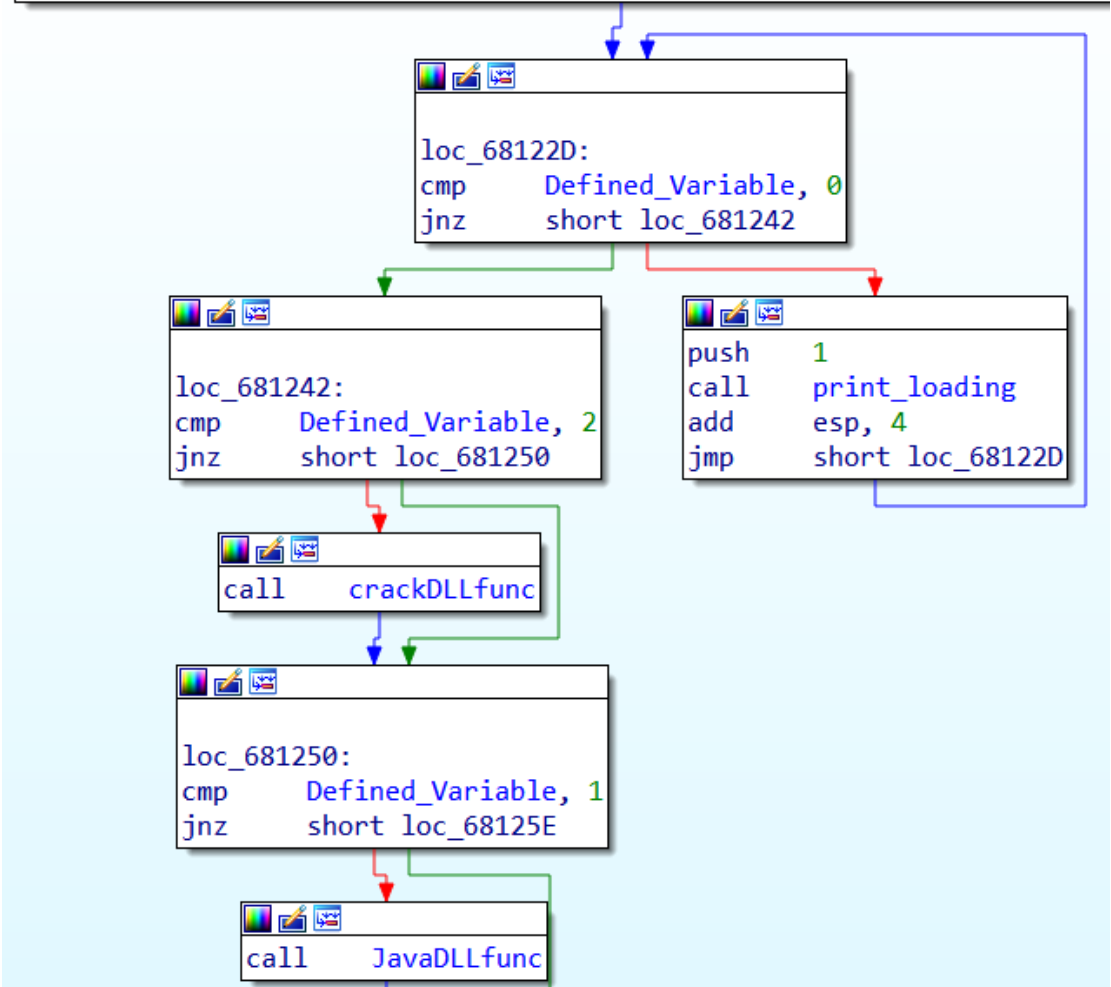


Final Assignment – Reversing 2024 – Roey Gross's Writeup

Python

I opened the exe in IDA, started debugging main and noticed that the flow depends on a variable that has been already miraculously defined, while before the debugging, it was undefined. The running of the program defines it, but where?

```
call    print_loading
add     esp, 4
```



To continue with the main flow and not get stuck at the printing function, we need to find the mysterious part of the program that defines the variable and make sure it defines it to something different from 0.

