

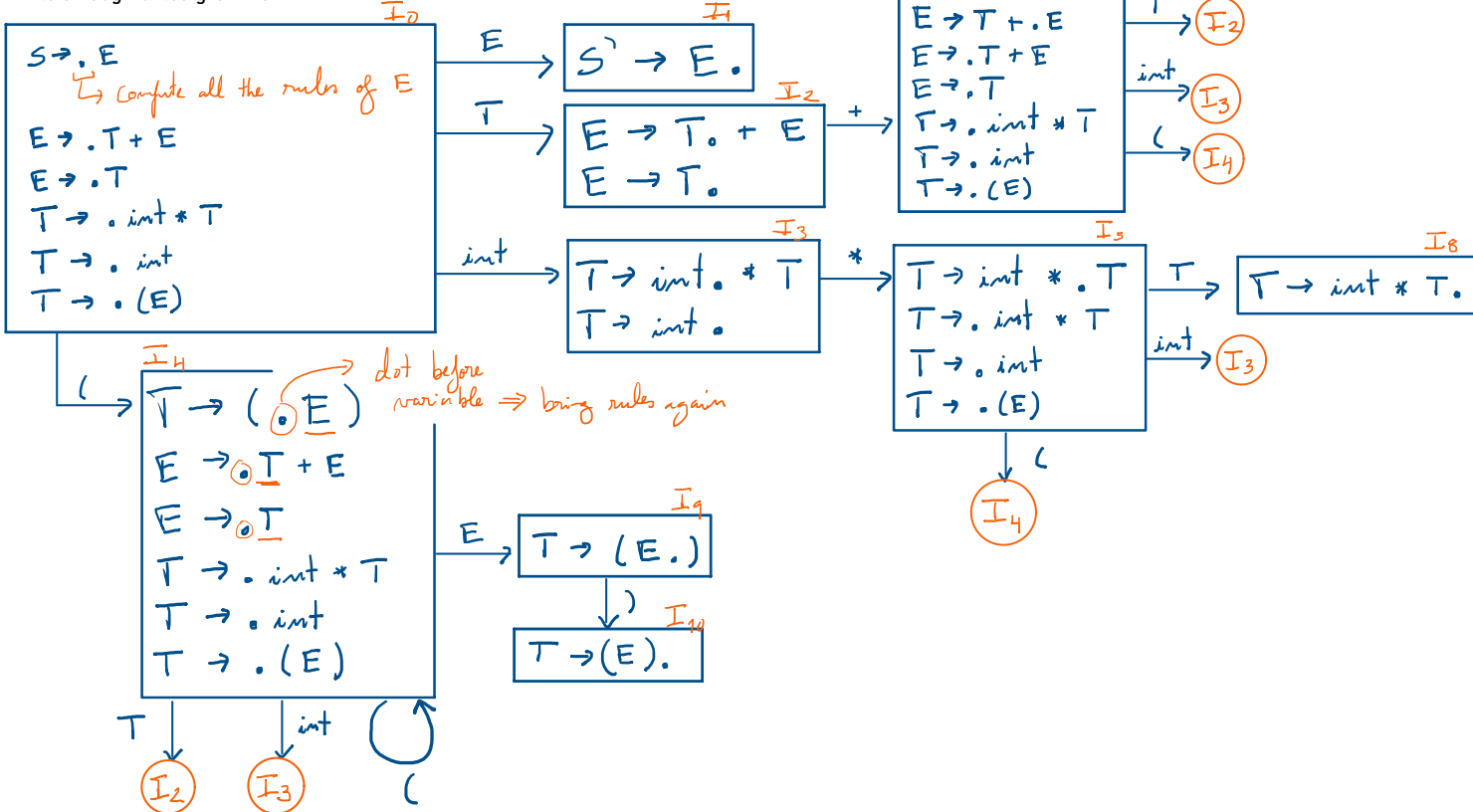
LR(0) & SLR(1) Automaton & Parsing Table

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1: $E \rightarrow T + E$
2: $E \rightarrow T$
3: $T \rightarrow \text{int} * T$
4: $T \rightarrow \text{int}$
5: $T \rightarrow (E)$

Não é obrigatório remover recursividade à esquerda (mas é melhor).

Write an augmented grammar:



SLR(1) Table

-	int	()	+	*	\$	E	T
0	S3	S4					1	2
1						AC		
2			R2	S6		R2		
3			R4	R4	S5	R4		
4	S3	S4					9	2
5	S3	S4						8
6	S3	S4					7	2
7			R1			R1		
8			R3	R3		R3		
9			S10					
10			R5	R5		R5		

$$\text{Follow}(T) = \{ +,), \$ \}$$

$$E \rightarrow T \cdot$$

= Reducing to 2

Follow(T) for "int" is terminal

$$\{ + \} \cup \text{Follow}(E) \cup \text{Follow}(T)$$

1: $E \rightarrow T + E$
2: $E \rightarrow T$
3: $T \rightarrow \text{int} * T$
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$$\Rightarrow \{ + \} \cup \{ \$,) \}$$

S \rightarrow Shift

R \rightarrow Reduce -- used in accepting states \Rightarrow shift cannot be used \Rightarrow use the follow

ex.: $\text{Follow}(E) = \{ \$,) \}$

AC: (accept) -- stops the parsing

LR(0) Table -- the reduce is in all the terminal symbols (lame)

-	int	()	+	*	\$	E	T
0	S3	S4					1	2
1						AC		
2	R2	R2	R2	R2/S6	R2	R2		
3	R4	R4	R4	R4	R4/S5	R4		
4	S3	S4					9	2
5	S3	S4						8
6	S3	S4					7	2
7	R1	R1	R1	R1	R1	R1		

Shift/Reduce Conflict!

8	R3	R3	R3	R3	R3	R3		
9			S10					
10	R5	R5	R5	R5	R5	R5		