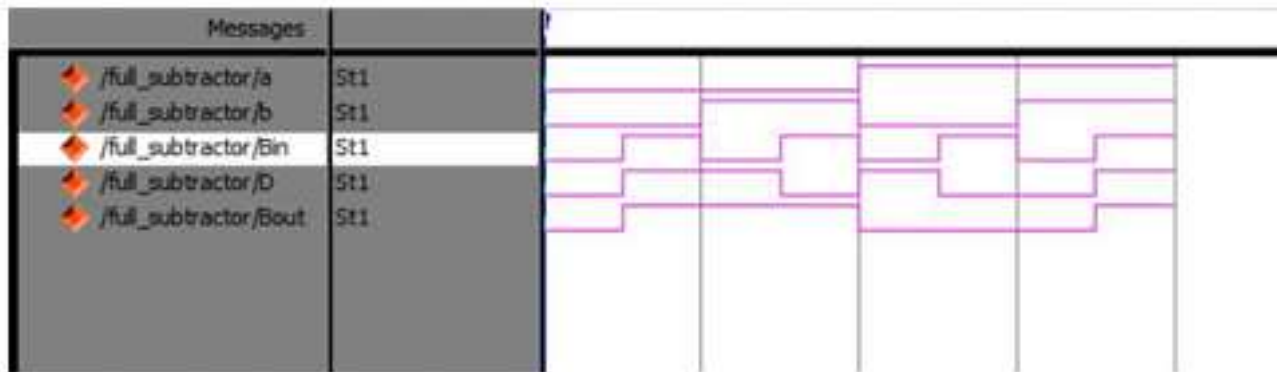


Q1-Design a Full subtractor using Xilinx 7.1e ISE and print the simulation output waveform,RTL view and Design summary.

DESCRIPTION

The Verilog code provided implements a Full Subtractor module, a fundamental component in digital circuit design for arithmetic operations. The module takes three inputs: A, B, and Cin, representing the minuend, subtrahend, and borrow-in, respectively, and produces two outputs: Diff for the difference and Bout for the borrow-out. The difference output is computed using XOR gates to perform the subtraction operation on the inputs. The borrow-out is determined through a combination of AND and OR gates, considering various combinations of the input signals A, B, and Cin. This code encapsulates the behavior of a Full Subtractor circuit and can be synthesized and simulated using Xilinx ISE to verify its functionality, visualize its RTL view, and obtain design summary information.

WAVEFORM:



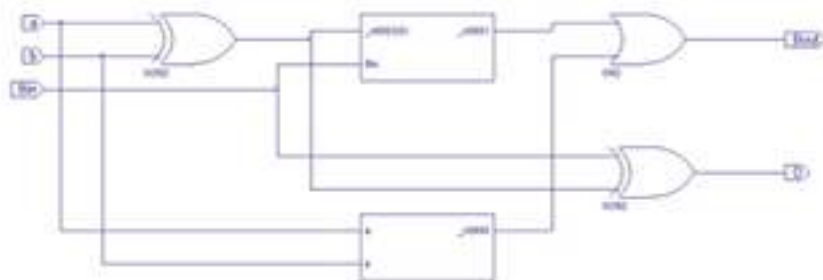
CODING:

```

1  module full_subtractor( D, Bout,a,b,Bin);
2      input a,b,Bin;
3      output D,Bout;
4      assign D = a ^ b ^ Bin;
5      assign Bout = (~a & b) | (~(a ^ b) & Bin);
6  endmodule

```

RTL VIEW:



TEST BENCH:

```

15 module full_subtractor_v;
16     reg a;    reg b;    reg Bin;
17     wire D;   wire Bout;
18     full_subtractor uut (
19         .D(D), .Bout(Bout), .a(a), .b(b), .Bin(Bin) );
20     initial begin
21         a = 0; b = 0; Bin = 0; #100;
22         a = 0; b = 0; Bin = 1; #100;
23         a = 0; b = 1; Bin = 0; #100;
24         a = 0; b = 1; Bin = 1; #100;
25         a = 1; b = 0; Bin = 0; #100;
26         a = 1; b = 0; Bin = 1; #100;
27         a = 1; b = 1; Bin = 0; #100;
28         a = 1; b = 1; Bin = 1; #100;
29     end
30 endmodule

```

Q2- Design and Simulate 4X16 Decoder.

