

Alexandre Rodrigues Baldé A70373  
Questão 1

- ① Variáveis livres:  $z$  ✓  
Variáveis ligadas:  $y, a, n$

②  $(\lambda y. \lambda a. (\lambda n. \lambda y. y n z) (y a)) ((\lambda n. \lambda y. n) (\lambda n. \lambda y. n))$

A expressão contém 3  $\beta$ -redexes.

$$(\lambda y. \lambda a. (\lambda n. \lambda y. y n z) (y a)) ((\lambda n. \lambda y. n) (\lambda n. \lambda y. n))$$

$$\rightarrow_{\beta} \lambda a. (\lambda n. \lambda y. y n z) (((\lambda n. \lambda y. n) (\lambda n. \lambda y. n)) a) \quad \checkmark$$

$$(\lambda y. \lambda a. (\lambda n. \lambda y. y n z) (y a)) ((\lambda n. \lambda y. n) (\lambda n. \lambda y. n))$$

$$\rightarrow_{\beta} (\lambda y. \lambda a. (\lambda u. u (y a)_z)) ((\lambda n. \lambda y. n) (\lambda n. \lambda y. n)) \quad \checkmark$$

(note-se a  $\alpha$ -conversão de  $y$  para  $u$ )

No caso do último redex:

$$(\lambda y. \lambda a. (\lambda n. \lambda y. y n z) (y a)) ((\lambda n. \lambda y. n) (\lambda n. \lambda y. n))$$

$$\rightarrow (\lambda y. \lambda a. (\lambda n. \lambda y. y n z) (y a)) (\lambda v. (\lambda n. \lambda y. n)) \quad \checkmark$$

No último termo note-se a  $\alpha$ -conversão da variável ligada  $y$  para  $v$ .

- ③ a)  $\lambda n. n$   
b)  $(\lambda n. n)_z$   
c)  $(\lambda n. \lambda y. y) ((\lambda n. n_n) (\lambda n. n_n))$   
d)  $((\lambda n. n_n) (\lambda n. n_n))$  ✓

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Questão 2

$$\textcircled{1} \quad = \underline{F \ B} \ (\underline{I \ B}) \ (\underline{K \ K \ B}) \ (\underline{K \ I}) \\ = (\underline{F \ B}) (\underline{I \ B}) \ (\underline{K \ K \ B}) (\underline{K \ I})$$

$$\rightarrow_p ((\lambda b. b) \ \underline{\underline{(I \ B)}}) (\underline{K \ K \ B}) (\underline{K \ I})$$

$$\rightarrow_p (\underline{I \ B}) (\underline{K \ K \ B}) (\underline{K \ I})$$

$$\rightarrow_p \underline{B} \ (\underline{K \ K \ B}) (\underline{K \ I})$$

$$\rightarrow_p (\lambda y. \lambda z. \underline{(\underline{K \ K \ B})} (\underline{y \ z})) (\underline{K \ I})$$

$$\rightarrow_p (\lambda z. (\underline{K \ K \ B}) ((K \ I) z))$$

$$\rightarrow_p (\lambda z. \underline{(\underline{K \ B})} ((K \ I) z)) \quad \times$$

$$\rightarrow_p (\lambda z. \underline{B} \ (\underline{(\underline{K \ I})} z)) \quad \times$$

$$\rightarrow_p (\lambda z. B \ (\underline{I} z))$$

$$\rightarrow_p (\lambda z. B z)$$

$$\dashv \alpha (\lambda u. B u) = (\lambda u. (\lambda n. \lambda y. \lambda z. n(y z)) u)$$

$$\rightarrow_p B \quad 0.5$$

B é uma forma normal

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$$\textcircled{2} \quad B(FI) = (\lambda x. \lambda y. \lambda z. u(yz)) ((\lambda a. \lambda b. b) (\lambda u. u)) .$$

Seja EXP o  $\lambda$ -termo anterior. Se for tipificável, então para qualquer  $X$  tal que  $\text{EXP} \xrightarrow[\beta]{}^* X$ ,  $X$  terá o mesmo tipo de EXP.

