

Licenciatura em Ciências da Computação

Semântica de Linguagens de Programação 2021/2022

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-

Tipos de dados para a linguagem While

```
In [ ]: :ext InstanceSigs
        :ext BangPatterns
```

```
In [ ]: import qualified Data.Map    as M
        import qualified Data.Maybe as Maybe

        import           Data.List (intercalate)

        -- For Natural Semantics
        import qualified Control.Monad.State.Strict as St
        import           Control.Monad                (when)
```

```
In [ ]: type Z = Integer

type Var = String

data Aexp
  = Num Z
  | Var Var
  | Aexp `Plus` Aexp
  | Aexp `Mul` Aexp
  | Aexp `Minus` Aexp
  deriving Eq
```

```
In [ ]: data Bexp
  = T
  | F
  | Aexp `Eq` Aexp
  | Aexp `Le` Aexp
  | Aexp `Lt` Aexp
  | Aexp `Ge` Aexp
  | Not Bexp
  | Bexp `And` Bexp
  | Bexp `Or` Bexp
  deriving Eq
```

```
In [ ]: type State = M.Map Var Z

getSt :: State -> Var -> Z
getSt st var = Maybe.fromMaybe 0 (M.lookup var st)

stUpdate :: State -> Var -> Z -> State
stUpdate st var v = M.insert var v st
```

```
In [ ]: subAexp :: Var -> Aexp -> Aexp -> Aexp
subBexp :: Var -> Aexp -> Bexp -> Bexp

arithEval :: Aexp -> (State -> Z)
boolEval :: Bexp -> (State -> Bool)
```

State Helpers

```
In [ ]: -- Igual ao tipo State, mas com instância de Show legível.
newtype State' = St {
  getState :: State
} deriving (Eq)

instance Show State' where
  show = showState . getState
```

```
In [ ]: data Stm
        = Var `Assign` Aexp
        | Skip
        | Stm `Comp` Stm
        | IfThenElse Bexp Stm Stm
        | WhileDo Bexp Stm
        deriving (Eq)
```

Example programs

Retirados da ficha 1.

```
In [ ]: ex1State :: State
ex1State = M.fromList [("n", 6), ("x", 3), ("y", 2)]

swap :: Stm
swap = Comp (Comp c1 c2) c3
  where
    c1 = Assign "n" (Var "x")
    c2 = Assign "x" (Var "y")
    c3 = Assign "y" (Var "n")

minProg :: Stm
minProg = IfThenElse b1 if2 if3
  where
    b1 = Lt (Var "x") (Var "y")
    b2 = Lt (Var "x") (Var "z")
    b3 = Lt (Var "y") (Var "z")
    if3 = IfThenElse b3 (Assign "m" (Var "y")) (Assign "m" (Var "
z"))
    if2 = IfThenElse b2 (Assign "m" (Var "x")) (Assign "m" (Var "
z"))

expProg :: Stm
expProg = Comp assgn while
  where
    assgn = Assign "r" (Num 1)
    while = WhileDo bexp whileStm
    bexp = Var "y" `Ge` Num 1
    whileStm = Comp
      (Assign "r" (Mul (Var "r") (Var "x")))
      (Assign "y" (Minus (Var "y") (Num 1)))

fact :: Stm
fact = Comp assgn while
  where
    assgn = Assign "f" (Num 1)
    while = WhileDo bexp whileStm
    bexp = Var "n" `Ge` Num 1
    whileStm = Comp
      (Assign "f" (Mul (Var "f") (Var "n")))
      (Assign "n" (Minus (Var "n") (Num 1)))
```

Natural Semantics

```
In [ ]: evalNS :: State -> Stm -> State
evalNS st stm = St.execState (helper stm) st
  where
    helper :: Stm -> St.State State ()
    ...
```

```
In [ ]: evalNS ex1State fact

fromList [("f",720),("n",0),("x",3),("y",2)]
```

```
In [ ]: evalNS ex1State expProg

fromList [("n",6),("r",9),("x",3),("y",0)]
```

