saving developer time and reducing error rates in the development process.

2.1.1 Types of Issues Detected by Linters

Many linters are configurable with a set of rules, which specify the categories of issues that the linter should detect. These rules can be enabled or disabled by users, allowing them to customise the linter to their needs. Rules are usually grouped by purpose: some rules are concerned with simply improving code style, while others are concerned with detecting suspicious code patterns indicative of potential bugs.

Style checks and code quality

Linters can suggest opportunities to improve code by utilising language features in a more idiomatic manner. Snippets of code that violate these stylistic rules are not necessarily incorrect, but should be fixed as they may be harder to read or maintain in the long term. Furthermore, many idiomatic practices exist to avoid common pitfalls that could lead to unintended behaviour. By highlighting good practices, linters can help users avoid these common mistakes that may cause bugs. For example, *ESLint*¹, one of the most popular JavaScript linters, warns against common JavaScript pitfalls such as using the regular equality operator `=` instead of its type-safe alternative `===`.

A well-designed linter can help programmers learn about useful language constructs by suggesting them in the context of their code, aiding them in adhering to best practices and common style conventions. This category of rules is therefore especially helpful as an educational tool for new users of a language, who may be unaware of these idioms. For example, the *Clippy*² linter for Rust [Li et al. 2023] categorises a collection of rules as `clippy::complexity` rules to detect code that does something simple in a complex way and suggests a simpler alternative. Fig. 2.1 provides an example of a similar rule in Haskell, from the *HLint*³ linter. The rule suggests to rewrite the function into an equivalent but more concise form via η -reduction, presented to the user as a code action that can be applied automatically.

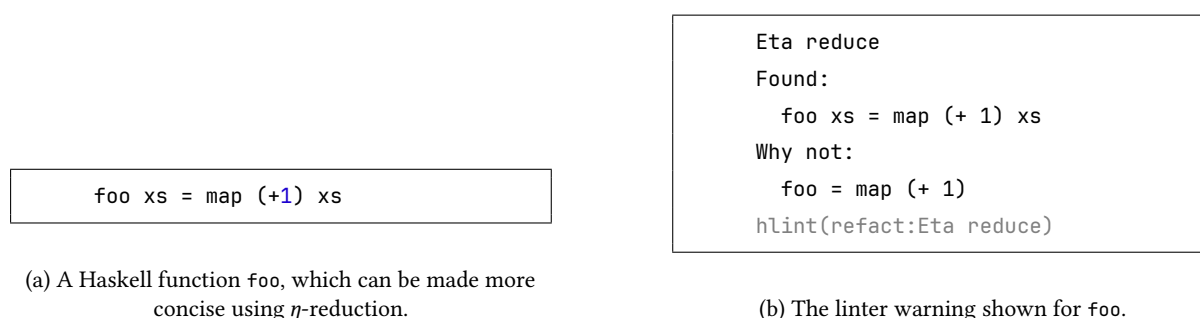


Fig. 2.1: An example of a warning from the Haskell linter *hlint*, suggesting a fix that a user can choose to automatically apply.

Domain-specific idioms A library or especially an embedded DSL may require a particular style of usage that is different from the host language [Hora et al. 2012]. The majority of linters are designed for general-purpose application domains, so they are unlikely to detect issues specific to a more specialised domain. Therefore, linters may be developed for specific libraries or DSLs, with their own set of domain-specific rules. In this case, the accompanying linter can benefit users and improve developer productivity in a similar manner to general-purpose linters: common misuses can be detected and sometimes automatically fixed, and users can be directed to relevant documentation to learn more about correct usage. For instance, the *xUnit.net* testing framework for C# is accompanied by the `xunit.analyzers`⁴ package which provides linting rules to enforce best practices specific to *xUnit*.

¹<https://eslint.org/docs/latest/rules/>

²<https://doc.rust-lang.org/clippy/>

³<https://hackage.haskell.org/package/hlint>

⁴<https://github.com/xunit/xunit.analyzers>

Code smells and opportunities for refactoring

Code refactoring is a well-established practice in software development. In his influential book *Refactoring: Improving the Design of Existing Code* [Fowler 2018], Fowler defines **refactoring** as “the process of changing a software system in such a way that it does not alter the external behaviour of the code yet improves its internal structure”. Refactoring may be employed to eliminate **code smells**, which are surface indications that could indicate deeper problems in the system. Code smells are not necessarily problematic on their own, however, they may lead to issues such as bugs or poor maintainability if left unchecked. They are conceptually similar to the stylistic issues mentioned earlier, however they may encompass higher-level structural and design-based problems that are not easily fixed by simple stylistic changes. Examples of code smells include duplicated code, which can be hard to update without introducing bugs, and long methods, which can be difficult to understand and maintain. Therefore, it is often productive to refactor code to eliminate code smells, even if the code is still correct and functional.

