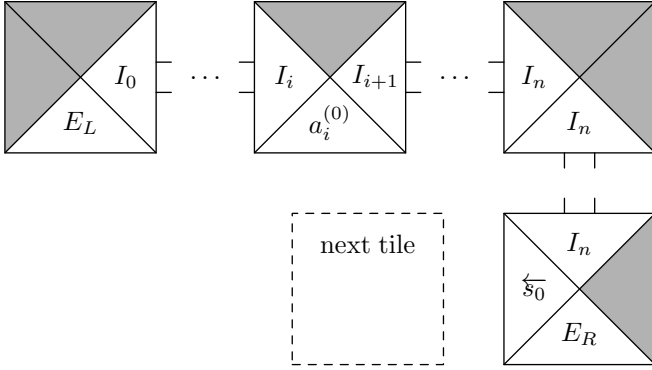
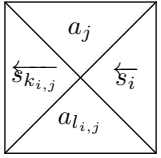


Assembly of input tape. E_L is left stop, $a_i^{(0)}$ is i -th symbol of input word $a_0^{(0)} \dots a_{n-1}^{(0)}$. Following assembly starts at a place denoted “next tile” while simulating a Turing machine reading $a_{n-1}^{(0)}$ and being in state s_0 . The arrow over s_0 means “comes from right”.



Coming from right, being in state s_i , reading tape symbol a_j .
Transition function says: write $a_{l_{i,j}}$, switch to state $s_{k_{i,j}}$ and go left.



Situation is like before with only difference: go right.
Now the rest of the tape must be copied by special tiles which thus exist for all pairs $a_m, s_{k_{i,j}}$ of tape symbol and state, respectively.