Structural Operational Semantics

```
In [ ]: -- Uma transição em semântica operacional estrutural:
-- * ou dá origem a um estado (Left)
-- * ou dá origem a um comando intermédio, juntamente com um novo estado (Right)
stepSOS :: State -> Stm -> Either State (Stm, State)

-- Devolve a transição após o número de passos de execução pedido, se for possível,
-- numa lista com todas as transições que lhe precederam.
nstepsSOS :: State -> Stm -> Integer -> [Either State' (Stm, State')]

-- Imprimir uma configuração numa string legível
-- e.g. mapM_ (putStr . helperSOS) $ nstepsSOS ex1State fact 10
helperSOS :: Either State' (Stm, State') -> String

evalSOS :: State -> Stm -> State'
```

```
In [ ]: evalSOS ex1State fact

f := 720; n := 0; x := 3; y := 2
```

```
In [ ]: evalSOS ex1State expProg

n := 6; r := 9; x := 3; y := 0
```

Abstract Machine 1

```

In [ ]: -- Mapping from variable names to positions in memory.
-- Used during "compilation" of While code to AM1 bytecode.
type Env = M.Map Var Z

getEnv :: Env -> Var -> Z
getEnv e var = e M.! var

type NextAddr = Z

data EnvStateAM1 = EnvSt {
  getEnvSt :: !Env,
  getNextAddr :: !NextAddr
} deriving (Eq, Show)

```

```

In [ ]: getEnv :: Env -> Var -> Z
getEnv e var = e M.! var

data AM1Instr
  = PUSH Z
  | ADD
  | MULT
  | SUB
  | TRUE
  | FALSE
  | EQUAL
  | LE
  | GE
  | LTHAN
  | AND
  | OR
  | NEG
  | PUT Z
  | GET Z
  | NOOP
  | BRANCH AM1Code AM1Code
  | LOOP AM1Code AM1Code
  deriving (Eq, Show)

type AM1Code = [AM1Instr]

```

```

In [ ]: aexpToAM1Code :: EnvStateAM1 -> Aexp -> (AM1Code, EnvStateAM1)
aexpToAM1Code m@(EnvSt e nxtAdr) a = case a of
    Num n -> ([PUSH n], m)
    Var var -> case M.lookup var e of
        Nothing -> ([GET nxtAdr], EnvSt (M.insert var nxtAdr e) (nxtAdr
+ 1))
        Just adr -> ([GET adr], m)
    ae `Plus` ae' ->
        let (code, m') = aexpToAM1Code m ae
            (code', m'') = aexpToAM1Code m' ae'
        in (concat [code', code, [ADD]], m'')
    ae `Mul` ae' ->
        ...

bexpToAM1Code :: EnvStateAM1 -> Bexp -> (AM1Code, EnvStateAM1)
bexpToAM1Code m@(EnvSt e nxtAdr) b = case b of
    T -> ([TRUE], m)
    F -> ([FALSE], m)
    ae `Eq` ae' ->
        let (code, m') = aexpToAM1Code m ae
            (code', m'') = aexpToAM1Code m' ae'
        in (concat [code', code, [EQUAL]], m'')
    ae `Le` ae' ->
        ...

```

```

In [ ]: whileToAM1 :: Stm -> (AM1Code, EnvStateAM1)
whileToAM1 stm = St.runState (helper stm) (EnvSt M.empty 0)
    where
        helper :: Stm -> St.State EnvStateAM1 AM1Code
        helper (var `Assign` aexp) = do
            ...
        helper Skip = return [NOOP]
        helper (c1 `Comp` c2) = do
            ...
        helper (IfThenElse b c1 c2) = do
            ...
        helper (WhileDo b c) = do
            memSt <- St.get
            let (predCode, memSt') = bexpToAM1Code memSt b
            St.put memSt'
            loopCode <- helper c
            return [LOOP predCode loopCode]

