

A Syntax of V^-

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We present a grammar of V^- :

$\langle \text{program} \rangle$	$::= \{ \langle \text{def} \rangle \}$
$\langle \text{def} \rangle$	$::= \text{val } \langle \text{name} \rangle \langle \text{exp} \rangle$ $\quad \quad \langle \text{exp} \rangle$
$\langle \text{name} \rangle$	$::=$ any token that is not an <i>int-lit</i> , does not contain whitespace, and is not a <i>value-constructor-name</i> or a reserved word. $\langle \text{value-constructor-name} \rangle ::=$ $\quad \text{cons } \langle \text{exp} \rangle ::= \langle \text{integer-literal} \rangle$ $\quad \quad \langle \text{name} \rangle$ $\quad \quad \langle \text{name} \rangle$ $\quad \quad \langle \text{guarded-if} \rangle$ $\quad \quad \langle \text{lambda} \rangle$
$\langle \text{integer-literal} \rangle$	$::=$ token composed only of digits, possibly prefixed with a + or -.
$\langle \text{lambda} \rangle$	$::= \lambda \{ \langle \text{name} \rangle \} . \langle \text{exp} \rangle$
$\langle \text{guarded-if} \rangle$	$::= \text{if } [\langle \text{guarded-exp} \rangle \{ [] \langle \text{guarded-exp} \rangle \}] \text{ fi}$
$\langle \text{guarded-exp} \rangle$	$::= \langle \text{exp} \rangle$ $\quad \quad \text{E } \{ \langle \text{logical-var} \rangle \} . \langle \text{guarded-exp} \rangle$ $\quad \quad \langle \text{exp} \rangle ; \langle \text{guarded-exp} \rangle$ $\quad \quad \langle \text{logical-var} \rangle = \langle \text{exp} \rangle ; \langle \text{guarded-exp} \rangle$ $\quad \quad \langle \text{exp} \rangle = \langle \text{exp} \rangle ; \langle \text{guarded-exp} \rangle$ $\quad \quad \langle \text{guarded-exp} \rangle \langle \text{guarded-exp} \rangle$ $\quad \quad \text{one}(\{ \langle \text{guarded-exp} \rangle \})$ $\quad \quad \text{all}(\{ \langle \text{guarded-exp} \rangle \})$
$\langle \text{logical-var} \rangle$	$::=$ a fresh name (cannot be lam- or E-bound in this scope).

Add patterns, value constructors (application), choice, one, all.

1 Forms of Judgement for V^- :

Metavariables:

v, v' : value

eq : equation

$?t$: a temporarily-stuck equation

fail : failure

$\rho, \hat{\rho}$: environment: $\text{name} \rightarrow \mathcal{V}_\perp$

\mathcal{T} : Context of all temporarily stuck equations

e : An expression

ge : A guarded expression

Forms of judgement on equations:

$$\langle \rho, eq \rangle \mapsto \hat{\rho} \text{ (EQUATIONSUCCESS)}$$

$$\langle \rho, eq \rangle \mapsto ?t \text{ (EQUATIONTEMPSTUCK)}$$

$$\langle \rho, eq \rangle \mapsto \mathbf{fail} \text{ (EQUATIONFAIL)}$$

Forms of judgement on expressions:

$$\begin{aligned} \langle \rho, e \rangle &\Downarrow v \text{ (EVALSUCC)} \\ \langle \rho, e \rangle &\Downarrow \mathbf{fail} \text{ (EVALFAIL)} \\ \langle \rho, \mathcal{T}, ge \rangle &\Downarrow v \text{ (EVALGESUCC)} \\ \langle \rho, \mathcal{T}, ge \rangle &\Downarrow \mathbf{fail} \text{ (EVALGEFAIL)} \end{aligned}$$

Other forms:

$$\langle \rho, ?t; \mathcal{T} \rangle \mapsto \langle \hat{\rho}, \mathcal{T} \rangle \text{ (CTXTORHO)}$$

In english:

An equation is either solved to produce bindings that extend an environment, gets temporarily stuck, or fails.

An expression either evaluates to produce a value or fails.

Other important guidelines (where do we put these?):

$$\begin{aligned} \rho &\subseteq \rho' \\ \text{when } \text{dom } \rho &\subseteq \text{dom } \rho' \\ \text{and } \forall x \in \text{dom } \rho : &\rho(x) \subseteq \rho'(x) \end{aligned}$$

When $\langle \rho, e \rangle \mapsto \rho'$, then $\rho \subseteq \rho'$.

2 Rules (Big-step Operational Semantics) for V^- :

$$\begin{aligned} \text{(GE-CTX-STUCK)} \quad & \frac{\langle \rho, eq \rangle \mapsto ?t \quad \langle \rho, ?t; \mathcal{T} \rangle \mapsto \langle \hat{\rho}, \mathcal{T} \rangle \quad \langle \rho, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, eq; \mathcal{T}, ge \rangle \Downarrow v} \\ \text{(GE-EQ-SUCC)} \quad & \frac{\langle \rho, eq \rangle \mapsto ?t \quad \langle \rho, ?t; \mathcal{T} \rangle \mapsto \langle \hat{\rho}, \mathcal{T} \rangle \quad \langle \rho, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, eq; \mathcal{T}, ge \rangle \Downarrow v} \text{ (GE-CTX-STUCK)} \\ & \frac{\langle \rho, eq; \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, \mathcal{T}, eq; ge \rangle \Downarrow v} \\ \text{(GE-CTX-SUCC)} \quad & \frac{\langle \rho, eq \rangle \mapsto \hat{\rho} \quad \langle \hat{\rho}, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, eq; \mathcal{T}, ge \rangle \Downarrow v} \\ \text{(GE-CTX-SUCC)} \quad & \frac{\langle \rho, eq \rangle \mapsto \hat{\rho} \quad \langle \hat{\rho}, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, eq; \mathcal{T}, ge \rangle \Downarrow v} \\ \text{(GE-EQ-SUCC)} \quad & \frac{\langle \rho, eq; \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, \mathcal{T}, eq; ge \rangle \Downarrow v} [Right] \\ \text{(GE-CTX-FAIL)} \quad & \frac{\langle \rho, eq \rangle \mapsto \mathbf{fail}}{\langle \rho, eq; \mathcal{T}, ge \rangle \Downarrow \mathbf{fail}} \\ \text{(GE-EQ-FAIL)} \quad & \frac{\langle \rho, eq \rangle \mapsto \mathbf{fail}}{\langle \rho, eq; \mathcal{T}, ge \rangle \Downarrow \mathbf{fail}} \text{ (GE-CTX-FAIL)} \\ & \frac{\langle \rho, eq; \mathcal{T}, ge \rangle \Downarrow \mathbf{fail}}{\langle \rho, \mathcal{T}, eq; ge \rangle \Downarrow \mathbf{fail}} \\ \text{(GE-EXP-FAIL)} \quad & \frac{\langle \rho, e \rangle \Downarrow \mathbf{fail}}{\langle \rho, \mathcal{T}, e; ge \rangle \Downarrow \mathbf{fail}} \\ \text{(GE-EQ-SUCC)} \quad & \frac{\langle \rho, e \rangle \Downarrow v' \quad \langle \rho, \mathcal{T}, ge \rangle \Downarrow v}{\langle \rho, \mathcal{T}, e; ge \rangle \Downarrow v} \end{aligned}$$