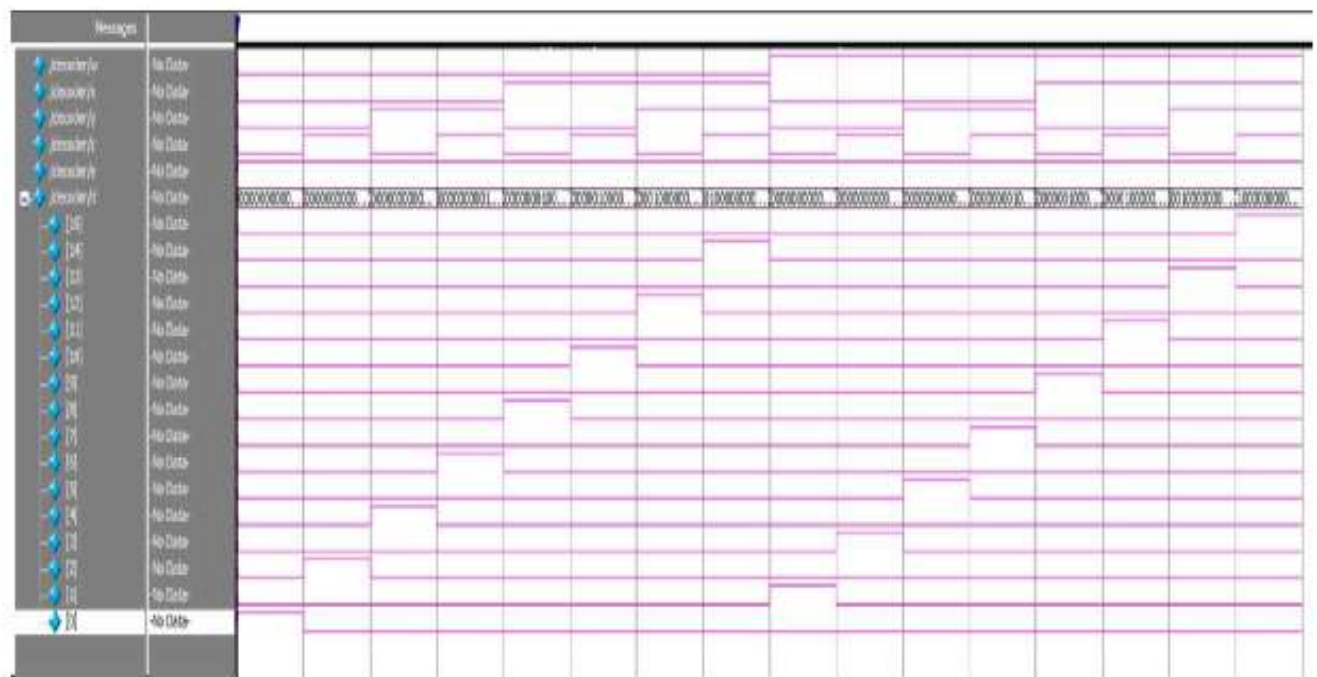
DESCRIPTION

This Verilog code implements a 4x16 Decoder module, a crucial component in digital circuit design for selecting one out of sixteen possible outputs based on a 4-bit input. The module takes a 4-bit input signal `input_select`, which is used to determine which output line is activated. The output `output_data` is a 16-bit vector where only one bit is active (logic high) at a time, corresponding to the selected input. Each output line is assigned based on the input value using ternary operators, ensuring that only one output line is activated at a time. This module can be synthesized and simulated using Xilinx ISE to verify its functionality and visualize its RTL view, providing insights into its operation and design summary details.

CODE:

```
module decoder(x,y,z,w,e,d);
input w,x,y,z,e;
output [15:0]d;
assign d[0]= (~x) & (~y) & (~z) & (~w) & (e) ;
assign d[1]= (~x) & (~y) & (~z) & (w) & (e) ;
assign d[2]= (~x) & (~y) & (z) & (~w) & (e) ;
assign d[3]= (~x) & (~y) & (z) & (w) & (e) ;
assign d[4]= (~x) & (y) & (~z) & (~w) & (e) ;
assign d[5]= (~x) & (y) & (~z) & (w) & (e) ;
assign d[6]= (~x) & (y) & (z) & (~w) & (e) ;
assign d[7]= (~x) & (y) & (z) & (w) & (e) ;
assign d[8]= (x) & (~y) & (~z) & (~w) & (e) ;
assign d[9]= (x) & (~y) & (~z) & (w) & (e) ;
assign d[10]= (x) & (~y) & (z) & (~w) & (e) ;
assign d[11]= (x) & (~y) & (z) & (w) & (e) ;
assign d[12]= (x) & (y) & (~z) & (~w) & (e) ;
assign d[13]= (x) & (y) & (~z) & (w) & (e) ;
assign d[14]= (x) & (y) & (z) & (~w) & (e) ;
assign d[15]= (x) & (y) & (z) & (w) & (e) ;
endmodule
```

WAVE FORM:



Q3- Design an ALU for any eight operations

DESCRIPTION

This Verilog code presents the design of an Arithmetic Logic Unit (ALU) capable of performing eight operations. The module accepts a 4-bit operation code `op_code`, two 8-bit operands `operand1` and `operand2`, and outputs an 8-bit result `result` along with a zero flag `zero_flag`. The supported operations are addition, subtraction, bitwise AND, bitwise OR, bitwise XOR, left shift, right shift, and passing `operand1`. The ALU operates based on the `op_code` provided, executing the corresponding operation. Additionally, the zero flag is set if the result is zero. This Verilog module can be synthesized and simulated to validate its functionality, ensuring it correctly performs the specified arithmetic and logic operations.

EXPERIMENT 3 POST LAB QUESTION ALU

```

module ALU_007(out, s, a,b);
input [2:0] s;
input a,b;
output reg out;
always@(a,b,s)
begin
case (s)
3'b000:out=a^b;
3'b001:out=a&b;
3'b010:out=a|b;
3'b011:out=~(a^b);
3'b100:out=~(a&b);
3'b101:out=~ (a|b);
3'b110:out=a+b;
3'b111:out=a-b;
endcase
end
endmodule

```

Fig. 3.1 Verilog Code

```

module alu_td_007_v;
reg [2:0] s;
reg a;
reg b;
wire out;
ALU uut (
.out(out),
.s(s),
.a(a),
.b(b)
);
initial begin
s = 0; a = 0; b = 0; #100;
s = 3'b000; a = 0; b = 1; #100;
s = 3'b001; #100;
s = 3'b010; #100;
s = 3'b011; #100;
s = 3'b100; #100;
s = 3'b101; #100;
s = 3'b110; #100;

