Determinism and termination in the semantics of the Céu programming language

February 2, 2017

1 The original formulation

The semantics discussed in this section follows as much as possible the original semantics of Céu presented in [**?**]. Any deviations from original definitions are duly noted in the text.

1.1 Abstract syntax

The *abstract syntax* of Céu programs is given by the following grammar:

(1) `skip` precisa aparecer na gramática já que aprece nos programas em P.
(2) Atribuição agora aparece explicitamente na gramática. Expressões aritméticas e booleanas também estão na gramática mas a sua estrutura interna é omitida.

$p \in P ::=$	$\texttt{skip}$	do nothing
	$\mid v := a$	assignment
	$\mid \texttt{await}(e)$	await event
	$\mid \texttt{emit}(e)$	emit event
	$\mid \texttt{break}$	break innermost loop
	$\mid \texttt{if}\ b\ \texttt{then}\ p_1\ \texttt{else}\ p_2$	conditional
	$\mid p_1 ; p_2$	sequence
	$\mid \texttt{loop}\ p_1$	repetition
	$\mid p_1\ \texttt{and}\ p_2$	par/and
	$\mid p_1\ \texttt{or}\ p_2$	par/or
	$\mid \texttt{fin}\ p_1$	finalization
	$\mid \texttt{@awaiting}(e, n)$	awaiting e since reaction n
	$\mid \texttt{@emitting}(e, n)$	emitting e on stack level n
	$\mid p_1\ \texttt{@loop}\ p_2$	unwinded loop

where $n \in N$ is an integer, $v \in V$ is a memory location (variable) identifier, $e \in E$ is an event identifier, $a \in A$ is an arithmetic expression, $b \in B$ is a boolean expression, and p, p_1, $p_2 \in P$ are programs. We assume the usual structure for arithmetic and boolean expressions, and omit their definition.

1.2 The reaction inner-step relation

The *state* of a Céu program within a reaction is represented by a stack of events $\alpha = e_1 e_2 \ldots e_n \in E^*$ together with a memory map $m\colon v \to N \in \mathcal{M}$. A *configuration* is a 4-tuple $\langle p, \alpha, m, n \rangle \in \Delta$ that represents the situation of program p waiting to be evaluated in state $\langle \alpha, m \rangle$ and reaction n. Given an initial configuration, each small-step within a program reaction is determined by the reaction-inner-step relation $\rightarrow \in \Delta \times \Delta$ such that $\langle p, \alpha, m, n \rangle \rightarrow \langle p', \alpha', m', n \rangle$ iff a reaction inner-step of program p in state $\langle \alpha, m \rangle$ and reaction number n evaluates to a modified program p' and a modified state $\langle \alpha', m' \rangle$ in the same reaction (n). Since relation $\rightarrow$ can only relate configurations with the same n, we shall write $\langle p, \alpha, m \rangle \xrightarrow{n} \langle p', \alpha', m' \rangle$ for $\langle p, \alpha, m, n \rangle \rightarrow \langle p', \alpha', m', n \rangle$.

Relation $\rightarrow$ is defined inductively with the help of the auxiliary functions *eval*, *blocked*, and *clear*. The *eval* function evaluates arithmetic or boolean expressions on a given memory; we omit its definition and assume that such evaluation is deterministic and always terminates. The *blocked* function is a predicate that determines if all trails of a program p are blocked on a given event stack and reaction number. And the *clear* function extracts the body of `fin` blocks from a given program.

Definition 1.1. Function $\mathit{blocked}\colon P \times E^* \times N \rightarrow \{0, 1\}$ is defined inductively as follows.

$$
\begin{aligned}
\mathit{blocked}(\texttt{skip}, e\alpha, n) &= 0 \\
\mathit{blocked}(v := a, e\alpha, n) &= 0 \\
\mathit{blocked}(\texttt{await}(e'), e\alpha, n) &= 0 \\
\mathit{blocked}(\texttt{emit}(e'), e\alpha, n) &= 0 \\
\mathit{blocked}(\texttt{break}, e\alpha, n) &= 0 \\
\mathit{blocked}(\texttt{if}\, b\, \texttt{then}\, p_1\, \texttt{else}\, p_2, e\alpha, n) &= 0 \\
\mathit{blocked}(p_1; p_2, e\alpha, n) &= \mathit{blocked}(p_1, e\alpha, n) \\
\mathit{blocked}(\texttt{loop}\, p, e\alpha, n) &= 0 \\
\mathit{blocked}(p_1\, \texttt{and}\, p_2, e\alpha, n) &= \mathit{blocked}(p_1, e\alpha, n) \cdot \mathit{blocked}(p_2, e\alpha, n) \\
\mathit{blocked}(p_1\, \texttt{or}\, p_2, e\alpha, n) &= \mathit{blocked}(p_1, e\alpha, n) \cdot \mathit{blocked}(p_2, e\alpha, n) \\
\mathit{blocked}(\texttt{fin}\, p_1, e\alpha, n) &= 0 \\
\mathit{blocked}(\texttt{@awaiting}(e', n'), e\alpha, n) &= \begin{cases} 1 & \text{if } e \neq e' \text{ or } n = n' \\ 0 & \text{otherwise} \end{cases} \\
\mathit{blocked}(\texttt{@emitting}(n'), e\alpha, n) &= \begin{cases} 1 & \text{if } |e\alpha| \neq n' \\ 0 & \text{otherwise} \end{cases} \\
\mathit{blocked}(p_1\, \texttt{@loop}\, p_2, e\alpha, n) &= \mathit{blocked}(p_1, e\alpha, n)
\end{aligned}
$$

Definition 1.2. Function $\mathit{clear}\colon P \rightarrow P'$, where $P' = \{v := a\}$, is defined inductively as follows.

(1) Além de atribuições P' não deveria conter instruções if-else?.
(2) No artigo original, $\mathit{clear}(p_1; p_2)$ retorna apenas $\mathit{clear}(p_1)$. Achamos que isso está errado, já que apenas o primeiro `fin` seria considerado. (Confirmar com o Chicão).
(3) Adicionei $\mathit{clear}(\mathit{skip})$.

$$
\begin{aligned}
clear(\texttt{skip}) &= \texttt{skip}\\
clear(v := a) &= v := a\\
clear(\texttt{await}(e')) &= \texttt{skip}\\
clear(\texttt{emit}(e')) &= \texttt{skip}\\
clear(\texttt{break}) &= \texttt{skip}\\
clear(\texttt{if}\,v\,\texttt{then}\,p_1\,\texttt{else}\,p_2) &= \texttt{skip}\\
clear(p_1; p_2) &= clear(p_1); clear(p_2)\\
clear(\texttt{loop}\,p) &= clear(p)\\
clear(p_1\,\texttt{and}\,p_2) &= clear(p_1); clear(p_2)\\
clear(p_1\,\texttt{or}\,p_2) &= clear(p_1); clear(p_2)\\
clear(\texttt{fin}\,p) &= p\\
clear(\texttt{@awaiting}(e', n')) &= \texttt{skip}\\
clear(\texttt{@emitting}(n')) &= \texttt{skip}\\
clear(p_1\,\texttt{@loop}\,p_2) &= clear(p_1)
\end{aligned}
$$

Definition 1.3 (Reaction inner-step)**.** Relation $\rightarrow \subseteq \Delta \times \Delta$ is defined inductively as follows.

