Lista de exercícios 01 - PCC116

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1 Item 1.5.7 - Prove os seguintes sequentes usando dedução natural. Tente demonstrá-los sem utilizar a regra RAA.

(a)
$$\{(P \land Q) \land R, \ S \land T\} \vdash Q \land S$$

$$\frac{\overline{(P \land Q) \land R}}{P \land Q} \{ \land_{EE} \} } \{ \land_{EE} \}$$

$$\frac{\overline{S \land T}}{Q} \{ \land_{EE} \} }{Q \land S} \{ \land_{I} \}$$

(b)
$$\{(P \wedge Q) \wedge R\} \vdash (P \wedge R) \vee Z$$

$$\frac{\overline{(P \wedge Q) \wedge R}}{P \wedge Q} \xrightarrow{\{\wedge_{EE}\}} \frac{\overline{(P \wedge Q) \wedge R}}{R} \xrightarrow{\{\wedge_{ED}\}}$$

$$\frac{P \wedge R^{1}}{(P \wedge R) \vee Z} \xrightarrow{\{\vee_{IE}\}^{1}}$$

(d)
$$\vdash (P \land Q) \rightarrow P$$

$$\frac{\overline{(P \land Q)^{1}} }{P} {}_{\{ \land_{EE} \}}$$

$$\frac{P}{(P \land Q) \rightarrow P} {}_{\{ \rightarrow_{I} \}^{1}}$$

(e)
$$\{P\} \vdash Q \to (P \land Q)$$

$$\frac{\overline{P} \ \{_{ID}\} \quad \overline{Q^1} \ \{_{ID}\}}{P \land Q \quad \{_{ID}\}}$$

$$\frac{P \land Q}{Q \to (P \land Q)} \ \{_{\to I}\}^1$$
 (f)

(f)
$$\{P\} \vdash (P \to Q) \to Q$$

$$\frac{P \to Q^1 \quad {}_{\{ID\}} \quad P \quad {}_{\{ID\}}}{Q}$$

$$\frac{Q}{(P \to Q) \to Q} \quad {}_{\{\to I\}^1}$$

(g)
$$\vdash (P \land Q) \rightarrow (P \lor Q)$$

$$\frac{\overline{P \land Q^{1}}}{P \lor Q} {}^{\{I_{D}\}} \qquad \overline{P \land Q^{1}}}_{\{V_{ID}\}} {}^{\{I_{D}\}}$$

$$\frac{P \land Q^{1}}{P \lor Q} {}^{\{\lor_{ID}\}}$$

$$(P \land Q) \rightarrow (P \lor Q)$$

(h)
$$\{Q \rightarrow (P \rightarrow R), \ \neg R, \ Q\} \vdash \neg P$$

$$\frac{\overline{Q \to (P \to R)} \quad {}^{\{I_D\}} \quad \overline{Q} \quad {}^{\{I_D\}}}{P \to R} \quad {}^{\{I_D\}} \quad {}^{\{I_$$

(i)
$$\{P\} \vdash Q \to (P \land Q)$$

$$\frac{P \quad Q^1}{P \land Q} \{ \land_1 \}$$

$$Q \to (P \land Q) \quad \{ \rightarrow_I \}^1$$

(j)
$$\{(P \to R) \land (Q \to R), \ P \land Q\} \vdash Q \land R$$

$$\frac{\overline{(P \to R) \land (Q \to R)}}{Q \to R} \xrightarrow{\{\land_{ED}\}} \frac{\overline{P \land Q}}{Q} \xrightarrow{\{\land_{ED}\}} \frac{P \land Q}{Q} \xrightarrow{\{\land_{ED}\}} \frac{P \land Q}{Q} \xrightarrow{\{\land_{ED}\}}$$

$$\frac{R}{Q \land R} \xrightarrow{Q \land R} \xrightarrow{\{\land_{ED}\}} \frac{P \land Q}{Q} \xrightarrow{\{\land_{ED}\}}$$