```

```

In [ ]: type Stack = [Either Z Bool]

-- Mapping from address positions to the values they contain.
-- Should look like:
-- 0 -> n_1
-- 1 -> n_2
-- 2 -> n_3,
-- ...
-- k -> n_k
-- and so on, where n_i are integers.
type Memory = M.Map Z Z

type AM1Config = (AM1Code, Stack, Memory)

```

```
In [ ]: stepAM1 :: AM1Config -> AM1Config
stepAM1 conf@([], stack, mem) = conf
stepAM1 (c : cs, stack, mem) = case c of
  PUSH n -> (cs, Left n : stack, mem)
  ADD -> case stack of
    Left z1 : Left z2 : stack' ->
      (cs, Left (z1 + z2) : stack, mem)
    _ -> error "ADD: invalid stack for operation!"
  ...
  BRANCH ins ins' -> case stack of
    Right b : stack' ->
      let instr = if b then ins else ins'
      in (instr, stack', mem)
    _ -> error "BRANCH: invalid stack for operator!"

  LOOP ins ins' -> (ins ++ [BRANCH (ins' ++ [LOOP ins ins']) [NOOP]] ++
cs, stack, mem)
```

```
In [ ]: initConfigAM1 :: State -> Stm -> (AM1Config, Env)
initConfigAM1 initSt stm =
  let code :: AM1Code
      envSt :: EnvStateAM1
      (code, envSt) = whileToAM1 stm

      environ = getEnvSt envSt

      memory :: Memory
      memory = M.fromList [(getEnv environ variable, getSt initSt variable) | variable <- M.keys environ]
  in ((code, [], memory), environ)

-- Dado um estado inicial e um comando da linguagem while, simula a sua execução
-- na máquina abstrata AM1.
-- Devolve as variáveis usadas no programa, e os valores que estavam nas respectivas
-- posições de memória aquando da terminação da execução.
-- Pode não terminar! (Halting problem).
runStmInAM1 :: State -> Stm -> M.Map Var Z
runStmInAM1 initSt stm =
  let (init@(initCode, initStack, initMemory), environ) = initConfigAM1 initSt stm
      run :: AM1Config -> AM1Config
      run !cfg =
        let cfg' = stepAM1 cfg
        in if cfg' == cfg then cfg else run cfg'
      (finalCode, finalStack, finalMemory) = run init

      varsToValues = M.fromList [(var, finalMemory M.! (environ M.! var)) | var <- M.keys environ]
  in varsToValues
```

```
In [ ]: runStmInAM1 ex1State minProg
runStmInAM1 ex1State swap
runStmInAM1 ex1State expProg
runStmInAM1 ex1State fact

fromList [("m",0),("x",3),("y",2),("z",0)]

fromList [("n",3),("x",2),("y",3)]

fromList [("r",9),("x",3),("y",0)]

fromList [("f",720),("n",0)]
```

Abstract Machine 2

```
In [ ]: -- Mapping from variable names to positions in memory.
-- Used during "compilation" of While code to AM1 bytecode.
type Env = M.Map Var Z

type NextAddr = Z

-- Program counter associated with each instruction.
-- Must be positive, starts at 1, each instruction has a unique PC value,
-- and strictly increases by 1 unit with every atomic instruction.
type ProgramCounter = Z

type Stack = [Either Z Bool]

-- Mapping from address positions to the values they contain.
-- Should look like:
-- 0 -> n_1
-- 1 -> n_2
-- 2 -> n_3,
-- ...
-- k -> n_k
-- and so on, where n_i are integers.
type Memory = M.Map Z Z
```

```
In [ ]: data AM2Instr
      = PUSH Z
      | ...
      | LABEL ProgramCounter
      | JUMP ProgramCounter
      | JUMPFALSE ProgramCounter
      deriving (Eq, Show)

type AM2Code = [AM2Instr]
```

```
In [ ]: type AM2Config = (ProgramCounter, AM2Code, Stack, Memory)

type AM2AnnotatedProgram = M.Map ProgramCounter AM2Instr
```



```
In [ ]: data EnvStateAM2 = EnvSt2 {
        getEnvSt  :: !Env,
        getNxtAdr :: !NextAddr,
        getInstrs :: AM2AnnotatedProgram,
        getNxtPC  :: ProgramCounter
      } deriving (Eq)
```