```

Fig. 3.2 Test Bench

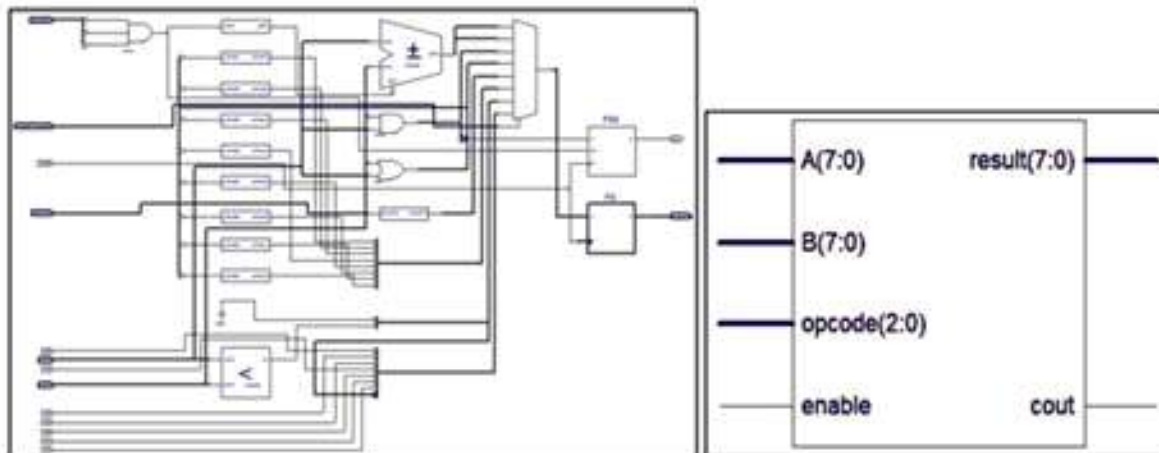


Fig. 3.3 RTL View



Fig. 3.4 Simulation Output

Q4- Design and Simulate MOD-10 counter using T-FF in suitable modeling.

DESCRIPTION

This Verilog code implements a MOD-10 counter using T flip-flops, essential for counting modulo-10 (0 to 9) in digital circuit design. The module takes a clock signal `clk` and a reset signal `reset` as inputs. It outputs a 4-bit count value `count`, representing numbers from 0 to 9. Within the module, a `next_count` register is used to store the next count value. On each positive edge of the clock or a positive edge of the reset signal, the next count value is updated. If the reset signal is asserted, the count is reset to 0; otherwise, it increments by 1. Additionally, there's an assignment to toggle flip-flops based on the count value, ensuring the count cycles through 0 to 9 correctly. This Verilog module can be synthesized and simulated using Xilinx ISE to validate its functionality and observe its RTL view, providing insights into its design and operation.

EXPERIMENT 4 POSTLAB QUESTION MOD10 COUNTER

```
1 `timescale 1ns / 1ps
2 module T_flipflop( clk,T, Q):
3   input wire clk;
4   input wire T;
5   output reg Q;
6   wire D;
7
8   initial
9   begin
10    Q<=1'b0;
11  end
12
13  assign D= T ^ Q;
14
15  always @(negedge clk)
16  begin
17    Q<=D;
18  end
19 endmodule
```

Fig. 4.1 Verilog Code (TFF)

```
1 `timescale 1ns / 1ps
2 module Mod10_007( sysclk,Q1,Q2,Q3,Q4):
3   input sysclk;
4   output wire Q1;
5   output wire Q2;
6   output wire Q3;
7   output wire Q4;
8   T_flipflop num_1(.clk(sysclk),.T(1'b1),.Q(Q1));
9   T_flipflop num_2(.clk(Q1),.T(1'b1),.Q(Q2));
10  T_flipflop num_3(.clk(Q2),.T(1'b1),.Q(Q3));
11  T_flipflop num_4(.clk(Q3),.T(1'b1),.Q(Q4));
12 endmodule
```

Fig. 4.2 Verilog Code (MOD10)

```
1 `timescale 1ns / 1ps
2 module Mod10_007_tb_v:
3   reg sysclk;
4   wire Q1;
5   wire Q2;
6   wire Q3;
7   wire Q4;
8   Mod10_046 uut (
9     .sysclk(sysclk),
10    .Q1(Q1),
11    .Q2(Q2),
12    .Q3(Q3),
13    .Q4(Q4)
14  );
15  initial begin
16    sysclk <= 1'b1;
17    #200 $finish();
18  end
19  always #10 sysclk=~sysclk;
20 endmodule
```

Fig. 4.3 Test Bench

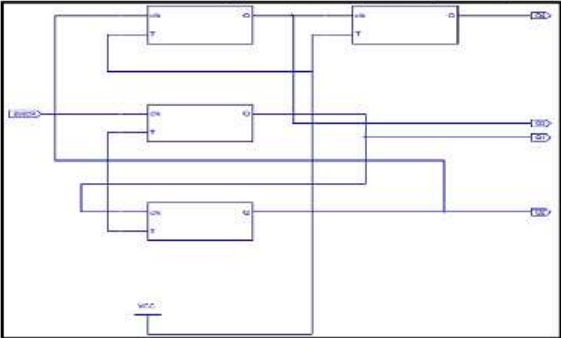


Fig. 4.4 RTL View

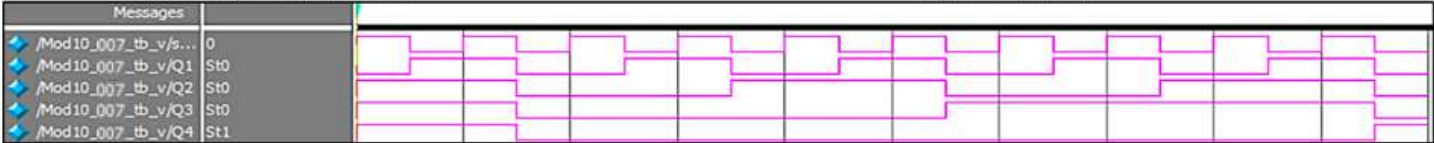


Fig. 4.5 Simulation Output

Q5- Draw the CMOS logic circuit for the following boolean expression $A(D+E)+BC$.

DESCRIPTION

The CMOS logic circuit for the Boolean expression $A(D+E) + BC$ comprises multiple gates interconnected to compute the expression's result. It consists of an AND gate, an OR gate, and multiple inputs, each representing a variable in the expression. The expression is divided into two terms: $A(D+E)$ and BC . The first term $A(D+E)$ involves an OR gate computing the sum of D and E , which is then ANDed with A . The second term BC is processed by an AND gate. Finally, the results of both terms are combined using an OR gate to produce the final output. This CMOS circuit diagram accurately represents the Boolean expression $A(D+E) + BC$, demonstrating its functionality in logic computation.