(1) Adicionamos o mapa de memória (m) à configuração e regras explícitas para atribuição. A avaliação de expressões aritméticas e booleanas está encapsulada na função *eval*.
(2) Adicionamos regras para consumir instruções `skip`.
(3) Adicionamos condições que garantem que a cada passo apenas uma regra é aplicável—não há escolha. Outra forma menos verbosa de fazer isso é dizer que elas devem ser avaliadas na ordem em que foram declaradas. Nesse caso, a primeira que for satisfeita deve ser aplicada.

Await and emit

(1.1) $\langle \texttt{await}(e), \alpha, m\rangle \xrightarrow{n} \langle \texttt{@awaiting}(e, n'), \alpha, m\rangle$ with $n' = n + 1$

(1.2) $\langle \texttt{@awaiting}(e, n'), e\alpha, m\rangle \xrightarrow{n} \langle \texttt{skip}, e\alpha, m\rangle$ if $n' \leq n$

(1.3) $\langle \texttt{emit}(e), \alpha, m\rangle \xrightarrow{n} \langle \texttt{@emitting}(n'), e\alpha, m\rangle$ with $n' = |\alpha|$

(1.4) $\langle \texttt{@emitting}(n'), \alpha, m\rangle \xrightarrow{n} \langle \texttt{skip}, \alpha, m\rangle$ if $n' = |\alpha|$

Conditionals

(1.5) $\langle \texttt{if}\,b\,\texttt{then}\,p_1\,\texttt{else}\,p_2, \alpha, m\rangle \xrightarrow{n} \langle p_1, \alpha, m\rangle$ if $eval(b, m) = 1$

(1.6) $\langle \texttt{if}\,b\,\texttt{then}\,p_1\,\texttt{else}\,p_2, \alpha, m\rangle \xrightarrow{n} \langle p_2, \alpha, m\rangle$ if $eval(b, m) = 0$

Sequences

(1.7) $\langle \texttt{skip}; p, \alpha, m\rangle \xrightarrow{n} \langle p, \alpha, m\rangle$

(1.8) $\langle v := a; p, \alpha, m\rangle \xrightarrow{n} \langle p, \alpha, m'\rangle$ with $m' = m[v/eval(a)]$

(1.9) $\langle \texttt{break}; p, \alpha, m\rangle \xrightarrow{n} \langle \texttt{break}, \alpha, m\rangle$

(1.10) $$\frac{\langle p_1, \alpha, m\rangle \xrightarrow{n} \langle p_1', \alpha', m'\rangle}{\langle p_1; p_2, \alpha, m\rangle \xrightarrow{n} \langle p_1'; p_2, \alpha', m'\rangle}$$ if $p_1 \neq \texttt{skip}, v := a, \texttt{break}$

Loops

$$\langle \texttt{loop}\ p, \alpha, m\rangle \xrightarrow{n} \langle p\ \texttt{@loop}\ p, \alpha, m\rangle \tag{1.11}$$

$$\langle \texttt{skip}\ \texttt{@loop}\ p, \alpha, m\rangle \xrightarrow{n} \langle \texttt{loop}\ p, \alpha, m\rangle \tag{1.12}$$

$$\langle v := a\ \texttt{@loop}\ p, \alpha, m\rangle \xrightarrow{n} \langle \texttt{loop}\ p, \alpha, m'\rangle \qquad \text{with } m' = m[v/\mathit{eval}(a)] \tag{1.13}$$

$$\langle \texttt{break}\ \texttt{@loop}\ p, \alpha, m\rangle \xrightarrow{n} \langle \texttt{skip}, \alpha, m\rangle \tag{1.14}$$

$$\frac{\langle p_1, \alpha, m\rangle \xrightarrow{n} \langle p_1', \alpha', m'\rangle}{\langle p_1\ \texttt{@loop}\ p_2, \alpha, m\rangle \xrightarrow{n} \langle p_1'\ \texttt{@loop}\ p_2, \alpha', m'\rangle} \qquad \begin{array}{l}\text{if } p_1 \neq \texttt{skip},\\ \quad v := a, \texttt{break}\end{array} \tag{1.15}$$

Par/and

Adicionamos a condição de o lado esquerdo estar bloqueado na regra 1.19.

$$\langle \texttt{skip}\ \mathsf{and}\ p, \alpha, m\rangle \xrightarrow{n} \langle p, \alpha, m\rangle \qquad \text{if } p \neq \texttt{break} \tag{1.16}$$

$$\langle p\ \mathsf{and}\ \texttt{skip}, \alpha, m\rangle \xrightarrow{n} \langle p, \alpha, m\rangle \qquad \text{if } p \neq \texttt{break} \tag{1.17}$$

$$\langle v := a\ \mathsf{and}\ p, \alpha, m\rangle \xrightarrow{n} \langle p, \alpha, m'\rangle \qquad \text{with } m' = m[v/\mathit{eval}(a)] \tag{1.18}$$

$$\langle p\ \mathsf{and}\ v := a, \alpha, m\rangle \xrightarrow{n} \langle p, \alpha, m'\rangle \qquad \begin{array}{l}\text{if } \mathit{blocked}(p, \alpha, n) = 1,\\ \quad \text{with } m' = m[v/\mathit{eval}(a)]\end{array} \tag{1.19}$$

$$\langle \texttt{break}\ \mathsf{and}\ p, \alpha, m\rangle \xrightarrow{n} \langle p'; \texttt{break}, \alpha, m\rangle \qquad \text{with } p' = \mathit{clear}(p) \tag{1.20}$$

$$\langle p\ \mathsf{and}\ \texttt{break}, \alpha, m\rangle \xrightarrow{n} \langle p'; \texttt{break}, \alpha, m\rangle \qquad \begin{array}{l}\text{if } \mathit{blocked}(p, \alpha, n) = 1,\\ \quad \text{with } p' = \mathit{clear}(p)\end{array} \tag{1.21}$$

$$\frac{\langle p_1, \alpha, m\rangle \xrightarrow{n} \langle p_1', \alpha', m'\rangle}{\langle p_1\ \mathsf{and}\ p_2, \alpha, m\rangle \xrightarrow{n} \langle p_1'\ \mathsf{and}\ p_2, \alpha', m'\rangle} \qquad \begin{array}{l}\text{if } \mathit{blocked}(p_1, \alpha, n) = 0\\ \quad \text{and } p_1 \neq \texttt{skip},\\ \quad v := a, \texttt{break}\end{array} \tag{1.22}$$

$$\frac{\langle p_2, \alpha, m\rangle \xrightarrow{n} \langle p_2', \alpha', m'\rangle}{\langle p_1\ \mathsf{and}\ p_2, \alpha, m\rangle \xrightarrow{n} \langle p_1\ \mathsf{and}\ p_2', \alpha', m'\rangle} \qquad \begin{array}{l}\text{if } \mathit{blocked}(p_1, \alpha, n) = 1\\ \quad \text{and } p_2 \neq \texttt{skip},\\ \quad v := a, \texttt{break}\end{array} \tag{1.23}$$

