

## **Experiment No.4**

### **Title: - First Sets**

**Aim:-** Write a program to find first() sets of given Grammar

#### **Finding First:-**

**First(x)** for all grammar symbols X

Apply following rules:

1. If X is terminal,  $\text{FIRST}(X) = \{X\}$ .
2. If  $X \rightarrow \epsilon$  is a production, then add  $\epsilon$  to  $\text{FIRST}(X)$ .
3. If X is a non-terminal, and  $X \rightarrow Y_1 Y_2 \dots Y_k$  is a production, and  $\epsilon$  is in all of  $\text{FIRST}(Y_1), \dots, \text{FIRST}(Y_k)$ , then add  $\epsilon$  to  $\text{FIRST}(X)$ .
4. If X is a non-terminal, and  $X \rightarrow Y_1 Y_2 \dots Y_k$  is a production, then add a to  $\text{FIRST}(X)$  if for some i, a is in  $\text{FIRST}(Y_i)$ , and  $\epsilon$  is in all of  $\text{FIRST}(Y_1), \dots, \text{FIRST}(Y_{i-1})$ .

Applying rules 1 and 2 is obvious. Applying rules 3 and 4 for  $\text{FIRST}(Y_1 Y_2 \dots Y_k)$  can be done as follows:

Add all the non- $\epsilon$  symbols of  $\text{FIRST}(Y_1)$  to  $\text{FIRST}(Y_1 Y_2 \dots Y_k)$ . If  $\epsilon \in \text{FIRST}(Y_1)$ , add all the non- $\epsilon$  symbols of  $\text{FIRST}(Y_2)$ . If  $\epsilon \in \text{FIRST}(Y_1)$  and  $\epsilon \in \text{FIRST}(Y_2)$ , add all the non- $\epsilon$  symbols of  $\text{FIRST}(Y_3)$ , and so on. Finally, add  $\epsilon$  to  $\text{FIRST}(Y_1 Y_2 \dots Y_k)$  if  $\epsilon \in \text{FIRST}(Y_i)$ , for all  $1 \leq i \leq k$ .

#### **Example:**

Consider the following grammar.

$$E \rightarrow E + T \mid T$$
$$T \rightarrow T * F \mid F$$
$$F \rightarrow (E) \mid \text{id}$$

Grammar after removing left recursion:

$E \rightarrow TX$

$X \rightarrow +TX \mid \varepsilon$

$T \rightarrow FY$

$Y \rightarrow *FY \mid \varepsilon$

$F \rightarrow (E) \mid id$

For the above grammar, following the above rules, the FIRST sets could be computed as follows:

$FIRST(E) = FIRST(T) = FIRST(F) = \{ (, id \}$

$FIRST(X) = \{ +, \varepsilon \}$

$FIRST(Y) = \{ *, \varepsilon \}$

**Exercise:-**

**Q1. The following grammar can be used to describe traveling schemes:**

$TS \rightarrow TS \text{ Time Time } TS \mid \text{Station}$

$\text{Station} \rightarrow \text{Identifier}$

$\text{Time} \rightarrow \text{Nat} : \text{Nat}$

Which of the following grammars is equivalent but no longer left-recursive? **(Tick Mark right answer)**

a)  $TS \rightarrow TS (\text{Time Time Station}) *$

$\text{Station} \rightarrow \text{Identifier}$

$\text{Time} \rightarrow \text{Nat} : \text{Nat}$

c)  $TS \rightarrow Z \text{ Time Time } TS \mid \text{Station}$

$Z \rightarrow TS \mid \#$

$\text{Station} \rightarrow \text{Identifier}$

b)  $TS \rightarrow \text{Station} \mid \text{Station } Z$

$Z \rightarrow \text{Time Time } TS \mid \text{Time Time } TS Z$

$\text{Station} \rightarrow \text{Identifier}$

$\text{Time} \rightarrow \text{Nat} : \text{Nat}$

d)  $TS \rightarrow \text{Station Time Time } Z$

$Z \rightarrow \text{Station} \mid TS$

$\text{Station} \rightarrow \text{Identifier}$

Time  $\rightarrow$  Nat : Nat

Time  $\rightarrow$  Nat : Nat

**Q2. Consider the grammar**

$S \rightarrow SX \mid SSb \mid XS \mid a$

$X \rightarrow Xb \mid Sa \mid b$

Eliminate left recursion and rewrite the grammar.

**Answer:-**

**Q3. Is the resulting grammar of Q2. suitable for top down parsing?**

**Answer:**