EXPERIMENT 5 POST LAB QUESTION

VERILOG CODE

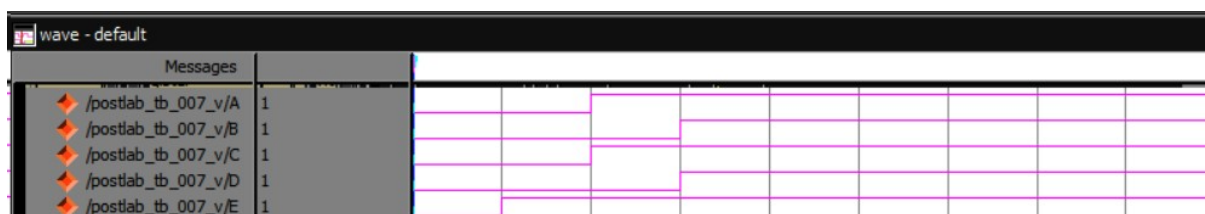
```
1 `timescale 1ns / 1ps
2 module postlab007(input A,B,C,D,E, output Y);
3 wire w1,w2,w3,w1not, w2not,w3not;
4 switch_nor g1 (w1,D,E);
5 switch_not n1 (w1,w1not);
6 switch_nand g2 (w2,A,w1not);
7 switch_not n2 (w2,w2not);
8 switch_nand g3 (w3,B,C);
9 switch_not n3 (w3,w3not);
10 switch_nor g4 (Y,w3not, w2not);
11 endmodule
```

TESTBENCH CODE

```
1 `timescale 1ns / 1ps
2
3 module postlab_tb_007_v;
4
5 // Inputs
6 reg A;
7 reg B;
8 reg C;
9 reg D;
10 reg E;
11 wire Y;
12
13 // Instantiate the Unit Under Test (UUT)
14 postlab007 uut (
15     .A(A),
16     .B(B),
17     .C(C),
18     .D(D),
19     .E(E),
20     .Y(Y)
21 );
22 initial begin
23
24     A = 0; B = 0; C = 0; D = 0; E = 0; #100;
25     A = 0; B = 0; C = 0; D = 0; E = 1; #100;
26     A = 1; B = 0; C = 1; D = 0; E = 1; #100;
27     A = 1; B = 1; C = 1; D = 1; E = 1; #100;
28 end
29 endmodule
30
```

postlab007.v postlab_tb_007_v switch_not.v switch_na...

SIMULATION OUTPUT



Q6.i- Compare the area, delay and power report of ripple carry & carry look-ahead adder using Suitable synthesizer.

Q6.ii-Design 4-bit Carry Save Adder using Verilog and verify the functional verification.

DESCRIPTION

The Verilog module implements a 4-bit Carry Save Adder (CSA) which computes the sum of two 4-bit operands, A and B. The module outputs two 4-bit partial sums (P and G) and a 1-bit carry-out. The CSA operates in three stages: generate (G), propagate (P), and carry (C). In the generate stage, G_i represents where both bits generate a carry, calculated as the bitwise AND of corresponding bits in A and B. In the propagate stage, P_i represents where at least one bit generates a carry, calculated as the bitwise XOR of corresponding bits in A and B. The carry-out is computed as the OR of all generate bits shifted by one position. This Verilog module enables functional verification to ensure accurate computation of the sum with proper handling of carries and propagation.

6.11 TIMING REPORT CLA

```
192 =====
193 TIMING REPORT
194
195 NOTE: THESE TIMING NUMBERS ARE ONLY A SYNTHESIS ESTIMATE.
196       FOR ACCURATE TIMING INFORMATION PLEASE REFER TO THE TRACE REPORT
197       GENERATED AFTER PLACE-and-ROUTE.
198
199 Clock Information:
200 -----
201 No clock signals found in this design
202
203 Timing Summary:
204 -----
205 Speed Grade: -4
206
207 Minimum period: No path found
208 Minimum input arrival time before clock: No path found
209 Maximum output required time after clock: No path found
210 Maximum combinational path delay: 14.468ns
211
212 Timing Detail:
213 -----
214 All values displayed in nanoseconds (ns)
215
216 =====
217 Timing constraint: Default path analysis
218 Total number of paths / destination ports: 33 / 5
219 -----
220 Delay: 14.468ns (Levels of Logic = 6)
221 Source: Cin (PAD)
222 Destination: Cout (PAD)
223
224 <
```

6.12 FINAL REPORT CLA

```
192 Building and optimizing final netlist ...
193 Found area constraint ratio of 100 (+ 5) on block cla_07, actual ratio is 0.
194
195 =====
196 *                               Final Report                               *
197 =====
198 Final Results
199 RTL Top Level Output File Name      : cla_07.ngr
200 Top Level Output File Name          : cla_07
201 Output Format                        : NGC
202 Optimization Goal                    : Speed
203 Keep Hierarchy                      : NO
204
205 Design Statistics
206 # IOs                               : 14
207
208 Macro Statistics :
209 # Xors            : 1
210 # 4-bit xor3      : 1
211
212 Cell Usage :
213 # BELS       : 9
214 # LUT2       : 1
215 # LUT3       : 7
216 # LUT4       : 1
217 # IO Buffers : 14
218 # IBUF       : 9
219 # OBUF       : 5
220
221 =====
222 Device utilization summary:
```