$$\{P \to Q \to R, \ P \to Q\} \vdash P \to R$$

$$\frac{P \to Q \to R}{P \to Q \to R} \xrightarrow{\{ID\}} \frac{P^1}{P^1} \xrightarrow{\{ID\}} \frac{P \to Q}{P \to Q} \xrightarrow{\{ID\}} \frac{P^1}{P^1} \xrightarrow{\{ID\}} \frac{P \to Q}{P \to R} \xrightarrow{\{D\}^1} \frac{P \to Q}{P \to R}$$

(1)
$$\{P \to Q, R \to S\} \vdash (P \lor R) \to (Q \lor S)$$

$$\frac{\overline{P \to Q}}{Q \lor S} {\{\iota_D\}} \frac{\overline{R \to S}}{Q \lor S} {\{\iota_D\}}$$

$$\frac{Q \lor S}{(P \lor R) \to (Q \lor S)} {\{\to_I\}}$$

(m)
$$\{Q \to R\} \vdash (P \to Q) \to (P \to R)$$

$$\frac{\overline{Q \to R}}{Q \to R} \xrightarrow{\{ID\}} \frac{\overline{P \to Q^1}}{Q} \xrightarrow{\{\to_E\}} \frac{R}{P \to R} \xrightarrow{\{\to_I\}} \frac{R}{(P \to Q) \to (P \to R)} \xrightarrow{\{\to_I\}^1}$$

(n)
$$\{(P \wedge Q) \vee (P \wedge R)\} \vdash P \wedge (Q \vee R)$$

$$\frac{\frac{(P \land Q) \lor (P \land R)}{P \land Q} \{_{\land_{IE}}\}}{\frac{P \land Q}{P} \{_{\land_{EE}}\}} \{_{\land_{IE}}\}} \frac{\overline{(P \land Q) \lor (P \land R)}}{\frac{P \land Q}{Q} \{_{\land_{EE}}\}} \{_{\land_{EE}}\}} \{_{\land_{EE}}\}}{\frac{P \land Q}{Q \lor R} \{_{\land_{IE}}\}} \{_{\land_{IE}}\}}$$

(o)
$$\{P \to (Q \land R)\} \vdash (P \to Q) \land (P \to R)$$

$$\frac{P \to (Q \land R) \xrightarrow{\{ID\}} \overline{P^1} \xrightarrow{\{ID\}}}{\{P \to Q \land R\}} \xrightarrow{\{ID\}} \frac{P \to (Q \land R) \xrightarrow{\{ID\}} \overline{P^1} \xrightarrow{\{ID\}}}{\{P \to Q \land R\}} \xrightarrow{\{A \to E\}^1} \frac{Q \land R}{R} \xrightarrow{\{A \to E\}^1} \xrightarrow{\{P \to Q\} \land \{P \to Q\} \land (P \to R)} \xrightarrow{\{A \to E\}^1} \xrightarrow{\{A \to$$

