

Operating System Assignment - 4

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1. Distributed Deadlocks Detection Simulation - → Fragments :

- $S_1 : P_1 \rightarrow P_2 \rightarrow P_3 \rightarrow P_4$
- $S_2 : P_2 \rightarrow P_3, P_3 \rightarrow P_4$
- $S_3 : P_4 \rightarrow P_1$

a) Global wait-for Graph (combined)
 $P_1 \rightarrow P_2 \rightarrow P_3 \rightarrow P_4 \rightarrow P_1$ (cycle) Also $P_3 \rightarrow P_4$ (separate)

b) Deadlock?

→ Yes Process involved in dead lock cycle
 P_1, P_2, P_3, P_4

c) Suggested distributed algorithm

Use the Chandy - Misra - Mass edge-chasing (Probe) algorithm for distributed deadlock detection each site send Probes along wait-for edges to detect cycles without centralized Graph assembly.

2. Distributed File System Performance :

Given : local = S ms, remote = 25 ms

Prob (remote) = 0.3

a) Expected access time :

$$E = 0.7 \times S \text{ ms} + 0.3 \times 25 \text{ ms} = 3S + 37.5 = 11 \text{ ms}$$

b) The caching strategy :

Client-side Read cache with LRU + TT2-based validation

→ Justification. Frequently - Read Remark File will be served locally reducing Remote access (0.3 fraction). LRU unicks des-used / items TTL keeps staleness bounded. Implements average latency while keeping consistency manageable.

3. Checkpointing Mix to Meet RPO = 1s

Sol Given,

Full = Perome, Incremental = some, RPO = 1s

a) Proposed mix (over 10s)

- Take one full checkpoint every 10s (at $t=0$ in the period)
 - Take incremental checkpoints every 2s (at $t=4, 2$)
- Total overhead (Per loss) : $1 \times 200ms + 9 \times 50ms$
 $= 200 + 450 = 650ms$

b) Reasoning

- with incremental every 1s, the maximum work lost on failure $\leq 1s \rightarrow$ meets \rightarrow RPO
- Full once per 10s bound Recovery time 1s by restoring an older full + a small no of incrementals is feasible)
- This mix minimize full checkpoint cost while keeping incrementals frequently high enough to meet RPO.

4. Case Study - Global E-commerce Platform

- a) Distributed scheduling challenges FRB flash sales.
- Massive sudden spikes in Request.

- Geographic distribution of latency data locality matters for latency & inventory correctness
- Heterogeneous nodes

b) Fault - tolerance Strategy (RTO & RPO):

- Active-active multiple-region development:
sol. Service Run concurrently in multiple regions
so failure is seamless.
(RTO = near-zero at service lead)

• Data Strategy:

- Critical transactional data: Use synchronous replication within the region to guarantee consistency & low RPO; cross-region replication but with frequent replication to keep RPO small
- Data log / less-critical data: Uses multi-region eventual consistency with frequent async replication & local caches.

- Recovery Plans: Automated failover by health-check & Global load balance, warm stand by replicas ready to accept traffic, & durable messaging to buffer incoming requests during failure
- Operational measures: chaos testing, backups, rollback & automated scaling to minimize RTO, use snapshots + incremental backups to band RPO

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