





Order-agnostic Binary Search (easy)

We'll cover the following

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Problem Statement

Given a sorted array of numbers, find if a given number 'key' is present in the array. Though we know that the array is sorted, we don't know if it's sorted in ascending or descending order. You should assume that the array can have duplicates.

Write a function to return the index of the 'key' if it is present in the array, otherwise return -1.

Example 1:

Input: [4, 6, 10], key = 10

Output: 2

Example 2:

Input: [1, 2, 3, 4, 5, 6, 7], key = 5

Output: 4

Example 3:

Input: [10, 6, 4], key = 10

Output: 0

Example 4:

Input: [10, 6, 4], key = 4

Output: 2

Try it yourself

Try solving this question here:



```
const binary_search = function(arr, key) {
   // TODO: Write your code here
   return -1;
};

console.log(binary_search([4, 6, 10], 10))
console.log(binary_search([1, 2, 3, 4, 5, 6, 7], 5))
console.log(binary_search([10, 6, 4], 10))
console.log(binary_search([10, 6, 4], 4))
Run

Run

Save
Reset []
```

Solution#

To make things simple, let's first solve this problem assuming that the input array is sorted in ascending order. Here are the set of steps for **Binary Search**:

1. Let's assume start is pointing to the first index and end is pointing to the last index of the input array (let's call it arr). This means:

```
int start = 0;
int end = arr.length - 1;
```

2. First, we will find the middle of start and end. An easy way to find the middle would be: middle = (start + end)/2. For **Java and C++**, this equation will work for most cases, but when start or end is large, this equation will give us the wrong result due to integer overflow. Imagine that end is equal to the maximum range of an integer (e.g. for Java: int end = Integer.MAX_VALUE). Now adding any positive number to end will result in an integer overflow. Since

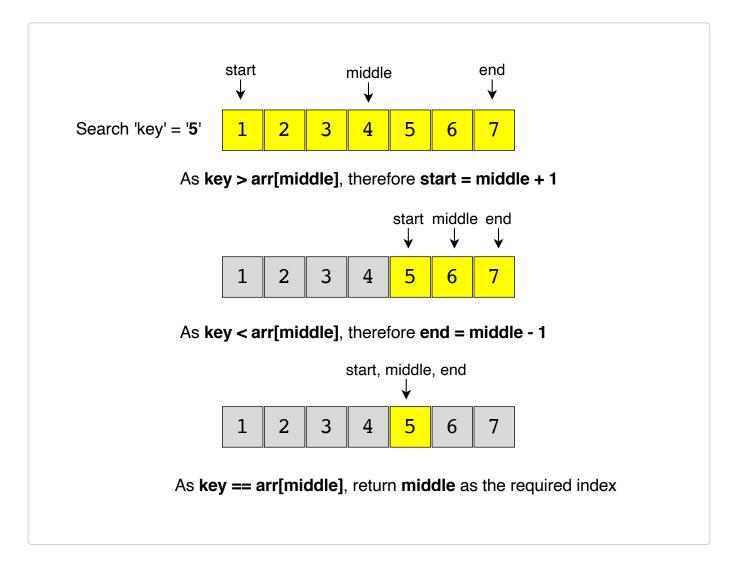
we need to add both the numbers first to evaluate our equation, an overflow might occur. The safest way to find the middle of two numbers without getting an overflow is as follows:

```
middle = start + (end-start)/2
```

The above discussion is not relevant for **Python**, as we don't have the integer overflow problem in pure Python.

- 3. Next, we will see if the 'key' is equal to the number at index middle. If it is equal we return middle as the required index.
- 4. If 'key' is not equal to number at index middle, we have to check two things:
 - o If key < arr[middle], then we can conclude that the key will be smaller than all the numbers after index middle as the array is sorted in the ascending order. Hence, we can reduce our search to end = mid − 1.
 - If key > arr[middle], then we can conclude that the key will be greater than all numbers before index middle as the array is sorted in the ascending order. Hence, we can reduce our search to start = mid + 1.
- 5. We will repeat steps 2-4 with new ranges of start to end. If at any time start becomes greater than end, this means that we can't find the 'key' in the input array and we must return '-1'.

Here is the visual representation of **Binary Search** for the Example-2:



If the array is sorted in the descending order, we have to update the step 4 above as:

- If key > arr[middle], then we can conclude that the key will be greater than all numbers after index middle as the array is sorted in the descending order. Hence, we can reduce our search to end = mid 1.
- If key < arr[middle], then we can conclude that the key will be smaller than all the numbers before index middle as the array is sorted in the descending order. Hence, we can reduce our search to start = mid + 1.

Finally, how can we figure out the sort order of the input array? We can compare the numbers pointed out by start and end index to find the sort order. If arr[start] < arr[end], it means that the numbers are sorted in ascending order otherwise they are sorted in the descending order.

Code

Here is what our algorithm will look like:



```
_{\perp}
function binary_search(arr, key) {
  let start = 0;
  end = arr.length - 1;
  isAscending = arr[start] < arr[end];</pre>
  while (start <= end) {</pre>
    // calculate the middle of the current range
    mid = Math.floor(start + (end - start) / 2);
    if (key === arr[mid]) {
      return mid;
    }
    if (isAscending) { // ascending order
      if (key < arr[mid]) {</pre>
        end = mid - 1; // the 'key' can be in the first half
      } else { // key > arr[mid]
        start = mid + 1; // the 'key' can be in the second half
    } else { // descending order
      if (key > arr[mid]) {
        end = mid - 1; // the 'key' can be in the first half
      } else { // key < arr[mid]</pre>
        start = mid + 1; // the 'key' can be in the second half
   }
  }
  return -1; // element not found
}
console.log(binary_search([4, 6, 10], 10));
console.log(binary_search([1, 2, 3, 4, 5, 6, 7], 5));
console.log(binary_search([10, 6, 4], 10));
console.log(binary_search([10, 6, 4], 4));
```

Run Save Reset []

Time complexity

Since, we are reducing the search range by half at every step, this means that the time complexity of our algorithm will be $O(\log N)$ where 'N' is the total elements in the given array.

Space complexity

(i)

The algorithm runs in constant space O(1).

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