



NagBody lectures: Complexity: analysis of algorithms

Mario Alberto Rodríguez-Meza

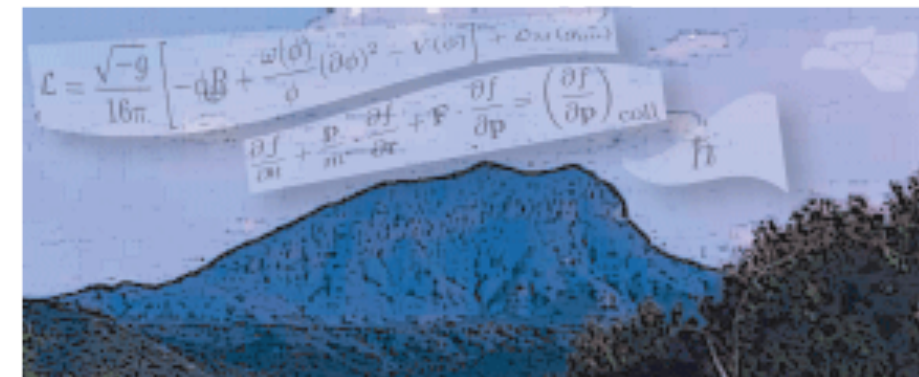
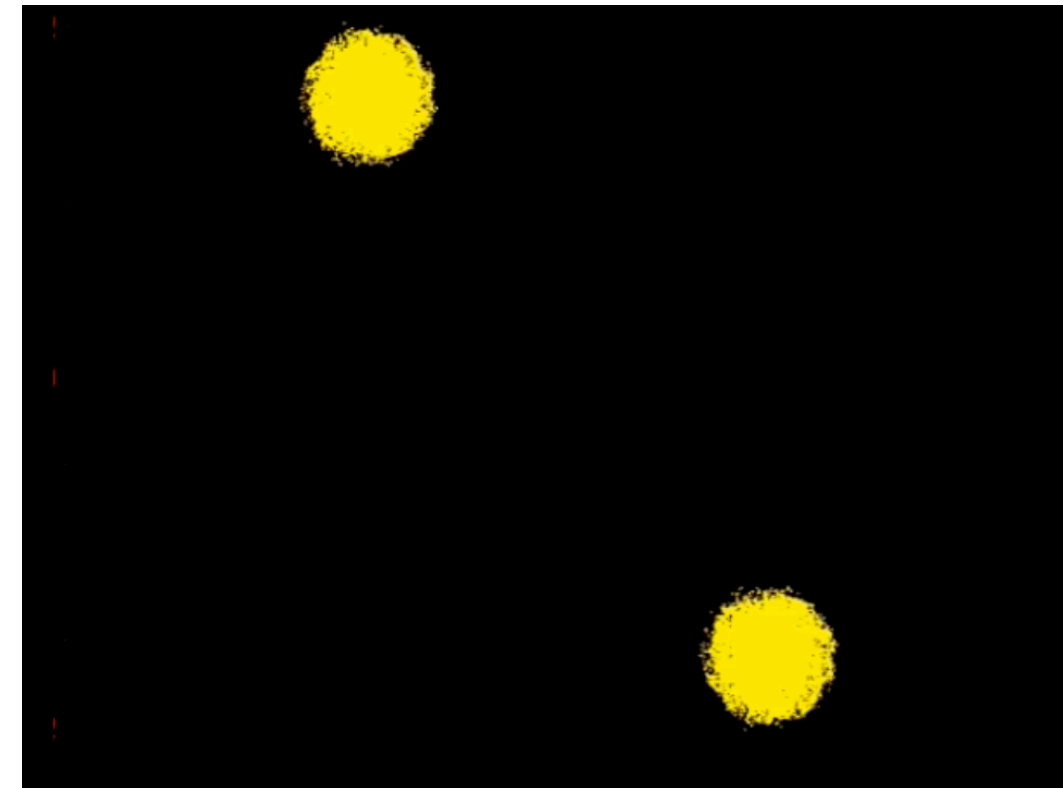
Instituto Nacional de Investigaciones Nucleares

