

Introduction to Named Data Networking

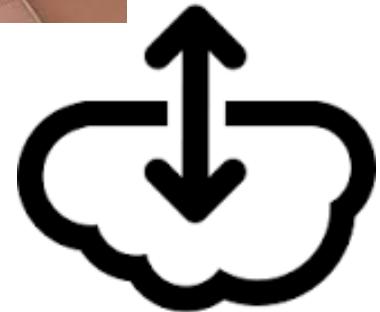
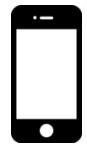
FYI, not HW or exam material

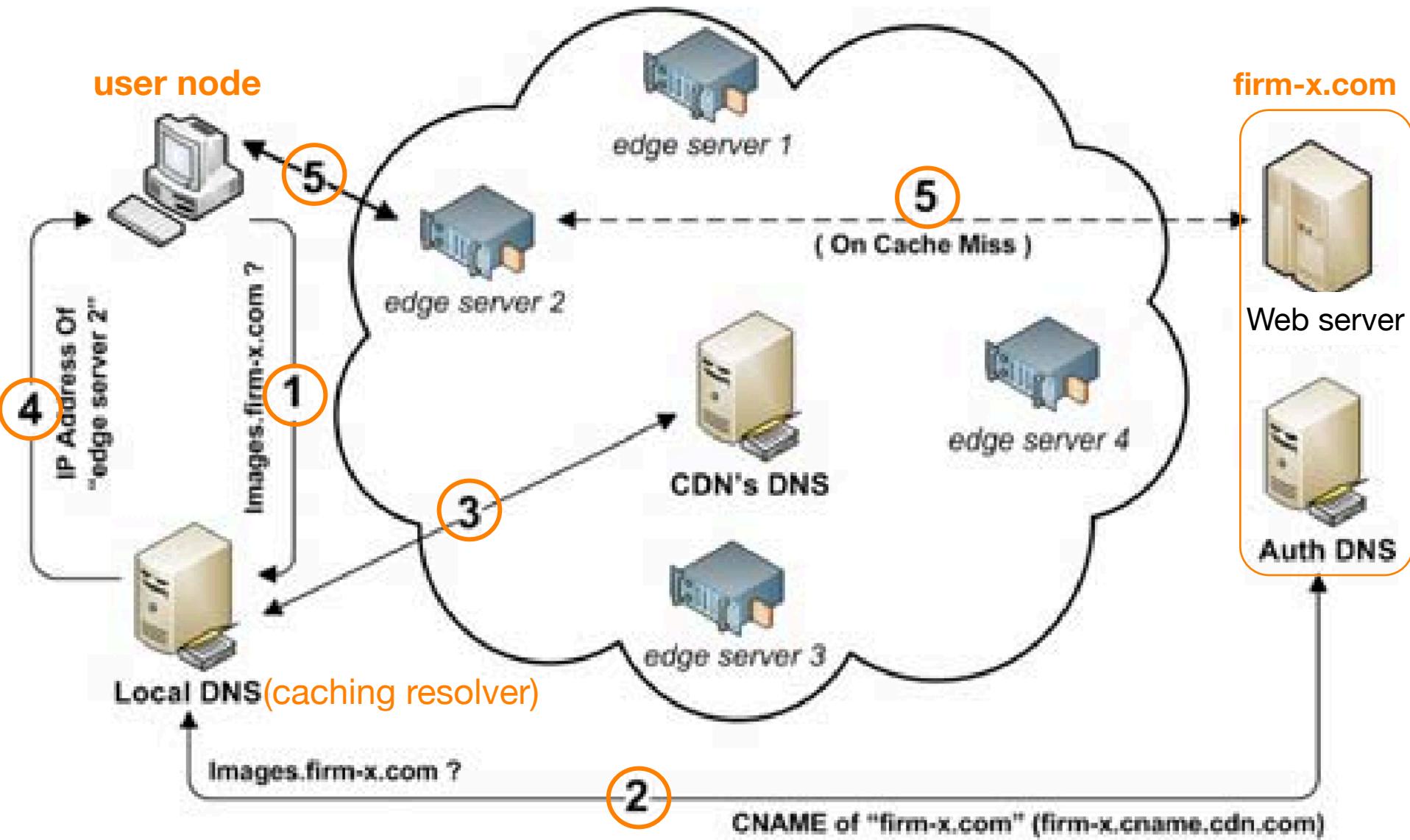
Today's Plan

- ◆ Why Internet even needs a new architecture
- ◆ Assuming it is needed,
 - How to figure out how this new architecture should look like?
 - When would Internet move to this new architecture?

A simple example

- ◆ Jon starts the morning at Starbucks, music streaming
 - TCP connection between his phone and server, thru WiFi
- ◆ When Jon walks out to the train station: WiFi fades → move Internet connectivity to cellular
 - TCP connection: interface specific
 - Re-establish TCP connection
- ◆ Gets on the train
 - Train offers free WiFi
 - the phone has to re-establish TCP connection (again)
 - Others listen to the same music, can they share the download?





A number of solutions developed over time

- ◆ To be able to use multiple interfaces on a phone: multipath-TCP developed
- ◆ To handle server load: CDN (content distribution network) developed
- ◆ To make CDN work: DNS hack developed
 - the phone's DNS lookup for itunes.apple.com
- ◆ Now DNS servers get new work to do: Make sure to pick a good IP address for a CDN box
 - Working?
 - Closest to the requester?
- ◆ Have not mentioned the complexity at data centers yet
- ◆ Have not mentioned **security** yet

```
[Lxia:~/papers/iotdi-2016] lixia% dig itunes.apple.com
; <>> DiG 9.8.3-P1 <>> itunes.apple.com
;; global options: +cmd
;; Got answer:
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 8328
;; flags: qr rd ra; QUERY: 1, ANSWER: 4, AUTHORITY: 8, ADDITIONAL: 8

;; QUESTION SECTION:
;itunes.apple.com.           IN  A

;; ANSWER SECTION:
itunes.apple.com.          2157  IN  CNAME  itunes-cdn.itunes-apple.com.akadns.net.
itunes-cdn.itunes-apple.com.akadns.net. 252  IN  CNAME  itunes.apple.com.edgekey.net.
itunes.apple.com.edgekey.net.        289  IN  CNAME  e673.e9.akamaiedge.net.
e673.e9.akamaiedge.net.          9  IN  A       104.68.98.172

;; AUTHORITY SECTION:
e9.akamaiedge.net.      2557  IN  NS    n5e9.akamaiedge.net.
e9.akamaiedge.net.      2557  IN  NS    a0e9.akamaiedge.net.
e9.akamaiedge.net.      2557  IN  NS    n3e9.akamaiedge.net.
e9.akamaiedge.net.      2557  IN  NS    n0e9.akamaiedge.net.
e9.akamaiedge.net.      2557  IN  NS    n1e9.akamaiedge.net.
e9.akamaiedge.net.      2557  IN  NS    n2e9.akamaiedge.net.
e9.akamaiedge.net.      2557  IN  NS    n4e9.akamaiedge.net.
e9.akamaiedge.net.      2557  IN  NS    a1e9.akamaiedge.net.

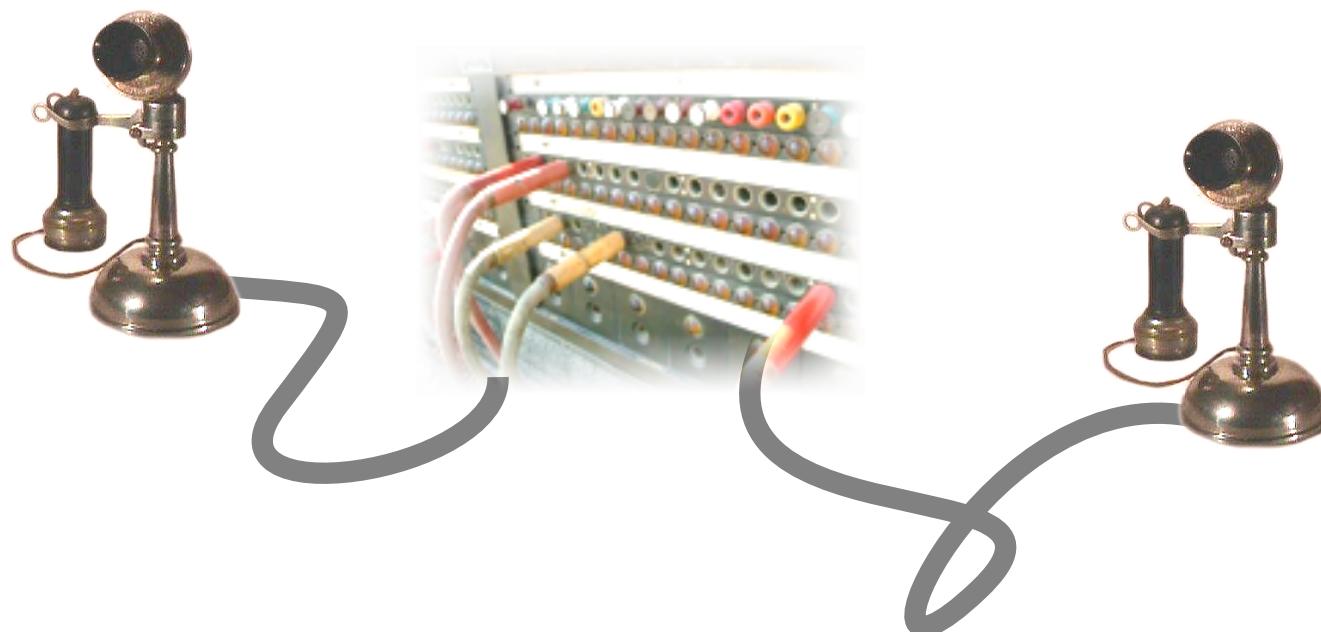
;; ADDITIONAL SECTION:
a0e9.akamaiedge.net.   4557  IN  AAAA 2600:1406:1c:f000:9029:a648:25fb:2656
a1e9.akamaiedge.net.   2557  IN  AAAA 2600:1406:1c:f000:902c:a648:25fb:2656
n0e9.akamaiedge.net.   2557  IN  A     63.141.195.111
n1e9.akamaiedge.net.   4557  IN  A     184.51.100.13
```

