

Formalizing Modal Embeddings of Call-by-Name and Call-by-Value

Floris Reuvers

Supervisor: dr. N.M. van der Weide

Second Reader: dr. E.G.M. Hubbers

June 17, 2025

Radboud University

How can the unification of call-by-name and call-by-value evaluation strategies using modal logic be formalised in Agda?

Introduction

Let f be defined as

$$f(x) = x * x$$

Call-by-name

Let f be defined as

$$f(x) = x * x$$

Call-by-name evaluation of $f(3 + 3)$

$$f(3 + 3) \rightarrow (3 + 3) * (3 + 3)$$

Call-by-name

Let f be defined as

$$f(x) = x * x$$

Call-by-name evaluation of $f(3 + 3)$

$$\begin{aligned} f(3 + 3) &\rightarrow (3 + 3) * (3 + 3) \\ &\rightarrow 6 * (3 + 3) \end{aligned}$$

Call-by-name

Let f be defined as

$$f(x) = x * x$$

Call-by-name evaluation of $f(3 + 3)$

$$\begin{aligned} f(3 + 3) &\rightarrow (3 + 3) * (3 + 3) \\ &\rightarrow 6 * (3 + 3) \\ &\rightarrow 6 * 6 \end{aligned}$$

Call-by-name

Let f be defined as

$$f(x) = x * x$$

Call-by-name evaluation of $f(3 + 3)$

$$\begin{aligned} f(3 + 3) &\rightarrow (3 + 3) * (3 + 3) \\ &\rightarrow 6 * (3 + 3) \\ &\rightarrow 6 * 6 \\ &\rightarrow 36 \end{aligned}$$

Let f be defined as

$$f(x) = x * x$$

Let f be defined as

$$f(x) = x * x$$

Call-by-value evaluation of $f(3 + 3)$

$$f(3 + 3) \rightarrow f(6)$$

Let f be defined as

$$f(x) = x * x$$

Call-by-value evaluation of $f(3 + 3)$

$$\begin{aligned} f(3 + 3) &\rightarrow f(6) \\ &\rightarrow 6 * 6 \end{aligned}$$

Let f be defined as

$$f(x) = x * x$$

Call-by-value evaluation of $f(3 + 3)$

$$f(3 + 3) \rightarrow f(6)$$

$$\rightarrow 6 * 6$$

$$\rightarrow 36$$

Unifying cbn and cbv

- Why unify cbn and cbv?
- Some approaches to unification:
 - Modal logic
 - Linear logic
 - Thunks

Background

Grammar Lambda Calculus (λ -calculus)

$$M, N, P, Q ::= x \mid \lambda x. M \mid MN$$

Lorem ipsum

Lorem ipsum

Lorem ipsum

Lorem ipsum

Lorem ipsum

Lorem ipsum

Lorem ipsum

Formalisation in Agda

Lorem ipsum

Propositions

Lorem ipsum

Lorem ipsum

Conclusion

Lorem ipsum