(p)
$$\{(P \to Q) \land (P \to R)\} \vdash P \to (Q \land R)$$

$$\frac{ P \rightarrow Q) \land (P \rightarrow R) }{P \rightarrow Q} {\{ \land_{EE} \}} \qquad \frac{ }{P_1} {\{ \downarrow_{ID} \}} \qquad \frac{ (P \rightarrow Q) \land (P \rightarrow R) }{P \rightarrow R} {\{ \land_{ED} \}} \qquad \frac{ }{P_1} {\{ \downarrow_{ID} \}} \qquad \frac{ P \rightarrow R \qquad \{ \land_{ED} \}}{P} \qquad \frac{ \{ \downarrow_{ID} \}}{\{ \rightarrow_{E} \}} \qquad \frac{ Q \land R }{P \rightarrow (Q \land R)} {\{ \rightarrow_{I} \}^1} \qquad \frac{ }{P \rightarrow R} \qquad \frac{ \{ \downarrow_{ID} \}}{P} \qquad \frac{ \{ \downarrow_{ID} \}}{\{ \rightarrow_{E} \}} \qquad \frac{ \{ \downarrow_{ID} \}}{P} \qquad \frac{ \{ \downarrow_{ID} \}}{\{ \rightarrow_{E} \}} \qquad \frac{ \{ \downarrow_{ID} \}}{P} \qquad \frac{ \{ \downarrow_{ID} \}}{\{ \rightarrow_{E} \}} \qquad \frac{ \{ \downarrow_{ID} \}}{P} \qquad \frac{ \{ \downarrow_{ID} \}}{\{ \rightarrow_{E} \}} \qquad \frac{ \{ \downarrow_{ID} \}}{P} \qquad \frac{ \{ \downarrow_{ID} \}}{\{ \rightarrow_{E} \}} \qquad \frac{ \{ \downarrow_{ID} \}}{P} \qquad \frac{ \{ \downarrow_{ID} \}}{\{ \rightarrow_{E} \}} \qquad \frac{ \{ \downarrow_{ID} \}}{\{ \downarrow_{ID} \}} \qquad \frac{ \{ \downarrow_{I$$

(q)
$$\{P \to Q\} \vdash ((P \land Q) \to P) \land (P \to (P \land Q))$$

$$\frac{P^{1}}{P^{1}} \xrightarrow{\{ID\}} \frac{P \to Q}{Q} \xrightarrow{\{A_{E}\}} \frac{P \to Q}{Q} \xrightarrow{\{A_{E}\}} \frac{P \land Q}{P \land Q} \xrightarrow{\{A_{E}\}} \frac{P \land Q}{P \to (P \land Q)} \xrightarrow{\{A_{E}\}} \frac{P \land Q}{P \to (P \land Q)} \xrightarrow{\{A_{E}\}} \frac{P \land Q}{P \to (P \land Q)} \xrightarrow{\{A_{E}\}} \frac{P \to Q}{P \to Q} \xrightarrow{\{A_{E}\}} \xrightarrow{\{A_{E}\}} \frac{P \to Q}{P \to Q} \xrightarrow{\{A_{E}\}} \xrightarrow{\{A_{E}\}} \frac{P \to Q}{P \to Q} \xrightarrow{\{A_{E}\}} \xrightarrow{\{A_{E}$$

(r)
$$\{P \to (Q \lor R), \ Q \to S, \ R \to S\} \vdash P \to S$$

$$\frac{\overline{P} \ \{_{ID}\} \quad \overline{P \to (Q \lor R)} \quad \{_{ID}\} }{Q \lor R} \quad \{_{ID}\} \quad Q \to S \quad R \to S \quad \{_{VE}\}$$

$$\frac{S}{P \to S} \quad \{_{ID}\} \quad \{_{ID}\}$$

(s)
$$\vdash \neg P \to P \to P \to Q$$

$$\frac{ \overline{\neg P^1} }_{\{ID\}} \frac{\{ID\}}{P^2} \underbrace{\{ID\}}_{\{\to E\}}$$

$$\frac{ \underline{\bot}_{\{CTR\}}}{P \to Q} \underbrace{\{\neg I\}}_{\{\to I\}^2}$$

$$\frac{ P \to Q}{P \to (P \to Q)} \underbrace{\{\neg I\}^2}_{\{\to I\}^1}$$

$$\frac{ \neg P \to (P \to Q)}{\neg P \to (P \to Q))} \underbrace{\{\neg I\}^1}_{\{\to I\}^1}$$