- ◆ Internet is facing (lots) problems today
- ◆ The root cause?

“The farther backward you can look, the farther forward you are likely to see.”

— Winston Churchill

Telephone Network as the 1st Communication system



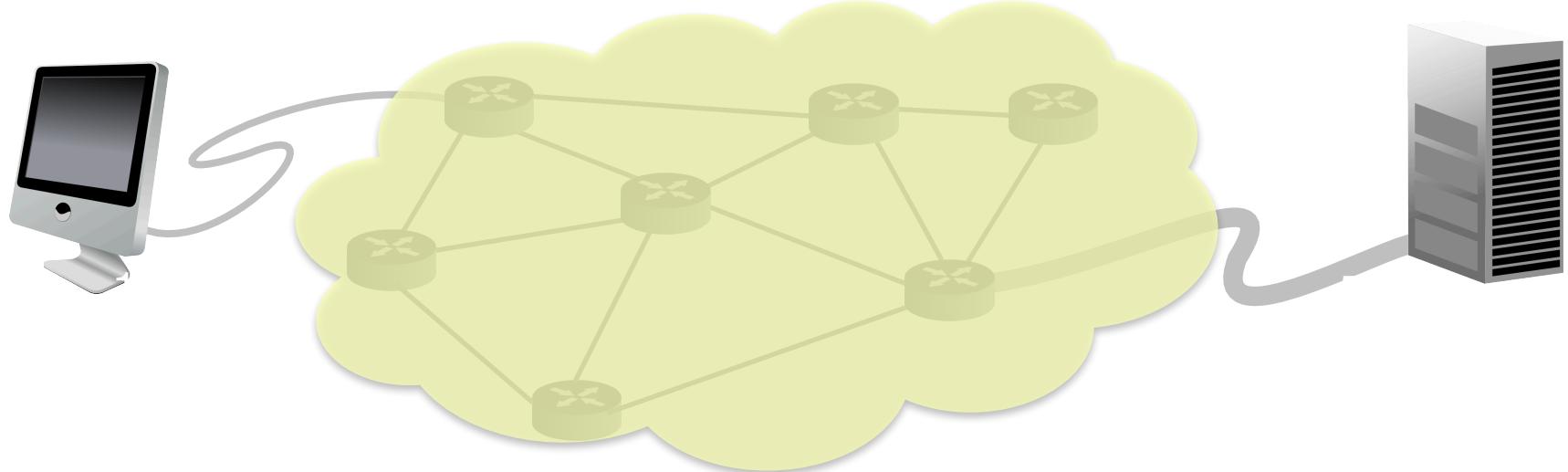
Focused on building the *wires* between 2 ends

The change over time

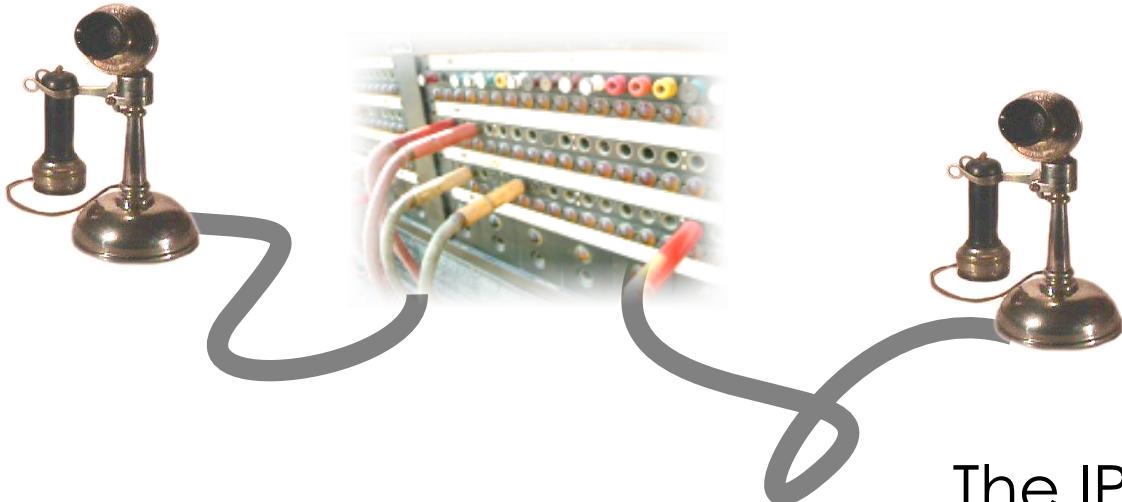
- ◆ Telephone networks last a long time
 - until computers were invented
- ◆ “On Distributed Communications Networks”
 - by Paul Baran, IEEE Transactions on Communications, March, 1964.
- ◆ Why a new design: Technology advanced, goals changed
 - need robust communication system for "second strike capability"
 - providing common services for a wide range of users & applications