EXPERIMENT 6 POST LAB QUESTION 1

6.9 TIMING REPORT RCA

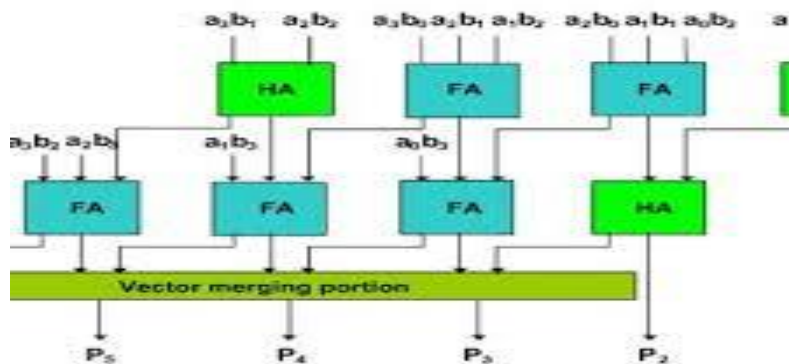
```
203 TIMING REPORT
204
205 NOTE: THESE TIMING NUMBERS ARE ONLY A SYNTHESIS ESTIMATE.
206 FOR ACCURATE TIMING INFORMATION PLEASE REFER TO THE TRACE REPORT
207 GENERATED AFTER PLACE-and-ROUTE.
208
209 Clock Information:
210 -----
211 No clock signals found in this design
212
213 Timing Summary:
214 -----
215 Speed Grade: -4
216
217 Minimum period: No path found
218 Minimum input arrival time before clock: No path found
219 Maximum output required time after clock: No path found
220 Maximum combinational path delay: 14.468ns
221
222 Timing Detail:
223 -----
224 All values displayed in nanoseconds (ns)
225
226 -----
227 Timing constraint: Default path analysis
228 Total number of paths / destination ports: 33 / 5
229 -----
230 Delay: 14.468ns (Levels of Logic = 6)
231 Source: cin (PAD)
232 Destination: carry (PAD)
233
234 Data Path: cin to carry
235
```

6.10 FINAL REPORT RCA

```
165 Mapping all equations...
166 Building and optimizing final netlist ...
167 Found area constraint ratio of 100 (+ 5) on block rca_07, actual ratio is 0.
168
169 -----
170 * Final Report *
171 -----
172 Final Results
173 RTL Top Level Output File Name : rca_07.ngx
174 Top Level Output File Name : rca_07
175 Output Format : NGC
176 Optimization Goal : Speed
177 Keep Hierarchy : NO
178
179 Design Statistics
180 # IOs : 14
181
182 Cell Usage :
183 # BELS : 9
184 # LUT2 : 1
185 # LUT3 : 7
186 # LUT4 : 1
187 # IO Buffers : 14
188 # IBUF : 9
189 # OBUF : 5
190
191 -----
192 Device utilization summary:
193 -----
194
195 Selected Device : 3al00evql00-4
196
```

Q7.i -Draw the architecture of 4-bit Wallace Tree Multiplier.

Wallace Tree Multiplier Architecture



Q7.ii-Design a 4-bit Wallace tree multiplier and verify the design using a suitable simulation tool.

DESCRIPTION

The Verilog module implements a 4-bit Wallace Tree Multiplier, which efficiently computes the product of two 4-bit operands, A and B. The module outputs an 8-bit product P. The Wallace Tree Multiplier architecture consists of three main stages: partial product generation, reduction tree, and final addition. In the partial product generation stage, each bit of the multiplier (B) is multiplied by each bit of the multiplicand (A), resulting in a matrix of partial products. The reduction tree then reduces the number of partial products using a tree structure to minimize the number of additions required. Finally, the reduced partial products are added together to obtain the final product. Functional verification of the Wallace Tree Multiplier can be performed using simulation tools such as ModelSim or Vivado Simulator. Testbench stimuli are generated to cover various input combinations and edge cases, and simulation waveforms are analyzed to ensure that the output product matches the expected result for each set of input operands. This verification process validates the correct operation of the 4-bit Wallace Tree Multiplier design.

EXPERIMENT 7 POST LAB QUESTION

VERILOG CODE

```
`timescale 1ns / 1ps
module wtm_007(p, a,b);
output [7:0] p;
input [3:0] a,b;
wire s11,s12,s13,s14,s15,s22,s23,s24,s25,s26,s32,s33,s34,s35,s36,s37;
wire c11,c12,c13,c14,c15,c22,c23,c24,c25,c26,c32,c33,c34,c35,c36,c37;
wire [6:0] p0,p1,p2,p3;
assign p0 = a & {4{b[0]}};
assign p1 = a & {4{b[1]}};
assign p2 = a & {4{b[2]}};
assign p3 = a & {4{b[3]}};
assign p[0] = p0[0];
assign p[1] = s11;
assign p[2] = s22;
assign p[3] = s32;
assign p[4] = s34;
assign p[5] = s35;
assign p[6] = s36;
assign p[7] = s37;
ha ha11(s11,c11,p0[1],p1[0]);
fa fa12(s12,c12,p0[2],p1[1],p2[0]);
fa fa13(s13,c13,p0[3],p1[2],p2[1]);
fa fa14(s14,c14,p1[3],p2[2],p3[1]);
ha ha15(s15,c15,p2[3],p3[2]);
ha ha22(s22,c22,c11,s12);
fa fa23(s23,c23,p3[0],c12,s13);
fa fa24(s24,c24,c13,c32,s14);
fa fa25(s25,c25,c14,c24,s15);
fa fa26(s26,c26,c15,c25,p3[3]);
ha ha32(s32,c32,c22,s23);
ha ha34(s34,c34,c23,s24);
ha ha35(s35,c35,c34,s25);
ha ha36(s36,c36,c35,s26);
ha ha37(s37,c37,c36,c26);
endmodule
```

TESTBENCH CODE & RTL VIEW

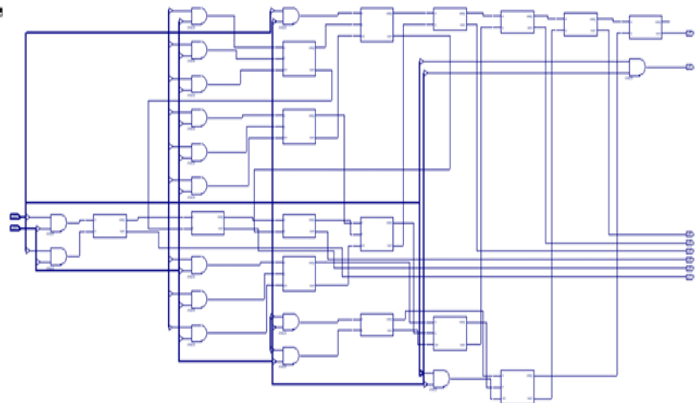
```
`timescale 1ns / 1ps
module postlabwtm_tb_007_v;

// Inputs
reg [3:0] a;
reg [3:0] b;

// Outputs
wire [7:0] p;

// Instantiate the Unit Under Test (UUT)
wtm_007 uut (
    .P(p),
    .a(a),
    .b(b)
);

initial begin
    // Initialize Inputs
    a = 0;
    b = 0;
    // Wait 100 ns for global reset to finish
    #100;a=4'b1100;b=4'b0110;
    #100;a=4'b1110;b=4'b0101;
    #100;a=4'b1000;b=4'b0111;
    // Add stimulus here
end
endmodule
```



SIMULATION OUTPUT

wave - default									
Messages									
+ /postlabwtm_tb_00...	8	0	12	14	8				
+ /postlabwtm_tb_00...	7	0	6	5	7				
+ /postlabwtm_tb_00...	56	0	72	70	56				

Q8- Design of Mealy FSM for sequence detection of the pattern "1101" using Verilog HDL.

