#### **CMPT 280**

Tutorial: Dijkstra's Algorithm

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```
Algoirthm diikstra(G. s)
Solves the single-source shortest paths problem.
G is a weighted graph with non-negative weights.
s is the start vertex.
Let V be the set of vertices in G.
For each v in V
    v.tentativeDistance = infinity
    v. visited = false
    v.predecessorNode = null
s.tentativeDistance = 0
while there is an unvisited vertex
    cur = the unvisited vertex with the smallest tentative distance.
    cur visited = true
    // update tentative distances for adjacent vertices if needed
    // note that w(i,j) is the cost of the edge from i to j.
    For each z adiacent to cur
        if (z is unvisited and z.tentativeDistance >
                               cur.tentativeDistance + w(cur.z) )
            z.tentativeDistance = cur.tentativeDistance + w(cur.z)
            z.predecessorNode = cur
```

5 6

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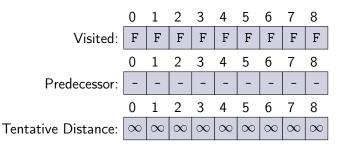
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## Implementing Dijkstra's Algorithm

- As we saw in class, while the algorithm as presented suggests that tentativeDistance, visited and predecessorNode are properties of nodes, it is not necessary to implement a specialized node object to contain this data.
- Since nodes are always mapped onto integer identifiers, we can store these values in a series of parallel arrays.

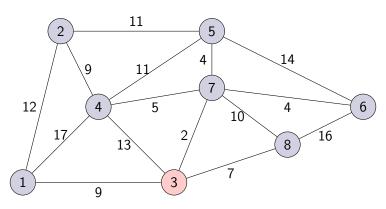


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## Tracing Dijkstra's Algorithm

Let's use Dijkstra's algorithm to find the shortest path from node 3 to every other node.



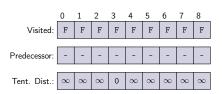
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## Dijkstra's Algorithm Initialization for Start Node 3

|                     | 0        | 1        | 2        | 3 | 4        | 5        | 6        | 7        | 8        |
|---------------------|----------|----------|----------|---|----------|----------|----------|----------|----------|
| Visited:            | F        | F        | F        | F | F        | F        | F        | F        | F        |
|                     | 0        | 1        | 2        | 3 | 4        | 5        | 6        | 7        | 8        |
| Predecessor:        | ı        | -        | -        | - | -        | -        | -        | -        | -        |
|                     | 0        | 1        | 2        | 3 | 4        | 5        | 6        | 7        | 8        |
| Tentative Distance: | $\infty$ | $\infty$ | $\infty$ | 0 | $\infty$ | $\infty$ | $\infty$ | $\infty$ | $\infty$ |

Note: Offset 0 of these arrays are not used as we use node IDs as array offsets. Nodes are numbered 1 to n as in 1ib280.



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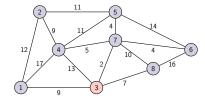
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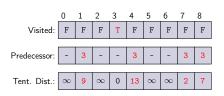
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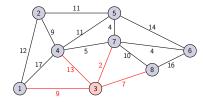
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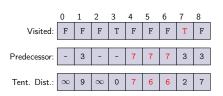


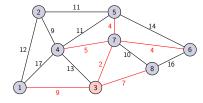
#### Unvisited node w/ Smallest tent. dist.: 3





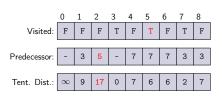
#### Processed Node 3, Next node to process: 7

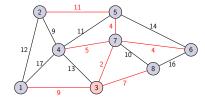




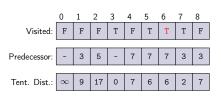
#### Processed Node 7, Next node to process: 5 or 6 (but we choose 5)

```
while there is an unvisited vertex
        cur = the unvisited vertex with the smallest tentative distance.
2
3
        cur.visited = true
5
        // update tentative distances for adjacent vertices if needed
        // note that w(i,j) is the cost of the edge from i to j.
6
7
        For each z adjacent to cur
            if (z is unvisited and z tentativeDistance >
8
                                    cur.tentativeDistance + w(cur,z) )
                z.tentativeDistance = cur.tentativeDistance + w(cur,z)
10
11
                z.predecessorNode = cur
```





#### Processed Node 5, Next node to process: 6



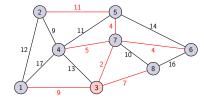
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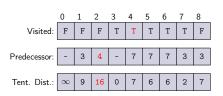
6 7

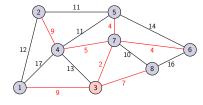
8

10 11



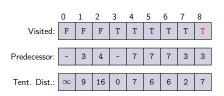
#### Processed Node 6, Next node to process: 4 or 8 (we choose 4)

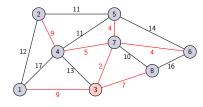




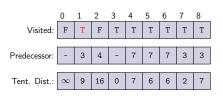
#### Processed Node 4, Next node to process: 8

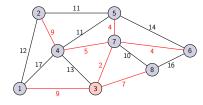
```
while there is an unvisited vertex
 2
        cur = the unvisited vertex with the smallest tentative distance.
3
        cur.visited = true
 5
        // update tentative distances for adjacent vertices if needed
        // note that w(i,j) is the cost of the edge from i to j.
6
7
        For each z adjacent to cur
            if (z is unvisited and z tentativeDistance >
8
                                    cur.tentativeDistance + w(cur,z) )
                z.tentativeDistance = cur.tentativeDistance + w(cur,z)
10
11
                z.predecessorNode = cur
```





#### Processed Node 8, Next node to process: 1

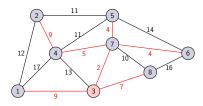




#### Processed Node 1, Next node to process: 2

```
while there is an unvisited vertex
 2
        cur = the unvisited vertex with the smallest tentative distance.
3
        cur.visited = true
 5
        // update tentative distances for adjacent vertices if needed
        // note that w(i,j) is the cost of the edge from i to j.
6
7
        For each z adjacent to cur
            if (z is unvisited and z tentativeDistance >
8
                                    cur.tentativeDistance + w(cur,z) )
                z.tentativeDistance = cur.tentativeDistance + w(cur,z)
10
11
                z.predecessorNode = cur
```

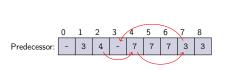


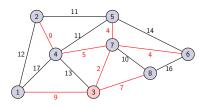


Processed Node 2, no more unvisited vertices, so we're done. Tentative distance array now contains the final lengths of the shortest paths from 3 to each node.

# Dijkstra's Algorithm Obtaining the Shortest Paths

To recover the shortest path from 3 to a node i, use the predecessor array starting at offset i, and follow the path back.





Recovering the shortest path from 3 to 2:

$$2 \leftarrow 4 \leftarrow 7 \leftarrow 3$$

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## Dijkstra's Algorithm More Examples

For more practice, trace through Dijkstra's algorithm again, but from a different start node, or, for an entirely different graph... maybe a directed one?