IP changed the way of building communication systems

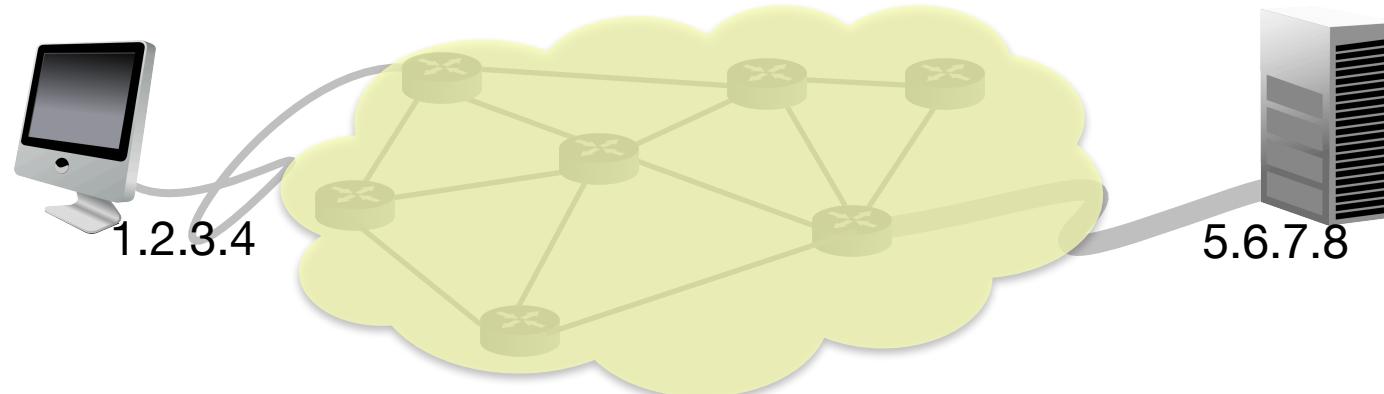
RFC791: Internet Protocol Specification



Focused on delivering packets to destination *node*

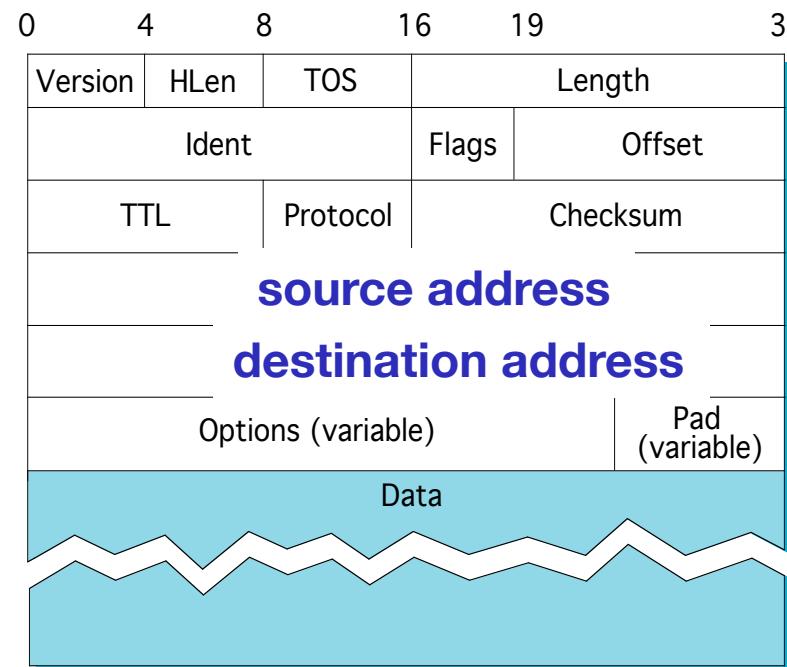


The IP model **separated** source-to-destination packet delivery from the specifics of which paths to take

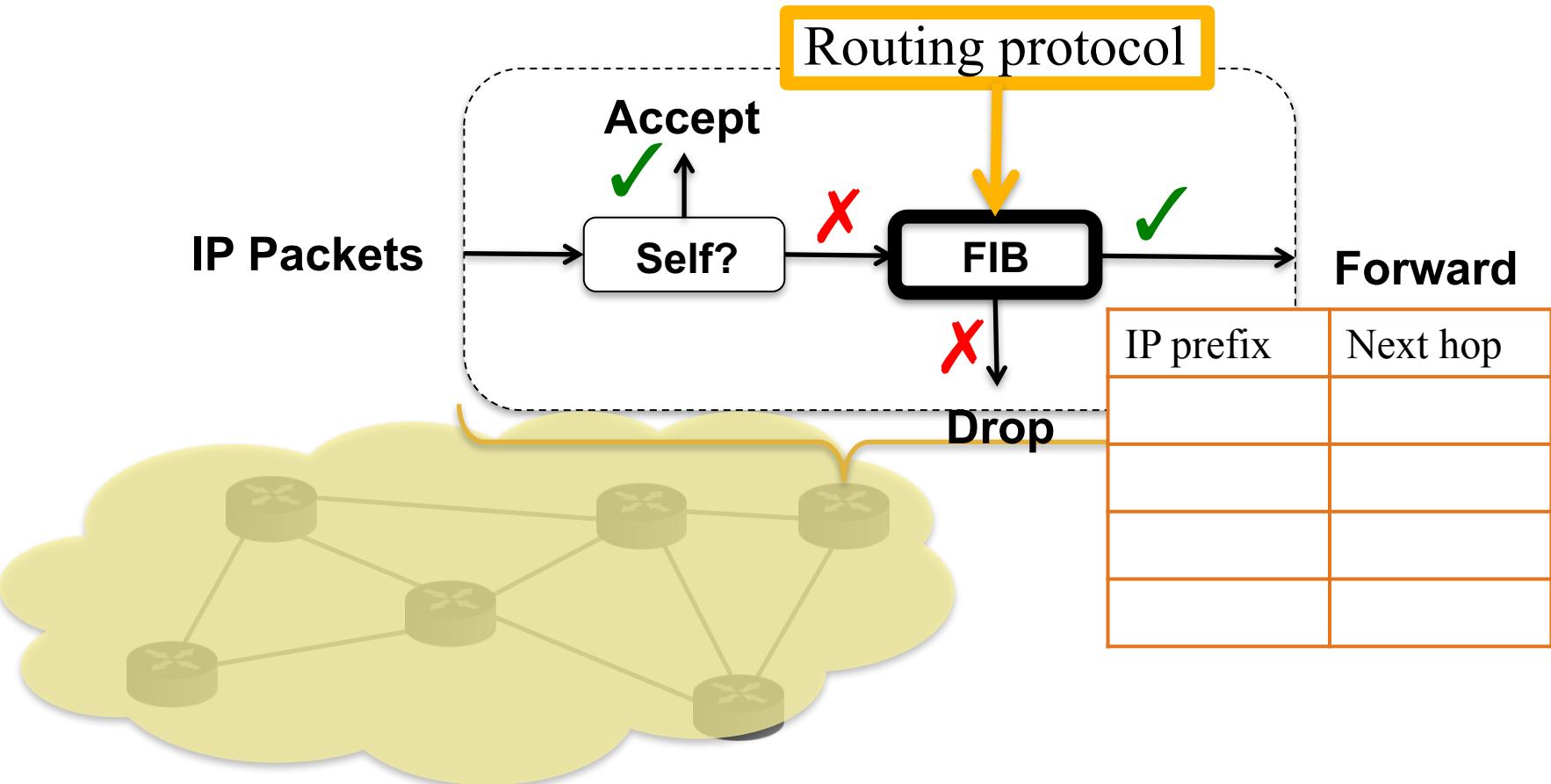


Today's Internet

- ◆ Created gizzalins of applications, a whole world of contents
- ◆ IP got us here, but was not designed for this reality
 - The central *abstraction* is a host identifier
 - The communication model: point-to-point conversation between a host pair