Correo Electrónico: marioalberto.rodriguez@inin.gob.mx

<http://bitbucket.org/rodriguezmeza>

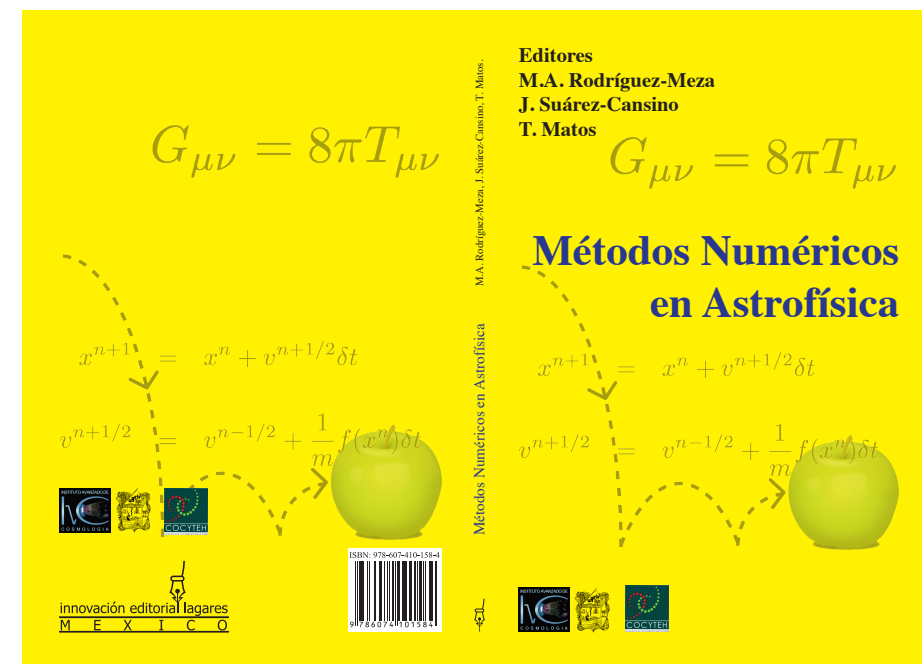
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quintessence
Group



References and material

- Cosmología numérica y estadística: NagBody kit (<http://bitbucket.org/rodriguezmeza>). Mario A. Rodríguez-Meza. And: https://github.com/rodriguezmeza/NagBody_lectures.git
- Métodos numéricos en astrofísica, capítulo I, Método de N-cuerpos en astrofísica. (https://www.researchgate.net/publication/316582859_Metodo_de_N-Cuerpos_en_Astrofisica)
- La estructura a gran escala del universo. Capítulo 22 en Travesuras cosmológicas de Einstein et al. https://www.researchgate.net/publication/316582400_La_estructura_a_gran_escaladel_universo_simulaciones_numericas
- https://www.researchgate.net/profile/Mario_Rodriguez-Meza
- https://www.researchgate.net/publication/314281416_Los_agujeros_negros_y_las_ondas_del_Dr_Einstein
- M.A. Rodríguez-Meza, Adv. Astron. 2012, 509682 (2012). arXiv: 1112.5201. (https://www.researchgate.net/publication/51967093_A_Scalar_Field_Dark_Matter_Model_and_Its_Role_in_the_Large-Scale_Structure_Formation_in_the_Universe)



Content: Complexity

- Worst-case analysis
- O -notation
- Computational complexity



Concept of Analysis of algorithms

- In computing it is important to understand how an algorithm perform.
- When we are designing an algorithm usually the aspect of most interest is how fast it will run.
- Storage is also important.
- Therefore we need a formal and a deterministic way of analyze algorithms.



Concept of

We will cover:

- **Worst-case analysis.** The metric by which most algorithms are compared. Other cases are: the average case and the best case.
- **O-notation.** It is a formal method to express the performance of an algorithm. It will be the upper bond of a function within a constant factor.
- **Computational complexity.** It is the growth rate of the resources (CPU time, storage...) an algorithm requires with respect to the size of the data it processes. The **O-notation** is a formal expression of an algorithm's complexity.



Concept of Worst-case analysis:

- Algorithms do not perform the same in all situations.
- There are three cases: *best* case, *worst* case and *average* case.
- Example: linear search. A natural but inefficient way of searching a list: traverse data from one end to the other.
- Best case: element we are looking is in the first place.
- Worst case: element we are logic is in the last place.
- Average case: element is somewhere in the middle.



Concept of Worst-case analysis:

- *Reasons why we chose Worst-case analysis:*
 - a. Many algorithms perform their worst case a large part of the time.
 - b. **Best case** is not very informative. For example, all searching algorithms behave almost the same in this case. Comparison among algorithms is an important matter.
 - c. Determine the **average-case** performance is not always easy. Or even what does it means “average”.
 - d. The **worst-case** gives us an upper bound on performance.



