PROVER-SKEPTIC GAMES AND LOGICAL PLURALISM

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A Problem for Logical

PLURALISTS

THE EXPLANATORY PROBLEM FOR LOGICAL PLURALISTS

A characterisation of a pluralist position needs to explain what it is to endorse all of the various consequence relations the pluralist accepts and how they relate to an intuitive notion of logical consequence.

—Keefe, 2014, p.1376

Two Central Questions

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- ► Why are these consequence relations relations of logical consequence?

THE BUILT-IN OPPONENT

CONCEPTION OF DEDUCTION

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MULTI-AGENT VIEW OF LOGICAL CONSEQUENCE

- ► Logical Consequence grounded in specialized discursive practices.
- View **Proofs** as semi-adversarial debates between a Prover and a Skeptic.
- ▶ Prover seeks to establish the conclusion.
- ► Skeptic forces Prover to make their reasons explicit.
- ► Deductive proofs are winning strategies for Prover in these debates.

Internalization of Skeptic

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- ► This is because Skeptic has become **Internalized** to the method, a very part of these practices.
- ► In effect their role is played 'offline' by Prover.
- ► This provides a bridge between Multi-Agent Dialogical and Mono-Agent Inferential Practices.

Introducing

PROVER-SKEPTIC GAMES

In the library two students, **Petunia** and **Samson** are arguing over the validity of the argument from $A \to B$ and $B \to C$ to $A \to C$. Petunia thinks that this argument is valid and is trying to convince Samson, who is skeptical.

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▶ **Petunia (1)**: I reckon that $A \to C$ follows from $A \to B$ and $B \to C$

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- ▶ **Petunia (1)**: I reckon that $A \to C$ follows from $A \to B$ and $B \to C$
- Samson (1): Yeah? Well if that's so then suppose I grant you A → B and B → C along with A, how are you mean to get C?

▶ Petunia (2): If you grant me B I can get C from B → C (which you just granted).

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- ► **Samson (2)**: But why should I grant you B?
- Petunia (3): Well if you were to grant me A then I could get B from A → B which you granted at the start.

- ▶ Petunia (2): If you grant me B I can get C from B → C (which you just granted).
- ► **Samson (2)**: But why should I grant you B?
- Petunia (3): Well if you were to grant me A then I could get B from A → B which you granted at the start.
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- ▶ Petunia (2): If you grant me B I can get C from B → C (which you just granted).
- ▶ **Samson (2)**: But why should I grant you B?
- Petunia (3): Well if you were to grant me A then I could get B from A → B which you granted at the start.
- ► **Samson (3)**: But why should I grant you A?
- ▶ **Petunia (4)**: Because you granted it to me at the start!

- ▶ Petunia (2): If you grant me B I can get C from B → C (which you just granted).
- ► **Samson (2)**: But why should I grant you B?
- Petunia (3): Well if you were to grant me A then I could get B from A → B which you granted at the start.
- ► **Samson (3)**: But why should I grant you A?
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Petunia leaves the library triumphantly.

Definition: A *dialogue* over the sequent $\Gamma \succ \phi$ is a possibly infinite sequence $\Sigma_1, \mathfrak{s}_1, \Sigma_2, \mathfrak{s}_2, \ldots$ where:

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$$\Sigma_1 = \{\Gamma \succ \phi\}$$

» Prover Begins by asserting the sequent under discussion.

DEFINITION: A *dialogue* over the sequent $\Gamma \succ \phi$ is a possibly infinite sequence $\Sigma_1, \mathfrak{s}_1, \Sigma_2, \mathfrak{s}_2, \ldots$ where:

- 1. $\Sigma_1 = \{\Gamma \succ \phi\}$
- 2. $\mathfrak{s}_i \Rightarrow_{(I)}^* \sigma$, for some non-axiomatic sequent $\sigma \in \Sigma_i$ where \mathfrak{s}_i is of the form $S_i \succ p_i$ for some propositional atom p_i .

» Skeptic challenges an (interesting) Prover assertion

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- 3. $\Sigma_{i+1} \Rightarrow_{(E)}^* \mathfrak{s}_i$
 - » Prover replies, offering sequents which derive the challenge.

Some Examples

REVISITING THE PREVIOUS EXAMPLE

Consider the following dialogue over $p \to q, q \to r \,{\succ}\, p \to r.$

P(1)
$$[p \rightarrow q, q \rightarrow r \succ p \rightarrow r]$$

» Prover Starts

REVISITING THE PREVIOUS EXAMPLE

Consider the following dialogue over $p \to q, q \to r \succ p \to r$.

