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CST

15

① for $j=1$ $i=1$
 $j=2$ $i=1+2=3$
 $j=3$ $i=1+2+3=6$
 $j=4$ $i=1+2+3+4=10$

$$1+2+3+4+\dots+m \text{ times} < n$$

$$\frac{m(m+1)}{2} < n$$

$$m \approx \sqrt{n}$$

$$\Rightarrow T(n) = O(\sqrt{n})$$

② $T(n) = T(n-1) + T(n-2) + O(1)$

$$T(n) \approx 2T(n-1) + 1 \text{ --- (1)}$$

put $n = n-1$ in eqn ①

$$T(n-1) \approx 2T(n-2) + 1 \text{ --- (2)}$$

put ② into ①

$$T(n) \approx 2[2T(n-2) + 1] + 1$$

$$T(n) \approx 4T(n-2) + 2 + 1 \text{ --- (3)}$$

put $n-2$ into ①

$$T(n-2) \approx 2T(n-3) + 1 \text{ --- (4)}$$

put ④ into ③

$$T(n) \approx 4[2T(n-3) + 1] + 2 + 1$$

$$T(n) \approx 8T(n-3) + 4 + 2 + 1$$

$$\Rightarrow T(n) = 2^K T(n-K) + 2^{K-1}$$

$$T(n) = 2^K (T(n-K) + 1/2)$$

$$\text{let } n-K = 1$$

$$= 2^K (1) \Rightarrow 2^K$$

$$T(n) = O(2^n) \quad \underline{\text{Ans}}$$

③ i) $n(\log n)$

```
void mergesort (int arr[], int l, int r)
{
    if (l < r)
    {
        int mid = l + (r-l)/2;
        mergesort (arr, l, mid);
        mergesort (arr, mid+1, r);
        merge (arr, l, mid, r);
    }
}
```

```
void merge (int arr[], int l, int mid, int r)
{
    int i = l;
    int t[arr.length];
    int i = l, j = mid+1, k = 0;
    while (i <= mid && j <= r)
    {
        if (arr[i] < arr[j])
            t[k++] = arr[i++];
        else t[k++] = arr[j++];
    }
    while (i <= mid)
        t[k++] = arr[i++];
    while (j <= r)
        t[k++] = arr[j++];
    for (int i = 0; i < arr.length + 1; i++)
    {
        arr[i] = t[i];
    }
}
```

③ (ii) n^3

```
for (int i = 0; i < n; i++)  
{  
    for (int j = 0; j < n; j++)  
    {  
        for (int k = 0; k < n; k++)  
            count++;  
    }  
}
```

④ (iii) $\log(\log n)$

```
for (i = 2; i < n; i = i * i)  
    c++;
```

$$Q(4) \quad T(n) = T(n/4) + T(n/2) + cn^2$$

$$T(n) = 1 \cdot T(n/2) + cn^2$$

master's method

$$a=1, b=2, d(n) = cn^2$$

$$c = \log_2 1 = 0$$

$$d(n) > (n)^0$$

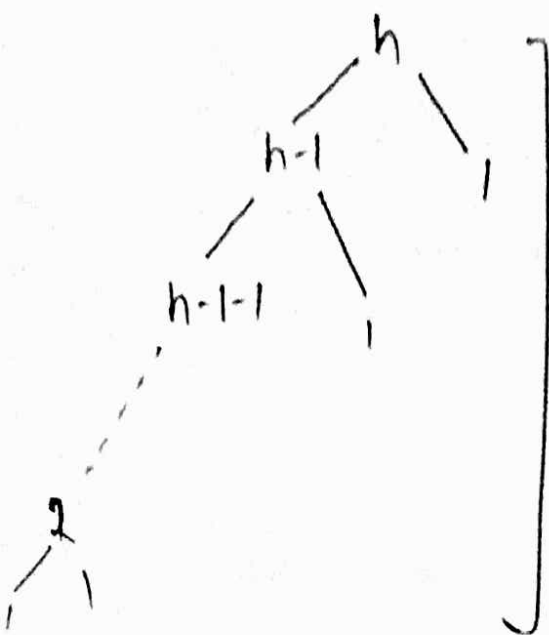
$$T(n) = O(n^2)$$

$$\begin{aligned} (5) \quad T(n) &= \sum_{i=1}^n \sum_{j=1}^{n-1} O(1) \\ &= \sum_{i=1}^n (n-1) \\ &= n(n-1) \\ &= n^2 - n \\ &= O(n^2) \quad \underline{\text{Ans}} \end{aligned}$$

$$\begin{aligned} (6) \quad i &= 2^1, 2^k, 2^{k^2}, 2^{k^3} \dots n \\ a_n &= 2^{k^m} < n \\ k^m &< \log n \\ m &< \log_k \log_2 n \end{aligned}$$

$$\sum_{k=1}^m 1 + 1 + 1 + \dots \text{m times} = O(\log_k \log_2 n)$$

① $T(n) = T(n-1) + O(1)$



$$T(n) = (T(n-1) + T(n-2) + \dots + T(1))$$

(1 + 1 + 1 + ... n times)

$$T(n) = n \times n$$

$$T(n) = O(n^2)$$

lowest level = 2

highest level = n

diff = n - 2

⑧

a) $100 < \log \log n < \log n < (\log n)^2 < \sqrt{n} < n < n \log n < \log(n!) < n^2 < 2^n < n^n < 2^{2^n}$

b) $1 < \log \log n < \sqrt{\log n} < \log n < \log_2 n < 2 \log n < n < n \log n < 2n < n^n < \log(n!) < n^2 < n < 2 \log_2 n < 5n$

c) $96 < \log_8 n < \log_2 n < 5n < n \log_8 n < n \log_2 n < \log(n!) < 3n^2 < 7n^3 < n! < 8^{2^n}$