Certain linting rules can aid in the refactoring process by broadly identifying code smells and candidate areas for refactoring, suggesting appropriate actions that the user can take. As an example, a linter may detect a fragment of code that is repeated in multiple places: this is a code smell, as discussed previously. The linter may then suggest a code action to automatically apply an *Extract Method* [Fowler 2018] refactoring to avoid code duplication: fig. 2.2 demonstrates how this automatic refactoring process can be performed in the IntelliJ IDEA⁵ IDE.

Possible bugs or errors

Linters may also directly attempt to detect more serious issues in code, such as possible logic errors. This can be helpful for even experienced users to avoid common pitfalls. For example, *Clippy* has `clippy::suspicious` and `clippy::correctness` rule categories to identify code that is very likely to be incorrect or useless. *ESLint* provides several rules to warn against code patterns that are likely to cause runtime errors, such as re-assigning a `const` variable.

Again, linters may attempt to provide auto-fixes for these issues where possible. However, these issues are usually more complex, which may limit the effectiveness or usefulness of auto-fixes: in the case of a suspicious code pattern, the programmer’s intent may not be clear, causing the linter to suggest a fix that does not align with the user’s expectations.

2.1.2 Implementing Linters

This is the old static analysis section - rework this

The vast majority of linters are implemented using static program analysis, analysing source code to extract information about its behaviour without executing the program itself. This is in contrast to dynamic analysis, which is performed on programs as they are run. Although dynamic analysis can yield more precise results as it observes the actual runtime behaviour of the program, it is also more heavyweight as a result.

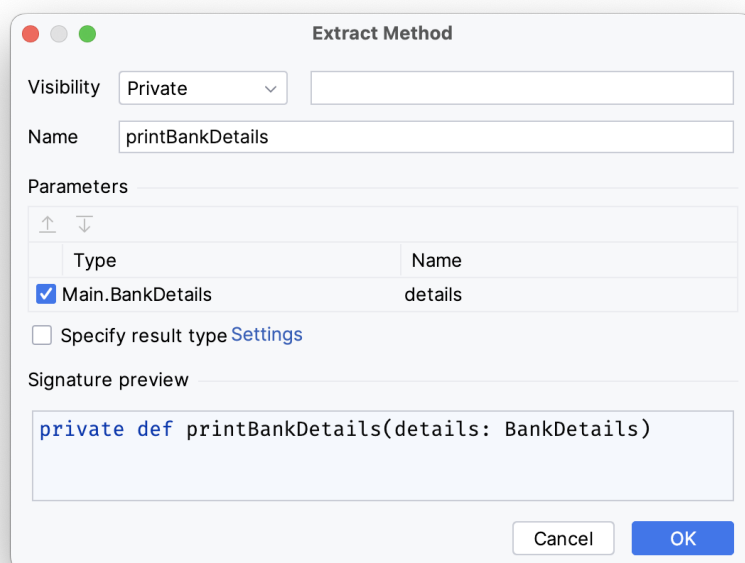
⁵<https://www.jetbrains.com/idea/>


```

1      object Main {
2          def main(args: Array[String]): Unit = {
3              val bankDetails = getBankDetails()
4              println(s"Account name: ${bankDetails.name}")
5              println(s"Account balance: ${bankDetails.balance}")
6          }
7      }

```

(a) A snippet of Scala code. A user may wish to extract the highlighted lines into a separate function.



(b) When a user selects the highlighted lines from fig. 2.2a in IntelliJ IDEA, choosing the *Extract Method* refactoring will open this dialogue to preview changes before applying them.

```

1      object Main {
2          def main(args: Array[String]): Unit = {
3              val bankDetails = getBankDetails()
4              printBankDetails(bankDetails)
5          }
6
7          private def printBankDetails(details: BankDetails): Unit = {
8              println(s"Account name: ${details.name}")
9              println(s"Account balance: ${details.balance}")
10         }
11     }

```

(c) The result of applying the *Extract Method* refactoring using the chosen parameters in fig. 2.2b.