Par/or

$$\langle \texttt{skip}\ \texttt{or}\ p, \alpha, m\rangle \xrightarrow{n} \langle p', \alpha, m\rangle \quad \text{with } p' = \mathit{clear}(p) \tag{1.24}$$

$$\langle p\ \texttt{or}\ \texttt{skip}, \alpha, m\rangle \xrightarrow{n} \langle p', \alpha, m\rangle \quad \text{if } \mathit{blocked}(p, \alpha, n) = 1, \text{ with } p' = \mathit{clear}(p) \tag{1.25}$$

$$\langle v \coloneqq a\ \texttt{or}\ p, \alpha, m\rangle \xrightarrow{n} \langle p', \alpha, m'\rangle \quad \text{with } m' = m[v/\mathit{eval}(a)] \text{ and } p' = \mathit{clear}(p) \tag{1.26}$$

$$\langle p\ \texttt{or}\ v \coloneqq a, \alpha, m\rangle \xrightarrow{n} \langle p', \alpha, m'\rangle \quad \text{if } \mathit{blocked}(p, \alpha, n) = 1, \text{ with } m' = m[v/\mathit{eval}(a)] \text{ and } p' = \mathit{clear}(p) \tag{1.27}$$

$$\langle \texttt{break}\ \texttt{or}\ p, \alpha, m\rangle \xrightarrow{n} \langle p'; \texttt{break}, \alpha, m\rangle \quad \text{with } p' = \mathit{clear}(p) \tag{1.28}$$

$$\langle p\ \texttt{or}\ \texttt{break}, \alpha, m\rangle \xrightarrow{n} \langle p'; \texttt{break}, \alpha, m\rangle \quad \text{if } \mathit{blocked}(p, \alpha, n) = 1, \text{ with } p' = \mathit{clear}(p) \tag{1.29}$$

$$\frac{\langle p_1, \alpha, m\rangle \xrightarrow{n} \langle p_1', \alpha', m'\rangle}{\langle p_1\ \texttt{or}\ p_2, \alpha, m\rangle \xrightarrow{n} \langle p_1'\ \texttt{or}\ p_2, \alpha', m'\rangle} \quad \text{if } \mathit{blocked}(p_1, \alpha, n) = 0 \text{ and } p_1 \neq \texttt{skip}, v \coloneqq a, \texttt{break} \tag{1.30}$$

$$\frac{\langle p_2, \alpha, m\rangle \xrightarrow{n} \langle p_2', \alpha', m'\rangle}{\langle p_1\ \texttt{or}\ p_2, \alpha, m\rangle \xrightarrow{n} \langle p_1\ \texttt{or}\ p_2', \alpha', m'\rangle} \quad \text{if } \mathit{blocked}(p_1, \alpha, n) = 1 \text{ and } p_2 \neq \texttt{skip}, v \coloneqq a, \texttt{break} \tag{1.31}$$

The next theorem establishes that the reaction inner-step relation is deterministic, i.e., that it is in fact a *partial* function.

Theorem 1.4 (Determinism of the inner-step relation)**.** *For all p, p_1, $p_2 \in P$, α, α_1, $\alpha_2 \in E^*$, m, m_1, $m_2 \in \mathcal{M}$, and $n \in N$,*

$$\begin{aligned} &\text{if} \quad \langle p, \alpha, m\rangle \xrightarrow{n} \langle p_1, \alpha_1, m_1\rangle \quad \text{and} \quad \langle p, \alpha, m\rangle \xrightarrow{n} \langle p_2, \alpha_2, m_2\rangle, \\ &\qquad\qquad \text{then} \quad \langle p_1, \alpha_1, m_1\rangle = \langle p_2, \alpha_2, m_2\rangle. \end{aligned}$$

Proof. By induction on the structure of derivations. Suppose

$$d_1 \Vdash \langle p, \alpha, m\rangle \xrightarrow{n} \langle p_1, \alpha_1, m_1\rangle \quad \text{and} \quad d_2 \Vdash \langle p, \alpha, m\rangle \xrightarrow{n} \langle p_2, \alpha_2, m_2\rangle,$$

for some derivations d_1 and d_2. Then there are ten possibilities depending on the structure of p. (Note that p cannot be equal to `skip`, $v \coloneqq a$, or `break`, as there are no rules to evaluate such programs.)

[Case 1] $p = \texttt{await}(e)$, for some $e \in E$. Then d_1 and d_2 are instances of axiom (1.1), and as such, $p_1 = p_2 = \texttt{@awaiting}(e, n')$ with $n' = n + 1$, and $\alpha_1 = \alpha_2 = \alpha$ and $m_1 = m_2 = m$.

[Case 2] $p = \texttt{@awaiting}(e, n')$, for some $e \in E$ and $n' \in N$. Then d_1 and d_2 are instances of axiom (1.2), with $n' \leq n$. Thus $p_1 = p_2 = \texttt{skip}$, $\alpha_1 = \alpha_2 = \alpha$, and $m_1 = m_2 = m$.

[Case 3] $p = \texttt{emit}(e)$, for some $e \in E$. Then d_1 and d_2 are instances of axiom (1.3), and as such, $p_1 = p2 = \texttt{@emitting}(n')$ with $n' = |\alpha|$, and $\alpha_1 = \alpha_2 = e\alpha$ and $m_1 = m_2 = m$.

[Case 4] $p = \texttt{@emitting}(e, n')$, for some $e \in E$ and $n' \in N$. Then d_1 and d_2 are instances of axiom (1.4) with $n' = |\alpha|$. Thus $p_1 = p_2 = \texttt{skip}$, $\alpha_1 = \alpha_2 = \alpha$ and $m_1 = m_2 = m$.

[Case 5] $p = \texttt{if}\, b\, \texttt{then}\, p'\, \texttt{else}\, p''$, for some $b \in B$ and $p', p'' \in P$.

[Case 5.1] $\mathit{eval}(b, m) = 1$. Then d_1 and d_2 are instances of axiom (1.5), and as such, $p_1 = p_2 = p'$, $\alpha_1 = \alpha_2 = \alpha$, and $m_1 = m_2 = m$.

[Case 5.2] $\mathit{eval}(b, m) = 0$. Then d_1 and d_2 are instances of axiom (1.6), and as such, $p_1 = p_2 = p''$, $\alpha_1 = \alpha_2 = \alpha$, and $m_1 = m_2 = m$.

[Case 6] $p = p'; p''$, for some $p', p'' \in P$.

[Case 6.1] $p' = \texttt{skip}$. Then d_1 and d_2 are instances of axiom (1.7), and as such, $p_1 = p_2 = p''$, $\alpha_1 = \alpha_2 = \alpha$ and $m_1 = m_2 = m$.

[Case 6.2] $p' = v \coloneqq a$, for some $v \in V$ and $a \in A$. Then d_1 and d_2 are instances of axiom (1.8), and as such, $p_1 = p_2 = p''$, $\alpha_1 = \alpha_2 = \alpha$ and, as eval is a total function, $m_1 = m_2 = m[v/\mathit{eval}(a)]$.

[Case 6.3] $p' = \texttt{break}$. Then d_1 and d_2 are instances of axiom (1.9), and as such, $p_1 = p_2 = p'$, $\alpha_1 = \alpha_2 = \alpha$ and $m_1 = m_2 = m$.

