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Configuration files

```
.vimrc
set nocp
set mouse=a
filetype plugin indent on
svntax on
set smd ls=2 nu sm is nohls bg=dark
set et sw=4 sts=4 ts=8 sta ai ci
set nowb nobk noswf
Template
#include <bits/stdc++.h>
using namespace std;
#define INF 0x3f3f3f3f
#define MOD 1000000007
#define PI M_PI
\#define mset(a, x) memset(a, x, sizeof(a))
#define pb push_back
#define mp make_pair
#define fi first
#define se second
typedef long long ll;
typedef unsigned long long ull;
typedef pair<int, int> pii;
const double inf = 1.0/0.0;
int cmp_double(double a, double b, double eps = 1e-9) {
    return a + eps > b? b + eps > a? 0:1:-1;
}
int main() {
```

Useful facts

Erdös-Gallai theorem: A sequence of non-negative integers $d_1 \geq \cdots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + \cdots + d_n$ is even and $\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i,k)$ holds for $1 \leq k \leq n$.

Split graph property: A split graph can be recognized solely from their degree sequence. Let the degree sequence of a graph G be $d_1 \geq \cdots \geq d_n$ and m is the largest value of i such that $d_i \geq i-1$. Then G if a split graph if and only if $\sum_{i=1}^m d_i = m(m-1) + \sum_{i=m+1}^n d_i$.

Stirling's approximation: $(n \ge 100)$ $\ln n! \approx n \ln n - n + \frac{1}{2} \ln(2\pi n)$

2-SAT: Algorithm for solving Boolean expression in 2-CNF form (example: $(A \lor B) \land (B \lor \sim C) \land (A \lor C) \land (B \lor D)$).

- 1) Transform each term of conjunctions $(A \lor B)$ into $(\sim A \to B) \land (\sim B \to A)$
- 2) Construct graph G = (V, E) such that each literal is a vertex and each implication is an edge
- 3) Run SCC algorithm. If there is a SCC such that A and $\sim A$ are in it, so the expression cannot be evaluated TRUE. Otherwise, it is possible.

Pick's Theorem: $A = i + \frac{b}{2} - 1$

Graph

}

```
Tarjan
Complexity: O(V+E)
int n, m;
vector<int> g[MAXN];
int lbl[MAXN], low[MAXN], idx, cnt_scc;
stack<int> st;
bool inSt[MAXN];
void dfs(int u) {
   lbl[u] = low[u] = idx++;
    st.push(u);
    inSt[u] = 1;
    for (int i = 0; i < q[u].size(); i++) {
        int v = q[u][i];
        if (lbl[v] == -1) {
            dfs(v);
            low[u] = min(low[u], low[v]);
        } else if (inSt[v]) {
            low[u] = min(low[u], lbl[v]);
        }
   if (low[u] == lbl[u]) {
        printf("%d -> ", ++cnt_scc);
        int v;
        do {
            v = st.top();
            st.pop();
            inSt[v] = 0;
            printf("%d; ", v);
        } while (v != u);
        putchar('\n');
}
void tarjan() {
    for (int i = 1; i <= n; i++) {
        lbl[i] = -1;
        inSt[i] = 0;
   }
   idx = cnt\_scc = 0;
    for (int i = 1; i <= n; i++)
        if (lbl\lceil i \rceil == -1)
            dfs(i);
```

Articulation

```
Complexity: O(V+E)
int n, m;
vector<int> g[MAXN];
int lbl[MAXN], low[MAXN], parent[MAXN], idx;
bool art[MAXN], has_art;
void dfs(int u) {
    int cnt = 0;
    lbl[u] = low[u] = idx++;
    for (int i = 0; i < g[u].size(); i++) {
        int v = g[u][i];
        if (lbl[v] == -1) {
            parent[v] = u;
            dfs(v);
            low[u] = min(low[u], low[v]);
            if (low[v] >= lbl[u])
                cnt++;
        } else if (v != parent[u]) {
            low[u] = min(low[u], lbl[v]);
        }
    }
    if (cnt > 1 || (lbl[u] != 0 && cnt > 0)) {
        art[u] = 1;
        has_art = 1;
void articulation() {
    for (int i = 1; i <= n; i++) {
        lbl[i] = -1;
        art[i] = 0;
    for (int i = 1; i <= n; i++) {
        if (lbl\lceil i \rceil == -1) {
            idx = 0;
            parent[i] = i;
            dfs(i);
        }
}
```

Bridge

```
Complexity: O(V+E)
int n, m;
vector<int> q[MAXN];
int lbl[MAXN], low[MAXN], parent[MAXN], idx;
bool has_bridge;
void dfs(int u) {
   lbl[u] = low[u] = idx++;
    bool parent_found = 0;
   for (int i = 0; i < g[u].size(); i++) {
        int v = g[u][i];
        if (lbl[v] == -1) {
            parent[v] = u;
            dfs(v);
            low[u] = min(low[u], low[v]);
            if (low[v] == lbl[v]) {
               printf("%d -> %d\n", u, v);
                has_bridge = 1;
        } else if (!parent_found && v == parent[u]) {
            parent_found = 1;
       } else {
           low[u] = min(low[u], lbl[v]);
}
void bridge() {
    for (int i = 1; i <= n; i++)
        lbl[i] = -1;
    for (int i = 1; i <= n; i++) {
        if (lbl[i] == -1) {
           idx = 0;
            parent[i] = i;
            dfs(i);
}
Hopcroft-Karp
Complexity: 0(E sqrt(V))
int n, m;
vector<int> g1[MAXN];
int pair_g1[MAXN], pair_g2[MAXM], dist[MAXN];
```

```
bool bfs() {
    queue<int> q;
    for (int u = 1; u <= n; u++) {
        dist[u] = INF;
        if (pair_g1[u] == 0) {
            dist[u] = 0;
            q.push(u);
       }
   }
   dist[0] = INF;
   while (!q.empty()) {
        int u = q.front(); q.pop();
        for (int i = 0; i < q1[u].size(); i++) {
           int v = q1[u][i];
            if (dist[pair_g2[v]] == INF) {
                dist[pair_g2[v]] = dist[u] + 1;
                q.push(pair_q2[v]);
       }
    return dist[0] != INF;
bool dfs(int u) {
   if (u == 0)
        return 1;
    for (int i = 0; i < g1[u].size(); i++) {
        int v = g1[u][i];
       if (dist[pair_q2[v]] == dist[u] + 1 && dfs(pair_q2[v])) {
            pair_q1[u] = v;
            pair_a2[v] = u;
            return 1;
       }
    dist[u] = INF;
    return 0;
int hk() {
   memset(pair_q1, 0, sizeof(pair_q1));
   memset(pair_q2, 0, sizeof(pair_q2));
   int matching = 0;
    while (bfs())
        for (int u = 1; u <= n; u++)
           if (pair_q1[u] == 0 \&\& dfs(u))
                matching++;
    return matching;
```