Tradução de expressões aritméticas e booleanas para "bytecode" AM2

```
In [ ]: aexpToAM2Code :: Aexp -> St.State EnvStateAM2 AM2Code

bexpToAM2Code :: Bexp -> St.State EnvStateAM2 AM2Code
bexpToAM2Code b = case b of
    ...
    be `Or` be' -> do
        -- Careful with the order with which this is done - whichever
        -- is done first
        -- puts its code on the stack first, so the second operand has
        -- to go first.
        code' <- bexpToAM2Code be'
        code <- bexpToAM2Code be
        St.modify' (\(EnvSt2 environ nxtAdr instrs nxtPC) -> EnvSt2 environ
            nxtAdr (M.insert nxtPC OR instrs) (nxtPC + 1))
        return $ concat [code', code, [OR]]
    ...
```

```
In [ ]: :ext FlexibleContexts
```

```

In [ ]: whileToAM2 :: Stm -> (AM2Code, EnvStateAM2)
whileToAM2 stm = St.runState (helper stm) (EnvSt2 M.empty 0 M.empty 1)
  where
    incrCounter = do
      EnvSt2 e nA is nxtPC <- St.get
      St.put $ EnvSt2 e nA is $ nxtPC + 1
      return nxtPC

    helper :: Stm -> St.State EnvStateAM2 AM2Code
    helper (var `Assign` aexp) = do
      code <- aexpToAM2Code aexp
      EnvSt2 environ nxtAdr instrs nxtPC <- St.get
      case M.lookup var environ of
        Nothing -> do
          let instr = PUT nxtAdr
          St.put $ EnvSt2 (M.insert var nxtAdr environ) (nxtAdr
+ 1) (M.insert nxtPC instr instrs) (nxtPC + 1)
          return $ code ++ [instr]
        Just n -> do
          let instr = PUT n
          St.put $ EnvSt2 environ nxtAdr (M.insert nxtPC instr
instrs) (nxtPC + 1)
          return $ code ++ [instr]
    helper Skip = do
      ...
    helper (c1 `Comp` c2) = do
      code1 <- helper c1
      code2 <- helper c2
      return $ code1 ++ code2

-- O código máquina gerado para o comando IfThenElse e para o com
-- ando WhileDo
-- é complexo porque deve primeiro gerar o código dos subcomandos
-- e predicados,
-- e só depois colocar as instruções de salto e labels - cujo pro
-- gram counter
-- terá de ser guardado antes da tradução dos subcomandos.
--
-- Ver incrCounter.
helper (IfThenElse b c1 c2) = do
  predCode <- bexpToAM2Code b
  jzProgCounter <- incrCounter
  thenCode <- helper c1
  elseProgCounter <- incrCounter
  elseCode <- helper c2
  afterIfProgCounter <- incrCounter
  let ifJump      = JUMPFALSE elseProgCounter
      elseLabel   = LABEL elseProgCounter
      jumpToRest  = JUMP afterIfProgCounter
      restLabel   = LABEL afterIfProgCounter
  EnvSt2 environ nxtAdr instrs _ <- St.get
  let jumps = M.fromList [(jzProgCounter, ifJump), (elseProgCou
nter, jumpToRest), (afterIfProgCounter, restLabel)]
  -- Incrementa-se o contador de código devido ao LABEL final,
  -- que apontará
  -- para o código depois do IfThenElse, se existir.
  St.put $ EnvSt2 environ nxtAdr (instrs `M.union` jumps) (afte