(t)
$$\{P \land Q \to R, \ R \to S, \ Q \land \neg S\} \vdash \neg P$$

$$\frac{Q \land \neg S}{\neg S} \stackrel{\{_{ID}\}}{\land ED\}} = \frac{P \land Q \rightarrow R}{P \land Q \rightarrow R} \stackrel{\{_{ID}\}}{\land ED} = \frac{Q \land \neg S}{Q} \stackrel{\{_{ID}\}}{\land EE}}{\land EB}$$

$$\frac{P \land Q \rightarrow R}{P \land Q} \stackrel{\{_{ID}\}}{\land EB} = \frac{P \land Q}{P \land Q} \stackrel{\{_{ID}\}}{\land EB}}{\land EB}$$

$$\frac{\bot}{\neg P} \{CTR\}^{1}$$

(u)
$$\{(P \to Q) \to R, S \to \neg P, T, \neg S \land T \to Q\} \vdash R$$

$$\frac{\frac{\neg S^{1}}{\neg P \to \neg S}}{\neg P \to \neg S}} \xrightarrow{\{\vdash I\}} \frac{\frac{\neg S \to \neg P}{\neg P \to E\}}}{\neg P \to E}} \xrightarrow{T} \xrightarrow{\{\vdash D\}} \frac{\neg S \land T \to Q}{\neg P}} \xrightarrow{\{\vdash D\}} \frac{\neg S \land T}{\neg P} \xrightarrow{\{\vdash D\}}} \frac{\neg S \land T}{\neg P} \xrightarrow{\{\vdash D\}}} \xrightarrow{T} \xrightarrow{\{\vdash D\}} \frac{\neg S \land T}{\neg S \land T}} \xrightarrow{\{\vdash D\}} \xrightarrow{\neg S \land T} \xrightarrow{\{\vdash D\}}} \frac{\neg S \land T}{\neg S \land T} \xrightarrow{\{\vdash D\}}} \xrightarrow{R} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{\vdash D\}}} \xrightarrow{\{\vdash D\}} \xrightarrow{\{$$

2 Item 2.9 - Prove os seguintes sequentes usando dedução natural:

(a)
$$\{\forall x. (P(x) \to Q(x))\} \vdash (\forall x. \neg Q(x)) \to (\forall x. \neg P(x))$$

$$\frac{\overline{\forall x. \neg Q(x)}}{\overline{\neg Q(a)}} \stackrel{\{_{ID}\}}{\{\forall E\}} \frac{}{P(a)} \stackrel{\{_{ID}\}}{\{\to E\}}$$

$$\frac{\underline{\bot}}{\overline{\neg P(a)}} \stackrel{\{\neg I\}}{\{\forall I\}}$$

$$\frac{\overline{\forall x. \neg P(x)}}{\{\forall X. \neg P(x)\}} \stackrel{\{\neg I\}}{\{\to I\}}$$