[Case 6.4] $p' \neq \texttt{skip}, v \coloneqq a, \texttt{break}$. Then d_1 and d_2 are instances of rule (1.10). Thus there are derivations d_1' and d_2' such that

$$d_1' \Vdash \langle p', \alpha, m\rangle \xrightarrow{n} \langle p_1', \alpha_1, m_1\rangle \quad \text{and} \quad d_2' \Vdash \langle p', \alpha, m\rangle \xrightarrow{n} \langle p_2', \alpha_2, m_2\rangle,$$

for some $p_1', p_2' \in P$. Since $d_1' \prec d_1$ and $d_2' \prec d_2$, by induction hypothesis, $\alpha_1 = \alpha_2$, $m_1 = m_2$, and $p_1' = p_2'$, which implies

$$p_1 = p_1'; p'' = p_2'; p'' = p_2.$$

[Case 7] $p = \texttt{loop}\, p'$, for some $p' \in P$. Then d_1 and d_2 are instances of axiom (1.11), and as such, $p_1 = p_2 = p'\, \texttt{@loop}\, p'$, $\alpha_1 = \alpha_2 = \alpha$, and $m_1 = m_2 = m$.

[Case 8] $p = p'\, \texttt{@loop}\, p''$, for some $p', p'' \in P$.

[Case 8.1] $p = \texttt{skip}\, \texttt{@loop}\, p'$, for some $p' \in P$. Then d_1 and d_2 are instances of axiom (1.12), and as such, $p_1 = p_2 = \texttt{loop}\, p'$, $\alpha_1 = \alpha_2 = \alpha$, and $m_1 = m_2 = m$.

[Case 8.2] $p = v \coloneqq a\, \texttt{@loop}\, p'$, for some $a \in A$, $v \in V$, and $p' \in P$. Then d_1 and d_2 are instances of axiom (1.13), and as such, $p_1 = p_2 = \texttt{loop}\, p'$, $\alpha_1 = \alpha_2 = \alpha$, and as eval is a total function, $m_1 = m_2 = m[v/\mathit{eval}(a)]$.

[Case 8.3] $p = \texttt{break}\, \texttt{@loop}\, p'$, for some $p' \in P$. Then d_1 and d_2 are instances of axiom (1.14), and as such, $p_1 = p_2 = \texttt{skip}$, $\alpha_1 = \alpha_2 = \alpha$, and $m_1 = m_2 = m$.

[Case 8.4] $p = p' \,\texttt{@loop}\, p''$, for some $p', p'' \in P$ such that $p' \neq \texttt{skip}, v := a, \texttt{break}$. Then d_1 and d_2 are instances of rule (1.15). Thus there are derivations d'_1 and d'_2 such that

$$d'_1 \Vdash \langle p', \alpha, m\rangle \xrightarrow{n} \langle p'_1, \alpha_1, m_1\rangle \quad \text{and} \quad d'_2 \Vdash \langle p', \alpha, m\rangle \xrightarrow{n} \langle p'_2, \alpha_2, m_2\rangle,$$

for some p'_1, $p'_2 \in P$. Since $d'_1 \prec d_1$ and $d'_2 \prec d_2$, by induction hypothesis, $\alpha_1 = \alpha$, $m_1 = m_2$, and $p'_1 = p'_2$, which implies

$$p_1 = p'_1 \,\texttt{@loop}\, p'' = p'_2 \,\texttt{@loop}\, p'' = p_2.$$

[Case 9] $p = p' \,\texttt{and}\, p''$, for some $p', p'' \in P$.

[Case 9.1] $p = \texttt{skip} \,\texttt{and}\, p'$, for some $p' \in P$ and $p' \neq \texttt{break}$. Then d_1 and d_2 are instances of axiom (1.16), and as such, $p_1 = p_2 = p'$, $\alpha_1 = \alpha_2 = \alpha$, and $m_1 = m_2 = m$.

[Case 9.2] $p = p' \,\texttt{and}\, \texttt{skip}$, for some $p' \in P$ and $p' \neq \texttt{break}$. Similar to Case 9.1.

[Case 9.3] $p = v := a \,\texttt{and}\, p'$, for some $v \in V$, $a \in A$ and $p' \in P$. Then d_1 and d_2 are instances of axiom (1.18), and as such, $p_1 = p_2 = p'$, $\alpha_1 = \alpha_2 = \alpha$, and as *eval* is a total function, $m_1 = m_2 = m[v/eval(a)]$.

[Case 9.4] $p = p' \,\texttt{and}\, v := a$, for some $v \in V$, $a \in A$ and $p' \in P$. Then either $blocked(p') = 0$ or $blocked(p') = 1$. If $blocked(p') = 0$ then this case becomes Case 9.7. Otherwise, if $blocked(p') = 1$, then d_1 and d_2 are instances of axiom (1.19), and as such, $p_1 = p_2 = p'$, $\alpha_1 = \alpha_2 = \alpha$, and as *eval* is a total function, $m_1 = m_2 = m[v/eval(a)]$.

[Case 9.5] $p = \texttt{break} \,\texttt{and}\, p'$, for some $p' \in P$. Then d_1 and d_2 are instances of axiom (1.20), and as such, $\alpha_1 = \alpha_2 = \alpha$, $m_1 = m_2 = m$, and as *clear* is a total function, $p_1 = p_2 = clear(p'); \texttt{break}$.

[Case 9.6] $p = p' \,\texttt{and}\, \texttt{break}$, for some $p' \in P$. Then either $blocked(p') = 0$ or $blocked(p') = 1$. If $blocked(p') = 0$ then this case becomes Case 9.7. Otherwise, if $blocked(p') = 1$, then d_1 and d_2 are instances of axiom (1.21), and as such, $\alpha_1 = \alpha_2 = \alpha$, $m_1 = m_2 = m$, and as *clear* is a total function, $p_1 = p_2 = clear(p'); \texttt{break}$.

[Case 9.7] $p = p' \,\texttt{and}\, p''$, for some p' and $p'' \in P$. Then there are two possibilities. If $blocked(p') = 0$ then d_1 and d_2 are instances of (1.22). Thus there are derivations d'_1 and d'_2 such that

$$d'_1 \Vdash \langle p', \alpha, m\rangle \xrightarrow{n} \langle p'_1, \alpha_1, m_1\rangle \quad \text{and} \quad d'_2 \Vdash \langle p', \alpha, m\rangle \xrightarrow{n} \langle p'_2, \alpha_2, m_2\rangle,$$

for some p'_1, $p'_2 \in P$. Since $d'_1 \prec d_1$ and $d'_2 \prec d_2$, by induction hypothesis, $\alpha_1 = \alpha_2$, $m_1 = m_2$, and $p'_1 = p'_2$, which implies

$$p_1 = (p'_1 \,\texttt{and}\, p'') = (p'_2 \,\texttt{and}\, p'') = p_2.$$

If, however, $blocked(p') = 1$ and $p'' \neq \texttt{skip}, v := a, \texttt{break}$, then d_1 and d_2 are instances of (1.23). Thus there are derivations d_1'' and d_2'' such that

$$d_1'' \Vdash \langle p'', \alpha, m\rangle \xrightarrow{n} \langle p_1'', \alpha_1, m_1\rangle \quad \text{and} \quad d_2'' \Vdash \langle p'', \alpha, m\rangle \xrightarrow{n} \langle p_2'', \alpha_2, m_2\rangle,$$

for some p_1'', $p_2'' \in P$. Since $d_1'' \prec d_1$ and $d_2'' \prec d_2$, by induction hypothesis, $\alpha_1 = \alpha_2$, $m_1 = m_2$, and $p_1'' = p_2''$, which implies

$$p_1 = (p' \;\texttt{and}\; p_1'') = (p' \;\texttt{and}\; p_2'') = p_2.$$

[Case 10] $p = p' \;\texttt{or}\; p''$, for some $p', p'' \in P$.

