

CMPT 280

Topic 9: Tree Traversals

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References

- Textbook, Chapter 9

Depth First Traversals

Recall the general algorithm for depth-first traversals:

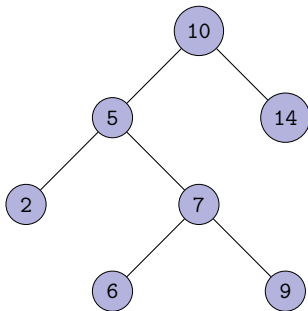
```
1 Algorithm depthFirstTraversal(N)
2 Parameters:
3     N is a tree node
4
5 visit node N
6 for each child T of N
7     depthFirstTraversal(T)
```

Exercise 1

- Write a method for `LinkedSimpleTree280<I>` that performs a depth-first traversal that prints out the contents of each node.
- Besides being a depth-first traversal, what kind of binary tree traversal is your solution?

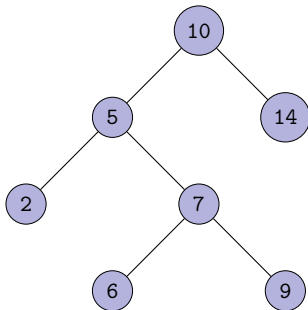
In-order Traversal of Binary Search Trees

- Recall the Binary Search Trees (from CMPT 145).
- What special property does the in-order traversal of a binary search tree have?



Exercise 2: In-order traversals

- Modify the method from Exercise 1 to create an in-order traversal.



Exercise 3: Post-order traversals

- Modify the method from Exercise 1 to create a post-order traversal.
- The sequence generated by a post-order traversal has no special significance for binary search trees.
- BUT... we can do some fairly interesting things with post-order traversal that don't involve printing the node contents.

Post-order traversals

We can accomplish the following with depth-first traversals:

- Count number of nodes in the tree.
- Determine height of the tree
- Evaluate expression tree.

Exercise 4: Counting number of nodes.

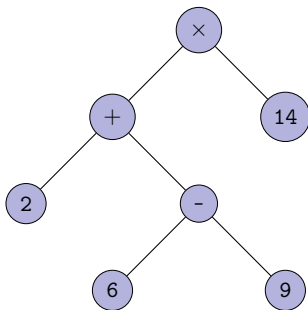
- Write a post-order traversal for `LinkedSimpleTree280<I>` that returns the number of nodes in the tree.
- Hint: If you know how many nodes are in the left and right subtrees, then the number of nodes in the tree is $1 + \# \text{ nodes in left subtree} + \# \text{ nodes in right subtree}$.
- What was the "visit" operation in the solution?

Exercise 5: Computing tree height.

- Modify the method from Exercise 4 to compute the height of the tree. It shouldn't take much since both are post-order traversals!
- Hint: If you know the height of the left and right subtrees, then the height of the tree is $1 + \max(\text{left subtree height}, \text{right subtree height})$
- What was the "visit" operation in the solution?

Postorder Traversal: Evaluate Expression Tree

- Why is evaluating an expression tree a post-order traversal?



Breadth-First Traversals

- Algorithm for BFT requires a queue.
- Idea:
 - After we visit a node, put all of it's children on a queue.
 - Keep visiting nodes as long as the queue is not empty.

```
1  Algorithm breathFirstTraversal(T)
2  T is a tree.
3
4  Let Q be a queue.
5  Q.add(T.rootNode)
6
7  while(Q is not empty)
8      p = Q.remove();
9      visit p
10     for each child s of p
11         Q.add(s)
```

Exercise 6

- Write a method for `LinkedSimpleTree280` that prints the contents of the nodes in level-order. The method should follow the BFT algorithm below with an appropriate "visit" operation.
- Hint: for a queue we could use any of `java.util.LinkedList`, `LinkedList280<I>`, or `LinkedQueue280<I>`.

```
1  Algorithm breathFirstTraversal(T)
2  T is a tree.
3
4  Let Q be a queue.
5  Q.add(T.rootNode)
6
7  while(Q is not empty)
8      p = Q.remove();
9      visit p
10     for each child s of p
11         Q.add(s)
```

Deep Clone: A depth-first algorithm.

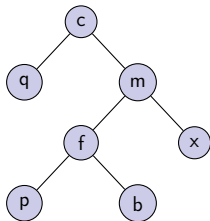
We can now see that deep cloning of an object is a depth-first traversal of the reference fields in an object. Note that the references in the object X and the objects those objects reference (etc.) form a tree (whether X is actually a tree ADT or not)!

```
1  Algorithm deepClone(X)
2  X is an object
3
4  Y = shallowClone(X)  // "Visit" operation is a shallow clone of
5                        // the current object.
6
7  for each non-null reference X.r in X
8      Y.r = deepClone(X.r);
9
10 return Y
```

Can we re-write this algorithm so that it is a bit more specific to cloning a `LinkedSimpleTree280<I>`?

Printing the Structure of Binary Trees

- Trees are hard to draw with root at the top.
- Trees are easy to draw with root at the left and descendants to the right:



1: c

2: q

3: x

2: m

4: b

3: f

4: p

- Right subtree of a node is above and indented; left subtree is below and indented.
- This can be achieved with a non-standard traversal:

right, root, left

Printing the Structure of Trees

(This is in `LinkedSimpleTree280`)

```
1  /** Form a string representation that includes level numbers when the tree
2  *   is at level i.
3  *   @param i   the level of this tree in the main tree, i.e., number of
4  *               indentations minus 1
5  *   @timing    Time = O(n), where n = number of items in this tree   */
6  protected String toStringByLevel(int i)
7  {
8      StringBuffer blanks = new StringBuffer((i - 1) * 5);
9      for (int j = 1; j < i; j++)
10         blanks.append("    ");
11     if (isEmpty())
12         return  "\n" + blanks + i + ": -";
13     else
14     {
15         boolean printSubtrees = !rootLeftSubtree().isEmpty() ||
16                                 !rootRightSubtree().isEmpty();
17
18         String result = "";
19         if (printSubtrees)
20             result += rootRightSubtree().toStringByLevel(i+1);
21
22         result += "\n" + blanks + i + ": " + rootItem();
23
24         if (printSubtrees)
25             result += rootLeftSubtree().toStringByLevel(i+1);
26
27         return result;
28     }
29 }
```


Next Class

- Next class reading: Chapter 10 — Dispensers