Dinic

```
Complexity: O(V^2 E)
struct Edge {
   int v, c, f, next;
   Edge() {}
   Edge(int v, int c, int f, int next) : v(v), c(c), f(f), next(next) {}
};
int n, m, head[MAXN], lvl[MAXN], src, snk, work[MAXN];
Edge e[MAXM];
void init(int _n, int _src, int _snk) {
    n = _n;
   m = 0;
    src = _src;
    snk = \_snk;
    memset(head. -1. sizeof(head)):
void addEdge(int u, int v, int c) {
    e[m] = Edge(v, c, 0, head[u]);
   head[u] = m++;
   e[m] = Edge(u, 0, 0, head[v]);
   head[v] = m++;
}
bool bfs() {
    queue<int> q;
    memset(lvl, -1, n * sizeof(int));
   lvl[src] = 0;
    q.push(src);
   while (!q.empty()) {
        int u = q.front(); q.pop();
        for (int i = head[u]; i != -1; i = e[i].next) {
            if (e[i].f < e[i].c && lvl[e[i].v] == -1) {
                [v][e[i].v] = [v][u] + 1;
                q.push(e[i].v);
                if (e[i].v == snk)
                    return 1;
           }
        }
    return 0;
```

```
int dfs(int u, int f) {
    if (u == snk)
        return f;
    for (int \&i = work[u]; i != -1; i = e[i].next) {
        if (e[i].f < e[i].c \&\& lvl[u] + 1 == lvl[e[i].v]) {
            int minf = dfs(e[i].v, min(f, e[i].c - e[i].f));
            if (minf > 0) {
                e[i].f += minf;
                e[i^1].f -= minf;
                return minf;
        }
    return 0;
int dinic() {
    int f, ret = 0;
    while (bfs()) {
        memcpy(work, head, n * sizeof(int));
        while (f = dfs(src, INF))
            ret += f;
    return ret;
Lowest Common Ancestor
Complexity: < O(N logN), O(logN) >
#define MAXN 50000
#define LOGMAXN 16
int n, m, u, v, w;
int ancestor[MAXN][LOGMAXN], parent[MAXN], level[MAXN], dist[MAXN];
vector<pair<int, int> > g[MAXN];
void dfs(int u) {
    for (int i = 0; i < g[u].size(); i++) {
        int v = q[u][i].first, w = q[u][i].second;
        if (v != parent[u]) {
            parent[v] = u;
            level[v] = level[u] + 1;
            dist[v] = dist[u] + w;
            dfs(v);
       }
}
```

```
void pre() {
    parent[0] = level[0] = dist[0] = 0;
    dfs(0);
    for (int i = 0; i < n; i++)
        ancestor[i][0] = parent[i];
    for (int j = 1; 1 << j < n; j++)
        for (int i = 0; i < n; i++)
            ancestor[i][j] = ancestor[ancestor[i][j-1]][j-1];
}
int lca(int u, int v) {
   if (level[u] < level[v])</pre>
        swap(u, v);
    int log;
    for (log = 1; 1 < log <= level[u]; log++);
    for (int i = log; i >= 0; i--)
        if (level[u] - (1 << i) >= level[v])
            u = ancestor[u][i]:
   if (u == v)
        return u;
    for (int i = log; i >= 0; i--)
        if (ancestor[u][i] != ancestor[v][i])
            u = ancestor[u][i], v = ancestor[v][i];
    return parent[u];
}
Min-Cost Max-Flow
Complexity: O(V E + f E \log V), f = maximum flow
struct Edge {
   int u, v, cap, flow, cost, next;
   Edge() {}
    Edge(int u, int v, int cap, int flow, int cost, int next)
        : u(u), v(v), cap(cap), flow(flow), cost(cost), next(next) {}
};
int n, m, head[MAXN], src, snk;
Edae e[MAXM]:
int pi[MAXN], dist[MAXN], path[MAXN], mincap[MAXN], vis[MAXN];
void init(int _n, int _src, int _snk) {
   n = _n;
   m = 0;
    src = _src;
    snk = \_snk;
    memset(head, -1, sizeof(head));
}
```

```
void addEdge(int u, int v, int cap, int cost) {
    e[m] = Edge(u, v, cap, 0, cost, head[u]);
    head[u] = m++;
    e[m] = Edge(v, u, 0, 0, -cost, head[v]);
    head[v] = m++;
}
int bellman_ford() {
    memset(pi, INF, sizeof(pi));
    pi[src] = 0;
    int flaq = 1;
    for (int i = 0; flag && i < n; i++) {
        flaa = 0:
        for (int j = 0; j < m; j++) {
            if (e[j].cap == e[j].flow) continue;
            if (pi[e[j].u] + e[j].cost < pi[e[j].v]) {
                pi[e[j].v] = pi[e[j].u] + e[j].cost;
                flaq = 1;
       }
    return flag;
int dijkstra() {
    priority_queue<pii, vector<pii>, greater<pii> > heap;
    memset(dist, INF, sizeof(dist));
    memset(vis, 0, sizeof(vis));
    dist[src] = 0;
    mincap[src] = INF;
    heap.push(mp(0, src));
    while (!heap.empty()) {
        int u = heap.top().se;
        heap.pop();
       if (vis[u]) continue;
        vis[u] = 1;
        for (int i = head[u]; i != -1; i = e[i].next) {
            int v = e[i].v;
            if (vis[v] | l e[i].flow == e[i].cap) continue;
            int w = dist[u] + e[i].cost + pi[u] - pi[v];
            if (w < dist[v]) {
                dist[v] = w:
                path[v] = i;
                mincap[v] = min(mincap[u], e[i].cap - e[i].flow);
                heap.push(mp(dist[v], v));
       }
   }
```

```
// update potencials
    for (int i = 0; i < n; i++)
        pi[i] += dist[i];
    return dist[snk] < INF;</pre>
}
pii mcmf() {
    // set potencials
    if (bellman_ford())
        return mp(-1, -1);
    int cost = 0, flow = 0;
    while (dijkstra()) {
        // augment path and update cost
        int f = mincap[snk];
        for (int v = snk; v != src; ) {
            int idx = path\lceil v \rceil;
            e[idx].flow += f;
            e[idx^1].flow -= f;
            cost += e[idx].cost * f;
            v = e[idx].u;
        flow += f;
    return mp(cost, flow);
}
Heavy-Light Decomposition
struct SegTree {
    vector<int> data, tree;
    int sz;
    SegTree(int tsz) : sz(1) {
        while (sz < tsz) sz *= 2;
        data.resize(sz):
        tree.resize(2*sz);
    }
    inline int left(int u) { return u << 1; }</pre>
    inline int right(int u) { return left(u) + 1; }
    void init(int u, int l, int r) {
        if (l == r) { tree[u] = data[l]; return; }
        int m = (l + r) >> 1;
        init(left(u), l, m):
        init(right(u), m+1, r);
        tree[u] = max(tree[left(u)], tree[right(u)]);
    void init() { init(1, 0, sz-1); }
```

```
int query(int u, int l, int r, int a, int b) {
        if (a <= 1 && r <= b) return tree[u];
        int m = (l + r) >> 1, ret = 0;
        if (a \le m) ret = query(left(u), l, m, a, b);
        if (m < b) ret = max(ret, query(right(u), m+1, r, a, b));
        return ret;
    int query(int a, int b) { return query(1, 0, sz-1, a, b); }
    void update(int u, int l, int r, int pos, int val) {
        if (l == r)  { tree[u] = val; return; }
        int m = (l + r) >> 1;
        if (pos <= m) update(left(u), l, m, pos, val);</pre>
        else update(right(u), m+1, r, pos, val);
        tree[u] = max(tree[left(u)], tree[right(u)]):
    void update(int pos, int val) { update(1, 0, sz-1, pos, val); }
};
struct edge {
    int u, v, w, next;
    edge() {}
    edge(int u, int v, int w, int next) : u(u), v(v), w(w), next(next) {}
};
int n, m, head[MAX];
edge e[2*MAX];
int dad[MAX], lvl[MAX], chd[MAX], sz[MAX], heavy[MAX];
int nump, psize[MAX], pfirst[MAX], path[MAX], offset[MAX];
vector<int> walk;
vector<SegTree> ptree;
void init() {
    m = 0;
    memset(head, -1, sizeof(head));
    memset(dad, 0, sizeof(dad));
void addEdge(int u, int v, int w, bool rev = false) {
    e[m] = edge(u, v, w, head[u]);
    head[u] = m++;
    if (!rev) addEdge(v, u, w, true);
```

```
void dfs(int u) {
    walk.push_back(u);
    sz[u] = 1;
    for (int i = head[u]; i != -1; i = e[i].next) {
        int v = e[i].v;
        if (!dad[v]) {
            dad[v] = u, lvl[v] = lvl[u] + 1, dfs(v);
            sz[u] += sz[v];
    for (int i = head[u]; i != -1; i = e[i].next) {
        int v = e[i].v;
        if (dad[v] == u \&\& 2*sz[v] >= sz[u]) {
            heavy[v] = 1;
            break:
        }
    heavy[u] = 0;
}
void hl_init() {
    walk.clear();
    dad[1] = 1, lvl[1] = 0, dfs(1);
   nump = 0;
    for (int i = 0; i < n; i++) {
        int u = walk[i];
        if (!heavy[u]) {
            offset\lceil u \rceil = 0;
            path[u] = nump;
            pfirst[nump] = u;
            psize[nump++] = 1;
        else {
            offset[u] = offset[dad[u]] + 1;
            path[u] = path[dad[u]];
            psize[path[u]]++;
    ptree.clear(); ptree.reserve(nump);
    for (int i = 0; i < nump; i++)
        ptree.push_back(SegTree(psize[i]));
    for (int i = 0; i < m; i += 2) {
        int u = e[i].u, v = e[i].v;
        if (u != dad[v]) swap(u, v);
        ptree[path[v]].data[offset[v]] = e[i].w;
    for (int i = 0; i < nump; i++)
        ptree[i].init();
}
```

```
int lca(int u. int v) {
    int fpu = pfirst[path[u]], fpv = pfirst[path[v]];
    while (fpu != fpv) {
        if (lvl[fpu] > lvl[fpv])
            u = dad[fpu], fpu = pfirst[path[u]];
        else
            v = dad[fpv], fpv = pfirst[path[v]];
    return lvl[u] < lvl[v] ? u : v;</pre>
void update(int idx, int val) {
    int u = e[idx].u, v = e[idx].v;
    if (u != dad[v]) swap(u, v);
    ptree[path[v]].update(offset[v], val);
}
int query(int u, int v, bool up = false) {
    if (!up) {
        int w = lca(u, v);
        return max(query(u, w, true), query(v, w, true));
    int ret = 0:
    int fpu = pfirst[path[u]], fpv = pfirst[path[v]];
    while (fpu != fpv) {
        ret = max(ret, ptree[path[u]].query(0, offset[u]));
        u = dad[fpu], fpu = pfirst[path[u]];
    if (u != v)
        ret = max(ret, ptree[path[u]].query(offset[v]+1, offset[u]));
    return ret;
```