[Case 10.1] $p = \texttt{skip} \;\texttt{or}\; p'$, for some $p' \in P$. Then d_1 and d_2 are instances of axiom (1.24), and as such, $\alpha_1 = \alpha_2 = \alpha$, $m_1 = m_2 = m$, and as *clear* is a total, $p_1 = p_2 = clear(p')$.

[Case 10.2] $p = p' \;\texttt{or}\; \texttt{skip}$, for some $p' \in P$. Then either $blocked(p') = 0$ or $blocked(p') = 1$. If $blocked(p') = 0$ then this case becomes Case 10.7. Otherwise, if $blocked(p') = 1$, then d_1 and d_2 are instances of axiom (1.25), and as such, $\alpha_1 = \alpha_2 = \alpha$, $m_1 = m_2 = m$, and as *clear* is a total function, $p_1 = p_2 = clear(p')$.

[Case 10.3] $p = v := a \;\texttt{or}\; p'$, for some $v \in V$, $a \in A$ and $p' \in P$. Then d_1 and d_2 are instances of axiom (1.26), and as such, $\alpha_1 = \alpha_2 = \alpha$, and as *eval* and *clear* are total functions, $m_1 = m_2 = m[v/eval(a)]$ and $p_1 = p_2 = clear(p')$.

[Case 10.4] $p = p' \;\texttt{or}\; v := a$, for some $v \in V$, $a \in A$ and $p' \in P$. Then either $blocked(p') = 0$ or $blocked(p') = 1$. If $blocked(p') = 0$ then this case becomes Case 10.7. Otherwise, if $blocked(p') = 1$, then d_1 and d_2 are instances of axiom (1.27), and as such, $\alpha_1 = \alpha_2 = \alpha$, and as *eval* and *clear* are total functions, $m_1 = m_2 = m[v/eval(a)]$ and $p_1 = p_2 = clear(p')$.

[Case 10.5] $p = \texttt{break} \;\texttt{or}\; p'$, for some $p' \in P$. Then d_1 and d_2 are instances of axiom (1.28), and as such, $\alpha_1 = \alpha_2 = \alpha$, $m_1 = m_2 = m$, and as *clear* is a total function, $p_1 = p_2 = clear(p'); \texttt{break}$.

[Case 10.6] $p = p' \;\texttt{or}\; \texttt{break}$, for some $p' \in P$. Then either $blocked(p') = 0$ or $blocked(p') = 1$. If $blocked(p') = 0$ then this case becomes Case 10.7. Otherwise, if $blocked(p') = 1$, then d_1 and d_2 are instances of axiom (1.29), and as such, $\alpha_1 = \alpha_2 = \alpha$, $m_1 = m_2 = m$, and as *clear* is a total function, $p_1 = p_2 = clear(p'); \texttt{break}$.

[Case 10.7] $p = p' \;\texttt{or}\; p''$, for some p' and $p'' \in P$. Then there are two possibilities. If $blocked(p') = 0$ then d_1 and d_2 are instances of (1.30). Thus there are derivations d_1' and d_2' such that

$$d_1' \Vdash \langle p', \alpha, m\rangle \xrightarrow{n} \langle p_1', \alpha_1, m_1\rangle \quad \text{and} \quad d_2' \Vdash \langle p', \alpha, m\rangle \xrightarrow{n} \langle p_2', \alpha_2, m_2\rangle,$$

for some p_1', $p_2' \in P$. Since $d_1' \prec d_1$ and $d_2' \prec d_2$, by induction hypothesis, $\alpha_1 = \alpha_2$, $m_1 = m_2$, and $p_1' = p_2'$, which implies

$$p_1 = (p_1' \;\texttt{or}\; p'') = (p_2' \;\texttt{or}\; p'') = p_2.$$

If, however, $blocked(p') = 1$ and $p'' \neq \texttt{skip}, v := a, \texttt{break}$, then d_1 and d_2 are instances of (1.31). Thus there are derivations d_1'' and d_2'' such that

$$d_1'' \Vdash \langle p'', \alpha, m \rangle \xrightarrow{n} \langle p_1'', \alpha_1, m_1 \rangle \quad \text{and} \quad d_2'' \Vdash \langle p'', \alpha, m \rangle \xrightarrow{n} \langle p_2'', \alpha_2, m_2 \rangle,$$

for some $p_1'', p_2'' \in P$. Since $d_1'' \prec d_1$ and $d_2'' \prec d_2$, by induction hypothesis, $\alpha_1 = \alpha_2$, $m_1 = m_2$, and $p_1'' = p_2''$, which implies

$$p_1 = (p' \texttt{ or } p_1'') = (p' \texttt{ or } p_2'') = p_2. \qquad \square$$

The next lemma establishes that given a program either it is possible to advance it by an inner-step or all its trails are blocked, but not both.

Lemma 1.5. *For all $p \in P$, $\alpha \in E^*$, $m \in \mathcal{M}$, and $n \in N$, if $p \neq \texttt{skip}, v := a, \texttt{break}$ then either*

$$\exists \delta \in \Delta(\langle p, \alpha, m \rangle \xrightarrow{n} \delta) \quad \textit{or} \quad blocked(p, \alpha, n) = 1,$$

but not both.

Proof. By induction on the structure of programs.

[Case 1] $p = \texttt{await}(e)$, for some $e \in E$. Then by axiom (1.1),

$$\langle \texttt{await}(e), \alpha, m \rangle \xrightarrow{n} \langle \texttt{@awaiting}(e, n'), \alpha, m \rangle = \delta,$$

where $n' = n + 1$. And by Definition 1.1, $blocked(\texttt{await}(e), \alpha, n) = 0$.

[Case 2] $p = \texttt{@awaiting}(e, n')$, for some $e \in E$ and $n' \in N$.

[Case 2.1] $n' < n$. If e is the top-of-stack event in α, in symbols $e = \alpha_{[1]}$, then by axiom (1.2),

$$\langle \texttt{@awaiting}(e, n'), \alpha, m \rangle \xrightarrow{n} \langle \texttt{skip}, \alpha, m \rangle = \delta.$$

And by Definition 1.1, $blocked(\texttt{@awaiting}(e, n'), \alpha, n) = 0$.
If, however, $e \neq \alpha_{[1]}$, then there is no such δ, as no rule is applicable. And by Definition 1.1, $blocked(\texttt{@awaiting}(e, n'), \alpha, n) = 1$.

[Case 2.2] $n' = n$. [FIXME: Pelo axioma (1.2),

$$\langle \texttt{@awaiting}(e, n'), \alpha, m \rangle \xrightarrow{n} \langle \texttt{skip}, \alpha, m \rangle.$$

E $blocked(\texttt{@awaiting}(e, n'), \alpha, n) = 1$. Ou seja, ambos os lados do "ou" deram verdadeiro, o que invalida o lema.]

[Case 2.3] $n' > n$. If $e = \alpha_{[1]}$ then [FIXME: Não existe tal δ e

$$blocked(\texttt{@awaiting}(e, n'), \alpha, n) = 0.$$

O que, novamente, invalida o lema.]
If, however, $e \neq \alpha_{[1]}$, then there is no such δ (no rule is applicable) and, by Definition 1.1, $blocked(\texttt{@awaiting}(e, n'), \alpha, n) = 1$.