IP' Node Model

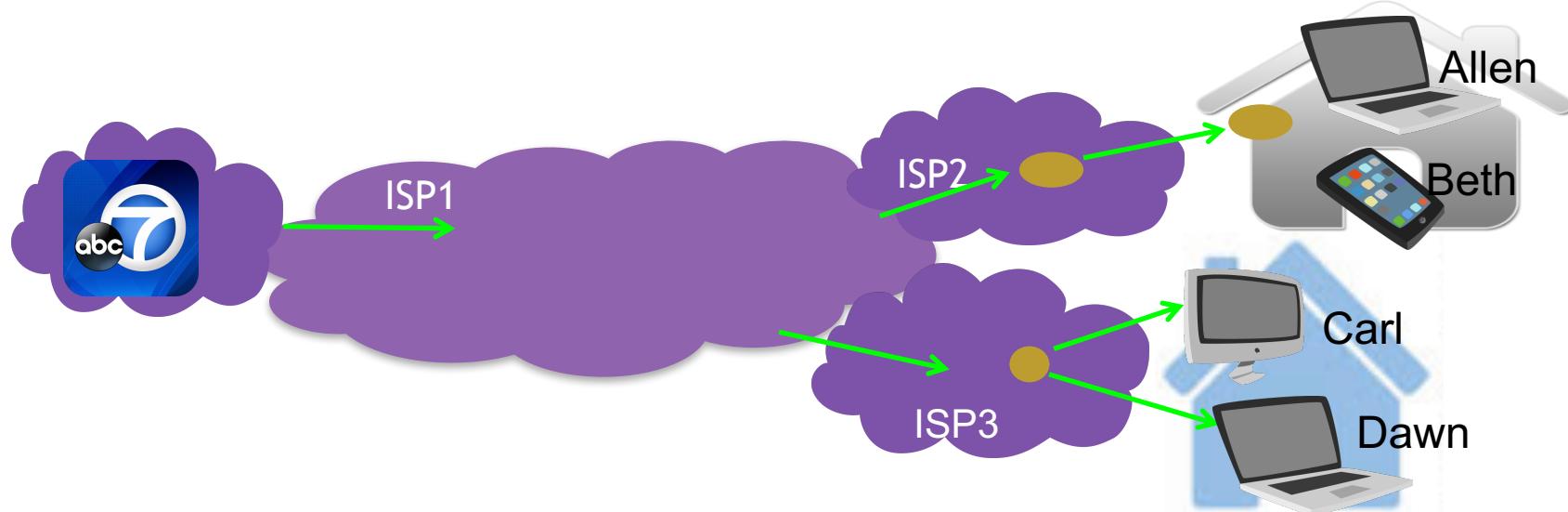


One-way traffic, stateless, no storage

How well IP serves today's applications

Example 1: content delivery

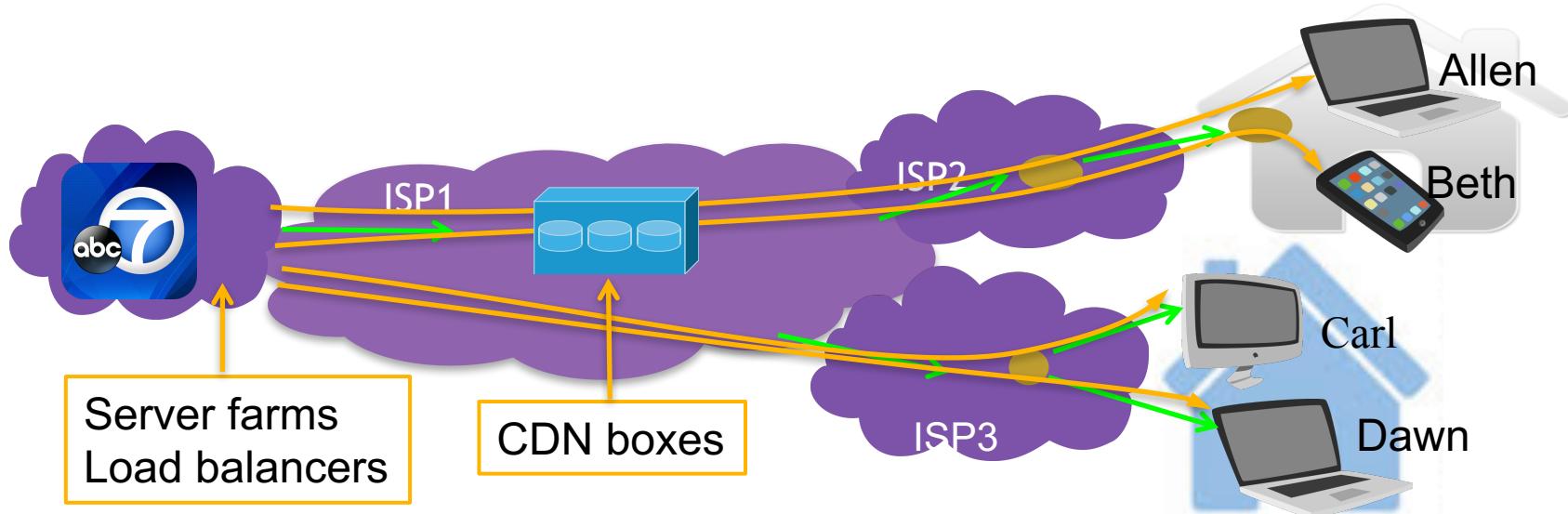
- Applications request data by names; network name packets by IP addresses



How well IP serves today's applications

Example 1: content delivery

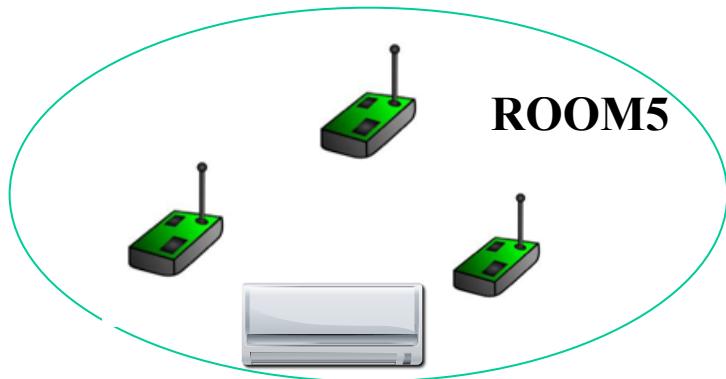
- ◆ Applications request data by names; network name packets by IP addresses
- ◆ IP delivers data between two end points
 - Multiple users may request the same data



How well IP serves today's applications

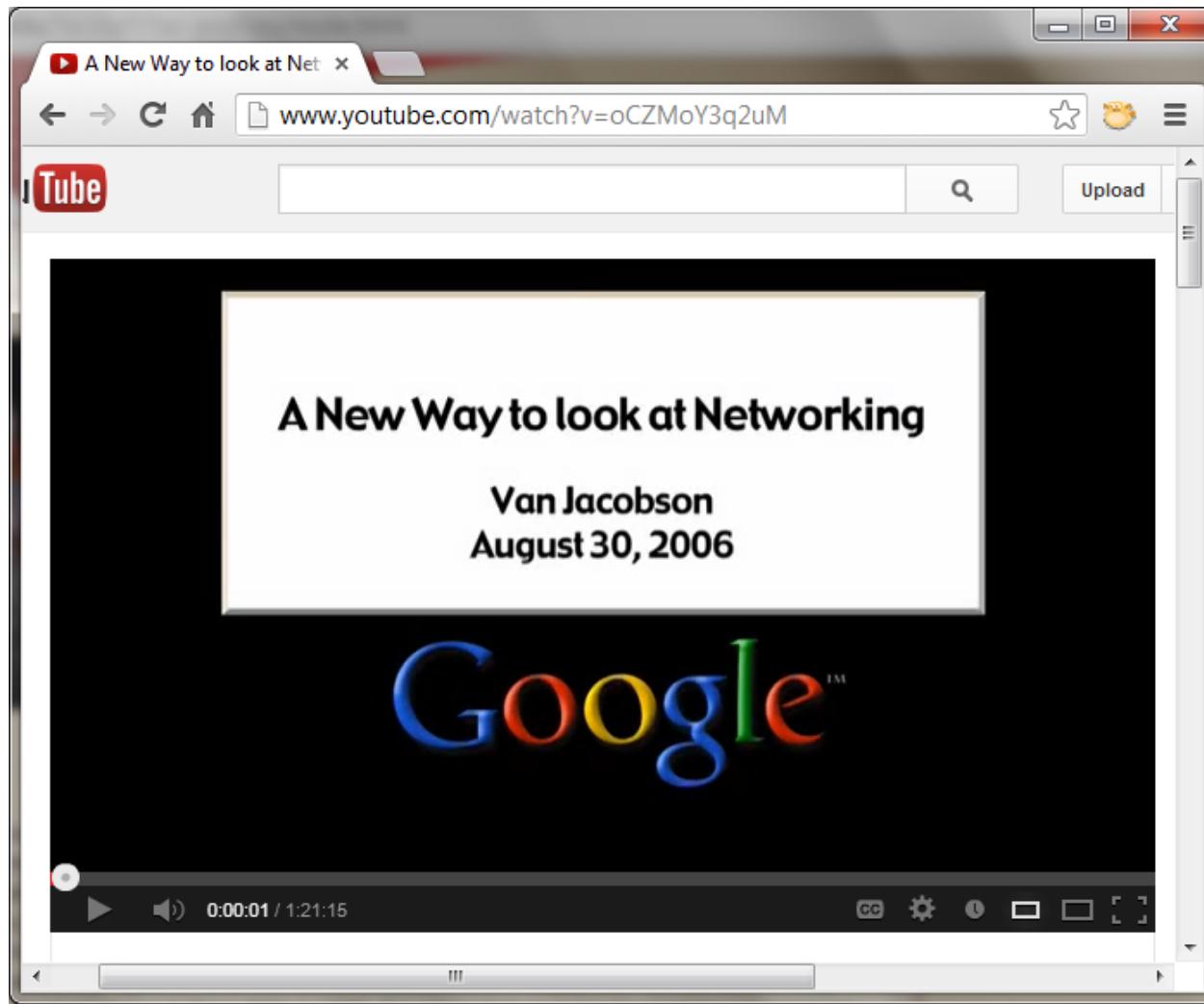
Example 2: emerging network applications

“ROOM5 temperature?”



“Turn on air conditioner”

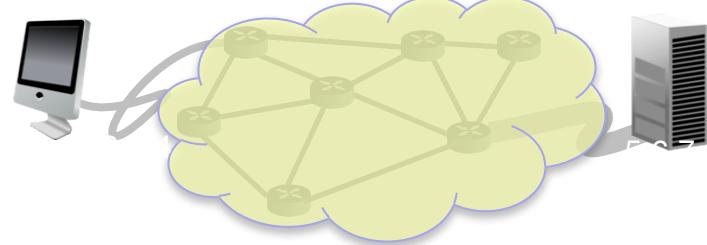




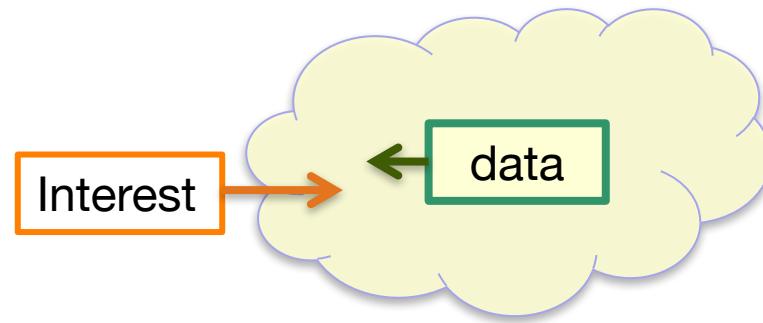
<https://www.youtube.com/watch?v=oCZMoY3q2uM>