$B(FI)$  (ordem applicativa)  
 $\rightarrow_B B(\lambda b.b)$   
 $\rightarrow_B \lambda y. \lambda z. (\lambda b.b)(yz)$   
 $\rightarrow_B \lambda y. \lambda z. \frac{yz}{y z} : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta$ , que se pode tipificar assim:  $(\lambda y : (\alpha \rightarrow \beta), \lambda z : \alpha,$   
 logo o tipo de EXP suaí EXP :  $(\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta$

Quais tipos nas Variáveis:

$$(\lambda x:\gamma. \lambda y:(\alpha \rightarrow \beta). \lambda z:\delta. n(y,z))((\lambda a:\delta. \lambda b:\beta. b) \\ (\lambda x:\gamma. n)) : (\alpha \rightarrow \beta) \rightarrow \alpha \rightarrow \beta$$

✓

onde  $\gamma = \beta \rightarrow \beta$

Büro 3

Questão 3  
 ① seja  $\alpha = \lambda n. \lambda l. \text{listang } l \text{ of } \{\langle n, 0 \rangle, \lambda h. \lambda t.$   
 $\quad \quad \quad \text{if } h > 0 \text{ then } \langle h, n \rangle \text{ else } \langle n, h \rangle\}$ .

let emp =  $\lambda$   
 par = 8  
 lis =  $q :: (5+2) :: \text{nil}$   
 in emp (( $\lambda y. y * y$ ) par) lis  
 $= (\lambda \text{emp}. \lambda \text{par}. \lambda \text{lis}. \text{emp} ((\lambda y. y * y) \text{par}) \text{lis}) \alpha 8 (q :: (5+2) :: \text{nil})$   
 $\xrightarrow{3} p \alpha ((\lambda y. y * y) 8) (q :: (5+2) :: \text{nil})$   
 $\xrightarrow{2} p \text{listcase } (q :: (5+2) :: \text{nil}) \text{ of } (((\lambda y. y * y) 8),)$   
 $\quad \lambda h. \lambda t. \text{if } h > 0 \text{ then } \{h, (\lambda y. y * y) 8\}$   
 $\quad \quad \text{else } ((\lambda y. y * y) 8, h)$   
 }

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$\rightarrow_p (\lambda h. \lambda t. \text{if } h > 0 \text{ then } \{ h, (\lambda y. y * y) 8 \} \\ \text{use } \{ ((\lambda y. y * y) 8, h) \})$   
 $q$   
 $(5+2) :: \text{nil}$

$\rightarrow_p^2 \text{ if } q > 0 \text{ then } \{ q, (\lambda y. y * y) 8 \} \\ \text{use } \{ ((\lambda y. y * y) 8, q) \}$

$\rightarrow_p (q, (\lambda y. y * y) 8)$

② Seja  $\pi_1 = l : \text{List Bool}, f : \text{Int} \rightarrow \text{List Bool} \rightarrow \text{Bool}$

1.  $\pi_1 \vdash \lambda n. \text{listcase } l \text{ of } (\text{True}, \lambda h. \lambda t. h \wedge (f \ w)) : \text{Int} \rightarrow \text{Bool}$   
1.1.  $\pi_2 \vdash \text{listcase } l \text{ of } (\text{True}, \lambda h. \lambda t. h \wedge (f \ w)) : \text{Bool}$   
1.1.1.  $\pi_2 \vdash l : \text{List Bool}$   
1.1.2.  $\pi_2 \vdash \text{True} : \text{Bool}$   
1.1.3.  $\pi_2 \vdash \lambda h. \lambda t. h \wedge (f \ w) : \text{Bool} \rightarrow \text{List Bool} \rightarrow \text{Bool}$   
1.1.3.1.  $\pi_3 \vdash \lambda t. h \wedge (f \ w) : \text{List Bool} \rightarrow \text{Bool}$   
1.1.3.1.1.  $\pi_4 \vdash h \wedge (f \ w) : \text{Bool}$   
1.1.3.1.1.1.  $\pi_4 \vdash h : \text{Bool}$   
1.1.3.1.1.2.  $\pi_4 \vdash (f \ w) : \text{Bool}$   
1.1.3.1.1.2.1.  $\pi_4 \vdash f \ w : \text{List Bool} \rightarrow \text{Bool}$   
1.1.3.1.1.2.1.1.  $\pi_4 \vdash f : \text{Int} \rightarrow \text{List Bool} \rightarrow \text{Bool}$   
1.1.3.1.1.2.1.2.  $\pi_4 \vdash w : \text{Int}$   
1.1.3.1.1.2.2.  $\pi_4 \vdash w : \text{List Bool}$