String

```
KMP
Complexity: O(N)
int pi[MAXS];
void kmp_table(char *pattern, int m) {
    pi[0] = -1;
    for (int i = 1, k = -1; i < m; i++) {
        while (k \ge 0 \& pattern[k+1] != pattern[i])
            k = pi[k];
        if (pattern[k+1] == pattern[i])
            k++;
        pi[i] = k;
}
void kmp_search(char *text, char *pattern) {
    int n = strlen(text), m = strlen(pattern);
    kmp_table(pattern, m);
    for (int i = 0, k = -1; i < n; i++) {
        while (k \ge 0 \&\& pattern[k+1] != text[i])
            k = pi[k];
        if (pattern[k+1] == text[i])
            k++;
        if (k+1 == m) {
            printf("Match at %d\n", i-k);
            k = pi[k];
   }
}
String Hashing
#define B 33
ull powB[MAX];
void init() {
    powB[0] = 1;
    for (int i = 1; i < MAX; i++)
        powB[i] = B*powB[i-1];
}
```

```
void calc_hash(char *str, ull *h) {
    h[0] = 0;
    for (int i = 0; str[i]; i++)
        h[i+1] = B*h[i] + str[i];
ull get_hash(ull *h, int l, int r) {
    return h[r] - h[l]*powB[r-l];
Manacher's Algorithm
Complexity: O(N)
char s[MAXN];
int p[2*MAXN]; // length of the palindrome centered at position (i-1)/2;
void manacher() {
    int m = 0;
    char t[2*MAXN];
    for (int i = 0; s[i]; i++) {
        t\lceil m++\rceil = '\#';
        t[m++] = s[i];
        p[i] = 0;
    t\lceil m++\rceil = '\#';
    int c = 0, r = 0;
    for (int i = 0; i < m; i++) {
        int i_{-} = 2 * c - i;
        p[i] = r > i ? min(r-i, p[i]) : i & 1;
        while (0 \le i-p[i]-1 \& i+p[i]+1 < m \& t[i-p[i]-1] == t[i+p[i]+1])
            p[i] += 2;
        if (i + p[i] > r) {
            c = i;
            r = i + p[i];
        }
```

```
Aho-Corasick
Complexity: < O(|S|), O(sum(|Si|)), O(|S|) >
struct Node {
    map<char, int> adj;
    int fail:
    pii match;
    int next;
    void init() {
        adj.clear();
        fail = -1;
       match = mp(-1, -1);
        next = -1:
   int getChild(const char &c) {
        map<char, int>::iterator it = adj.find(c);
        if (it != adj.end())
            return it->second;
        return -1;
};
int qntNodes, qntPatts;
Node trie[MAX];
void init() {
    trie[0].init();
    qntNodes = 1;
    antPatts = 0:
}
void addWord(const char *word) {
    int node = 0, aux = -1;
    for (int i = 0; word[i]; i++) {
        aux = trie[node].getChild(word[i]);
        if (aux == -1) {
            trie[qntNodes].init();
            aux = qntNodes++;
            trie[node].adj[word[i]] = aux;
        node = aux;
    trie[node].match = mp(qntPatts++, strlen(word));
}
```

```
void build() {
    queue<int> q:
    map<char, int>::iterator it;
    trie[0].fail = -1;
    a.push(0);
    while (!q.empty()) {
        int u = q.front();
        q.pop();
        for (it = trie[u].adj.begin(); it != trie[u].adj.end(); it++) {
            int v = it->second;
            char c = it->first;
            q.push(v);
            int f = trie[u].fail;
            while (f \ge 0 \& trie[f].getChild(c) == -1)
                f = trie[f].fail;
            f = f >= 0 ? trie[f].getChild(c) : 0;
            trie[v].fail = f;
            trie[v].next = trie[f].match.fi >= 0 ? f : trie[f].next;
       }
void search(const char *text) {
    int node = 0;
    for (int i = 0; text\lceil i \rceil; i++) {
        while (node >= 0 && trie[node].getChild(text[i]) == -1)
            node = trie[node].fail;
        node = node >= 0 ? trie[node].getChild(text[i]) : 0;
        int aux = node:
        while (aux >= 0) {
            if (trie[aux].match.fi >= 0) {
                // do something with the match
                printf("patt: %d, pos: %d\n",
                        trie[aux].match.fi,
                        i - trie[aux].match.se + 1);
            aux = trie[aux].next;
       }
```

```
Suffix Array and Longest Common Prefix
Complexity: < O(N logN), O(N) >
//Output:
// pos = The suffix array. Contains the n suffixes of str sorted in
          lexicographical order. Each suffix is represented as a
//
          single integer (the position of str where it starts).
// rank = The inverse of the suffix array.
          rank[i] = the index of the suffix <math>str[i..n) in the pos array.
//
          (In other words, pos[i] = k \ll rank[k] = i)
          With this array, you can compare two suffixes in O(1):
//
          Suffix str[i..n) is smaller than str[i..n) iff rank[i] < rank[i]
//
int n: // length of the string
char str[MAXN];
int rank[MAXN], pos[MAXN], cnt[MAXN], next[MAXN];
bool bh[MAXN], b2h[MAXN];
bool cmp(int a, int b) {
    return str[a] < str[b];</pre>
}
void suffix_array() {
    for (int i = 0; i < n; i++)
        pos[i] = i;
    sort(pos, pos+n, cmp);
    for (int i = 0; i < n; i++) {
        bh[i] = (i == 0 \mid | str[pos[i]] != str[pos[i-1]]);
        b2h[i] = 0;
   }
    for (int h = 1; h < n; h <<= 1) {
        int buckets = 0;
        for (int i = 0, j; i < n; i = j) {
            i = i + 1;
            while (j < n \&\& !bh[j])
                j++;
            next[i] = j;
            buckets++;
        if (buckets == n)
            break;
        for (int i = 0; i < n; i = next[i]) {
            cnt[i] = 0;
            for (int j = i; j < next[i]; j++)
                rank[pos[j]] = i;
        }
```

```
cnt[rank[n-h]]++;
        b2h[rank[n-h]] = 1;
        for (int i = 0; i < n; i = next[i]) {
            for (int j = i; j < next[i]; j++) {
                int s = pos[j] - h;
                if (s >= 0) {
                    int head = rank[s];
                    rank[s] = head + cnt[head]++;
                    b2h[rank[s]] = 1;
            for (int j = i; j < next[i]; j++) {
                int s = pos[j] - h;
                if (s >= 0 \&\& b2h[rank[s]]) {
                    for (int k = rank[s] + 1; !bh[k] \&\& b2h[k]; k++)
                        b2h\lceil k\rceil = 0;
            }
        }
        for (int i = 0; i < n; i++) {
            pos[rank[i]] = i;
            bh[i] |= b2h[i];
        }
    for (int i = 0; i < n; i++)
        rank[pos[i]] = i;
int height[MAXN];
void getHeight() {
    height[0] = 0;
    for (int i = 0, h = 0; i < n; i++) {
        if (rank[i] > 0) {
            int j = pos[rank[i] - 1];
            while (i + h < n \&\& j + h < n \&\& str[i+h] == str[j+h])
                h++;
            height[rank[i]] = h;
            if (h > 0)
                h--;
       }
```