Core Idea

Today's Internet is based on host-to-host connections



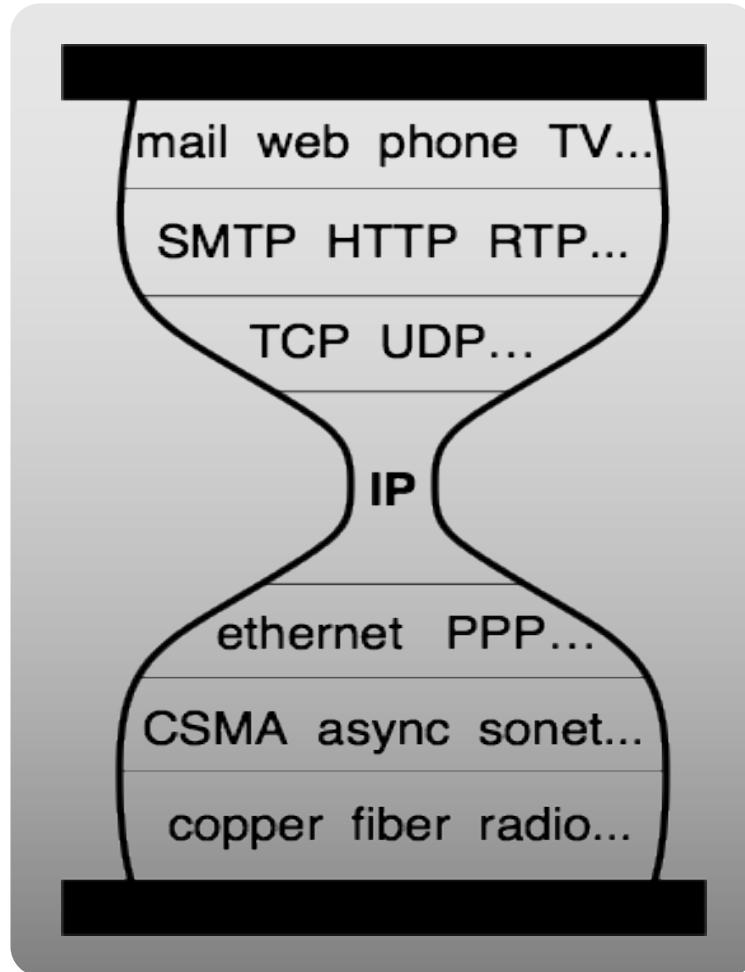
Modern communication consists of requests for named data



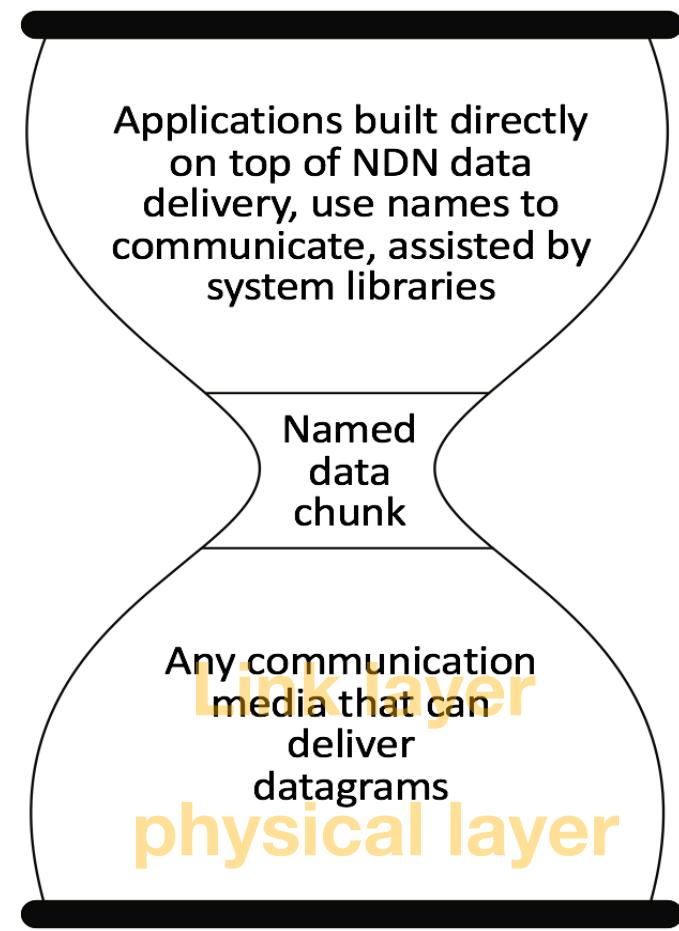
NDN: a general-purpose network protocol supporting requests for named data

NDN: Preserving Hourglass Shape

TCP/IP



NDN



Two Packet Types

Interest packet

Data name

(may carry a few optional parameters)

Data packet

Data name

meta information

Data

Crypto signature

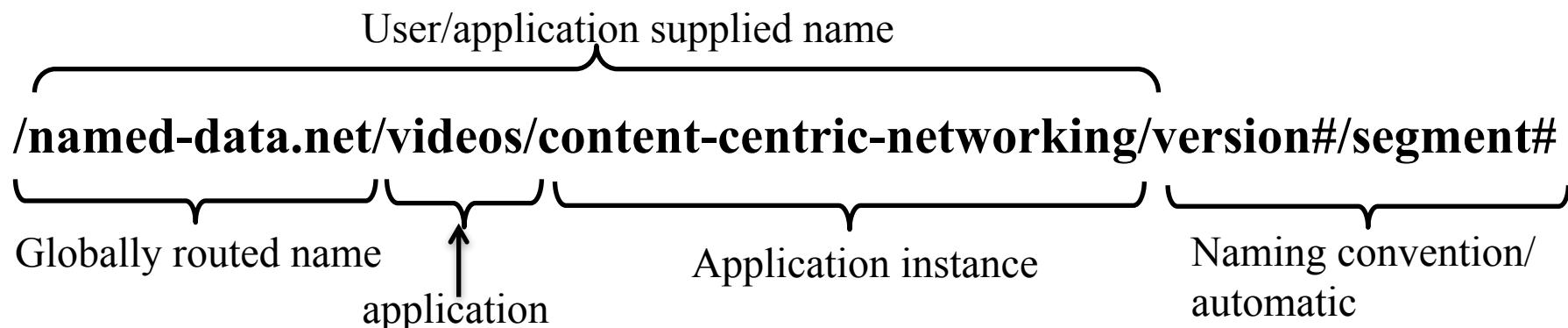
Data consumers send Interest packets

Whoever has the matching Data packet can reply

- ◆ Publisher binds name to data; receivers verify
- ◆ No address !

Content Naming

- Names are generated by applications, opaque to the network
 - Packet granularity
 - Hierarchical
 - identify content relationship & facilitate aggregation
 - Every data packet carries a signature, binding the name to the content