Concept of O -notation

- It is the **most common notation** used to express algorithms' performance.
- It is an **upper bound** of a function within a constant factor:
- If $g(n)$ is an upper bound of $f(n)$, then for some constant c it is possible to find a value of n , called it n_0 , for which any value of $n \geq n_0$ will result in $f(n) \leq c g(n)$.
- We **express performance of an algorithm** as a function of the size of the data it processes. Then $f(n)$ will be such a function and n the size of the data.
- We are only interested in the **algorithm's growth rate**, or the **order of the growth**. It will describe **how efficient an algorithm is** for arbitrary input.
- And **O -notation** is such a function.



Simple rules for *O*-notation

- We can ignore constant terms because as the value of n becomes larger and larger, eventually constant terms become insignificant: $T(n) = n + 50$.
- We can ignore constant multipliers of terms because they too will become insignificant as values of n increases. $T_1(n) = n^2$ and $T_2(n) = 10n$. Value of n has only be greater than 10 for T_1 to becomes greater than T_2 .
- We only consider the highest-order therm because, again as n increases, higher-order terms quickly outweigh the lower-order term ones. For example, $T(n) = n^2 + n$, for $n=1024$, the lesser order term of this expression constitutes less than 0.1% of the running time.



Simple rules for O -notation

- These ideas are formalized in the following rules:
 - a. Constant terms are expressed as $O(1)$. Regardless of the size of the data. formally: $O(c) = O(1)$.
 - b. Multiplicative constants are omitted. Formally stated, for a constant c : $O(cT) = cO(T) = O(1)$. If three tasks each run in time $T(n)$, the result is $O(3n)$.
 - c. Addition is performed by taking the maximum. Formally stated: $O(T_1) + O(T_2) = \text{Max}(O(T_1), O(T_2))$. For example: $T_1(n)=n$ and $T_2(n)=n^2$ describe two task executed sequentially, the result is $O(n) + O(n^2)$. Which simplifies to $O(n^2)$.
 - d. Multiplication is not changed but often is rewritten more compactly. Formally: $O(T_1) O(T_2) = O(T_1 T_2)$. For example, in a nested loop where the outer loop are described by T_1 and whose inner iterations by T_2 and if they both are equal to n , then the result is n^2 .



Concept of Computational complexity:

Homework: analyse nbody_n2 and gbsph codes..
Compare computing time.

- When speaking of the performance of an algorithm usually the aspect of interest is its *complexity*. The *grow rate* of the resources (time and storage: memory and disk) with respect to the size of the data.
- *O-notation* describes algorithm's complexity.
- We frequently can describe *worst-case complexity* by simply *inspecting its overall structure*. Other times it is helpful to employ techniques involving recurrences and summation formulas.
- To understand complexity it is important to *look at the way to surmise the resources* an algorithm will require.
- *CPU time* can be measured at each one of the major pieces that form the computational flow.



Conclusions:

Recursion. Part II

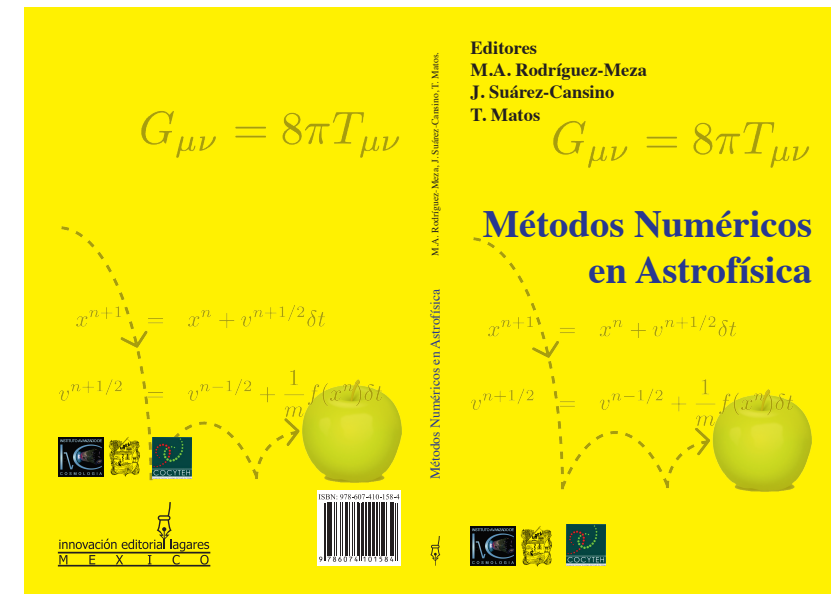
We have seen:

- Worst-case analysis. Against best and average cases.
- O -notation. Simple rules, examples and why it works.
- Computational complexity. CPU time, memory consuming and hard-disk storage.



References and material

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See you!