Dynamic Programming

```
Optimal Binary Search Tree
Complexity: O(N^3)
int n, p[MAXN];
int c[MAXN][MAXN], f[MAXN][MAXN], r[MAXN][MAXN];
void obst() {
    for (int i = 1; i <= n; i++)
        c[i][i-1] = 0;
    c[n+1][n] = 0;
    for (int i = 1; i <= n; i++) {
        c[i][i] = p[i];
        f[i][i] = p[i];
        r[i][i] = i;
   }
    for (int d = 1; d < n; d++) {
        for (int i = 1; i \le n-d; i++) {
            int j = i+d;
            c[i][j] = INF;
            f[i][j] = f[i][j-1] + p[j];
            int rmin = r[i][j-1], rmax = r[i+1][j];
            for (int k = rmin; k \leftarrow rmax; k++) {
                int t = c[i][k-1] + c[k+1][j];
               if (t < c[i][j]) {
                    c[i][j] = t;
                    r[i][j] = k;
               }
           c[i][j] += f[i][j];
   }
}
```

```
Longest Increasing Subsequence
Complexity: O(N logN)
int n, m, a[MAXN], b[MAXN], p[MAXN];
void lis() {
    int u, v;
    b\lceil m++\rceil = 0;
    for (int \ i' = 1; \ i < n; \ i++) {
        if (a\lceil b\lceil m-1\rceil\rceil < a\lceil i\rceil) {
            p[i] = b[m-1];
            b[m++] = i;
             continue;
        for (u = 0, v = m-1; u < v;)
            int c = (u + v)/2;
            if (a[b[c]] < a[i])
                 u = c + 1;
            else
                 V = C;
        if (a[i] < a[b[u]]) {
            if (u > 0)
                 p[i] = b[u-1];
            b[u] = i;
        }
    for (u = m, v = b[m-1]; u--; v = p[v]) {
        b[u] = v;
```

