

Syntax and Semantics of V^-

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1 Syntax

NR: I recommend macros for optional and sequence, not literal brackets.

rab: Despite my best efforts, such a macro does not work in grammar mode.

We present a grammar of V^- :

$\langle \text{program} \rangle$	$::= \{ \langle \text{def} \rangle \}$
$\langle \text{def} \rangle$	$::= \text{val } \langle \text{name} \rangle \langle \text{exp} \rangle$ $\langle \text{exp} \rangle$
$\langle \text{exp} \rangle$	$::= \langle \text{integer-literal} \rangle$ $\langle \text{name} \rangle$ $\langle \text{guarded-if} \rangle$ $\langle \text{lambda} \rangle$ $\langle \text{value-constructor-name} \rangle \{ \langle \text{exp} \rangle \}$ $\langle \text{exp} \rangle \mathbin{\ } \langle \text{exp} \rangle$
$\langle \text{lambda} \rangle$	$::= \lambda \{ \langle \text{name} \rangle \} . \langle \text{exp} \rangle$
$\langle \text{guarded-if} \rangle$	$::= \text{if } [\langle \text{guarded-exp} \rangle \{ [] \langle \text{guarded-exp} \rangle \}] \text{ fi}$
$\langle \text{guarded-exp} \rangle$	$::= \rightarrow \langle \text{exp} \rangle$ $\text{E } \{ \langle \text{name} \rangle \} . \langle \text{guarded-exp} \rangle$ $\langle \text{exp} \rangle ; \langle \text{guarded-exp} \rangle$ $\langle \text{name} \rangle = \langle \text{exp} \rangle ; \langle \text{guarded-exp} \rangle$ $\langle \text{exp} \rangle = \langle \text{exp} \rangle ; \langle \text{guarded-exp} \rangle$ $\langle \text{guarded-exp} \rangle \mathbin{\ } \langle \text{guarded-exp} \rangle$ $\text{one}(\{ \langle \text{guarded-exp} \rangle \})$ $\text{all}(\{ \langle \text{guarded-exp} \rangle \})$
$\langle \text{name} \rangle$	$::=$ any token that is not an <i>int-lit</i> , does not contain whitespace, and is not a <i>value-constructor-name</i> or a reserved word.
$\langle \text{value-constructor-name} \rangle$	$::= \text{cons} \mid [] \mid$ any token that begins with a capital letter or a colon
$\langle \text{integer-literal} \rangle$	$::=$ token composed only of digits, possibly prefixed with a + or -.

2 Refinement ordering on environments

$$\begin{aligned} \rho \subseteq \rho' \text{ when } & \text{dom } \rho \subseteq \text{dom } \rho' \\ & \text{and } \forall x \in \text{dom } \rho : \rho(x) \subseteq \rho'(x) \end{aligned}$$

3 Forms of Judgement for V^- :

<i>Metavariables</i>	
v, v'	value
eq	equation
$?t . \{eq\}$	a temporarily-stuck equation
fail	failure
$\rho, \hat{\rho}$	environment: $name \rightarrow \mathcal{V}_\perp$
\mathcal{T}	Context of all temporarily stuck equations (a sequence)
e	An expression
ge	A guarded expression

<i>Sequences</i>	
ε	the empty sequence
$S_1 \cdot S_2$	Concatenate sequence S_1 and sequence S_2
$x \cdot S_2$	Cons x onto sequence S_2

Equations

An *equation* is added to the context of equations temporarily stuck equations or produces **reject**.

$$\langle \rho, eq \rangle \mapsto ?t . \{eq\} \text{ (EQUATIONTEMPSTUCK)}$$

$$\langle \rho, eq \rangle \mapsto \mathbf{reject} \text{ (EQUATIONREJECT)}$$

Expressions

An expression either evaluates to produce a sequence of values or fails. Failure is indicated by the empty sequence ε .

