

“Why is my heuristic behaving this way?!”

True story of a curious student

A good opportunity to do some revisions on A* and heuristics!

Updated on 8th August 2020

Found this message in my email inbox

*“Me and my assignment partner have been steadily working through the assignment and one night he decided to spend all night trying to get a heuristic working and **just off of a whiff when he returned the cost he multiplied it by 100.** After he did this each of the puzzle **times were halved** and we have no idea why,*

...

If you can, can you please explain to us why this is and if we are allowed to keep using it.”

Natural questions



- Assume h is an admissible heuristic
- If we use $5 \cdot h$ instead of h , will A^* find a goal?
- If we use $64 \cdot h$ instead of h , will A^* find a goal?

What is your gut feeling?

- How does the returned solution (if any) compares with an optimal solution in the worst case?



- Let $h' = 5 * h$ where h admissible

[5 is arbitrary, our reasoning would work with any other positive constant]

- Let $n' = \operatorname{argmin}_{x \text{ in Frontier}} g(x) + h'(x)$
- Let n be a node of the frontier on an optimal path between the root node and a goal node



Recall

- $h' = 5 * h$ where h admissible
- Let $n' = \underset{x \text{ in Frontier}}{\operatorname{argmin}} g(x) + h'(x)$
- Let n be a node of the frontier on an optimal path to a goal

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- $g(n') + h'(n') \leq g(n) + h'(n)$
 - $g(n) + h'(n) = g(n) + 5 * h(n)$
 - $g(n) + 5 * h(n) \leq 5 * g(n) + 5 * h(n)$
 - $5 * g(n) + 5 * h(n) \leq 5 * g(n) + 5 * h^*(n)$
 - Therefore, $g(n') + h'(n') \leq 5 * \text{optimal_cost}$

Can you justify each step?



Recall

- $h' = 5 * h$ where h admissible
- Let $n' = \underset{x \text{ in Frontier}}{\operatorname{argmin}} g(x) + h'(x)$
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- $g(n') + h'(n') \leq g(n) + h'(n)$ By definition of n'
 - $g(n) + h'(n) = g(n) + 5 * h(n)$ By definition of h'
 - $g(n) + 5 * h(n) \leq 5 * g(n) + 5 * h(n)$ $1 \leq 5$
 - $5 * g(n) + 5 * h(n) \leq 5 * g(n) + 5 * h^*(n)$ h is admissible
 - $g(n') + h'(n') \leq 5 * \text{optimal_cost}$ n is on an optimal path

So what?!

- If we use $g+5*h$ instead of $g+h$ to pop a node from the frontier the path returned will cost at most

$$5 * \text{optimal_cost}$$

- A* with $g+5*h$ is behaving more like a greedy search and might find a goal more quickly at the cost of a loss of optimality guarantee

Conclusion

- This scenario could happen in any company
 - A software engineer modifies an existing algorithm to tackle a problem
 - Before adopting the modification, any responsible person will want to know about
 - impact on convergence guarantee
 - Impact on performance guarantee
- Software engineers and computer scientists have to wear a mathematician hat sometimes!