Longest Common Increasing Subsequence

```
Complexity: O(N^2)
int n, m, a[MAXN], b[MAXN];
int c[MAXN], prev[MAXN], seq[MAXN];
void lcis() {
    for (int j = 0; j < m; j++)
        c[j] = 0;
    for (int i = 0; i < n; i++) {
        int actual = 0, last = -1;
        for (int j = 0; j < m; j++) {
            if (a[i] == b[j] \&\& actual+1 > c[j]) {
                c[j] = actual+1;
                prev[j] = last;
            } else if (a[i] > b[j] && actual < c[j]) {</pre>
                actual = c[j];
                last = j;
        }
    int length = 0, index = -1;
    for (int j = 0; j < m; j++) {
        if (c[j] > length) {
            length = c[j];
            index = j;
    int len = length;
    while (index != -1) {
        seq[--len] = b[index];
        index = prev[index];
    printf("length: %d\n", length);
    for (int i = 0; i < length; i++)
        printf("%d ", seq[i]);
    printf("\n");
}
```

```
Weighted Activity Selection
```

```
Complexity: O(N logN)
#include <cstdio>
#include <algorithm>
using namespace std;
#define MAXN 10005
struct Event {
    int b, e, w;
    Event () {}
    Event (int b, int e, int w) : b(b), e(e), w(w) {}
    bool operator< (const Event& o) const {</pre>
        return e != o.e ? e < o.e ? b < o.b;
};
int n;
Event e[MAXN];
int dp[MAXN];
int main() {
    scanf("%d", &n);
    e[0] = Event(0, 0, 0);
    for (int i = 1; i <= n; i++)
        scanf("%d %d %d", &e[i].b, &e[i].e, &e[i].w);
    sort(e+1, e+n+1);
    dp[0] = 0;
    for (int i = 1; i <= n; i++) {
        int lo = 0, hi = i-1;
        while (lo < hi) {</pre>
            int mid = (lo + hi + 1)/2;
            if (e[mid].e > e[i].b)
                hi = mid - 1;
            else
                lo = mid:
        dp[i] = max(dp[i-1], e[i].w + dp[lo]);
    printf("Max weight: %d\n", dp[n]);
```

Data Structure

```
Segment Tree with Lazy Propagation
Complexity: < O(N), O(logN) >
#define left(x) ((x) \ll 1)
#define right(x) (left(x) + 1)
ll tree[4*MAXN], lazy[4*MAXN];
void propagate(int node, int lo, int hi) {
    tree[node] += lazy[node] * (hi-lo+1);
   if (lo != hi) {
        lazy[left(node)] += lazy[node];
        lazy[right(node)] += lazy[node];
    lazy[node] = 0:
}
void update(int node, int lo, int hi, int i, int j, int val) {
   if (i <= lo && hi <= j) {
        lazy[node] += val;
        return;
   int mid = (lo + hi)/2;
   if (i <= mid)
        update(left(node), lo, mid, i, j, val);
   if (j > mid)
        update(right(node), mid+1, hi, i, j, val);
    propagate(left(node), lo, mid);
    propagate(right(node), mid+1, hi);
    tree[node] = tree[left(node)] + tree[right(node)];
}
ll query(int node, int lo, int hi, int i, int j) {
    propagate(node, lo, hi);
   if (i <= lo && hi <= j)
        return tree[node];
   ll ret = 0:
    int mid = (lo + hi)/2:
   if (i <= mid)
        ret = query(left(node), lo, mid, i, j);
   if (i > mid)
        ret += query(right(node), mid+1, hi, i, j);
    return ret;
}
```