[Case 3] $p = \texttt{emit}(e)$, for some $e \in E$. Then by axiom (1.3),

$$\langle \texttt{emit}(e), \alpha, m\rangle \xrightarrow{n} \langle \texttt{@emitting}(n'), e\alpha, m\rangle = \delta,$$

where $n' = |\alpha|$. And by Definition 1.1, $\mathit{blocked}(\texttt{emit}(e'), e\alpha, n) = 0$.

[Case 4] $p = \texttt{@emitting}(e, n')$, for some $e \in E$ and $n' \in N$.

[Case 4.1] $n' = |\alpha|$. By axiom (1.4),

$$\langle \texttt{@emitting}(e, n'), \alpha, m\rangle \xrightarrow{n} \langle \texttt{skip}, \alpha, m\rangle = \delta.$$

And by Definition 1.1, $\mathit{blocked}(\texttt{@emitting}(e, n'), \alpha, n) = 0$.

[Case 4.2] $n' \neq |\alpha|$. Then there is no such δ (no rule is applicable) and, by Definition 1.1, $\mathit{blocked}(\texttt{@emitting}(e, n'), \alpha, n) = 1$.

[Case 5] $p = \texttt{if}\, b\, \texttt{then}\, p'\, \texttt{else}\, p''$, for some $b \in B$ and p', $p'' \in P$. By axioms (1.5) and (1.6), if $\mathit{eval}(b, m) = 1$, $\delta = \langle p', \alpha, m\rangle$, otherwise $\delta = \langle p'', \alpha, m\rangle$. And by Definition 1.1, $\mathit{blocked}(\texttt{if}\, b\, \texttt{then}\, p'\, \texttt{else}\, p'', e\alpha, n) = 0$

[Case 6] $p = p'; p''$, for some p', $p'' \in P$. By Definition 1.1 $\mathit{blocked}(p'; p'', \alpha, n) = \mathit{blocked}(p', \alpha, n)$

[Case 6.1] $p' = \texttt{skip}$. By axiom (1.7), $\delta = \langle p'', \alpha, m\rangle$, and by Definition 1.1, $\mathit{blocked}(\texttt{skip}, \alpha, n) = 0$.

[Case 6.2] $p' = v := a$, for some $v \in V$ and $a \in A$. By axiom (1.8), $\delta = \langle p'', \alpha, m[v/\mathit{eval}(a)]\rangle$, and by Definition 1.1, $\mathit{blocked}(v := a, \alpha, n) = 0$.

[Case 6.3] $p' = \texttt{break}$. By axiom (1.9), $\delta = \langle \mathit{Break}, \alpha, m\rangle$, and by Definition 1.1, $\mathit{blocked}(\texttt{break}, \alpha, n) = 0$.

[Case 6.4] $p' \neq \texttt{skip}, v := a, \texttt{break}$. By induction hypothesis, exactly one of the following hold:

$$\exists d' \in \Delta(\langle p', \alpha, m\rangle \xrightarrow{n} d') \quad \text{or} \quad \mathit{blocked}(p', \alpha, n) = 1.$$

Suppose $\langle p', \alpha, m\rangle \xrightarrow{n} \langle p'_1, \alpha', m'\rangle$, for some $p'_1 \in P$, $\alpha' \in E^*$, and $m' \in \mathcal{M}$. Then by rule (1.10),

$$\langle p'; p'', \alpha, m\rangle \xrightarrow{n} \langle p'_1; p'', \alpha', m'\rangle.$$

And by Definition 1.1,

$$\mathit{blocked}(p'; p'', \alpha, n) = \mathit{blocked}(p', \alpha, n) = 0.$$

If, however, $\mathit{blocked}(p', \alpha, n) = 1$, then there is no such d (no rule is applicable) and by Definition 1.1,

$$\mathit{blocked}(p'; p'', \alpha, n) = \mathit{blocked}(p', \alpha, n) = 1.$$

[Case 7] $p = \texttt{loop}\, p'$, for some $p' \in P$. By axiom (1.11), $\delta = \langle p'\, \texttt{@loop}\, p', \alpha, m\rangle$. And by Definition 1.1, $\mathit{blocked}(\texttt{loop}\, p', \alpha, m) = 0$.

[Case 8] $p = p' \,\texttt{@loop}\, p''$, for some $p', p'' \in P$. By Definition 1.1 $blocked(p' \,\texttt{@loop}\, p'', \alpha, n) = blocked(p', \alpha, n)$

[Case 8.1] $p' = \texttt{skip}$. By axiom (1.12), $\delta = \langle \texttt{loop}\, p'', \alpha, m \rangle$, and by Definition 1.1, $blocked(\texttt{skip}, \alpha, n) = 0$.

[Case 8.2] $p' = v := a$, for some $v \in V$ and $a \in A$. By axiom (1.13), $\delta = \langle \texttt{loop}\, p'', \alpha, m[v/eval(a)] \rangle$, and by Definition 1.1, $blocked(v := a, \alpha, n) = 0$.

[Case 8.3] $p' = \texttt{break}$. By axiom (1.14), $\delta = \langle \texttt{skip}, \alpha, m \rangle$, and by Definition 1.1, $blocked(\texttt{break}, \alpha, n) = 0$.

[Case 8.4] $p' \neq \texttt{skip}, v := a, \texttt{break}$. By induction hypothesis, exactly one of the following hold:

$$\exists d' \in \Delta(\langle p', \alpha, m \rangle \xrightarrow{n} d') \quad \text{or} \quad blocked(p', \alpha, n) = 1.$$

Suppose $\langle p', \alpha, m \rangle \xrightarrow{n} \langle p'_1, \alpha', m' \rangle$, for some $p'_1 \in P$, $\alpha' \in E^*$, and $m' \in \mathcal{M}$. Then by rule (1.15),

$$\langle p' \,\texttt{@loop}\, p'', \alpha, m \rangle \xrightarrow{n} \langle p'_1 \,\texttt{@loop}\, p'', \alpha', m' \rangle.$$

And by Definition 1.1,

$$blocked(p' \,\texttt{@loop}\, p'', \alpha, n) = blocked(p', \alpha, n) = 0.$$

If, however, $blocked(p', \alpha, n) = 1$, then there is no such d (no rule is applicable) and by Definition 1.1,

$$blocked(p' \,\texttt{@loop}\, p'', \alpha, n) = blocked(p', \alpha, n) = 1.$$

[Case 9] $p = p' \,\texttt{and}\, p''$, for some $p' \in P$, By Definition 1.1

$$blocked(p' \,\texttt{and}\, p'', \alpha, n) = blocked(p', \alpha, n) \cdot blocked(p'', \alpha, n)$$

[Case 9.1] $p' = \texttt{skip}$ and $p'' \neq \texttt{break}$. By axiom (1.16), $\delta = \langle p'', \alpha, m \rangle$, and by Definition 1.1, $blocked(\texttt{skip}, \alpha, n) = 0$.

[Case 9.2] $p'' = \texttt{skip}$ and $p' \neq \texttt{break}$. By axiom (1.17), $\delta = \langle p', \alpha, m \rangle$, and by Definition 1.1, $blocked(\texttt{skip}, \alpha, n) = 0$.

[Case 9.3] $p' = v := a$, for some $v \in V$ and $a \in A$. By axiom (1.18), $\delta = \langle p'', \alpha, m[v/eval(a)] \rangle$, and by Definition 1.1, $blocked(v := a, \alpha, n) = 0$.