Fig. 2.2: An example of the *Extract Method* refactoring in IntelliJ IDEA.

Given their need to be scalable and efficient enough to handle entire codebases, linters are designed to be fast and lightweight. In most practical cases, therefore, dynamic analysis is too slow and resource-intensive to be used for the purpose of linting. However, dynamic languages such as JavaScript present challenges for static analysis due to their dynamic runtime nature. This has led to the exploration of linters that incorporate dynamic analysis techniques to identify issues that static methods cannot detect [Gong et al. 2015]. Despite these developments, state-of-the-art JavaScript linters utilised in industry, such as *ESLint*, still solely rely on static analysis.

* parsley is a scala library * scala is a static language, strongly typed * this section will therefore focus on the traditional static analysis approach, since dynamic analysis is not as relevant here

<https://scalacenter.github.io/scalafix/docs/users/related-projects.html> Lint on compile vs lint after compile?

AST stuff

* how does HLint work? seems like its sometimes prone to false positives

Static analysis tools can reason about how to safely refactor code in an automated manner, performing refactorings as source-to-source transformations. These transformations may be implemented as simple text-based replacements or more robust rewrite rules that operate on the abstract syntax tree (AST) of the source code.

These tools can perform a variety of tasks, ranging from detecting possible bugs [Johnson 1978; Hovemeyer and Pugh 2004] to formal software verification of program properties [Blanchet et al. 2003].

2.2 Static Analysis for Scala

DSL linting is hard but luckily Parsley is an eDSL so we can just use Scala metaprogramming utilities

2.2.1 Choice of Tooling

The goal of parsley-garnish is to provide linting and refactoring capabilities for the parsley parser combinator library. Since parsley is a Scala library, this project must be implemented using a tool capable of statically analysing Scala code. This section will therefore discuss and evaluate the choices available for implementing parsley-garnish.

Scala compiler plugins The most powerful approach would be to implement parsley-garnish as a compiler plugin [Pickering, Wu, and Németh 2019]. Using the low-level compiler API, it is possible to perform arbitrary code transformations at any step of the compilation process. Compiler plugins therefore offer full freedom to extend the Scala compiler with extra functionality, such as extra passes for code analysis and emitting lint warnings as diagnostics or even compiler errors.

However, this approach has several drawbacks. Firstly, compiler plugins are tightly coupled with the compiler itself, and therefore not portable across major compiler versions. For instance, plugins written for the Scala 3 compiler, known as dotty, are completely incompatible with Scala 2 plugins [LAMP/EPFL 2022]. Additionally, developing compiler plugins requires a deep understanding of arcane and poorly documented compiler internals. Exposing the full compiler API permits unsafe operations that may violate undocumented invariants assumed

by the compiler, leading to exceptions during compilation or even malformed bytecode [Sherwany, Zaza, and Nystrom 2015]. The lack of higher-level abstractions also makes it difficult to implement even trivial tasks such as renaming a field.

For these reasons, it would be preferable to explore other tools that may use compiler plugins themselves but provide a higher-level interface for implementing code analysis and transformations.

Scalameta *Scalameta*⁶ is a metaprogramming framework for Scala that provides a unified interface for performing common metaprogramming tasks. It provides a high-level syntactic API for transforming and pretty-printing Scala source code, as well as a semantic API providing access to semantic information such as type inference and name resolution. *make this clearer about the unification of runtime and compile-time metaprogramming? scalameta does more than scala-reflect* Scalameta is the successor of the earlier `scala.reflect` metaprogramming framework, which parsed source code into lossy trees that discarded syntactic information such as comments and whitespace [Burmako 2017]. *this is called hygiene right?* On the other hand, Scalameta trees are lossless and preserve all syntactic details, a key feature that allows code transformations and refactorings to preserve formatting details.

Scalameta’s semantic API is powered by *SemanticDB*, a compiler-agnostic data model for semantic information in Scala programs. This allows Scalameta to extract semantic information via compiler plugins that emit data in the SemanticDB format. Thus, Scalameta can work with any compiler that supports SemanticDB, rather than being tied to a specific compiler implementation.