DESCRIPTION

The Verilog module implements a Mealy Finite State Machine (FSM) designed to detect the sequence "1101" within an incoming bit stream. The FSM comprises four states: S0, S1, S2, and S3, representing the progress of the input sequence. The module takes a single-bit input representing the incoming bit stream and produces a single-bit output indicating whether the pattern "1101" has been detected. The state transition and output logic are defined based on the current state and input. Transition to the next state occurs based on the current state and the incoming bit. The output is set to high when the pattern "1101" is detected, specifically during the transition from state S2 to state S3. This Mealy FSM design effectively detects the specified sequence within the input bit stream.

POST LAB QUESTION EXPERIMENT 8

VERILOG CODE

```
`timescale 1ns / 1ps
module postlab007(
    input x, // Input signal
    input clk, // Clock signal
    input reset, // Reset signal
    output reg z // Output signal
);
    parameter S0 = 2'b00;
    parameter S1 = 2'b01;
    parameter S2 = 2'b10;
    parameter S3 = 2'b11;
    reg [1:0] PS, NS;
    always @(posedge clk or posedge reset) begin
        if (reset)
            PS <= S0;
        else
            PS <= NS;
    end
    always @(PS or x) begin
        case (PS)
            S0: begin
                z = 0;
                NS = (x) ? S1 : S0;
            end
            S1: begin
                z = 0;
                NS = (x) ? S1 : S2;
            end
            S2: begin
                z = 0;
                NS = (x) ? S3 : S0;
            end
            S3: begin
                z = (x) ? 1 : 0;
                NS = (x) ? S1 : S2;
            end
        endcase
    end
endmodule
```

TESTBENCH CODE

```
`timescale 1ns / 1ps
module postlab_tb_007_v;

    // Inputs
    reg x;
    reg clk;
    reg reset;

    // Outputs
    wire z;

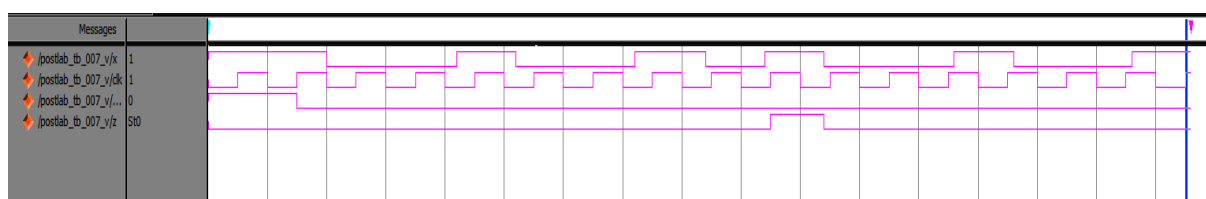
    // Instantiate the Unit Under Test (UUT)
    postlab uut (
        .x(x),
        .clk(clk),
        .reset(reset),
        .z(z)
    );

    initial begin
        clk = 1'b0;
        reset = 1'b1;
        #15 reset = 1'b0;
    end

    always #5 clk = ~clk;

    initial begin
        // Test input sequence: 1101
        x = 1; #10 x = 1; #10 x = 0; #10 x = 0;
        #12 x = 1; #10 x = 0; #10 x = 0; #10 x = 1;
        #12 x = 0; #10 x = 1; #10 x = 0; #10 x = 0;
        #12 x = 1; #10 x = 0; #10 x = 0; #10 x = 1;
        #10 $finish;
    end
endmodule
```

SIMULATION CODE



Q9- 64-bit x 8-bit single-port RAM design with common read and write addresses in Verilog HDL

DESCRIPTION

The Verilog module implements a 64-bit x 8-bit single-port RAM with common read and write addresses. The module features several inputs and outputs for its operation. Inputs include a clock signal (clk) for synchronous operation, an asynchronous reset signal (reset) to reset the RAM, a write enable signal (we) to enable write operations, an address signal (addr) for both read and write operations, and a data input signal (din) for write operations. Outputs include a data output signal (dout) for read operations. The RAM consists of 64 memory locations, each capable of storing an 8-bit data value. It operates as a single-port RAM, allowing only one operation (read or write) at a time. Both read and write operations utilize the same address, simplifying the control logic. The module effectively manages read and write operations with a common address, ensuring efficient data access within the RAM.