$$\langle \rho, e \rangle \Downarrow v \text{ (EVALSUCC)}$$

$$\langle \rho, e \rangle \Downarrow \varepsilon \text{ (EVALFAIL)}$$

$$\langle \rho, \mathcal{T}, ge \rangle \Downarrow v \text{ (EVALGESUCC)}$$

$$\langle \rho, \mathcal{T}, ge \rangle \Downarrow \varepsilon \text{ (EVALGEFAIL)}$$

Temporarily stuck equations

$$\langle \rho, ?t . \{eq\} \cdot \mathcal{T} \rangle \mapsto \langle \hat{\rho}, \mathcal{T} \rangle \text{ (CTXTORHO)}$$

4 Sequences

The trivial sequence is ε . Sequences can be concatenated with infix \cdot . In an appropriate context, a value like x stands for the singleton sequence containing x .

$$\varepsilon \cdot ys \equiv ys$$

$$ys \cdot \varepsilon \equiv ys$$

$$(xs \cdot ys) \cdot zs \equiv xs \cdot (ys \cdot zs)$$

5 Rules (Big-step Operational Semantics) for V^- :

$$\text{(GE-CTX-STUCK)} \frac{\langle \rho, eq \rangle \mapsto ?t . \{eq\} \quad \langle \rho, ?t . \{eq\} \cdot \mathcal{T} \rangle \mapsto \langle \hat{\rho}, \mathcal{T} \rangle \quad \langle \hat{\rho}, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, \mathcal{T} \cdot eq \cdot \mathcal{T}, ge \rangle \Downarrow v}$$

NR: The notation below has multiple horizontal lines. That makes it a derivation, not a rule. I'm having trouble figuring out what's being said here.

$$\text{(GE-EQ-SUCC)} \frac{\frac{\langle \rho, eq \rangle \mapsto ?t . \{eq\} \quad \langle \rho, ?t . \{eq\} \cdot \mathcal{T} \rangle \mapsto \langle \hat{\rho}, \mathcal{T} \rangle \quad \langle \hat{\rho}, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, \mathcal{T} \cdot eq \cdot \mathcal{T}, ge \rangle \Downarrow v} \text{(GE-CTX-STUCK)}}{\langle \rho, \mathcal{T}, eq; ge \rangle \Downarrow v}$$

NR: If you're trying to pluck an equation out of a list of things, try " $\mathcal{T} \cdot eq \cdot \mathcal{T}'$."

$$\text{(GE-CTX-SUCC)} \frac{\langle \rho, eq \rangle \mapsto \hat{\rho} \text{REMOVETHIS} \quad \langle \hat{\rho}, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, \mathcal{T} \cdot eq \cdot \mathcal{T}, ge \rangle \Downarrow v}$$

$$\text{(GE-EQ-SUCC)} \frac{\frac{\langle \rho, eq \rangle \mapsto \hat{\rho} \text{REMOVETHIS} \quad \langle \hat{\rho}, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, \mathcal{T} \cdot eq \cdot \mathcal{T}, ge \rangle \Downarrow v} \text{(GE-CTX-SUCC)}}{\langle \rho, \mathcal{T}, eq; ge \rangle \Downarrow v}$$

$$\text{(GE-CTX-FAIL)} \frac{\langle \rho, eq \rangle \mapsto \mathbf{reject}}{\langle \rho, eq \cdot \mathcal{T}, ge \rangle \Downarrow \mathbf{fail}}$$

$$\text{(GE-EQ-FAIL)} \frac{\frac{\langle \rho, eq \rangle \mapsto \mathbf{reject}}{\langle \rho, eq \cdot \mathcal{T}, ge \rangle \Downarrow \mathbf{fail}} \text{(GE-CTX-FAIL)}}{\langle \rho, \mathcal{T}, eq; ge \rangle \Downarrow \mathbf{fail}}$$

$$\text{(GE-EXP-FAIL)} \frac{\langle \rho, e \rangle \Downarrow \mathbf{fail}}{\langle \rho, \mathcal{T}, e; ge \rangle \Downarrow \mathbf{fail}}$$

$$\text{(GE-EQ-SUCC)} \frac{\langle \rho, e \rangle \Downarrow v' \quad \langle \rho, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, \mathcal{T}, e; ge \rangle \Downarrow v}$$