Since Scalameta provides a high-level interface for manipulating syntactic and semantic information, it is a promising choice for this project. Being able to access semantic information is especially helpful for implementing more complex code analyses. However, Scalameta’s primary focus is on providing a general metaprogramming framework and therefore lacks API support specifically for implementing linting and refactoring rules. For example, the Scalameta tree transformation utilities do not preserve formatting details when pretty-printed, despite the underlying trees containing this information.

Scalafix *Scalafix*⁷ is a refactoring and linting tool built on top of Scalameta. It specifically provides an API for implementing fine-grained code transformations that preserve comments and formatting details. Scalafix provides a framework for implementing linting rules to emit warnings, as well as rewrite rules to perform automated code transformations [Geirsson 2017]. Since it is built on Scalameta, a major advantage of Scalafix is that it is also compiler-agnostic and could be integrated into any IDE if a plugin is developed for it.

Originally, Scalafix was designed to help automate the process of migrating code from Scala 2 to 3, which involved many breaking changes to the language [Geirsson 2016]. However, Scalafix has since evolved into a general-purpose tool for implementing code transformations and analyses, utilising the powerful syntactic and semantic APIs provided by Scalameta. Scalafix rules can be either syntactic or semantic, depending on whether they require semantic information to perform their analysis. Syntactic rules are faster to run, since they operate purely on the AST without requiring compilation to extract semantic information, but are more limited in the accuracy of analyses they can perform.

⁶<https://scalameta.org/>

⁷<https://scalacenter.github.io/scalafix/>

Scalafix is growing to become the de-facto modern successor to earlier refactoring tools such as Abide⁸ and scala-refactoring⁹. scala-refactoring used `scala.reflect` to implement code transformations, with much extra work utilising the Scala Presentation Compiler AST to preserve formatting details lost by `scala.reflect`. As a result, maintaining the library became difficult and the project was abandoned in favour of a clean implementation using Scalameta, which was designed in part to address the shortcomings of `scala.reflect`.

A drawback of Scalafix is that it is primarily a command-line tool, and therefore by default does not provide a user-friendly interface for interactive usage. However, this can be rectified in the future by integrating Scalafix into the Metals LSP server for Scala, which would allow it to be integrated into any IDE that supports the LSP.

Overall, Scalafix emerges as the most favorable choice for implementing parsley-garnish. It provides high-level APIs specifically for implementing linting and rewrite rules without necessitating extensive knowledge of compiler internals. Scalafix is also being actively developed and maintained, with good basic documentation and a growing number of examples of usage in the wild.

Other tools considered The main alternate contender to Scalafix is the IntelliJ Scala Plugin¹⁰. However, while the plugin offers superior interactive usage within the IntelliJ IDEA IDE, it is tied to the IntelliJ Scala compiler and therefore not portable across other compilers. To maintain flexibility and not tie parsley-garnish to a particular compiler or code editor, Scalafix is a preferable option. Furthermore, documentation is less clear on how to write a Scala plugin for IntelliJ compared to the Scalafix documentation.

WartRemover¹¹ is a linter implemented as a compiler plugin, with support for writing custom rules. However, it only can emit warnings or errors and does not support auto-fixes, making it less suitable for parsley-garnish's goals.

ScalaStyle¹² is primarily a style checker which also supports custom rules. However, it is only able to perform syntactic analyses and does not have access to semantic information, restricting the types of analyses it can perform.

2.3 Parser Combinators

Parsing is the process of extracting structured information from a flat, unstructured representation of the data. Parsers are programs that perform this process, using a specified grammar to determine the structure of the data. They are utilised in a variety of applications such as compilers, interpreters, and processing of data storage formats such as JSON and XML.

Traditionally, parsers have either been written by hand or by using parser generator frameworks such as ANTLR [Parr 2013]. Hand-rolling a parser is a tedious process, requiring the programmer to manually implement the parsing algorithm for the grammar. However, this approach is the most powerful and flexible and can provide excellent performance. Alternatively, parser generators lift the burden of implementing the parsing algorithm, instead requiring the programmer to specify the grammar in the format of a domain-specific language (DSL)

⁸<https://contributors.scala-lang.org/t/whats-the-status-of-abide/>

⁹<https://github.com/scala-ide/scala-refactoring>

¹⁰<https://github.com/JetBrains/intellij-scala>

¹¹<https://www.wartremover.org/>

¹²<http://www.scalastyle.org/>

similar to a high-level grammar. The grammar is then compiled by the parser generator tool to produce a parser in a target language. This approach is less flexible but can be more convenient and less error-prone.