I have realized that something must be activated before the entry point of main. I looked for some clues and went to the strings window. I found unique strings from a function called TlsCallback_0.

| Address | Length | Type | String |
|------------------|----------|----------|--------------------|
| .rdata:004031... | 00000814 | C | -----:.... |
| .rdata:00403... | 00000009 | C | Java.dll |
| .rdata:00403... | 00000024 | C (16... | SOFTWARE\\GoPython |

This function, creates a new thread, and runs a code that defines with the mysterious variable from the main function:

The screenshot shows assembly code in a debugger. A code window displays the following instructions:

```

push    offset SubKey    ; "SOFTWARE\\GoPython"
push    80000001h        ; hKey
call    ds:RegOpenKeyExW
mov     [ebp+var_20], eax
cmp     [ebp+var_20], 0
jnz     loc_401139

```

Below this, two code windows show the logic of the jump:

- Left window (taken if the key is not found):


```

mov     Defined_Variable, 1
mov     edx, 1
imul    eax, edx, 0

```
- Right window (loc_401139, taken if the key is found):


```

loc_401139:
mov     Defined_Variable, 2

```

BINGO! – Here's where the "predefined" variable is defined! At the Tls_Callback function! The variable is changed based on the existence of a key named "GoPython" inside the registry at the "SOFTWARE" file.

I noticed that there is no use of this registry key at the rest of the program, so I decided to patch the condition.

The screenshot shows the assembly code from the previous block. The instruction `jnz loc_401139` is highlighted. An "Assemble instruction" dialog box is open, showing the previous line's address as `0x1 : 0x4010B1` and the instruction being patched as `jz loc_401139`. Below the dialog, the code window shows the patched instruction:

```

loc_401139:
mov     Defined_Variable, 2

```

Adding keys can be dangerous in some situations, while patching the condition when possible - is safer. Now the defined variable is equal to 1.

Parallely, I wondered how the Tls_Callback function was running before the main, and why my debugger didn't notice it.

I found this recommended article from the digital whisper magazine about TLS callbacks: <https://www.digitalwhisper.co.il/files/Zines/0x62/DW98-2-TLSCallBacks.pdf>

TLS callbacks are created for thread-specific initialization and cleanup tasks. A program can use a TLS callback to execute code before the main entry point of the program. It enables the program to operate covertly without the debugger

notice, as the debugger will typically start at the main entry point and may not be aware of the TLS callback!

Back to the function. The flow inside the TLS callback continues with a decryption algorithm that turns a ciphertext into a plaintext called "Java.dll".

```
.data:00E77018 LibFileName db 'Mfqf)ckk',0  
.data:007A7018 LibFileName db 'Java.dll',0
```

At the end of this decryption, we go into a function that creates and opens for writing a dll file called Java.dll. (If the dll already exists and has content, the function overrides it all. If the dll doesn't exist it creates a new one.)

```
push    offset Mode      ; "wb"  
push    offset FileName  ; "Java.dll"  
lea     eax, [ebp+Stream]  
push    eax              ; Stream  
call    ds:fopen_s
```

Then it writes to the Java.dll a very big number of bytes and closes the stream.

```
mov     ecx, [ebp+Stream]  
push    ecx              ; Stream  
push    2400h            ; ElementCount  
push    1                ; ElementSize  
push    offset unk_683D68 ; Buffer  
call    ds:fwrite  
add     esp, 10h  
mov     [ebp+var_10], eax  
mov     edx, [ebp+Stream]  
push    edx              ; Stream  
call    ds:fclose
```

That's THE END of the TlsCallback function.

To sum up, I found that TLS Callbacks can run before the main function, then to continue without touching the registry I patched a condition, so now the defined variable equals 1. Then it decrypts a ciphertext to "Java.dll" string. Lastly, it creates a file named Java.dll and writes inside huge number of bytes.

