

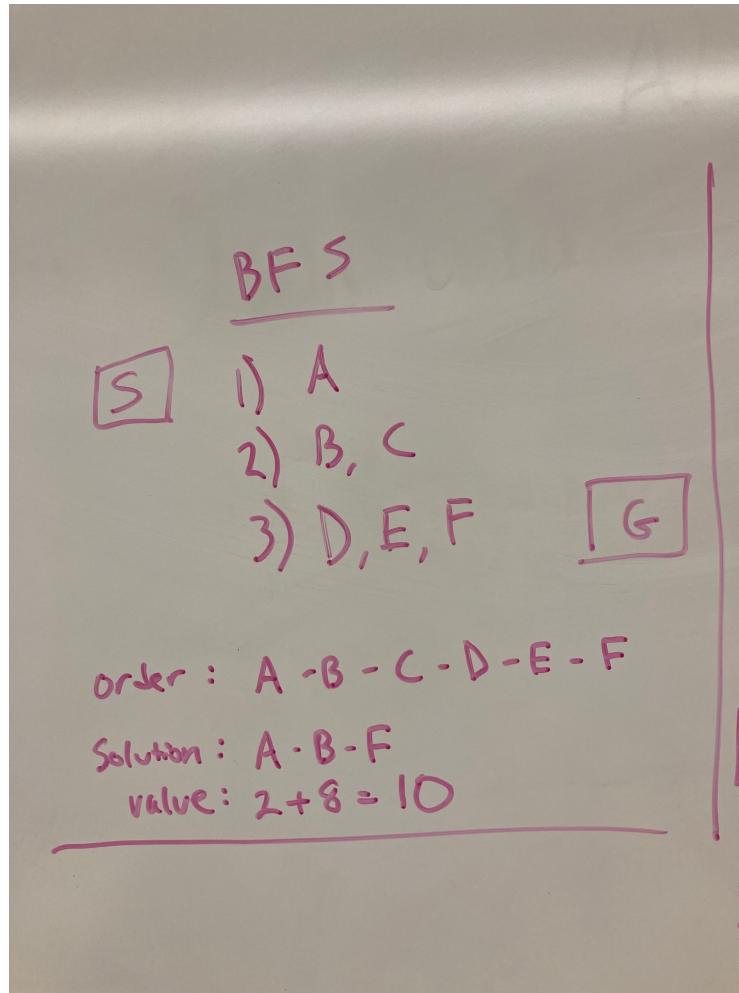
# Midterm AI 131

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# Q1

- a) False, if you hit a goal node and it is the same or better cost as other locations and no negative costs are allowed you can stop based on L $\rightarrow$ R convention/tie breaker preference. You must continue if a goal node cost is > frontier nodes.
- b) True, this won't yield a better solution and violates how A\* works. In Q2, once we hit N=9, we stop and don't look at (F, 10) at the front of the priority queue.
- c) True, for NW02 calculating all neighbors and intersections does work, but it is generally a lot worse for development if the user does not go off rules but rather manually writes out combos.
- d) False  $KB \models \alpha$   
If we don't know that  $KB$  entails  $\alpha$  then it is not complete. Completeness only happens when models of  $\alpha$  are entailed w/in  $KB$ .

# Q2 BFS



# Q2 A\* Star

A\* Star

[S]

- 1) A
- 2) (B, 2)
- 3) (C, 5)
- 4) (G, 7)
- 5) (H, 8)
- 6) (E, 9)
- 7) (I, 9)
- 8) (N, 9)

Order: A-B-C-G-H-E-I-N

Solution: A-C-I-N

Value: 9

$$g + h = f$$

$$B: 2 + 1 = 3$$

$$C: 1 + 4 = 5$$

$$\bar{g}_P + \bar{g} + \bar{h} = \bar{f}$$

$$D: 2 + (13+2) = 17$$

$$E: 2 + (2+5) = 9$$

$$* F: 2 + (8+0) = 10$$

B is low

PQ

[C, 5]

— — —

C is low

— — —

PQ

[ (E, 9), (F, 10), (D, 17) ]

$$G: 1 + (2+4) = 7 \quad G \text{ is low}$$

$$H: 1 + (4+3) = 8$$

$$I: 1 + (5+3) = 9$$

PQ

[ (H, 8), (E, 7), (I, 9), (F, 10), (D, 17) ]

$$L: \infty \quad H \text{ is low}$$

PQ

[ (E, 9), (I, 9), (F, 10), (D, 17), (L, \infty) ]

$$M: \infty \quad E \text{ is low}$$

[ (I, 9), (F, 10), (D, 17), (L, \infty), (M, \infty) ]

$$K: \infty \quad I \text{ is low}$$

[ (F, 10), (D, 17), (K, \infty), (L, \infty), (M, \infty) ]

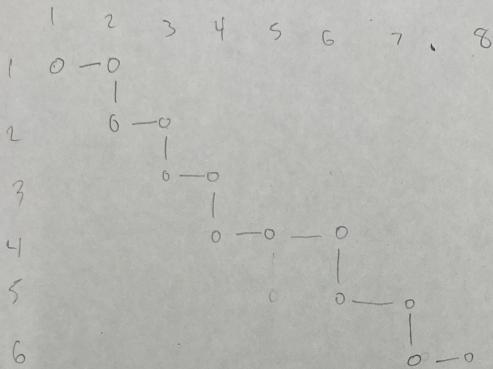
$$N: 6 + (3+0) = 9$$

9 < anything in PQ ... done

Q3

3)

$\leftarrow 8 \rightarrow$   
48 nodes  
 $\downarrow$



12

Totally out of time here

# Q4 CSP

Variables:

Timeslots {1...t}

Classes {1...c}

Classrooms {1...r}

Instructors {1...i}

$$\text{total variables} = t * c * r * i$$

Constraints

Given a specific timeslot and classroom, there may only be 1 instructor and 1 class

$$t_1, c_1, r_1, i_1 \neq t_1, r_1, c_{2..n}, i_{2..n}$$

$$t_1, c_1, r_2, i_2 \neq t_1, r_1, c_{1,3..n}, i_{1,3..n}$$

Keeping timeslot and classroom as controls, it is not possible to have any other instructor or any other class in the same classroom at the same time

Given a specific class and time, instructors can't match

$$@ (t_1, c_1) i_1 \neq i_{2..n}$$

Instructor 1 can't be any other instructor

Given a specific class and time, classrooms can't match

$$@ (t_1, c_1) r_1 \neq r_{2..n}$$

classroom 1 can't be any other classroom

At each timeslot we must have enough classrooms for all the classes taking place

$$@ t_{1..n} (r_{1..n} \geq c_{1..n})$$

r indicates # of classrooms

c indicates # of classes

There must be enough instructors at each timeslot to teach each class

$$@ t_{1..n} (i_{1..n} \geq c_{1..n})$$

i indicates # of instructors

c indicates # of classes

# Q5 Logic

- i.  $C \Rightarrow M$
  - ii.  $D \Rightarrow M$
  - iii.  $M \Rightarrow W$
  - iv.  $\neg(C \Rightarrow D) \vee C$

C = cat ...  
D = dog ...  
M = mammal ...  
W = warm ...