Parser combinators [Hutton 1992], which stem from a functional programming background, are a middle ground between the two approaches. They take the form of an embedded DSL written directly in a general-purpose language, rather than the parser generator approach where the DSL is a separate language. With a parser generator, the provided DSL is often limited in its expressiveness. This is not the case with parser combinators, as the full power of the host language is available to the programmer. This approach also reduces boilerplate code: for example, the programmer does not need to convert between the AST produced by the parser generator and their own AST.

A downside of parser combinators, however, is that they are unstandardised compared to parser generators. Across different implementations, parser combinator APIs can vary significantly, making it difficult to transfer knowledge between different libraries. Experienced users of parser combinators may approach a new library with prior knowledge of general concepts but may have misconceptions about the specifics of the API which can lead to confusion and frustration. This is another motivating reason for the development of parsley-garnish, to lower the barrier of entry for new users of the parsley library.

2.3.1 Parsley

Parsley [Willis and Wu 2018] is a parser combinator library for Scala that provides an API inspired by the parsec [Leijen and Meijer 2001] style of parser combinators. This section provides an illustrative example of a simple expression parser to demonstrate what a parser written in parsley looks like.

Consider the EBNF grammar for a simple expression language shown in fig. 2.3a. The parser in fig. 2.4 will parse an expression into the AST represented by the Scala datatype in fig. 2.3b.

Notice how the parser closely resembles the high-level EBNF grammar. The main differences of note include the use of:

- `map` to transform the result of a parser to help construct tree nodes consisting of a single value.
- `zipped` to combine the results of two parsers to help construct tree nodes consisting of multiple values.
- `←` and `→` operators to guide the direction of parsers.

Except for the possibly cryptic-looking implementation of `num` to parse a series of digits into an integer, the parser is relatively straightforward to understand.

2.3.2 Design Patterns for Parsley

(This background section is a work-in-progress, and will likely expand to include more information about the specific problems and design patterns I choose to explore in the project.)

Willis and Wu [Willis and Wu 2022] describe several design patterns for writing maintainable parsers using parser combinators in Scala. They identified common problems and anti-patterns in parser design, and proposed solutions in the form of design patterns. This provides a guideline for writing idiomatic parsley code for practical parser design, which enables opportunities for the development of linting and refactoring rules.

```
ident ::= "x" | "y" | "z"
num   ::= digit+
expr  ::= factor "+" expr
factor ::= atom "*" factor
atom  ::= ident | num | "(" expr ")"
```

(a) The grammar in EBNF.

```
sealed trait Expr
case class Ident(name: String) extends Expr
case class Num(x: Int) extends Expr
case class Add(x: Expr, y: Expr) extends Expr
case class Mul(x: Expr, y: Expr) extends Expr
```

(b) The Scala AST to parse into.

Fig. 2.3: The grammar and AST for our simple expression language.

```
val ident = "x" | "y" | "z"
val num: Parsley[Int] = digit.foldLeft1(0)((n, d) => n * 10 + d.asDigit)

lazy val expr: Parsley[Expr] = (factor, "+" ~> expr).zipped(Add)
lazy val factor: Parsley[Expr] = (atom, "*" ~> factor).zipped(Mul)
lazy val atom: Parsley[Expr]
  = ident.map(Ident) | num.map(Num) | "(" ~> expr ~> ")"
```

Fig. 2.4: A parser for our simple expression language.

This thesis hopes to explore how these common problems can be formalised into code smells and suspicious code patterns that can be automatically detected using linting rules. Some of the design patterns are also theoretically amenable to automated refactoring, which we hope to explore and implement in `parsley-garnish`.

Not sure how to order things: lint rules first, or implementation?

Chapter 3

Lint Rules in parsley-garnish

Catalogue of lint rules implemented.

Categorise these – but also somehow split into the “simple” rules and the “complex” rules. Simple rules can consist of a single heading, containing: * Explanation of the rule * Simple example to show a diagnostic, and a before and after if it’s fixable * How it’s implemented in the code * Proof (if applicable) * Limitations

Simple rule ideas: * Overly complex parser definitions * Manually calling `implicitSymbol` instead of using the `implicit`

Not sure how to lay out the complex rules yet – so far this is just the left-recursion removal rule. The other complex rule(s) will likely share implementation details with the `Parser/Func` representation, so work from there.