P(1)
$$[p \rightarrow q, q \rightarrow r \rightarrow p \rightarrow r]$$

S(1)
$$p \rightarrow q, q \rightarrow r, p \rightarrow r$$

» Skepticchallenges thesole Proverassertion.

Consider the following dialogue over $p \rightarrow q$, $q \rightarrow r > p \rightarrow r$.

P(1)
$$[p \rightarrow q, q \rightarrow r \succ p \rightarrow r]$$

S(1)
$$p \rightarrow q, q \rightarrow r, p \succ r$$

P(2)
$$[q \rightarrow r \rightarrow q \rightarrow r, p \rightarrow q, p \rightarrow q]$$

» Prover replies, as these sequents derive the challenge

sequent.

Consider the following dialogue over $p \rightarrow q$, $q \rightarrow r > p \rightarrow r$.

$$\begin{array}{lll} P(1) & [p \rightarrow q, q \rightarrow r \succ p \rightarrow r] & \text{$>$} & \textit{Skeptic} \\ S(1) & p \rightarrow q, q \rightarrow r, p \succ r & \textit{challenges the} \\ P(2) & [q \rightarrow r \succ q \rightarrow r, p \rightarrow q, p \succ q] & \textit{only} \\ S(2) & p \rightarrow q, p \succ q & \textit{non-axiomatic} \\ & & \textit{sequent.} \end{array}$$

Consider the following dialogue over $p \to q, q \to r \succ p \to r$.

P(1)
$$[p \rightarrow q, q \rightarrow r \rightarrow p \rightarrow r]$$

S(1)
$$p \rightarrow q, q \rightarrow r, p \succ r$$

P(2)
$$[q \rightarrow r \succ q \rightarrow r, p \rightarrow q, p \succ q]$$

S(2)
$$p \rightarrow q, p \succ q$$

P(3)
$$[p \rightarrow q \succ p \rightarrow q, p \succ p]$$

» Prover replies, and wins the dialogue.

Consider the following dialogue over $p \to q, q \to r \,{\succ}\, p \to r.$

P(1)
$$[p \rightarrow q, q \rightarrow r \rightarrow p \rightarrow r]$$

S(1)
$$p \rightarrow q, q \rightarrow r, p \succ r$$

P(2)
$$[q \rightarrow r \succ q \rightarrow r, p \rightarrow q, p \succ q]$$

S(2)
$$p \rightarrow q, p \succ q$$

P(3)
$$[p \rightarrow q \succ p \rightarrow q, p \succ p]$$

Consider the dialogue for the sequent $p \to r \,{\succ}\, p \to (q \to r)$

P(1)
$$[p \rightarrow r \succ p \rightarrow (q \rightarrow r)]$$

» Prover Starts

Consider the dialogue for the sequent $p \to r \,{\succ}\, p \to (q \to r)$

P(1)
$$[p \rightarrow r > p \rightarrow (q \rightarrow r)]$$

S(1)
$$p \rightarrow r, p, q \succ r$$

» Skeptic challenges, saying what would need to be shown to convince them.

Consider the dialogue for the sequent $p \to r \,{\succ}\, p \to (q \to r)$

P(1)
$$[p \rightarrow r > p \rightarrow (q \rightarrow r)]$$

S(1)
$$p \rightarrow r, p, q \succ r$$

P(2)
$$[p \rightarrow r \rightarrow p \rightarrow r, p \rightarrow p]$$

» From these two sequents Prover can derive the challenge by one application of $(\rightarrow E)$ and then (K)to weaken in q.

Consider the dialogue for the sequent $p \to r \,{\succ}\, p \to (q \to r)$

P(1)
$$[p \rightarrow r > p \rightarrow (q \rightarrow r)]$$

S(1)
$$p \rightarrow r, p, q \succ r$$

P(2)
$$[p \rightarrow r \rightarrow p \rightarrow r, p \rightarrow p]$$

» This wins the dialogue for Prover in the presence of (K).

Consider the dialogue for the sequent $p \to r \,{\succ}\, p \to (q \to r)$

P(1)
$$[p \rightarrow r > p \rightarrow (q \rightarrow r)]$$

S(1)
$$p \rightarrow r, p, q \succ r$$

P(2)
$$[p \rightarrow r > p \rightarrow r, p, q > p]$$

» The 'best' move Prover

has available without

(K) is the following

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Consider the dialogue for the sequent $p \to r \,{\succ}\, p \to (q \to r)$

P(1)
$$[p \rightarrow r > p \rightarrow (q \rightarrow r)]$$

S(1)
$$p \rightarrow r, p, q \succ r$$

P(2)
$$[p \rightarrow r \succ p \rightarrow r, p, q \succ p]$$

S(2)
$$p, q \succ p$$

» Skeptic can only challenge the rightmost sequent.