Applications fetch data by names, today

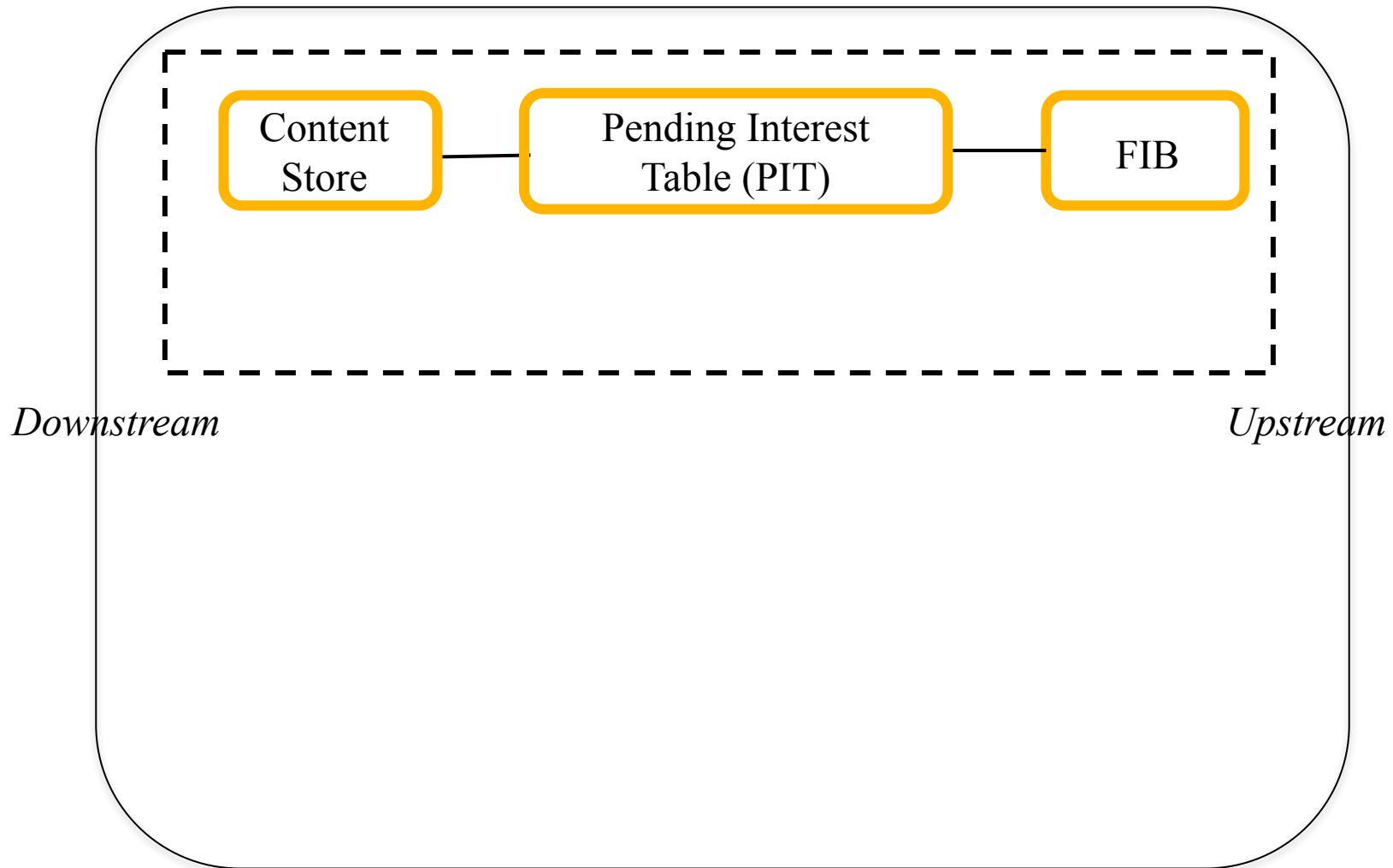
- ◆ My browser: get nytimes.com
- ◆ What happens today:
 - Translate that name to IP addresses
 - identify packets by piecing together parameters from multiple protocol layers
 - nytimes.com → destination IP address
 - packets for nytimes.com identified by a combination of [senderIP, recvIP, Protocol-ID, send-port#, recv-port#, seq#]

Doing the same in NDN: a much simpler picture

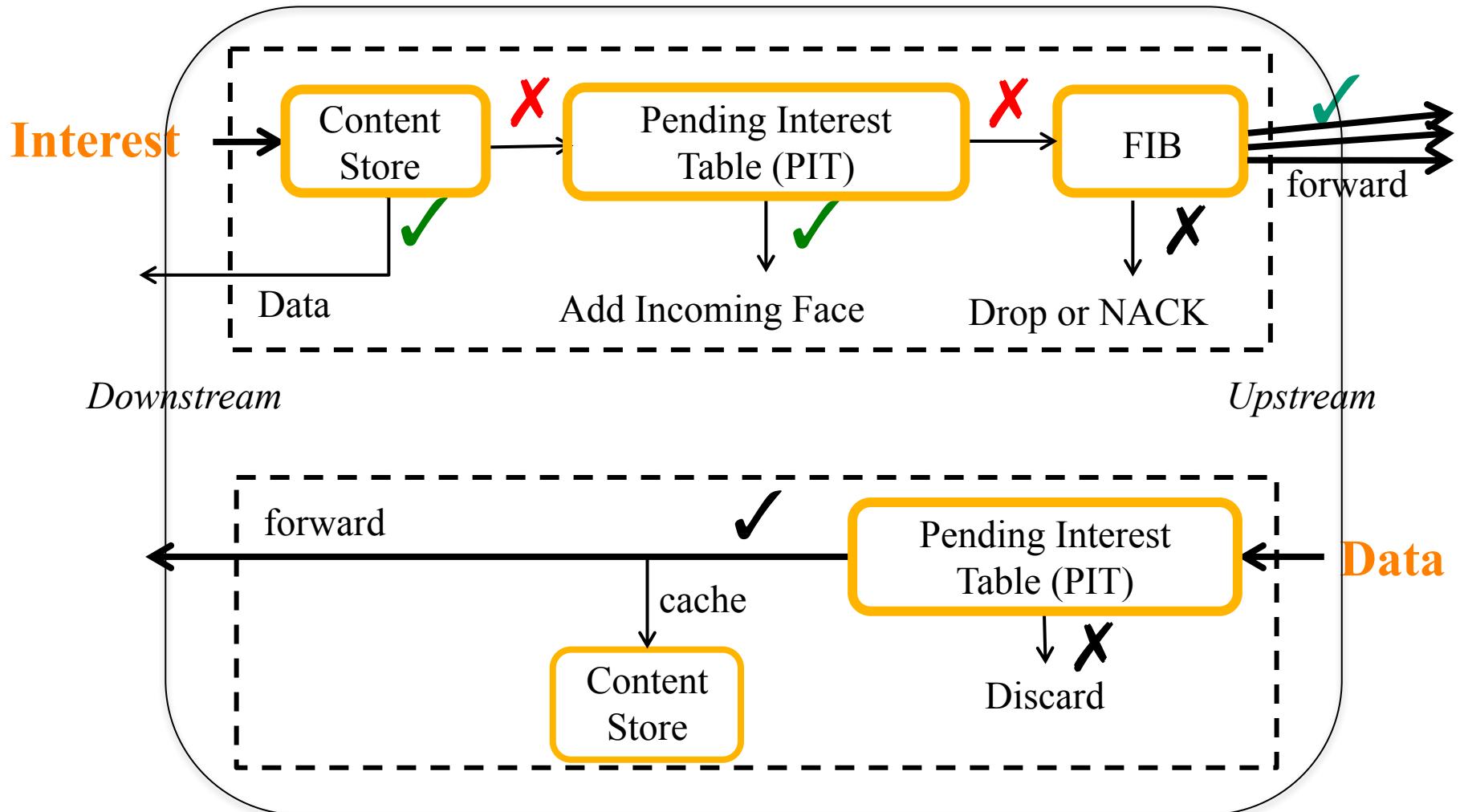
To retrieve front page of **nytimes.com**

- ◆ My browser: Send an Interest with name
/nytimes.com/web/index.html
- ◆ NewYorkTimes Web server: listening for any incoming Interest with the name prefix
/nytimes.com/web/
 - Receive the Interest packet
 - Create data packet, *add signature*
 - Send out to the incoming interface of the Interest

