```
For Divide & Conquer
                  Master Theorem: If a >0, b >1 and of ≥0 are constants the
      solutions of a recurrence of the form T(n) = (a T(\frac{n}{6}) + O(\frac{n}{6})) n > n_0 for some
                                                                                                            constant > 0 if n \le no

\begin{cases}
\frac{O(n^d \log n)}{O(n^{\log n})} & \text{if } a/b^d < 1 \\
\frac{O(n^d \log n)}{O(n^{\log n})} & \text{if } a/b^d > 1
\end{cases}

                                                                                                                                    DIF not, re is at most a const
                 T(n) = \int O(\Pi^d)
                                                                                                                                          factor away from a power
                                                                                                                                        1 of b. ⇒ No effect on T(n) in O()
                                                                                                                                       Assume: The is a power
                                                                                                                                                   of b i.e. n = b^{t}
                Let T(n) = a T(1/b) + c nd for some constant c > 0
                                                                                                                                                  for some t
                                                                                                                                       Also, let T(1) = C, for
                                                                                                                                                 some constant 0,>0
        \Rightarrow T(n) = a T(\frac{\pi}{b}) + c n^d
                         = \alpha \left[ \alpha T \left( \frac{\Gamma}{b^2} \right) + c \left( \frac{\Gamma}{b} \right)^d \right] + c n^d = \alpha^2 T \left( \frac{\Gamma}{b^2} \right) + c n^d \left( 1 + \frac{\alpha}{b^d} \right)
                        =a^{2}\left(aT\left(\frac{n}{b^{3}}\right)+c\left(\frac{n}{b^{2}}\right)^{d}\right)+cn^{d}\left(1+\frac{a}{b^{d}}\right)=a^{3}T\left(\frac{n}{b^{3}}\right)+cn^{d}\left(1+\frac{a}{b^{d}}+\left(\frac{a}{b^{d}}\right)^{2}\right)
                         = a^{\frac{1}{b}} T\left(\frac{n}{b^{\frac{1}{b}}}\right) + cn^{\frac{1}{b}} \left(1 + \frac{a}{b^{\frac{1}{d}}} + \left(\frac{a}{b^{\frac{1}{d}}}\right)^{2} + \cdots + \left(\frac{a}{b^{\frac{1}{d}}}\right)^{\frac{1}{d}}\right)
    When n = b^{t} is t = \log_{b} n. a^{t} T(\frac{n}{b^{t}}) = c, a^{t} = O(a^{t}) = O(a^{\log_{b} n}) = O(n^{\log_{b} a})
     Case | a/b^d < 1 is \log_b a < d \Rightarrow a^t T(\frac{n}{b^t}) = \theta(n^{\log_b a}) = O(n^d) - 0
                                           \operatorname{Cn}^{d}\left(1+\frac{a}{b^{d}}+\left(\frac{a}{b^{d}}\right)^{2}+\cdots\right)\leq\operatorname{Cn}^{d}\frac{1}{1-\frac{a}{1d}}=O\left(\operatorname{n}^{d}\right)-2
                                       From P. 2 T(n) = O(nd)
   Case 2: a/b^d=1 is \log_b a=d \Rightarrow a^t T(\frac{\pi}{b^t})=\Theta(\pi^{\log_b a})=\Theta(\pi^d)
                          \operatorname{Cnd}\left(1+\frac{a}{b^{a}}+\left(\frac{a}{b^{a}}\right)^{2}+\cdots+\left(\frac{a}{b^{a}}\right)^{t-1}\right)=O\left(\operatorname{nd}\left(\log n\right)-2\right)
                                      From O. D T(n) = O(ndlogn) *
   Case 3: a/b^d > 1 is \log_b a > d \Rightarrow a^t T(\frac{\pi}{b^t}) = O(\pi^{\log_b a}) - D
                         \operatorname{Cnd}\left(1+\frac{a}{b^{a}}+\left(\frac{a}{b^{d}}\right)^{2}+\cdots+\left(\frac{a}{b^{d}}\right)^{t-1}\right)=\operatorname{Cnd}\left(\frac{a}{b^{d}}\right)^{t}-1=\operatorname{O}\left(\operatorname{nd}\left(\frac{a^{t}}{b^{d}}\right)^{t}-1\right)
                   From O O T(n) = O(n^{\log_b q}) \otimes = O(n^d \cdot \frac{a^t}{n^d}) = O(a^t) = O(a^{\log_b n}) = O(n^{\log_b q}) = O(n^{\log_b q})
Note: a^{\log_b n} = a^{\left(\frac{\log n}{\log_a n}, \frac{\log a}{\log_b n}\right)} = a^{\left(\log_a n\right)\left(\log_b a\right)} = \left(a^{\log_a n}\right)^{\log_b a} = n^{\log_b a}
```