Consider the dialogue for the sequent $p \to r \,{\succ}\, p \to (q \to r)$

P(1)
$$[p \rightarrow r > p \rightarrow (q \rightarrow r)]$$

S(1)
$$p \rightarrow r, p, q \succ r$$

P(2)
$$[p \rightarrow r \succ p \rightarrow r, p, q \succ p]$$

S(2)
$$p, q > p$$

P(3)
$$[q \succ p, p \succ p \rightarrow p]$$

» This is Prover's 'best' response, and things will quite quickly devolve from here...

Consider the dialogue for the sequent $p \to (p \to q) \succ p \to q$)

P(1)
$$[p \rightarrow (p \rightarrow q) \succ p \rightarrow q]$$

» Prover starts.

Consider the dialogue for the sequent $p \rightarrow (p \rightarrow q) \succ p \rightarrow q$)

P(1)
$$[p \rightarrow (p \rightarrow q) \succ p \rightarrow q]$$

S(1)
$$p \rightarrow (p \rightarrow q), p \rightarrow q$$

» Skeptic challenges, saying what would need to be shown to convince them.

Consider the dialogue for the sequent $p \to (p \to q) \succ p \to q$)

P(1)
$$[p \rightarrow (p \rightarrow q) > p \rightarrow q]$$

S(1)
$$p \rightarrow (p \rightarrow q), p \rightarrow q$$

$$P(2) \hspace{0.2cm} [p \rightarrow (p \rightarrow q) \! \succ \! p \rightarrow (p \rightarrow q), p \! \succ \! p, p \! \succ \! p]$$

» Prover responds, deriving the challenge via two applications of $(\rightarrow E)$ and one of (W).

Consider the dialogue for the sequent $p \to (p \to q) \succ p \to q$)

P(1)
$$[p \rightarrow (p \rightarrow q) \succ p \rightarrow q]$$

S(1)
$$p \rightarrow (p \rightarrow q), p \rightarrow q$$

$$P(2) \quad [p \to (p \to q) \succ p \to (p \to q), p \succ p, p \succ p]$$

This wins the dialogue for Prover in the presence of (W).

Consider the dialogue for the sequent $p \rightarrow (p \rightarrow q) \succ p \rightarrow q$)

P(1)
$$[p \rightarrow (p \rightarrow q) \succ p \rightarrow q]$$

S(1)
$$p \rightarrow (p \rightarrow q), p \rightarrow q$$

P(2)
$$[p \rightarrow (p \rightarrow q) \succ p \rightarrow q, p \succ p]$$

» In the absence of(W) this is the'best' moveavailable toProver.

Consider the dialogue for the sequent $p \to (p \to q) \succ p \to q$)

P(1)
$$[p \rightarrow (p \rightarrow q) > p \rightarrow q]$$

S(1)
$$p \rightarrow (p \rightarrow q), p \rightarrow q$$

P(2)
$$[p \rightarrow (p \rightarrow q) \succ p \rightarrow q, p \succ p]$$

S(2)
$$p \rightarrow (p \rightarrow q), p \succ q$$

» Skeptic must challenge

$$p \to (p \to q) {\succ} p \to q$$

Consider the dialogue for the sequent $p \to (p \to q) \succ p \to q$)

P(1)
$$[p \rightarrow (p \rightarrow q) \succ p \rightarrow q]$$

S(1)
$$p \rightarrow (p \rightarrow q), p \rightarrow q$$

P(2)
$$[p \rightarrow (p \rightarrow q) \succ p \rightarrow q, p \succ p]$$

S(2)
$$p \rightarrow (p \rightarrow q), p \succ q$$

» But this just puts us back in the position we were in at S(1)

STRUCTURAL RULES AND

NORMS ON EXPLANATION

PROVER-SKEPTIC GAMES FOR

SUB-STRUCTURAL LOGICS

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ADEQUACY OF

PROVER-SKEPTIC GAMES

ADEQUACY: A SKETCH

▶ Winning Strategies correspond to proof sketches.

ADEQUACY: A SKETCH

- ▶ Winning Strategies correspond to proof sketches.
- ► Constructing a Winning Strategy given a proof is a relatively simple operation.

