

# Time Complexity of Prime Checking Algorithm

Rohan Thomas

## 1 Problem

Given any number  $n \in \mathbb{N}$  such that  $n > 1$ , check if  $n$  is a **prime** or not

## 2 Algorithm

The intuition is to start from  $k = 2 \in \mathbb{N}$  and iterate  $k$  by checking if  $k$  divides  $n$ . but before we formally state the algorithm we should answer the question, **till where should we check ?**

That is to find the **worst case scenario** of prime checking algorithm

**Theorem:** For any composite number  $n > 2 \in \mathbb{N}$ , there exist a prime divisor less than or equal to  $\sqrt{n}$ .

**Proof:** Since  $n$  is a composite number, there exists  $a, b < n \in \mathbb{N}$  such that  $n = ab$   
suppose  $a, b > \sqrt{n} \implies n = ab > n$  which is a contradiction.  
 $\implies$  is either  $a$  or  $b$  is less than or equal to  $\sqrt{n}$

Assume WLOG,  $a \leq \sqrt{n}$ , by fundamental theorem of arithmetic, there exist a prime  $p \leq a$  such that  $p \mid a$ . but  $n = ab \implies a \mid n$ . by transitivity of divisibility relation  
 $\implies p \mid n \quad \square$

Notice, we now we can formally state the algorithm,

---

**Algorithm 1** Checking if a number is prime

---

```
 $n, flag \leftarrow 0$ 
for  $i = 2$  to  $\sqrt{n}$  do
  if  $n \bmod i = 0$  then
     $flag \leftarrow 1$ 
  end if
end for
if  $flag = 1$  then
  not a prime
else
  prime
end if
```

---

### 3 Graphing the algorithm

We can plot the algorithm in number of input  $n$  versus number of iterations.  
notice the graph is bounded by the function  $\sqrt{n}$