3.1 Avoid Redefining Existing Parsers

* Catch cases when user manually writes out a parser that is already defined in the library

3.2 Simplify Complex Parsers

* Apply parser laws, re-using `Parser` and `Func` representations to do cool things <- should this be a separate rule?

3.3 Ambiguous Implicit Conversions

Heroin is just one letter away from heroine, and implicit conversions are the heroine we don’t deserve.

Jamie Willis, 2024

Implicit conversions are a powerful feature in Scala, allowing for automatic type coercion between types. However, they are often considered a double-edged sword, as they can be easily abused, leading to unexpected behaviour and making code harder to reason about.

* Implicit conversions are controversial, but less evil when it comes to dsls * from the man `odersky` himself: <https://contributors.scala-lang.org/t/can-we-wean-scala-off-implicit-conversions/4388> * Cite Jamie’s design pattern from the scala paper - pattern 0 and 2c * See his discussion in phd thesis (nothing has been elaborated on in the masters thesis)

* have to use `syntacticRule` because we want to be able to provide diagnostics even if compilation fails, due to the exact problem of clashing implicits * so actually this would’ve been better as `lint-on-compile`, to annotate the compiler error message at the exact location - we’d also crucially have more information at that stage * but because the only other option for `scalafix` is `semanticRules` that are completely after compilation time, we can’t

gather even this partial information * wartremover actually has a wart for this [^], but it only targets the definition of implicit conversions, so it's not actually what we'd want out of a lint-on-compile rule

3.4 Remove Explicit Usage of Implicit Conversions

3.5 Refactor to use Parser Bridges

* This would be cool, idk if I have time though, but this should also piggyback off of Func * the pos bridges don't actually exist, so we can ignore that case and just say its too much code synthesis * shouldn't be too bad? idk * indicate limitations that this will only work if the ADT is defined in the same file, in order to extend it

3.6 Left Recursion Removal

Chapter 4

Implementation

Non-terminal detection. This may get reworked/renamed since it's pretty specialised for leftrec rn, and in reality it's just trying to grab all the parsers.

Other util things? ACTUALLY NEED TO DO: import combinators if they aren't already imported

4.1 Parser Representation

Representation of Parsley combinators in parsley-garnish. Compare with approach in Scala Parsley, take cues from the 2018 paper. * Approach to composites? Need to think about this. * For LeftRec: Parse ASTs into a small group of core combinators, but we also need to represent composite combinators as their own case classes – recombine/"simplify" after analysis is concluded, it doesn't really matter if we completely change what combinators are used as long as semantic meaning is preserved. * For others: probably need to parse directly into composite combinators, since we don't want to destructively modify what combinators have been used. * Optimisations: for us, the goal is human readability, so this is interesting to compare to the paper. Lots of similar stuff actually, like top-down peephole optimisations utilising parser laws (I think I do it this way? Need to double check). * For cleanliness to isolate boilerplate: <https://blog.sumtypeofway.com/posts/introduction-to-recursion-schemes.html> – we don't have a generic traversal, but we can decouple the recursive application of a given partial function from the actual pf itself (I've called it .transform for the Parser class)

4.2 Function Representation

Abstraction built over scalafix/meta ASTs to represent functions. Allows us to statically evaluate function composition/flipping etc, so it doesn't turn into one big mess – again, human readability of the transformed output is the goal. Abstraction is again an ADT as a lambda calculus, but with parameter lists so not everything is curried. idk, this is still a work-in-progress. Seems that there might not be enough time to uncurry the leftrec analysis so this design decision might not be super important. Representation as a lambda calc has allocation overhead, but greatly simplifies function evaluation via beta reduction, instead of having to deal with high-level representations of compose/id (not too bad tbh) and flip (annoying). Also attempted to make it typed but that didn't go so well with Scala's limitations on type inference.

* Extracting method arguments (alongside their types) is very painful * Need to unify information from signature (within symbolinformation) and synthetics * synthetics exist in certain cases: .apply methods, showing the concrete type of a generic argument, implicit conversions * from <https://scalacenter.github.io/scalafix/docs/developers/semantic-tree.html>: SemanticTree is a sealed data structure that encodes tree nodes that are generated by the compiler from inferred type parameters, implicit arguments, implicit conversions, inferred .apply and for-comprehensions.