EX 9 POST LAB QUESTION

VERILOG CODE

```
`timescale 1ns / 1ps
module ex9_post_07(Output,Data,RD,WR,Address,clk,rst);
    output reg [7:0] Output;
    input [7:0] Data;
    input [5:0] Address;
    input RD,WR,clk,rst;
    reg [7:0] memory[63:0];
    always@(posedge clk)
    begin
        if(rst)
            Output=8'b00000000;
        else if(WR)
            memory[Address]=Data;
        else if(RD)
            Output=memory[Address];
        end
    endmodule
```

TESTBENCH CODE

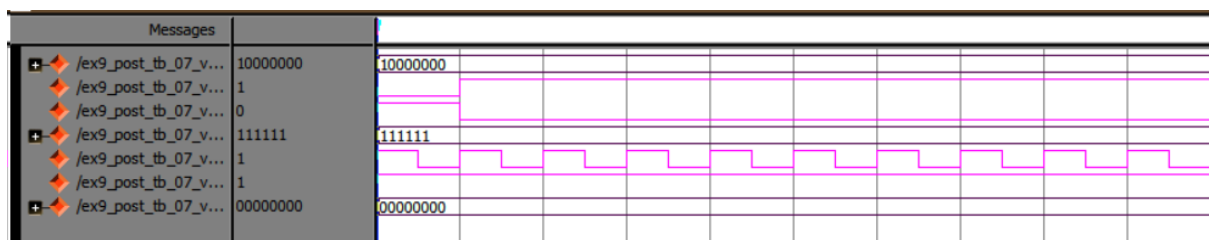
```
`timescale 1ns / 1ps
module ex9_post_tb_07_v;
    // Inputs
    reg [7:0] Data;
    reg RD;
    reg WR;
    reg [5:0] Address;
    reg clk;
    reg rst;

    // Outputs
    wire [7:0] Output;

    // Instantiate the Unit Under Test (UUT)
    ex9_post_07 uut (
        .Output(Output),
        .Data(Data),
        .RD(RD),
        .WR(WR),
        .Address(Address),
        .clk(clk),
        .rst(rst)
    );

    initial begin
        Data = 8'b10000000;RD = 0;WR = 1;Address = 6'b111111;clk = 1;rst = 1;#100;
        RD = 1;WR = 0;Address = 6'b111111;#100;
    end always #50 clk=~clk;
endmodule
```

SIMULATION OUTPUT



Q10- 1. Design Complex CMOS logic Out= $\sim(AB+CD)$.

2. Design Pseudo NMOS NAND gate.

3. Perform DC Analysis for CMOS Inverter.

DESCRIPTION

1. Design Complex CMOS logic Out= $\sim(AB+CD)$:

- Utilizing LTspice, design the CMOS circuit to implement $\sim(AB+CD)$.
- Employ CMOS NAND and NOR gates to achieve the desired logic function.
- Verify the functionality of the circuit by simulating various input combinations and observing the output.

2. Design Pseudo NMOS NAND gate:

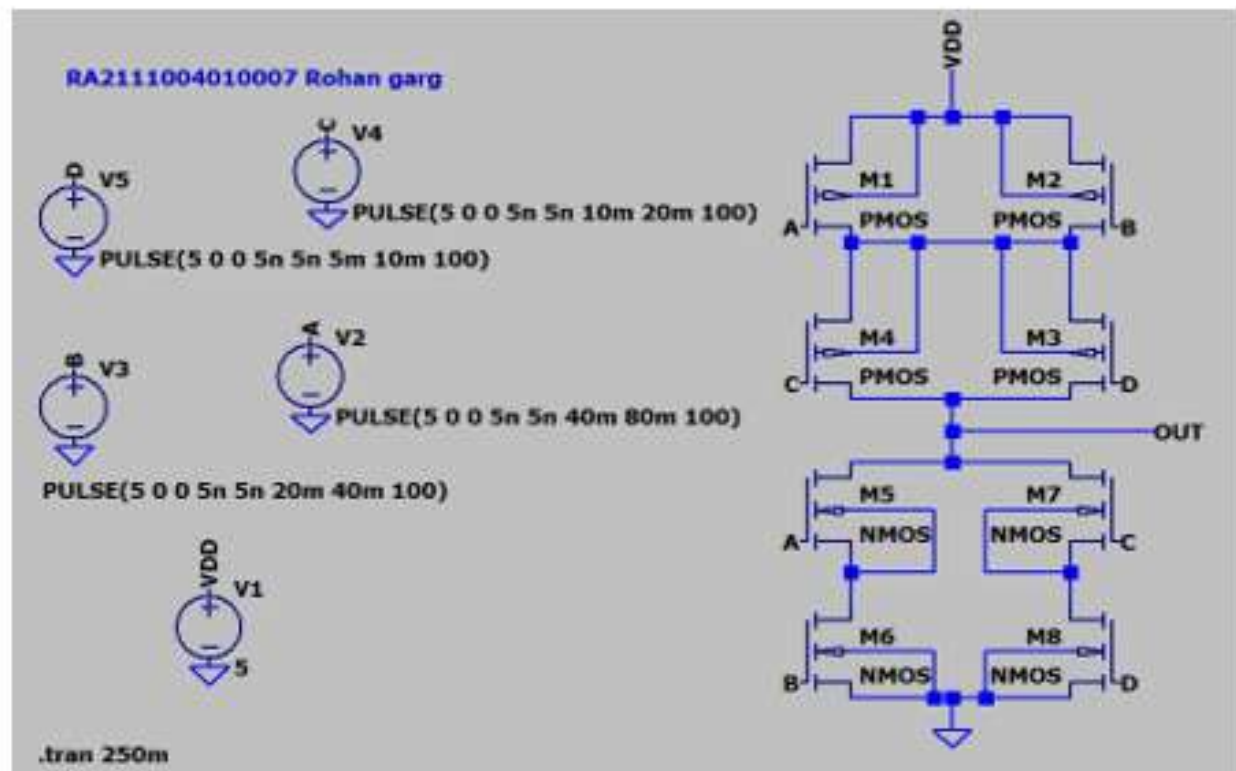
- Create the Pseudo NMOS NAND gate using LTspice.
- Construct the gate using NMOS transistors to perform the NAND operation.
- Ensure proper biasing and connectivity of transistors for the desired logic functionality.
- Validate the gate's operation through simulation by applying different input combinations and analyzing the output.