$$\begin{array}{c} \textbf{(b)} \\ \{ \forall x. (P(x) \rightarrow \neg Q(x)) \} \vdash \neg (\exists x. (P(x) \land Q(x))) \} \\ \hline \\ \frac{P(t) \land Q(t)}{P(t) \land Q(t)} \{ \land E \} & \frac{\forall x. (P(x) \rightarrow \neg Q(x))}{P(t) \rightarrow \neg Q(t)} \{ \land E \} & \frac{P(t) \land Q(t)}{Q(t)} \{ \land E \} \\ \hline \\ \frac{P(t)}{-Q(t)} & \frac{\bot}{-Ax. (P(x) \land Q(x))} \{ \land E \} & \frac{P(t) \land Q(t)}{Q(t)} \{ \land E \} \\ \hline \\ \frac{\bot}{-Ax. (P(x) \land Q(x))} \{ \land E \} & \frac{\bot}{-Ax} (P(x) \land Q(x)) \\ \hline \\ \textbf{(c)} \\ \hline \\ \frac{\forall x. (A(x) \rightarrow (B(x) \lor C(x))), \ \forall x. \neg B(x) \} \vdash (\forall x. A(x)) \rightarrow (\forall x. C(x)) \\ \hline \\ \frac{A(t)}{A(t)} & \forall E & \frac{\bot}{A(t) \rightarrow (B(t) \lor C(t))} \rightarrow E & \frac{\bot}{-B(t)} & \frac{\bot}{B(t)} & \frac{\bot}{B(t)} \\ \hline \\ \frac{B(t) \lor C(t)}{A(t)} & \frac{\bot}{A(t) \rightarrow (B(t) \lor C(t))} \rightarrow E & \frac{\bot}{-B(t)} & \frac{\bot}{B(t)} & \frac{\bot}{B(t)} \\ \hline \\ \textbf{(d)} \\ \hline \\ \frac{A(t)}{A(t)} & \frac{\bot}{A(t) \rightarrow (B(t) \lor C(t))} \rightarrow E & \frac{\bot}{-B(t)} & \frac{\bot}{B(t)} & \frac{\bot}{B(t)} \\ \hline \\ \textbf{(d)} \\ \hline \\ \frac{A(t)}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t) \rightarrow (B(t) \lor C(t))} \rightarrow E & \frac{\bot}{-B(t)} & \frac{\bot}{B(t)} & \frac{\bot}{B(t)} & \frac{\bot}{B(t)} \\ \hline \\ \textbf{(d)} \\ \hline \\ \frac{A(t)}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} \\ \hline \\ \textbf{(d)} \\ \hline \\ \frac{A(t)}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} & \frac{\bot}{A(t)} \\ \hline \\ \textbf{(d)} & \frac{\bot}{A(t)} \\ \hline \\ \textbf{(d)} & \frac{\bot}{A(t)} &$$

(g)
$$\{\exists x. \neg P(x)\} \vdash \neg \forall x. P(x)$$

$$\frac{ \frac{}{\neg P(t)} \text{ [hipótese]} \quad \frac{\forall x. P(x)}{P(t)} \text{ $\forall E$}}{ \frac{\bot}{\neg \forall x. P(x)}} \text{ $\exists E$}$$

3 Item 5.2.1 - Prove os seguintes teoremas

(a). Para todo $n \in \mathbb{N}$,

$$\sum_{i=0}^{n} 3^{i} = \frac{3^{n+1} - 1}{2}$$

Base da indução: Para n = 0, temos:

$$\sum_{i=0}^{0} 3^i = 3^0 = 1$$

е

$$\frac{3^{0+1}-1}{2} = \frac{3-1}{2} = \frac{2}{2} = 1$$

Logo, a base da indução é válida.

Passo indutivo: Suponha que a fórmula seja válida para um certo $n \in \mathbb{N}$, ou seja,

$$\sum_{i=0}^{n} 3^i = \frac{3^{n+1} - 1}{2}$$

Vamos provar que:

$$\sum_{i=0}^{n+1} 3^i = \frac{3^{n+2} - 1}{2}$$

De fato:

$$\sum_{i=0}^{n+1} 3^i = \left(\sum_{i=0}^n 3^i\right) + 3^{n+1}$$

$$= \frac{3^{n+1} - 1}{2} + 3^{n+1}$$

$$= \frac{3^{n+1} - 1 + 2 \cdot 3^{n+1}}{2}$$

$$= \frac{3 \cdot 3^{n+1} - 1}{2}$$

$$= \frac{3^{n+2} - 1}{2}$$

Portanto, a fórmula vale para n+1. Pelo princípio da indução matemática, a identidade é válida para todo $n \in \mathbb{N}$.