* map, lift (implicit and explicit), zipped, (.as perhaps?) – these should surely boil down into two cases: (x, y).xxx(f) and xxx(f, x, y) * named function literals (val) * named method literals (def) * anonymous functions i.e. lambdas * functions with placeholder syntax * apply methods of case classes - symbol will tell its a class signature

so we use this as a clue to look at synthetics??? * generic bridges – I reckon the information will probably show up in synthetics again

I think look at symbol signature first, then look at synthetics based on some heuristics (e.g. if no symbol sig – if this happens will there even be synthetics?, if class signature)

Chapter 5

Project Plan

5.1 Proposed deliverables and milestones

This project is rather exploratory in nature, and as such, these milestones are highly subject to change as the project progresses. However, I have attempted to outline some key deliverables which I believe to be achievable, as well as some optional deliverables to explore if time permits.

5.1.1 Key deliverables

- Simple lint rules for non-idiomatic combinator usage which can be simplified with an automated refactoring. For example, `endBy(p, sep)` is the idiomatic equivalent to `many(p <* sep)`.
- Detection of left-recursive parsers, which are a bug pattern since they will cause infinite recursion at runtime. It should be possible to perform a static analysis to produce a code transformation to utilise the `postfix` family of combinators to eliminate left-recursion and resolve the bug.
- Lint rules to detect ambiguity in parsers, which can be a bug pattern since the parser may not behave as expected. It may be possible to provide an automated fix in some cases, but more research has to be done to determine the feasibility of this.

Conversely, it may also be possible to detect overuse of the `atomic` combinator. Users may overuse this combinator to avoid ambiguity, but overuse is a code smell as it can negatively impact error messages reported by the parser.

- Proofs that all code transformations preserve the semantics of the original parser.

5.1.2 Optional deliverables

- Lint rules that warn against side-effecting parsers. Parsley employs aggressive parser optimisations which may assume that parsers are pure. The lint rule would remind users to use the `impure` combinator to avoid parsley from optimising away any intended behaviour.
- Lint rules to detect when a parser using `chain/infix/postfix` combinators can be rewritten to use a precedence table, and the automated refactoring to accompany this rule. Further lints/refactorings related to precedence tables may also be possible, such as looking at type information to determine the most appropriate shape (`0ps/S0ps/G0ps`).
- An editor plugin or extension to provide integration of parsley-garnish into the code editor. As explored in the background section, it is possible to provide this integration for Scalafix, but the difficulty of this task is unknown.
- Exploration of possible pedagogical applications of parsley-garnish. For example, lint rules may be useful for marking submissions for the Second Year WACC compilers project.

5.2 Current Progress

Most of the work so far has been on research and experimentation with Scalafix. Although there was good documentation to get started, Scalafix is an active project and the documentation is not always fully comprehensive or up-to-date. I still have some concerns about whether the Scalafix API is expressive enough to implement all of the analyses I have in mind, but I will continue to explore this as the project progresses.

From my experimentation, I have a relatively good idea of how to implement most of the “simple” lint rules.

At the time of writing, I have also implemented a proof-of-concept rule for transforming left-recursive parsers, on a small subset of parsley’s combinators. This has been encouraging but much work remains to be done to be able to apply the rule on more complex parsers. The current results of the transformations are also sometimes too complex, so I will need to investigate methods such as defunctionalisation and/or alternate methods of representing terms to simplify the output statically. Furthermore, the transformation is likely to be correct but not proven yet.

5.3 Project Timeline

My immediate plan is to continue work on the left-recursion transformation, while also working on the simpler lint rules. I hope to complete the left-recursion rule by early March, although unforeseen difficulties may arise (such as if output simplification proves to be more difficult than expected). If it is not complete or close to complete by then, I will re-evaluate the feasibility of the rest of the milestones and adjust accordingly. I have already communicated to my supervisor that this left-recursion rule will likely be the most involved deliverable of the project, although there is also a chance that many of the techniques employed could be reused for other transformations.

The next two months will be spent exploring other milestones. I plan on approaching this by alternating research and implementation, as I have found that this has been a good way to make progress so far, while also guaranteeing that I will make progress on the final deliverable.

By the beginning of summer term, I would also like to have a very rough draft of the final report completed, based on the work completed so far. This will allow me to communicate with my supervisor on the state of the project, as well as leaving me ample time to edit and refine the report. The remaining milestones to cover will be decided during this meeting; at this point, having implemented a number of rules, I hope to have a much better grasp on how long each deliverable could take. This remaining time will then be divided between writing the report and implementing the agreed-upon milestones.

