Correctness of selection Sort algorithm det us find the smallet element in the unsorted part of the array and swap it with the first element in the part of array for 1:0 to n-1 } = content think of the sole of the flore of the " เกลา อมิเ**วณ์หน้าพบาก สม**ัสธา กลี" เกอน (อ เกาแอ 10 Jor j= 1+1-10 no 500 of Camility and Infinity So minimum = j If (i=i) { Swap (ith element, minimum element)

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the transfer to

Correctness of the selection sort algorithm can be proven by using a loop invariant and demonstrating that it holds

true during the collision of algorithm.

Intralization:

At the start of Each Pteration of the outlet of loop the subarray before the outlet index "i' is sorted, and all element in this subarray are smaller than or in Equal to any element in the unsolved portion of array.

for "= 0 to n-1 of

minimum = i

1

Maintnance:

Before the execution of inner loop, all the elements of before i are sorted. The inner loop identifies the minimum element in the array and swaps the values accordingly

for (j=i+1 ton) {

if (ar [j] (ar [minimum]) {
minimum = j

2

Swap [arr [i], arr [minimom]]

The algorithm sorts the element at Endea P which belongs at the index minimum. The outer loop runs from 0 to n-1, iterating through array.

Termination.

Once the outer loop Completes the algorithm terminates, the entire array is sorted, and here are no intinite loops of conditions that would present the Completion.