(b). Para todo $n \in \mathbb{N}$,

$$\sum_{i=0}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

Base: Para n = 0,

$$\sum_{i=0}^{0} i^2 = 0, \quad e \quad \frac{0(0+1)(2\cdot 0+1)}{6} = 0$$

Passo indutivo: Suponha a fórmula válida para n,

$$\sum_{i=0}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

Prove para n+1:

$$\sum_{i=0}^{n+1} i^2 = \sum_{i=0}^n i^2 + (n+1)^2 = \frac{n(n+1)(2n+1)}{6} + (n+1)^2$$

$$= \frac{n(n+1)(2n+1) + 6(n+1)^2}{6} = \frac{(n+1)\left[n(2n+1) + 6(n+1)\right]}{6}$$

$$= \frac{(n+1)(2n^2 + 7n + 6)}{6} = \frac{(n+1)(n+2)(2n+3)}{6}$$

$$\Rightarrow \sum_{i=0}^{n+1} i^2 = \frac{(n+1)(n+2)(2n+3)}{6}$$

Logo, a fórmula vale para todo $n \in \mathbb{N}$.

(c). Para todo $n \in \mathbb{N}$,

$$\sum_{i=1}^{n} (8i - 5) = 4n^2 - n$$

Base: Para n = 1,

$$8(1) - 5 = 3$$
, $4(1)^2 - 1 = 3$

Passo indutivo: Suponha que

$$\sum_{i=1}^{n} (8i - 5) = 4n^2 - n$$

então

$$\sum_{i=1}^{n+1} (8i - 5) = \sum_{i=1}^{n} (8i - 5) + (8(n+1) - 5)$$
$$= 4n^{2} - n + 8n + 3 = 4n^{2} + 7n + 3 = 4(n+1)^{2} - (n+1)$$

(d). Para todo $n \geq 1$,

$$5 \mid (n^5 - n)$$

Base: Para n = 1, temos $1^5 - 1 = 0$, e $5 \mid 0$

Passo indutivo: Suponha $5 \mid n^5 - n$, ou seja, $n^5 - n = 5k$

Verificar para n+1:

$$(n+1)^5 - (n+1) = n^5 + 5n^4 + 10n^3 + 10n^2 + 5n + 1 - n - 1 = (n^5 - n) + 5(n^4 + 2n^3 + 2n^2 + n)$$

Como 5 | $n^5 - n$ e o restante é múltiplo de 5, segue que 5 | $(n+1)^5 - (n+1)$

(e). Para todo $n \geq 1$,

$$6 \mid (7^n - 1)$$

Provar que 6 | $(7^n - 1)$ para todo $n \ge 1$

"'latex **Base:** Para $n = 1, 7^1 - 1 = 6 \Rightarrow 6 \mid 6$

Passo indutivo: Suponha 6 | $(7^n - 1) \Rightarrow 7^n \equiv 1 \mod 6$

Multiplicando ambos os lados por 7:

$$7^{n+1} = 7 \cdot 7^n \equiv 7 \cdot 1 = 7 \equiv 1 \mod 6 \Rightarrow 6 \mid (7^{n+1} - 1)$$

(f). Para todo
$$n \ge 1$$
,

$$6 \mid (n^3 + 5n)$$

Provar que $6 \mid (n^3 + 5n)$

"'latex Base: Para $n = 1, 1 + 5 = 6 \Rightarrow 6 \mid 6$

Passo indutivo: Suponha $6 \mid (n^3 + 5n)$

$$(n+1)^3 + 5(n+1) = n^3 + 3n^2 + 3n + 1 + 5n + 5 = (n^3 + 5n) + 3n^2 + 3n + 6$$

Como 6 | (n^3+5n) e os demais termos somam múltiplo de 6, temos 6 | $(n+1)^3+5(n+1)$

(g). Para todo $n \in \mathbb{N}$,

$$2 | (n^2 + n)$$

Provar que $2 \mid (n^2 + n)$

"'latex Base: Para $n=0,\,0^2+0=0\Rightarrow 2\mid 0$

Passo indutivo: Suponha $2 \mid n^2 + n$

$$(n+1)^2 + (n+1) = n^2 + 2n + 1 + n + 1 = (n^2 + n) + 2n + 2$$

Como 2 | (n^2+n) e 2n+2é par, então 2 | $(n+1)^2+(n+1)$