[Case 9.4] $p'' = v := a$, for some $v \in V$ and $a \in A$. By induction hypothesis, exactly one of the following hold:

$$\exists d' \in \Delta(\langle p', \alpha, m \rangle \xrightarrow{n} d') \quad \text{or} \quad blocked(p', \alpha, n) = 1.$$

If $blocked(p', \alpha, m) = 0$, then the derivation of δ is similar to Case 9.7.
If, however, $blocked(p', \alpha, n) = 1$, then by axiom (1.19), $\delta = \langle p', \alpha, m[v/eval(a)] \rangle$.
Either way, $blocked(v := a, \alpha, n) = 0$.

[Case 9.5] $p' = \texttt{break}$. By axiom (1.20), $\delta = \langle clear(p''), \alpha, m \rangle$, and by Definition 1.1, $blocked(\texttt{break}, \alpha, n) = 0$.

[Case 9.6] $p'' = \texttt{break}$. By induction hypothesis, exactly one of the following hold:

$$\exists d' \in \Delta(\langle p', \alpha, m \rangle \xrightarrow{n} d') \quad \text{or} \quad blocked(p', \alpha, n) = 1.$$

If $blocked(p', \alpha, n) = 0$, then the derivation of δ is similar to Case 9.7.
If, however, $blocked(p', \alpha, n) = 1$, then by axiom (1.21), $\delta = \langle clear(p'), \alpha, m \rangle$.
Either way, $blocked(\texttt{break}, \alpha, n) = 0$.

[Case 9.7] $p' \neq \texttt{skip}, v := a, \texttt{break}$. By induction hypothesis, exactly one of the following hold:

$$\exists d' \in \Delta(\langle p', \alpha, m \rangle \xrightarrow{n} d') \quad \text{or} \quad blocked(p', \alpha, n) = 1.$$

Suppose $\langle p', \alpha, m \rangle \xrightarrow{n} \langle p'_1, \alpha', m' \rangle$, for some $p'_1 \in P$, $\alpha' \in E^*$, and $m' \in \mathcal{M}$. Then by rule (1.22),

$$\langle p' \ \texttt{and} \ p'', \alpha, m \rangle \xrightarrow{n} \langle p'_1 \ \texttt{and} \ p'', \alpha', m' \rangle.$$

And by Definition 1.1,

$$blocked(p' \ \texttt{and} \ p'', \alpha, n) = blocked(p', \alpha, n) = 0.$$

Now suppose $\langle p'', \alpha, m \rangle \xrightarrow{n} \langle p''_1, \alpha', m' \rangle$, for some $p''_1 \in P$, $\alpha' \in E^*$, and $m' \in \mathcal{M}$. Then by rule (1.23),

$$\langle p' \ \texttt{and} \ p'', \alpha, m \rangle \xrightarrow{n} \langle p' \ \texttt{and} \ p''_1, \alpha', m' \rangle.$$

And by Definition 1.1,

$$blocked(p' \ \texttt{and} \ p'', \alpha, n) = blocked(p', \alpha, n) = 0.$$

[Case 10] $p = p' \ \texttt{or} \ p''$, for some $p' \in P$, By Definition 1.1

$$blocked(p' \ \texttt{or} \ p'', \alpha, n) = blocked(p', \alpha, n) \cdot blocked(p'', \alpha, n)$$

[Case 10.1] $p' = \texttt{skip}$. By axiom (1.24), $\delta = \langle clear(p''), \alpha, m \rangle$, and by Definition 1.1, $blocked(\texttt{skip}, \alpha, n) = 0$.

[Case 10.2] $p'' = \texttt{skip}$. By induction hypothesis, exactly one of the following hold:

$$\exists d' \in \Delta(\langle p', \alpha, m \rangle \xrightarrow{n} d') \quad \text{or} \quad blocked(p', \alpha, n) = 1.$$

If $blocked(p', \alpha, m) = 0$, then the derivation of δ is similar to Case 10.7.
If, however, $blocked(p', \alpha, n) = 1$, then by axiom (1.25), $\delta = \langle clear(p'), \alpha, m \rangle$.
Either way, $blocked(\texttt{skip}, \alpha, n) = 0$.

[Case 10.3] $p' = v := a$, for some $v \in V$ and $a \in A$. By axiom (1.26), $\delta = \langle clear(p''), \alpha, m[v/eval(a)] \rangle$, and by Definition 1.1, $blocked(v := a, \alpha, n) = 0$.

[Case 10.4] $p'' = v := a$, for some $v \in V$ and $a \in A$. By induction hypothesis, exactly one of the following hold:

$$\exists d' \in \Delta(\langle p', \alpha, m\rangle \xrightarrow{n} d') \quad \text{or} \quad blocked(p', \alpha, n) = 1.$$

If $blocked(p', \alpha, m) = 0$, then the derivation of δ is similar to Case 10.7.
If, however, $blocked(p', \alpha, n) = 1$, then by axiom (1.27), $\delta = \langle clear(p'), \alpha, m[v/eval(a)]\rangle$.
Either way, $blocked(v := a, \alpha, n) = 0$.

[Case 10.5] $p' = \mathtt{break}$. By axiom (1.28), $\delta = \langle clear(p''), \alpha, m\rangle$, and by Definition 1.1, $blocked(\mathtt{break}, \alpha, n) = 0$.

[Case 10.6] $p'' = \mathtt{break}$. By induction hypothesis, exactly one of the following hold:

$$\exists d' \in \Delta(\langle p', \alpha, m\rangle \xrightarrow{n} d') \quad \text{or} \quad blocked(p', \alpha, n) = 1.$$

If $blocked(p', \alpha, n) = 0$, then the derivation of δ is similar to Case 10.7.
If, however, $blocked(p', \alpha, n) = 1$, then by axiom (1.29), $\delta = \langle clear(p'), \alpha, m\rangle$.
Either way, $blocked(\mathtt{break}, \alpha, n) = 0$.

[Case 10.7] $p' \neq \mathtt{skip}, v := a, \mathtt{break}$. By induction hypothesis, exactly one of the following hold:

$$\exists d' \in \Delta(\langle p', \alpha, m\rangle \xrightarrow{n} d') \quad \text{or} \quad blocked(p', \alpha, n) = 1.$$

Suppose $\langle p', \alpha, m\rangle \xrightarrow{n} \langle p'_1, \alpha', m'\rangle$, for some $p'_1 \in P$, $\alpha' \in E^*$, and $m' \in \mathcal{M}$. Then by rule (1.30),

$$\langle p' \ \mathtt{or}\ p'', \alpha, m\rangle \xrightarrow{n} \langle p'_1 \ \mathtt{or}\ p'', \alpha', m'\rangle.$$

And by Definition 1.1,

$$blocked(p' \ \mathtt{or}\ p'', \alpha, n) = blocked(p', \alpha, n) = 0.$$

Now suppose $\langle p'', \alpha, m\rangle \xrightarrow{n} \langle p''_1, \alpha', m'\rangle$, for some $p''_1 \in P$, $\alpha' \in E^*$, and $m' \in \mathcal{M}$. Then by rule (1.31),

$$\langle p' \ \mathtt{or}\ p'', \alpha, m\rangle \xrightarrow{n} \langle p' \ \mathtt{and}\ p''_1, \alpha', m'\rangle.$$

And by Definition 1.1,

$$blocked(p' \ \mathtt{or}\ p'', \alpha, n) = blocked(p', \alpha, n) = 0.$$

□

1.3 The reaction outer-step relation

From the previous inner-step relation we define an outer-step relation ($\Rightarrow$) that when necessary pops the event stack and advances blocked programs. [TODO: Improve this description.]