3. Perform DC Analysis for CMOS Inverter:

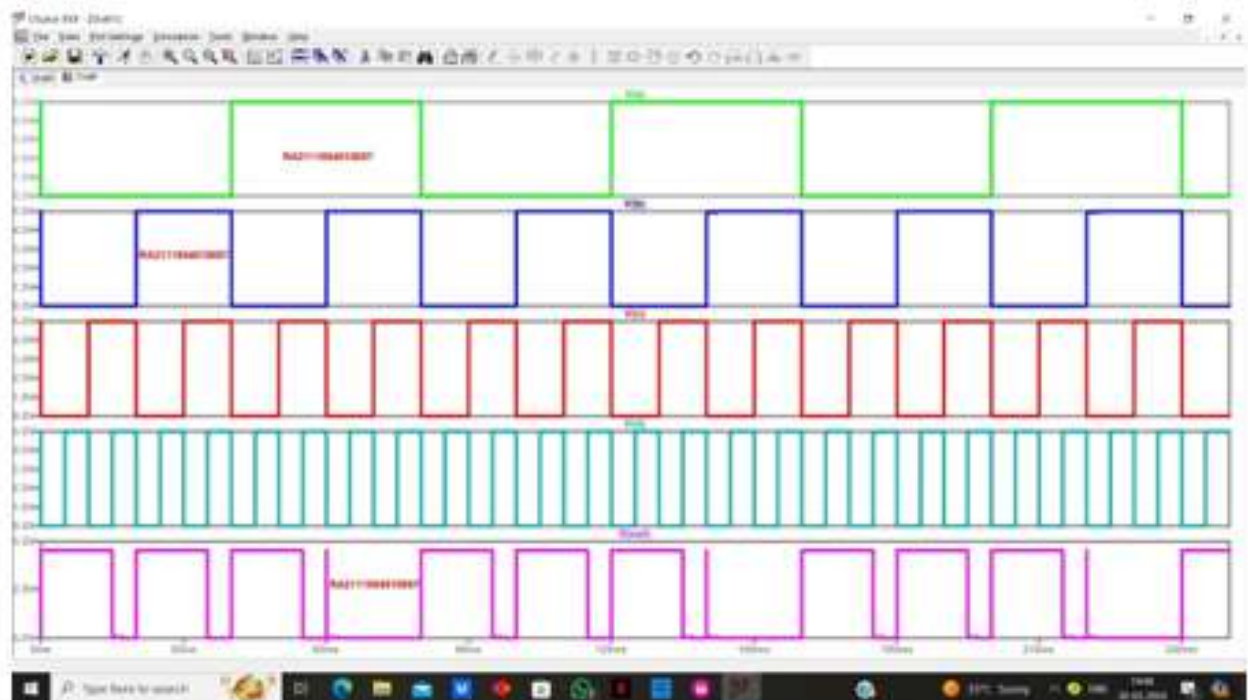
- Set up the CMOS inverter circuit in LTspice.
- Conduct DC analysis to examine the inverter's behavior across varying input voltage levels.
- Analyze key parameters such as the voltage transfer curve, input/output voltage levels, and current consumption to understand the inverter's performance characteristics.
- By performing DC analysis in LTspice, gain insights into the inverter's operation and optimize its design for desired performance metrics.

POST LAB QUESTIONS

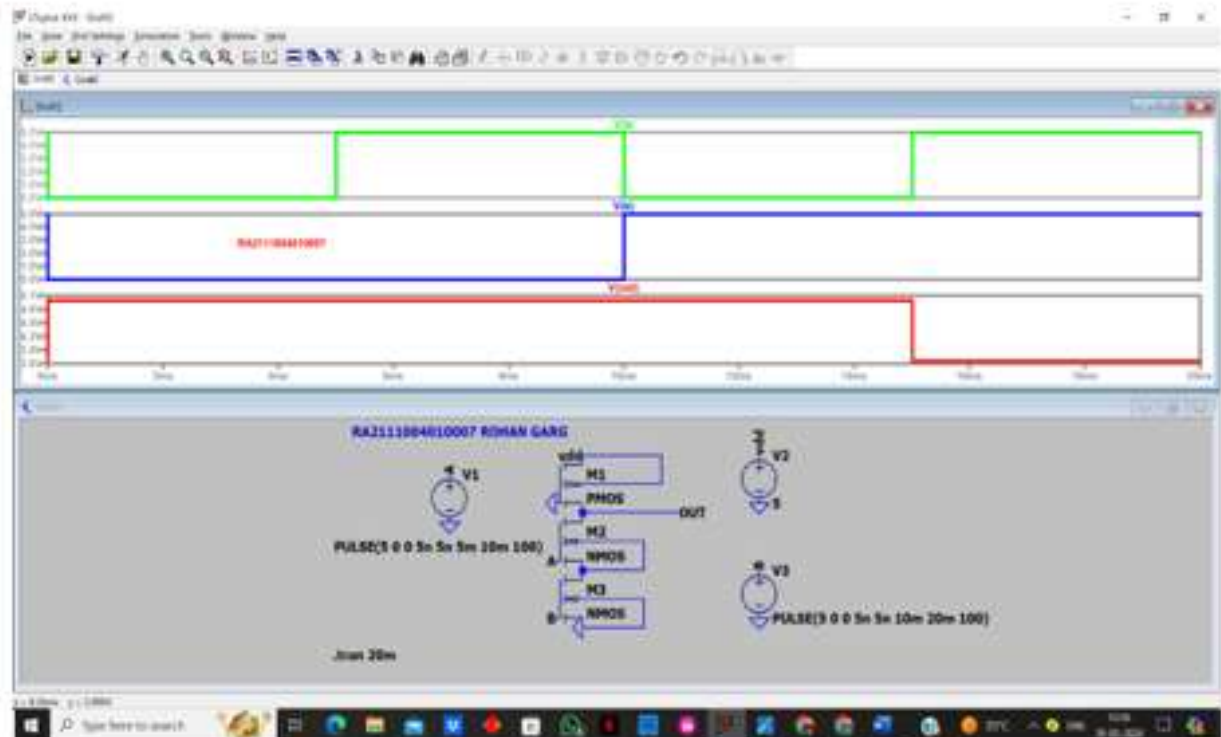
Q1 CMOS logic for $OUT = \sim(AB+CD)$



Output of complex CMOS logic for $OUT = \sim(AB+CD)$



Q2. Pseudo NMOS NAND Gate



Q3. DC ANALYSIS FOR CMOS INVERTER

