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# Intervals Intersection (medium)

#### We'll cover the following

- Problem Statement
- Try it yourself
- Solution
- Code
  - Time complexity
  - Space complexity

## Problem Statement #

Given two lists of intervals, find the **intersection of these two lists**. Each list consists of **disjoint intervals sorted on their start time**.

#### Example 1:

Input: arr1=[[1, 3], [5, 6], [7, 9]], arr2=[[2, 3], [5, 7]]

Output: [2, 3], [5, 6], [7, 7]

Explanation: The output list contains the common intervals betwe

en the two lists.

#### Example 2:

```
Input: arr1=[[1, 3], [5, 7], [9, 12]], arr2=[[5, 10]]
Output: [5, 7], [9, 10]
Explanation: The output list contains the common intervals betwe en the two lists.
```

# Try it yourself #

Try solving this question here:

```
Pvthon3
👙 Java
                                     C++
                         JS JS
    class Interval {
 1
       constructor(start, end) {
 2
 3
         this.start = start;
 4
         this.end = end;
 5
       }
 6
 7
       print_interval() {
         process.stdout.write(`[${this.start}, ${this.end}]`);
 8
       }
 9
10
    }
11
12
    const merge = function(intervals_a, intervals_b) {
      let result = [];
13
      // TODO: Write your code here
14
15
       return result;
16
    };
17
    process.stdout.write('Intervals Intersection: ');
18
19
    let result = merge([new Interval(1, 3), new Interval(5, 6), new Interval(7, 6)]
20
    for (i = 0; i < result.length; i++) {
21
       result[i].print_interval();
22
23
    console.log();
24
25
    process.stdout.write('Intervals Intersection: ');
```

```
26 result = merge([new Interval(1, 3), new Interval(5, 7), new Interval(9, 12 27 for (i = 0; i < result.length; i++) {
28    result[i].print_interval();

Run

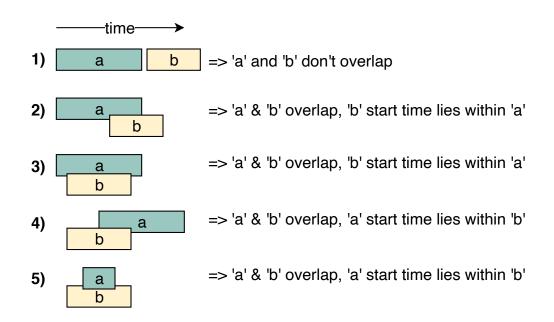
Save
Reset []
```

## Solution #

This problem follows the Merge Intervals

(https://www.educative.io/collection/page/5668639101419520/56714648543 55968/5652017242439680/) pattern. As we have discussed under Insert Interval

(https://www.educative.io/collection/page/5668639101419520/56714648543 55968/5718314357620736/), there are five overlapping possibilities between two intervals 'a' and 'b'. A close observation will tell us that whenever the two intervals overlap, one of the interval's start time lies within the other interval. This rule can help us identify if any two intervals overlap or not.



Now, if we have found that the two intervals overlap, how can we find the overlapped part?

Again from the above diagram, the overlapping interval will be equal to:

```
start = max(a.start, b.start)
end = min(a.end, b.end)
```

That is, the highest start time and the lowest end time will be the overlapping interval.

So our algorithm will be to iterate through both the lists together to see if any two intervals overlap. If two intervals overlap, we will insert the overlapped part into a result list and move on to the next interval which is finishing early.

## Code #

Here is what our algorithm will look like:

```
Python3
                         G C++
👙 Java
                                      JS JS
    class Interval {
 2
       constructor(start, end) {
 3
         this.start = start;
         this.end = end;
 5
       }
 6
       print_interval() {
         process.stdout.write(`[${this.start}, ${this.end}]`);
 8
 9
       }
     }
10
11
```

```
function merge(intervals_a, intervals_b) {
12
13
                                            let result = [],
14
                                                          i = 0,
15
                                                          j = 0;
16
17
                                         while (i < intervals_a.length && j < intervals_b.length) {</pre>
                                                        // check if intervals overlap and intervals_a[i]'s start time lies with
18
19
                                                        a_overlaps_b = intervals_a[i].start >= intervals_b[j].start && intervals_b[j].
20
21
                                                        // check if intervals overlap and intervals_a[j]'s start time lies wit
22
                                                         b_overlaps_a = intervals_b[j].start >= intervals_a[i].start && intervals_a[i].
23
24
                                                        // store the the intersection part
25
                                                         if (a_overlaps_b || b_overlaps_a) {
                                                                       result.push(new Interval(Math.max(intervals_a[i].start, intervals_b
26
27
                                                                                    Math.min(intervals a[i].end, intervals b[j].end)));
28
                                                        }
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                       Run
                                                                                                                                                                                                                                                                                                                                                                            Save
                                                                                                                                                                                                                                                                                                                                                                                                                                            Reset
```

## Time complexity #

As we are iterating through both the lists once, the time complexity of the above algorithm is O(N+M), where 'N' and 'M' are the total number of intervals in the input arrays respectively.

### Space complexity #

Ignoring the space needed for the result list, the algorithm runs in constant space O(1).



Next -

Insert Interval (medium)

Conflicting Appointments (medium)



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(https://discuss.educative.io/tag/intervals-intersection-medium\_\_pattern-mergeintervals\_grokking-the-coding-interview-patterns-for-coding-questions? open=true&ctag=grokking-the-coding-interview-patterns-for-codingquestions\_\_design-

gurus&aid=5668639101419520&cid=5671464854355968&pid=6518042546667520)