1 Formal Semantics

1.1 Abstract Syntax

• Simple rules

```
 \begin{array}{lll} \langle p \rangle ::= & \operatorname{mem}(\operatorname{varid}) & (\operatorname{memory\ access\ to\ `varid'}) \\ | & \operatorname{await}(\operatorname{evtid}) & (\operatorname{await\ event\ `evtid'}) \\ | & \operatorname{nop} & (\operatorname{no\ operation}) \end{array}
```

• Compound expressions

• Semantic rules

1.2 Operational Semantics

Relation-inner

$$\langle S, p \rangle \xrightarrow{n} \langle S', p' \rangle$$

where:

$$S,S'\in evtid$$
 (sequence of event identifiers)
$$p,p'\in P$$
 (programs following the abstract syntax)
$$n\in\mathbb{N}$$
 (unique identifier for the reaction chain)

Rules:

$$\langle S, mem(id) \rangle \xrightarrow{n} \langle S, nop \rangle$$
 (mem)

$$\langle S, await(id) \rangle \xrightarrow{n} \langle S, awaiting(id, n+1) \rangle \tag{await}$$

$$\langle id : S, awaiting(id, m) \rangle \xrightarrow{n} \langle id : S, nop \rangle, M < n$$
 (awake)

$$\langle S, emit(id) \rangle \xrightarrow{n} \langle id : S, emitting(|S|) \rangle$$
 (emit)

$$\langle S, emitting(|S|) \rangle \xrightarrow{n} \langle S, nop \rangle$$
 (pop)

Compound expressions:

$$\frac{mem(varid, n) \neq 0}{\langle S, \text{if mem(varid) then p else q} \rangle \xrightarrow{n} \langle S, p \rangle}$$
 (if-true)

$$\frac{val(varid, n) = 0}{\langle S, \text{if mem(varid) then pelse q} \rangle \xrightarrow{n} \langle S, q \rangle}$$
 (if-false)

$$\frac{\langle S, p \rangle \xrightarrow{n} \langle S', p' \rangle}{\langle S, (p; q) \rangle \xrightarrow{n} \langle S', (p'; q) \rangle}$$
(seq-adv)

$$\langle S, (nop; q) \rangle \xrightarrow{n} \langle S, q \rangle$$
 (seq-nop)

$$\langle S, (loop \ p) \rangle \xrightarrow{n} \langle S, (p @ loop \ p) \rangle$$
 (loop-expd)

$$\frac{\langle S, p \rangle \xrightarrow{n} \langle S', p' \rangle}{\langle S, (p @ loop q) \rangle \xrightarrow{n} \langle S', (p' @ loop q) \rangle}$$
(loop-adv)

$$\langle S, (nop @ loop p) \rangle \xrightarrow{n} \langle S, loop p \rangle$$
 (loop-nop)

$$\langle S, (break; p @ loop p) \rangle \xrightarrow{n} \langle S, nop \rangle$$
 (loop-break)

$$\frac{\langle S, p \rangle \xrightarrow{n} \langle S', p' \rangle}{\langle S, (p \ AND \ q) \rangle \xrightarrow{n} \langle S', (p' \ AND \ q) \rangle}$$
(and-adv1)

$$\frac{isBlocked(n, S, p) \quad \langle S, q \rangle \xrightarrow{n} \langle S', q' \rangle}{\langle S, (p \ AND \ q) \rangle \xrightarrow{n} \langle S', (p \ AND \ q' \rangle}$$
(and-adv2)

$$\langle S, (loop \ p \ AND \ q) \rangle \xrightarrow{n} \langle S, (p \bullet loop \ p \ AND \ q) \rangle$$
 (and-adv-loop1)

$$\frac{isBlocked(n, S, p)}{\langle S, (p \ AND \ loop \ q) \rangle \xrightarrow{n} \langle S, (p \ AND \ q \bullet loop \ q) \rangle}$$
 (and-adv-loop2)

$$\langle S, (break; p \bullet loop \ p \ AND \ q) \rangle \xrightarrow{n} \langle S, (clear(q); nop) \rangle$$
 (and-brk1)

$$\frac{isBlocked(n,S,p)}{\langle S,(p \ AND \ break; \ q \bullet loop \ q)\rangle \xrightarrow{n} \langle S,(clear(p);nop)\rangle} \qquad \text{(and-brk2)}$$

$$\frac{\langle S,p\rangle \xrightarrow{n} \langle S',p'\rangle}{\langle S,(p \bullet loop \ q)\rangle \xrightarrow{n} \langle S',(p'\bullet loop \ q)\rangle} \qquad \text{(loop-comp-adv)}$$

$$\langle S,(nop \bullet loop \ p)\rangle \xrightarrow{n} \langle S,p \bullet loop \ p\rangle \qquad \text{(loop-comp-nop)}$$

$$\langle S,(break; \ p \bullet loop \ p)\rangle \xrightarrow{n} \langle S,break\rangle \qquad \text{(loop-comp-break)}$$

$$\frac{\langle S,p\rangle \xrightarrow{n} \langle S',p'\rangle}{\langle S,(p \ OR \ q)\rangle \xrightarrow{n} \langle S',(p' \ OR \ q)\rangle} \qquad \text{(or-adv1)}$$

$$\frac{isBlocked(n,S,p) \quad \langle S,q\rangle \xrightarrow{n} \langle S',q'\rangle}{\langle S,(p \ OR \ q)\rangle \xrightarrow{n} \langle S',(p \ OR \ q')\rangle} \qquad \text{(or-adv2)}$$

$$\langle S,(p \ OR \ nop)\rangle \xrightarrow{n} \langle S,p\rangle \qquad \text{(and-nop1)}$$

$$\langle S,(p \ AND \ nop)\rangle \xrightarrow{n} \langle S,p\rangle \qquad \text{(and-nop2)}$$

$$\langle S,(break; p @ loop \ p \ AND \ q)\rangle \xrightarrow{n} \langle S,clear(q)\rangle \qquad \text{(and-brk1)}$$

$$\langle S,(nop \ OR \ q)\rangle \xrightarrow{n} \langle S,nop\rangle \qquad \text{(or-nop2)}$$