FROM PROOFS TO WINNING STRATEGIES

$$\begin{array}{c} \frac{p \times p \quad p \to q \times p \to q}{p, p \to q \times q} \xrightarrow{(\to E)} \frac{p \to (q \to r) \times p \to (q \to r) \quad p \times p}{p, p \to (q \to r) \times q \to r} \xrightarrow{(\to E)} \\ \frac{p \to (q \to r), p \to q, p, p \times r}{(K)} \xrightarrow{(K)} \\ \frac{p \to (q \to r), p \to q, p \times r}{p \to (q \to r), p \to q \times p \to r} \xrightarrow{(\to I)} \xrightarrow{(\to I)} \\ p \to (q \to r) \times (p \to q) \to (p \to r) \end{array}$$

FROM PROOFS TO WINNING STRATEGIES

$$\begin{array}{c} \frac{p \mathrel{\succ} p \quad p \mathrel{\rightarrow} q \mathrel{\succ} p \mathrel{\rightarrow} q}{p, p \mathrel{\rightarrow} q \mathrel{\succ} q} \xrightarrow{(\mathrel{\rightarrow} E)} & \frac{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r) \mathrel{\succ} p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r)}{p, p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r) \mathrel{\succ} q \mathrel{\rightarrow} r} \xrightarrow{(\mathrel{\rightarrow} E)} \\ \\ \frac{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r), p \mathrel{\rightarrow} q, p, p \mathrel{\succ} r}{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r), p \mathrel{\rightarrow} q, p \mathrel{\succ} r} \xrightarrow{(\mathrel{\rightarrow} I)} \xrightarrow{(\mathrel{\rightarrow} I)} \\ \\ \frac{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r), p \mathrel{\rightarrow} q, p \mathrel{\succ} r}{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r), p \mathrel{\rightarrow} q \mathrel{\succ} p \mathrel{\rightarrow} r} \xrightarrow{(\mathrel{\rightarrow} I)} \xrightarrow{(\mathrel{\rightarrow} I)} \end{array}$$

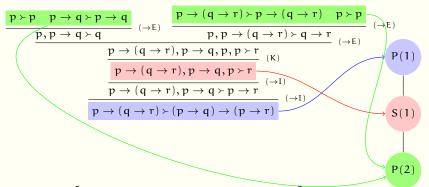
$$P(1) = \left[p \to (q \to r) \succ (p \to q) \to (p \to r) \right]$$

FROM PROOFS TO WINNING STRATEGIES

$$\begin{array}{c|c} \underline{p \mathrel{\succ} p \quad p \mathrel{\rightarrow} q \mathrel{\succ} p \mathrel{\rightarrow} q} \\ \underline{p, p \mathrel{\rightarrow} q \mathrel{\succ} q} \\ \underline{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r), p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r) \mathrel{\succ} p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r)}_{(\mathrel{\rightarrow} E)} \\ \underline{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r), p \mathrel{\rightarrow} q, p, p \mathrel{\succ} r}_{(\mathrel{\leftarrow} I)} \\ \underline{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r), p \mathrel{\rightarrow} q, p \mathrel{\succ} r}_{(\mathrel{\rightarrow} I)} \\ \underline{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r), p \mathrel{\rightarrow} q \mathrel{\succ} p \mathrel{\rightarrow} r}_{(\mathrel{\rightarrow} I)} \\ \underline{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r), p \mathrel{\rightarrow} q \mathrel{\succ} p \mathrel{\rightarrow} r}_{(\mathrel{\rightarrow} I)} \\ \underline{p \mathrel{\rightarrow} (q \mathrel{\rightarrow} r) \mathrel{\succ} (p \mathrel{\rightarrow} q) \mathrel{\rightarrow} (p \mathrel{\rightarrow} r)}_{(\mathrel{\rightarrow} I)} \\ \end{array}} \begin{array}{c} (\mathrel{\rightarrow} E) \\ ($$

$$P(1) = \left[p \to (q \to r) \succ (p \to q) \to (p \to r) \right]$$

From Proofs to Winning Strategies



$$P(1) = \left[p \to (q \to r) \succ (p \to q) \to (p \to r) \right]$$

$$> S(1) = p \rightarrow (q \rightarrow r), p \rightarrow q, p \succ r$$

▶
$$P(2) =$$

$$[p \to (q \to r) \succ p \to (q \to r) , p \succ p , p \to q \succ p \to q]$$



THANK YOU!