Geometry

```
Template
struct Point {
    double x, y;
    Point() {}
    Point(double x, double y) : x(x), y(y) {}
    Point operator+ (const Point &o) const { return Point(x + o.x, y + o.y); }
    Point operator- (const Point &o) const { return Point(x - o.x, y - o.y); }
    Point operator* (const double &o) const { return Point(x * o, y * o); }
    Point operator/ (const double &o) const { return Point(x / o, y / o); }
    double operator* (const Point &o) const { return x * o.x + y * o.y; }
    double operator% (const Point &o) const { return x * o.y - o.x * y; }
    bool operator== (const Point &o) const {
        return cmp_double(x, o.x) == 0 \& cmp_double(y, o.y) == 0;
    bool operator< (const Point &o) const {</pre>
        return x != o.x ? x < o.x : y < o.y;
};
typedef Point Vector;
double abs(Point p) {
    return sqrt(p * p);
}
Vector norm(Vector v) {
    return v / abs(v);
double ccw(Point p, Point q, Point r) {
    return (q - p) \% (r - p);
}
struct Seament {
    Point p, q;
    Segment() {}
    Segment(Point p, Point q) : p(p), q(q) {}
}:
```

```
bool in seament(Point p. Seament s) {
    double t;
    Vector v = s.q - s.p;
   if (cmp\_double(v.x, 0) != 0)
        t = (p.x - s.p.x) / v.x:
    else
        t = (p.y - s.p.y) / v.y;
   return cmp_double(t, 0) >= 0 && cmp_double(t, 1) <= 0 && s.p + v * t == p;
}
struct Line {
   Vector v:
   Point p;
   int a, b, c;
   void init() {
        a = -v.y;
        b = v.x;
        c = a * p.x + b * p.y;
        int d = abs(\_qcd(a, \__qcd(b, c)));
        if (d != 1)
            a /= d, b /= d, c /= d;
        if (a < 0)
            a = -a, b = -b, c = -c;
        else if (a == 0 \&\& b < 0)
            b = -b, c = -c;
   }
   Line() {}
    Line(Point p, Point q) : v(q-p), p(p) {
        init();
    Line(Segment s) : v(s.q-s.p), p(p) {
        init();
   Point operator() (double t) const { return p + v * t; }
   Vector normal() {
        return Vector(-v.y, v.x);
};
pair<double, double> line_intersection(Line a, Line b) {
    double den = a.v \% b.v;
   if (den == 0)
        return make_pair(inf, inf);
    double t = -(b.v \% (b.p - a.p)) / den;
    double s = -(a.v \% (b.p - a.p)) / den;
    return make_pair(t, s);
}
```

```
Point segment_intersection(Segment a, Segment b) {
    Line la = Line(a), lb = Line(b);
    pair<double, double> pdd = line_intersection(la, lb);
    double t = pdd.first. s = pdd.second:
    if (t == inf) {
        if (in_segment(b.p, a))
            return b.p:
       if (in_segment(b.q, a))
            return b.q;
       if (in_segment(a.p, b))
           return a.p;
       if (in_segment(a.q, b))
           return a.a;
        return Point(inf, inf);
    if (cmp\_double(t, 0) < 0 \mid | cmp\_double(t, 1) > 0)
        return Point(inf, inf);
    if (cmp\_double(s, 0) < 0 \mid | cmp\_double(s, 1) > 0)
        return Point(inf, inf);
    return la(t);
}
double distPointToLine(Point p, Line l) {
    Vector n = 1.normal();
    return (l.p - p) * n / abs(n);
struct Circle {
    Point p;
    double r;
    Circle() {}
    Circle(Point p, double r) : p(p), r(r) {}
};
bool in_circle(const Circle &c, const Point &p) {
    return cmp_double(abs(c.p - p), c.r) <= 0;
Point circumcenter(Point p, Point q, Point r) {
    Point a = p - r, b = a - r, c = Point(a*(p+r)/2, b*(a+r)/2):
    return Point(c % Point(a.y, b.y), Point(a.x, b.x) % c)/(a % b);
Point incenter(Point p, Point q, Point r) {
    double a = abs(r - q), b = abs(r - p), c = abs(q - p);
    return (p * a + a * b + r * c) / (a + b + c);
```

```
Monotone Chain Convex Hull
Complexity: O(N loaN)
int n. k:
Point p[MAXN], h[MAXN];
void convex_hull() {
    sort(p, p+n);
    k = 0;
    h[k++] = p[0];
    for (int i = 1; i < n; i++) {
        if (i != n-1 && ccw(p[0], p[n-1], p[i]) >= 0) continue;
        while (k > 1 \&\& ccw(h[k-2], h[k-1], p[i]) <= 0) k--;
        h \lceil k + + \rceil = p \lceil i \rceil:
    for (int i = n-2, lim = k; i >= 0; i--) {
        if (i != 0 && ccw(p[n-1], p[0], p[i]) >= 0) continue;
        while (k > \lim \& ccw(h\lceil k-2\rceil, h\lceil k-1\rceil, p\lceil i\rceil) \le 0) k--;
        h[k++] = p[i];
}
Smallest Enclosing Circle
Complexity: O(N^2)
int n;
Point p[MAXN];
Circle spanning_circle() {
    random_shuffle(p, p+n);
    Circle c(Point(), -1);
    for (int i = 0; i < n; i++) if (!in_circle(c, p[i])) {
        c = Circle(p[i], 0);
        for (int j = 0; j < i; j++) if (!in_circle(c, p[j])) {
            c = Circle((p[i] + p[j])/2, abs(p[i] - p[j])/2);
            for (int k = 0; k < j; k++) if (!in_circle(c, p[k])) {
                 Point o = circumcenter(p[i], p[j], p[k]);
                 c = Circle(o, abs(o - p[k]));
        }
    return c;
```

```
Closest Pair of Points
Complexity: O(N loaN)
#include <cstdio>
#include <cmath>
#include <alaorithm>
#include <set>
using namespace std;
struct Point {
    int x, y;
    Point(int x = 0, int y = 0) : x(x), y(y) {}
    Point operator- (const Point &o) const { return Point(x - o.x, y - o.y); }
    int operator* (const Point &o) const { return x * o.x + v * o.v: }
    bool operator< (const Point &o) const {</pre>
        return y != o.y ? y < o.y : x < o.x;
};
bool cmpx(const Point &p, const Point &q) {
    return p.x != q.x ? p.x < q.x : p.y < q.y;
double abs(const Point &p) {
    return sart(p * p);
int main() {
    int n;
    Point pnts[MAXN];
    set<Point> box;
    set<Point>::iterator it;
    scanf("%d", &n);
    for (int i = 0; i < n; i++)
        scanf("%d %d", &pnts[i].x, &pnts[i].y);
    sort(pnts, pnts+n, cmpx);
    double best = inf;
    box.insert(pnts[0]):
    for (int i = 1, j = 0; i < n; i++) {
        while (j < i && pnts[i].x - pnts[j].x > best)
            box.erase(pnts[j++]);
        for (it = box.lower_bound(Point(pnts[i].x-best, pnts[i].y-best));
            it != box.end() && it->y <= pnts[i].y + best; it++) {
            best = min(best, abs(pnts[i] - *it));
        box.insert(pnts[i]);
    printf("%.2lf\n", best);
```

Math

```
Sieve, primality, factorization, phi
int np, p[MAXN];
int lp[MAXN];
void sieve(int n) {
    for (int i = 2; i < n; i++)
        lp[i] = i;
    for (int i = 4; i < n; i += 2)
        lp[i] = 2;
    for (int i = 3; i*i < n; i += 2)
        if (lp[i] == i)
            for (int j = i*i; j < n; j += i)
                lp[j] = i;
    np = 0;
    p\lceil np++\rceil = 2;
    for (int i = 3; i < n; i += 2)
        if (lp[i] == i)
            p[np++] = i;
}
int nf, f[MAXN], e[MAXN];
void factor(int n) {
    nf = 0;
    for (int i = 0; n != 1 && p[i]*p[i] <= n; i++) {
        if (n \% p[i] == 0) {
            f[nf] = p[i];
            e[nf] = 1;
            n /= p[i];
            while (n \% p[i] == 0) {
                e[nf]++;
                n \neq p[i];
            nf++;
        }
    if (n != 1) {
        f[nf] = n;
        e[nf] = 1;
        nf++;
   }
}
```

```
int _phi(int n) {
    int ret = 1;
    for (int i = 0; n != 1 && p[i]*p[i] <= n; i++) {
        if (n \% p[i] == 0) {
            int pk = p[i];
            n \neq p[i];
            while (n \% p[i] == 0) {
                pk *= p[i];
                n \neq p[i];
           ret *= pk - pk/p[i];
       }
    if (n != 1)
        ret *= n-1;
    return ret;
int phi[MAXN];
void build_phi(int n) {
    for (int i = 0; i < n; i++)
        phi[i] = i;
    for (int i = 2; i < n; i++) if (phi[i] == i)
        for (int j = i; j < n; j += i)
            phi[j] = phi[j] / i * (i-1);
```

Chinese Remainder Algorithm

```
#include <cstdio>
#include <algorithm>
using namespace std;
const int MAXN = 100010;
typedef pair<int, int> tpii;
struct teq {
   int r, n; // x = r (mod n)
int ant:
teq eqs[MAXN];
tpii eqcd(int a, int b) {
   int x = 0, last x = 1, auxx;
   int y = 1, lasty = 0, auxy;
   while (b) {
        int q = a / b, r = a % b;
        a = b, b = r;
        auxx = x;
        x = lastx - q*x, lastx = auxx;
        auxy = y;
       y = lasty - q*y, lasty = auxy;
    return make_pair(lastx, lasty);
}
int chinese_remainder_algorithm() {
    int beta, sum = 0, n = 1;
    for (int i = 0; i < qnt; i++)
        n *= eqs[i].n;
   for (int i = 0; i < qnt; i++) {
        beta = egcd(eqs[i].n, n/eqs[i].n).second;
        while (beta < 0)
            beta += eqs[i].n;
        sum += (eqs[i].r * beta * n/eqs[i].n) % n;
    return sum;
}
int main() {
    scanf("%d", &qnt);
    for (int i = 0; i < qnt; i++)
        scanf("%d %d", &eqs[i].r, &eqs[i].n);
    printf("%d\n", chinese_remainder_algorithm());
}
```

