Lecture 5 Worksheet: Complexity Analysis

A **step** is any unit of work with bounded execution time (it doesn't keep getting slower with growing input size).

We classify algorithm complexity by classifying the **order of growth** of a function f(N), where f gives the number of steps the algorithm must perform for a given input size.

Big O definition: if **$f(N) \le C * g(N)$** for large **N** values and some fixed constant **C**, then **$f(N) \in O(g(N))$**

1

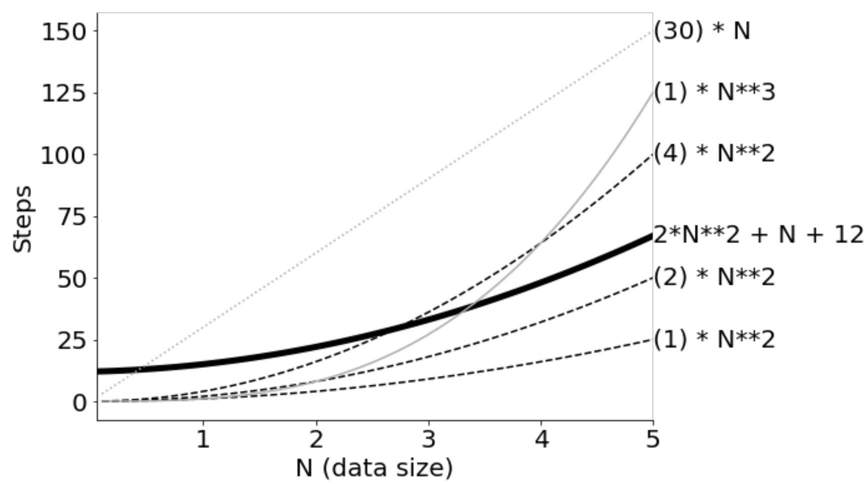
Let $f(N) = 2N^2 + N + 12$

If we want to show $f(N) \in O(N^3)$, what is a good lower bound on N? Let's have C=1.

If we want to show $f(N) \in O(N^2)$, do we pick 1, 2, or 4 for the C? After picking C, should we choose for N's lower bound?

What is more informative to show?
$f(N) \in O(N^3)$ or $f(N) \in O(N^2)$?

Somebody claims $f(N) \in O(N)$, offering C=30 and N>0. Suggest an N value to counter their claim.



2

Each of the following list operations are either O(1) or O(N), where N is len(L). Circle those you think are O(N).

```
L.insert(0, x)    L.pop(0)     x = L[0]        x = max(L)     x = len(L)

L.append(x)       L.pop(-1)    L2.extend(L)    x = sum(L)     found = X in L
```

3

```
def search(L, target):
   for x in L:
      if x == target: #line A
         return True
   return False
```

assume this is asked unless otherwise stated

Let **f(N)** be the number of times line A executes, with `N=len(L)`. What is **f(N)** in each case?

Worst Case (target is at end of list): f(N) = __________

Best Case (target is at beginning of list): f(N) = __________

Average Case (target in middle of list): f(N) = __________

4

```
def selection_sort(L):
    for i in range(len(L)):
        idx_min = i
        for j in range(i, len(L)):
            if L[j] < L[i]:
                idx_min = j
        L[idx_min], L[i] = L[i], L[idx_min] # swap values at i and idx_min

nums = [2, 4, 3, 1]
selection_sort(nums)
print(nums)
```

if this runs **f(N)** times, where **N=len(L)**,

then **f(N)** = ______________________

The complexity of selection sort is

O(____________)

5

```
# assume L is already sorted, N=len(L)
def binary_search(L, target):
    left_idx = 0 # inclusive
    right_idx = len(L) # exclusive
    while right_idx - left_idx > 1:
        mid_idx = (right_idx + left_idx) // 2
        mid = L[mid_idx]
        if target >= mid:
            left_idx = mid_idx
        else:
            right_idx = mid_idx

    return right_idx > left_idx and L[left_idx] == target
```

how many times does this step run when N = 1? N = 2? N = 4? N = 8?

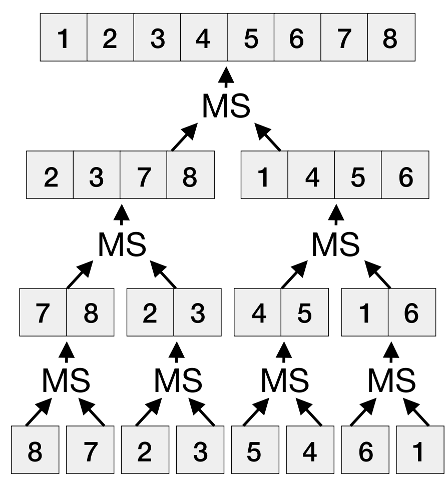
If f(N) is the number of times this step runs, then f(N) = ____________

The complexity of binary search is O(____________)

6

```
def merge(L1, L2):
  rv = []
  idx1 = 0
  idx2 = 0

  while True:
    done1 = idx1 == len(L1)
    done2 = idx2 == len(L2)

    if done1 and done2:
      return rv

    choose1 = False
    if done2:
      choose1 = True
    elif not done1 and L1[idx1] < L2[idx2]:
      choose1 = True

    if choose1:
      rv.append(L1[idx1])
      idx1 += 1
    else:
      rv.append(L2[idx2])
      idx2 += 1

  return rv
```

```
def merge_sort(L):
  if len(L) < 2:
    return L
  mid = len(L) // 2
  left = L[:mid]
  right = L[mid:]
  left = merge_sort(left)
  right = merge_sort(right)
  return merge(left, right)

merge_sort([4, 1, 2, 3])
```



`merge([1, 3], [2, 4])` will return ______________.

`merge(L1, L2)` implements an O(N) algorithm. But how can we measure the size of the input? N = ______________.

If we double the list size, there will be ___ more level(s). Level count grows O(______). Work per level is O(______). merge_sort complexity: O(_________)

7

```
nums = [...]

first100sum = 0

for x in nums[:100]:
    first100sum += x
print(x)
```

If we increase the size of nums from 20 items to 100 items, the code will probably take _______ times longer to run.

If we increase the size of nums from 100 to 1000, will the code take longer? Yes / No

The complexity of the code is O(_____), with N=len(nums).