While debugging I noticed that the program is multithreaded - Randomly, IDA decided to jump to the beginning and run from there, then it switched back and forth between parts of the program. I noticed that every time it returns to the beginning - it creates a new thread. To disable this mess, I had to suspend every new thread that returned me to the start:

| Threads | | | |
|---------|-------|-----------|--------------|
| Decimal | Hex | State | Name |
| 383296 | 5D940 | Ready | //BD5930 |
| 385596 | 5E23C | Suspended | StartAddress |
| 386348 | 5E52C | Ready | StartAddress |
| 223712 | 369E0 | Suspended | StartAddress |

Another thing that slowed down my debugging process was the graphic printing. It has many Sleep functions that caused the thread to wait a lot:

```

push    64h ; 'd'           ; dwMilliseconds
call    ds:Sleep
push    offset asc_593D00 ; "\b"
push    offset aS_0       ; "%s"
call    printf

```

I patched the condition before the printing, to disable the prints and sleeps.

The screenshot shows a debugger interface with three main components:

- Assembly instruction window (top right):** Displays the previous line at address 0x1: 0x59128C with the instruction `j| loc_591347`. An "OK" button is visible.
- Code window (top left):** Shows assembly code for `loc_591286:`:


```

loc_591286:
mov     ecx, [ebp+var_4]
cmp     ecx, [ebp+arg_0]
jge     loc_591347
      
```
- Code window (bottom left):** Shows assembly code for a function:


```

push    offset asc_593C7C ; "-"
push    offset aS        ; "%s"
call    printf
add     esp, 8
push    64h ; 'd'         ; dwMilliseconds
call    ds:Sleep
      
```
- Code window (bottom right):** Shows assembly code for `loc_591347:`:


```

loc_591347:
mov     esp, ebp
pop     ebp
retn
print_loading endp
      
```

Arrows indicate control flow: a red arrow points from the `jge` instruction in the top-left window to the bottom-left window, and a green arrow points from the `jge` instruction to the bottom-right window.

Java

Back to main. As I mentioned before, the variable changed from 0 to 1. So, the main leads us to a new function I called the "JavaDLL" function.


```
CC 09
XOR
81 53
=
4D 5A <=> M Z
```

BINGO! The first two bytes are decrypted to MZ if the key is 83-51! I found the key for decrypting the DLL!

I wrote a C program that opens Java.dll, decrypts it with the key, and writes back the decryption inside a new DLL called DecJava.dll .

```
int main() {
    xor_decrypt("Java.dll", "DecJava.dll");
    return 0; }
```

"xor_decrypt" calls for the opening of the DLLs, and allocates a buffer for reading from Java.dll. It reads from Java.dll to the buffer and then decrypts the buffer.

This is the important part of the decryption:

```
// Decryption
// XOR each first byte with 0x81 and each second byte with 0x53
char first, second;
for (int i = 0; i < bytes_read; i += 2) {
    first=buffer[i];
    first ^= 0x81;
    buffer[i] = first;
    second = buffer[i+1];
    second ^= 0x53;
    buffer[i + 1] = second;
}
```

Explanation: The Encryption is repeatedly XOR with key of 0x8153. Thus, the decryption is - for every two bytes from the Java.dll, the first one is XORed with 0x81, and the second one is XORed with 0x53.

Then, it writes to the DecJava.dll the decrypted buffer.

```
// Write decrypted buffer to DecJava.dll
fwrite(buffer, 1, bytes_read, output_file);
```

After the decryption DecJava.dll looks normal, with the MZ at the beginning, and repeating zeros.


```

push    offset Format    ; "Password:\n"
movsw
movsb
call    fprintf
add     esp, 4
lea     eax, [esp+70h+Arglist]
push    15h
push    eax              ; Arglist
push    offset aS        ; "%s"
call    fscanf

```

Afterwards, there is a validation checking of the password that has two conditions:

1. The password's length is not 0.
2. The ASCII sum of the characters is equal to 1337.

After the validation there is a building of an URL!!!!

```

movzx   eax, [esp+70h+var_2D]
mov     [esp+70h+ArgList], al
movzx   eax, [esp+70h+var_37]
mov     [esp+70h+ArgList+1], al
movzx   eax, [esp+70h+var_4E]
mov     [esp+70h+ArgList+2], al
movzx   eax, [esp+70h+var_49]
mov     [esp+70h+ArgList+3], al
movzx   eax, [esp+70h+var_5F]
mov     byte ptr [esp+70h+var_20+2], al
lea     eax, [esp+70h+ArgList]
push    eax              ; ArgList
push    offset aBitLyS   ; "bit.ly/%s\n"
call    fprintf