NDN's node model



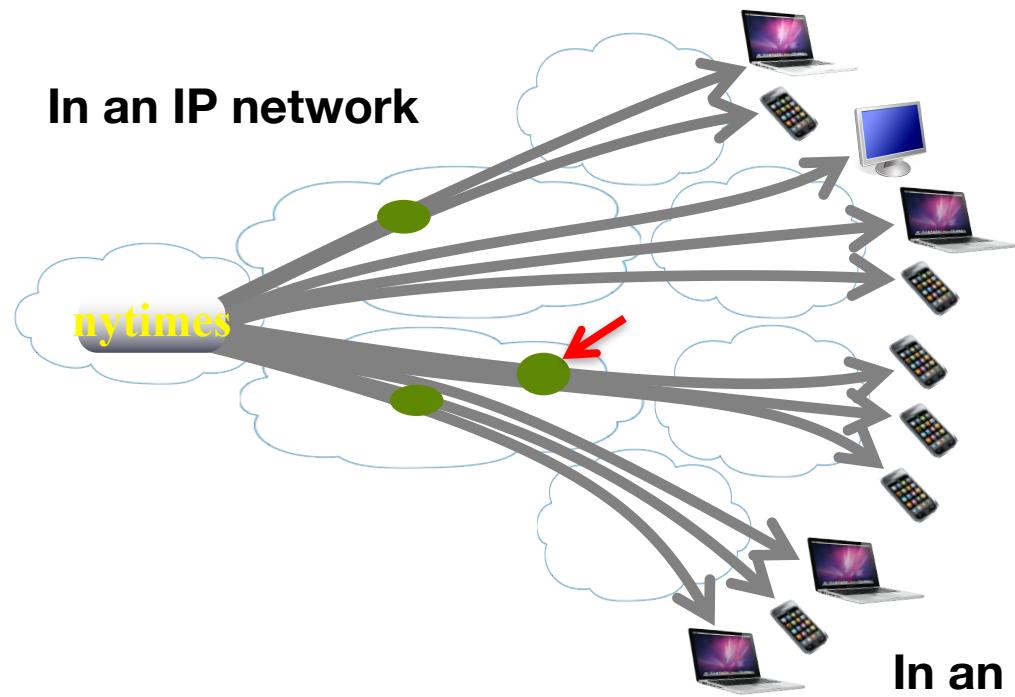
NDN's node model



2-way packet flow, fetch, stateful, with storage

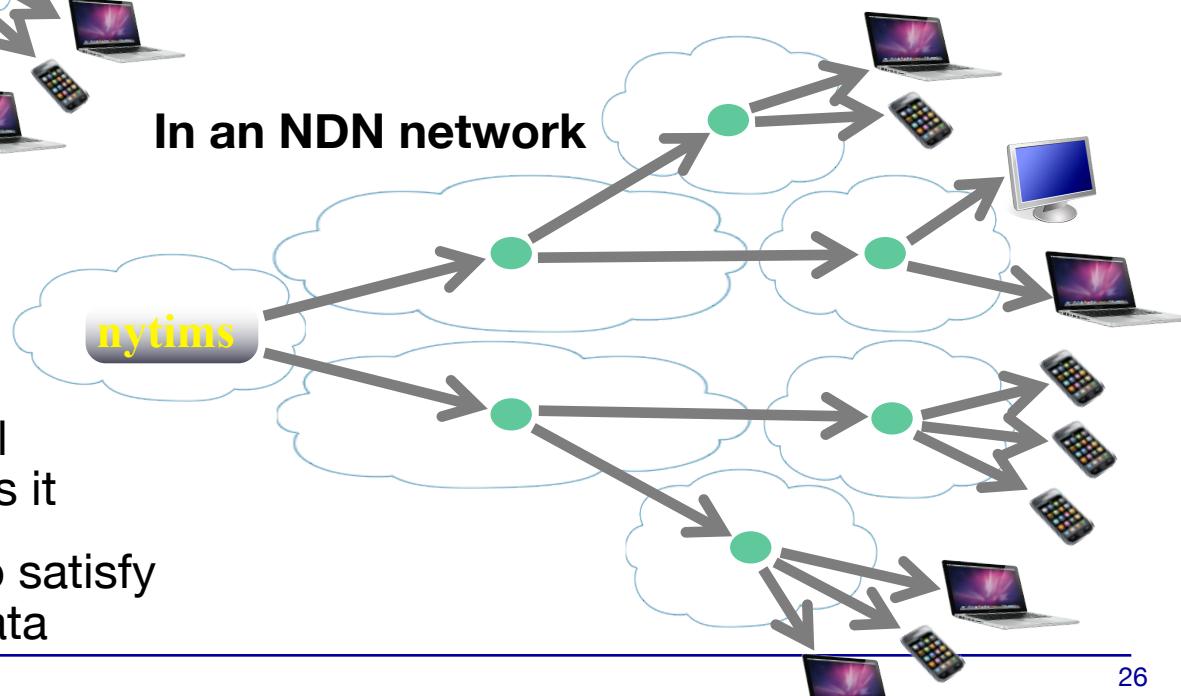
What's the difference to the network

In an IP network



- Router forwards each packet towards its destination
- The packet is useless after been forwarded

In an NDN network

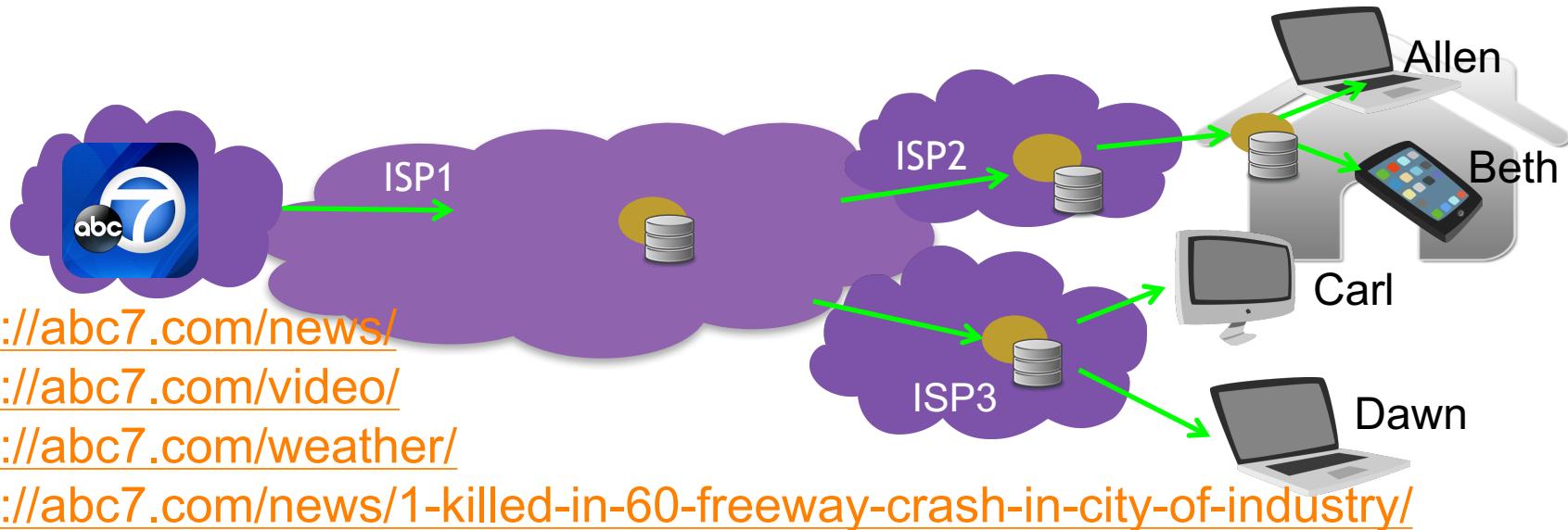


- Each data packet is meaningful independent from who requests it
- Data packets can be cached to satisfy future requests for the same data

How well NDN can serve applications

Example 1: content delivery

- ◆ Network uses app. data names for delivery
- ◆ Multiple users request the same data: net can retrieve from nearby copy
 - Name + data-signature enables in-network storage



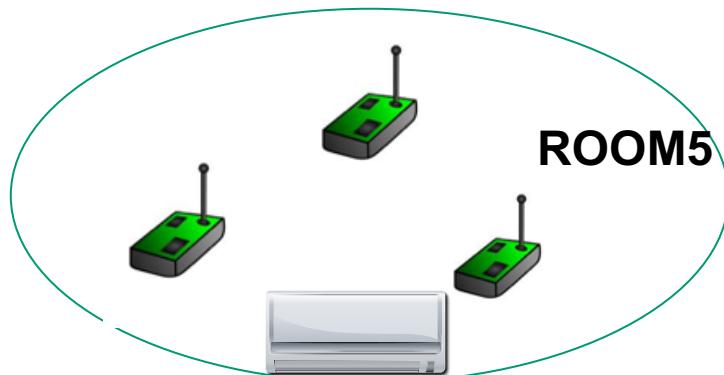
How well NDN can serve applications

Example 2: emerging network applications

“ROOM5 temperature?”

INTEREST(/ucla/bldg#/room5/temp) →

← DATA (name|data|signature)



“Turn on air conditioner”

INTEREST(/ucla/bldg#/room5/AC-on/sig) →

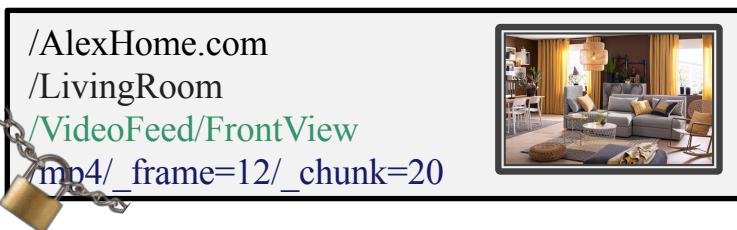
← DATA (name|ACK|signature)

INTEREST(/traffic/LA/HW405/location) →
← DATA (name|data|signature)