FFT

```
typedef complex<long double> pt;
pt tmp[1<<20];
void fft(pt *in, int sz, bool inv) {
    if (sz == 1)
        return;
    for (int i = 0, j = 0, h = sz >> 1; i < sz; i += 2, j++) {
        in[j] = in[i];
        tmp[h+j] = in[i+1];
    for (int i = sz >> 1; i < sz; i++)
        in[i] = tmp[i];
    sz >>= 1;
    pt *even = in, *odd = in + sz;
    fft(even, sz, inv);
    fft(odd, sz, inv);
    long double p = (inv ? -1 : 1) * M_PI / sz;
    pt w = pt(cosl(p), sinl(p)), w_i = 1;
    for (int i = 0; i < sz; i++) {
        pt conv = w_i * odd[i];
        odd[i] = even[i] - conv;
        even[i] += conv;
        w_i *= w;
   }
}
```

Polynomial

```
struct Poly {
    int n;
    double a[MAXN];
    Poly(int n = \emptyset): n(n) { memset(a, \emptyset, sizeof(a)); }
    Poly(const Poly &o): n(o.n) { memcpy(a, o.a, sizeof(a)); }
    const double& operator[] (int i) const { return a[i]; }
    double& operator[] (int i) { return a[i]; }
    double operator() (double x) const {
        double ret = 0:
        for (int i = n; i >= 0; i--)
            ret = ret * x + a[i];
        return ret;
    Poly operator+ (const Poly &o) const {
        Poly ret = o;
        for (int i = 0; i \le n; i++)
            ret[i] += a[i];
        ret.n = max(n, o.n);
        return ret;
    Polv operator- (const Polv &o) const {
        Poly ret = o;
        for (int i = 0; i \le n; i++)
            ret[i] -= a[i];
        ret.n = max(n, o.n);
        return ret;
    Poly operator* (const Poly &o) const {
        Poly ret(n + o.n);
        for (int i = 0; i <= n; i++)
            for (int j = 0; j <= o.n; j++)
                ret[i+j] += a[i] * o[j];
        return ret;
    }
};
```

```
Poly fastMult(const Poly &p, const Poly &q) {
    int sz = 1 \ll (32 - \_builtin\_clz(p.n + q.n + 1));
    pt pin[sz], qin[sz];
    for (int i = 0; i < sz; i++) {
        if (i \ll p.n)
            pin[i] = p[i];
        else
            pin[i] = 0;
        if (i \ll q.n)
            qin[i] = q[i];
        else
            qin[i] = 0;
    fft(pin, sz, 0);
    fft(qin, sz, 0);
    for (int i = 0; i < sz; i++)
        pin[i] *= qin[i];
    fft(pin, sz, 1);
    Poly ret(p.n + q.n);
    for (int i = 0; i \leftarrow ret.n; i++)
        ret[i] = pin[i].real() / sz;
    while (ret.n > 0 && cmp(ret[ret.n], 0) == 0)
        ret.n--:
    return ret;
Poly diff(const Poly &p) {
    Poly ret(p.n-1);
    for (int i = 1; i \le p.n; i++)
        ret[i-1] = i * p[i];
    return ret;
}
pair<Poly, double> ruffini(const Poly &p, double x) {
    if (p.n == 0)
        return make_pair(Poly(), 0);
    Poly ret(p.n-1);
    for (int i = p.n; i > 0; i--)
        ret[i-1] = ret[i] * x + p[i];
    return make_pair(ret, ret[0] * x + p[0]);
```

```
* Find a root in range [lo, hi] assuming that exists only one root in [lo,
 * pair::second is true if exists a root in the given range or false otherwise
 * pair::first is the root if pair::second is true or 0 if false
pair<double, int> findRoot(const Poly &p, double lo, double hi) {
   if (cmp(p(lo), 0) == 0)
        return make_pair(lo, 1);
   if (cmp(p(hi), 0) == 0)
        return make_pair(hi, 1);
   if (cmp(p(lo), 0) = cmp(p(hi), 0))
        return make_pair(0, 0);
   if (cmp(p(lo), p(hi)) > 0)
        swap(lo, hi);
   while (cmp(lo, hi) != 0) {
        double mid = (lo + hi) / 2;
        double val = p(mid);
        if (cmp(val, 0) == 0)
            lo = hi = mid;
        else if (cmp(val, 0) < 0)
            lo = mid;
        else
            hi = mid;
    return make_pair(lo, 1);
}
/**
 * Return a vector of all real roots with their multiplicity in ascending
order
*/
vector<double> roots(const Poly &p) {
   vector<double> ret;
   if (p.n == 1) {
        ret.push_back(-p[0] / p[1]);
   else {
        vector<double> r = roots(diff(p));
        r.push_back(-MAXX);
        r.push_back(MAXX);
        sort(r.begin(), r.end());
        for (int i = 0, j = 1; j < (int) r.size(); i++, j++) {
            pair<double, int> pr = findRoot(p, r[i], r[j]);
            if (pr.second)
                ret.push_back(pr.first);
       }
    return ret;
}
```