```



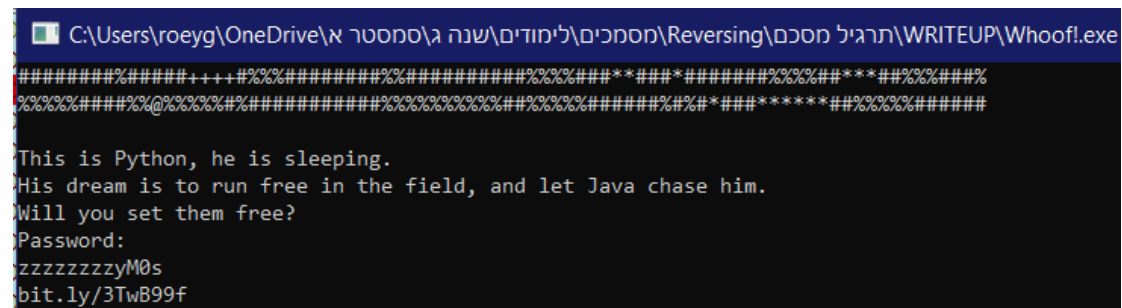
Back to Whoof!.exe main. As mentioned before, it loads the Java.dll, and runs the DLL's XorForever function.

Now, the main loads the old Java.dll. I want it to load the new DLL so I patched the file name to the decrypted one:

```
push    offset NewJava ; "DecJava.dll"
call    ds:LoadLibraryA
mov     [ebp+hModule], eax
cmp     [ebp+hModule], 0
```

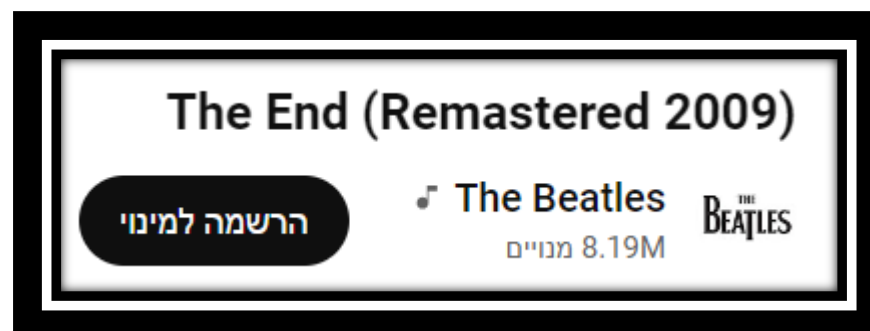
Now we can run main. It loads DecJava.dll and runs the XorForever function that requires a password.

Let's enter a password that applies to all conditions:

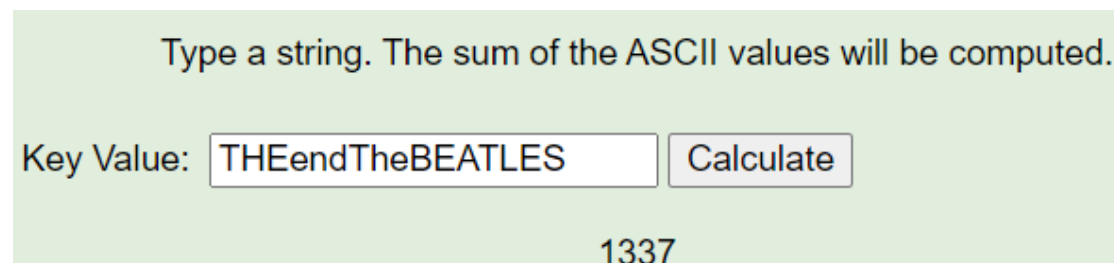


The password passed the checking, the link has been built and printed!

The link leads to the end! BINGO!



Then, I guessed some passwords and realised where 1337 came from...



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