Chapter 6

Evaluation Plan

A large corpus of student WACC parsers is available to evaluate the real-world applicability of `parsley-garnish`. This can be utilised to assess the tool’s ability to achieve quantitative and qualitative metrics.

The large number of parsers available to test can help evaluate the correctness and accuracy of each lint rule. We might find that certain rules are more prone to false positives than others, or that some automated refactorings may introduce new issues such as compilation errors. These quantitative metrics can be counted and used to evaluate rules against each other, as well as the overall effectiveness of the tool.

We can also estimate the proportion of issues that `parsley-garnish` is able to detect automatically. This can be done by comparing the number of issues detected automatically against the number of issues found manually by a code reviewer.

Qualitative metrics such as the quality of automated fixes can also be assessed by manually reviewing the output of the tool. For example, an issue I pointed out in the project plan is that the output of the left-recursion transformation may be in a form that is “ugly” when it could be simplified to a much more human-readable form. The quality of fixes would affect the practicality of the `parsley-garnish`, since users may be reluctant to apply automated fixes if the output quality can be poor.

Furthermore, the performance of `parsley-garnish` can be evaluated in terms of execution time and resource consumption. There are no direct competitors to benchmark against, but we can still evaluate the tool’s performance against a baseline of acceptable performance, or perhaps a collection of other Scala static analysis tools deemed to be of similar complexity. I am thinking of benchmarking the tool during development as well, since it may affect design decisions such as implementing some rules as purely syntactic rather than semantic rules. I am not sure how much resource overhead a semantic rule would incur, so this could be insignificant or a major tradeoff (between speed, correctness, and accuracy) to consider.

Finally, it is possible to gather feedback from students interested in using `parsley-garnish` to improve their WACC parsers. User surveys and interviews would provide insights into the tool’s usability, helpfulness, and areas for improvement. Although the timing of the WACC Lab does not align with the project timeline, I am hoping that some students may still be interested in trying out the tool on their final WACC submissions or possibly personal projects using `parsley`.

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Appendix A

Proofs of Parser Laws

The approach taken to prove the following parser laws for `parsley` is via equational reasoning on `gigaparsec` semantics, under the assumption that their semantics are equivalent. While there is no formal proof of this equivalence at the present, `gigaparsec` was designed to have semantics equivalent to `parsley`'s.

A.1 Left absorption for `fmap`

```

f <$> empty
=   { applicative functor law }
    pure f <*> empty
=   { definition of <*> }
    liftA2 ($) (pure f) empty
=   { semantics of liftA2 }
    Parsec $ \st ok err ->
      let ok' x st' = (unParsec empty) st' (ok . (x $)) err
      in (unParsec $ pure f) st ok' err
=   { semantics of empty }
    Parsec $ \st ok err ->
      let ok' x st' = (unParsec $ raise (`emptyErr` 0)) st' (ok . (x $)) err
      in (unParsec $ pure f) st ok' err
=   { semantics of raise }
    Parsec $ \st ok err ->
      let ok' x st' = (unParsec $ Parsec $ \st'' _ bad ->
        useHints bad (emptyErr st'' 0) st') st' (ok . (x $)) err
      in (unParsec $ pure f) st ok' err
=   {  $\beta$ -reduction }
    Parsec $ \st ok err ->
      let ok' x st' = useHints err (emptyErr st' 0) st'
      in (unParsec $ pure f) st ok' err
=   { semantics of pure }
    Parsec $ \st ok err ->
      let ok' x st' = useHints err (emptyErr st' 0) st'
      in (unParsec $ Parsec $ \st'' ok'' _ -> ok'' f st'') st ok' err
=   {  $\beta$ -reduction }
    Parsec $ \st ok err ->
      let ok' x st' = useHints err (emptyErr st' 0) st'

```

```
      in ok' f st
=    { inline ok' }
      Parsec $ \st ok err -> useHints err (emptyErr st 0) st
=    { rearrange and  $\alpha$ -conversion }
      Parsec $ \st _ bad -> useHints bad ((`emptyErr` 0) st) st
=    { fold definition of raise }
      raise (`emptyErr` 0)
=    { fold definition of empty }
      empty
```