```

```

rIfProgCounter + 1)
    return $ predCode ++ [ifJump] ++ thenCode ++ [jumpToRest] ++
elseCode ++ [restLabel]

    helper (WhileDo b c) = do
        boolTestCounter <- incrCounter
        predCode <- bexpToAM2Code b
        jzProgCounter <- incrCounter
        loopCode <- helper c
        jumpCounter <- incrCounter
        afterWhileCounter <- incrCounter
        let whileLabel = LABEL boolTestCounter
            whileJump = JUMPFALSE afterWhileCounter
            loopJump = JUMP boolTestCounter
            restLabel = LABEL afterWhileCounter
        EnvSt2 environ nxtAdr instrs _ <- St.get
        let jumps = M.fromList [(boolTestCounter, whileLabel), (jzPro
gCounter, whileJump), (jumpCounter, loopJump), (afterWhileCounter, restLa
bel)]
        St.put $ EnvSt2 environ nxtAdr (instrs `M.union` jumps) (afte
rWhileCounter + 1)

        return $ [whileLabel] ++ predCode ++ [whileJump] ++ loopCode
++ [loopJump] ++ [restLabel]

```

```

In [ ]: -- Given an AM2 configuration, execute a single instruction
-- and transition into the next configuration.
stepAM2 :: AM2Config -> AM2AnnotatedProgram -> AM2Config
stepAM2 conf@(_, [], stack, mem) _ = conf
stepAM2 (pc, c : cs, stack, mem) ann = case c of
    PUSH n -> (pc', cs, Left n : stack, mem)
    ADD -> case stack of
        Left z1 : Left z2 : stack' ->
            (pc', cs, Left (z1 + z2) : stack, mem)
        _ -> error "ADD: invalid stack for operation!"
    ...
    LABEL lab -> (pc', cs, stack, mem)
    JUMP lab -> case M.lookup lab ann of
        Nothing -> error "JUMP: invalid label!"
        Just instr ->
            let instrs = M.elems $ M.dropWhileAntitone (<= lab) ann
            in (lab, instr : instrs, stack, mem)
    JUMPFALSE lab -> case stack of
        Right b : stack' -> if b
            then (pc', cs, stack', mem)
            else case M.lookup lab ann of
                Nothing -> error "JUMPFALSE: invalid label!"
                Just instr ->
                    let instrs = M.elems $ M.dropWhileAntitone (<= la
b) ann
                    in (lab, instr : instrs, stack', mem)
        _ -> error "JUMPFALSE: invalid stack for operation"

    where
        pc' = pc + 1

```

```

In [ ]: -- A configuração inicial de um programa para AM2 precisa vir acompanhada
-- de
-- um Map com a associação entre cada instrução e o seu program counter,
-- porque no caso das instruções de salto em que é possível "regredir" no
-- programa, usar só uma lista para instruções não o permitirá.
initConfigAM2 :: State -> Stm -> (AM2Config, Env, AM2AnnotatedProgram)
initConfigAM2 initSt stm =
    let code :: AM2Code
        envSt :: EnvStateAM2
        (code, envSt) = whileToAM2 stm

        environ = getEnvSt envSt
        annotatedByteCode = getInstrs envSt--M.fromList $ zip (M.keys . g
etInstrs $ envSt) code

        memory :: Memory
        memory = M.fromList [(getEnv environ variable, getSt initSt varia
ble) | variable <- M.keys environ]
    in ((1, code, [], memory), environ, annotatedByteCode)

-- Dado um estado inicial e um comando da linguagem while, simula a sua e
xecução
-- na máquina abstrata AM2.
-- Devolve as variáveis usadas no programa, e os valores que estavam nas
respetivas
-- posições de memória aquando da terminação da execução.
-- Pode não terminar! (Halting problem).
runStmInAM2 :: State -> Stm -> M.Map Var Z

```

```

In [ ]: runStmInAM2 ex1State expProg
runStmInAM2 ex1State fact

```

```

fromList [("r",9),("x",3),("y",0)]

```

```

fromList [("f",720),("n",0)]

```