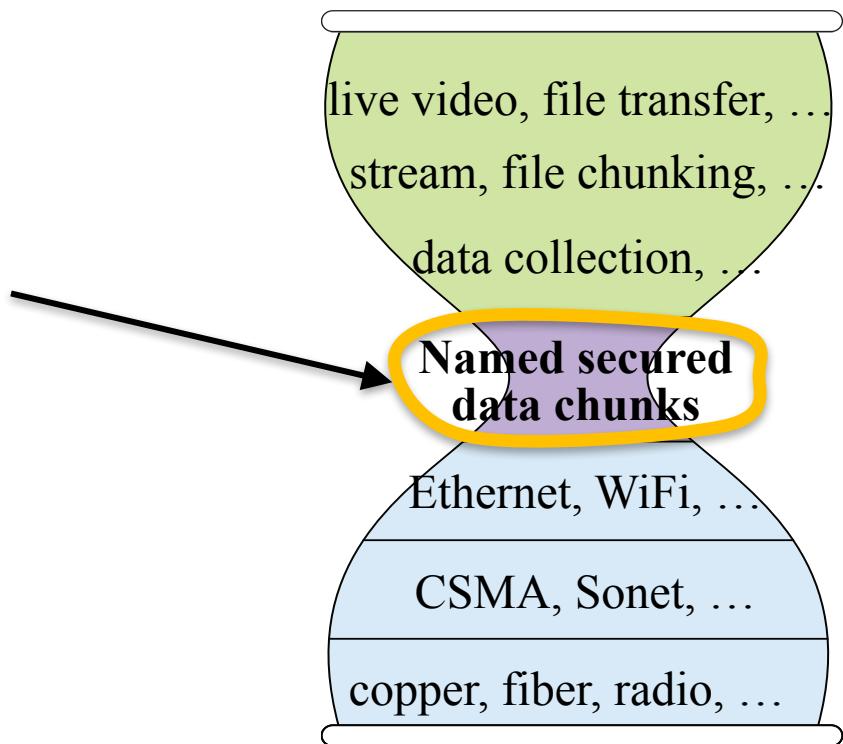


Named Data Networking: Built-in Security

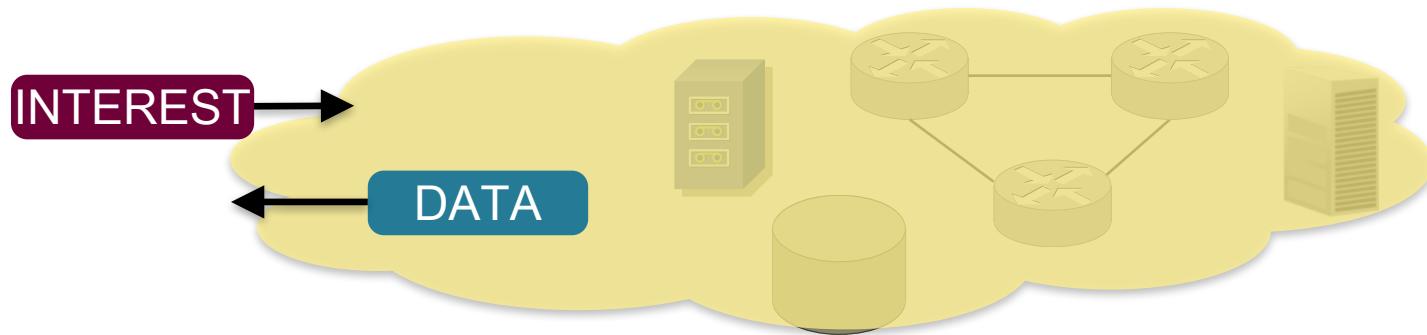
- ◆ Hierarchically structured names, shared between application and network layers
- ◆ Security
 - Built-in into the networking layer



- ◆ Focus on application data
 - Data secured in motion and at rest
- ◆ Universal mechanism
 - Same security mechanisms for networking, transport, and application layers



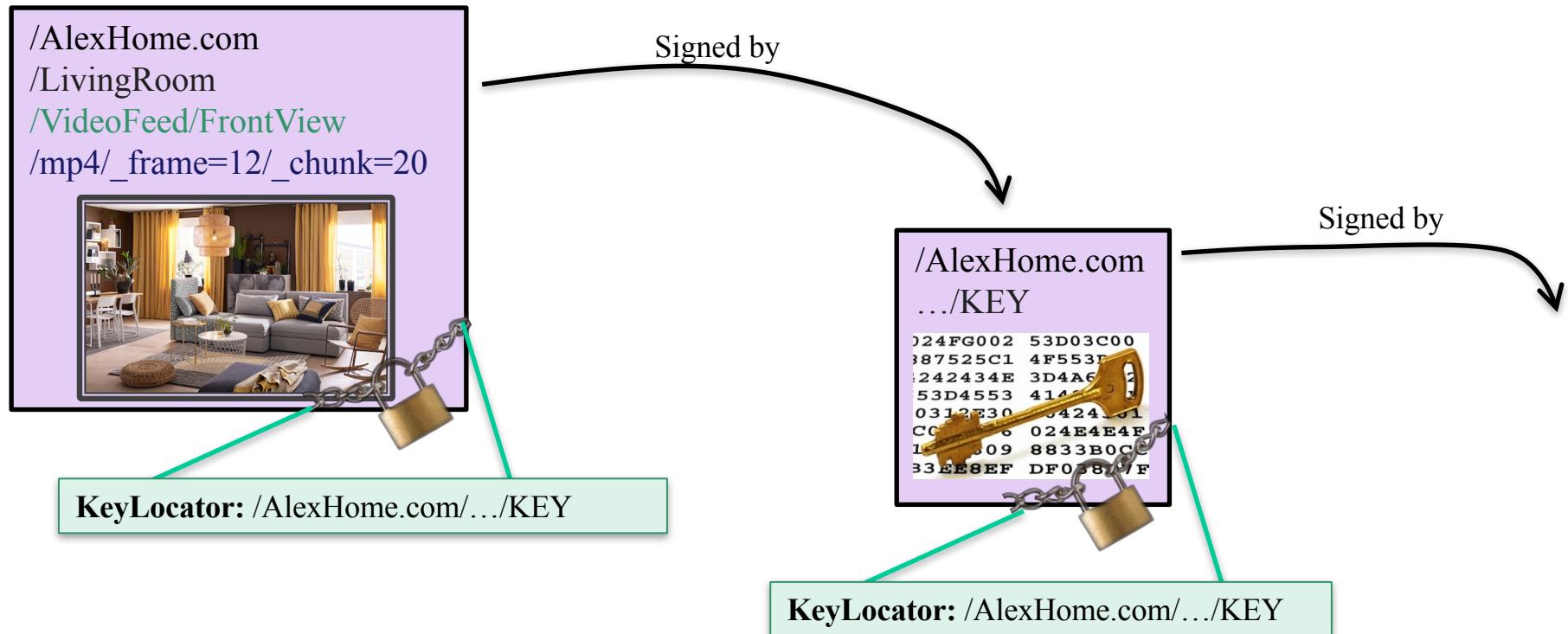
Security: built into the narrow waist



- ◆ Every data packet is signed
 - encrypted whenever needed
 - Data always secured, whether in motion or at rest
- ◆ Keys retrieved in the same way as any other content objects

How to verify a key: trust management

How NDN's Data-Centric Authenticity Works?



Not Just One Key

/AlexHome.com/LivingRoom/Video
Feed

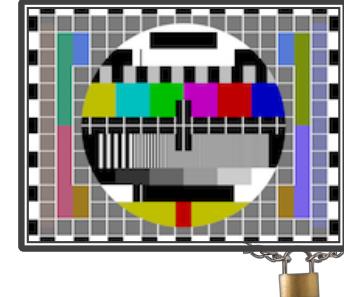
/FrontView/mp4/_frame=12/_chunk
=20



/AlexHome.com/Camera/KEY

/AlexHome.com/LivingRoom/Video
Feed

/FrontView/mp4/_frame=42/_chunk
=1



/AlexHome.com/TV/KEY

A frame from a camera
I have installed in my
living room



TV incorrectly trying to
publish living room
feed



NDN Transport Function

- ◆ What TCP does today
 - Demultiplexing
 - Reliability (including sequencing)
 - Congestion control
 - More recently: multipath TCP
- ◆ In NDN
 - Demultiplexing: by names directly
 - Reliability: controlled by applications directly
 - Congestion control: by network layer
 - Multipath forwarding: by network layer

NDN: Just 3 Simple Ideas

1. Naming data

→ Routing & forwarding by names

2. Securing every data packet

→ removing dependency on transport security

3. Per Interest, per hop forwarding state

→ Creating closed feedback loop

- ▲ measure performance, detect failures

→ Enabling multi-path forwarding

- ▲ Add a strategy module to assist the forwarding decisions

The road to a new architecture

- ◆ Application-driven development
 - Running code, useful apps, testbed with real traffic
 - tackling emerging environments and applications where no good IP-based solutions exist
- ◆ Incremental Deployment
 - NDN runs on everything, and everything runs on NDN
 - Start as an overlay, the same way as IP did

What to take home

- ◆ Future of networking lies in recognizing the right communication abstraction
- ◆ IP conceptualizes communication as between nodes
- ◆ NDN directly focuses on the outcome: retrieving data