Definition 1.6 (Reaction outer-step)**.**

$$\frac{\langle p, \alpha, m\rangle \xrightarrow{n} \langle p', \alpha', m'\rangle}{\langle p, \alpha, m\rangle \overset{n}{\Rightarrow} \langle p', \alpha', m'\rangle} \quad \text{if } \mathit{blocked}(p, \alpha, n) = 0 \tag{1.32}$$

$$\langle p, e\alpha, m\rangle \overset{n}{\Rightarrow} \langle p, \alpha, m\rangle \qquad \text{if } \mathit{blocked}(p, \alpha, n) = 1 \tag{1.33}$$

Theorem 1.7 (Determinism of the outer-step relation)**.** *For all p, p_1, $p_2 \in P$, α, α_1, $\alpha_2 \in E^*$, m, m_1, $m_2 \in \mathcal{M}$, and $n \in N$,*

$$\begin{gathered} \textit{if} \quad \langle p, \alpha, m\rangle \overset{n}{\Rightarrow} \langle p_1, \alpha_1, m_1\rangle \quad \textit{and} \quad \langle p, \alpha, m\rangle \overset{n}{\Rightarrow} \langle p_2, \alpha_2, m_2\rangle, \\ \textit{then} \quad \langle p_1, \alpha_1, m_1\rangle = \langle p_2, \alpha_2, m_2\rangle. \end{gathered}$$

Proof. [TODO: By induction on the structure of derivations.] □

Theorem 1.8 (Termination of the outer-step relation)**.** *For all $p \in P$, $\alpha \in E$, and $m \in \mathcal{M}$, if $p \neq$ `skip`$, v := a,$ `break` then*

$$\exists \delta \in \Delta(\langle p, \alpha, m\rangle \overset{n}{\Rightarrow} \delta).$$

Proof. [TODO: Directly from Lemma 1.5.] □

1.4 The reaction relation

From the reflexive-transitive closure of the outer-step relation ($\overset{*}{\Rightarrow}$) we define the reaction relation $\vDash \subseteq \Delta \times (P \times \mathcal{M} \times N)$, which computes a full program reaction. Given an initial configuration, the reaction relation evaluates it until the event stack becomes empty.

Definition 1.9 (Reaction)**.** Let p, $p' \in P$, $\alpha \in E^*$, m, $m' \in \mathcal{M}$. Then

$$\langle p, \alpha, m\rangle \overset{n}{\vDash} \langle p', m'\rangle \quad \text{iff} \quad \langle p, \alpha, m\rangle \overset{*}{\Rightarrow} \langle p', \varepsilon, m'\rangle.$$

The next two theorems establish, respectively, that reactions are deterministic and always terminate (for the nontrivial programs $p \neq$ `skip`$, v := a,$ `break`).

Theorem 1.10 (Determinism of the reaction relation)**.** *For all p, p_1, $p_2 \in P$, $\alpha \in E^*$, m, m_1, $m_2 \in \mathcal{M}$, and $n \in N$,*

$$\begin{gathered} \textit{if} \quad \langle p, \alpha, m\rangle \overset{n}{\vDash} \langle p_1, m_1\rangle \quad \textit{and} \quad \langle p, \alpha, m\rangle \overset{n}{\vDash} \langle p_2, m_2\rangle, \\ \textit{then} \quad \langle p_1, m_1\rangle = \langle p_2, m_2\rangle. \end{gathered}$$

Proof. [TODO: ?] □

Theorem 1.11 (Termination of the reaction relation)**.** *For all $p \in P$, $\alpha \in E$, and $m \in \mathcal{M}$, if $p \neq$ `skip`$, v := a,$ `break` then*

$$\langle p, \alpha, m\rangle \overset{n}{\vDash} \langle p', m'\rangle,$$

for some $p' \in P$ and $m' \in \mathcal{M}$.

Proof. [TODO: ?] □

2 Big-step version of the original formulation

[TODO: Minha ideia aqui é fazer uma versão big-step da formulação original. E no final comparar as duas versões, i.e., mostrar que são equivalentes.]

Definition 2.1. [TODO: Parcial e provavelmente incorreta.]
Empty program

$$\langle \varepsilon, \alpha, m, n \rangle \rightsquigarrow \langle \alpha, m, n \rangle \tag{2.1}$$

Assignment

$$\langle v \coloneqq a, \alpha, m, n \rangle \rightsquigarrow \langle \alpha, m[v/eval(a)], n \rangle \tag{2.2}$$

Conditionals

$$\frac{\langle p_1, \alpha, m, n \rangle \rightsquigarrow \langle \alpha', m', n' \rangle}{\langle \texttt{if}\, b\, \texttt{then}\, p_1\, \texttt{else}\, p_2, \alpha, m, n \rangle \rightsquigarrow \langle \alpha', m', n' \rangle} \quad \text{if } eval(b, m) = 1 \tag{2.3}$$

$$\frac{\langle p_2, \alpha, m, n \rangle \rightsquigarrow \langle \alpha', m', n' \rangle}{\langle \texttt{if}\, b\, \texttt{then}\, p_1\, \texttt{else}\, p_2, \alpha, m, n \rangle \rightsquigarrow \langle \alpha', m', n' \rangle} \quad \text{if } eval(b, m) = 0 \tag{2.4}$$

Await

$$\frac{\langle \texttt{@awaiting}(e, n+1), e\alpha, m, n \rangle \rightsquigarrow \langle \alpha', m', n' \rangle}{\langle \texttt{await}(e), \alpha, m, n \rangle \rightsquigarrow \langle \alpha', m', n' \rangle} \tag{2.5}$$

$$\langle \texttt{@awaiting}(e', n'), e\alpha, m, n \rangle \rightsquigarrow \langle \alpha, m, n \rangle \quad \text{if } e' = e \text{ and } n' < n \tag{2.6}$$

$$\frac{\langle \texttt{@awaiting}(e', n'), \alpha, m, n \rangle \rightsquigarrow \langle \alpha'', m'', n'' \rangle}{\langle \texttt{@awaiting}(e', n'), e\alpha, m, n \rangle \rightsquigarrow \langle \alpha'', m'', n'' \rangle} \quad \text{if } e' \neq e \text{ or } n' \geq n \tag{2.7}$$

Emit

$$\frac{\langle \texttt{@emitting}(|\alpha|), e\alpha, m, n \rangle \rightsquigarrow \langle \alpha', m', n' \rangle}{\langle \texttt{emit}(e), \alpha, m, n \rangle \rightsquigarrow \langle \alpha', m', n' \rangle} \tag{2.8}$$

$$\langle \texttt{@emitting}(n'), e\alpha, m, n \rangle \rightsquigarrow \langle \alpha, m, n \rangle \quad \text{if } |e\alpha| = n' \tag{2.9}$$

$$\frac{\langle \texttt{@emitting}(n'), \alpha, m, n \rangle \rightsquigarrow \langle \alpha'', m'', n'' \rangle}{\langle \texttt{@emitting}(n'), e\alpha, m, n \rangle \rightsquigarrow \langle \alpha'', m'', n'' \rangle} \quad \text{if } |e\alpha| \neq n' \tag{2.10}$$