Bignum

```
#include <cstring>
#include <algorithm>
#include <limits>
using namespace std;
typedef long long 11;
typedef unsigned long long ull;
const int MAXD = 1005. DIG = 9. BASE = 10000000000:
const ull BOUND = numeric_limits <ull> :: max() - (ull) BASE * BASE;
struct bignum
    int D, digits[MAXD / DIG + 2];
    int sign;
    inline void trim () {
        while (D > 1 \&\& digits[D - 1] == 0)
            D--;
    inline void init (ll x) {
        memset(digits, 0, sizeof(digits));
        D = 0;
        if (x < 0) {
            sign = -1;
            X = -X;
        }
        else {
            sign = 1;
        }
        do {
            digits\Gamma D++1 = x \% BASE;
            x /= BASE:
        } while (x > 0);
    inline bignum (ll x) {
        init(x);
    inline bignum (int x = 0) {
        init(x);
```

```
inline bianum (char *s) {
    memset(digits, 0, sizeof(digits));
    if (s[0] == '-') {
        sign = -1;
        S++;
    else {
        sign = 1;
    int len = strlen(s), first = (len + DIG - 1) % DIG + 1;
    D = (len + DIG - 1) / DIG;
    for (int i = 0: i < first: i++)
        digits[D - 1] = digits[D - 1] * 10 + s[i] - '0';
    for (int i = first, d = D - 2; i < len; i += DIG, d--)
        for (int j = i; j < i + DIG; j++)
            digits[d] = digits[d] * 10 + s[j] - '0';
    trim();
inline char *str () {
    trim();
    char *buf = new char [DIG * D + 2];
    int pos = 0, d = digits[D - 1];
    if (sign == -1)
        buf[pos++] = '-':
    do {
        buf[pos++] = d \% 10 + '0';
        d /= 10;
    } while (d > 0);
    reverse(buf + (sign == -1 ? 1 : 0), buf + pos);
    for (int i = D - 2; i >= 0; i--, pos += DIG)
        for (int j = DIG - 1, t = digits[i]; <math>j >= 0; j--) {
            buf[pos + i] = t \% 10 + '0':
            t /= 10;
        }
    buf[pos] = '\0':
    return buf;
```

```
inline bool operator < (const bignum &o) const {</pre>
    if (sign != o.sign)
        return sign < o.sign;</pre>
   if (D != o.D)
        return sign * D < o.sign * o.D;
    for (int i = D - 1; i >= 0; i--)
        if (digits[i] != o.digits[i])
            return sign * digits[i] < o.sign * o.digits[i];
    return false;
inline bool operator > (const bignum &o ) const {
    if (sign != o.sign)
        return sign > o.sign;
   if (D != o.D)
        return sign * D > o.sign * o.D;
    for (int i = D - 1; i >= 0; i--)
        if (digits[i] != o.digits[i])
            return sign * digits[i] > o.sign * o.digits[i];
    return false;
inline bool operator == (const bignum &o) const {
   if (sign != o.sign)
        return false;
   if (D != o.D)
        return false;
    for (int i = 0; i < D; i++)
        if (digits[i] != o.digits[i])
            return false;
    return true;
```

```
inline bignum operator << (int p) const {</pre>
    bignum temp;
    temp.D = D + p;
    for (int i = 0; i < D; i++)
        temp.digits[i + p] = digits[i];
    for (int i = 0; i < p; i++)
        temp.digits[i] = 0;
    return temp;
}
inline bignum operator >> (int p) const {
    bianum temp:
    temp.D = D - p;
    for (int i = 0; i < D - p; i++)
        temp.digits[i] = digits[i + p];
    for (int i = D - p; i < D; i++)
        temp.digits[i] = 0;
    return temp;
}
inline bignum range (int a, int b) const {
    bianum temp = 0:
    temp.D = b - a;
    for (int i = 0; i < temp.D; i++)
        temp.digits[i] = digits[i + a];
    return temp;
inline bignum abs () const {
    bignum temp = *this;
    temp.sign = 1;
    return temp;
}
```

```
inline bignum operator + (const bignum &o) const {
    if (sign != o.sign) {
        if (sign == 1)
            return *this - o.abs();
            return o - this->abs();
    bignum sum = o;
    int carry = 0;
    for (sum.D = 0; sum.D < D \mid | carry > 0; sum.D++) {
        sum.digits[sum.D] += (sum.D < D ? digits[sum.D] : 0) + carry;</pre>
        carrv = 0:
        if (sum.digits[sum.D] >= BASE) {
            sum.digits[sum.D] -= BASE;
            carry = 1;
    }
    sum.D = max(sum.D, o.D);
    sum.trim();
    return sum;
inline bignum operator - (const bignum &o) const {
    if (sign != o.sign) {
       if (sign == 1)
            return *this + o.abs();
            return -(this->abs() + o);
    else if (sign == -1) {
        return o.abs() - this->abs();
   }
    bignum diff, temp;
    if (o > *this) {
        diff = o;
        diff.sian = -1:
        temp = *this;
   }
    else {
        diff = *this:
        temp = o;
   }
```

```
for (int i = 0, carry = 0; i < temp.D \mid\mid carry > 0; i++) {
        diff.digits[i] -= (i < temp.D ? temp.digits[i] : 0) + carry;</pre>
        carrv = 0:
        if (diff.digits[i] < 0) {
            diff.digits[i] += BASE;
            carry = 1;
        }
    }
    diff.trim();
    return diff;
inline bianum operator - () const {
    bignum temp = *this;
    temp.sign = -temp.sign;
    return temp;
}
inline bignum operator * (const bignum &o) const {
    bignum prod = 0;
    ull sum = 0, carry = 0;
    for (prod.D = 0; prod.D < D + o.D - 1 || carry > 0; prod.D++) {
        sum = carry % BASE;
        carry /= BASE;
        for (int j = max(prod.D-o.D+1, \emptyset); j \le min(D-1, prod.D); j++) {
            sum += (ull) digits[j] * o.digits[prod.D - j];
            if (sum >= BOUND) {
                carry += sum / BASE;
                sum %= BASE;
        }
        carry += sum / BASE;
        prod.digits[prod.D] = sum % BASE;
    prod.sign = sign * o.sign;
    prod.trim();
    return prod;
}
```

```
inline double double div (const bianum &o) const {
   double val = 0, oval = 0;
   int num = 0, onum = 0;
    for (int i = D - 1; i >= max(D - 3, 0); i --, num++)
       val = val * BASE + digits[i];
   for (int i = 0.D - 1; i >= max(0.D - 3, 0); i--, onum++)
       oval = oval * BASE + o.digits[i];
   return sign * o.sign * val / oval * (D - num > o.D - onum ? BASE : 1);
inline pair<br/>bignum, bignum> divmod (const bignum &o) const {
   if (sian != o.sian) {
       pair<bignum, bignum> p = (this->abs()).divmod(o.abs());
       p.first.sign = -1;
       p.second.sign = sign;
       return p;
   }
   else if (sign == -1) {
       pair<bignum, bignum> p = (this->abs()).divmod(o.abs());
       p.second.sign = sign;
       return p;
   }
   bignum quot = 0, rem = *this, temp;
   for (int i = D - o.D; i >= 0; i--) {
       temp = rem.range(i, rem.D);
       int div = (int) temp.double_div(o);
       bignum mult = o * div;
       while (div > 0 && temp < mult) {
            mult = mult - o;
            div--:
       }
       while (div + 1 < BASE \&\& !(temp < mult + o)) {
            mult = mult + o;
            div++;
       rem = rem - (o * div << i):
       if (div > 0) {
            quot.digits[i] = div;
            quot.D = max(quot.D, i + 1);
   }
```

```
quot.trim();
        rem.trim();
        return make_pair(quot, rem);
    inline bignum operator / (const bignum &o) const {
        return divmod(o).first;
    inline bignum operator % (const bignum &o) const {
        return divmod(o).second;
    inline bignum power (int exp) const {
        bignum p = 1, temp = *this;
        while (exp > 0) {
            if (exp \& 1) p = p * temp;
            if (exp > 1) temp = temp * temp;
            exp >>= 1;
        return p;
    }
};
inline bignum gcd (bignum a, bignum b) {
    bignum t;
    while (!(b == 0)) {
        t = a \% b;
        a = b;
        b = t;
    return a;
}
```