onde  $\pi_2 = \pi_1, n : \text{Int}$

$\pi_3 = \pi_2, h : \text{Bool}$

$\pi_4 = \pi_3, + : \text{List Bool}$

③ a)  $\text{letrec fun} = \lambda l. \text{listcase } l \text{ of } (\text{nil}, \lambda h. \lambda t. \\ \text{listcase } t \text{ of } \{ n :: \text{nil}, \lambda h1. \lambda t1. \\ h :: \text{fun}((h+h1)::t1) \})$

in fun

$\langle \text{exp} \rangle ::= \text{empty} \mid \text{node } \langle \text{exp} \rangle \langle \text{exp} \rangle$   
 $\mid \text{truecase } \langle \text{exp} \rangle \text{ of } ((\langle \text{exp} \rangle, \langle \text{exp} \rangle) \mid \dots)$

$\langle \text{type} \rangle ::= \dots$   
 $\mid \text{RoseTree } \langle \text{type} \rangle$

m de inferência de tipos

$$\frac{\text{empty} : \text{RoseTree } \theta}{T \vdash e : \theta} \quad \frac{T \vdash e_1 : \theta \quad T \vdash e_2 : \text{list } \theta}{T \vdash (\text{node } e_1, e_2) : \text{RoseTree } \theta}$$
$$\frac{T \vdash e : \text{RoseTree } \theta \quad T \vdash e_1 : \theta' \quad T \vdash e_2 : \theta \rightarrow \text{RoseTree } \theta \rightarrow \theta'}{T \vdash \text{truecase } e \text{ of } (e_1, e_2) : \theta'}$$

$\langle \text{efun} \rangle ::= \dots$   
 $\mid \text{empty} \mid \text{node } \langle \text{efun} \rangle \langle \text{efun} \rangle$

m de atalaiaõ

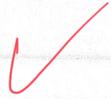
$$\frac{e \Rightarrow z \quad e_1 \Rightarrow z_1 \quad e \Rightarrow \text{nil} \quad e_1 \Rightarrow z}{e \text{ node } z z_1 \quad \text{truecase } e \text{ of } (e_1, e_2) \Rightarrow z}$$
$$\frac{e \Rightarrow \text{node } z z_1 \quad e_2 \not\Rightarrow z_1 \Rightarrow z_2}{\text{truecase } e \text{ of } (e_1, e_2) \Rightarrow z_2}$$

true

$$\text{sumTree} \equiv \lambda t. \text{truecase } t \text{ of } (0, \lambda n. \text{ftree.} \\ (\text{sumTree dist ts}) + n)$$
$$\text{sumTree dist} \equiv \lambda l. \text{listcase } l \text{ of } (0, \lambda t. \text{ftree.} \\ (\text{sumTree f1}) + \text{sumTree dist fts})$$

sumTree

④ empty = @1 ()  
 node =  $\lambda z. \lambda z_1. @2(z, z_1)$



trucase  $e$  of  $(e_1, e_2) = \text{sumcase } (\lambda () . e_1, \lambda (n, y) . e_2 n y)$

RoseTree <type> = Unit + <type> × List (RoseTree <type>)

### Questão 3

b) letrec map =  $\lambda f. \lambda l. \text{listcase } l \text{ of } (\text{nil}, \lambda h. \lambda t. f h :: \text{map } f t)$   
 $f \text{ un} = \dots (*)$   
 in fun  $(\text{rec } (\lambda l. l :: \text{map } (\lambda n. n + 1) l))$



ft) Ver Questão 3 a) numa folha anterior, peço  
 desculpa pela confusão.